

# EGGHUNTING WORKSHOP

## IT SECURITY MEETUP KASSEL / NORDHESSEN

DENNIS KNIEP - 27.05.2017

EGG HUNTERS  
Welcome



Holiday Inn



# AGENDA

- 0x01 Buffer Overflow
- 0x02 **Egghunting**
- 0x03 Linux (Egghunting)
- 0x04 Fuzzing
- 0x05 Windows (Fuzzing & Egghunting)

# BUFFER OVERFLOW



# **BASICS**

## **MEAT BEFORE MILK**

# PROGRAM EXECUTION

- CPU
- Memory

# CPU

CPU loads instruction (code) from memory and  
executes it

# CPU INSTRUCTIONS

- load information from memory into registers (mov)
- process data from memory or registers (comp, add, sub)
- change next executed instruction (jmp, call, int)
- etc.



# CPU REGISTERS

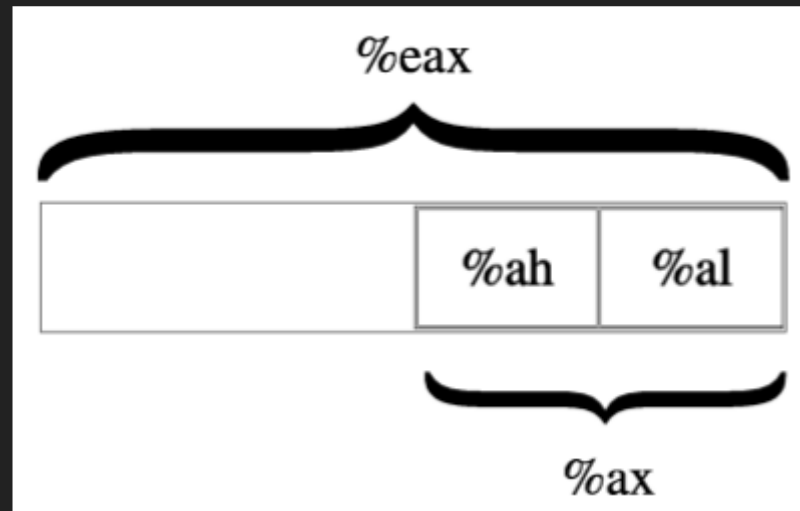
Most information is stored in memory, brought in to the registers for processing, and then put back into memory when the processing is completed.

Registers are what the computer uses for computation.

# GENERAL-PURPOSE REGISTERS

- eax
- ebx
- ecx
- edx
- edi
- esi

# REGISTERS SIZE



4 bytes / 32 bit

0xFFFFFFFF

# SPECIAL-PURPOSE REGISTERS

- ebp - base pointer register
- esp - address of the value on the top of the stack
- eip - instruction pointer
- eflags

