cartierSolution6.v

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Proph

https://gitee.com/0001337777/cartier/blob/master/cartierSolution6.v
https://gitee.com/0001337777/cartier/blob/master/cartierSolution6.v.pdf

solves half of some question of Cartier which is how to program grammatical polymorph non-contextual ("1-weighted") 2-fold ("2-higher") pairing-projections ("product") ... (this multi-folding is the foundation of homotopy "algebraic topology"/"fibre functor")

SHORT ::

The ends is to do polymorph mathematics which is 2-folded/enriched ref some indexer which is made of all the graphs as indexes and all the graph-morphisms as arrows . Such indexer admits the generating-views ("generators") subindexer (4.5.17.h) made of : the singleton-object $\{0\}$ graph (for « morphisms » , possibly interdependent with « transformations of morphisms » via non-grammatical "Yoneda" ...) , and the singleton-morphism-between-two-distinct-objects $\{0 \sim\!\!> 1\}$ graph (for « transformations of morphisms » , for « left-whisk » composition , for « right-whisk » composition , for « inner » composition along some tight/strict or lax « cut-adherence ») , and the two structural-dividing ("boundary") $\{0\}$ |- $\{0 \sim\!\!> 1\}$ graph-morphisms (for « domain-codomain-morphisms-of-each-transformation » type-indexes) , and the structural-multiplying ("degeneracy") $\{0 \sim\!\!> 1\}$ |- $\{0\}$ graph-morphism (for « unit-transformation-on-each-morphism » type-constructor)

The 2-conversion-for-transformations relation shall convert across two transformations whose domain-codomain-morphisms-computation arguments are not syntactically/grammatically-the-same. But oneself does show that , by logical-deduction [convTransfCoMod_convMorCoMod_dom] [convTransfCoMod_convMorCoMod_cod] , these two domain-codomain-morphisms are indeed 1-convertible ("soundness lemma") .

Finally, some linear total/asymptotic grade is defined on the morphisms/transformations and the tactics-automated degradation lemma shows that each of the conversion indeed degrades the redex morphism/transformation .

For instant first impression , the 2-conversion-relation-for-transformations constructor which says that the first projection morphism/transformation is natural/polyarrowing (commutativity along these structure-arrow-actions: the structural-multiplying-arrow ("degeneracy") action which is the unit-transformation-on-each-morphism ['UnitTransfCoMod] type-constructor, and the two structural-dividing-arrow ("boundary") actions which are the domain-codomain-morphisms-of-each-transformation ['transfCoMod] type-indexes), is written as:

KEYWORDS :: 1337777.000 ; COQ ; cut-elimination ; 2-fold functors ; non-contextual 2-fold pairing-projections ; polymorph metafunctors-grammar ; modos

OUTLINE ::

- Indexer metalogic for 2-fold-enrichment
 - Generating data 2-folded-enriched ref the generating-views subindexer
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```
HINT :: free master-engineering-thesis ; program this grammatical polymorph generated-functoralong-reindexing ( "Kan extension" ) : generatedFunc ( I : IndexerCat ) ( G : GeneratorsCat ) := { R : ReIndexerCat & { f : G \sim generatingFunc R | p : reIndexingFunc R | - I } }
```

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1 Indexer metalogic for 2-fold-enrichment

The ends is to do polymorph mathematics which is 2-folded/enriched ref some indexer which is made of all the graphs as indexes and all the graph-morphisms as arrows . Such indexer admits the generating-views ("generators") subindexer (4.5.17.h) made of : the singleton-object $\{0\}$ graph (for « morphisms » , possibly interdependent with « transformations of morphisms » via nongrammatical "Yoneda" ...) , and the singleton-morphism-between-two-distinct-objects $\{0 \sim\!\!> 1\}$ graph (for « transformations of morphisms » , for « left-whisk » composition , for « right-whisk » composition , for « inner » composition along some tight/strict or lax « cut-adherence ») , and the two structural-dividing ("boundary") $\{0\}$ |- $\{0 \sim\!\!> 1\}$ graph-morphisms (for « domain-codomain-morphisms-of-each-transformation » type-indexes) , and the structural-multiplying ("degeneracy") $\{0 \sim\!\!> 1\}$ |- $\{0\}$ graph-morphism (for « unit-transformation-on-each-morphism » type-constructor)

Again : The ends is to do polymorph mathematics which is 2-folded/enriched ref some indexer (symmetric-associative-monoidal metalogic/metacategory) which is made of all the graphs as indexes and all the graph-morphisms as arrows . Such indexer admits the generating-views ("generators") subindexer made of : the singleton-object {0} graph , and the singleton-morphism-between-twodistinct-objects $\{0 \sim 1\}$ graph , and the two structural-dividing ("boundary") $\{0\}$ |- $\{0 \sim 1\}$ graph-morphims , and the structural-multiplying ("degeneracy") $\{0 \sim 1\} \mid -\{0\}$ graph-morphism . Primo this infers , for the material (as contrasted from metalogical) mathematics , that the morphisms can no longer be touched individually but many morphisms shall be touched at the same time via some indexing/multiplier/shape : when the shape is the singleton-morphism-between-twodistinct-objects {0 ~> 1} graph such touched-morphisms will be named « transformation of morphisms » ; when the shape is the singleton-object {0} graph such touched-morphisms will be named « morphism » . Secondo this infers that the two structural-dividing-arrows ("boundary") actions via the domain-codomain-morphisms-of-each-transformation , and that represented structural-multiplying-arrow ("degeneracy") action is represented via the unit-transformationon-each-morphism . Tertio this infers , regardless that the common operations on the touched-morphisms are multifold/multiplicative , that oneself can avoid the multiplicative/outer/material ("horizontal") composition of transformation-next-transformation (whose output multiplicity is subindexer and instead it is sufficient multiplicative/outer/material composition of transformation-next-morphism (« right-whisk ») and the multiplicative/outer/material composition of morphism-next-transformation (« left-whisk ») (whose output multiplicity is the shape $\{0 \sim 1\}$ inside the subindexer) .

1.1 Indexer metalogic admits some generating-views subindexer , and is non-contextual ("1-weighted")

As common for the more-general multifold-enriched polymorh mathematics , the indexer metalogic/metacategory is symmetric associative monoidal ; but for the 2-fold polymorh mathematics there are 3 contrasts .

Primo contrast: because of the presence of the generating-views subindexer for the 2-fold polymorh mathematics, then the presentation of this subindexer metalogic is blended with the presentation of its action on the material mathematics. This infers that the two structural-dividing-arrows ("boundary") actions are represented via the domain-codomain-morphisms-of-each-transformation type-indexes of the type-family [transfCoMod], and that the structural-multiplying-arrow ("degeneracy") action is represented via the unit-transformation-on-each-morphism [UnitTransfCoMod] type-constructor (the constructor [UnitTransfCoMod] of the type-family [transfCoMod], which is elsewhere also hidden/blended in the outer left/right-whisk cut constructors [TransfCoMod_PolyMorCoMod_PolyMorCoMod_Post]).

Secondo contrast : here there are none parameter/customized-arrow action (non-structural reindexing , customized boundary-or-degeneracy) ; is it possible to make sense of such ?

Tertio contrast: for now , oneself shall only describe non-contextual ("1-weighhed") pairing-projections , this infers that there is no grammatically-distinguished context constructor in the metalogic and consequently-for-the-material-mathematics that each projection outputs (as minimum-factor-weight as) some singleton-morphism and that the pairing ingets/inputs (as minimum-factor-weight as) two singleton-morphisms . In the future , the description of contextual pairing-projections would require some interdependence between the presentation of morphisms and the presentation of transformations , BUT THIS DEPENDENCE NEED-NOT BE GRAMMATICAL! (inductive-inductive types) , as long as the generating morphisms-data are actually polymorph generating-views ... This dependence could be expressed via the sense-decoding ("Yoneda") of the grammatical transformations .

```
From mathcomp
Require Import ssreflect ssrfun ssrbool eqtype ssrnat seq choice fintype tuple.
Require Omega Psatz. (* Omega.omega is too weak for degradeMor at
Pairing_Mor_morphism , also degradeTransf *)
Require Coq.Logic.Eqdep_dec.

Module TWOFOLD.

Set Implicit Arguments.
Unset Strict Implicit.
Unset Printing Implicit Defensive.

Arguments Nat.sub : simpl nomatch.
Arguments Nat.add : simpl nomatch.
Delimit Scope poly_scope with poly.
Open Scope poly.
```

1.2 Generating data 2-folded-enriched ref the generating-views subindexer

As common , oneself shall start from some generating data which is 2-folded/enriched ref the generating-views subindexer . But because this subindexer is non-contextual ("1-weighted") there will be no-surprise and no-contrast from the more-general multifold-enriched polymorph mathematics ; therefore this part is not described for now .

In the future , the description of contextual pairing-projections would require some interdependence between the presentation of morphisms and the presentation of transformations , BUT THIS DEPENDENCE NEED-NOT BE GRAMMATICAL ! (inductive-inductive types) , as long as the generating morphisms-data are actually polymorph generating-views ... This dependence could be expressed via the sense-decoding ("Yoneda") of the grammatical transformations .

```
Parameter obCoMod_Gen : <u>Type</u>.
Parameter morCoMod_Gen : <u>forall</u> (F G : obCoMod_Gen), <u>Type</u>.
Parameter transfCoMod_Gen : <u>forall</u> (F G : obCoMod_Gen) (g g' : morCoMod_Gen F G), <u>Type</u>.
```

2 Grammatical presentation of objects and touched-morphisms 2-folded/enriched ref the generating-views subindexer

For 2-folded/enriched polymorph mathematics , each object can be touched individually but the morphisms can no longer be touched individually , and many morphisms shall be touched at the same time via some indexing/multiplier/shape : when the shape is the singleton-morphism-between-two-distinct-objects $\{0 \sim 1\}$ graph such touched-morphisms will be named « transformation of morphisms » ; when the shape is the singleton-object $\{0\}$ graph such touched-morphisms will be named « morphism » .

Each decoding ("Yoneda") of some index-for-touched-morphisms which encodes all the touched-morphisms-at-some-domain-codomain is some metafunctor-on-the-subindexer , which is therefore programmed by some inductive-family-presentation [morCoMod] for the shape $\{0\}$ (morphisms) together with some inductive-family-presentation [transfCoMod] for the shape $\{0\}$ ~> 1} (transformations-of-morphisms) , which could possibly be interdependent (non-grammatically "Yoneda" ...) . Now the inductive-family-presentation [transfCoMod] has some additional/embedded type-indexes and type-constructor : the domain-codomain-morphisms-of-each-transformation [transfCoMod] type-indexes to represent the two structural-dividing-arrow ("boundary") actions, and the unit-transformation-on-each-morphism [UnitTransfCoMod] type-constructor to represent the structural-multiplying-arrow ("degeneracy") action .

Each decoding ("Yoneda") of the whatever-is-interesting arrows between the indexes-for-touched-morphisms are metatransformations which are programmed as some grammatical-constructors of the inductive-family-presentations [morCoMod] and [transfCoMod] .

Memo that the functoriality ("arrows-action") of each metafunctor (decoded index-for-touched-morphisms) and the naturality ("arrows-action") of each metatransformation (decoded arrow-between-indexes) is signified via the additional/embedded type-indexes of [transfCoMod] and type-

constructor [UnitTransfCoMod] of [transfCoMod] . All this is effected via the two conversion relations [convMorCoMod] [convTransfCoMod] which relate those grammatical-touched-morphisms : [convMorCoMod] is for morphisms and [convTransfCoMod] is for transformations .

For 2-folded-enriched polymorph mathematics , the common operations on the touched-morphisms are multiplicative ; this contrast from internal polymorph mathematics where many morphisms are touched at the same time and moreover many objects are touched at the same time and moreover the common operations on the objects or touched-morphisms are coordinatewise/dimensional/pointwise . Memo that here the (multiplicative) outer/material ("horizontal") composition [PolyMorCoMod] [TransfCoMod_PolyMorCoMod_PolyMorCoMod_Post] is some common operation , but there is also some uncommon operation [PolyTransfCoMod] which is the (coordinatewise/dimensional/pointwise) inner/(structure-logical) ("vertical") composition of transformation-later-transformation (along some tight/strict or lax « cut-adherence ») inside each enrichment/indexer-graph ; and both compositions cut-constructors shall be eliminated/erased

Memo that , for the material mathematics , the decidable equality [obCoMod_eq] on the objects will enable to do any logical-inversion of the very-dependently-typed propositional-equality-across-any-two-morphisms [Inversion_Project1] [Inversion_Exfalso] , and will also enable to do the logical-inversion of any morphism whose domain-codomain-objects are the same [Inversion_domEqcod] [Inversion toPolyMor] .

```
Inductive obCoMod : Type :=
(** | ObCoMod_Gen : obCoMod_Gen -> obCoMod *)
| Pair : obCoMod -> obCoMod -> obCoMod .
Module ObCoMod_eq.
Definition obCoMod_eq : \underline{\text{forall }} F G : \text{obCoMod}, \{F = G\} + \{ \sim F = G\} .
Proof.
  (** decide equality. *)
  induction F.
  <u>destruct</u> G.
  destruct (IHF1 G1).
  - { destruct (IHF2 G2).
       + <u>left</u>. <u>clear</u> IHF1 IHF2. <u>subst</u>; reflexivity.
       + <u>right</u>. <u>clear</u> IHF1 IHF2. abstract (<u>subst</u>; <u>simplify_eq</u>; done).
  - <u>right</u>. <u>clear</u> IHF1 IHF2. abstract (<u>subst</u>; <u>simplify eq</u>; done).
Defined.
Definition obCoMod_eqP : forall F : obCoMod, obCoMod_eq F F = left (Logic.eq_refl F).
Proof. <u>induction</u> F. <u>simpl</u>. <u>rewrite</u> IHF1 IHF2. reflexivity. Qed.
Definition Eqdep_dec_inj_pair2_eq_dec
  : \underline{forall} (P: obCoMod -> \underline{Type}) (p: obCoMod) (x y: P p),
     \frac{1}{\text{existT}} P p x = \text{existT} P p y -> x = y
  := Eqdep_dec.inj_pair2_eq_dec _ ObCoMod_eq.obCoMod_eq.
End ObCoMod eq.
Reserved Notation "''morCoMod' (0 F' ~> F )0"
          (at level 0, format "''morCoMod' (0 F' \sim > F )0").
Inductive morCoMod : obCoMod -> obCoMod -> Type :=
| PolyMorCoMod : forall (F F' : obCoMod),
       'morCoMod(0 \overline{F'} \sim F)0 \rightarrow \underline{forall} (F'' : obCoMod),
          'morCoMod(0 F'' ~> F' )0 -> 'morCoMod(0 F'' ~> F )0
| UnitMorCoMod : <u>forall</u> (F : obCoMod),
     'morCoMod(0 F \sim> F)0
(** | MorCoMod Gen : forall (F G : obCoMod Gen),
    morCoMod_Gen F G -> 'morCoMod(0 (ObCoMod_Gen F) ~> (ObCoMod_Gen G) )0 *)
```

```
| Project1_Mor : <u>forall</u> (F1 F2 : obCoMod) (Z1 : obCoMod),
    'morCoMod(0 F1 ~> Z1 )0 ->
    'morCoMod(0 (Pair F1 F2) ~> Z1 )0
| Project2 Mor : forall (F1 F2 : obCoMod) (Z2 : obCoMod),
    'morCoMod(0 F2 ~> Z2 )0 ->
    'morCoMod(0 (Pair F1 F2) ~> Z2 )0
| Pairing_Mor : forall (L : obCoMod) (F1 F2 : obCoMod),
    'morCoMod(0 L ~> F1 )0 -> 'morCoMod(0 L ~> F2 )0 ->
    'morCoMod(0 L ~> (Pair F1 F2) )0
where "''morCoMod' (0 F' ~> F )0" := (@morCoMod F' F) : poly_scope.
Notation "ff o>CoMod ff'" :=
  (@PolyMorCoMod ___ ff' _ ff_) (<u>at</u> level 40 , ff' <u>at</u> next level) : poly_scope.
Notation "@ ''UnitMorCoMod' F" := (@UnitMorCoMod F)
                                      (<u>at</u> level 10, only parsing) : poly_scope.
Notation "''UnitMorCoMod'" := (@UnitMorCoMod _) (<u>at</u> level 0) : poly_scope.
(** Notation "''MorCoMod_Gen' ff" :=
      (@MorCoMod_Gen _ _ _ ff) (at level 3) : poly_scope. **)
(* @ in ~_1 @
                  says argument *)
Notation "~ 1 @ F2 o>CoMod z1" :=
  (@Project1_Mor _ F2 _ z1) (<u>at</u> level 4, F2 <u>at</u> next level) : poly_scope.
Notation "\sim 1 o>CoMod z1" :=
  (@Project1_Mor _ _ z1) (<u>at</u> level 4) : poly_scope.
Notation "~ 2 @ F1 o>CoMod z2" :=
  (@Project2_Mor F1 _ _ z2) (<u>at</u> level 4, F1 <u>at</u> next level) : poly_scope.
Notation "\sim 2 o>CoMod z2" :=
  (@Project2_Mor _ _ z2) (<u>at</u> level 4) : poly_scope.
Notation "<< f1 , CoMod f2 >>" :=
  (@Pairing_Mor _ _ _ f1 f2) (at level 4, f1 at next level, f2 at next level,
                                format "<< f1 , CoMod f2 >>" ) : poly scope.
```

3 Solution morphisms

As common, the purely-grammatical polymorphism cut-constructors , for (multiplicative) outer/material composition [PolyMorCoMod] [TransfCoMod_PolyMorCoMod_Pre] [TransfCoMod_PolyMorCoMod_Post] and (coordinatewise) inner/structural composition [PolyTransfCoMod] , are not part of the solution terminology .

3.1 Solution morphisms without polymorphism

```
Module Sol.
Section Section1.
Delimit Scope sol_scope with sol.
Open Scope sol_scope.

Inductive morCoMod : obCoMod -> obCoMod -> Type :=

| UnitMorCoMod : forall (F : obCoMod),
    'morCoMod(0 F ~> F )0

| Project1_Mor : forall (F1 F2 : obCoMod) (Z1 : obCoMod),
    'morCoMod(0 F1 ~> Z1 )0 ->
```

```
'morCoMod(0 (Pair F1 F2) ~> Z1 )0
| Project2 Mor : forall (F1 F2 : obCoMod) (Z2 : obCoMod),
    'morCoMod(0 F2 \sim> Z2 )0 \rightarrow
    'morCoMod(0 (Pair F1 F2) ~> Z2 )0
| Pairing Mor : forall (L : obCoMod) (F1 F2 : obCoMod),
    'morCoMod(0 L ~> F1 )0 -> 'morCoMod(0 L ~> F2 )0 ->
    'morCoMod(0 L ~> (Pair F1 F2) )0
where "''morCoMod' (0 F' \sim> F )0" := (@morCoMod F' F) : sol_scope.
End Section1.
Module Export Ex Notations.
  Delimit Scope sol_scope with sol.
  Notation "''morCoMod' (0 F' \sim> F )0" := (@morCoMod F' F) : sol scope.
  Notation "@ ''UnitMorCoMod' F" := (@UnitMorCoMod F)
                                         (<u>at</u> level 10, only parsing) : sol_scope.
  Notation "''UnitMorCoMod'" := (@UnitMorCoMod _) (<u>at</u> level 0) : sol_scope.
  (* @ in ~_1 @ says argument *)
  Notation "\sim \overline{1} @ F2 o>CoMod z1" :=
    (@Project1_Mor _ F2 _ z1) (<u>at</u> level 4, F2 <u>at</u> next level) : sol_scope.
  Notation "~ 1 o>CoMod z1" :=
    (@Project1_Mor _ _ z1) (<u>at</u> level 4) : sol_scope.
  Notation "~ 2 @ F1 o>CoMod z2" :=
    (@Project2_Mor F1 ___ z2) (<u>at</u> level 4, F1 <u>at</u> next level) : sol scope.
  Notation "~ 2 o>CoMod z2" :=
    (@Project2_Mor _ _ z2) (<u>at</u> level 4) : sol_scope.
  Notation "<< f1 , CoMod f2 >>" :=
    (@Pairing_Mor _ _ _ f1 f2) (at level 4, f1 at next level, f2 at next level,
                                  format "<< f1 , CoMod f2 >>" ) : sol scope.
End Ex_Notations.
Fixpoint toPolyMor (F G : obCoMod) (G : 'morCoMod(0 F \sim> G )0 %sol)
         \{ struct \ g \} : 'morCoMod(0 F \sim G )0 %poly .
Proof.
  <u>refine</u>
    match g with
    | ( @'UnitMorCoMod F )%sol => ( @'UnitMorCoMod F )%poly
    | ( \sim_1 @ F2 o>CoMod z1 )%sol => ( \sim_1 @ F2 o>CoMod (toPolyMor _ _ z1) )%poly
    | ( ~_2 @ F1 o>CoMod z2 )%sol => ( ~_2 @ F1 o>CoMod (toPolyMor _ _ z2) )%poly
    | ( << f1 ,CoMod f2 >> )%sol =>
      ( << (toPolyMor _ _ f1) ,CoMod (toPolyMor _ _ f2) >> )%poly
    end.
Defined.
```

3.2 Inversion of morphisms with same domain-codomain objects

In contrast to some dependent-destruction of morphisms , this dependent-inversion of morphisms with same domain-codomain objects , is logical/propositional , therefore it is only usable during deductions/proofs . But memo that it is also possible to program some nondependent-destruction of morphisms with same domain-codomain objects which is usable during both programming/data and deductions/proofs .

```
Module Inversion domEqcod.
Inductive morCoMod domEqcod : \underline{forall} (F : obCoMod), 'morCoMod(0 F \sim F )0 %sol -> \underline{Prop} :=
| UnitMorCoMod : forall (F : obCoMod),
    morCoMod domEqcod (@'UnitMorCoMod F)%sol
| Project1 Mor : forall (F1 F2 : obCoMod), forall (z1 : 'morCoMod(0 F1 ~> Pair F1 F2 )0%sol),
      morCoMod_domEqcod ( ~_1 @ F2 o>CoMod z1 )%sol
| Project2 Mor : forall (F1 F2 : obCoMod), forall (z2 : 'morCoMod(0 F2 ~> Pair F1 F2 )0%sol),
      morCoMod_domEqcod ( ~_2 @ F1 o>CoMod z2 )%sol
| Pairing Mor : forall (F1 F2 : obCoMod) (f1 : 'morCoMod(0 (Pair F1 F2) ~> F1 )0 %sol)
                    (f2 : 'morCoMod(0 (Pair F1 F2) ~> F2 )0 %sol),
    morCoMod domEqcod ( << f1 ,CoMod f2 >> )%sol .
Definition morCoMod domEqcodP Type
            (F G : obCoMod) (g : 'morCoMod(0 F <math>\sim G )0 %sol ) : <u>Type</u>.
Proof.
  destruct (ObCoMod_eq.obCoMod_eq F G).
  - <u>destruct</u> e. <u>refine</u> (morCoMod_domEqcod g).
  - intros; refine (unit : Type).
Defined.
Lemma morCoMod domEgcodP
  : forall (FG: obCoMod) (g: 'morCoMod(0 F \sim G )0 %sol), morCoMod_domEqcodP_Type g .
Proof.
  <u>intros</u>. <u>case</u>: F G / g.
    intros F. unfold morCoMod_domEqcodP_Type. simpl.
    <u>rewrite</u> ObCoMod eq.obCoMod eqP.
    constructor 1.
    intros ? ? Z1. destruct Z1 as [Z1_1 Z1_2]. intros.
    <u>unfold</u> morCoMod_domEqcodP_Type. <u>simpl</u>.
    { destruct (ObCoMod_eq.obCoMod_eq F1 Z1_1).
      * { destruct (ObCoMod_eq.obCoMod_eq F2 Z1_2).

    simpl. subst. simpl. constructor 2.

           intros; exact: tt.
        }
      * <u>intros</u>; exact: tt.
    }
    <u>intros</u> ? ? Z2 * . <u>destruct</u> Z2 <u>as</u> [Z2_1 Z2_2].
    <u>unfold</u> morCoMod_domEqcodP_Type. <u>simpl</u>.
    { <u>destruct</u> (ObCoMod_eq.obCoMod_eq F1 Z2_1).
      * { destruct (ObCoMod eq.obCoMod eq F2 Z2 2).

    simpl. subst. simpl. constructor 3.

           - <u>intros</u>; exact: tt.
        <u>intros</u>; exact: tt.
  - intros L *. destruct L as [L1 L2]. unfold morCoMod domEqcodP Type. simpl.
    { <u>destruct</u> (ObCoMod_eq.obCoMod_eq L1 F1).
      + { destruct (ObCoMod eq.obCoMod eq L2 F2).
           - <u>simpl</u>. <u>subst</u>. <u>simpl</u>. <u>constructor</u> 4.
           intros; exact: tt.
      + intros; exact: tt.
Qed.
End Inversion domEqcod.
```

3.3 Destruction of morphisms with inner-instantiation of object-indexes

For the [morCoMod] inductive-family-presentation , there are no extra-argument/parameter (for example , the domain-codomain morphisms in [transfCoMod]) beyond the domain-codomain-objects , therefore this is the common dependent-destruction of morphisms with inner-instantiation of object-indexes

```
Module Destruct_domPair.
Inductive morCoMod domPair
: <u>forall</u> (F1 F2 : obCoMod), <u>forall</u> (G : obCoMod),
       'morCoMod(0 (Pair F1 F2) ~> G )0 %sol -> <u>Type</u> :=
| UnitMorCoMod : <u>forall</u> (F1 F2 : obCoMod),
    morCoMod_domPair ( @'UnitMorCoMod (Pair F1 F2) )%sol
| Project1 Mor : forall (F1 F2 : obCoMod),
    forall (Z1 : obCoMod) (z1 : 'morCoMod(0 F1 ~> Z1 )0 %sol),
      morCoMod_domPair ( ~_1 @ F2 o>CoMod z1 )%sol
| Project2_Mor : forall (F1 F2 : obCoMod),
    forall (Z2 : obCoMod) (z2 : 'morCoMod(0 F2 ~> Z2 )0 %sol),
      morCoMod_domPair ( ~_2 @ F1 o>CoMod z2 )%sol
| Pairing Mor :
    forall (M M' : obCoMod) (F1 F2 : obCoMod) (f1 : 'morCoMod(0 (Pair M M') ~> F1 )0 %sol)
      (f2 : 'morCoMod(0 (Pair M M') ~> F2 )0 %sol),
    morCoMod_domPair ( << f1 ,CoMod f2 >> )%sol .
Definition morCoMod_domPairP_Type
           (F G : obCoMod) (g : 'morCoMod(0 F ~> G )0 %sol ) :=
  ltac:( destruct F; [ (*intros; refine (unit : Type)
                          | *) refine (morCoMod_domPair g) ] ).
Lemma morCoMod domPairP
  : forall (F G : obCoMod) (g : 'morCoMod(0 F \sim> G )0 %sol), morCoMod domPairP Type g .
Proof.
  intros. case: F G / g.
  - destruct F; [ (*intros; exact: tt |*) ].
    constructor 1.

    constructor 2.

  - <u>constructor</u> 3.
   destruct L; [ (* intros; exact: tt | *) ].
    constructor 4.
Defined.
End Destruct_domPair.
Module Destruct_codPair.
Inductive morCoMod codPair
: <u>forall</u> (F : obCoMod), <u>forall</u> (G1 G2 : obCoMod),
       'morCoMod(0 F ~> (Pair G1 G2) )0 %sol -> <u>Type</u> :=
| UnitMorCoMod : <u>forall</u> (F1 F2 : obCoMod),
    morCoMod_codPair ( @'UnitMorCoMod (Pair F1 F2) )%sol
| Project1 Mor : <u>forall</u> (F1 F2 : obCoMod),
    <u>forall</u> (Z1 Z1' : obCoMod) (z1 : 'morCoMod(0 F1 ~> (Pair Z1 Z1') )0 %sol),
      morCoMod_codPair ( ~_1 @ F2 o>CoMod z1 )%sol
| Project2 Mor : forall (F1 F2 : obCoMod),
    forall (Z2 Z2' : obCoMod) (z2 : 'morCoMod(0 F2 ~> (Pair Z2 Z2') )0 %sol),
      morCoMod_codPair ( ~_2 @ F1 o>CoMod z2 )%sol
| Pairing Mor : forall (L : obCoMod) (F1 F2 : obCoMod),
    <u>forall</u> (f1 : 'morCoMod(0 L ~> F1 )0 %sol) (f2 : 'morCoMod(0 L ~> F2 )0 %sol),
      morCoMod codPair ( << f1 ,CoMod f2 >> )%sol .
```

```
Definition morCoMod codPairP_Type
  (\textit{F G}: obCoMod) \ (\textit{g}: 'morCoMod(0 F \sim> G )0 %sol ) := \\ \textbf{ltac}: ( \ \underline{destruct} \ G; \ [ \ \textit{(*intros; refine (unit : Type)} ]
                               | *) refine (morCoMod codPair g) ] ).
Lemma morCoMod codPairP
  : forall (F G : obCoMod) (g : 'morCoMod(0 F \sim G )0 %sol ), morCoMod codPairP Type g .
Proof.
  intros. case: F G / g.
  - destruct F; [ (*intros; exact: tt |*) ].
     constructor 1.
  - <u>destruct</u> Z1; [ (*intros; exact: tt |*) ].
     constructor 2.

    destruct Z2; [ (*intros; exact: tt |*) ].

     constructor 3.
    constructor 4.
Defined.
End Destruct codPair.
End Sol.
```

4 Grammatical 2-conversion of transformations , which infer the 1-conversions of their domain-codomain morphisms

As common , the grammatical 1-conversions-for-morphisms [convMorCoMod] ans 2-conversions-for-transformations [convTransfCoMod] are classified into : the total/(multi-step) conversions , and the congruences conversions , and the constant conversions which are used in the polymorphism/cut-elimination lemma , and the constant conversions which are only for the wanted sense of pairing-projections-grammar , and the constant conversions which are only for the confluence lemma , and the constant conversions which are derivable by using the finished cut-elimination lemma .

In contrast , because of the structural-multiplying-arrow ("degeneracy") action which is the unit-transformation-on-each-morphism [UnitTransfCoMod] type-constructor (which is elsewhere also hidden/blended in the outer left/right-whisk cut constructors [TransfCoMod_PolyMorCoMod_Pre] [TransfCoMod_PolyMorCoMod_Post]) , then the 2-conversions-for-transformations [convTransfCoMod] depends/uses of the 1-conversions-for-morphisms [convMorCoMod] , via the conversion-constructors [UnitTransfCoMod_cong] [TransfCoMod_PolyMorCoMod_Pre_cong] [TransfCoMod_PolyMorCoMod_Post_cong] .

Also in contrast , because of the embedded/computed domain-codomain morphisms extra-argument/parameter in the inductive-family-presentation of the transformations , the 2-conversion-for-transformations relation shall convert across two transformations whose domain-codomain-morphisms-computation arguments are not syntactically/grammatically-the-same . But oneself does show that , by logical-deduction [convTransfCoMod_convMorCoMod_dom] [convTransfCoMod_convMorCoMod_cod] , these two domain-codomain-morphisms are indeed 1-convertible ("soundness lemma") .

Finally , some linear total/asymptotic morphism-grade [gradeMor] is defined on the morphisms and another linear total/asymptotic transformation-grade [gradeTransf] , which depends/uses of the morphism-grade [gradeMor] , is defined on the transformations ; and the tactics-automated degradation lemmas shows that each of the 1-conversions-for-morphisms or 2-conversions-for-transformations indeed degrades the redex morphism or transformation . (ERRATA: Memo that this new grade function is simplified in comparison from earlier attempts , because strict-degrading-of-the-conversions is not really required but some form of strict-degrading occurs during the computational/total/asymptotic cut-elimination ...)

4.1 Grammatical 1-conversions of morphisms

```
Reserved Notation "g' <~~1 g" (at level 70).
Inductive convMorCoMod :
  forall (F G : obCoMod) (g g' : 'morCoMod(0 F ~> G )0 %poly), Prop :=
```