# CPU CYCLE WITH EIP

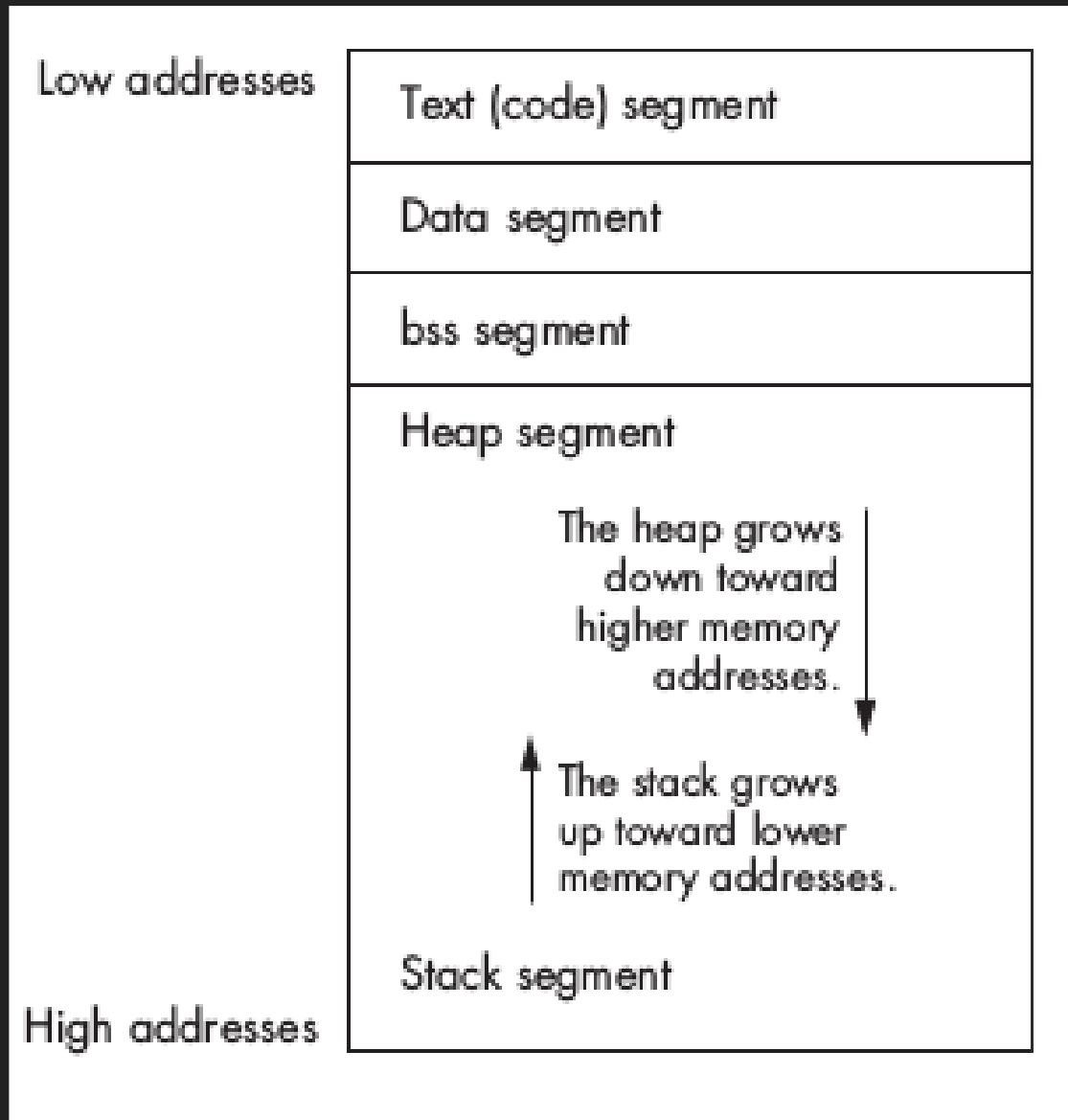
- EIP points to next instruction in memory
- CPU loads OpCode from memory
- Increments EIP
- Executes OpCode

# MEMORY

**0X0 - 0xFFFFFFFF**

Jeder Prozess arbeitet virtuell mit dem gesamten  
Memory

# A COMPILED PROGRAM'S MEMORY





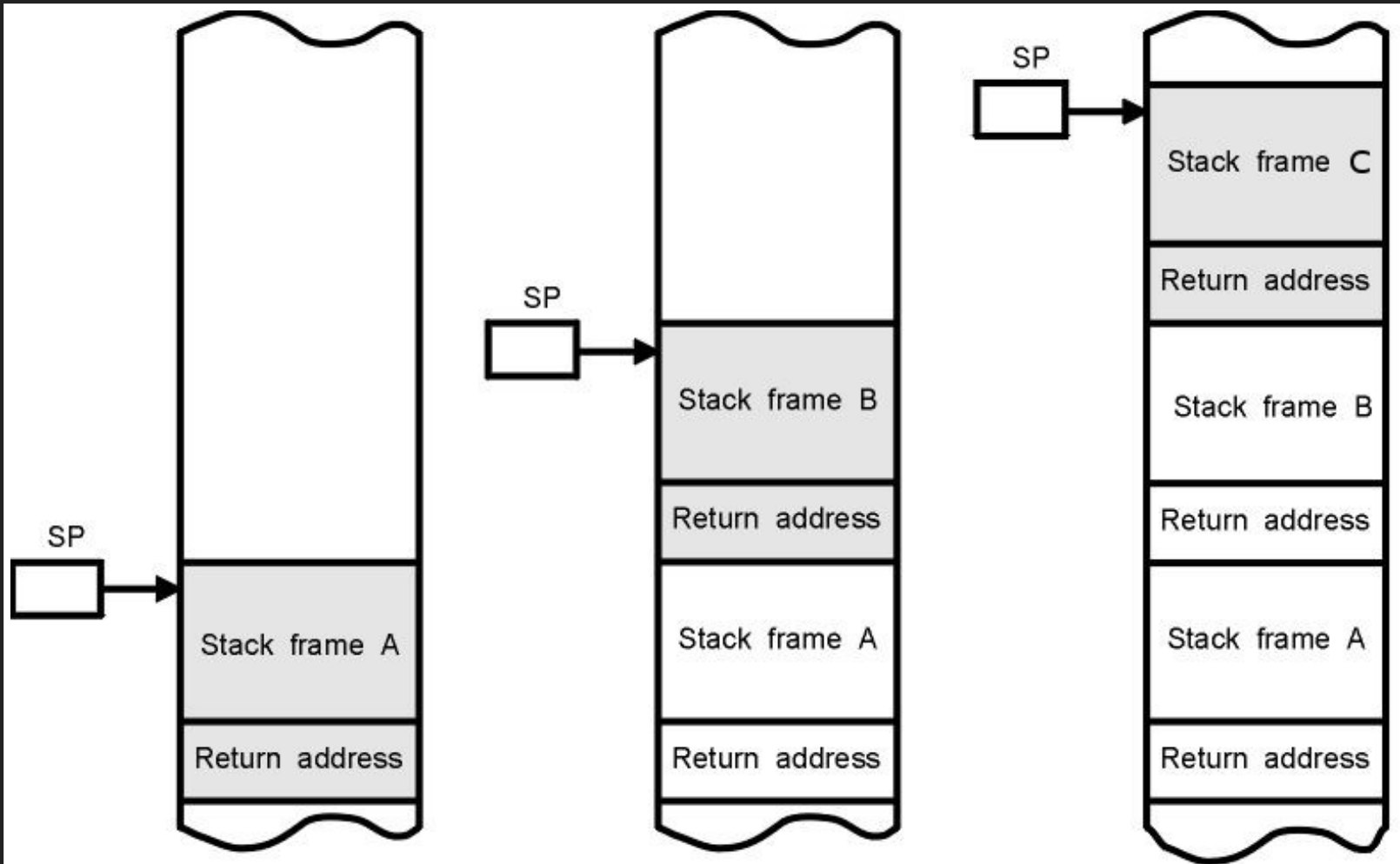
# STACK

Used for storing the data for the current executed function chain.

# STACKFRAMES

```
void C(){  
  
}  
  
int B(){  
    C();  
    return 2;  
}  
  
int A(int x){  
    int b = B();  
    return x + b;  
}  
  
int main(int argCount, char *args[]){  
    return A(15);  
}
```