```
(** ---- the total/(multi-step) conversions ---- **)
| convMorCoMod Refl : <u>forall</u> (F G : obCoMod) (g : 'morCoMod(0 F \sim G )0 ),
    g <~~1 g
| convMorCoMod Trans : forall (F G : obCoMod) (g : 'morCoMod(0 F \sim G )0 )
    q00 <\sim 1 uTrans -> q00 <\sim 1 q
(** ---- the congruences conversions ---- **)
| PolyMorCoMod cong : forall (F F' : obCoMod) (f' f'\theta : 'morCoMod(0 F' \sim F )0),
    \underline{\text{forall}} (F'': obCoMod) (f_{\underline{}} f_{\underline{}}\theta: 'morCoMod(\theta F'' \sim> F')\theta),
      f'0 <-~1 f' -> f 0 <-~1 f -> ( f 0 o>CoMod f'0 ) <-~1 ( f o>CoMod f' )
| Project1_Mor_cong : forall (F1 F2 : obCoMod) (Z1 : obCoMod)
                          (z1 z1' : 'morCoMod(0 F1 ~> Z1 )0),
  z1' \leftarrow 1 z1 \rightarrow (\sim 1 @ F2 o \sim CoMod z1') \leftarrow 1 (\sim 1 @ F2 o \sim CoMod z1)
| Project2_Mor_cong : forall (F1 F2 : obCoMod) (Z2 : obCoMod)
                          (z2 z2' : 'morCoMod(0 F2 ~> Z2 )0),
  z2' <~~1 z2 -> ( ~_2 @ F1 o>CoMod z2' ) <~~1 ( ~_2 @ F1 o>CoMod z2 )
Pairing_Mor_cong : <u>forall</u> (L : obCoMod) (F1 F2 : obCoMod),
    forall (f1 f1': 'morCoMod(0 L ~> F1 )0) (f2 f2': 'morCoMod(0 L ~> F2 )0),
  f1' \leftarrow 1 f1 \rightarrow f2' \leftarrow 1 f2 \rightarrow ( << f1' , CoMod f2' >> ) <\sim 1 ( << f1 , CoMod f2 >> )
(** ---- the constant conversions which are used during the polyarrowing
elimination ---- **)
(** here there are none parameter/customized-arrow action ( non-structural
reindexing , customized boundary-or-degeneracy ) ; is it possible to make sense of
such ? *)
(** ---- the constant conversions which are used during the polymorphism
elimination ---- **)
| UnitMorCoMod_morphismMor_Pre :
    forall (F F'' : obCoMod), forall (f : 'morCoMod(0 F'' \sim F )0),
         ( f ) <~~1 ( f o>CoMod ( @'UnitMorCoMod F ) )
| UnitMorCoMod_morphismMor_Post :
    forall (F F' : obCoMod), forall (f' : 'morCoMod(0 F \sim F' )0),
        (f') < \sim 1 ((G'UnitMorCoMod F) o > CoMod f')
| Project1_Mor_morphism : <u>forall</u> (F1 F2 : obCoMod) (Z1 : obCoMod),
  \underline{forall} (z1: 'morCoMod(0 F1 \sim Z1 )0), \underline{forall} (Y1: obCoMod) (y: 'morCoMod(0 Z1 \sim Y1 )0),
      ( \sim_1 @ F2 o>CoMod (z1 o>CoMod y) )
        <\sim\sim1 ( ( \sim_1 @ F2 o>CoMod z1 ) o>CoMod y )
| Project2_Mor_morphism : forall (F1 F2 : obCoMod) (Z2 : obCoMod),
  <u>forall</u> (z2 : 'morCoMod(0 F2 ~> Z2 )0), <u>forall</u> (Y2 : obCoMod) (y : 'morCoMod(0 Z2 ~> Y2 )0),
      (\sim 2 \oplus F1 \circ CoMod (z2 \circ CoMod y))
        <\sim 1 ( ( \sim 2 @ F1 o>CoMod z2 ) o>CoMod y )
(**memo: Pairing Mor morphism derivable below *)
| Pairing Mor morphism : forall (L1 L2 : obCoMod) (F1 F2 : obCoMod)
    (f1: 'morCoMod(0 Pair L1 L2 \sim> F1 )0) (f2: 'morCoMod(0 Pair L1 L2 \sim> F2 )0),
    <u>forall</u> (M: obCoMod) (11: 'morCoMod(0 M \sim> L1 )0) (12: 'morCoMod(0 M \sim> L2 )0),
      ( << ( ( << l1 ,CoMod l2 >> ) o>CoMod f1 )
         ,CoMod ( ( << l1 ,CoMod l2 >> ) o>CoMod f2 ) >> )
        <-~1 ( ( << l1 ,CoMod l2 >> ) o>CoMod ( << f1 ,CoMod f2 >> ) )
| Pairing Mor Project1 Mor : forall (L : obCoMod) (F1 F2 : obCoMod)
    (f1 : \text{'morCoMod}(0 L \sim F1))) (f2 : \text{'morCoMod}(0 L \sim F2))
```

```
forall (Z1 : obCoMod) (z1 : 'morCoMod(0 F1 ~> Z1 )0 ),
      ( fl o>CoMod zl )
        <\sim\sim1 ( ( << f1 ,CoMod f2 >> ) o>CoMod ( \sim 1 @ F2 o>CoMod z1 )
              : 'morCoMod(0 L ~> Z1 )0 )
| Pairing Mor Project2 Mor : forall (L : obCoMod) (F1 F2 : obCoMod)
    (f1 : \text{'morCoMod}(0 \ L \sim F1))) (f2 : \text{'morCoMod}(0 \ L \sim F2)))
    forall (Z2 : obCoMod) (z2 : 'morCoMod(0 F2 ~> Z2 )0 ),
      (f2 o>CoMod z2)
        <\sim\sim1 ( ( << f1 ,CoMod f2 >> ) o>CoMod ( \sim 2 @ F1 o>CoMod z2 )
              : 'morCoMod(0 L ~> Z2 )0 )
(** ---- the constant conversions which are only for the wanted sense of
pairing-projections-grammar ---- **)
(** Attention : for non-contextual ( "1-weigthed" ) pairing-projections , none of
    such thing as [Project1_Mor_Project2_Mor_Pairing_Transf] for transformations
    instead of [Project1 Mor Project2 Mor Pairing Mor] for morphisms **)
| Project1 Mor Project2 Mor Pairing Mor : forall (F1 F2 : obCoMod),
    ( @'UnitMorCoMod (Pair F1 F2) )
      <\sim\sim1 ( << ( \sim_{\_}1 @ F2 o>CoMod ( @'UnitMorCoMod F1 ) )
            ,CoMod ( \sim_2 @ F1 o>CoMod ( @'UnitMorCoMod F2 ) ) >> )
(** ---- the constant conversions which are only for the confluence lemma ---- **)
| Pairing Mor morphism Project1 Mor :
    forall (L : obCoMod) (F1 F2 : obCoMod)
      (f1 : 'morCoMod(0 L ~> F1 )0) (f2 : 'morCoMod(0 L ~> F2 )0) (H : obCoMod),
      ( \sim 1 @ H o > CoMod ( << f1 , CoMod f2 >> ) )
         <~~1 ( << ( ~_1 @ H o>CoMod f1 )
              ,CoMod (\sim 1 @ H o>CoMod f2 ) >> )
| Pairing Mor morphism Project2 Mor :
    forall (L : obCoMod) (F1 F2 : obCoMod)
      (f1 : \text{'morCoMod}(0 L \sim F1)0) (f2 : \text{'morCoMod}(0 L \sim F2)0) (H : obCoMod),
      ( \sim_2 @ H o>CoMod ( << f1 ,CoMod f2 >> ) )
        <~~1 ( << ( ~_2 @ H o>CoMod f1 )
              ,CoMod ( \sim 2 @ H o>CoMod f2 ) >> )
(** ---- the constant conversions which are derivable by using the finished
cut-elimination lemma ---- **)
(*TODO: COMMENT *)
| PolyMorCoMod morphism Pre : forall (F F' : obCoMod) (f' : 'morCoMod(0 F' ~> F )0),
    forall (F'' : obCoMod) (f_' : 'morCoMod(\theta F'' \sim> F' )\theta), forall (F''' : obCoMod) (f__ : 'morCoMod(\theta F''' \sim> F'' )\theta),
              o>CoMod f' ) o>\overline{Co}Mod f' )
        <\sim\sim 1 ( f_ o>Co\overline{M}od ( f_' o>CoMod f' ) )
(*TODO: COMMENT *)
| PolyMorCoMod_morphism_Post : forall (F F' : obCoMod) (f : 'morCoMod(0 F' ~> F )0),
    forall (F'' : obCoMod) (f' : 'morCoMod(0 F'' \sim> F')0),
    forall (F''' : obCoMod) (f'' : 'morCoMod(0 F''' ~> F'' )0),
      (f'' o > CoMod (f' o > CoMod f))
        <\sim\sim1 ( ( f'' o>CoMod f' ) o>CoMod f )
**)
(** ---- the constant conversions which are derivable immediately without the
finished cut-elimination lemma ---- **)
(*TODO: COMMENT *)
| Pairing_Mor_morphism_derivable : forall (L : obCoMod) (F1 F2 : obCoMod)
    (f1 : 'morCoMod(0 L \sim F1)0) (f2 : 'morCoMod(0 L \sim F2)0),
    forall (L' : obCoMod) (l : 'morCoMod(0 L' ~> L )0),
```

4.2 Linear total/asymptotic morphism-grade and the degradation lemma

```
Notation max m n := ((Nat.add m (Nat.sub n m))%coq nat).
Definition gradeOb (F: obCoMod) : nat := 0 .
Fixpoint gradeMor (F G : obCoMod) (g : 'morCoMod(0 F \sim G)0) {struct g} : nat .
Proof.
  case : F G / g .
   - <u>intros</u> ? ? f' ? f .
                                        _ f' + gradeMor _ _ f_)%coq_nat))%coq_nat .
     exact: (2 * (S (gradeMor
  - (* memo that the unit-transformation-on-each-morphism [ UnitTransfCoMod ]
         type-constructor is some form of structural-multiplying-arrow (
         "degeneracy" ) action ... now: the unit-morphism-on-each-object [
         UnitMorComod ] can also be seen as some form of action , whose argument is
         this object F *)
     <u>intros</u> F .
     exact: (S ( grade0b F (* = 0 *) )).
  - <u>intros</u> ? ? ? z1 .
     exact: (S (S (gradeMor _ _ z1))).
  - <u>intros</u> ? ? ? z2 .
     exact: (S (S (gradeMor _ z2))).
     <u>intros</u> ? ? ? f1 f2 .
     <u>refine</u> (S (S (max (gradeMor _ _ f1) (gradeMor _ _ f2)))).
Defined.
Lemma gradeMor_gt0 : forall (F G : obCoMod) (g : 'morCoMod(0 F \sim G )0 ),
      ((S 0) <= (gradeMor g))%coq_nat.</pre>
Proof. intros; case : g; intros; apply/leP; intros; simpl; auto. Qed.
Ltac tac gradeMor gt0 :=
  match goal with
               g2 : 'morCoMod(0 _ ~> _ )0 ,
  | [ g1 : 'morCoMod(0 ~>
                      'morCoMod(0 _ ~> _ )0 ,
g3 : 'morCoMod(0 _ ~> _ )0 ,
g4 : 'morCoMod(0 _ ~> _ )0 |- _
     \underline{move} \; : \; (@gradeMor\_gt0 \; \_ \; g1) \; (@gradeMor\_gt0 \; \_ \; g2)
              (@gradeMor_gt0 _ _ g3) (@gradeMor_gt0 _ _ g4)
               morCoMod(0 _ ~> _ )0 ,
g2 : 'morCoMod(0 _ ~> _
  | [ g1 : 'morCoMod(0 ~>
                      \begin{array}{ll} \operatorname{morCoMod}(0 \ \_ \ \sim \ \_ \ )0 \ , \\ \operatorname{g3} : \ '\operatorname{morCoMod}(0 \ \_ \ \sim \ ) \end{array}
                              morCoMod(0 _ ~> _ )0 ,
g4 : 'morCoMod(0 _ ~> _
                                                           )0 |- ] =>
     \begin{array}{c} \underline{\mathsf{move}} \ : \ ( @\mathsf{gradeMor\_gt0} \ \_ \ \_ \ g1 ) \ ( @\mathsf{gradeMor\_gt0} \ \_ \ g2 ) \\ ( @\mathsf{gradeMor\_gt0} \ \_ \ g3 ) \ ( @\mathsf{gradeMor\_gt0} \ \_ \ g4 ) \end{array}
  | [ g1 : 'morCoMod(0 \_ \sim> \_ )0 ,
               )0 |-
                                                              _ ] =>
     move : (@gradeMor_gt0 _ _ g1) (@gradeMor_gt0 _ _ g2) (@gradeMor_gt0 _ _ g3)
  | [ g1 : 'morCoMod(0 _ ~>
                                    _ )0 ,
               g2 : 'morCo\overline{M}od(0 \_ \sim> )
                                            ) 0 | -
     move : (@gradeMor_gt0 _ _ g1) (@gradeMor_gt0 _ _ g2)
  | [ g1 : 'morCoMod(0 _ ~> _ )0 |- _ ] =>
     move : (@gradeMor_gt0 _ _ g1)
```

```
end.
Lemma degradeMor
  : forall (F G : obCoMod) (g g' : 'morCoMod(0 F \sim G )0 ),
       <~~1 g -> ( gradeMor g' <= gradeMor g )%coq nat .
  intros until g'. intros red_g.
elim : F G g g' / red_g;
    try solve [ simpl; rewrite ?/grade0b; intros;
                 abstract Psatz.nia ].
  (*memo: Omega.omega too weak at Pairing_Mor_morphism *)
  (*erase associativities conversions then Qed. *)
Qed.
Ltac tac degradeMor H gradeMor :=
  intuition idtac;
  repeat <u>match</u> goal <u>with</u>
          | [ Hred : ( _ <~~1
                                ) |-
            move : (degradeMor Hred) ; clear Hred
          <u>end</u>;
  move: H_gradeMor; clear; simpl; intros;
  try tac_gradeMor_gt0; intros; Omega.omega.
```

5 Polymorphism/cut-elimination by computational/total/asymptotic/reduction/(multi-step) resolution

For 2-folded polymorph mathematics , this resolution is made of some 1-resolution-for-morphisms [solveMorCoMod] and some 2-resolution-for-transformations [solveTransfCoMod] which depends/uses of this 1-resolution-for-morphisms .

The 1-resolution-for-morphisms [solveMorCoMod] is common , but has more attention into clearly separating the computational data-content function [solveMorCoMod] of the resolution from the derived logical properties [solveMorCoModP] which are satisfied by this function [solveMorCoMod] . In other words , because the 1-resolution-for-morphisms will be used during the 2-resolution-for-transformations [solveTransfCoMod] , then this 1-resolution-for-morphisms must compute somehow (without blockage from opaque logical interference). How compute ? by definitional-metaconversions or by propositional-equations ? Now , because this 1-resolution-for-morphisms function is not programmed by morphisms-structural recursion but instead is programmed by grade-structural recursion , then it is not easily-immediately usable by the 2-resolution-for-transformations , which indeed uses the instantiated 1-resolution-for-morphisms function [solveMorCoMod0] . Therefore oneself primo shall derive the propositional-equations [solveMorCoMod0_rewrite'] corresponding to the definitional-metaconversions of the 1-resolution-for-morphisms function [solveMorCoMod0] .

As always , this COQ program and deduction is mostly-automated !

```
(* len is (S len) *)
- <u>move</u> => F F' f'Sol F'' f_Sol gradeMor_g; <u>destruct</u> f_Sol <u>as</u>
             [ _F (* @'UnitMorCoMod _F *)
             F1 F2 Z1 z1Sol (* ~_1@ F2 o>CoMod z1Sol *)
             | F1 F2 Z2 Z2Sol (* ~ 2 @ F1 o>CoMod z2Sol *)
             L F1 F2 f1Sol f2Sol (* << f1Sol ,CoMod f2Sol >> *) ] .
  (* g is f_ o>CoMod f' , to (f_Sol o>CoMod f'Sol) , is (@'UnitMorCoMod _F
 o>CoMod f'Sol) *)
  * refine (f'Sol)%sol .
  (* g is f o>CoMod f' , to (f Sol o>CoMod f'Sol) , is ( ( ~ 1 @ F2 o>CoMod
 z1Sol ) o>CoMod f'Sol) *)
  * have [:blurb] z1Sol_o_f'Sol :=
                                          f'Sol z1Sol blurb);
      (@solveMorCoMod PolyMorCoMod len
        first by clear - gradeMor g; abstract tac degradeMor gradeMor g .
    refine ( ~_1 @ F2 o>CoMod z1Sol_o_f'Sol )%sol .
  (* g is f_o) o>CoMod f' , to (f_s) o>CoMod f'$01) , is ( ( \sim_s2 @ F1 o>CoMod
 z2Sol ) o>CoMod f'Sol) *)
  * have [:blurb] z2Sol_o_f'Sol :=
      (@solveMorCoMod_PolyMorCoMod len _ _ f'Sol _ z2Sol blurb);
        first by <u>clear</u> - gradeMor_g; abstract tac_degradeMor gradeMor_g .
    refine ( ~_2 @ F1 o>CoMod z2Sol_o_f'Sol )%sol .
  (* g is f_ o>CoMod f' , to (f_Sol o>CoMod f'Sol) , is ( << f1Sol ,CoMod f2Sol
 >> o>CoMod f'Sol ) *)
  * <u>move</u>: (Sol.Destruct_domPair.morCoMod_domPairP f'Sol) => f'Sol_domPairP.
    { destruct f'Sol_domPairP as
          [ F1 F2 (* (@'UnitMorCoMod (Pair F1 F2) )%sol *)
          | F1 F2 Z1 z1 (* ( ~_1 @ F2 o>CoMod z1 )%sol *)
| F1 F2 Z2 z2 (* ( ~_2 @ F1 o>CoMod z2 )%sol *)
| M M' F1 F2 f1 f2 (* ( << f1 ,CoMod f2 >> )%sol *) ] .
      (* g is f_ o>CoMod f' , to (f_Sol o>CoMod f'Sol) , is ( << f1Sol ,CoMod
      f2Sol >> o>CoMod @'UnitMorCoMod (Pair F1 F2) ) *)
      - refine ( << f1Sol ,CoMod f2Sol >> )%sol .
      (* g is f o>CoMod f' , to (f Sol o>CoMod f'Sol) , is ( << f1Sol ,CoMod
      f2Sol >> o>CoMod ~_1 @ F2 o>CoMod z1 *)
      - <u>have</u> [:blurb] f1Sol_o_z1 :=
          (@solveMorCoMod_PolyMorCoMod len _ _ z1 _ f1Sol blurb);
            first by clear - gradeMor_g; abstract tac_degradeMor gradeMor_g .
        refine ( f1Sol_o_z1 )%sol .
      (* g is f_ o>CoMod f' , to (f_Sol o>CoMod f'Sol) , is ( << f1Sol ,CoMod
      f2Sol >> o>CoMod ~_2 @ F1 o>CoMod z2 *)
      - have [:blurb] f2Sol_o_z2 :=
          (@solveMorCoMod_PolyMorCoMod len
                                               z2 _ f2Sol blurb);
            first by <u>clear</u> - gradeMor_g; abstract tac_degradeMor gradeMor_g .
        refine ( f2Sol_o_z2 )%sol .
      (* g is f_ o>CoMod f' , to (f_Sol o>CoMod f'Sol) , is ( << f1Sol ,CoMod
      f2Sol >> o>CoMod << f1 ,CoMod f2 >> *)
      - have [:blurb] f_Sol_o_f1 :=
          first by <u>clear</u> - gradeMor_g; abstract tac_degradeMor gradeMor_g .
        have [:blurb] f_Sol_o_f2 :=
                                               f2 _
          (@solveMorCoMod_PolyMorCoMod len _
                                        ( << f1Sol ,CoMod f2Sol >> %sol) blurb);
            first by <u>clear</u> - gradeMor_g; abstract tac_degradeMor gradeMor_g .
```

```
refine ( << f_Sol_o_f1 ,CoMod f_Sol_o_f2 >> )%sol .
Defined.
Arguments solveMorCoMod_PolyMorCoMod !len _ _ _ !f_Sol _ : simpl nomatch .
Notation "ff_ o>CoMod ff' @ gradeMor_g" :=
                                        ff'
  (@solveMorCoMod_PolyMorCoMod
                                              ff gradeMor g)
    (at level 40 , ff' at next level) : sol scope.
Lemma solveMorCoMod PolyMorCoMod len :
  <u>forall</u> len, <u>forall</u> (F F' : obCoMod) (f'Sol : 'morCoMod(0 F' <math>\sim F) 0 %sol)
            (F'' : obCoMod) (f_Sol : 'morCoMod(0 F'' ~> F' )0 %sol),
      forall gradeMor_g_len : (gradeMor ((Sol.toPolyMor f_Sol)
                                        o>CoMod (Sol.toPolyMor f'Sol)) <= len)%coq_nat,
      forall len', forall gradeMor_g_len' : (gradeMor ((Sol.toPolyMor f_Sol)
                                       o>CoMod (Sol.toPolyMor f'Sol)) <= len')%coq_nat,
           (@solveMorCoMod_PolyMorCoMod len _ _ f'Sol _ f_Sol gradeMor_g_len)
= (@solveMorCoMod_PolyMorCoMod len' _ _ f'Sol _ f_Sol gradeMor_g_len') .
Proof.
  <u>induction</u> len <u>as</u> [ | len ].
  - ( move => ? ? ? ? gradeMor_g_len ); exfalso; clear -gradeMor_g_len;
      by abstract tac_degradeMor gradeMor_g_len.
  - intros. destruct len'.
    + <u>exfalso</u>; <u>clear</u> -gradeMor <u>g</u> len'; by abstract tac degradeMor gradeMor <u>g</u> len'.
    + destruct f_Sol; simpl.
      * reflexivity.
      * erewrite IHlen. reflexivity.
      * <u>erewrite</u> IHlen. reflexivity.
      * { destruct (Sol.Destruct_domPair.morCoMod_domPairP f'Sol); simpl.

    reflexivity.

           - erewrite IHlen. reflexivity.
           - <u>erewrite</u> IHlen. reflexivity.

    <u>erewrite</u> IHlen. <u>rewrite</u> {1}[solveMorCoMod_PolyMorCoMod]lock .

             <u>erewrite</u> IHlen. <u>rewrite</u> -lock. reflexivity.
        }
Qed.
Fixpoint solveMorCoMod_PolyMorCoModP len {struct len} :
  forall (F F' : obCoMod) (f'Sol : 'morCoMod(0 F' ~> F )0 %sol)
    (F'' : obCoMod) (f_Sol : 'morCoMod(0 F'' ~> F' )0 %sol),
  forall gradeMor_g : (gradeMor ((Sol.toPolyMor f_Sol))
                                 o>CoMod (Sol.toPolyMor f'Sol)) <= len)%coq_nat,
                                                            _f'Sol _ f_Sol gradeMor_g))
    (Sol.toPolyMor (@solveMorCoMod_PolyMorCoMod len
      <~~1 (Sol.toPolyMor f_Sol o>CoMod Sol.toPolyMor \overline{f}'Sol) .
Proof.
  <u>case</u> : len => [ | len ].
  (* len is 0 *)
  - ( move => ? ? ? ? gradeMor_g ); exfalso; clear -gradeMor_g;
      by abstract tac_degradeMor gradeMor_g.
  (* len is (S len) *)
  - move => F F' f'Sol F'' f Sol gradeMor q; destruct f Sol as
                F (* @'UnitMorCoMod F *)
                 | F1 F2 Z1 z1Sol (* ~_1 @ F2 o>CoMod z1Sol *)
                 | F1 F2 Z2 z2Sol (* ~_2 @ F1 o>CoMod z2Sol *)
                | L F1 F2 f1Sol f2Sol (* << f1Sol ,CoMod f2Sol >> *) ] .
    (* g is f_o>CoMod\ f' , to (f_Sol\ o>CoMod\ f'Sol) , is (@'UnitMorCoMod \_F
    o>CoMod f'Sol) *)
    * clear; abstract tac_reduce.
    (* g is f_o) o>CoMod f' , to (f_oSol o>CoMod f'Sol) , is ( ( \sim_1 @ F2 o>CoMod
    z1Sol ) o>CoMod f'Sol) *)
    * move: (@solveMorCoMod_PolyMorCoModP len _ _ f'Sol _ z1Sol);
```

```
clear; abstract tac_reduce.
    (* g is f o>CoMod f' , to (f Sol o>CoMod f'Sol) , is ( ( ~ 2 @ F1 o>CoMod
    z2Sol ) o>CoMod f'Sol) *)
    * move: (@solveMorCoMod_PolyMorCoModP len _ _ f'Sol _ z2Sol);
        clear; abstract tac reduce.
    (* g is f_ o>CoMod f' , to (f_Sol o>CoMod f'Sol) , is ( << f1Sol ,CoMod f2Sol
    >> o>CoMod f'Sol ) *)
    * move: (Sol.Destruct domPair.morCoMod domPairP f'Sol) => f'Sol domPairP.
      { destruct f'Sol_domPairP as
             [ F1 F2 (* (@'UnitMorCoMod (Pair F1 F2) )%sol *)
             | F1 F2 Z1 z1 (* ( ~_1 @ F2 o>CoMod z1 )%sol *)
| F1 F2 Z2 z2 (* ( ~_2 @ F1 o>CoMod z2 )%sol *)
             | M M' F1 F2 f1 f2 (* ( << f1 ,CoMod f2 >> )%sol *) ] .
        (* g is f_o>CoMod\ f' , to (f_Sol\ o>CoMod\ f'Sol) , is ( << f1Sol ,CoMod
        f2Sol >> o>CoMod @'UnitMorCoMod (Pair F1 F2) ) *)

    <u>clear</u>; abstract tac_reduce.

         (* g is f_ o>CoMod f' , to (f_Sol o>CoMod f'Sol) , is ( << f1Sol ,CoMod
        f2Sol >> o>CoMod ~_1 @ F2 o>CoMod z1 *)
        - move: (@solveMorCoMod_PolyMorCoModP len _ _ z1 _ f1Sol);
            clear; abstract tac_reduce.
         (* g is f_ o>CoMod f' , to (f_Sol o>CoMod f'Sol) , is ( << f1Sol ,CoMod
        f2Sol >> o>CoMod ~_2 @ F1 o>CoMod z2 *)
         - move: (@solveMorCoMod_PolyMorCoModP len _ _ z2 _ f2Sol);
            clear; abstract tac_reduce.
        (* g is f_ o>CoMod f' , to (f_Sol o>CoMod f'Sol) , is ( << f1Sol ,CoMod
        f2Sol >> o>CoMod << f1 , CoMod f2 >> *)
        - move: (@solveMorCoMod_PolyMorCoModP len _ f1 _ ( << f1Sol ,CoMod f2Sol >> %sol))
                   (@solveMorCoMod_PolyMorCoModP len _ _ f2 ]
                                                   ( << f1Sol ,CoMod f2Sol >> %sol));
            clear; abstract tac_reduce.
Defined.
Fixpoint solveMorCoMod len {struct len} :
  forall (F G : obCoMod) (g : 'morCoMod(0 F <math>\sim G )0 ),
  forall gradeMor_g : (gradeMor g <= len)%coq_nat,</pre>
    'morCoMod(0 F \sim> G)0 %sol.
Proof.
  <u>case</u> : len => [ | len ].
  (* len is 0 *)
  - ( <u>move</u> => F G g gradeMor_g ); <u>exfalso</u>; abstract tac_degradeMor gradeMor_g.
  (* len is (S len) *)
  - \underline{move} => F G g; \underline{case} : F G / g =>
    F' f' F' f (* f o>CoMod f' *)
    | F (* @'UnitMorCoMod F *)
    | F1 F2 Z1 z1 (* ~_1 @ F2 o>CoMod z1 *)
    F1 F2 Z2 z2 (* ~_2 @ F1 o>CoMod z2 *)
    | L F1 F2 f1 f2 (* << f1 ,CoMod f2 >> *)
    ] gradeMor_g .
    (* g is f_ o>CoMod f' *)
    + all: cycle 1.
    (* q is @'UnitMorCoMod F *)
    + <u>refine</u> (@'UnitMorCoMod F)%sol.
    (* g is ~ 1 @ F2 o > CoMod z1 *)
    + have [:blurb] z1Sol := (solveMorCoMod len _ z1 blurb);
```

```
first by clear -gradeMor g; abstract tac degradeMor gradeMor g.
      refine ( ~ 1 @ F2 o>CoMod z1Sol )%sol.
    (* g is ~_2 @ F1 o>CoMod z2 *)
    + have [:blurb] z2Sol := (solveMorCoMod len z2 blurb);
          first by clear -gradeMor_g; abstract tac_degradeMor gradeMor_g.
      refine ( ~_2 @ F1 o>CoMod z2Sol )%sol.
    (* g is << f1 ,CoMod f2 >> *)
    + have [:blurb] f1Sol := (solveMorCoMod len f1 blurb);
        first by <u>clear</u> -gradeMor_g; abstract (tac_degradeMor_gradeMor_g) .
      have [:blurb] f2Sol := (solveMorCoMod len f2 blurb);
        first by clear -gradeMor g; abstract (tac degradeMor gradeMor g) .
      refine ( << f1Sol ,CoMod f2Sol >> )%sol.
    (* g is f o>CoMod f' *)
    + have [:blurb] f_Sol := (solveMorCoMod len _ _ f_ blurb);
        first by clear -gradeMor_g; abstract tac_degradeMor gradeMor_g.
      have [:blurb] f'Sol := (solveMorCoMod len _ _ f' blurb);
        first by <u>clear</u> -gradeMor g; abstract tac degradeMor gradeMor g.
      have [:blurb] f Sol o f'Sol :=
         (@solveMorCoMod PolyMorCoMod (gradeMor ((Sol.toPolyMor f Sol)
                           o>CoMod (Sol.toPolyMor f'Sol))) _ _ f'Sol _ f_Sol blurb);
           first by <u>clear</u>; abstract reflexivity.
      refine ( f_Sol_o_f'Sol )%sol.
Defined.
Arguments solveMorCoMod !len _ _ !g _ : clear implicits , simpl nomatch .
Lemma solveMorCoMod_len :
  forall len, forall (F G : obCoMod) (g : 'morCoMod(0 F <math>\sim G) (0)),
      forall gradeMor_g_len : (gradeMor g <= len)%coq_nat,</pre>
      forall len', forall gradeMor_g_len' : (gradeMor g <= len')%coq_nat,</pre>
           (@solveMorCoMod len _ _ g gradeMor_g_len)
= (@solveMorCoMod len' _ _ g gradeMor_g_len') .
Proof.
  induction len as [ | len ].
    ( move => ? ? gradeMor_g_len ); exfalso; clear -gradeMor_g_len;
      by abstract tac_degradeMor gradeMor_g_len.
  - <u>intros</u>. <u>destruct</u> len'.
    + <u>exfalso</u>; <u>clear</u> -gradeMor_g_len'; by abstract tac_degradeMor gradeMor_g_len'.
    + <u>destruct</u> g; <u>cbn</u> -[solveMorCoMod_PolyMorCoMod]; cycle 1.
      * reflexivity.
      * <u>erewrite</u> IHlen. reflexivity.
      * <u>erewrite</u> IHlen. reflexivity.
      * erewrite IHlen. rewrite {1}[solveMorCoMod]lock . erewrite IHlen.
        <u>rewrite</u> -lock. reflexivity.
      * <u>erewrite</u> IHlen. <u>rewrite</u> {1}[solveMorCoMod]lock . <u>erewrite</u> IHlen.
        <u>rewrite</u> -lock. reflexivity.
Qed.
Fixpoint solveMorCoModP len {struct len} :
  <u>forall</u> (F G : obCoMod) (g : 'morCoMod(0 F <math>\sim G )0 ),
  forall gradeMor_g : (gradeMor g <= len)%coq_nat,</pre>
    (Sol.toPolyMor (@solveMorCoMod len _ _ g gradeMor_g)) <~~1 g.</pre>
Proof.
  <u>case</u> : len => [ | len ].
```

```
(* len is 0 *)
  - ( <u>move</u> => F G g gradeMor g ); <u>exfalso</u>; abstract tac degradeMor gradeMor g.
  (* len is (S len) *)
  - <u>move</u> => F G g; <u>case</u> : F G / g =>
[ F F' f' F'' f_ (* f_ o>CoMod f' *)
      F (* @'UnitMorCoMod F *)
    | F1 F2 Z1 z1 (* ~_1 @ F2 o>CoMod z1 *)
| F1 F2 Z2 z2 (* ~_2 @ F1 o>CoMod z2 *)
    | L F1 F2 f1 f2 (* << f1 ,CoMod f2 >> *)
    ] gradeMor_g .
    (* g is f_ o>CoMod f' *)
    + all: cycle 1.
    (* g is @'UnitMorCoMod F *)
    + clear; abstract tac_reduce.
    (* g is ~ 1 @ F2 o > CoMod z1 *)
    + <u>move</u>: (@solveMorCoModP len _ _ z1); <u>clear</u>; abstract tac_reduce.
    (* g is ~ 2 @ F1 o>CoMod z2 *)
    + <u>move</u>: (@solveMorCoModP len _ _ z2); <u>clear</u>; abstract tac_reduce.
    (* g is << f1 ,CoMod f2 >> *)
    + move: (@solveMorCoModP len _ _ f1) (@solveMorCoModP len _ _ f2);
        clear; abstract tac_reduce.
    (* g is f_ o>CoMod f' *)
    + move: (@solveMorCoModP len _ _ f_) (@solveMorCoModP len _ _ f') .
      move: solveMorCoMod_PolyMorCoModP.
      clear; abstract tac_reduce.
Qed.
Definition solveMorCoMod PolyMorCoMod0 :
  forall (F F' : obCoMod) (f'Sol : 'morCoMod(0 F' ~> F )0 %sol)
    (F'' : obCoMod) (f_Sol : 'morCoMod(0 F'' ~> F' )0 %sol),
    'morCoMod(0 F'' \sim> F)0 %sol.
Proof.
  intros; apply: (@solveMorCoMod PolyMorCoMod (gradeMor ((Sol.toPolyMor f Sol)
                                         o>CoMod (Sol.toPolyMor f'Sol)))); constructor.
Defined.
Notation "ff o>CoMod ff'" :=
  (@solveMorCoMod_PolyMorCoMod0 _
                                      ff'
    (at level 40 , ff' at next level) : sol_scope.
Lemma solveMorCoMod PolyMorCoMod0 len :
  forall (F F' : obCoMod) (f'Sol : 'morCoMod(0 F' ~> F )0 %sol)
    (F'' : obCoMod) (f\_Sol : 'morCoMod(0 F'' \sim> F')0 %sol),
  forall len', forall gradeMor_g_len' : (gradeMor ((Sol.toPolyMor f_Sol)
                                        o>CoMod (Sol.toPolyMor f'Sol)) <= len')%coq_nat,
      (solveMorCoMod_PolyMorCoMod0 f'Sol f_Sol)
      = (@solveMorCoMod_PolyMorCoMod len' _ _ f'Sol _ f_Sol gradeMor_g_len') .
Proof. intros. apply: solveMorCoMod_PolyMorCoMod_len . Qed.
Lemma solveMorCoMod PolyMorCoMod0 UnitMorCoMod:
  (\underline{forall} (F F' : obCoMod) (f'Sol : 'morCoMod(0 F' <math>\sim F)0),
       (@'UnitMorCoMod F') o>CoMod f'Sol = f'Sol )%sol.
Proof. <u>intros</u>. reflexivity. Qed.
  mma solveMorCoMod_PolyMorCoMod0___Project1_Mor :
( forall (F Z1 : obCoMod) (f'Sol : 'morCoMod(0 Z1 ~> F )0),
Lemma solveMorCoMod PolyMorCoMod0
      <u>forall</u> (F1 F2 : obCoMod) (z1Sol : 'morCoMod(0 F1 ~> Z1 )0),
         ( ~_1 @ F2 o>CoMod z1Sol) o>CoMod f'Sol
         = \sim 1 @ F2 o>CoMod (z1Sol o>CoMod f'Sol) )%sol .
```