# STACKFRAMES

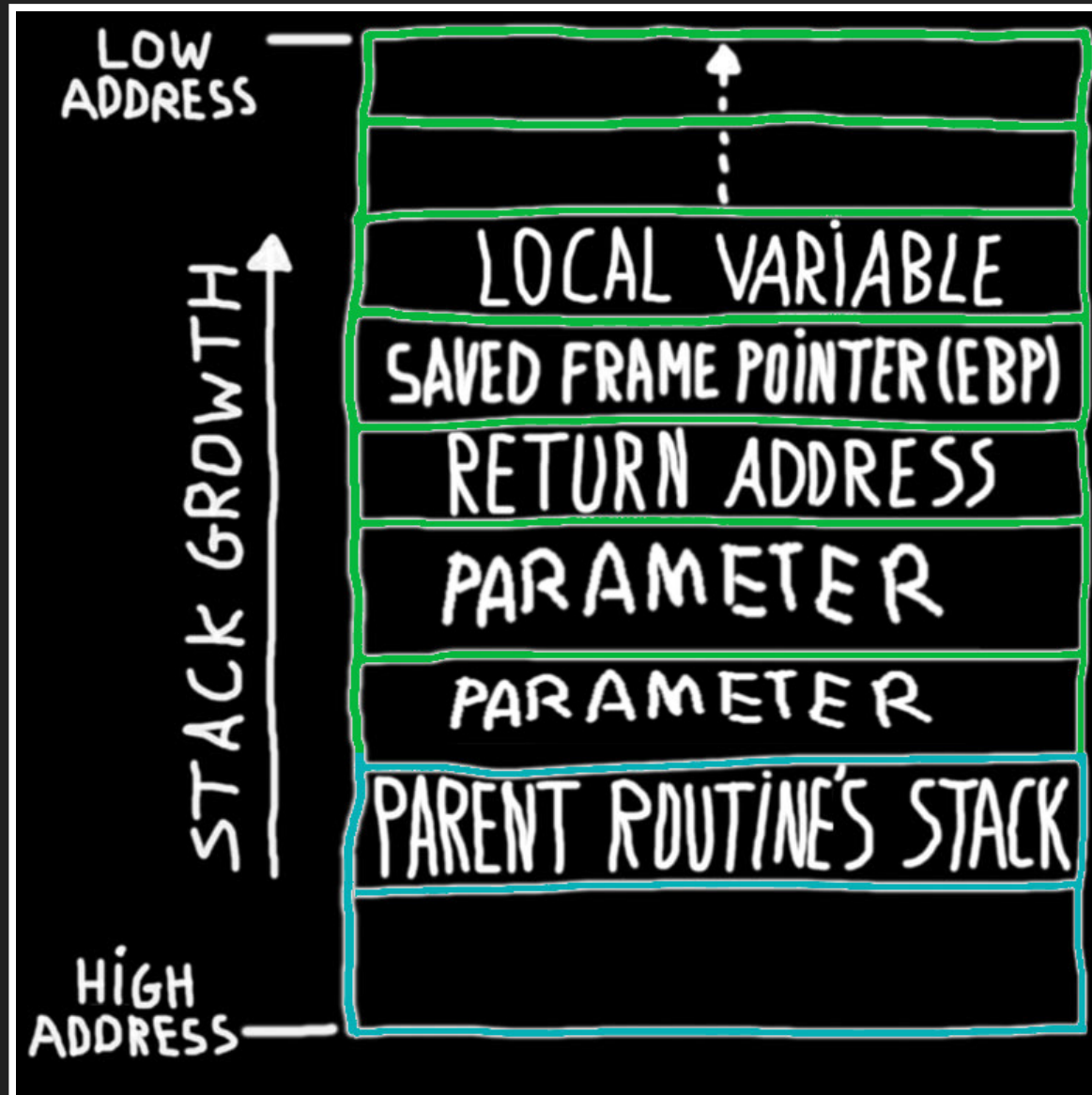


# FUNKTIONSAUFRUF SCHEMA

```
A(15);
```

- (Save registers to stack)
- Add Functionparameter to Stack
- Call the Function  
(Stores Return EIP, Set EIP to Function)
- Save calling functions ebp
- Set current functions ebp

# STACKFRAME



# BASEPOINTER

constant reference to the stack frame

Access all of the data in the function by using base pointer addressing using different offsets from ebp!

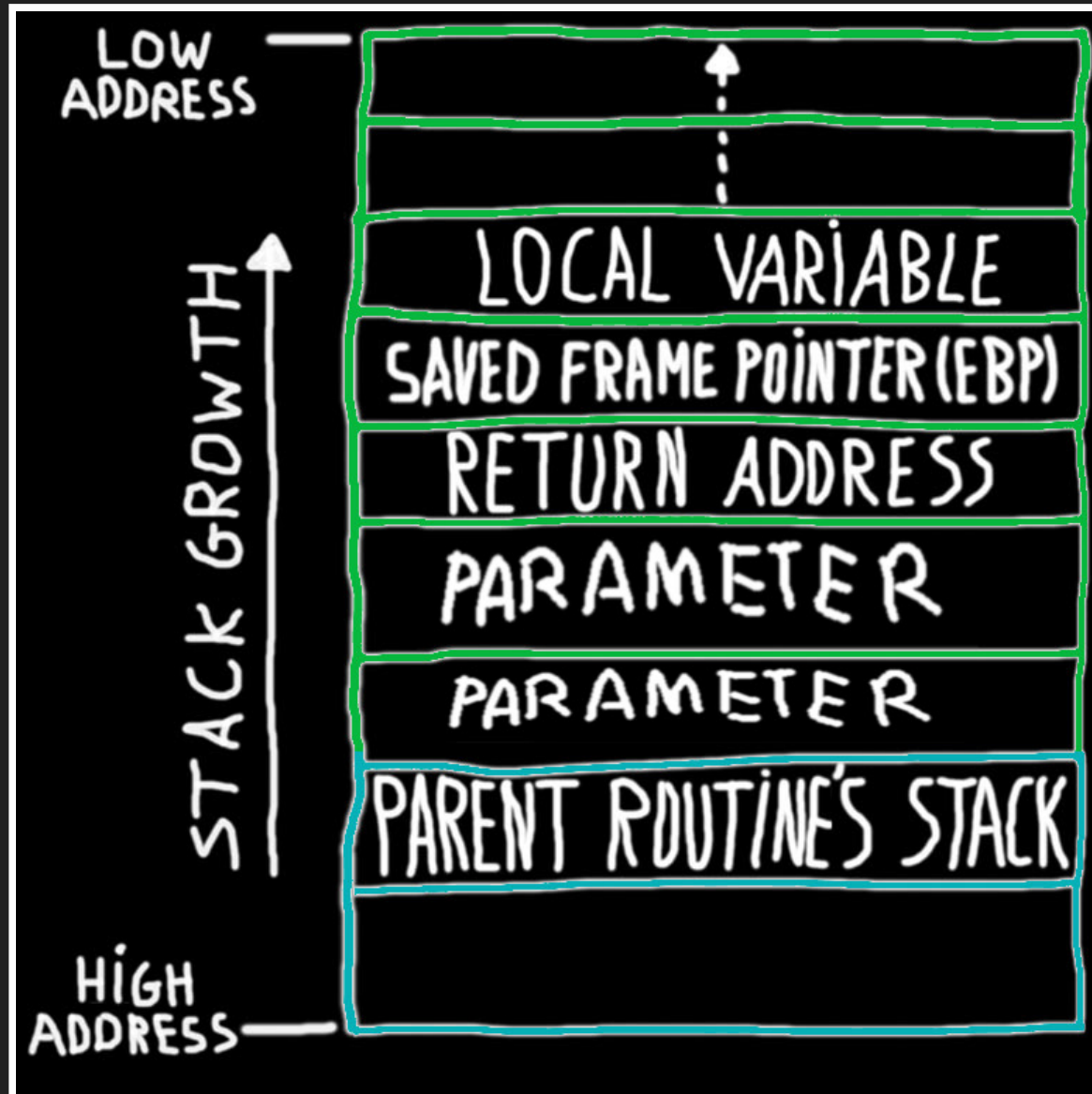
- Local Var 2: [ebp-8]
- Local Var 1: [ebp-4]
- Parameter 1: [ebp+8]
- Parameter 2: [ebp+12]

# FUNKTIONSABBAU SCHEMA

```
return;
```

- Store return value (in eax)
- Reset Stack to current ebp
- Restore calling functions ebp
- Leave Function (Set EIP to Return EIP)
- Remove the Parameters from Stack
- (Pop the saved registers from the Stack)
- Process return value

# STACKFRAME





# STACKFRAMES

## KONVENTION, KEINE PFLICHT FÜR CPU

- Compiler übersetzt den C Code so, dass er Stackframes benutzt
- Stack könnte aber aus Sicht der CPU völlig anders genutzt werden
- Es gibt keinen Supervisor