```
Proof.
  intros. rewrite [solveMorCoMod_PolyMorCoMod0 in LHS]lock.
  erewrite solveMorCoMod PolyMorCoMod0 len.
  <u>rewrite</u> -lock /solveMorCoMod PolyMorCoMod0. reflexivity.
Qed.
Lemma solveMorCoMod PolyMorCoMod0
                                       Project2 Mor :
  ( \underline{forall} (F Z2 : obCoMod) (f'Sol : 'morCoMod(0 Z2 \sim> F )0),
      forall (F1 F2 : obCoMod) (z2Sol : 'morCoMod(0 F2 ~> Z2 )0),
         ( ~ 2 @ F1 o>CoMod z2Sol) o>CoMod f'Sol
         = \sim 2 @ F1 o>CoMod (z2Sol o>CoMod f'Sol) )%sol .
Proof.
  <u>intros</u>. <u>rewrite</u> [solveMorCoMod_PolyMorCoMod0 <u>in</u> LHS]lock.
  erewrite solveMorCoMod_PolyMorCoMod0_len.
  <u>rewrite</u> -lock /solveMorCoMod PolyMorCoMod0. reflexivity.
Qed.
Lemma solveMorCoMod_PolyMorCoMod0_UnitMorCoMod_Pairing_Mor :
  (\underline{forall} (L F1 : obCoMod) (f1Sol : 'morCoMod(0 L <math>\sim F1 )0) (F2 : obCoMod)
      (f2Sol : 'morCoMod(0 L ~> F2 )0),
      << f1Sol ,CoMod f2Sol >> o>CoMod (@'UnitMorCoMod (Pair F1 F2))
                 = << f1Sol ,CoMod f2Sol >> )%sol .
Proof. <u>intros</u>. reflexivity. Qed.
Lemma solveMorCoMod_PolyMorCoMod0_Project1_Mor_Pairing_Mor :
  ( forall (L F2 : obCoMod) (f2Sol : 'morCoMod(0 L ~> F2 )0) (F1 : obCoMod)
    (f1Sol : morCoMod(0 L \sim F1)0) (Z1 : obCoMod) (Z1 : morCoMod(0 F1 \sim F1)0),
      << f1Sol ,CoMod f2Sol >> o>CoMod (~_1 o>CoMod z1)
                 = f1Sol o>CoMod z1 )%sol .
Proof.
  <u>intros</u>. <u>rewrite</u> [solveMorCoMod_PolyMorCoMod0 <u>in</u> LHS]lock.
  erewrite solveMorCoMod PolyMorCoMod0 len.
  rewrite -lock /solveMorCoMod_PolyMorCoMod0. reflexivity.
Lemma solveMorCoMod PolyMorCoMod0_Project2_Mor_Pairing_Mor :
  ( forall (L F1 : obCoMod) (f1Sol : 'morCoMod(0 L ~> F1 )0) (F2 : obCoMod)
    (f2Sol : 'morCoMod(0 L ~> F2 )0) (Z2 : obCoMod) (z2 : 'morCoMod(0 F2 ~> Z2 )0),
      << f1Sol ,CoMod f2Sol >> o>CoMod (~_2 o>CoMod z2)
                 = f2Sol o>CoMod z2 )%sol .
Proof.
  <u>intros</u>. <u>rewrite</u> [solveMorCoMod_PolyMorCoMod0 <u>in</u> LHS]lock.
  erewrite solveMorCoMod_PolyMorCoMod0_len.
  rewrite -lock /solveMorCoMod_PolyMorCoMod0. reflexivity.
Qed.
Lemma solveMorCoMod_PolyMorCoMod0_Pairing_Mor_Pairing_Mor :
  ( <u>forall</u> (L M : obCoMod) (f1Sol : 'morCoMod(0 L ~> M )0) (M' : obCoMod)
      (f2Sol : 'morCoMod(0 L \sim> M')0) (F1 F2 : obCoMod)
      (f1 : 'morCoMod(0 Pair M M' ~> F1 )0) (f2 : 'morCoMod(0 Pair M M' ~> F2 )0),
      << f1Sol ,CoMod f2Sol >> o>CoMod ( << f1 ,CoMod f2 >> )
                 = ( << ( << f1Sol ,CoMod f2Sol >> o>CoMod f1 )
                      ,CoMod ( << f1Sol ,CoMod f2Sol >> o>CoMod f2 ) >> ) %sol .
Proof.
  <u>intros</u>. <u>rewrite</u> [solveMorCoMod PolyMorCoMod0 <u>in</u> LHS]lock.
  do 2 <u>erewrite</u> solveMorCoMod_PolyMorCoMod0 len.
  rewrite -lock /solveMorCoMod_PolyMorCoMod0. reflexivity.
Qed.
Lemma solveMorCoMod PolyMorCoMod0P :
  forall (F F' : obCoMod) (f'Sol : 'morCoMod(0 F' ~> F )0 %sol)
    (F'': obCoMod) (f\_Sol: 'morCoMod(0 F'' \sim> F')0 %sol), (Sol.toPolyMor(@solveMorCoMod_PolyMorCoMod0 _ _ f'Sol_
      <~~1 (Sol.toPolyMor f_Sol o>CoMod Sol.toPolyMor f'Sol) .
Proof. intros; apply: solveMorCoMod_PolyMorCoModP. Qed.
Definition solveMorCoMod0 :
```

```
forall (F G : obCoMod) (g : 'morCoMod(0 F \sim> G)0),
    'morCoMod(0 F \sim> G)0 %sol.
Proof.
  intros; apply: (@solveMorCoMod (gradeMor g)); constructor.
Defined.
Lemma solveMorCoMod0 len :
  <u>forall</u> (F G : obCoMod) (g : 'morCoMod(0 F \sim G )0 ),
  forall len, forall gradeMor_g : (gradeMor g <= len)%coq_nat,</pre>
      (solveMorCoMod0 g) = (@solveMorCoMod len _ _ g gradeMor_g).
Proof. <u>intros</u>. <u>apply</u>: solveMorCoMod len . Qed.
Lemma solveMorCoMod0 UnitMorCoMod:
  forall (F : obCoMod)
    solveMorCoMod0 (@'UnitMorCoMod F) = (@'UnitMorCoMod F)%sol.
Proof. <u>intros</u>. reflexivity. Qed.
Lemma solveMorCoMod0_Project1_Mor :
  forall (F1 F2 Z1 : obCoMod) (z1 : 'morCoMod(0 F1 ~> Z1 )0),
    solveMorCoMod0 ( ~_1 @ F2 o>CoMod z1)
    = ( \sim_1 @ F2 o \sim CoMod (solveMorCoMod0 z1) )%sol.
Proof.
  <u>intros</u>. <u>rewrite</u> [solveMorCoMod0 <u>in</u> LHS]lock. <u>erewrite</u> solveMorCoMod0_len.
  rewrite -lock /solveMorCoMod0. reflexivity.
Lemma solveMorCoMod0_Project2_Mor :
  <u>forall</u> (F1 F2 Z2 : obCoMod) (z2 : 'morCoMod(0 F2 ~> Z2 )0),
    solveMorCoMod0 ( ~_2 @ F1 o>CoMod z2)
    = ( \sim_2 @ F1 o>CoMod (solveMorCoMod0 z2) )%sol.
  intros. rewrite [solveMorCoMod0 in LHS]lock. erewrite solveMorCoMod0 len.
  rewrite -lock /solveMorCoMod0. reflexivity.
Qed.
Lemma solveMorCoMod0_Pairing_Mor :
 forall (L F1 F2 : obCoMod) (f1 : 'morCoMod(0 L \sim> F1 )0) (f2 : 'morCoMod(0 L \sim> F2 )0),
    solveMorCoMod0 ( << f1 ,CoMod f2 >> )
    = ( << solveMorCoMod0 f1 ,CoMod solveMorCoMod0 f2 >> ) %sol.
Proof.
  <u>intros</u>. <u>rewrite</u> [solveMorCoMod0 <u>in</u> LHS]lock.
  do 2 <u>erewrite</u> solveMorCoMod0_len.
  rewrite -lock /solveMorCoMod0. reflexivity.
Qed.
Lemma solveMorCoMod0_PolyMorCoMod :
  \underline{forall} (FF': obCoMod) (f': 'morCoMod(0 F' \sim F )0)
    (F'': obCoMod) (f_: 'morCoMod(0 F'' \sim> F')0),
    solveMorCoMod0 ( f_ o>CoMod f' )%poly
    = ( (solveMorCoMod0 f_) o>CoMod (solveMorCoMod0 f') )%sol.
Proof.
  <u>intros</u>. <u>rewrite</u> [solveMorCoMod0 <u>in</u> LHS]lock.
  do 2 <u>erewrite</u> solveMorCoMod0 len. <u>erewrite</u> solveMorCoMod PolyMorCoMod0 len.
  rewrite -lock /solveMorCoMod0. reflexivity.
Qed.
Lemma solveMorCoModOP :
  forall (F G : obCoMod) (g : 'morCoMod(0 F <math>\sim S G) (0)),
                                        _ _ g)) <~~1 g.
    (Sol.toPolyMor (@solveMorCoMod0
Proof. intros. apply: solveMorCoModP . Qed.
Definition solveMorCoMod_PolyMorCoMod0_rewrite :=
  (solveMorCoMod_PolyMorCoMod0___UnitMorCoMod,
   solveMorCoMod_PolyMorCoMod0___Project1_Mor,
   solveMorCoMod_PolyMorCoMod0___Project2_Mor,
```

```
solveMorCoMod PolyMorCoMod0 UnitMorCoMod Pairing Mor,
   solveMorCoMod PolyMorCoMod0 Project1_Mor_Pairing_Mor,
   solveMorCoMod PolyMorCoMod0 Project2 Mor Pairing Mor,
   solveMorCoMod PolyMorCoMod0 Pairing Mor Pairing Mor).
Definition solveMorCoMod0 rewrite :=
  ( ( solveMorCoMod0 PolyMorCoMod, solveMorCoMod PolyMorCoMod0 rewrite ) ,
    solveMorCoMod0_UnitMorCoMod, solveMorCoMod0_Project1_Mor,
    solveMorCoMod0_Project2_Mor, solveMorCoMod0_Pairing_Mor ).
(*TODO: NOT USED , COMMENT , derivable immediately without the finished
cut-elimination lemma *)
Lemma solveMorCoMod PolyMorCoMod0 Pairing Mor:
  ( forall (L : obCoMod) (F1 F2 : obCoMod)
      (f1 : 'morCoMod(0 L \sim F1)0) (f2 : 'morCoMod(0 L \sim F2)0),
      forall (L' : obCoMod) (l : 'morCoMod(0 L' ~> L )0),
        l o>CoMod ( << f1 ,CoMod f2 >> )
        = << ( l o>CoMod f1 ) ,CoMod ( l o>CoMod f2 ) >> )%sol .
Abort.
**)
Definition solveMorCoMod0_toPolyMor :
  \underline{forall} (F G : obCoMod) (g : 'morCoMod(0 F <math>\sim G) 0 %sol),
    solveMorCoMod0 (Sol.toPolyMor g) = g .
  induction g; simpl in *; rewrite ?solveMorCoMod0 rewrite; intros;
    repeat match goal with
           | [ Hred : ( solveMorCoMod0 _ = _ ) |- _ ] =>
             rewrite Hred; clear Hred
           end; reflexivity.
Qed.
Definition solveMorCoMod0_rewrite' :=
  ( ( solveMorCoMod0_PolyMorCoMod, solveMorCoMod_PolyMorCoMod0_rewrite ) ,
    solveMorCoMod0_UnitMorCoMod, solveMorCoMod0_Project1_Mor,
    solveMorCoMod0_Project2_Mor, solveMorCoMod0_Pairing_Mor,
    (** solveMorCoMod_PolyMorCoMod0_Pairing_Mor, **) solveMorCoMod0 toPolyMor ) .
Ltac tac_reduce_solveMorCoMod0 :=
  rewrite /= ?solveMorCoMod0 rewrite' /= ;
  <u>intuition</u> (<u>subst</u>; repeat <u>match</u> goal <u>with</u>
                            | [ Hred : ( solveMorCoMod0 _ = _ ) |- _ ] =>
                              <u>rewrite</u> Hred
                            <u>end</u>;
             rewrite /= ?solveMorCoMod0 rewrite' /= ;
             try congruence;
             eauto).
End Resolve.
```

6 Grammatical presentation of transformations

Now the inductive-family-presentation [transfCoMod] has some additional/embedded type-indexes and type-constructor : type-indexes to represent the domain-codomain-morphisms and type-constructor [UnitTransfCoMod] to represent the unit-transformation .

Each decoding ("Yoneda") of the whatever-is-interesting arrows between the indexes-for-touched-morphisms are metatransformations which are programmed as some grammatical-constructors of the inductive-family-presentations [morCoMod] and [transfCoMod] .

Memo that the functoriality ("arrows-action") of each metafunctor (decoded index-for-touched-morphisms) and the naturality ("arrows-action") of each metatransformation (decoded arrow-between-indexes) is signified via the additional/embedded type-indexes of [transfCoMod] and type-constructor [UnitTransfCoMod] of [transfCoMod] . All this is effected via the two conversion relations [convMorCoMod] [convTransfCoMod] which relate those grammatical-touched-morphisms : [convMorCoMod] is for morphisms and [convTransfCoMod] is for transformations .

Memo that here the (multiplicative) outer/material ("horizontal") composition [PolyMorCoMod] [TransfCoMod_PolyMorCoMod_Pre] [TransfCoMod_PolyMorCoMod_Post] is some common operation , but there is also some uncommon operation [PolyTransfCoMod] which is the (coordinatewise/dimensional/pointwise) inner/structural ("vertical") composition of transformation-later-transformation (along some tight/strict or lax « cut-adherence ») inside each enrichment/indexer-graph ; and both compositions cut-constructors shall be eliminated/erased

Attention: the formulation of the inner/structural composition [PolyTransfCoMod] cut-constructor must relate the codomain-morphism of the prefix-input transformation in-relation-with the domain-morphism of the postfix-input transformation, by propositional-equality. In other words, this [PolyTransfCoMod] constructor have some extra argument/parameter, named « cut-adherence/adhesive », which behold this propositional-equality. Such formulation is because of this fact: the 2-conversion-for-transformations do convert across two transformations whose domain-codomain-morphisms-computation arguments are not syntactically/grammatically-the-same, and therefore, during the recursion step which is the 2-resolution of the inner/structural composition [PolyTransfCoMod] cut-constructor, oneself lacks to know that the codomain-morphism of the recursively 2-resolved prefix-output transformation is grammatically-same (or propositionally-equal ...) as the domain-morphism of the recursively 2-resolved postfix-output transformation; the presence of the cut-adherence in the input will infer some adherence in the output.

During the programming of the (inner/structural) « resolved cut » [solveTransfCoMod_PolyTransfCoMod], which ingets two morphisms together with some adhesive, the precedence is such that primo each of the two input morphisms is destructed, then secondo the input adhesive is inverted such to exfalso/exclude/impossible the preceding/outer case or to obtain some subadhesive for the recursive call.

6.1 Inversion of the cut-adherence (here propositional-equality)

For the material mathematics , the decidable equality [obCoMod_eq] on the objects enables to do any logical-inversion of the cut-adherence , which here is some very-dependently-typed propositional-equality-across-any-two-morphisms . In the alternative formulation where the cut-adherence is the more-lax (instead of tight/strict) polymorphism-conversion , the lemma [convMorCoMod_toPolyMorP'] would be some confluence lemma-corollary , instead of simply being the inversion lemma below .

```
Module Transf.
Module EqMorCoMod.
Definition convMorCoMod (F G : obCoMod) (g g' : 'morCoMod(0 F \sim G )0 %poly)
: Prop := @eq ( 'morCoMod(0 F ~> G )0 ) g' g .
Module Export Ex_Notations.
  Notation "g' < >1 g" := (@convMorCoMod _ _ g g') (at level 70): poly_scope .
End Ex Notations.
Lemma convMorCoMod_eq : forall (F G : obCoMod) (g g' : 'morCoMod(0 F ~> G )0 %poly),
    g' < >1 g -> g' = g.
Proof. trivial. Qed.
Lemma eq_convMorCoMod : forall (F G : obCoMod) (g g' : 'morCoMod(0 F ~> G )0 %poly),
    g' = g -> g' <\sim>1 g.
Proof. trivial. Qed.
Lemma convMorCoMod_sym : forall (F G : obCoMod) (g g' : 'morCoMod(0 F ~> G )0 %poly),
    g' <\sim>1 g -> g <\sim>1 g'.
Proof. symmetry. trivial. Qed.
Module Inversion_Project1.
Lemma convMorCoMod Project1P'
: <u>forall</u> (F1 F2 : obCoMod) (Z1 : obCoMod) (z1 z1' : 'morCoMod(0 F1 ~> Z1 )0),
    ((\sim 1 @ F2 o) CoMod z1) <\sim 1 (\sim 1 @ F2 o) CoMod z1')) -> (z1 <\sim 1 z1').
Proof.
  <u>intros</u> <u>until</u> z1'. <u>intros</u> H. <u>inversion</u> H.
```

```
match goal with
  | [ H1 : existT
                          = existT
                                           do 2 <u>apply</u> (ObCoMod_eq.Eqdep_dec_inj_pair2_eq_dec) <u>in</u> H1
  <u>end</u>; <u>subst</u>; <u>constructor</u>.
Qed.
End Inversion Project1.
Module Inversion Project2.
Lemma convMorCoMod_Project2P'
: <u>forall</u> (F1 F2 : obCoMod) (Z2 : obCoMod) (Z2 Z2' : 'morCoMod(0 F2 ~> Z2 )0),
    ( ( ~ 2 @ F1 o>CoMod z2 ) <~>1 ( ~ 2 @ F1 o>CoMod z2' ) ) -> (z2 <~>1 z2') .
Proof.
  <u>intros</u> <u>until</u> z2'. <u>intros</u> H. <u>inversion</u> H.
  match goal with
                          = existT
  | [ Hl : existT
    do 2 <u>apply</u> (ObCoMod eq.Eqdep dec inj pair2 eq dec) <u>in</u> H1
  end; subst; constructor.
Qed.
End Inversion_Project2.
Module Inversion_Pairing.
Lemma convMorCoMod_PairingP'
: <u>forall</u> (L : obCoMod) (F1 F2 : obCoMod) (f1 : 'morCoMod(0 L ~> F1 )0)
    (f2 : 'morCoMod(0 L ~> F2 )0),
    forall (g1 : morCoMod(0 L \sim F1)0) (g2 : morCoMod(0 L \sim F2)0),
       ( << f1 ,CoMod f2 >> <\sim>1 <math><< g1 ,CoMod g2 >> )
       -> (f1 <~>1 g1) /\ (f2 <~>1 g2).
Proof.
  <u>intros</u> <u>until</u> g2. <u>intros</u> H. <u>inversion</u> H.
  do 2 match goal with
        | [ H1 : existT
                               _ = existT <sub>_</sub>
                                                . |- .
          do 2 <u>apply</u> (ObCoMod_eq.Eqdep_dec_inj_pair2_eq_dec) <u>in</u> H1
       end; subst; split; constructor.
Qed.
End Inversion_Pairing.
Module Inversion_Exfalso.
Lemma convMorCoMod_ExfalsoP_Project1_Project2
: forall (F1 F2 : obCoMod) (Z1 : obCoMod) (Z1 : 'morCoMod(0 F1 \sim> Z1 )0)
    (z2 : 'morCoMod(0 F2 ~> Z1 )0),
     ( ( \sim 1 @ F2 o>CoMod z1 ) <\sim>1 ( \sim 2 @ F1 o>CoMod z2 ) ) -> False .
Proof. intros until z2. intros H. inversion H. Qed.
Lemma convMorCoMod_ExfalsoP_UnitMorCoMod_Project1
  : forall (F1 F2 : obCoMod) (z1 : 'morCoMod(0 F1 ~> Pair F1 F2 )0),
    ( @'UnitMorCoMod (Pair F1 F2 ) ) <~>1 ( ~ 1 @ F2 o>CoMod z1 ) -> <u>False</u>.
Proof. intros until z1. intros H. inversion H. Qed.
Lemma convMorCoMod ExfalsoP Project1 Pairing
  : <u>forall</u> (F1 F2 : obCoMod) (Z1 Z1' : obCoMod) (z1 : 'morCoMod(0 F1 ~> Pair Z1 Z1' )0)
    (f1 : 'morCoMod(0 Pair F1 F2 ~> Z1 )0) (f2 : 'morCoMod(0 Pair F1 F2 ~> Z1' )0),
    ( ( \sim 1 @ F2 o > CoMod z1 ) < \sim 1 ( << f1 , CoMod f2 >> ) ) -> False .
Proof. <u>intros</u> <u>until</u> f2. <u>intros</u> H. <u>inversion</u> H. Qed.
Lemma convMorCoMod ExfalsoP Project2 Pairing
  : <u>forall</u> (F1 F2 : obCoMod) (Z2 Z2' : obCoMod) (z2 : 'morCoMod(0 F2 ~> Pair Z2 Z2' )0)
    (f1 : 'morCoMod(0 Pair F1 F2 ~> Z2 )0) (f2 : 'morCoMod(0 Pair F1 F2 ~> Z2' )0),
    ( ( \sim_2 @ F1 o > CoMod z2 ) <\sim>1 ( << f1 , CoMod f2 >> ) ) -> False .
Proof. <u>intros</u> <u>until</u> f2. <u>intros</u> H. <u>inversion</u> H. Qed.
End Inversion_Exfalso.
```

```
Module Inversion toPolyMor.
Import Sol.Ex Notations.
Local Definition Sol convMorCoMod (F G : obCoMod)
               (g g' : 'morCoMod(0 F ~> G )0 %sol) : Prop :=
     Qeq ( 'morCoMod(0 F \sim> G )0 %sol ) g' g .
Local Notation "g' \sim 1 g'' := (@Sol_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_convMorCoMod_co
                                                                              (<u>at</u> level 70): sol scope.
Lemma convMorCoMod toPolyMorP' :
     <u>forall</u> (F G : obCoMod) (g' g : 'morCoMod(0 F <math>\sim G) 0 %sol),
          ( Sol.toPolyMor q' <~>1 Sol.toPolyMor q )%poly -> ( q' ~~~1 q )%sol .
     <u>induction</u> g' <u>as</u> [ | ? ? ? ? IHg' | ? ? ? ? IHg' | ? ? ? ? IHg'1 ? IHg'2 ];
          simpl; intros g H.

    { move: (Sol.Inversion_domEqcod.morCoMod_domEqcodP g).

              unfold Sol.Inversion_domEqcod.morCoMod_domEqcodP_Type.
               rewrite (ObCoMod_eq.obCoMod_eqP). intros g_morCoMod_domEqcodP.
              <u>destruct</u> g morCoMod domEqcodP.
               - reflexivity.

    <u>exfalso</u>; <u>inversion</u> H.

             (* apply (Inversion_Exfalso.convMorCoMod_ExfalsoP_UnitMorCoMod_Project1 H). *)

    <u>exfalso</u>; <u>inversion</u> H.

    exfalso; inversion H.

     - { <u>destruct</u> (Sol.Destruct domPair.morCoMod domPairP g); <u>simpl</u> in * .
               - <u>exfalso</u>; <u>inversion</u> H.
               - move: (Inversion_Project1.convMorCoMod_Project1P' H);
                        move /IHg' -> ; reflexivity.
               - <u>exfalso</u>; <u>inversion</u> H.
               - exfalso; inversion H.
        { <u>destruct</u> (Sol.Destruct domPair.morCoMod domPairP g); <u>simpl</u> <u>in</u> * .

    <u>exfalso</u>; <u>inversion</u> H.

               - <u>exfalso</u>; <u>inversion</u> H.
               - move: (Inversion_Project2.convMorCoMod_Project2P' H);
                        move /IHg' -> ; reflexivity.

    exfalso; inversion H.

     - { destruct (Sol.Destruct codPair.morCoMod codPairP g); simpl in * .
              - <u>exfalso</u>; <u>inversion</u> H.
               - <u>exfalso</u>; <u>inversion</u> H.
               - <u>exfalso</u>; <u>inversion</u> H.
               - move: (Inversion_Pairing.convMorCoMod_PairingP' H)
                   => [ ] /IHg'1 -> /IHg'2 -> ; reflexivity.
Qed.
End Inversion_toPolyMor.
End EgMorCoMod.
```

6.2 Outer ("horizontal") left-whisk cut , outer ("horizontal") right-whisk cut , and inner ("vertical") composition cut with cut-adhesive

The common operations on the touched-morphisms are multifold/multiplicative; but because of the generating-views subindexer, then oneself can avoid the (multiplicative) outer/material ("horizontal") composition of transformation-next-transformation (whose output multiplicity is outside the subindexer) and instead it is sufficient to describe the (multiplicative) outer/material composition of transformation-next-morphism ("right-whisk") and the (multiplicative) outer/material composition of morphism-next-transformation ("left-whisk") (whose output multiplicity is the shape $\{0 \sim > 1\}$ inside the subindexer) .

Moreover, there is also inner/(structure-logical) ("vertical") composition of transformation-later-transformation inside each enrichment/indexer-graph , whose formulation has some cut-adhesive parameter between the codomain of the prefix-transformation and the domain of the infix-transformation .

```
Import EqMorCoMod.Ex_Notations.
Reserved Notation "''transfCoMod' (0 g \sim g' )0"
          (at level 0, format "''transfCoMod' (0 g \sim g' )0").
(** here there are none parameter/customized-arrow action ( non-structural
reindexing , customized boundary-or-degeneracy ) ; is it possible to make sense of
such ? *)
Inductive transfCoMod : forall (F G : obCoMod), morCoMod F G -> morCoMod F G -> Type :=
| PolyTransfCoMod : forall (F G : obCoMod), forall (g g' : 'morCoMod(0 F <math>\sim G )0),
      \underline{\text{forall}} (g'g: 'transfCoMod(0 g' \sim g )0), \underline{\text{forall}} (g'\theta g'': 'morCoMod(0 F \sim G )0 ),
           forall (g''g' : 'transfCoMod(0 g'' ~> g'0 )0),
           forall (eqMor : g'0 <~>1 g' ) (** cut-adherence **) ,
              'transfCoMod(0 g'' ~> g )0
| TransfCoMod PolyMorCoMod Pre :
 \underline{forall} (F G : obCoMod), \underline{forall} (g g' : 'morCoMod(0 F \sim G )0),
 \underline{\textbf{forall}} \ (\textbf{\textit{g'g}} : \ \texttt{'transfCoMod}(0 \ \texttt{g'} \ \texttt{\sim} \ \texttt{g} \ )0), \ \underline{\textbf{forall}} \ (\textbf{\textit{E}} : \ \texttt{obCoMod}) \ (\textbf{\textit{f}} : \ \texttt{'morCoMod}(0 \ \texttt{E} \ \texttt{\sim} \ \texttt{F} \ )0),
      'transfCoMod(0 f o>CoMod g' ~> f o>CoMod g )0
| TransfCoMod PolyMorCoMod Post :
 forall (G H : obCoMod) (h : 'morCoMod(0 G ~> H )0),
 <u>forall</u> (F: obCoMod) (g g': 'morCoMod(0 F \sim G )0) (g'g: 'transfCoMod(0 g' \sim > g )0),
   'transfCoMod(0 g' o>CoMod h ~> g o>CoMod h )0
(** unit-transformation-on-each-morphism **)
| UnitTransfCoMod : forall (F G : obCoMod), forall (g : 'morCoMod(0 F ~> G )0),
       'transfCoMod(0 g (*memo: same g *) \sim g (*memo: same g *) )0
| TransfCoMod Gen : forall (F G : obCoMod Gen),
    forall (g g' : morCoMod_Gen F G), transfCoMod_Gen g g' ->
       'transfCoMod(0 (MorCoMod_Gen g) ~> (MorCoMod_Gen g') )0
**)
| Project1_Transf : forall (F1 F2 : obCoMod) (Z1 : obCoMod),
    <u>forall</u> (z1 z1': 'morCoMod(0 F1 \sim Z1 )0) (z1z1': 'transfCoMod(0 z1 \sim z1' )0),
       'transfCoMod(0 \sim 1 @ F2 o>CoMod z1 \sim> \sim 1 @ F2 o>CoMod z1' )0
| Project2 Transf : forall (F1 F2 : obCoMod) (Z2 : obCoMod),
    <u>forall</u> (z2 z2' : 'morCoMod(0 F2 ~> Z2 )0) (z2z2' : 'transfCoMod(0 z2 ~> z2' )0),
       'transfCoMod(0 \sim 2 @ F1 o>CoMod z2 \sim> \sim 2 @ F1 o>CoMod z2' )0
| Pairing_Transf : forall (L : obCoMod) (F1 F2 : obCoMod),
    transfCoMod(0 \ll f1 ,CoMod f2 >> ~> \ll f1' ,CoMod f2' >> )0
where "''transfCoMod' (0 g \sim g' )0" := (@transfCoMod _ _ g g') : poly_scope.
Notation "g''g' o^CoMod g'g # egMor" :=
  (@PolyTransfCoMod _ _ _ g'g _ g''g' eqMor)
    (<u>at</u> level 40 , g'g <u>at</u> next level) : poly_scope.
Notation "f o > CoMod^ g'g" :=
    TransfCoMod_PolyMorCoMod_Pre _ _ _ g'g _ f)
(<u>at</u> level 40 , g'g <u>at</u> next level) : poly_scope.
  (@TransfCoMod PolyMorCoMod Pre
Notation "q'q ^o>CoMod h" :=
  (@TransfCoMod_PolyMorCoMod_Post _ _ h _ _ g'g)
```

```
(at level 40 , h at next level) : poly scope.
Notation "@ ''UnitTransfCoMod' g" :=
  (@UnitTransfCoMod _ _ g) (<u>at</u> level 10, only parsing) : poly_scope.
Notation "''UnitTransfCoMod'" :=
  (@UnitTransfCoMod _ _ _) (<u>at</u> level 0) : poly_scope.
(** Notation "''TransfCoMod Gen' ff" :=
       (@TransfCoMod_Gen _ _ _ ff) (at level 3) : poly_scope. **)
(* @ in ~ 1 @ says argument *)
Notation "\sim \overline{1} @ F2 o>CoMod^ z1z1'" :=
  (@Project1_Transf _ F2 _ _ z1z1') (<u>at</u> level 4, F2 <u>at</u> next level) : poly_scope.
Notation "~ 1 o>CoMod^ z1z1'" :=
  (@Project1_Transf _ _ _ _ z1z1') (<u>at</u> level 4) : poly_scope.
(* @ in ~_2 @ says argument *)
Notation "~_2 @ F1 _o>CoMod^ z2z2'" :=
  (@Project2_Transf F1 _ _ _ z2z2') (<u>at</u> level 4, F1 <u>at</u> next level) : poly_scope.
Notation "~_2 _o>CoMod^ z2z2'" :=
  (@Project2_Transf \_ \_ \_ \_ z2z2') (<u>at</u> level 4) : poly_scope.
Notation "<< f1f1' ,^CoMod f2f2' >>" :=
     Pairing_Transf _ _ _ _ f1f1' _ _ f2f2')
(at level 4, f1f1' at next level, f2f2' at next level, format "<< f1f1' ,^CoMod f2f2' >>"): poly_scope.
  (@Pairing_Transf _ _ _ _ f1f1'
```

7 Solution transformations

As common, the purely-grammatical polymorphism cut-constructors , for (multiplicative) outer/material composition [PolyMorCoMod] [TransfCoMod_PolyMorCoMod_Pre] [TransfCoMod_PolyMorCoMod_Post] and (coordinatewise) inner/structural composition [PolyTransfCoMod] , are not part of the solution terminology .

7.1 Solution transformations without polymorphism

```
Module Sol.
Export TWOFOLD.Sol.
Section Section1.
Delimit Scope sol scope with sol.
Open Scope sol scope.
Inductive transfCoMod : forall (F G : obCoMod), morCoMod F G -> morCoMod F G -> Type :=
| UnitTransfCoMod : forall (F G : obCoMod), forall (g : 'morCoMod(0 F ~> G )0),
      'transfCoMod(0 g ~> g )0
(**
| TransfCoMod_Gen : forall (F G : obCoMod_Gen), forall (g g' : morCoMod_Gen F G),
    transfCoMod_Gen g g' -> 'transfCoMod(0 (MorCoMod_Gen g) ~> (MorCoMod_Gen g') )0
| Project1_Transf : forall (F1 F2 : obCoMod) (Z1 : obCoMod),
    forall (z1 z1' : \text{morCoMod}(0 F1 \sim Z1))) (z1z1' : \text{transfCoMod}(0 z1 \sim z1')),
      'transfCoMod(0 ~_1 @ F2 o>CoMod z1 ~> ~_1 @ F2 o>CoMod z1' )0
| Project2_Transf : forall (F1 F2 : obCoMod) (Z2 : obCoMod),
    <u>forall</u> (z2 z2' : 'morCoMod(0 F2 ~> Z2 )0) (z2z2' : 'transfCoMod(0 z2 ~> z2' )0),
      'transfCoMod(0 \sim 2 @ F1 o>CoMod z2 \sim> \sim 2 @ F1 o>CoMod z2')0
```

```
| Pairing Transf : forall (L : obCoMod) (F1 F2 : obCoMod),
    \underline{forall} (f1 f1': 'morCoMod(0 L \sim F1 )0) (f1f1': 'transfCoMod(0 f1 \sim f1')0),
    forall (f2 f2': 'morCoMod(0 L ~> F2 )0) (f2f2': 'transfCoMod(0 f2 ~> f2')0),
       'transfCoMod(0 << f1 ,CoMod f2 >> ~> << f1' ,CoMod f2' >> )0
where "''transfCoMod' (0 g \sim g')0" := (@transfCoMod _ _ g g') : sol_scope.
End Section1.
Module Export Ex Notations.
Export TWOFOLD.Sol.Ex_Notations.
Delimit Scope sol scope with sol.
Notation "''transfCoMod' (\theta \ g \sim g')\theta" := (\theta transfCoMod _ _ g \ g') : sol_scope.
Notation "@ ''UnitTransfCoMod' g" := (@UnitTransfCoMod
                                          (<u>at</u> level 10, only parsing) : sol_scope.
Notation "''UnitTransfCoMod'" := (@UnitTransfCoMod _ _ _) (at level 0) : sol_scope.
(* ^CoMod in _o>CoMod^ says vertical transformation *)
Notation "\sim_1 @ \overline{F}2 _o>CoMod^ z1z1'" :=
  (@Project1_Transf _ F2 _ _ z1z1') (<u>at</u> level 4, F2 <u>at</u> next level) : sol_scope.
Notation "~_1 _o>CoMod^ z1z1'" :=
  (@Project1_Transf _ _ _ _ zlz1') (<u>at</u> level 4) : sol_scope.
(* ^CoMod in o>CoMod^ says vertical transformation *)
Notation "~ 2 @ \overline{F}1 o>CoMod^ z2z2'" :=
  (@Project2_Transf F1 _ _ _ z2z2') (at level 4, F1 at next level) : sol_scope.
Notation "~ 2 o>CoMod^ z2z2'" :=
  (@Project2_Transf \_ \_ \_ \_ z2z2') (<u>at</u> level 4) : sol_scope.
(* ^CoMod in ,^CoMod says vertical transformation *)
Notation "<< f1f1' , ^CoMod f2f2' >>" :=

(@Pairing_Transf ____ f1f1' ___ f2f2')
    (at level 4, f1f1' at next level, f2f2' at next level,
     format "<< f1f1' ,^CoMod f2f2' >>") : sol_scope.
End Ex Notations.
Fixpoint toPolyTransf (F G : obCoMod) (g g' : 'morCoMod(0 F <math>\sim G )0 %sol)
          (gg' : 'transfCoMod(0 g ~> g' )0 %sol) {struct gg'} :
  'transfCoMod(0 toPolyMor g ~> toPolyMor g' )0 %poly .
Proof.
  <u>refine</u>
    match gg' with
    | ( @'UnitTransfCoMod f )%sol => ( @'UnitTransfCoMod (toPolyMor f) )%poly
    | ( \sim_1 @ F2 _o>CoMod^ z1z1' )%sol =>
      ( ~_1 @ F2 _o>CoMod^ (toPolyTransf _ _ _ z1z1') )%poly
    | ( ~_2 @ F1 _o>CoMod^ z2z2' )%sol =>
    ( ~_2 @ F1 _o>CoMod^ (toPolyTransf _ _ _ z2z2') )%poly
| ( << f1f1' ,^CoMod f2f2' >> )%sol =>
   ( << (toPolyTransf _ _ _ f1f1') ,^CoMod (toPolyTransf _ _ _ f2f2') >> )%poly
    end.
Defined.
```

7.2 Destruction of transformations with inner-instantiation of morphism-indexes or object-indexes

Regardless the domain-codomain-morphisms extra-argument/parameter in the inductive-family-presentation [transfCoMod] of transformations, oneself easily still-gets the common dependent-destruction of transformations with inner-instantiation of object-indexes or inner-instantiation of morphism-indexes or inner-instantiation of both morphism-indexes and object-indexes. Memo that because the cut-adhesive will be inverted such to exfalso/exclude/impossible some couplings of

prefix-transformation and postfix-transformation or to obtain some subadhesive for some recursive call to resolution , therefore this below [Destruct_domProject1_Mor] and similar destruction lemmas will not be used ...

```
Module Destruct domPair.
Inductive transfCoMod domPair
: forall (F1 F2 G : obCoMod) (g g' : 'morCoMod(0 Pair F1 F2 ~> G )0 %sol),
    'transfCoMod(0 g ~> g' )0 %sol -> <u>Type</u> :=
| UnitTransfCoMod : ( forall (F1 F2 G : obCoMod),
                         forall (g : 'morCoMod(0 Pair F1 F2 ~> G )0),
                           transfCoMod_domPair (@'UnitTransfCoMod g ) )%sol
| Project1 Transf :
    ( forall (F1 F2 : obCoMod) (Z1 : obCoMod),
        forall (z1 z1' : 'morCoMod(0 F1 \sim Z1 )0) (z1z1' : 'transfCoMod(0 z1 \sim z1' )0),
          transfCoMod_domPair ( ~_1 @ F2 _o>CoMod^ z1z1' ) )%sol
| Project2 Transf :
    ( forall (F1 F2 : obCoMod) (Z2 : obCoMod),
        forall (z2 z2' : 'morCoMod(0 F2 ~> Z2 )0) (z2z2' : 'transfCoMod(0 z2 ~> z2' )0),
          transfCoMod_domPair ( ~_2 @ F1 _o>CoMod^ z2z2' ) )%sol
| Pairing_Transf :
    ( forall (L L' : obCoMod) (F1 F2 : obCoMod),
 forall (f1 f1' : 'morCoMod(0 Pair L L' ~> F1 )0) (f1f1' : 'transfCoMod(0 f1 ~> f1' )0),
 forall (f2 f2' : 'morCoMod(0 Pair L L' ~> F2 )0) (f2f2' : 'transfCoMod(0 f2 ~> f2' )0),
   transfCoMod domPair ( << f1f1' ,^CoMod f2f2' >> ) )%sol .
Definition transfCoMod_domPairP_Type
           (F G : obCoMod) (g g' : 'morCoMod(0 F ~> G )0 %sol)
           (gg' : 'transfCoMod(0 g ~> g' )0 %sol) :=
  ltac:( destruct F; refine (transfCoMod_domPair gg') ).
Lemma transfCoMod domPairP
  : forall (F G: obCoMod) (g g': 'morCoMod(0 F \sim G )0 %sol)
      (gg' : 'transfCoMod(0 g ~> g' )0 %sol),
    transfCoMod_domPairP_Type gg' .
Proof.
  intros. case: F G g g' / gg' .

    destruct F;

      constructor 1.
  - <u>constructor</u> 2.

    constructor
    3.

   <u>destruct</u> L; <u>constructor</u> 4.
Defined.
End Destruct domPair.
Module Destruct codPair.
Inductive transfCoMod codPair
: <u>forall</u> (F G1 G2 : obCoMod) (g g' : 'morCoMod(0 F ~> Pair G1 G2 )0 %sol),
    'transfCoMod(0 g ~> g' )0 %sol -> <u>Type</u> :=
| UnitTransfCoMod :
    ( forall (F G1 G2 : obCoMod),
        forall (g : 'morCoMod(0 F ~> Pair G1 G2 )0),
          transfCoMod_codPair (@'UnitTransfCoMod g ) )%sol
| Project1 Transf :
    ( forall (F1 F2 : obCoMod) (Z1 Z1' : obCoMod),
        forall (z1 z1' : 'morCoMod(0 F1 ~> Pair Z1 Z1' )0)
          (z1z1' : 'transfCoMod(0 z1 \sim> z1')0),
          transfCoMod_codPair ( ~_1 @ F2 _o>CoMod^ z1z1' ) )%sol
```

```
| Project2 Transf :
    ( forall (F1 F2 : obCoMod) (Z2 Z2' : obCoMod),
        forall (z2 z2' : 'morCoMod(0 F2 ~> Pair Z2 Z2' )0)
           (z2z2' : 'transfCoMod(0 z2 ~> z2' )0),
           transfCoMod codPair ( ~ 2 @ F1 o>CoMod^ z2z2' ) )%sol
| Pairing Transf :
    ( forall (L : obCoMod) (F1 F2 : obCoMod),
        forall (f1 f1': 'morCoMod(0 L \sim F1 )0) (f1f1': 'transfCoMod(0 f1 \sim f1')0),
        forall (f2 f2' : 'morCoMod(0 L ~> F2 )0) (f2f2' : 'transfCoMod(0 f2 ~> f2' )0),
           transfCoMod codPair ( << f1f1' ,^CoMod f2f2' >> ) )%sol .
Definition transfCoMod_codPairP_Type
            (\mathbf{F} \ \mathbf{G} : \mathsf{obCoMod}) \ (\mathbf{g} \ \mathbf{g}^{\mathsf{T}} : \mathsf{'morCoMod}(0 \ \mathsf{F} \sim \mathsf{S} \ )0 \ %sol)
            (gg' : 'transfCoMod(0 g ~> g' )0 %sol) :=
  ltac:( destruct G; refine (transfCoMod_codPair gg') ).
Lemma transfCoMod_codPairP
  : forall (F G: obCoMod) (g g': 'morCoMod(0 F \sim G )0 %sol)
      (gg' : 'transfCoMod(0 g ~> g' )0 %sol),
    transfCoMod_codPairP_Type gg' .
Proof.
  intros. case: F G g g' / gg'
  - <u>intros</u> ? G. <u>destruct</u> G; <u>constructor</u> 1.
  - intros ? ? Z1. destruct Z1. constructor 2.
  - intros ? ? Z2. destruct Z2. constructor 3.
  - constructor 4.
Defined.
End Destruct_codPair.
(** because of the cut-adhesive , therefore this below [Destruct_domProject1_Mor]
and similar destruction lemmas will not be used ... *)
Module Destruct_domProject1_Mor.
Inductive transfCoMod_domProject1_Mor
: <u>forall</u> (F1 F2 : obCoMod) (Z1 : obCoMod) (z1 : 'morCoMod(0 F1 ~> Z1 )0 %sol),
    forall (g : 'morCoMod(0 Pair F1 F2 ~> Z1 )0 %sol),
       'transfCoMod(0 (Project1_Mor F2 z1) ~> g )0 %sol -> <u>Type</u> :=
| UnitTransfCoMod :
    forall (F1 F2 : obCoMod) (Z1 : obCoMod) (Z1 : 'morCoMod(0 F1 \sim Z1 )0 %sol),
     transfCoMod_domProject1_Mor ( @'UnitTransfCoMod ( ~_1 @ F2 o>CoMod z1 ) )%sol
| Project1_Transf : forall (F1 F2 : obCoMod) (Z1 : obCoMod),
    <u>forall</u> (z1 z1' : 'morCoMod(0 F1 ~> Z1 )0 %sol)
      (z1z1' : 'transfCoMod(0 z1 ~> z1' )0 %sol),
      transfCoMod_domProject1_Mor ( ~_1 @ F2 _o>CoMod^ z1z1' )%sol .
Definition transfCoMod_domProject1_MorP_Type
            (F G : obCoMod) (g g' : 'morCoMod(0 F ~> G )0 %sol)
            (gg' : 'transfCoMod(0 g ~> g' )0 %sol) :=
  ltac:( destruct g; [ intros; refine (unit : Type)
                       | refine (transfCoMod domProject1 Mor gg')
                       | intros; refine (unit : Type)
                       | intros; refine (unit : Type) ] ).
Lemma transfCoMod_domProject1_MorP
  : \underline{forall} (F G : obCoMod) (g g' : 'morCoMod(0 F \sim G )0 %sol)
      (gg' : 'transfCoMod(0 g ~> g' )0 %sol),
    transfCoMod_domProject1_MorP_Type gg' .
  intros. case: F G g g' / gg' .
  - <u>destruct</u> g; [ <u>intros</u>; exact: tt | | <u>intros</u>; exact: tt ..] ;
    constructor 1.

    constructor 2.
```

```
- intros; exact: tt.
- intros; exact: tt.
Defined.

End Destruct_domProject1_Mor.
End Sol.
```

8 Grammatical 2-conversion of transformations , which infer the 1-conversions of their domain-codomain morphisms

As common , the grammatical 1-conversions-for-morphisms [convMorCoMod] ans 2-conversions-for-transformations [convTransfCoMod] are classified into : the total/(multi-step) conversions , and the congruences conversions , and the constant conversions which are used in the polymorphism/cut-elimination lemma , and the constant conversions which are only for the wanted sense of pairing-projections-grammar , and the constant conversions which are only for the confluence lemma , and the constant conversions which are derivable by using the finished cut-elimination lemma .

In contrast , because of the structural-multiplying-arrow ("degeneracy") action which is the unit-transformation-on-each-morphism [UnitTransfCoMod] type-constructor (which is elsewhere also hidden/blended in the outer left/right-whisk cut constructors [TransfCoMod_PolyMorCoMod_Pre] [TransfCoMod_PolyMorCoMod_Post]) , then the 2-conversions-for-transformations [convTransfCoMod] depends/uses of the 1-conversions-for-morphisms [convMorCoMod] , via the conversion-constructors [UnitTransfCoMod_cong] [TransfCoMod_PolyMorCoMod_Pre_cong] [TransfCoMod_PolyMorCoMod_Post_cong] .

8.1 Grammatical 2-conversions of transformations

```
Reserved Notation "g' g <~~2 g'g" (at level 70).
Inductive convTransfCoMod :
  forall (F G : obCoMod) (g g' : 'morCoMod(0 F ~> G )0 %poly)
    (gg' : 'transfCoMod(0 g \sim g') 0 %poly) (g_g'_ : 'morCoMod(0 F \sim G) 0 %poly)
    (g g' : 'transfCoMod(0 g \sim g' )0 %poly), Prop :=
(** ---- the total/(multi-step) conversions ---- **)
| convTransfCoMod Refl : forall (F G : obCoMod) (g : 'morCoMod(0 F ~> G )0 )
                              (gg : 'transfCoMod(0 g ~> g )0) ,
    gg <~~2 gg
| convTransfCoMod_Trans :
    forall (F G : obCoMod) (g g' : 'morCoMod(0 F ~> G )0 %poly)
  (gg' : 'transfCoMod(0 g ~> g' )0 %poly)
  (g0 g'0 : 'morCoMod(0 F ~> G )0 ) (uTrans : 'transfCoMod(0 g0 ~> g'0 )0 ),
      uTrans <~~2 gg' ->
  <u>forall</u> (g00 g'00 : morCoMod(0 F \sim G)0) (g00g'00 : transfCoMod(0 g00 <math>\sim g'00)0),
    g00g'00 <~~2 uTrans -> g00g'00 <~~2 gg'
(** ---- the congruences conversions ---- **)
| PolyTransfCoMod_cong : forall (F G : obCoMod),
    <u>forall</u> (g g' : \text{'morCoMod}(0 F \sim G))) (g'g : \text{'transfCoMod}(0 g' \sim g)),
    forall (g'0 g'' : 'morCoMod(0 F ~> G )0)
      (g''g'': 'transfCoMod(0 g'' \sim> g'0 )0) eqMor,
    (g''\_g'\_: 'transfCoMod(0 g''\_ \sim> g'\_0 )0) eqMor0, g''\_g'\_ <\sim 2 g''g' -> g'\_g\_ <\sim 2 g'g ->
      (g''_g'_o^{coMod} g'_g_# eqMor0) < \sim 2 (g''g' o^{coMod} g'g # eqMor)
| TransfCoMod_PolyMorCoMod_Pre_cong : forall (F G : obCoMod),
    <u>forall</u> (g g' : \text{'morCoMod}(0 F \sim G))) (g'g : \text{'transfCoMod}(0 g' \sim g))
    <u>forall</u> (g_g'_g'_: morCoMod(0 F \sim> G)0) (g'_g_: transfCoMod(0 g'_ <math>\sim> g_)0),
    forall (E: obCoMod) (f f_: 'morCoMod(0 E \sim> F )0),
```

```
f_ <~~1 f -> g'_g_ <~~2 g'g -> ( f__o>CoMod^ g'_g_ ) <~~2 ( f_o>CoMod^ g'g )
| TransfCoMod PolyMorCoMod Post cong :
    forall (G H : obCoMod) (h h : 'morCoMod(0 G \sim> H)0), forall (F : obCoMod),
        h_ <-~1 h -> g'_g_ <-~2 g'g -> ( g'_g_ ^o>CoMod_ h_ ) <-~2 ( g'g ^o>CoMod_ h )
| UnitTransfCoMod_cong : <u>forall</u> (F G : obCoMod), <u>forall</u> (g g' : 'morCoMod(0 F ~> G )0),
      g' <~~1 g -> (@'UnitTransfCoMod g') <~~2 (@'UnitTransfCoMod g)
| Project1 Transf cong : forall (F1 F2 : obCoMod) (Z1 : obCoMod),
    \underline{\text{forall}} (z1 z1 : 'morCoMod(0 F1 \sim Z1 )0) (z1z1' : 'transfCoMod(0 z1 \sim z1' )0),
 forall (z1 z1' : morCoMod(0 F1 <math>\sim Z1)0) (z1 z1' : transfCoMod(0 z1 <math>\sim z1' )0),
   z1 z1' <~~2 z1z1' ->
   ( ~_1 @ F2 _o>CoMod^ z1_z1'_ ) <~~2 ( ~_1 @ F2 _o>CoMod^ z1z1' )
| Project2_Transf_cong : forall (F1 F2 : obCoMod) (Z2 : obCoMod),
forall (z2 z2' : 'morCoMod(0 F2 ~> Z2 )0) (z2z2' : 'transfCoMod(0 z2 ~> z2' )0),
forall (z2_ z2'_ : 'morCoMod(0 F2 ~> Z2 )0) (z2_z2'_ : 'transfCoMod(0 z2_ ~> z2'_ )0),
   z2 z2' <~~2 z2z2' ->
   ( ~_2 @ F1 _o>CoMod^ z2_z2'_ ) <~~2 ( ~_2 @ F1 _o>CoMod^ z2z2' )
| Pairing_Transf_cong : forall (L : obCoMod) (F1 F2 : obCoMod),
    \underline{\text{forall}} (f1 f1': 'morCoMod(0 L ~> F1 )0) (f1f1': 'transfCoMod(0 f1 ~> f1')0),
  forall (f2 f2' : 'morCoMod(0 L ~> F2 )0) (f2f2' : 'transfCoMod(0 f2 ~> f2' )0),
forall (f2 f2' : 'morCoMod(0 L ~> F2 )0) (f2f2' : 'transfCoMod(0 f2 ~> f2' )0),
    f1_f1'_ <~~2 f1f1' -> f2_f2'_ <~~2 f2f2' ->
    ( << f1_f1'_, ^CoMod f2_f2'_>> ) <\sim 2 ( << f1f1', ^CoMod f2f2'>> )
(** ---- the constant conversions which are used during the polyarrowing
elimination ---- **)
(** here there are none parameter/customized-arrow action ( non-structural
reindexing , customized boundary-or-degeneracy ) ; is it possible to make sense of
such ? *)
(** ---- the constant conversions which are used during the polymorphism
elimination ---- **)
(* in other words : this polymorphisms also as structural-polyarrowing along the (
"degeneracy" ) structural-multiplying-arrow {0 -> 1} |- {0} *)
| UnitTransfCoMod morphismMor Pre :
    \underline{forall} (F G : obCoMod), \underline{forall} (g : 'morCoMod(0 F <math>\sim G )0),
        forall (E : obCoMod) (f : 'morCoMod(0 E ~> F )0),
          ( @'UnitTransfCoMod (f o>CoMod g)
            : 'transfCoMod(0 f o>CoMod g ~> f o>CoMod g )0 )
            <~~2 ( f o>CoMod^ ( @'UnitTransfCoMod q )
                 : 'transfCoMod(0 f o>CoMod g ~> f o>CoMod g )0 )
(* in other words : this polymorphisms also as structural-polyarrowing along the (
"degeneracy" ) structural-multiplying-arrow {0 -> 1} |- {0} *)
| UnitTransfCoMod_morphismMor_Post :
    forall (G H : obCoMod) (h : 'morCoMod(0 G ~> H )0),
    <u>forall</u> (F: obCoMod), <u>forall</u> (g: 'morCoMod(0 F \sim G )0),
        (@'UnitTransfCoMod (g o>CoMod h)
          : 'transfCoMod(0 g o>CoMod h ~> g o>CoMod h )0 )
          <~~2 ( (@'UnitTransfCoMod g ) ^o>CoMod h
                 'transfCoMod(0 g o>CoMod h ~> g o>CoMod h )0 )
| UnitTransfCoMod morphismTransf Pre :
    <u>forall</u> (F G : obCoMod), <u>forall</u> (g' : 'morCoMod(0 F <math>\sim S G))),
( g''g'0 ) <~~2 ( ( g''g'0 o^CoMod ( g'UnitTransfCoMod g' ) # eqMor )
                   : 'transfCoMod(0 g'' ~> g' )0 )
```

```
| UnitTransfCoMod morphismTransf Post : forall (F G : obCoMod),
    forall (g g' : "morCoMod(0 F \sim G))) (g'g : "transfCoMod(0 g' \sim g)))
      (g'\theta : 'morCoMod(0 F \sim G)0) eqMor,
      ( g'g ) <~~2 ( (@'UnitTransfCoMod g'0 ) o^CoMod g'g # egMor
                   : 'transfCoMod(0 g'0 ~> g )0 )
| UnitMorCoMod morphismTransf Pre :
<u>forall</u> (F G : obCoMod) (g g' : 'morCoMod(0 F \sim G)0) (g'g : 'transfCoMod(0 g' \sim G)0),
  ( g'g
    : 'transfCoMod(0 g' ~> g )0 )
    <~~2 ( g'g ^o>CoMod_ ( @'UnitMorCoMod G )
         : 'transfCoMod(0 g' o>CoMod 'UnitMorCoMod ~> g o>CoMod 'UnitMorCoMod )0 )
| UnitMorCoMod morphismTransf Post :
<u>forall</u> (F G : obCoMod) (g g' : 'morCoMod(0 F \sim G)0) (g'g : 'transfCoMod(0 g' \sim G)0),
  ( g'g
    : 'transfCoMod(0 g' ~> g )0 )
    <~~2 ( (@'UnitMorCoMod F ) _o>CoMod^ g'g
         : 'transfCoMod(0 'UnitMorCoMod o>CoMod g' ~> 'UnitMorCoMod o>CoMod g )0 )
| Project1_Mor_morphismTransf :
    forall (F1 F2 : obCoMod) (Z1 : obCoMod) (z1 : 'morCoMod(0 F1 ~> Z1 )0),
    forall (Y1 : obCoMod) (y' y : 'morCoMod(0 Z1 \sim Y1 )0)
      (y'y : 'transfCoMod(0 y' \sim y)0),
 ( \sim 1 @ F2 _o>CoMod^ (z1 _o>CoMod^ y'y)
   : 'transfCoMod(0 \sim 1 o>CoMod (z1 o>CoMod y') \sim \sim 1 o>CoMod (z1 o>CoMod y) )0 )
   <\sim\sim2 ( ( \sim 1 @ F2 o>CoMod z1 ) o>CoMod^ y'y
: 'transfCoMod(0 ( \sim 1 o>CoMod z1 ) o>CoMod y' \sim ( \sim 1 o>CoMod z1 ) o>CoMod y )0 )
| Project2 Mor morphismTransf :
    forall (F1 F2 : obCoMod) (Z2 : obCoMod) (z2 : 'morCoMod(0 F2 ~> Z2 )0),
forall (Y2 : obCoMod) (y' y : 'morCoMod(0 Z2 \sim Y2 )0) (y'y : 'transfCoMod(0 y' \sim y )0),
  ( ~ 2 @ F1 _o>CoMod^ (z2 _o>CoMod^ y'y)
   : transfCoMod(0 \sim 2 o>CoMod (z2 o>CoMod y') \sim 2 o>CoMod (z2 o>CoMod y) )0 )
| Project1_Transf_morphismMor : forall (F1 F2 : obCoMod) (Z1 : obCoMod),
    <u>forall</u> (z1 z1' : 'morCoMod(0 F1 ~> Z1 )0) (z1z1' : 'transfCoMod(0 z1 ~> z1' )0),
    forall (Y1 : obCoMod) (y : 'morCoMod(0 Z1 ~> Y1 )0),
( ~_1 @ F2 _o>CoMod^ (z1z1' ^o>CoMod_ y)
  : 'transfCoMod(0 ~ 1 o>CoMod (z1 o>CoMod y) ~> ~ 1 o>CoMod (z1' o>CoMod y) )0 )
  <-~2 ( ( ~_1 @ F2 _o>CoMod^ z1z1' ) ^o>CoMod_ y
: 'transfCoMod(0 ( \sim_1 o>CoMod z1 ) o>CoMod y \sim> ( \sim_1 o>CoMod z1' ) o>CoMod y )0 )
| Project2_Transf_morphismMor : <u>forall</u> (F1 F2 : obCoMod) (Z2 : obCoMod),
    forall (z2 z2' : 'morCoMod(0 F2 ~> Z2 )0) (z2z2' : 'transfCoMod(0 z2 ~> z2' )0),
    forall (Y2 : obCoMod) (y : 'morCoMod(0 Z2 \sim Y2 )0),
( \sim_2 @ F1 _o>CoMod^ (z2z2' ^o>CoMod_ y)
  : 'transfCoMod(0 \sim_2 o>CoMod (z2 o>CoMod y) \sim> \sim_2 o>CoMod (z2' o>CoMod y) )0 )
  <~~2 ( ( ~_2 @ F1 _o>CoMod^ z2z2' ) ^o>CoMod_ y
: 'transfCoMod(0 ( \sim 2 o>CoMod z2 ) o>CoMod y \sim ( \sim 2 o>CoMod z2' ) o>CoMod y )0 )
| Project1_Transf_morphismTransf : forall (F1 F2 : obCoMod) (Z1 : obCoMod),
    <u>forall</u> (z1 z1': 'morCoMod(0 F1 \sim Z1 )0) (z1z1': 'transfCoMod(0 z1 \sim z1' )0),
forall (z1'0 z1'' : 'morCoMod(0 F1 ~> Z1 )0) (z1'z1'' : 'transfCoMod(0 z1'0 ~> z1'' )0),
forall (eqMor_param : ( (~ 1 o>CoMod z1') <~>1 (~ 1 o>CoMod z1'0) ))
  (eqMor\ ex\ :\ (z1' <\sim>1\ z1'0)),
  ( ~_1 @ F2 _o>CoMod^ (z1z1' o^CoMod z1'z1'' # eqMor_ex)
   : 'transfCoMod(0 ~ 1 o>CoMod z1 \sim ~ 1 o>CoMod z1'' )0 )
| Project2_Transf_morphismTransf : forall (F1 F2 : obCoMod) (Z2 : obCoMod),
    <u>forall</u> (z2 z2' : 'morCoMod(0 F2 ~> Z2 )0) (z2z2' : 'transfCoMod(0 z2 ~> z2' )0),
forall (z2'0 z2'' : 'morCoMod(0 F2 ~> Z2 )0) (z2'z2'' : 'transfCoMod(0 z2'0 ~> z2'' )0),
```

```
forall egMor param egMor ex,
  (\sim 2 \oplus F1 \quad o>CoMod^ (z2z2' o^CoMod z2'z2'' \# eqMor ex)
    : 'transfCoMod(0 ~ 2 o>CoMod z2 \sim ~ 2 o>CoMod z2'' )0 )
   : 'transfCoMod(0 ~ 2 \overline{0}>CoMod z2 ~> ~ 2 \overline{0}>CoMod z\overline{2}'' )0 )
(**memo: Pairing_Mor_morphismTransf derivable below *)
| Pairing_Mor_morphismTransf :
    forall (L1 L2 : obCoMod) (F1 F2 : obCoMod) (f1 : 'morCoMod(0 Pair L1 L2 ~> F1 )0)
      (f2 : 'morCoMod(0 Pair L1 L2 ~> F2 )0),
   forall (M : obCoMod) (l1 l1' : 'morCoMod(0 M ~> L1 )0)
      (l1l1' : 'transfCoMod(0 l1 ~> l1' )0)
      (l2 l2' : 'morCoMod(0 M ~> L2 )0) (l2l2' : 'transfCoMod(0 l2 ~> l2' )0),
      ( << ( ( << l1l1' ,^CoMod l2l2' >> ) ^o>CoMod f1 )
           : 'transfCoMod(0 << ( ( << l1 ,CoMod l2 >> ) o>CoMod f1)
                      ,CoMod ( ( << l1 ,CoMod l2 >> ) o>CoMod f2) >> \sim>
                             << ( ( << l1' ,CoMod l2' >> ) o>CoMod f1)
                      ,CoMod ( ( << l1' ,CoMod l2' >> ) o>CoMod f2) >> )0 )
       <~~2 ( ( << l1l1' ,^CoMod l2l2' >> ) ^o>CoMod_ ( << f1 ,CoMod f2 >> )
            : 'transfCoMod(0 ( << l1 ,CoMod l2 >> ) o>CoMod << f1 ,CoMod f2 >> ~>
                         ( << l1' ,CoMod l2' >> ) o>CoMod << f1 ,CoMod f2 >> )0 )
(**memo: Pairing_Transf_morphismMor_derivable below *)
 Pairing Transf morphismMor : forall (L1 L2 : obCoMod) (F1 F2 : obCoMod),
<u>forall</u> (f\overline{1} f1': 'morCoMod(0 Pair L1 L2 \sim F1 )0) (f1f1': 'transfCoMod(0 f1 \sim f1')0),
forall (f2 f2' : 'morCoMod(0 Pair L1 L2 ~> F2 )0) (f2f2' : 'transfCoMod(0 f2 ~> f2' )0),
<u>forall</u> (M : obCoMod) (11 : 'morCoMod(0 M ~> L1 )0) (12 : 'morCoMod(0 M ~> L2 )0),
  ( << ( ( << l1 ,CoMod l2 >> ) _o>CoMod^ f1f1' )
       ,^CoMod ( ( << l1 ,CoMod l2 >> ) _o>CoMod^ f2f2' ) >>
    : 'transfCoMod(0 << ( ( << l1 ,CoMod l2 >> ) o>CoMod f1)
                  ,CoMod ( ( << l1 ,CoMod l2 >> ) o>CoMod f2) >> \sim>
                         << ( ( << l1 ,CoMod l2 >> ) o>CoMod f1')
                  ,CoMod ( ( << l1 ,CoMod l2 >> ) o>CoMod f2') >> )0 )
   ( << l1 ,CoMod l2 >> ) o>CoMod << f1' ,CoMod f2' >> )0 )
| Pairing_Transf_morphismTransf : <u>forall</u> (L : obCoMod) (F1 F2 : obCoMod),
   forall (f1 f1': 'morCoMod(0 L ~> F1 )0) (f1f1': 'transfCoMod(0 f1 ~> f1')0),
    forall (f2 f2': 'morCoMod(0 L ~> F2 )0) (f2f2': 'transfCoMod(0 f2 ~> f2')0),
forall (f1'0 f1'' : 'morCoMod(0 L ~> F1 )0) (f1'f1'' : 'transfCoMod(0 f1'0 ~> f1'' )0),
forall (f2'0 f2'' : 'morCoMod(0 L ~> F2 )0) (f2'f2'' : 'transfCoMod(0 f2'0 ~> f2'' )0),
forall eqMor1_ex eqMor2_ex eqMor_param,
( << f1f1' o^CoMod f1'f1'' # eqMor1_ex ,^CoMod f2f2' o^CoMod f2'f2'' # eqMor2_ex >>
  : 'transfCoMod(0 << f1 ,CoMod f2 >> ~> << f1'' ,CoMod f2'' >> )0 )
  <\sim\sim2 ( ( << f1f1' ,^CoMod f2f2' >> )
        o^CoMod ( << f1'f1'' ,^CoMod f2'f2'' >> ) # eqMor_param
       : 'transfCoMod(0 << f1 ,CoMod f2 >> \sim << f1'' ,CoMod f2'' >> )0 )
| Pairing Transf Project1 Mor : forall (L : obCoMod) (F1 F2 : obCoMod),
    \underline{forall} (f1 f1': 'morCoMod(0 L \sim F1 )0) (f1f1': 'transfCoMod(0 f1 \sim f1')0),
    forall (f2 f2' : 'morCoMod(0 L ~> F2 )0) (f2f2' : 'transfCoMod(0 f2 ~> f2' )0),
   forall (Z1 : obCoMod) (Z1 : 'morCoMod(0 F1 \sim Z1 )0),
      (flf1' ^o>CoMod z1
       : 'transfCoMod(0 f1 o>CoMod z1 ~> f1' o>CoMod z1 )0 )
       <~~2 ( ( << f1f1' ,^CoMod f2f2' >> ) ^o>CoMod_ ( \sim_1 @ F2 o>CoMod z1 )
            : 'transfCoMod(0 << f1 ,CoMod f2 >> o>CoMod ~ 1 o>CoMod z1 ~>
                             << f1' , CoMod f2' >> o>CoMod \sim 1 o>CoMod z1 )0 )
| Pairing Transf Project2 Mor : forall (L : obCoMod) (F1 F2 : obCoMod),
   <u>forall</u> (Z2 : obCoMod) (z2 : 'morCoMod(0 F2 ~> Z2 )0),
      ( f2f2' ^o>CoMod_ z2
        : 'transfCoMod(0 f2 o>CoMod z2 ~> f2' o>CoMod z2 )0 )
```

```
<-~2 ( ( << f1f1' ,^CoMod f2f2' >> ) ^o>CoMod ( ~ 2 @ F1 o>CoMod z2 )
              : 'transfCoMod(0 << f1 ,CoMod f2 >> o>CoMod ~ 2 o>CoMod z2 ~>
                                << f1' ,CoMod f2' >> o>CoMod ~ 2 o>CoMod z2 )0 )
| Pairing Mor Project1 Transf :
    forall (L: obCoMod) (F1 F2: obCoMod) (f1: 'morCoMod(0 L \sim> F1 )0)
       (f2 : 'morCoMod(0 L ~> F2 )0),
    \underline{\text{forall}} (Z1 : obCoMod) (Z1 Z1' : 'morCoMod(0 F1 \sim Z1 )0)
      (z1z1' : 'transfCoMod(0 z1 ~> z1' )0),
      ( f1 o>CoMod^ z1z1'
        : 'TransfCoMod(0 fl o>CoMod zl ~> fl o>CoMod zl' )0 )
        <\sim 2 ( ( << f1 ,CoMod f2 >> ) o>CoMod^ ( \sim 1 @ F2 o>CoMod^ zlz1' )
              : 'transfCoMod(0 << f1 ,CoMod f2 >> o>\overline{CoMod} \sim \overline{1} o>CoMod z1 \sim>
                                << f1 , CoMod f2 >> o>CoMod ~ 1 o>CoMod z1')0)
| Pairing_Mor_Project2_Transf :
    forall (L : obCoMod) (F1 F2 : obCoMod) (f1 : 'morCoMod(0 L <math>\sim F1 )0)
      (f2 : 'morCoMod(0 L ~> F2 )0),
    forall (Z2 : obCoMod) (z2 z2' : 'morCoMod(0 F2 ~> Z2 )0)
      (z2z2' : 'transfCoMod(0 z2 ~> z2' )0),
      ( f2 _o>CoMod^ z2z2'
        : 'transfCoMod(0 f2 o>CoMod z2 ~> f2 o>CoMod z2' )0 )
        <\sim\sim2 ( ( << f1 ,CoMod f2 >> ) _o>CoMod^ ( \sim_2 @ F1 _o>CoMod^ z2z2' )
              : 'transfCoMod(0 << f1 ,CoMod f2 >> o>CoMod ~_2 o>CoMod z2 ~>
                                << f1 ,CoMod f2 >> o>CoMod ~2 o>CoMod z2')0)
(** ---- the constant conversions which are only for the wanted sense of
pairing-projections-grammar ---- **)
(** Attention : for non-contextual ( "1-weigthed" ) pairing-projections , none of
    such thing as [Project1_Mor_Project2_Mor_Pairing_Transf] for transformations
    instead of [Project1_Mor_Project2_Mor_Pairing_Mor] for morphisms **)
(* structural-polyarrowing along the ( "degeneracy" ) structural-multiplying-arrow
\{0 \rightarrow 1\} \mid -\{0\} *\}
(*TODO: shall reverse this conversion ? *)
| Project1_UnitTransfCoMod :
    forall (F1 F2 Z1 : obCoMod) (z1 : 'morCoMod(0 F1 ~> Z1 )0),
      ( ~_1 @ F2 _o>CoMod^ ( @'UnitTransfCoMod z1 : 'transfCoMod(0 z1 ~> z1 )0 )
         : 'transfCoMod(0 ~_1 @ F2 o>CoMod z1 ~>
                 ( ~ 1 @ F2 o>CoMod z1 : 'morCoMod(0 Pair F1 F2 ~> Z1 )0 ) )0 )
        <\sim\sim2 ( @'UnitTransfCoMod ( \sim_1 @ F2 o>CoMod z1 ) )
(* structural-polyarrowing along the ( "degeneracy" ) structural-multiplying-arrow
\{0 \rightarrow 1\} \mid -\{0\} *\}
(*TODO: shall reverse this conversion ? *)
| Project2_UnitTransfCoMod : <u>forall</u> (F1 F2 : obCoMod) (Z2 : obCoMod)
    (z2: 'morCoMod(0 F2 ~> Z2 )0),
    ( ~_2 @ F1 _o>CoMod^ ( @'UnitTransfCoMod z2 )
   : 'transfCoMod(0 ~_2 o>CoMod z2 ~> ~_2 o>CoMod z2 )0 )
      <~~2 ( @'UnitTransfCoMod ( ~_2 @ F1 o>CoMod z2 )
            : 'transfCoMod(0 \sim_2 o>CoMod z2 \sim> \sim_2 o>CoMod z2 )0 )
(* structural-polyarrowing along the ( "degeneracy" ) structural-multiplying-arrow
\{0 \rightarrow 1\} \mid -\{0\} *\}
(*TODO: shall reverse this conversion ?*)
| Pairing UnitTransfCoMod : forall (L : obCoMod) (F1 F2 : obCoMod),
    \underline{forall} (f1 : 'morCoMod(0 L \sim F1 )0), \underline{forall} (f2 : 'morCoMod(0 L \sim F2 )0),
        ( << ( @'UnitTransfCoMod f1 ) ,^CoMod ( @'UnitTransfCoMod f2 ) >>
            : 'transfCoMod(0 << f1 ,CoMod f2 >> ~> << f1 ,CoMod f2 >> )0 )
           <~~2 ( @'UnitTransfCoMod << f1 ,CoMod f2 >>
                  'transfCoMod(0 << f1 ,CoMod f2 >> ~> << f1 ,CoMod f2 >> )0 )
(** ---- the constant conversions which are only for the confluence lemma --- **)
| Pairing Transf morphism Projectl Transf : forall (L : obCoMod) (F1 F2 : obCoMod),
    <u>forall</u> (f1 \ f1': 'morCoMod(0 L \sim F1 )0) (f1f1': 'transfCoMod(0 f1 \sim f1' )0),
```

```
forall (f2 \ f2': 'morCoMod(0 L \sim F2 )0) (f2f2': 'transfCoMod(0 f2 \sim f2' )0),
    forall (H : obCoMod),
      ( \sim_1 @ H _o>CoMod^ ( << f1f1' ,^CoMod f2f2' >> )
        : 'transfCoMod(0 ~ 1 o>CoMod ( << f1 ,CoMod f2 >> ) ~>
                           1 o>CoMod ( << f1' ,CoMod f2' >> ) )0 )
      <-~2 ( << ( ~_1 @ H _o>CoMod^ f1f1' )
           ,^CoMod ( ~_1 @ H _o>CoMod^ f2f2' ) >> : 'transfCoMod(0 << ( ~_1 o>CoMod f1 ) ,CoMod ( ~_1 o>CoMod f2 ) >> ~>
                           << ( \sim 1 o>CoMod f1' ) ,CoMod ( \sim 1 o>CoMod f2' ) >> )0 )
| Pairing Transf_morphism_Project2_Transf : forall (L : obCoMod) (F1 F2 : obCoMod),
    \underline{\text{forall}} (f1 f1': 'morCoMod(0 L \sim F1 )0) (f1f1': 'transfCoMod(0 f1 \sim f1')0),
    forall (f2 f2': 'morCoMod(0 L ~> F2 )0) (f2f2': 'transfCoMod(0 f2 ~> f2')0),
    forall (H : obCoMod),
      ( \sim_2 @ H _o>CoMod^ ( << f1f1' ,^CoMod f2f2' >> )
        : 'transfCoMod(0 ~ 2 o>CoMod ( << f1 ,CoMod f2 >> ) ~>
                           __2 o>CoMod ( << f1' ,CoMod f2' >> ) )0 )
      <~~2 ( << ( ~_2 @ H _o>CoMod^ f1f1' )
               ,^{\circ}CoMod ( ^{\circ}Z @ H _o>CoMod^{\circ} f2f2' ) >>
              'transfCoMod(0 << ( ~ 2 o>CoMod f1 ) ,CoMod ( ~ 2 o>CoMod f2 ) >> ~>
                            << ( \sim_2 o>CoMod f1' ) ,CoMod ( \sim_\overline{2} o>CoMod f2' ) >> )0 )
(** ---- the constant conversions which are derivable by using the finished
cut-elimination lemma ---- **)
(**
(*TODO: COMMENT *)
| PolyTransCoMod morphism : forall (F G : obCoMod),
    forall (g g' : 'morCoMod(0 F \sim> G)0) (g'g : 'transfCoMod(0 g' \sim> g)0),
    forall (g'0 g'' : 'morCoMod(0 F ~> G )0)
      (g''g': 'transfCoMod(0 g'' ~> g'0 )0) eqMor_param_in eqMor_ex_out,
  forall (g''0 g''' : 'morCoMod(0 F \sim> G )0) (g'''g'' : 'transfCoMod(0 g''' <math>\sim> g''0 )0)
    eqMor_param_out eqMor_ex_in,
    ( ( g'''g'' o^CoMod g''g' # eqMor_ex_in ) o^CoMod g'g # eqMor_ex_out
: 'transfCoMod(0 g''' ~> g )0 )
   <\sim\sim2 ( g'''g'' o^CoMod ( g''g' o^CoMod g'g # eqMor_param_in ) # eqMor_param_out
        : 'transfCoMod(0 g''' ~> g )0 )
(*TODO: COMMENT *)
| TransfCoMod_PolyMorCoMod_Pre_morphism_Pre : forall (F G : obCoMod),
    forall (g g' : 'morCoMod(0 F \sim> G)0) (g'g : 'transfCoMod(0 g' \sim> g)0),
    forall (E : obCoMod) (f : 'morCoMod(0 E ~> F )0),
    forall (D : obCoMod) (e : 'morCoMod(0 D ~> E )0),
      ( ( e o>CoMod f ) _o>CoMod^ g'g
        : 'transfCoMod(θ (e o>CoMod f) o>CoMod g' ~> (e o>CoMod f) o>CoMod g )θ )
        <\sim\sim2 ( e o>CoMod^ ( f o>CoMod^ g'g )
          : 'transfCoMod(0 e o>\overline{C}oMod (f o>CoMod g') ~> e o>CoMod (f o>CoMod g) )0 )
(*TODO: COMMENT *)
| TransfCoMod_PolyMorCoMod_Pre_morphism_Post : forall (F G : obCoMod),
    forall (g g' : 'morCoMod(0 F \sim> G)0) (g'g : 'transfCoMod(0 g' \sim> g)0),
    forall (E : obCoMod) (f : 'morCoMod(0 E ~> F )0),
    forall (H : obCoMod) (h : 'morCoMod(0 G ~> H )0),
      ( f _o>CoMod^ ( g'g ^o>CoMod_ h )
        : 'transfCoMod(0 f o>CoMod (g' o>CoMod h) ~> f o>CoMod (g o>CoMod h) )0 )
        <\sim\sim2 ( ( f o>CoMod^ g'g ) ^o>CoMod h
          : 'transfCoMod(θ (f o>CoMod g') o>CoMod h ~> (f o>CoMod g) o>CoMod h )θ )
(*TODO: COMMENT *)
| TransfCoMod PolyMorCoMod Post morphism Pre :
    forall (G H : obCoMod) (h : 'morCoMod(0 G ~> H )0), forall (F : obCoMod),
        forall (g \ g' : 'morCoMod(0 \ F \sim> G )0) \ (g'g : 'transfCoMod(0 \ g' \sim> g )0),
        forall (E : obCoMod) (f : 'morCoMod(0 E ~> F )0),
          ( ( f _o>CoMod^ g'g ) ^o>CoMod_ h
          : 'transfCoMod(θ (f o>CoMod g') o>CoMod h ~> (f o>CoMod g) o>CoMod h )θ )
            <~~2 ( f _o>CoMod^ ( g'g ^o>CoMod_ h )
          : 'transfCoMod(0 f o>CoMod (g' o>CoMod h) ~> f o>CoMod (g o>CoMod h) )0 )
```

```
(*TODO: COMMENT *)
| TransfCoMod PolyMorCoMod Post morphism Post :
    forall (G \ H : obCoMod) (h : 'morCoMod(0 \ G \sim> H )0), forall (F : obCoMod),
        forall (g g' : 'morCoMod(0 F \sim> G)0) (g'g : 'transfCoMod(0 g' \sim> g)0),
        forall (I : obCoMod) (i : 'morCoMod(0 H ~> I )0),
           ( g'g ^o>CoMod ( h o>CoMod i )
           : 'transfCoMod(0 g' o>CoMod (h o>CoMod i) ~> g o>CoMod (h o>CoMod i) )0 )
             <~~2 ( ( g'g ^o>CoMod_ h ) ^o>CoMod_ i
           : 'transfCoMod(0 (g' o>CoMod h) o>CoMod i ~> (g o>CoMod h) o>CoMod i )0 )
**)
(*TODO: COMMENT *)
| TransfCoMod PolyMorCoMod Pre morphismTransf : forall (F G : obCoMod),
    forall (g g' : 'morCoMod(0 F <math>\sim S G) ) ) (g'g : 'transfCoMod(0 g' <math>\sim S G) ) ),
    forall (E : obCoMod) (f : 'morCoMod(0 E ~> F )0),
    forall (g'0 g'' : 'morCoMod(0 F <math>\sim S G)0) (g''g' : 'transfCoMod(0 g'' <math>\sim S G)0),
    forall eqMor_param eqMor_ex,
      ( f _o>CoMod^ ( g''g' o^CoMod g'g # eqMor_ex )
   : 'transfCoMod(0 f o>CoMod g'' ~> f o>CoMod g )0 )
        <~~2 ( ( f _o>CoMod^ g''g' ) o^CoMod ( f _o>CoMod^ g'g ) # eqMor param
              : 'transfCoMod(0 f o>CoMod g'' ~> f o>CoMod g )0 )
(*TODO: COMMENT *)
| TransfCoMod_PolyMorCoMod_Post_morphismTransf :
    forall (GH: obCoMod) (h: 'morCoMod(0 G \sim> H)0), forall (F: obCoMod),
        <u>forall</u> (g g' : \text{'morCoMod}(0 F \sim G))) (g'g : \text{'transfCoMod}(0 g' \sim g))
        forall (g'0 g'' : \text{'morCoMod}(0 F \sim G))) (g''g' : \text{'transfCoMod}(0 g'' \sim g'0)),
        forall eqMor_param eqMor_ex,
           ( ( g''g' o^CoMod g'g # eqMor_ex ) ^o>CoMod_ h
             : 'transfCoMod(0 g'' o>CoMod h ~> g o>CoMod h )0 )
             <\sim\sim2 ( ( g''g' ^{\circ}o>CoMod_ h ) o^{\circ}CoMod ( g'g ^{\circ}o>CoMod_ h ) # eqMor_param
                  : 'transfCoMod(0 g'' o>CoMod h ~> g o>CoMod h )0 )
(** ---- the constant conversions which are derivable immediately without the
finished cut-elimination lemma ---- **)
(*TODO: COMMENT *)
| Pairing_Mor_morphismTransf_derivable :
    forall (L : obCoMod) (F1 F2 : obCoMod) (f1 : 'morCoMod(0 L ~> F1 )0)
      (f2 : morCoMod(0 L \sim F2)0),
    forall (L' : obCoMod) (l l0 : 'morCoMod(0 L' ~> L )0)
      (ll0 : 'transfCoMod(0 l ~> l0 )0),
      ( << ( ll0 ^o>CoMod_ f1 ) ,^CoMod ( ll0 ^o>CoMod_ f2 ) >>
        : 'transfCoMod(θ << (l o>CoMod f1) ,CoMod (l o>CoMod f2) >> ~>
                           << (l0 o>CoMod f1) ,CoMod (l0 o>CoMod f2) >> )0 )
        <~~2 ( ll0 ^o>CoMod ( << f1 ,CoMod f2 >> )
              : 'transfCoMod(0 l o>CoMod << f1 ,CoMod f2 >> ~>
                                10 o>CoMod << f1 ,CoMod f2 >> )0 )
(*TODO: COMMENT *)
| Pairing Transf morphismMor derivable : forall (L : obCoMod) (F1 F2 : obCoMod),
    forall\ (f1\ f1': 'morCoMod(0\ L \sim> F1\ )0)\ (f1f1': 'transfCoMod(0\ f1 \sim> f1'\ )0),
    forall (f2 f2' : 'morCoMod(0 L ~> F2 )0) (f2f2' : 'transfCoMod(0 f2 ~> f2' )0),
    forall (L': obCoMod) (l: 'morCoMod(0 L' \sim> L)0),
      ( << ( l o>CoMod^ f1f1' ) ,^CoMod ( l o>CoMod^ f2f2' ) >>
        : 'transfCoMod(θ << (l o>CoMod f1) ,CoMod (l o>CoMod f2) >> ~>
                           << (l o>CoMod f1') ,CoMod (l o>CoMod f2') >> )0 )
        <~~2 ( l o>CoMod^ ( << f1f1' ,^CoMod f2f2' >> )
              : 'transfCoMod(0 l o>CoMod << f1 ,CoMod f2 >> ~>
                                l o>CoMod << f1' ,CoMod f2' >> )0 )
**)
where "g_g'_ <~~2 gg'" := (@convTransfCoMod _ _ _ gg' _ gg'_) .
Hint Constructors convTransfCoMod.
```

8.2 1-convertibility of the domain/codomain morphisms for 2-convertible transformations

Because of the embedded/computed domain-codomain morphisms extra-argument/parameter in the inductive-family-presentation of the transformations , the 2-conversion-for-transformations relation shall convert across two transformations whose domain-codomain-morphisms-computation arguments are not syntactically/grammatically-the-same . But oneself does show that , by logical-deduction [convTransfCoMod_convMorCoMod_dom] [convTransfCoMod_convMorCoMod_cod] , these two domain-codomain-morphisms are indeed 1-convertible ("soundness lemma") .

```
Section NotUsed.
Lemma convTransfCoMod_convMorCoMod_dom :
  forall (F G : obCoMod) (g g' : 'morCoMod(0 F <math>\sim G) 0 %poly)
    (gg' : 'transfCoMod(0 g \sim g') 0 %poly) (g_g' : 'morCoMod(0 F \sim G) 0 %poly)
    (g_g'_: 'transfCoMod(0 g_ \sim> g'_ )0 %poly),
   g_g' < \sim 2 gg' -> g_ < \sim 1 g.
  induction 1 ; try solve [eauto];
   match goal with
    | [eqMor : ( _ <~>1 _ )%poly |- _] =>
     move: (EqMorCoMod.convMorCoMod_eq eqMor)
   <u>intros</u>; <u>subst</u>; <u>eauto</u>.
0ed.
Lemma convTransfCoMod_convMorCoMod_cod :
  forall (F G : obCoMod) (g g' : 'morCoMod(0 F ~> G )0 %poly)
    g_g'_ <\sim 2 gg' -> g'_ <\sim 1 g'.
Proof.
  induction 1 ; try solve [eauto];
   match goal with
    | [eqMor : ( _ <~>1 _ )%poly |- _] =>
     move: (EqMorCoMod.convMorCoMod eq eqMor)
    intros; subst; eauto.
Qed.
End NotUsed.
```

8.3 Linear total/asymptotic transformation-grade and the degradation lemma

```
- intros ? ? ? ? f1f1' ? ? f2f2' .
refine ($ ($ (may /* ) - )
    refine (S (S (max (gradeTransf _ _ _ f1f1') (gradeTransf _ _ _ f2f2')))).
Defined.
Lemma gradeTransf_gt0 : forall (F G : obCoMod) (g g' : 'morCoMod(0 F \sim G )0 ),
    <u>forall</u> (gg': 'transfCoMod(0 g ~> g')0),
     ((S 0) <= (gradeTransf gg'))%coq nat.</pre>
Proof. intros; case : gg'; intros; apply/leP; intros; simpl; auto. Qed.
Ltac tac_gradeTransf_gt0 :=
  match goal with
  \overline{\mid [gg\bar{1} : 'transfCoMod(0 \_ \sim> \_)0 ,
            gg2 : 'transfCoMod(0 ~>
                        gg4 : 'transfCoMod(0 _ ~> _ )0 ,
                  gg3 : 'transfCoMod(0 ~>
                                                    )0 |- ] =>
    (@gradeTransf_gt0 _ _ _ gg4)
  _ )0 ,
                         'transfCoMod(0 _ ~> _ )0 ,
gg4 : 'transfCoMod(0 _ ~> _
                  gg3 : 'transfCoMod(0
                                                    _ )0 |- _ ] =>
    \underline{move} \; : \; (@gradeTransf\_gt0 \; \_ \; \_ \; \_ \; gg1) \; (@gradeTransf\_gt0 \; \_ \; \_ \; \_ \; gg2)
                                           (@gradeTransf_gt0 _ _ _ gg3)
                                           (@gradeTransf_gt0 _ _ _ gg4)
  | [ gg1 : 'transfCoMod(0 \_ \sim > \_ )0 ,
            move : (@gradeTransf_gt0 _ _ _ gg1) (@gradeTransf_gt0 _ _ _ gg2)
                                           (@gradeTransf_gt0 _ _ _ gg3)
  | [ gg1 : 'transfCoMod(0 = \sim  )0 ,
            gg2 : 'transfCoMod(0 \_ \sim > \_ )0 |- \_ ] =>
    move : (@gradeTransf_gt0 _ _ _ gg1) (@gradeTransf_gt0 _ _ _ gg2)
  move : (@gradeTransf_gt0 _ _ _ gg1)
  <u>end</u> .
Lemma degradeTransf :
  forall (F G : obCoMod) (g g' : 'morCoMod(0 F <math>\sim G) 0 %poly)
    (gg' : 'transfCoMod(0 g \sim g') 0 \poly) (g_g' : 'morCoMod(0 F \sim G) 0 \poly)
    (g_g'_: 'transfCoMod(0 g_ \sim> g'_ )0 %poly),
    g_g'_<\sim\sim2 gg' -> ( gradeTransf g_g'_<= gradeTransf gg' )%coq_nat .
Proof.
  intros until g_g'_ . intros red_gg'.
elim : F G g g' gg' g_ g'_ g_g'_ / red_gg';
    try solve [ simpl; intros;
                try <u>match</u> goal <u>with</u>
                    | [ Hred : ( _ <~~1 _ ) |-
                                                 _ ] =>
                      move : (degradeMor Hred) ; clear Hred
                    <u>end</u>;
                intros; abstract Psatz.nia ].
(*memo: Omega.omega too weak at Pairing_Mor_morphismTransf
  Pairing Transf morphismMor *)
  (* erase associativities conversions then Qed. *)
0ed.
Ltac tac_degradeTransf H_gradeTransf :=
  <u>intuition</u> <u>idtac</u>;
  repeat <u>match</u> goal <u>with</u>
         | [ Hred : ( _ <~~1 _ ) |- _
                                      _ ] =>
           move : (degradeMor Hred) ; clear Hred
         | [ Hred : ( _ <~~2 _ ) |- _ ] =>
           move : (degradeTransf Hred) ; clear Hred
         <u>end</u>;
```

move: H_gradeTransf; clear; simpl; intros;
try tac_gradeMor_gt0; try tac_gradeTransf_gt0; intros; Omega.omega.

9 Polymorphism/cut-elimination by computational/total/asymptotic/reduction/(multi-step) resolution

For 2-folded polymorph mathematics , this resolution is made of some 1-resolution-for-morphisms [solveMorCoMod] and some 2-resolution-for-transformations [solveTransfCoMod] which depends/uses of this 1-resolution-for-morphisms .

In contrast from 1-resolution-for-morphisms , the 2-resolution-for-transformations has some almost-computational data-content function [solveTransfCoMod] of the resolution and some derived logical properties [solveTransfCoModP] which are satisfied by this function [solveTransfCoMod] . In other words : the programmation of the function [solveTransfCoMod] cannot be purelycomputational because of this fact : the 2-conversion-for-transformations do convert across two domain-codomain-morphisms-computation arguments transformations whose syntactically/grammatically-the-same , and therefore , during the recursion step which is the 2-resolution of the inner/structural ("vertical") composition [PolyTransfCoMod] cut-constructor , oneself lacks to know that the codomain-morphism of the recursively 2-resolved prefix-output transformation is grammatically-same (or propositionally-equal ...) as the domain-morphism of the recursively 2-resolved postfix-output transformation . In short : the programmation of the 2shall memorize the function [solveTransfCoMod] logical property that domain/codomain of the 2-resolved output transformation is propositionally-equal to the 1resolution of the domain/codomain of the input transformation .

Also in contrast , regardless that oneself may extract/derive the propositional-equations [solveTransfCoMod0_Project1_Transf] [solveTransfCoMod0_PolyTransfCoMod] corresponding to the purely-data part/component of the definitional-metaconversions of the 2-resolution-for-transformations function [solveTransfCoMod] , the fact that these propositional-equations are across ([existsT]) dependent-pairs (in other words , are dependent-equalities) , will make them not very-practical/usable . This suggests that the expected presentation of the 2-resolution (and 1-resolution) is not by this ongoing computational-resolution but instead shall be by some logical-resolution (logical congruent-rewrite-style cut-elimination resolution) . Ultimately , the future confluence lemmas will require this logical-resolution .

Memo that some more-lax (instead of tight/strict) alternative formulation of the inner/structural ("vertical") composition [PolyTransfCoMod] cut-constructor is possible . Such new formulation of the [PolyTransfCoMod] cut-constructor again will have some cut-adherence parameter which will relate the codomain-morphism of the prefix-input transformation in-relation-with the domainmorphism of the postfix-input transformation , by the (symmetrized , with associativity-conversion) 1-polymorphism-conversion-for-morphisms , instead of by propositional-equality . Consequently the new 2-resolution [solveTransfCoMod_requireConfluenceOf_solveMorCoMod] shall be formulated , by using 1-conversions instead of propositional-equality , from the domain/codomain of the input transformation to the domain/codomain of the 2-resolved output transformation . Now memo that the 1-confluence lemma for morphisms says that two morphisms are 1-polymorphismconvertible if and only if their 1-resolution is 1-solution-convertible . Finally such new formulations of [PolyTransfCoMod] and [solveTransfCoMod_requireConfluenceOf_solveMorCoMod] together with the finished 1-confluence lemma for morphisms , will enable to sucessfully do the recursion step which is the 2-resolution of the [PolyTransfCoMod] cut-constructor . BUT , even in such new formulation , for simplicity (to avoid any recursively-schematic 2-conversions-fortransformations ...) , the full 1-resolution of morphisms to their final solution will be required such that the 1-solution-conversion-for-morphism is easier and coincides with tight/strict propositional-equality (only congruences ...). Elsewhere , memo that , even in this alternative formulation , the "soundness lemma" [convTransfCoMod_convMorCoMod_dom] will not be immediatelyused during this alternative 2-resolution , but will only confirm after-the-fact the properties of this alternative 2-resolution .

As always , this COQ program and deduction is mostly-automated !

Module Resolve.
Import TWOFOLD.Resolve.

```
Export Sol.Ex_Notations.
Fixpoint solveTransfCoMod requireConfluenceOf solveMorCoMod len {struct len} :
  \underline{\text{forall}} (F G : obCoMod) (g g' : 'morCoMod(0 F \sim> G )0 )
    (gg' : 'transfCoMod(0 g ~> g' )0 ),
  forall gradeTransf_gg' : (gradeTransf gg' <= len)%coq nat,</pre>
    { g_ : 'morCoMod(0 F ~> G )0 %sol & { g'_ : 'morCoMod(0 F ~> G )0 %sol &
       { gg'Sol : 'transfCoMod(0 g_ ~> g'_ )0 %sol |
( ( Sol.toPolyMor g_ <~~1 g ) * ( Sol.toPolyMor g'_ <~~1 g' ) )%type } } } .
Abort.
Fixpoint solveTransfCoMod PolyTransfCoMod len {struct len} :
  forall (F G : obCoMod) (gSol g'Sol : 'morCoMod(0 F ~> G )0%sol)
    (g'Sol_gSol: 'transfCoMod(0 g'Sol ~> gSol )0%sol)
    (g'Sol0 g''Sol : 'morCoMod(0 F ~> G )0%sol)
    (g''Sol_g'Sol : 'transfCoMod(0 g''Sol ~> g'Sol0 )0%sol)
    (eqMor : Sol.toPolyMor g'Sol0 <~>1 Sol.toPolyMor g'Sol)
    (gradeTransf_gg' : (gradeTransf ((Sol.toPolyTransf g''Sol_g'Sol)
                   o^CoMod (Sol.toPolyTransf g'Sol_gSol) # eqMor) <= len)%coq_nat),
    \{g\_: \text{'morCoMod}(0 \ F \sim G )0\%sol \& \{g'\_: \text{'morCoMod}(0 \ F \sim G )0\%sol \& G \}
                                                  'transfCoMod(0 g_ ~> g'_ )0%sol } } .
Proof.
  <u>case</u> : len => [ | len ].
  (* len is 0 *)
  - ( move => F G gSol g'Sol g'Sol g'Sol0 g''Sol g''Sol g''Sol
                eqMor gradeTransf gg');
      exfalso; abstract tac_degradeTransf gradeTransf_gg'.
  (* len is (S len) *)
  (* gg' is g''g' o^CoMod g'g , to (g''Sol_g'Sol o^CoMod g'Sol_gSol) *)
  - move => F G gSol g'Sol g'Sol_gSol g'Sol0 g''Sol g''Sol_g'Sol
               eqMor gradeTransf_gg' .
    { destruct g''Sol_g'Sol as
           [ F G g'Sol0 (* @'UnitTransfCoMod g'Sol *)
             F1 F2 Z1 z1Sol z1'Sol z1Solz1'Sol (* ~_1 @ F2 _o>CoMod^ z1Solz1'Sol *)
F1 F2 Z2 z2Sol z2'Sol z2Solz2'Sol (* ~_2 @ F1 _o>CoMod^ z2Solz2'Sol *)
            L F1 F2 f1Sol f1'Sol f1Solf1'Sol f2Sol f2'Sol f2Solf2'Sol
(* << f1Solf1'Sol ,^CoMod f2Solf2'Sol >> *) ] .
      (* gg' is g''g' o^CoMod g'g , to (g''Sol g'Sol o^CoMod g'Sol gSol) , is
      (@'UnitTransfCoMod g'Sol o^CoMod g'Sol gSol) *)
      * <u>eexists</u>. <u>eexists</u>. <u>refine</u> (g'Sol_gSol)%sol .
      (* gg' is g''g' o^CoMod g'g , to (g''Sol\_g'Sol o^CoMod g'Sol\_gSol) , is ( (
      ~_1 @ F2 _o>CoMod^ z1Solz1'Sol ) o^CoMod g'Sol_gSol) *)
      * <u>move</u>: (Sol.Destruct_domPair.transfCoMod_domPairP g'Sol_gSol)
         => g'Sol_gSol_domPairP.
         { destruct g'Sol_gSol_domPairP as
               [ F1 F2 G g (* ( @'UnitTransfCoMod g %sol *) 
| F1 F2 Z1 _z1Sol _z1'Sol _z1Sol_z1'Sol
               (* ~_1 @ F2 _o>CoMod^ _z1Sol_z1'Sol *)
               | F1 F2 Z2 z2Sol z2'Sol z2Solz2'Sol
               (* ~_2 @ F1 _o>CoMod^ z2Solz2'Sol *)
               | L L' F1 F2 f1Sol f1'Sol f1Solf1'Sol f2Sol f2'Sol f2Solf2'Sol
(* << f1Solf1'Sol ,^CoMod f2Solf2'Sol >> *) ] .
           (* gg' is g''g' o^CoMod g'g , to (g''Sol\_g'Sol o^CoMod g'Sol\_gSol) , is
           ( (~_1 @ F2 _o>CoMod^ z1Solz1'Sol ) o^CoMod ( @'UnitTransfCoMod ( g ) ) *)
           - <u>eexists</u>. <u>eexists</u>. <u>refine</u> ( ~_1 @ F2 _o>CoMod^ z1Solz1'Sol )%sol .
           (* gg' is g''g' o^CoMod g'g , to (g''Sol\_g'Sol o^CoMod g'Sol\_gSol) , is
           ( ( \sim_1 @ F2 _o>CoMod^ z1Solz1'Sol ) o^CoMod ( \sim_1 @ F2 _o>CoMod^
           _z1Sol_z1'Sol ) *)
           - have [:blurb] z1Sol_z1' :=
               (projT2 (projT2 (solveTransfCoMod PolyTransfCoMod len
```

```
z1Sol z1'Sol
                                  z1Solz1'Sol
           (EqMorCoMod. Inversion Project1.convMorCoMod Project1P' eqMor) blurb )));
                first by clear -gradeTransf gg';
                abstract tac_degradeTransf gradeTransf_gg'
            eexists. eexists. refine ( ~_1 @ F2 _o>CoMod^ z1Sol_z1' )%sol .
          (* gg' is g''g' o^CoMod g'g , to (g''Sol\_g'Sol o^CoMod g'Sol\_gSol) , is
          ( ( ~_1 @ F2 _o>CoMod^ z1Solz1'Sol ) o^CoMod ( ~_2 @ F1 _o>CoMod^
          z2Solz2'Sol ) ) *)

    <u>exfalso</u>. <u>clear</u> -eqMor. <u>apply</u>:

    ( EqMorCoMod.Inversion_Exfalso.convMorCoMod_ExfalsoP_Project1_Project2_eqMor ).
          (* gg' is g''g' o^CoMod g'g , to (g''Sol_g'Sol o^CoMod g'Sol_gSol) , is
          ( ( ~_1 @ F2 _o>CoMod^ z1Solz1'Sol ) o^CoMod ( << f1Solf1'Sol ,^CoMod
          f2Solf2'Sol >> ) ) *)
           exfalso. clear -eqMor. apply:
     ( EqMorCoMod.Inversion_Exfalso.convMorCoMod_ExfalsoP_Project1_Pairing eqMor ).
      (* gg' is g''g' o^CoMod g'g , to (g''Sol\_g'Sol o^CoMod g'Sol\_gSol) , is ( (
      ~_2 @ F1 _o>CoMod^ z2Solz2'Sol ) o^CoMod g'Sol_gSol) *)
      * move: (Sol.Destruct_domPair.transfCoMod_domPairP g'Sol_gSol)
        => g'Sol_gSol_domPairP.
        { destruct g'Sol_gSol_domPairP as
              [ F1 F2 G g (* ( @'UnitTransfCoMod g %sol *)
              | F1 F2 Z1 z1Sol z1'Sol z1Solz1'Sol
              (* ~_1 @ F2 _o>CoMod^ z1Solz1'Sol *)
              | F1 F2 Z2 _z2Sol _z2'Sol _z2Sol z2'Sol
              (* ~_2 @ F1 _o>CoMod^ _z2Sol_z2'Sol *)
              L L' F1 F2 f1Sol f1'Sol f1Solf1'Sol f2Sol f2'Sol f2Solf2'Sol
(* << f1Solf1'Sol ,^CoMod f2Solf2'Sol >> *) ] .
          (* gg' is g''g' o^CoMod g'g , to (g''Sol_g'Sol o^CoMod g'Sol_gSol) , is
          ( ( ~_2 @ F1 _o>CoMod^ z2Solz2'Sol ) o^CoMod ( @'UnitTransfCoMod ( g ) )
          ) *)
          - \underline{eexists}. \underline{eexists}. \underline{refine} ( \sim_2 @ F1 \_o>CoMod^ z2Solz2'Sol )%sol .
          (* gg' is g''g' o^CoMod g'g , to (g''Sol\_g'Sol o^CoMod g'Sol\_gSol) , is
          ( ( ~_2 @ F1 _o>CoMod^ z2Solz2'Sol ) o^CoMod ( ~_1 @ F2 _o>CoMod^
          z1Solz1'Sol ) ) *)

    exfalso. <u>clear</u> -eqMor. <u>apply</u>:

            ( EqMorCoMod.Inversion Exfalso.convMorCoMod ExfalsoP Project1 Project2
                ( EqMorCoMod.convMorCoMod_sym eqMor) ).
          (* gg' is g''g' o^CoMod g'g , to (g''Sol_g'Sol o^CoMod g'Sol_gSol) , is
          ( (~_2 @ F1 _o>CoMod^ z2Solz2'Sol ) o^CoMod ( ~_2 @ F1 _o>CoMod^
          _z2Sol_z2'Sol<sup>_</sup>) ) *)
          - have [:blurb] z2Sol z2' :=
              (projT2 (projT2 (solveTransfCoMod_PolyTransfCoMod len _ _ _ _
                 (EqMorCoMod.Inversion_Project2.convMorCoMod_Project2P' eqMor) blurb )));
                first by <u>clear</u> -gradeTransf_gg';
                abstract tac_degradeTransf gradeTransf_gg'
            eexists. eexists. refine ( ~ 2 @ F1 _o>CoMod^ z2Sol_z2' )%sol .
          (* gg' is g''g' o^CoMod g'g , to (g''Sol_g'Sol o^CoMod g'Sol_gSol) , is
          ( ( ~_2 @ F1 _o>CoMod^ z2Solz2'Sol ) o^CoMod ( << f1Solf1'Sol ,^CoMod
          f2Solf2'Sol >> ) ) *)
           - <u>exfalso</u>. <u>clear</u> -eqMor. <u>apply</u>:
     ( EqMorCoMod.Inversion_Exfalso.convMorCoMod_ExfalsoP_Project2_Pairing eqMor ).
      (* gg' is g''g' o^CoMod g'g , to (g''Sol_g'Sol o^CoMod g'Sol_gSol) , is ( (
      << f1Solf1'Sol ,^CoMod f2Solf2'Sol >> ) o^CoMod g'Sol_gSol) *)
      * move: (Sol.Destruct_codPair.transfCoMod_codPairP g'Sol_gSol)
        => g'Sol_gSol_codPairP.
        { <u>destruct</u> g'Sol_gSol_codPairP <u>as</u>
```

```
[ F G1 G2 q (* ( @'UnitTransfCoMod (g) %sol *)
               | F1 F2 Z1 Z1' z1Sol z1'Sol z1Solz1'Sol
              (* \sim 1 @ F2 o > CoMod^ z1Solz1'Sol *)
              | F1 F2 Z2 Z2' z2Sol z2'Sol z2Solz2'Sol
              (* ~_2 @ F1 _o>CoMod^ z2Solz2'Sol *)
| L F1 F2 _f1Sol _f1'Sol _f1Sol_f1'Sol _f2Sol _f2'Sol _f2Sol_f2'Sol
(* << f1Sol f1'Sol ,^CoMod f2Sol f2'Sol >> *) ] .
          (* gg' is g''g' o^CoMod g'g , to (g''Sol\_g'Sol o^CoMod g'Sol\_gSol) , is
          ( ( << f1Solf1'Sol ,^CoMod f2Solf2'Sol >> ) o^CoMod ( @'UnitTransfCoMod
          (g) ) ) *)
          - eexists. eexists. refine ( << f1Solf1'Sol ,^CoMod f2Solf2'Sol >> )%sol .
          (* gg' is g''g' o^CoMod g'g , to (g''Sol_g'Sol o^CoMod g'Sol gSol) , is
          ( ( << f1Solf1'Sol ,^CoMod f2Solf2'Sol >> ) o^CoMod ( ~ 1 @ F2 o>CoMod^
          z1Solz1'Sol ) ) *)
           exfalso. clear -eqMor. apply:
            ( EqMorCoMod.Inversion Exfalso.convMorCoMod ExfalsoP Project1 Pairing
                 (EqMorCoMod.convMorCoMod_sym eqMor) ).
          (* gg' is g''g' o^CoMod g'g , to (g''Sol_g'Sol o^CoMod g'Sol_gSol) , is
          ( ( << f1Solf1'Sol ,^CoMod f2Solf2'Sol >> ) o^CoMod ( _2 @ F1 _o>CoMod^
          z2Solz2'Sol ) ) *)
          - <u>exfalso</u>. <u>clear</u> -eqMor. <u>apply</u>:
             ( EqMorCoMod.Inversion Exfalso.convMorCoMod ExfalsoP Project2 Pairing
                 (EqMorCoMod.convMorCoMod_sym eqMor) ).
          (* gg' is g''g' o^CoMod g'g , to (g''Sol g'Sol o^CoMod g'Sol gSol) , is
          ( ( << f1Solf1'Sol ,^CoMod f2Solf2'Sol >> ) o^CoMod ( << f1Sol f1'Sol
          ,^CoMod _f2Sol_f2'Sol >> ) ) *)
          - <u>simpl</u> <u>in</u> eqMor , gradeTransf_gg' .
            have [:blurb] f1Sol f1' :=
              (projT2 (projT2 (solveTransfCoMod_PolyTransfCoMod len _ _ _ _
                         f1Sol_f1'Sol _ _ f1Solf1'Sol
     (proj1 (EqMorCoMod_Inversion_Pairing.convMorCoMod_PairingP' eqMor)) blurb )));
                first by abstract tac_degradeTransf gradeTransf_gg' .
            have [:blurb] f2Sol_f2' :=
              (projT2 (projT2 (solveTransfCoMod_PolyTransfCoMod len _ _ _ _
                                 _f2Sol_f2'Sol _ _ f2Solf2'Sol
     (proj2 (EqMorCoMod.Inversion Pairing.convMorCoMod PairingP' eqMor)) blurb )));
                first by abstract tac_degradeTransf gradeTransf_gg'
            eexists. eexists. refine ( << f1Sol f1' ,^CoMod f2Sol f2' >> )%sol .
        }
    }
Defined.
Arguments solveTransfCoMod_PolyTransfCoMod !len _ _ _ _ !g''Sol_g'Sol _ _
  : <u>simpl</u> nomatch .
Notation "g''g' o^CoMod g'g # eqMor @ gradeTransf_gg'" :=
  (@solveTransfCoMod_PolyTransfCoMod _ _ _ _ g'g _ g''g' eqMor gradeTransf_gg') (at level 40 , g'g at next level, eqMor at next level) : sol_scope.
Lemma solveTransfCoMod_PolyTransfCoMod_len :
  forall len (F G : obCoMod) (gSol g'Sol : 'morCoMod(0 F ~> G )0%sol)
    (g'Sol gSol : 'transfCoMod(0 g'Sol ~> gSol )0%sol)
    (g'Sol0 g''Sol : 'morCoMod(0 F ~> G )0%sol)
    (g''Sol_g'Sol : 'transfCoMod(0 g''Sol ~> g'Sol0 )0%sol)
    (eqMor: Sol.toPolyMor g'Sol0 <~>1 Sol.toPolyMor g'Sol)
    (gradeTransf_gg'_len : (gradeTransf ((Sol.toPolyTransf g''Sol_g'Sol)
                    o^CoMod (Sol.toPolyTransf g'Sol gSol) # eqMor) <= len)%coq nat),
  o^CoMod (Sol.toPolyTransf g'Sol_gSol) \# eqMor') <= len')%coq_nat),
    ( g''Sol_g'Sol o^CoMod g'Sol_gSol # eqMor @ gradeTransf_gg'_len
      = g''Sol_g'Sol o^CoMod g'Sol_gSol # eqMor' @ gradeTransf_gg'_len' )%sol .
Proof.
```

```
induction len as [ | len ].
  - ( move => ? ? ? ? ? ? ? ? gradeTransf gg' len ); exfalso;
       clear -gradeTransf gg' len;
         by abstract tac_degradeTransf gradeTransf_gg'_len.

    intros. destruct len'.

    + <u>exfalso</u>; <u>clear</u> -gradeTransf_gg'_len';
         by abstract tac degradeTransf gradeTransf gg' len'.
    + destruct g''Sol g'Sol .
       * reflexivity.
       * {            <u>destruct</u> (Sol.Destruct domPair.transfCoMod domPairP g'Sol gSol);            <u>simpl</u>.
            - reflexivity.
           - erewrite IHlen. reflexivity.

    exfalso. apply

       (EqMorCoMod.Inversion_Exfalso.convMorCoMod_ExfalsoP_Project1_Project2_eqMor).
            - <u>exfalso</u>. <u>apply</u>
       (EqMorCoMod.Inversion Exfalso.convMorCoMod ExfalsoP Project1 Pairing eqMor).

    reflexivity.

    exfalso. apply

         (EqMorCoMod.Inversion Exfalso.convMorCoMod ExfalsoP Project1 Project2
                                   (EqMorCoMod.convMorCoMod_sym eqMor)).
           - erewrite IHlen. reflexivity.
           - <u>exfalso</u>. <u>apply</u>
        (EqMorCoMod.Inversion_Exfalso.convMorCoMod_ExfalsoP_Project2_Pairing_eqMor).
       * { <a href="mailto:destruct">destruct</a> (Sol.Destruct_codPair.transfCoMod_codPairP g'Sol_gSol); <a href="mailto:simpl">simpl</a>.
           - reflexivity.
           - <u>exfalso</u>. <u>apply</u>
                (EqMorCoMod.Inversion_Exfalso.convMorCoMod_ExfalsoP_Project1_Pairing
                    (EqMorCoMod.convMorCoMod_sym eqMor)).
           - <u>exfalso</u>. <u>apply</u>
                (EqMorCoMod.Inversion Exfalso.convMorCoMod ExfalsoP Project2 Pairing
                    (EqMorCoMod.convMorCoMod sym eqMor)).
           - <u>unfold</u> ssr_have. <u>simpl</u>.
              erewrite (IHlen _ _ _ f1f1' _ _ g''Sol_g'Sol1).
erewrite (IHlen _ _ _ f2f2' _ _ g''Sol_g'Sol2). reflexivity.
         }
Qed.
Definition solveTransfCoMod PolyTransfCoMod0 :
  forall (F G : obCoMod) (gSol g'Sol : 'morCoMod(0 F ~> G )0%sol)
     (g'Sol_gSol : 'transfCoMod(0 g'Sol ~> gSol )0%sol)
     (g'Sol0 g''Sol : 'morCoMod(0 F ~> G )0%sol)
     (g''Sol_g'Sol : 'transfCoMod(0 g''Sol ~> g'Sol0 )0%sol)
     (eqMor: Sol.toPolyMor g'Sol0 <~>1 Sol.toPolyMor g'Sol),
     \{g\_: \mathsf{'morCoMod}(0\ \mathsf{F} \leadsto \mathsf{G}\ )0\%\mathsf{sol}\ \&\ \{g'\_: \mathsf{'morCoMod}(0\ \mathsf{F} \leadsto \mathsf{G}\ )0\%\mathsf{sol}\ \&\ \}
                                                    'transfCoMod(0 g_ \sim g'_ )0%sol } } .
Proof.
  intros; apply: (@solveTransfCoMod_PolyTransfCoMod
                       (gradeTransf ((Sol.toPolyTransf g''Sol_g'Sol)
                               o^CoMod (Sol.toPolyTransf g'Sol_gSol) # eqMor))%coq_nat
          F G gSol g'Sol_gSol g'Sol0 g''Sol g''Sol_g'Sol eqMor); constructor.
Defined.
Notation "g''g' o^CoMod g'g # eqMor" :=
  (@solveTransfCoMod_PolyTransfCoMod0 _ _
    solveTransfCoMod_PolyTransfCoMod0 _ _ _ _ g'g _ _ g''g' eqMor) (<u>at</u> level 40 , g'g <u>at</u> next level) : sol_scope.
Lemma solveTransfCoMod PolyTransfCoMod0 len :
  <u>forall</u> len (F G : obCoMod) (gSol g'Sol : 'morCoMod(0 F <math>\sim G )0%sol)
     (g'Sol_gSol : 'transfCoMod(0 g'Sol ~> gSol )0%sol)
    (g'Sol0 g''Sol : 'morCoMod(0 F ~> G )0%sol)
(g''Sol_g'Sol : 'transfCoMod(0 g''Sol ~> g'Sol0 )0%sol)
     (eqMor: Sol.toPolyMor g'Sol0 <~>1 Sol.toPolyMor g'Sol)
     (gradeTransf_gg'_len : (gradeTransf ((Sol.toPolyTransf g''Sol_g'Sol)
                       o^CoMod (Sol.toPolyTransf g'Sol_gSol) # eqMor) <= len)%coq_nat),
```

```
forall (eqMor': Sol.toPolyMor q'Sol0 <~>1 Sol.toPolyMor q'Sol),
    ( g''Sol g'Sol o^CoMod g'Sol gSol # egMor'
     = g''Sol g'Sol o^CoMod g'Sol gSol # eqMor @ gradeTransf gg' len )%sol .
Proof. intros. erewrite solveTransfCoMod_PolyTransfCoMod_len. reflexivity. Qed.
Lemma solveTransfCoMod PolyTransfCoMod0 UnitTransfCoMod:
  forall (F G : obCoMod) (gSol g'Sol : 'morCoMod(0 F ~> G )0%sol)
    (g'Sol_gSol : 'transfCoMod(0 g'Sol ~> gSol )0%sol)
    (g'Sol0 : 'morCoMod(0 F ~> G )0%sol)
    (eqMor : Sol.toPolyMor g'Sol0 <~>1 Sol.toPolyMor g'Sol),
   projT2 (projT2 ('UnitTransfCoMod o^CoMod g'Sol_gSol # eqMor)%sol)
   = (g'Sol gSol)%sol.
Proof. reflexivity. Qed.
Lemma solveTransfCoMod PolyTransfCoMod0 UnitTransfCoMod Project1 Transf :
  forall (F1 G : obCoMod) (z1Sol z1'Sol : 'morCoMod(0 F1 ~> G )0%sol)
    (z1Solz1'Sol : 'transfCoMod(0 z1Sol ~> z1'Sol )0%sol)
    (F2 : obCoMod) (g : 'morCoMod(0 Pair F1 F2 ~> G )0%sol)
    (eqMor : ~_1 o>CoMod (Sol.toPolyMor z1'Sol) <~>1 Sol.toPolyMor g),
   projT2 (projT2 (~ 1 o>CoMod^ z1Solz1'Sol o^CoMod 'UnitTransfCoMod # eqMor)%sol)
   = ( ~_1 @ F2 _o>CoMod^ z1Solz1'Sol )%sol.
Proof. reflexivity. Qed.
Lemma solveTransfCoMod_PolyTransfCoMod0_Project1_Transf_Project1_Transf_dom :
  <u>forall</u> (F1 Z1 : obCoMod) (z1Sol z1'Sol : 'morCoMod(0 \overline{F1} \sim Z1 )0\%sol)
    (z1Solz1'Sol : 'transfCoMod(0 z1Sol ~> z1'Sol )0%sol)
    (F2: obCoMod) ( z1Sol: 'morCoMod(0 F1 ~> Z1 )0%sol)
    (eqMor : ~_1 o>CoMod (Sol.toPolyMor z1'Sol)
                        <~>1 Sol.toPolyMor (~_1 o>CoMod _z1Sol)%sol)
    (_z1'Sol : 'morCoMod(0 F1 ~> Z1 )0%sol)
    (_z1Sol_z1'Sol : 'transfCoMod(0 _z1Sol ~> _z1'Sol )0%sol),
 # EqMorCoMod.Inversion_Project1.convMorCoMod_Project1P' eqMor)))%sol.
Proof.
  <u>intros</u>. <u>simpl</u>. <u>erewrite</u> solveTransfCoMod_PolyTransfCoMod0_len. reflexivity.
Qed. (*TIME: 44 sec *)
Lemma solveTransfCoMod PolyTransfCoMod0 Project1 Transf Project1 Transf :
  forall (F1 Z1 : obCoMod) (z1Sol z1'Sol : 'morCoMod(0 F1 ~> Z1 )0%sol)
    (z1Solz1'Sol : 'transfCoMod(0 z1Sol ~> z1'Sol )0%sol) (F2 : obCoMod)
    ( z1Sol : 'morCoMod(0 F1 ~> Z1 )0%sol)
    (eqMor : ~_1 o>CoMod (Sol.toPolyMor z1'Sol)
                        <~>1 Sol.toPolyMor (~_1 o>CoMod _z1Sol)%sol)
    (_z1'Sol : 'morCoMod(0 F1 ~> Z1 )0%sol)
    (_z1Sol_z1'Sol : 'transfCoMod(0 _z1Sol ~> _z1'Sol )0%sol)
    (z1Sol_z1' := (z1Solz1'Sol o^CoMod _z1Sol_z1'Sol
            # (EqMorCoMod.Inversion_Project1.convMorCoMod_Project1P' eqMor))%sol),
    ( ( ( \sim_1 _o>CoMod^ z1Solz1'Sol )
          o^CoMod ( ~_1 _o>CoMod^ _z1Sol_z1'Sol ) # eqMor )%sol)
   = existT
            _ (~_1 o>CoMod (projT1 z1Sol_z1'))%sol
            (existT _ (~_1 o>CoMod (projT1 (projT2 z1Sol_z1' )))%sol
                    Proof.
  intros. subst z1Sol z1'. rewrite [solveTransfCoMod PolyTransfCoMod0 in LHS]lock.
 erewrite solveTransfCoMod PolyTransfCoMod0 len. unlock. reflexivity.
Qed. (*TIME: LONG 412 sec /!\ *)
(**ETC : ... *)
Fixpoint solveTransfCoMod PolyTransfCoModP len {struct len} :
  forall (F G : obCoMod) (gSol g'Sol : 'morCoMod(0 F ~> G )0%sol)
    (g'Sol_gSol : 'transfCoMod(0 g'Sol ~> gSol )0%sol)
    (g'Sol0 g''Sol : 'morCoMod(0 F ~> G )0%sol)
    (g''Sol_g'Sol : 'transfCoMod(0 g''Sol ~> g'Sol0 )0%sol)
    (eqMor : Sol.toPolyMor g'Sol0 <~>1 Sol.toPolyMor g'Sol)
    (gradeTransf_gg' : (gradeTransf ((Sol.toPolyTransf g''Sol_g'Sol)
```

```
o^CoMod (Sol.toPolyTransf g'Sol gSol) # egMor) <= len)%cog nat)
    (eqMor_param : Sol.toPolyMor g'Sol0 <~>1 Sol.toPolyMor g'Sol),
    ( ( (g''Sol = (projT1 (g''Sol_g'Sol o^CoMod g'Sol_gSol
                                           # eqMor @ gradeTransf gg' )%sol)) *
        (qSol = (projT1 (projT2 (q''Sol q'Sol o^CoMod q'Sol qSol
                                           # eqMor @ gradeTransf gg' )%sol))) ) *
      ( Sol.toPolyTransf (projT2 (projT2 (g''Sol_g'Sol o^CoMod g'Sol_gSol
                                           # eqMor @ gradeTransf_gg' )%sol))
       <-~2 (Sol.toPolyTransf g''Sol_g'Sol o^CoMod Sol.toPolyTransf g'Sol_gSol
                                                   # eqMor_param )%poly ) )%type .
Proof.
 <u>case</u> : len => [ | len ].
  (* len is 0 *)
  - ( move => F G gSol g'Sol g'Sol g'Sol g'Sol g''Sol g''Sol g'Sol
               eqMor gradeTransf_gg' eqMor_param );
      exfalso; abstract tac_degradeTransf gradeTransf_gg'.
  (* len is (S len) *)
  (* gg' is g''g' o^CoMod g'g , to (g''Sol_g'Sol o^CoMod g'Sol_gSol) *)
  - move => F G gSol g'Sol g'Sol_gSol g'Sol0 g''Sol g''Sol_g'Sol
            eqMor gradeTransf_gg' eqMor_param.
    { destruct g''Sol_g'Sol as
          [ F G g'Solo (* @'UnitTransfCoMod g'Solo *)
          | F1 F2 Z1 z1Sol z1'Sol z1Solz1'Sol (* ~_1 @ F2 _o>CoMod^ z1Solz1'Sol *)
          F1 F2 Z2 z2Sol z2'Sol z2Solz2'Sol (* ~_2 @ F1 _o>CoMod^ z2Solz2'Sol *)
          L F1 F2 f1Sol f1'Sol f1Solf1'Sol f2Sol f2'Sol f2Solf2'Sol
(* << f1Solf1'Sol ,^CoMod f2Solf2'Sol >> *) ] .
      (* gg' is g''g' o^CoMod g'g , to (g''Sol_g'Sol o^CoMod g'Sol_gSol) , is
      (@'UnitTransfCoMod g'Sol0 o^CoMod g'Sol gSol) *)
      * clear; move:
                 (EqMorCoMod.Inversion toPolyMor.convMorCoMod toPolyMorP' eqMor);
        abstract tac reduce solveMorCoMod0.
      (* gg' is g''g' o^CoMod g'g , to (g''Sol\_g'Sol o^CoMod g'Sol\_gSol) , is ( (
      ~_1 @ F2 _o>CoMod^ z1Solz1'Sol ) o^CoMod g'Sol_gSol) *)
      * move: (Sol.Destruct_domPair.transfCoMod_domPairP g'Sol_gSol)
        => g'Sol_gSol_domPairP.
        { destruct g'Sol_gSol_domPairP as
              [ F1 F2 G \overline{g} (* ( @'UnitTransfCoMod g %sol *)
              | F1 F2 Z1 _z1Sol _z1'Sol _z1Sol_z1'Sol
              (* ~ 1 @ F2 _o>CoMod^ _z1Sol_z1'Sol *)
              | F1 F2 Z2 z2Sol z2'Sol z2Solz2'Sol
              (* ~_2 @ F1 _o>CoMod^ z2Solz2'Sol *)
              L L' F1 F2 f1Sol f1'Sol f1Solf1'Sol f2Sol f2'Sol f2Solf2'Sol
(* << f1Solf1'Sol ,^CoMod f2Solf2'Sol >> *) ] .
          (* gg' is g''g' o^CoMod g'g , to (g''Sol_g'Sol o^CoMod g'Sol_gSol) , is
      ( ( ~_1 @ F2 _o>CoMod^ z1Solz1'Sol ) o^CoMod ( @'UnitTransfCoMod ( g ) ) ) *)
          - clear; move:
                   (EqMorCoMod.Inversion toPolyMor.convMorCoMod toPolyMorP' eqMor);
            abstract tac_reduce_solveMorCoMod0.
          (* gg' is g''g' o^CoMod g'g , to (g''Sol\_g'Sol o^CoMod g'Sol\_gSol) , is
          ( ( ~ 1 @ F2 o>CoMod^ z1Solz1'Sol ) o^CoMod ( ~ 1 @ F2 o>CoMod^
          _z1Sol_z1'Sol<sup>_</sup>) ) *)
          - <u>simpl</u>; <u>set</u> same_blurb := (_ gradeTransf_gg' : ( _ <= len )%coq_nat) .
           clear; abstract tac_reduce_solveMorCoMod0.
          (* gg' is g''g' o^CoMod g'g , to (g''Sol_g'Sol o^CoMod g'Sol_gSol) , is
          ( ( ~_1 @ F2 _o>CoMod^ z1Solz1'Sol ) o^CoMod ( ~_2 @ F1 _o>CoMod^
          z2Solz2'Sol ) ) *)
```

```
exfalso. clear -eqMor. apply:
    ( EqMorCoMod.Inversion_Exfalso.convMorCoMod_ExfalsoP_Project1_Project2_eqMor ).
          (* gg' is g''g' o^CoMod g'g , to (g''Sol_g'Sol o^CoMod g'Sol_gSol) , is
          ( ( ~_1 @ F2 _o>CoMod^ z1Solz1'Sol ) o^CoMod ( << f1Solf1'Sol ,^CoMod
          f2Solf2'Sol >> ) ) *)

    exfalso. clear -eqMor. apply:

     ( EqMorCoMod.Inversion Exfalso.convMorCoMod ExfalsoP Project1 Pairing eqMor ).
      (* gg' is g''g' o^CoMod g'g , to (g''Sol_g'Sol o^CoMod g'Sol_gSol) , is ( (
      ~ 2 @ F1 o>CoMod^ z2Solz2'Sol ) o^CoMod g'Sol gSol) *)
      * <u>move</u>: (Sol.Destruct_domPair.transfCoMod_domPairP g'Sol_gSol)
        => g'Sol_gSol_domPairP.
        { destruct g'Sol_gSol_domPairP as
              [ F1 F2 G g (* ( @'UnitTransfCoMod g %sol *)
              | F1 F2 Z1 z1Sol z1'Sol z1Solz1'Sol
              (* ~_1 @ F2 _o>CoMod^ z1Solz1'Sol *)
              | F1 F2 Z2 _z2Sol _z2'Sol _z2Sol_z2'Sol

(* ~ 2 @ F1 _o>CoMod^ _z2Sol_z2'Sol *)

| L L' F1 F2 f1Sol f1'Sol f1Solf1'Sol f2Sol f2'Sol f2Solf2'Sol
(* << f1Solf1'Sol ,^CoMod f2Solf2'Sol >> *) ] .
          (* gg' is g''g' o^CoMod g'g , to (g''Sol_g'Sol o^CoMod g'Sol_gSol) , is
      ( ( ~_2 @ F1 _o>CoMod^ z2Solz2'Sol ) o^CoMod ( @'UnitTransfCoMod ( g ) ) ) *)
            move: (EqMorCoMod.Inversion toPolyMor.convMorCoMod toPolyMorP' eqMor);
            abstract tac_reduce_solveMorCoMod0.
          (* gg' is g''g' o^CoMod g'g , to (g''Sol_g'Sol o^CoMod g'Sol_gSol) , is
          ( ( ~_2 @ F1 _o>CoMod^ z2Solz2'Sol ) o^CoMod ( ~_1 @ F2 _o>CoMod^
          z1Solz1'Sol ) ) *)
          - exfalso. clear -eqMor. apply:
            ( EqMorCoMod.Inversion Exfalso.convMorCoMod ExfalsoP Project1 Project2
                                           (EqMorCoMod.convMorCoMod_sym eqMor) ).
          (* gg' is g''g' o^CoMod g'g , to (g''Sol_g'Sol o^CoMod g'Sol_gSol) , is
          ( ( ~_2 @ F1 _o>CoMod^ z2Solz2'Sol ) o^CoMod ( ~_2 @ F1 _o>CoMod^
          _z2Sol_z2'Sol<sup>-</sup>) ) *)
          - <u>simpl</u>; <u>set</u> same_blurb := (_ gradeTransf_gg' : ( _ <= len )%coq nat) .
            move: (solveTransfCoMod PolyTransfCoModP len
                              _z2Sol_z2'Sol _ _ z2Solz2'Sol _ same_blurb
             (EqMorCoMod.Inversion_Project2.convMorCoMod_Project2P' eqMor_param) );
              clear; abstract tac_reduce_solveMorCoMod0.
          (* gg' is g''g' o^CoMod g'g , to (g''Sol_g'Sol o^CoMod g'Sol_gSol) , is
          ( ( \sim_2 @ F1 _o>CoMod^ z2Solz2'Sol ) o^CoMod ( << f1Solf1'Sol ,^CoMod
          f2Solf2'Sol >> ) ) *)

    exfalso. <u>clear</u> -eqMor. <u>apply</u>:

     ( EqMorCoMod.Inversion Exfalso.convMorCoMod ExfalsoP Project2 Pairing eqMor ).
      (* gg' is g''g' o^CoMod g'g , to (g''Sol_g'Sol o^CoMod g'Sol_gSol) , is ( (
      << f1Solf1'Sol ,^CoMod f2Solf2'Sol >> ) o^CoMod g'Sol_gSol) *)
      * move: (Sol.Destruct codPair.transfCoMod codPairP g'Sol gSol)
        => g'Sol gSol codPairP.
        { destruct g'Sol_gSol_codPairP as
              [ F G1 G2 g (* ( @'UnitTransfCoMod (g) %sol *)
              | F1 F2 Z1 Z1' z1Sol z1'Sol z1Solz1'Sol
              (* ~_1 @ F2 _o>CoMod^ z1Solz1'Sol *)
              | F1 F2 Z2 Z2' z2Sol z2'Sol z2Solz2'Sol
              (* ~_2 @ F1 _o>CoMod^ z2Solz2'Sol *)
               L F1 F2 _f1Sol _f1'Sol _f1Sol_f1'Sol _f2Sol _f2'Sol _f2Sol_f2'Sol
(* << _f1Sol_f1'Sol ,^CoMod _f2Sol_f2'Sol >> *) ] .
          (* gg' is g''g' o^CoMod g'g , to (g''Sol_g'Sol o^CoMod g'Sol_gSol) , is
          ( ( << f1Solf1'Sol ,^CoMod f2Solf2'Sol >> ) o^CoMod ( @'UnitTransfCoMod
```

```
clear;
             move: (EgMorCoMod.Inversion toPolyMor.convMorCoMod toPolyMorP' egMor);
             abstract tac_reduce_solveMorCoMod0.
           (* gg' is g''g' o^CoMod g'g , to (g''Sol g'Sol o^CoMod g'Sol gSol) , is
           ( ( << f1Solf1'Sol ,^CoMod f2Solf2'Sol >> ) o^CoMod ( ~_1 @ F2 _o>CoMod^
           z1Solz1'Sol ) ) *)
            <u>exfalso</u>. <u>clear</u> -eqMor. <u>apply</u>:
             (EqMorCoMod.Inversion Exfalso.convMorCoMod ExfalsoP Project1 Pairing
                (EqMorCoMod.convMorCoMod_sym eqMor)).
           (* gg' is g''g' o^CoMod g'g , to (g''Sol_g'Sol o^CoMod g'Sol_gSol) , is
           ( ( << f1Solf1'Sol ,^CoMod f2Solf2'Sol >> ) o^CoMod ( 2 @ F1 o>CoMod^
           z2Solz2'Sol ) ) *)
             exfalso. clear -eqMor. apply:
             (EqMorCoMod.Inversion_Exfalso.convMorCoMod_ExfalsoP_Project2_Pairing
                (EqMorCoMod.convMorCoMod_sym eqMor)).
           (* gg' is g''g' o^CoMod g'g , to (g''Sol_g'Sol o^CoMod g'Sol_gSol) , is
           ( ( << f1Solf1'Sol ,^CoMod f2Solf2'Sol >> ) o^CoMod ( << _f1Sol_f1'Sol</pre>
           ,^CoMod _f2Sol_f2'Sol >> ) ) *)
           - simpl; set same_blurb1 := (_ gradeTransf_gg' : ( _ <= len )%coq_nat) .</pre>
             move: (solveTransfCoMod_PolyTransfCoModP len
                                       f1Sol_f1'Sol___f1Solf1'Sol__same_blurb1
         (proj1 (EqMorCoMod.Inversion_Pairing.convMorCoMod_PairingP' eqMor_param))).
             <u>set</u> same blurb2 := ( gradeTransf gg' : ( <= len )%coq nat) .
             move: (solveTransfCoMod_PolyTransfCoModP len
                                        f2Sol f2'Sol
                                                         f2Solf2'Sol same blurb2
         (proj2 (EqMorCoMod.Inversion Pairing.convMorCoMod PairingP' eqMor param))).
             clear; abstract tac_reduce_solveMorCoMod0.
        }
Qed.
Fixpoint solveTransfCoMod TransfCoMod PolyMorCoMod Pre len {struct len} :
  \underline{\text{forall}} (F G : obCoMod) (g'Sol \underline{\text{gSol}} : 'morCoMod(0^{-}F \sim> G )0%sol)
    (g'Sol_gSol : 'transfCoMod(0 g'Sol ~> gSol )0%sol)
    (E : obCoMod) (fSol : 'morCoMod(0 E ~> F )0%sol)
    (gradeTransf_gg' : (gradeTransf (Sol.toPolyMor fSol
                              o>CoMod^ Sol.toPolyTransf g'Sol gSol) <= len)%coq nat),
    {q : 'morCoMod(0 E ~> G )0%sol &
           {g'_: 'morCoMod(0 E ~> G )0%sol &
                  'transfCoMod(0 g_ \sim> g'_ )0%sol } }.
Proof.
  <u>case</u> : len => [ | len ].
  (* len is 0 *)
  - ( move => F G g'Sol gSol g'Sol gSol E fSol gradeTransf gg' );
      exfalso; abstract tac_degradeTransf gradeTransf_gg'.
  (* len is (S len) *)
  (* gg' is f _o>CoMod^ g'g , to (fSol _o>CoMod^ g'Sol_gSol) *)
  - <u>move</u> => F G g'Sol gSol g'Sol_gSol E fSol gradeTransf_gg'.
    { <u>destruct</u> g'Sol_gSol <u>as</u>
           [ F G gSol (* @'UnitTransfCoMod gSol *)
            F1 F2 Z1 z1Sol z1'Sol z1Solz1'Sol (* ~_1 @ F2 _o>CoMod^ z1Solz1'Sol *)
            F1 F2 Z2 z2Sol z2'Sol z2Solz2'Sol (* ~ 2 @ F1 _o>CoMod^ z2Solz2'Sol *)
           | L F1 F2 f1Sol f1'Sol f1Solf1'Sol f2Sol f2'Sol f2Solf2'Sol
(* << f1Solf1'Sol ,^CoMod f2Solf2'Sol >> *) ] .
      (* gg' is f _o>CoMod^ g'g , to (fSol _o>CoMod^ g'Sol_gSol) , is (fSol
      _o>CoMod^ @'UnitTransfCoMod gSol) *)
* have fSol_o_gSol_prop :=
      (@solveMorCoModOP _ _ ( (Sol.toPolyMor fSol) o>CoMod (Sol.toPolyMor gSol) )).
    set fSol_o_gSol := (@solveMorCoModO _ _ ( (Sol.toPolyMor fSol)
                                 o>CoMod (Sol.toPolyMor gSol) )) in fSol_o_gSol_prop.
```

(g))) *)

```
eexists. eexists. refine (@'UnitTransfCoMod fSol o gSol )%sol .
(* gg' is f o>CoMod^ g'g , to (fSol o>CoMod^ g'Sol gSol) , is (fSol
o>CoMod^ ~_1 @ F2 _o>CoMod^ z1Solz1'Sol) *)
* move: (Sol.Destruct_codPair.morCoMod_codPairP fSol) => fSol_codPairP.
  { destruct fSol codPairP as
         [ F1 F2 (* (@'UnitMorCoMod (Pair F1 F2) )%sol *)
         | _F1 _F2 _Z1 _Z1' _z1 (* ( ~_1 @ _F2 o>CoMod _z1 )%sol *)
| _F1 _F2 _Z2 _Z2' _z2 (* ( ~_2 @ _F1 o>CoMod _z2 )%sol *)
| L F1 F2 f1 f2 (* ( << f1 ,CoMod f2 >> )%sol *) ] .
    (* gg' is f o>CoMod^ g'g , to (fSol o>CoMod^ g'Sol gSol) , is
   ( @'UnitMorCoMod (Pair F1 F2) _o>CoMod^ ~_1 @ _ _o>CoMod^ z1Solz1'Sol ) *)
     eexists. eexists. refine ( ~ 1 @ F2 o>CoMod^ z1Solz1'Sol )%sol .
    (* gg' is f _o>CoMod^ g'g , to (fSol _o>CoMod^ g'Sol_gSol) , is ( ( \sim\!\!\!\_1
    @ _F2 o>CoMod _z1 ) _o>CoMod^ ( ~_1 @ _ _o>CoMod^ z1Solz1'Sol ) ) *)
- have [:blurb] _z1_o_g'Sol__z1_o_gSol :=
         (projT2 (@solveTransfCoMod_TransfCoMod_PolyMorCoMod_Pre
                  _ _ _ ( ~_1 @ _ _o>CoMod^ z1Solz1'Sol )%sol _ _z1 blurb)));
           first by abstract tac_degradeTransf gradeTransf_gg' .
      <u>eexists</u>. <u>eexists</u>.
      <u>refine</u> ( ~_1 @ _F2 _o>CoMod^ _z1_o_g'Sol__z1_o_gSol )%sol .
    (* gg' is f _o>CoMod^ g'g , to (fSol _o>CoMod^ g'Sol_gSol) , is ( ( ~_2
    @ _F1 o>CoMod _z2 ) _o>CoMod^ ( ~_1 @ _ _o>CoMod^ z1Solz1'Sol ) ) *)
    - have [:blurb] z2 o g'Sol z2 o gSol :=
         (projT2 (projT2 (@solveTransfCoMod_TransfCoMod_PolyMorCoMod_Pre
                          _ ( ~_1 @ _ _o>CoMod^ z1Solz1'Sol )%sol _ _z2 blurb)));
           first by abstract tac_degradeTransf gradeTransf gg' .
      <u>eexists</u>. <u>eexists</u>.
      <u>refine</u> ( ~_2 @ _F1 _o>CoMod^ _z2_o_g'Sol__z2_o_gSol )%sol .
    (* gg' is f _o>CoMod^ g'g , to (fSol _o>CoMod^ g'Sol_gSol) , is ( << f1 \,
    ,CoMod f2 >> _o>CoMod^ ( ~_1 @ _ _o>CoMod^ z1Solz1'Sol ) ) *)
- have [:blurb] f1_o_z1Sol_f1_o_z1'Sol :=
         (projT2 (@solveTransfCoMod_TransfCoMod_PolyMorCoMod_Pre
                               len _ _ _ z1Solz1'Sol _ f1 blurb)));
           first by abstract tac_degradeTransf gradeTransf_gg' .
      eexists. eexists. refine ( f1 o z1Sol f1 o z1'Sol )%sol .
  }
(* gg' is f _o>CoMod^ g'g , to (fSol _o>CoMod^ g'Sol_gSol) , is (fSol
_o>CoMod^ ~_2 @ F1 _o>CoMod^ z2Solz2'\(\overline{S}ol\) *)
* move: (Sol.Destruct_codPair.morCoMod_codPairP fSol) => fSol_codPairP.
  { destruct fSol_codPairP as
         [ F1 F2 (* (@'UnitMorCoMod (Pair F1 F2) )%sol *)
         | _F1 _F2 _Z1 _Z1' _z1 (* ( ~_1 @ _F2 o>CoMod _z1 )%sol *)
| _F1 _F2 _Z2 _Z2' _z2 (* ( ~_2 @ _F1 o>CoMod _z2 )%sol *)
| L F1 F2 f1 f2 (* ( << f1 ,CoMod f2 >> )%sol *) ] .
    (* gg' is f _o>CoMod^ g'g , to (fSol _o>CoMod^ g'Sol_gSol) , is
  ( @'UnitMorCoMod (Pair F1 F2) _o>CoMod^ ~_2 @ _ _o>CoMod^ z2Solz2'Sol ) *)
    - <u>eexists</u>. <u>eexists</u>. <u>refine</u> ( ~_2 @ F1 _o>CoMod^ z2Solz2'Sol )%sol .
    (* gg' is f _o>CoMod^ g'g , to (fSol _o>CoMod^ g'Sol_gSol) , is ( ( ~_1
    @ _F2 o>CoMod _z1 ) _o>CoMod^ ( ~_2 @ _ _o>CoMod^ z2Solz2'Sol ) ) *)
    - have [:blurb] _z1_o_g'Sol__z1_o_gSol :=
         (projT2 (projT2 (@solveTransfCoMod_TransfCoMod_PolyMorCoMod_Pre
           len _ _ _ ( ~_2 @ _ _o>CoMod^ z2Solz2'Sol )%sol _ _z1 blurb)));
first by abstract tac_degradeTransf gradeTransf_gg' .
      <u>eexists</u>. <u>eexists</u>.
      <u>refine</u> ( ~_1 @ _F2 _o>CoMod^ _z1_o_g'Sol__z1_o_gSol )%sol .
    (* gg' is f _o>CoMod^ g'g , to (fSol _o>CoMod^ g'Sol_gSol) , is ( ( ~_2
    @ _F1 o>CoMod _z2 ) _o>CoMod^ ( ~_2 @ _ _o>CoMod^ z2Solz2'Sol ) ) *)
    - have [:blurb] z2 o g'Sol z2 o gSol :=
```

```
(projT2 (projT2 (@solveTransfCoMod TransfCoMod PolyMorCoMod Pre
                   len _ _ _ ( ~_2 @ _ _o>CoMod^ z2Solz2'Sol )%sol _ _z2 blurb )));
                  first by abstract tac_degradeTransf gradeTransf_gg' .
             <u>eexists</u>. <u>eexists</u>.
             refine ( ~_2 @ _F1 _o>CoMod^ _z2_o_g'Sol__z2_o_gSol )%sol .
           (* gg' is f _o>CoMod^ g'g , to (fSol _o>CoMod^ g'Sol_gSol) , is ( << f1
           , CoMod f2 >> _o>CoMod^ ( ~_2 @ _ _o>CoMod^ z2Solz2'Sol ) ) *)
- have [:blurb] f2_o_z2Sol_f2_o_z2'Sol :=
                (projT2 (projT2 (@solveTransfCoMod TransfCoMod PolyMorCoMod Pre
                                     len _ _ _ z2Solz2'Sol _ f2 blurb)));
                  first by abstract tac degradeTransf gradeTransf gg' .
             eexists. eexists. refine (f2 o z2Sol f2 o z2'Sol)%sol.
         }
       (* gg' is f _o>CoMod^ g'g , to (fSol _o>CoMod^ g'Sol_gSol) , is (fSol
       o>CoMod^ << f1Solf1'Sol ,^CoMod f2Solf2'Sol >> ) *)
       * { <u>destruct</u> fSol <u>as</u>
                [ F (* @'UnitMorCoMod F *)
                | _F1 _F2 Z1 z1Sol (* ~_1 @ F2 o>CoMod z1Sol *)
| _F1 _F2 Z2 z2Sol (* ~_2 @ F1 o>CoMod z2Sol *)
                | L _F1 _F2 _f1Sol _f2Sol (* << _f1Sol ,CoMod _f2Sol >> *) ] .
           (* gg' is f _o>CoMod^ g'g , to (fSol _o>CoMod^ g'Sol_gSol) , is
           ( @'UnitMorCoMod F _o>CoMod^ << f1Solf1'Sol ,^CoMod f2Solf2'Sol >> ) *)
           - <u>eexists</u>. <u>eexists</u>. <u>refine</u> ( << f1Solf1'Sol ,^CoMod f2Solf2'Sol >> )%sol.
           (* gg' is f _o>CoMod^ g'g , to (fSol _o>CoMod^ g'Sol_gSol) , is ( ( ~_1
         @ F2 o>CoMod zISol ) o>CoMod^ << fISolf1'Sol ,^CoMod f2Solf2'Sol >> ) *)
            - <u>have</u> [:blurb] z1Sol_o_g'Sol_z1Sol_o_gSol :=
                (projT2 (projT2 (@solveTransfCoMod_TransfCoMod_PolyMorCoMod_Pre
                  __ ( << f1Solf1'Sol ,^CoMod f2Solf2'Sol >> )%sol _ z1Sol blurb)));
first by abstract tac_degradeTransf gradeTransf_gg' .
             <u>eexists</u>. <u>eexists</u>.
             refine ( ~_1 @ _F2 _o>CoMod^ z1Sol_o_g'Sol_z1Sol_o_gSol )%sol .
     (* gg' is f _o>CoMod^ g'g , to (fSol _o>CoMod^ g'Sol_gSol) , is ( ( ~_2 @ _F1 o>CoMod z2Sol ) _o>CoMod^ ( << f1Solf1'Sol ,^CoMod f2Solf2'Sol >> ) ) *)
            - have [:blurb] z2Sol_o_g'Sol_z2Sol_o_gSol :=
                (projT2 (projT2 (@solveTransfCoMod_TransfCoMod_PolyMorCoMod_Pre
                    ( << f1Solf1'Sol ,^CoMod f2Solf2'Sol >> )%sol _ z2Sol blurb)));
                  first by abstract tac_degradeTransf gradeTransf_gg'
             <u>eexists</u>. <u>eexists</u>.
             refine ( ~_2 @ _F1 _o>CoMod^ z2Sol_o_g'Sol_z2Sol_o_gSol )%sol .
   (* gg' is f _o>CoMod^ g'g , to (fSol _o>CoMod^ g'Sol_gSol) , is ( <<
_f1Sol ,CoMod _f2Sol >> _o>CoMod^ ( << f1Solf1'Sol ,^CoMod f2Solf2'Sol >> ) ) *)
             * have [:blurb] fSol_o_f1Sol_fSol_o_f1'Sol :=
                  (\texttt{projT2} \ (\texttt{@solveTransfCoMod\_TransfCoMod\_PolyMorCoMod\_Pre})
             len _ _ _ f1Solf1'Sol _ ( << _f1Sol ,CoMod _f2Sol \rightarrow %sol) blurb)));
                    first by abstract tac_degradeTransf gradeTransf_gg' .
                have [:blurb] fSol_o_f2Sol_fSol_o_f2'Sol :=
                  (projT2 (@solveTransfCoMod_TransfCoMod_PolyMorCoMod_Pre
                      f2Solf2'Sol ( << f1Sol ,CoMod f2Sol >> %sol) blurb)));
                    first by abstract tac degradeTransf gradeTransf gg' .
                eexists. eexists. refine ( << fSol_o_f1Sol_fSol_o_f1'Sol</pre>
                                           ,^CoMod fSol_o_f2Sol_fSol_o_f2'Sol >> )%sol .
         }
Defined.
Arguments solveTransfCoMod_TransfCoMod_PolyMorCoMod_Pre
           !len _ _ _ !g'Sol_gSol _ _ _ : <u>simpl</u> nomatch .
Notation "f _o>CoMod^ g'g @ gradeTransf_gg'" :=
 (@solveTransfCoMod_TransfCoMod_PolyMorCoMod_Pre _ _ _ _ g'g _ f gradeTransf_gg')
   (<u>at</u> level 40 , g'g <u>at</u> next level) : sol_scope.
```

```
Fixpoint solveTransfCoMod TransfCoMod PolyMorCoMod PreP len {struct len} :
  forall (F G : obCoMod) (g'Sol gSol : 'morCoMod(0 F ~> G )0%sol)
    (g'Sol_gSol : 'transfCoMod(0 g'Sol ~> gSol )0%sol)
    (E : obCoMod) (fSol : 'morCoMod(0 E ~> F )0%sol)
    (gradeTransf gg' : (gradeTransf (Sol.toPolyMor fSol
                                o>CoMod^ Sol.toPolyTransf g'Sol gSol) <= len)%coq nat),
    ( ( (solveMorCoMod0 (Sol.toPolyMor fSol o>CoMod Sol.toPolyMor g'Sol)
          = projT1(fSol _o>CoMod^ g'Sol_gSol @ gradeTransf_gg')%sol) *
         (solveMorCoMod0 (Sol.toPolyMor fSol o>CoMod Sol.toPolyMor gSol)
          = projT1(projT2(fSol _o>CoMod^ g'Sol gSol @ gradeTransf_gg')%sol)) ) *
(Sol.toPolyTransf (projT2(projT2(fSol o>CoMod^ g'Sol @Sol @ gradeTransf gg')%sol))
     <~~2 (Sol.toPolyMor fSol_o>CoMod^ Sol.toPolyTransf g'Sol_gSol)%poly ) )%type.
Proof.
  <u>case</u> : len => [ | len ].
  (* len is 0 *)
  - ( move => F G g'Sol gSol g'Sol_gSol E fSol gradeTransf_gg' );
      exfalso; abstract tac_degradeTransf gradeTransf_gg'.
  (* len is (S len) *)
  (* gg' is f _o>CoMod^ g'g , to (fSol _o>CoMod^ g'Sol_gSol) *)
  - move => F G g'Sol gSol g'Sol_gSol E fSol gradeTransf_gg' .
    { destruct g'Sol_gSol as
           [ F G gSol (* @'UnitTransfCoMod gSol *)
           | F1 F2 Z1 z1Sol z1'Sol z1Solz1'Sol (* ~_1 @ F2 _o>CoMod^ z1Solz1'Sol *)
           F1 F2 Z2 z2Sol z2'Sol z2Solz2'Sol (* ~_2 @ F1 _o>CoMod^ z2Solz2'Sol *)
           L F1 F2 f1Sol f1'Sol f1Solf1'Sol f2Sol f2'Sol f2Solf2'Sol
(* << f1Solf1'Sol ,^CoMod f2Solf2'Sol >> *) ] .
      (* gg' is f _o>CoMod^ g'g , to (fSol _o>CoMod^ g'Sol_gSol) , is (fSol
      _o>CoMod^ @'UnitTransfCoMod gSol) *)
* have fSol_o_gSol_prop :=
      (@solveMorCoMod0P _ _ ( (Sol.toPolyMor fSol) o>CoMod (Sol.toPolyMor gSol) )).
   set fSol_o_gSol := (@solveMorCoMod0 _ _
         ( (Sol.toPolyMor fSol) o>CoMod (Sol.toPolyMor gSol) )) in fSol o gSol prop.
        subst fSol_o_gSol; move: fSol_o_gSol_prop;
           clear ; abstract tac_reduce_solveMorCoMod0.
      (* gg' is f _o>CoMod^ g'g , to (fSol _o>CoMod^ g'Sol_gSol) , is (fSol
       o>CoMod^ ~ 1 @ F2 o>CoMod^ z1Solz1'Sol) *)
      * <u>move</u>: (Sol.Destruct_codPair.morCoMod_codPairP fSol) => fSol_codPairP.
         { destruct fSol codPairP as
               [ F1 F2 (* (@'UnitMorCoMod (Pair F1 F2) )%sol *)
                | _F1 _F2 _Z1 _Z1' _z1 (* ( ~_1 @ _F2 o>CoMod _z1 )%sol *)
| _F1 _F2 _Z2 _Z2' _z2 (* ( ~_2 @ _F1 o>CoMod _z2 )%sol *)
                L F1 F2 f1 f2 (* ( << f1 ,CoMod f2 >> )%sol *) ] .
          (* gg' is f _o>CoMod^ g'g , to (fSol _o>CoMod^ g'Sol_gSol) , is
( @'UnitMorCoMod (Pair F1 F2) _o>CoMod^ ~_1 @ _ _o>CoMod^ z1Solz1'Sol ) *)

    <u>clear</u>; abstract tac_reduce_solveMorCoMod0.

           (* gg' is f _o>CoMod^ g'g , to (fSol _o>CoMod^ g'Sol_gSol) , is ( ( ~_1
           @ _F2 o>CoMod _z1 ) _o>CoMod^ ( ~_1 @ _ _o>CoMod^ z1Solz1'Sol ) ) *)
           - <u>simpl</u>; <u>set</u> same_blurb := ( <u>_ gradeTransf_gg'</u> : ( <u>_</u> <= len )%coq_nat) .
             move: (@solveTransfCoMod_TransfCoMod_PolyMorCoMod_PreP
                          _ ( ~_1 @ _ _o>CoMod^ z1Solz1'Sol )%sol _ _z1 same_blurb ) .
             clear ; abstract tac_reduce_solveMorCoMod0.
           (* gg' is f _o>CoMod^ g'g , to (fSol _o>CoMod^ g'Sol_gSol) , is ( ( ~_2
           @ _F1 o>CoMod _z2 ) _o>CoMod^ ( ~_1 @ _ _o>CoMod^ z1Solz1'Sol ) ) *)
            simpl; set same_blurb := ( _ gradeTransf_gg' : ( _ <= len )%coq_nat) .
move: (@solveTransfCoMod_TransfCoMod_PolyMorCoMod_PreP</pre>
                         _ _ ( \sim_1 @ _ _o>CoMod^ z1\overline{S}olz1'Sol )%\overline{S}ol _ _z2 same_blurb) .
             clear ; abstract tac_reduce_solveMorCoMod0.
           (* gg' is f _o>CoMod^ g'g , to (fSol _o>CoMod^ g'Sol_gSol) , is ( << f1
```

```
,CoMod f2 >> _o>CoMod^ ( ~_1 @ _ _o>CoMod^ z1Solz1'Sol ) ) *)
       - <u>simpl</u>; <u>set</u> same_blurb := ( <u>_ gradeTransf_gg'</u> : ( <u>_</u> <= len )%coq_nat) .
         move: (@solveTransfCoMod TransfCoMod PolyMorCoMod PreP
                    len \underline{\phantom{a}} \underline{\phantom{a}} \underline{\phantom{a}} z1Solz1'Sol \underline{\phantom{a}} f1 same blurb).
         clear; abstract tac reduce solveMorCoModO.
    }
  (* gg' is f _o>CoMod^ g'g , to (fSol _o>CoMod^ g'Sol_gSol) , is (fSol _o>CoMod^ ~_2 @ F1 _o>CoMod^ z2Solz2'Sol) *)
  * move: (Sol.Destruct_codPair.morCoMod_codPairP fSol) => fSol_codPairP.
     { <u>destruct</u> fSol_codPairP <u>as</u>
            [ F1 F2 (* (@'UnitMorCoMod (Pair F1 F2) )%sol *)
            | _F1 _F2 _Z1 _Z1' _z1 (* ( ~_1 @ _F2 o>CoMod _z1 )%sol *)
| _F1 _F2 _Z2 _Z2' _z2 (* ( ~_2 @ _F1 o>CoMod _z2 )%sol *)
            L F1 F2 f1 f2 (* ( << f1 ,CoMod f2 >> )%sol *) ] .
       (* gg' is f _o>CoMod^ g'g , to (fSol _o>CoMod^ g'Sol_gSol) , is
      (@'UnitMorCoMod (Pair F1 F2) _o>CoMod^ ~_2 @ _ _o>CoMod^ z2Solz2'Sol ) *)

    <u>clear</u>; abstract tac_reduce_solveMorCoMod0.

       (* gg' is f _o>CoMod^ g'g , to (fSol _o>CoMod^ g'Sol_gSol) , is ( ( ~_1
       @ _F2 o>CoMod _z1 ) _o>CoMod^ ( ~_2 @ _ _o>CoMod^ z2Solz2'Sol ) ) *)
       - <u>simpl</u>; <u>set</u> same_blurb := ( <u>_ gradeTransf_gg'</u> : ( <u>_</u> <= len )%coq_nat) .
         move: (@solveTransfCoMod_TransfCoMod_PolyMorCoMod_PreP
             len _ _ _ ( ~_2 @ _ _o>CoMod^ z2Solz2'Sol )%sol _ _z1 same_blurb).
         clear ; abstract tac_reduce_solveMorCoMod0.
       (* gg' is f _o>CoMod^ g'g , to (fSol _o>CoMod^ g'Sol_gSol) , is ( ( ~_2
       @ _F1 o>CoMod _z2 ) _o>CoMod^ ( ~_2 @ _ _o>CoMod^ z2Solz2'Sol ) ) *)
       - <u>simpl</u>; <u>set</u> same_blurb := ( <u>_ gradeTransf_gg'</u> : ( <u>_</u> <= len )%coq_nat) .
         move: (@solveTransfCoMod_TransfCoMod_PolyMorCoMod_PreP
                   _ _ _ ( ~_2 @ _ _o>CoMod^ z2Solz2'Sol )%sol _ _z2 same_blurb ).
         clear ; abstract tac_reduce_solveMorCoMod0.
       (* gg' is f _o>CoMod^ g'g , to (fSol _o>CoMod^ g'Sol_gSol) , is ( << f1
       ,CoMod f2 >> _o>CoMod^ ( ~_2 @ _ _o>CoMod^ z2Solz2'Sol ) ) *)
       - <u>simpl</u>; <u>set</u> same_blurb := ( <u>_ gradeTransf_gg'</u> : ( _ <= len )%coq_nat) .
         move: (@solveTransfCoMod_TransfCoMod_PolyMorCoMod_PreP
                              _ z2Solz2'Sol _ f2 same blurb).
         clear ; abstract tac_reduce_solveMorCoMod0.
    }
  (* gg' is f _o>CoMod^ g'g , to (fSol _o>CoMod^ g'Sol_gSol) , is (fSol
   _o>CoMod^ << f1Solf1'Sol ,^CoMod f2Solf2'Sol >> ) *)
  * { <u>destruct</u> fSol <u>as</u>
            [ F (* @'UnitMorCoMod F *)
               _F1 _F2 Z1 z1Sol (* ~_1 @ F2 o>CoMod z1Sol *)
_F1 _F2 Z2 z2Sol (* ~_2 @ F1 o>CoMod z2Sol *)
            | _F1
            | L _F1 _F2 _f1Sol _f2Sol (* << _f1Sol ,CoMod _f2Sol >> *) ] .
       (* gg' is f _o>CoMod^ g'g , to (fSol _o>CoMod^ g'Sol_gSol) , is
       ( @'UnitMorCoMod F _o>CoMod^ << f1Solf1'Sol ,^CoMod f2Solf2'Sol >> ) *)

    <u>clear</u>; abstract tac_reduce_solveMorCoMod0.

       (* qq' is f o>CoMod^ q'q , to (fSol o>CoMod^ q'Sol qSol) , is ( ( ~ 1
    @ F2 o>CoMod z1Sol ) o>CoMod^ << f1Solf1'Sol ,^CoMod f2Solf2'Sol >> ) *)
       - <u>simpl</u>; <u>set</u> same_blurb := ( <u>_ gradeTransf_gg'</u> : ( <u>_</u> <= len )%coq_nat) .
         move: (@solveTransfCoMod_TransfCoMod_PolyMorCoMod_PreP
len \ \_ \ \_ \ \_ \ ( << f1Solf1'Sol \ , ^CoMod \ f2Solf2'Sol >> ) \\ %sol \ \_ \ z1Sol \ same\_blurb) \ .
         clear ; abstract tac_reduce_solveMorCoMod0.
 (* gg' is f _o>CoMod^ g'g , to (fSol _o>CoMod^ g'Sol_gSol) , is ( ( ~_2 @ _F1 o>CoMod z2Sol ) _o>CoMod^ ( << f1Solf1'Sol ,^CoMod f2Solf2'Sol >> ) ) *)
       - simpl; set same_blurb := ( _ gradeTransf_gg' : ( _ <= len )%coq_nat) .
move: (@solveTransfCoMod_TransfCoMod_PolyMorCoMod_PreP</pre>
len \ \_ \ \_ \ \_ \ ( << f1Solf1'Sol \ , ^CoMod \ f2Solf2'Sol >> ) \\ %sol \ \_ \ z2Sol \ same\_blurb) \ .
         clear ; abstract tac_reduce_solveMorCoMod0.
```

```
(* gg' is f _o>CoMod^ g'g , to (fSol _o>CoMod^ g'Sol_gSol) , is ( <<
   f1Sol ,CoMod f2Sol >> o>CoMod^ ( << f1Solf1'Sol ,^CoMod f2Solf2'Sol >> ) ) *)
             * <u>simpl</u>; <u>set</u> same_blurbl := ( _ gradeTransf_gg' : ( _ <= len )%coq_nat).
               move: (@solveTransfCoMod_TransfCoMod_PolyMorCoMod_PreP
         _ f2Solf2'Sol _ ( << _f1Sol ,CoMod _f2Sol >> %sol) same_blurb2).
               clear ; abstract tac reduce solveMorCoMod0.
        }
0ed.
Fixpoint solveTransfCoMod TransfCoMod PolyMorCoMod Post len {struct len} :
  forall (G H : obCoMod) (hSol : 'morCoMod(0 G ~> H )0%sol)
    (F : obCoMod) (gSol g'Sol : 'morCoMod(0 F <math>\sim G )0%sol)
    (g'Sol_gSol : 'transfCoMod(0 g'Sol ~> gSol )0%sol)
    (gradeTransf_gg' : (gradeTransf (Sol.toPolyTransf g'Sol_gSol
                                       ^o>CoMod Sol.toPolyMor hSol) <= len)%coq nat),
    {g_ : 'morCoMod(0 F ~> H )0%sol &
           \{g'_{}: \text{'morCoMod(0 F} \sim> \text{H })0\%\text{sol }\&
                  'transfCoMod(0 g_ ~> g'_ )0%sol } }.
Proof.
  <u>case</u> : len => [ | len ].
  (* len is 0 *)
  - ( move => G H hSol F gSol g'Sol g'Sol gSol gradeTransf_gg' );
      exfalso; abstract tac_degradeTransf gradeTransf_gg'.
  (* len is (S len) *)
  (* gg' is g'g ^o>CoMod_ h , to (g'Sol_gSol ^o>CoMod_ hSol) *)
  - move => G H hSol F gSol g'Sol g'Sol_gSol gradeTransf_gg'.
    { <u>destruct</u> g'Sol_gSol <u>as</u>
           [ F G gSol (* @'UnitTransfCoMod gSol *)
            F1 F2 Z1 z1Sol z1'Sol z1Solz1'Sol (* ~_1 @ F2 _o>CoMod^ z1Solz1'Sol *)
F1 F2 Z2 z2Sol z2'Sol z2Solz2'Sol (* ~_2 @ F1 _o>CoMod^ z2Solz2'Sol *)
           L F1 F2 f1Sol f1'Sol f1Solf1'Sol f2Sol f2'Sol f2Solf2'Sol
(* << f1Solf1'Sol ,^CoMod f2Solf2'Sol >> *) ] .
      (* gg' is g'g ^o>CoMod_ h , to (g'Sol_gSol ^o>CoMod_ hSol) , is
      (@'UnitTransfCoMod gSol ^o>CoMod hSol) *)
      * have gSol_o_hSol_prop :=
        gsolveMorCoMod0P _ _ ( (Sol.toPolyMor gSol) o>CoMod (Sol.toPolyMor hSol) )).
set gSol_o_hSol := (@solveMorCoMod0 _ _
      (@solveMorCoMod0P
        ( (Sol.toPolyMor gSol) o>CoMod (Sol.toPolyMor hSol) )) in gSol o hSol prop.
        eexists. eexists. refine ( @'UnitTransfCoMod gSol_o hSol )%sol .
      (* gg' is g'g ^o>CoMod_ h , to (g'Sol_gSol ^o>CoMod_ hSol) , is (~_1 @ F2 _o>CoMod^ z1Solz1'Sol ^o>CoMod_ hSol) *)
      * have [:blurb] z1Sol_o_hSol_z1'Sol_o_hSol :=
           (projT2 (@solveTransfCoMod_TransfCoMod_PolyMorCoMod_Post
                               len _ _ hSol _ _ _ z1Solz1'Sol blurb )));
             first by abstract tac_degradeTransf gradeTransf_gg' .
        <u>eexists</u>. <u>eexists</u>.
        refine ( ~_1 @ F2 _o>CoMod^ z1Sol_o_hSol_z1'Sol_o_hSol )%sol .
      (* gg' is g'g ^o>CoMod_ h , to (g'Sol_gSol ^o>CoMod_ hSol) , is (~_2 @ F1
      o>CoMod^ z2Solz2'Sol \(^o>CoMod \quad hSol) \(^*\)
      * <u>have</u> [:blurb] z2Sol_o_hSol_z2'Sol_o_hSol :=
           (projT2 (projT2 (@solveTransfCoMod TransfCoMod PolyMorCoMod Post
                               len _ _ hSol _ _ _ z2Solz2'Sol blurb)));
            first by abstract tac_degradeTransf gradeTransf_gg' .
        eexists. eexists. refine ( ~_2 @ F1 _o>CoMod^ z2Sol_o_hSol_z2'Sol_o_hSol )%sol .
      (* gg' is g'g ^o>CoMod_ h , to (g'Sol_gSol ^o>CoMod_ hSol) , is ( <<
```

```
f1Solf1'Sol ,^CoMod f2Solf2'Sol >> ^o>CoMod hSol ) *)
      * move: (Sol.Destruct domPair.morCoMod domPairP hSol) => hSol domPairP.
        { destruct hSol domPairP as
               [ F1 F2 (* ( @'UnitMorCoMod (Pair F1 F2) )%sol *)
               | F1 F2 Z1 z1 (* ( ~_1 @ F2 o>CoMod z1 )%sol *)
| F1 F2 Z2 z2 (* ( ~_2 @ F1 o>CoMod z2 )%sol *)
| M M' F1 F2 f1 f2 (* ( << f1 ,CoMod f2 >> )%sol *) ] .
          (* gg' is g'g ^o>CoMod_ h , to (g'Sol\_gSol ^o>CoMod_ hSol) , is ( <<
      f1Solf1'Sol ,^CoMod f2Solf2'Sol >> ^o>CoMod_ @'UnitMorCoMod (Pair F1 F2) ) *)
          - eexists. eexists. refine ( << f1Solf1 Sol ,^CoMod f2Solf2'Sol >> )%sol .
          (* gg' is g'g ^o>CoMod_ h , to (g'Sol_gSol ^o>CoMod_ hSol) , is ( <<
          f1Solf1'Sol ,^CoMod f2Solf2'Sol >> ^o>CoMod_ ~_1 @ F2 o>CoMod z1 *)
          - have [:blurb] f1Sol o z1 f1'Sol o z1 :=
               (projT2 (projT2 (@solveTransfCoMod_TransfCoMod PolyMorCoMod Post
                                   len _ _ z1 _ _ _ f1Solf1'Sol blurb)));
                 first by abstract tac_degradeTransf gradeTransf_gg' .
            eexists. eexists. refine ( f1Sol_o_z1_f1'Sol_o_z1 )%sol .
          (* gg' is g'g ^o>CoMod_ h , to (g'Sol_gSol ^o>CoMod_ hSol) , is ( <<
          f1Solf1'Sol ,^CoMod f2Solf2'Sol >> ^o>CoMod_ ~_2 @ F1 o>CoMod z2 *)
          - have [:blurb] f2Sol_o_z2_f2'Sol_o_z2 :=
               (\texttt{projT2} \ (\texttt{@solveTransfCoMod\_TransfCoMod\_PolyMorCoMod\_Post})
                                   len _ _ z2 _ _ _ f2Solf2'Sol blurb )));
                 first by abstract tac_degradeTransf gradeTransf_gg' .
            eexists. eexists. refine ( f2Sol_o_z2_f2'Sol_o_z2 )%sol .
          (* gg' is g'g ^o>CoMod_ h , to (g'Sol_gSol ^o>CoMod_ hSol) , is ( <<
          f1Solf1'Sol ,^CoMod f2Solf2'Sol >> ^o>CoMod << f1 ,CoMod f2 >> *)
          - have [:blurb1] g'Sol_o_f1_gSol_o_f1 :=
               (projT2 (projT2 (@solveTransfCoMod_TransfCoMod_PolyMorCoMod_Post
        len _ _ f1 _ _ _ ( << f1Solf1'Sol ,^CoMod f2Solf2'Sol >> \$sol ) blurb1 )));
                 first by abstract tac degradeTransf gradeTransf gg' .
            have [:blurb2] g'Sol_o_f2_gSol_o_f2 :=
               (projT2 (@solveTransfCoMod TransfCoMod PolyMorCoMod Post
                       _ _ ( << f1Solf1'Sol ,^CoMod f2Solf2'Sol >> %sol ) blurb2 )));
                first by abstract tac_degradeTransf gradeTransf_gg' .
            <u>eexists</u>. <u>eexists</u>.
            <u>refine</u> ( << g'Sol o f1 gSol o f1 ,^CoMod g'Sol o f2 gSol o f2 >> )%sol.
        }
Defined.
Arguments solveTransfCoMod_TransfCoMod_PolyMorCoMod_Post
          !len _ _ _ _ !g'Sol_gSol _ : <u>simpl</u> nomatch .
Notation "g'g ^o>CoMod_ h @ gradeTransf_g'g" :=
  (@solveTransfCoMod_TransfCoMod_PolyMorCoMod_Post
    ___ h __ g'g gradeTransf_g'g)
(at level 40 , h at next level) : sol_scope.
Fixpoint solveTransfCoMod_TransfCoMod_PolyMorCoMod_PostP len {struct len} :
  forall (G H : obCoMod) (hSol : 'morCoMod(0 G ~> H )0%sol)
    (F : obCoMod) (gSol g'Sol : 'morCoMod(0 F ~> G )0%sol)
    (g'Sol_gSol: 'transfCoMod(0 g'Sol ~> gSol )0%sol)
    (gradeTransf_gg' : (gradeTransf (Sol.toPolyTransf g'Sol_gSol
                                       ^o>CoMod_ Sol.toPolyMor hSol) <= len)%coq_nat),
    ( ( (solveMorCoMod0 (Sol.toPolyMor g'Sol o>CoMod Sol.toPolyMor hSol)
         = projT1(g'Sol_gSol ^o>CoMod_ hSol @ gradeTransf_gg')%sol) *
        (solveMorCoMod0 (Sol.toPolyMor gSol o>CoMod Sol.toPolyMor hSol)
         = projT1(projT2(g'Sol_gSol ^o>CoMod_ hSol @ gradeTransf_gg')%sol)) ) *
      ( Sol.toPolyTransf (projT2(projT2(g'Sol_gSol ^o>CoMod_ hSol
                                                       @ gradeTransf_gg')%sol))
     <~~2 (Sol.toPolyTransf g'Sol_gSol ^o>CoMod_ Sol.toPolyMor hSol)%poly ) )%type.
Proof.
```

```
case : len => [ | len ].
  (* len is 0 *)
  - ( move => G H hSol F qSol q'Sol q'Sol qSol qradeTransf qq' );
      exfalso; abstract tac degradeTransf gradeTransf gg'.
  (* len is (S len) *)
  (* gg' is g'g ^o>CoMod_ h , to (g'Sol_gSol ^o>CoMod_ hSol) *)
  - move => G H hSol F gSol g'Sol g'Sol gSol gradeTransf_gg'.
    { destruct g'Sol_gSol as
          [ F G gSol (* @'UnitTransfCoMod gSol *)
           | F1 F2 Z1 z1Sol z1'Sol z1Solz1'Sol (* ~_1 @ F2 _o>CoMod^ z1Solz1'Sol *)
           F1 F2 Z2 z2Sol z2'Sol z2Solz2'Sol (* ~ 2 @ F1 _o>CoMod^ z2Solz2'Sol *)
           | L F1 F2 f1Sol f1'Sol f1Solf1'Sol f2Sol f2'Sol f2Solf2'Sol
(* << f1Solf1'Sol ,^CoMod f2Solf2'Sol >> *) ] .
      (* gg' is g'g ^o>CoMod h , to (g'Sol gSol ^o>CoMod hSol) , is
      (@'UnitTransfCoMod gSol ^o>CoMod hSol) *)
      * have gSol_o_hSol_prop :=
      (@solveMorCoMod0P _ _ ( (Sol.toPolyMor gSol) o>CoMod (Sol.toPolyMor hSol) )).
        o>CoMod (Sol.toPolyMor hSol) )) <u>in</u> gSol_o_hSol_prop.
        subst gSol_o_hSol; move: gSol_o_hSol_prop;
          clear ; abstract tac_reduce_solveMorCoMod0.
      (* gg' is g'g ^o>CoMod_ h , to (g'Sol_gSol ^o>CoMod_ hSol) , is (~_1 @ F2
       o>CoMod^ z1Solz1'Sol ^o>CoMod hSol) *)
      * <u>simpl</u>; <u>set</u> same_blurb := ( <u>_</u> gradeTransf_gg' : ( _ <= len )%coq_nat) .
        move: (@solveTransfCoMod_TransfCoMod_PolyMorCoMod_PostP
                  len _ _ hSol _ _ _ z1Solz1'Sol same_blurb );
          clear ; abstract tac_reduce_solveMorCoMod0.
      (* gg' is g'g ^o>CoMod_ h , to (g'Sol_gSol ^o>CoMod_ hSol) , is (~_2 @ F1
      _o>CoMod^ z2Solz2'Sol ^o>CoMod_ hSol) *)
* simpl; set same_blurb := ( _ gradeTransf_gg' : ( _ <= len )%coq_nat) .
        move: (@solveTransfCoMod TransfCoMod PolyMorCoMod PostP
                  len _ _ hSol _ _ _ z2Solz2'Sol same_blurb).
        clear ; abstract tac_reduce_solveMorCoMod0.
      (* gg' is g'g ^o>CoMod_ h , to (g'Sol_gSol ^o>CoMod_ hSol) , is ( <<
      f1Solf1'Sol ,^CoMod f2Solf2'Sol >> ^o>CoMod hSol ) *)
      * move: (Sol.Destruct domPair.morCoMod domPairP hSol) => hSol domPairP.
        { destruct hSol domPairP as
               [ F1 F2 (* (@'UnitMorCoMod (Pair F1 F2) )%sol *)
               | F1 F2 Z1 z1 (* ( ~_1 @ F2 o>CoMod z1 )%sol *)
| F1 F2 Z2 z2 (* ( ~_2 @ F1 o>CoMod z2 )%sol *)
| M M' F1 F2 f1 f2 (* ( << f1 ,CoMod f2 >> )%sol *) ] .
          (* gg' is g'g ^o>CoMod_ h , to (g'Sol_gSol ^o>CoMod_ hSol) , is ( <<
      f1Solf1'Sol ,^CoMod f2Solf2'Sol >> ^o>CoMod_ @'UnitMorCoMod (Pair F1 F2) ) *)

    <u>clear</u>; abstract tac_reduce_solveMorCoMod0.

          (* gg' is g'g ^o>CoMod_ h , to (g'Sol\_gSol ^o>CoMod_ hSol) , is ( <<
          f1Solf1'Sol ,^CoMod f2Solf2'Sol >> ^o>CoMod_ ~_1 @ F2 o>CoMod z1 *)
          - <u>simpl</u>; <u>set</u> same_blurb := ( <u>_ gradeTransf_gg'</u> : ( _ <= len )%coq_nat) .
            move: (@solveTransfCoMod_TransfCoMod_PolyMorCoMod_PostP
                       len _ z1 _ _ f1Solf1'Sol same_blurb).
            clear ; abstract tac_reduce_solveMorCoMod0.
          (* gg' is g'g ^o>CoMod_ h , to (g'Sol_gSol ^o>CoMod_ hSol) , is ( <<
          f1Solf1'Sol ,^CoMod f2Solf2'Sol >> ^o>CoMod_ ~_2 @ F1 o>CoMod z2 *)
- simpl; set same_blurb := ( _ gradeTransf_gg' : ( _ <= len )%coq_nat) .
move: (@solveTransfCoMod_TransfCoMod_PolyMorCoMod_PostP</pre>
                       len _ z2 _ _ f2Solf2'Sol same_blurb ).
            clear ; abstract tac_reduce_solveMorCoMod0.
          (* gg' is g'g ^o>CoMod_ h , to (g'Sol_gSol ^o>CoMod_ hSol) , is ( <<
```

```
f1Solf1'Sol ,^CoMod f2Solf2'Sol >> ^o>CoMod << f1 ,CoMod f2 >> *)
                     - <u>simpl</u>; <u>set</u> same blurb1 := ( <u>gradeTransf_gg'</u> : ( <u><=</u> len )%coq nat) .
                        move: (@solveTransfCoMod TransfCoMod PolyMorCoMod PostP
                        ______fl____( << flSolf1'Sol ,^CoMod f2Solf2'Sol >> %sol ) same_blurb1).
____simpl; set same_blurb2 := ( __ gradeTransf_gg' : ( __ <= len )%coq_nat) .
____move: (@solveTransfCoMod_TransfCoMod_PolyMorCoMod_PostP
            len \_ f2 \_ \_ ( << f1Solf1'Sol ,^CoMod f2Solf2'Sol >> %sol ) same blurb2).
                        <u>clear</u>; abstract tac reduce solveMorCoMod0.
0ed.
Fixpoint solveTransfCoMod len {struct len} :
    forall (F G : obCoMod) (g g' : 'morCoMod(0 F <math>\sim G )0 )
        (gg' : 'transfCoMod(0 g \sim> g' )0 ),
    forall gradeTransf_gg' : (gradeTransf gg' <= len)%coq_nat,</pre>
        { g_g'_: \{ g_: \text{'morCoMod(0 F } \sim \text{S }) \text{0 } \text{%sol } \& \text{}
                            _ : 'morCoMod(0 F ~> G )0 %sol & 'transfCoMod(0 g_ ~> g'_ )0 %sol } }
        | ( ( solveMorCoMod0 g = projT1 g_g'_ ) *
                ( solveMorCoMod0 g' = projT1 (projT2 g_g'_) ) )%type }.
Proof.
   <u>case</u> : len => [ | len ].
    (* len is 0 *)
    - ( <u>move</u> => F G g g' gg' gradeTransf_gg' ); <u>exfalso</u>; abstract tac_degradeTransf gradeTransf_gg'.
    (* len is (S len) *)
    - <u>move</u> => F G g g' gg'; <u>case</u> : F G g g' / gg' =>
        F G g g' g'g g'0 g'' g''g' eqMor (* g''g' o^CoMod g'g *)
        | F G g g' g'g E f (* f _o>CoMod^ g'g *)
         | G H h F g g ' g 'g (* g 'g ^o>CoMod h *)
        | F G g (* @'UnitTransfCoMod g *)
        | F1 F2 Z1 z1 z1' z1z1' (* ~_1 @ F2 _o>CoMod^ z1z1' *)
| F1 F2 Z2 z2 z2' z2z2' (* ~_2 @ F1 _o>CoMod^ z2z2' *)
| L F1 F2 f1 f1' f1f1' f2 f2' f2f2' (* << f1f1' ,^CoMod f2f2' >> *)
        ] gradeTransf gg' .
        (* gg' is g''g' o^CoMod g'g *)
        + all: cycle 1.
        (* gg' is f _o>CoMod^ g'g *)
        + all: cycle 1.
        (* gg' is g'g ^o>CoMod_ h *)
        + all: cycle 1.
        (* gg' is @'UnitTransfCoMod g *)
        + have gSol prop := (@solveMorCoModOP
            <u>set</u> gSol := (@solveMorCoMod0 _ _ g) \overline{in} gSol_prop.
            unshelve <a href="mailto:eexists">eexists</a>. <a href="mailto:eexists
            <u>clear</u>; abstract tac reduce solveMorCoMod0.
        (* gg' is ~_1 @ F2 _o>CoMod^ z1z1' *)
        + have [:blurb] z1Sol_z1'Sol_prop :=
                (proj2_sig (solveTransfCoMod len _
                                                                                                z1z1' blurb));
                    first by abstract tac_degradeTransf gradeTransf_gg'.
                                                                          _ _ z1z1' blurb) z1Sol_z1'Sol_prop
            move: (solveTransfCoMod len _ _
            => z1Sol_z1'Sol z1Sol_z1'Sol_prop.
            unshelve <u>eexists</u>. <u>eexists</u>. <u>eexists</u>.
            refine ( ~_1 @ F2 _o>CoMod^ (projT2 (projT2 (proj1_sig z1Sol_z1'Sol))))%sol.
            move: z1Sol_z1'Sol prop; clear; abstract tac reduce solveMorCoMod0.
        (* gg' is ~_2 @ F1 _o>CoMod^ z2z2' *)
        + have [:blurb] z2Sol_z2'Sol_prop :=
                    proj2_sig (solveTransfCoMod len _ _ _ z2z2' blurb)); first by abstract tac_degradeTransf gradeTransf_gg'.
                (proj2_sig (solveTransfCoMod len _ _
            move: (solveTransfCoMod len _ _ _ z2z2' blurb) z2Sol_z2'Sol_prop
            => z2Sol_z2'Sol z2Sol_z2'Sol_prop.
            unshelve <u>eexists</u>. <u>eexists</u>. <u>eexists</u>.
```

```
refine ( ~ 2 @ F1 o>CoMod^ (projT2 (projT2 (proj1 sig z2Sol z2'Sol))))%sol.
   move: z2Sol_z2'Sol prop; clear; abstract tac reduce solveMorCoMod0.
 (* gg' is << f1f1' ,^CoMod f2f2' >> *)
 + have [:blurb1] f1Sol f1'Sol prop :=
       roj2_sig (solveTransfCoMod len _ _ _ f1f1' blurb1));
first by abstract tac_degradeTransf gradeTransf_gg'.
     (proj2_sig (solveTransfCoMod len _ _
   move: (solveTransfCoMod len _ _ _ f1f1' blurb1) f1Sol_f1'Sol_prop
=> f1Sol_f1'Sol_f1'Sol_prop .
   have [:blurb] f2Sol f2'Sol prop :=
     (proj2_sig (solveTransfCoMod len_
                                                 f2f2' blurb));
       first by abstract tac degradeTransf gradeTransf gg'
   move: (solveTransfCoMod len _ _ _ f2f2' blurb) f2Sol_f2'Sol_prop
   => f2Sol f2'Sol f2Sol_f2'Sol_prop .
   unshelve <u>eexists</u>. <u>eexists</u>. <u>eexists</u>.
   refine ( << (projT2 (projT2 (proj1_sig f1Sol_f1'Sol)))</pre>
                 ,^CoMod (projT2 (projT2 (proj1_sig f2Sol_f2'Sol))) >> )%sol.
   move: f1Sol_f1'Sol_prop f2Sol_f2'Sol_prop;
     clear; abstract tac_reduce_solveMorCoMod0.
 (* gg' is g''g' o^CoMod g'g *)
 + <u>have</u> solveTransfCoMod len := solveTransfCoMod len. <u>clear</u> solveTransfCoMod.
   Definition solveTransfCoMod_sub_PolyTransfCoMod :
     forall (len : nat) (F G : obCoMod) (g g' : 'morCoMod(0 F ~> G )0)
       (g'g : 'transfCoMod(0 g' \sim g )0) (g'0 g'' : 'morCoMod(0 F \sim G )0)
       (g''g' : 'transfCoMod(0 g'' \sim g'0 )0) (eqMor : g'0 <\sim 1 g')
       (gradeTransf_gg' : (gradeTransf (g''g' o^CoMod g'g # eqMor)
                             <= len.+1)%coq nat)
       (solveTransfCoMod\_len : forall (F G : obCoMod) (g g' : 'morCoMod(0 F <math>\sim G )0)
                                    (gg' : 'transfCoMod(0 g \sim> g' )0),
            (gradeTransf gg' <= len)%coq_nat ->
            {g_g'_ : {g_ : 'morCoMod(0 F ~> G )0%sol &
                            \{g'\_: \mathsf{'morCoMod(0 F} \sim G)0\%sol &
                                    'transfCoMod(0 g_ ~> g'_ )0%sol}} |
             ((solveMorCoMod0 g = projT1 g_g'_) *
              (solveMorCoMod0 g' = projT1 (projT2 g_g'_)))%type}),
       \{g\_g'\_: \{g\_: \text{'morCoMod}(0 F \sim> G )0\%sol &
                        \{g'_{-}: \text{'morCoMod}(0 \text{ F} \sim> \text{G})0\%\text{sol }\&
                                'transfCoMod(0 g_ ~> g'_ )0%sol}} |
         ((solveMorCoMod0 g'' = projT1 g g')^{-}
          (solveMorCoMod0 g = projT1 (projT2 g_g'_)))%type}.
   Proof.
     intros; have [:blurb_] g''Sol_prop :=
         st ( ( (proj2_sig (solveTransfCoMod_len _ _ _ g''g' blurb_))))); first by abstract tac_degradeTransf gradeTransf_gg'.
       (fst ( ( (proj2_sig (solveTransfCoMod_len )
     move: (proj1_sig ( ((solveTransfCoMod_len _ _ _ g''g' blurb_))))
g''Sol_prop (snd ( (proj2_sig (solveTransfCoMod_len _ _ _ g''g' blurb_)))))
     => g''Sol_g'Sol g''Sol_prop g'Sol0_prop.
     have [:blurb'] g'Sol_prop :=
       (fst ( ( (proj2_sig (solveTransfCoMod_len _ _ _
                                                             _ g'g blurb')))));
         first by abstract tac_degradeTransf gradeTransf_gg'.
     move: (proj1_sig ( ( (solveTransfCoMod_len _ _ _ g'g blurb'))))
    g'Sol_prop (snd ( ( (proj2_sig (solveTransfCoMod_len _ _ _ g'g blurb')))))
     => g'Sol gSol g'Sol prop gSol prop.
     clear solveTransfCoMod len.
     have eqMor' : (Sol.toPolyMor (projT1 (projT2 g''Sol_g'Sol))
                                     <~>1 (Sol.toPolyMor (projT1 g'Sol_gSol))%poly)
       by abstract
             ( apply: EqMorCoMod.eq_convMorCoMod;
               rewrite -g'Sol0_prop -g'Sol_prop ;
               congr Sol.toPolyMor; congr solveMorCoMod0;
               apply: EqMorCoMod.convMorCoMod eq;
               <u>apply</u>: eqMor
             )
     clear g'Sol0_prop g'Sol_prop .
     have [:blurb] g''Sol_g'Sol_o'_g'Sol_gSol_prop :=
```

```
(fst (@solveTransfCoMod PolyTransfCoModP
               (gradeTransf ((Sol.toPolyTransf (projT2 (projT2 g''Sol_g'Sol)))
              o^CoMod (Sol.toPolyTransf (projT2 (projT2 g'Sol gSol))) # eqMor'))
                       (projT2 (projT2 g'Sol_gSol))
      (*memo: this argument is same eqMor' as earlier argument *));
        first by <u>clear</u>; abstract reflexivity.
    move: (@solveTransfCoMod_PolyTransfCoMod
              (gradeTransf ((Sol.toPolyTransf (projT2 (projT2 g''Sol_g'Sol)))
              o^CoMod (Sol.toPolyTransf (projT2 (projT2 g'Sol_gSol))) # eqMor'))
                     (projT2 (projT2 g'Sol_gSol))
                  (projT2 (projT2 g''Sol g'Sol)) eqMor' blurb )
             g''Sol_g'Sol_o'_g'Sol_gSol_prop
    => g''Sol_g'Sol_o'_g'Sol_gSol_g'Sol_o'_g'Sol_gSol_prop.
    exists g''Sol_g'Sol_o'_g'Sol_gSol.
    move: g''Sol prop gSol prop g''Sol g'Sol o' g'Sol gSol prop;
      clear; abstract tac_reduce_solveMorCoMod0.
  Defined.
  apply: (solveTransfCoMod_sub_PolyTransfCoMod gradeTransf_gg'
                                                  solveTransfCoMod_len).
(* gg' is f _o>CoMod^ g'g , to (fSol _o>CoMod^ g'Sol_gSol) *)
+ have solveTransfCoMod_len := solveTransfCoMod len. clear solveTransfCoMod.
  Definition solveTransfCoMod_sub_TransfCoMod_PolyMorCoMod Pre :
    <u>forall</u> (len : nat) (F G : obCoMod) (g g' : 'morCoMod(0 F \sim G )0)
      (g'g : 'transfCoMod(0 g' \sim g )0) (E : obCoMod)
      (f: 'morCoMod(0 E ~> F )0)
      (gradeTransf_gg' : (gradeTransf (f _o>CoMod^ g'g) <= len.+1)%coq_nat)</pre>
      (solveTransfCoMod\_len : forall (F G : obCoMod) (g g' : 'morCoMod(0 F ~> G )0)
                                  (gg' : 'transfCoMod(0 g ~> g' )0),
           (gradeTransf gg' <= len)%coq_nat ->
           {g_g'_ : {g_ : 'morCoMod(0 F ~> G )0%sol &
                           \{g'\_: \mathsf{'morCoMod}(0 \ \mathsf{F} \sim \mathsf{G} \ )0\%\mathsf{sol} \ \&
                                  'transfCoMod(0 g_ ~> g'_ )0%sol}} |
            ((solveMorCoMod0 g = projT1 g_g'_) *
            (solveMorCoMod0 g' = projT1 (projT2 g_g'_)))%type}),
      \{g\_g'\_: \{g\_: \text{'morCoMod}(0 E \sim> G)0\%sol &
                       {g'_: 'morCoMod(0 E ~> G )0%sol &
                               'transfCoMod(0 g_ ~> g'_ )0%sol}} |
       ((solveMorCoMod0 (f o>CoMod g') = proj\overline{T}1 g_{g'}) *
        (solveMorCoMod0 (f o>CoMod g) = projT1 (projT2 g_g'_)))%type}.
    intros. have [:blurb] g'Sol_gSol_prop :=
               (proj2_sig (solveTransfCoMod_len _ _
                 proj2_sig (solveTransfCoMod_len _ _ _ g'g blurb));
first by abstract tac_degradeTransf gradeTransf_gg'.
    move: (solveTransfCoMod_len _ _ _ g'g blurb) g'Sol_gSol_prop
    => g'Sol_gSol g'Sol_gSol_prop.
    have fSol_prop := (@solveMorCoMod0P
    set fSol := (@solveMorCoMod0 _ _ f) in fSol_prop.
    have [:blurb'] fSol_o_g'Sol_gSol_prop :=
      (fst (@solveTransfCoMod_TransfCoMod_PolyMorCoMod_PreP)
               (gradeTransf ((Sol.toPolyMor fSol)
        _o>CoMod^ (Sol.toPolyTransf (projT2 (projT2 (proj1_sig g'Sol_gSol))))))
                     _ (projT2 (projT2 (proj1_sig g'Sol_gSol)))                                fSol blurb'));
        first by <u>clear</u>; abstract reflexivity.
    move: (@solveTransfCoMod_TransfCoMod_PolyMorCoMod_Pre
              (gradeTransf ((Sol.toPolyMor fSol)
       _o>CoMod^ (Sol.toPolyTransf (projT2 (projT2 (proj1_sig g'Sol_gSol))))))
                      (projT2 (projT2 (proj1_sig g'Sol_gSol)))
               fSol blurb') fSol_o_g'Sol_gSol_prop
    => fSol_o_g'Sol_gSol fSol_o_g'Sol_gSol_prop.
    exists fSol_o_g'Sol_gSol .
    subst fSol; move: g'Sol_gSol_prop fSol_o_g'Sol_gSol_prop;
      clear; abstract tac_reduce_solveMorCoMod0.
  Defined.
```

```
apply: (solveTransfCoMod_sub_TransfCoMod_PolyMorCoMod_Pre
                 gradeTransf gg' solveTransfCoMod len).
    (* gg' is g'g ^o>CoMod h *)
    + have solveTransfCoMod len := solveTransfCoMod len. clear solveTransfCoMod.
      Definition solveTransfCoMod_sub_TransfCoMod_PolyMorCoMod_Post :
         forall (len : nat) (G H : obCoMod) (h : 'morCoMod(0 G ~> H )0)
           (F : obCoMod) (g g' : 'morCoMod(0 F ~> G )0)
           (g'g : 'transfCoMod(0 g' \sim> g)0)
           (gradeTransf gg' : (gradeTransf (g'g ^o>CoMod h) <= len.+1)%coq nat)
           (solveTransfCoMod\_len : forall (F G : obCoMod) (g g' : 'morCoMod(0 F <math>\sim S G)0)
                                        (gg' : 'transfCoMod(0 g ~> g' )0),
               (gradeTransf gg' <= len)%coq nat ->
               {g_g'_ : {g_ : 'morCoMod(0 F ~> G )0%sol &
                                {g'_ : 'morCoMod(0 F ~> G )0%sol &
                                        "transfCoMod(0 g\_ \sim> g"\_ )0\%sol\} \ |
                ((solveMorCoMod0 g = projT1 g_g'_) *
                 (solveMorCoModO g' = projT1 (projT2 g_g'_)))%type}),
           \{g\_g'\_: \{g\_: \text{'morCoMod(0 F } \sim \text{H })0\%\text{sol } \& \}
                           \{g'_{-}: \text{'morCoMod}(0 F \sim> H)0\%sol &
                                    'transfCoMod(0 g_ ~> g'_ )0%sol}} |
            ((solveMorCoMod0 (g' o>CoMod h) = projT1 g_g'_) *
             (solveMorCoMod0 (g o>CoMod h) = projT1 (projT2 g_g'_)))%type}.
      Proof.
        <u>intros</u>; <u>have</u> [:blurb] g'Sol gSol prop :=
           (proj2_sig (solveTransfCoMod_len _ _ _ g'g blurb));
  first by abstract tac_degradeTransf gradeTransf_gg'.
        move: (solveTransfCoMod_len _ _ _ g'g blurb) g'Sol_gSol_prop
         => g'Sol_gSol g'Sol_gSol_prop.
        have hSol_prop := (@solveMorCoMod0P
                                                   h).
        set hSol := (@solveMorCoMod0 _ _ h) in hSol_prop.
        have [:blurb'] g'Sol_gSol_o_hSol_prop :=
  (fst (@solveTransfCoMod_TransfCoMod_PolyMorCoMod_PostP)
           (gradeTransf ((Sol.toPolyTransf (projT2 (projT2 (proj1_sig g'Sol_gSol))))
                           ^o>CoMod_ (Sol.toPolyMor hSol)))
                          _ (projT2 (projT2 (proj1_sig g'Sol_gSol))) blurb'));
               hSol
             first by <u>clear</u>; abstract reflexivity.
         move: (@solveTransfCoMod_TransfCoMod_PolyMorCoMod_Post
           (gradeTransf ((Sol.toPolyTransf (projT2 (projT2 (proj1_sig g'Sol_gSol))))
                          ^o>CoMod_ (Sol.toPolyMor hSol))) _ _ hSol
           (projT2 (projT2 (proj1_sig g'Sol_gSol))) blurb') g'Sol_gSol_o_hSol_prop
        => g'Sol_gSol_o_hSol g'Sol_gSol_o_hSol_prop.
         exists g'Sol_gSol_o_hSol .
        move: g'Sol_gSol_prop g'Sol_gSol_o_hSol_prop;
           clear; abstract tac_reduce_solveMorCoMod0.
      Defined.
      apply: (solveTransfCoMod_sub_TransfCoMod_PolyMorCoMod_Post
                 gradeTransf_gg' solveTransfCoMod_len).
Defined.
Arguments solveTransfCoMod !len _ _ _ !gg' _ : <u>clear</u> implicits , <u>simpl</u> nomatch .
Lemma solveTransfCoMod len :
  forall len (F G : obCoMod) (g G' : 'morCoMod(0 F <math>\sim G )0)
    (gg' : 'transfCoMod(0 g ~> g' )0 )
    (gradeTransf_gg'_len : (gradeTransf gg' <= len)%coq_nat),
  forall len' (gradeTransf_gg'_len' : (gradeTransf gg' <= len')%coq_nat),</pre>
    @solveTransfCoMod len _ _ _ gradeTransf_gg'_len
= @solveTransfCoMod len' _ _ _ gradeTransf_gg'_len' .
  <u>induction</u> len <u>as</u> [ | len ].
  - ( move => ? ? ? ? gradeTransf_gg'_len ); exfalso;
     clear -gradeTransf_gg'_len; by abstract tac_degradeTransf gradeTransf_gg'_len.
  - intros. destruct len'.
```

```
+ exfalso; clear -gradeTransf gg' len';
         by abstract tac_degradeTransf gradeTransf_gg' len'.
    + destruct gg'; cycle 3.
       * reflexivity.
       * <u>simpl</u>. <u>erewrite</u> IHlen. reflexivity.
       * <u>simpl</u>. <u>erewrite</u> IHlen. reflexivity.
        simpl. unfold ssr_have.
         erewrite (IHlen _ _ _ gg'l).
       erewrite (IHlen _ _ _ gg'2). reflexivity.
* simpl. unfold ssr_have. unfold solveTransfCoMod_sub_PolyTransfCoMod.
         unfold ssr_have. erewrite (IHlen
                                                      gg'1).
                             _{-} _{-} gg'2). reflexivity.
         <u>erewrite</u> (IHlen
       * <u>simpl</u>. <u>unfold</u> ssr_have.
         <u>rewrite</u> [in LHS]/solveTransfCoMod_sub_TransfCoMod_PolyMorCoMod_Pre.
         <u>unfold</u> ssr have. <u>erewrite</u> IHlen. reflexivity.
       * <u>simpl</u>. <u>unfold</u> ssr_have.
         rewrite [in LHS]/solveTransfCoMod_sub_TransfCoMod_PolyMorCoMod_Post.
         <u>unfold</u> ssr_have. <u>erewrite</u> IHlen. reflexivity.
Qed.
Definition solveTransfCoMod0 :
  { g_g'_: \{ g_: \text{'morCoMod}(0 F \sim G) 0 \% \text{sol } \& \}
                       { g'_ : 'morCoMod(0 F ~> G )0 %sol &
                                'transfCoMod(0 g_ \sim> g'_ )0 %sol } 
    | ( ( solveMorCoMod0 g = projT1 g_g'_ ) *
         ( solveMorCoMod0 g' = projT1 (projT2 g_g'_) ) )%type }.
  intros; apply: (@solveTransfCoMod (gradeTransf gg') F G g g' gg'); constructor.
Defined.
Lemma solveTransfCoMod0 len :
  \underline{forall} len (F G : obCoMod) (g g' : 'morCoMod(0 F <math>\sim > G )0)
    (gg' : 'transfCoMod(0 g ~> g' )0 )
    (gradeTransf_gg'_len : (gradeTransf gg' <= len)%coq_nat),
    @solveTransfCoMod0
    @solveTransfCoMod0 _ _ _ gg' = @solveTransfCoMod len _ _ _ _
                                        _ gradeTransf_gg'_len.
Proof. intros. erewrite solveTransfCoMod_len. reflexivity. Qed.
Lemma solveTransfCoMod0_Project1_Transf :
  <u>forall</u> (F1 F2 Z1 : obCoMod) (z1 z1' : 'morCoMod(0 F1 ~> Z1 )0)
    (z1z1' : 'transfCoMod(0 z1 ~> z1' )0),
    (proj1_sig (solveTransfCoMod0 ( ~_1 @ F2 _o>CoMod^ z1z1' )%poly))
= (existT _ (~_1 o>CoMod (projT1 (proj1_sig (solveTransfCoMod0 z1z1'))))%sol
           (~ 1 o>CoMod (projT1 (projT2 (proj1_sig (solveTransfCoMod0 z1z1')))))%sol
    (\sim 1 _{o}\text{-CoMod}^{\circ} (\text{projT2} (\text{projT2} (\text{proj1}_{sig} (\text{solveTransfCoMod0} z1z1')))))%sol)).
Proof.
  <u>intros</u>. <u>rewrite</u> [solveTransfCoMod0 <u>in</u> LHS]lock.
  erewrite solveTransfCoMod0 len.
  rewrite -lock /solveTransfCoMod0. simpl. reflexivity.
Qed.
Lemma solveTransfCoMod0 PolyTransfCoMod :
  <u>forall</u> (F G : obCoMod) (g g' : 'morCoMod(0 F \sim G)0) (g'g : 'transfCoMod(0 g' \sim G)0)
    (q'0 q'' : \text{'morCoMod}(0 F \sim G)0) (q''q' : \text{'transfCoMod}(0 q'' \sim q'0)0)
    (eqMor: g'0 <~>1 g') eqMor',
    proj1 sig (solveTransfCoMod0 (g''g' o^CoMod g'g # eqMor)%poly)
    = ((projT2 (projT2 (proj1_sig (solveTransfCoMod0 g''g'))))
      o^CoMod (projT2 (projT2 (proj1_sig (solveTransfCoMod0 g'g)))) # eqMor')%sol.
Proof.
  <u>intros</u>. <u>rewrite</u> [solveTransfCoMod0 <u>in</u> LHS]lock. <u>move</u>: eqMor'.
  do 2 <u>erewrite</u> solveTransfCoMod0 len.
  <u>intros</u>. <u>erewrite</u> solveTransfCoMod_PolyTransfCoMod0_len.
  rewrite -lock /solveTransfCoMod0. simpl. reflexivity.
Qed.
(**ETC : ... *)
```

```
Fixpoint solveTransfCoModP len {struct len} :
<u>forall</u> (F G : obCoMod) (g g' : 'morCoMod(0 F \sim G) 0) (gg' : 'transfCoMod(0 g \sim G') 0),
forall gradeTransf_gg' : (gradeTransf gg' <= len)%coq_nat,</pre>
   (Sol.toPolyTransf (projT2 (projT2 (proj1 sig
                     (@solveTransfCoMod len _ _ _ gg' gradeTransf_gg'))))) <~~2 gg'.
Proof.
  <u>case</u> : len => [ | len ].
   (* len is 0 *)
   - ( move => F G g g' gg' gradeTransf_gg' );
       exfalso; abstract tac degradeTransf gradeTransf gg'.
  (* len is (S len) *)
   - move => F G g g' gg'; case : F G g g' / gg' =>
[ F G g g' g'g g'0 g'' g''g' eqMor (* g''g' o^CoMod g'g *)
| F G g g' g'g E f (* f _o>CoMod^ g'g *)
| G H h F g g' g'g (* g'g ^o>CoMod_ h *)
     | F G g (* @'UnitTransfCoMod g *)
     | F1 F2 Z1 z1 z1' z1z1' (* ~_1 @ F2 _o>CoMod^ z1z1' *)
| F1 F2 Z2 z2 z2' z2z2' (* ~_2 @ F1 _o>CoMod^ z2z2' *)
| L F1 F2 f1 f1' f1f1' f2 f2' f2f2' (* << f1f1' ,^CoMod f2f2' >> *)
     ] gradeTransf_gg' .
     (* gg' is g''g' o^CoMod g'g *)
     + all: cycle 1.
     (* gg' is f o>CoMod^ g'g *)
     + all: cycle 1.
     (* gg' is g'g ^o>CoMod_ h *)
     + all: cycle 1.
     (* gg' is @'UnitTransfCoMod g *)
     + have gSol_prop := (@solveMorCoMod0P _ _ g).
       \underline{\mathsf{set}} \ \mathsf{gSol} := \ (\mathsf{@solveMorCoMod0} \ \_ \ \underline{\mathsf{g}}) \ \underline{\mathsf{in}} \ \mathsf{gSol\_prop}.
       clear -gSol_prop; abstract tac_reduce_solveMorCoMod0.
     (* gg' is ~_1 @ F2 _o>CoMod^ z1z1' *)
     + <u>simpl</u>; <u>set</u> same_blurb := (<u>_</u> gradeTransf_gg' : ( _ <= len )%coq_nat) .
       move: (solveTransfCoModP len _ _ _ z1z1' same_blurb).
       clear; abstract tac_reduce_solveMorCoMod0.
     (* gg' is ~_2 @ F1 _o>CoMod^ z2z2' *)
     + <u>simpl</u>; <u>set</u> same_blurb := (<u>_</u> gradeTransf_gg' : ( _ <= len )%coq_nat) .
       move: (solveTransfCoModP len _ _ _ z2z2' same_blurb).
       clear; abstract tac_reduce_solveMorCoMod0.
     (* gg' is << f1f1' ,^CoMod f2f2' >> *)
     + <u>simpl</u>; <u>set</u> same_blurb1 := (<u>_ gradeTransf_gg'</u> : ( _ <= len )%coq_nat) .
       move: (solveTransfCoModP len _ _ _ f1f1' same_blurb1).
set same_blurb2 := ( ( _ gradeTransf_gg' _ _ ) ) : ( _ <= len)%coq_nat) .</pre>
       move: (solveTransfCoModP len _ _ _ f2f2' same_blurb2).
       clear; abstract tac_reduce_solveMorCoMod0.
     (* gg' is g''g' o^CoMod g'g *)
     + simpl; set same_blurb' := (_ gradeTransf_gg' : ( _ <= len )%coq_nat) .</pre>
       move: (solveTransfCoModP len _ _ _ g'g same_blurb').
set same_blurb := (_ gradeTransf_gg' : ( _ <= len )%coq_nat) .</pre>
       same_eqMor' same_blurb_refl same_eqMor' (*memo: same*) )).
       clear; abstract tac reduce solveMorCoMod0.
     (* gg' is f _o>CoMod^ g'g , to (fSol _o>CoMod^ g'Sol_gSol) *)
     + simpl; set same_blurb := (_ gradeTransf_gg' : ( _ <= len )%coq_nat) .</pre>
       move: (solveTransfCoModP len _ _ _ g'g same_blurb).
```

```
have fSol_prop := (@solveMorCoMod0P _ _ f).
       set fSol := (@solveMorCoMod0 _ _ f) in fSol_prop.
      move: fSol prop.
      set same_blurb_refl := (_ f _ : ( _ <= _ )%coq_nat) .
move: (@solveTransfCoMod_TransfCoMod_PolyMorCoMod_PreP</pre>
                                                                      same blurb refl).
       clear; abstract tac reduce solveMorCoMod0.
    (* gg' is g'g ^o>CoMod_ h *)
    + <u>simpl</u>; <u>set</u> same_blurb := (_ gradeTransf_gg' : ( _ <= len )%coq_nat) .
      move: (solveTransfCoModP len _ _ _ g'g same_blurb).
have hSol_prop := (@solveMorCoModOP _ _ h).
      set hSol := (@solveMorCoMod0 _ _ h) in hSol_prop.
      move: hSol prop.
                                              _ : ( _ <= _ )%coq nat) .
      set same_blurb_refl := (_ h _
      <u>move</u>: (@solveTransfCoMod_TransfCoMod_PolyMorCoMod_PostP _ _
                                                                       same_blurb_refl).
      clear; abstract tac_reduce_solveMorCoMod0.
Qed.
Lemma solveTransfCoMod0P :
forall (F G : obCoMod) (g g' : 'morCoMod(0 F \sim G )0 ) (gg' : 'transfCoMod(0 g \sim g' )0 ),
  (Sol.toPolyTransf (projT2 (projT2 (proj1_sig
                                        (@solveTransfCoMod0 _ _ _ gg'))))) <~~2 gg' .
Proof. intros. apply: solveTransfCoModP . Qed.
End Resolve.
End Transf.
End TWOFOLD.
```

Voila.