Wiki Calling Convention

# LOKALE VARIABLEN

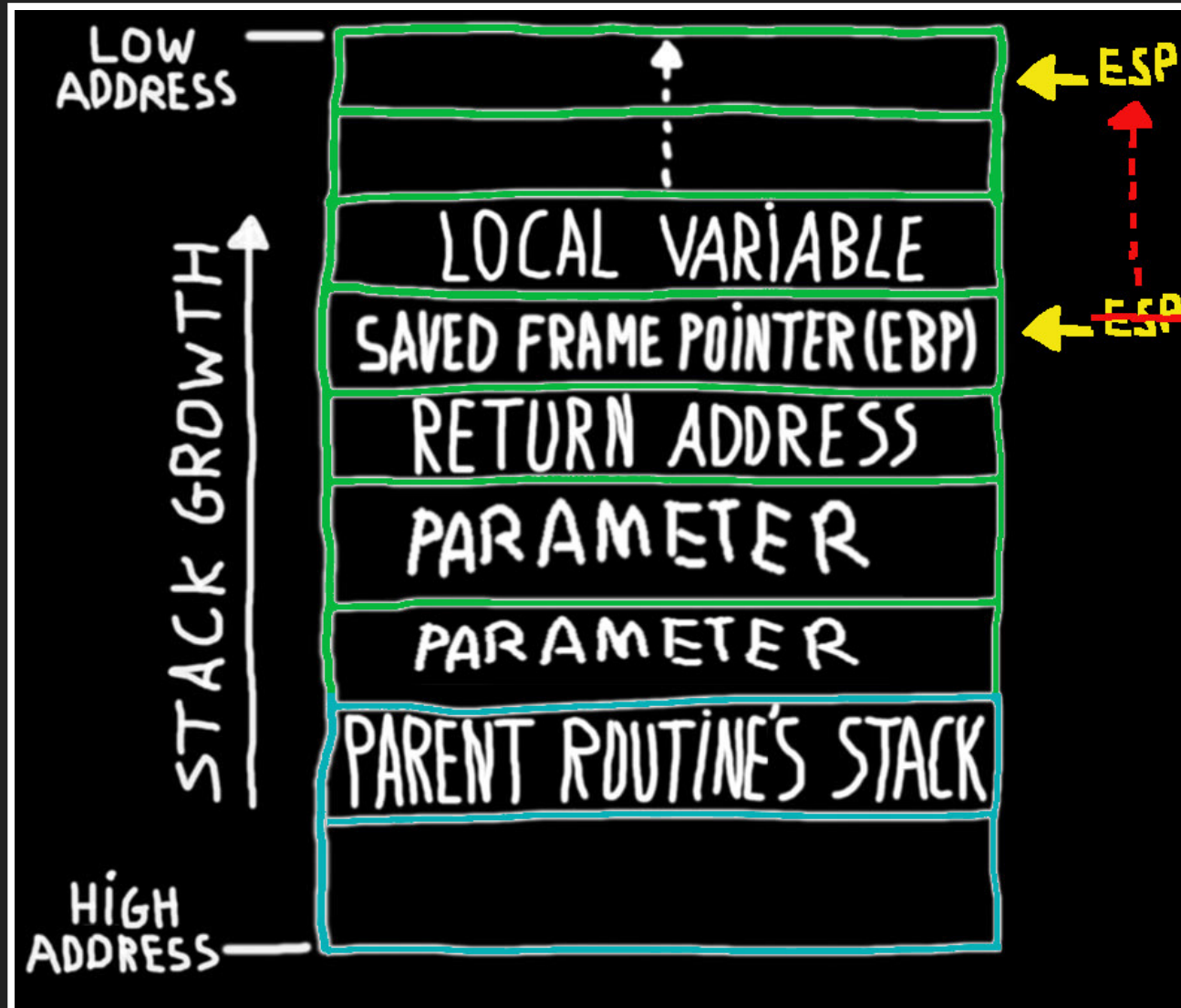
Lokale Variablen werden im Stackframe gespeichert.

```
int main(int argCount, char *args[]){  
    char input[100];  
    return 0;  
}
```

```
sub    esp, 0x70
```

# LOKALE VARIABLEN

Platz für lokale Variable geschaffen.



# LOKALE VARIABLEN

Es wurde Platz für 100 byte geschaffen.

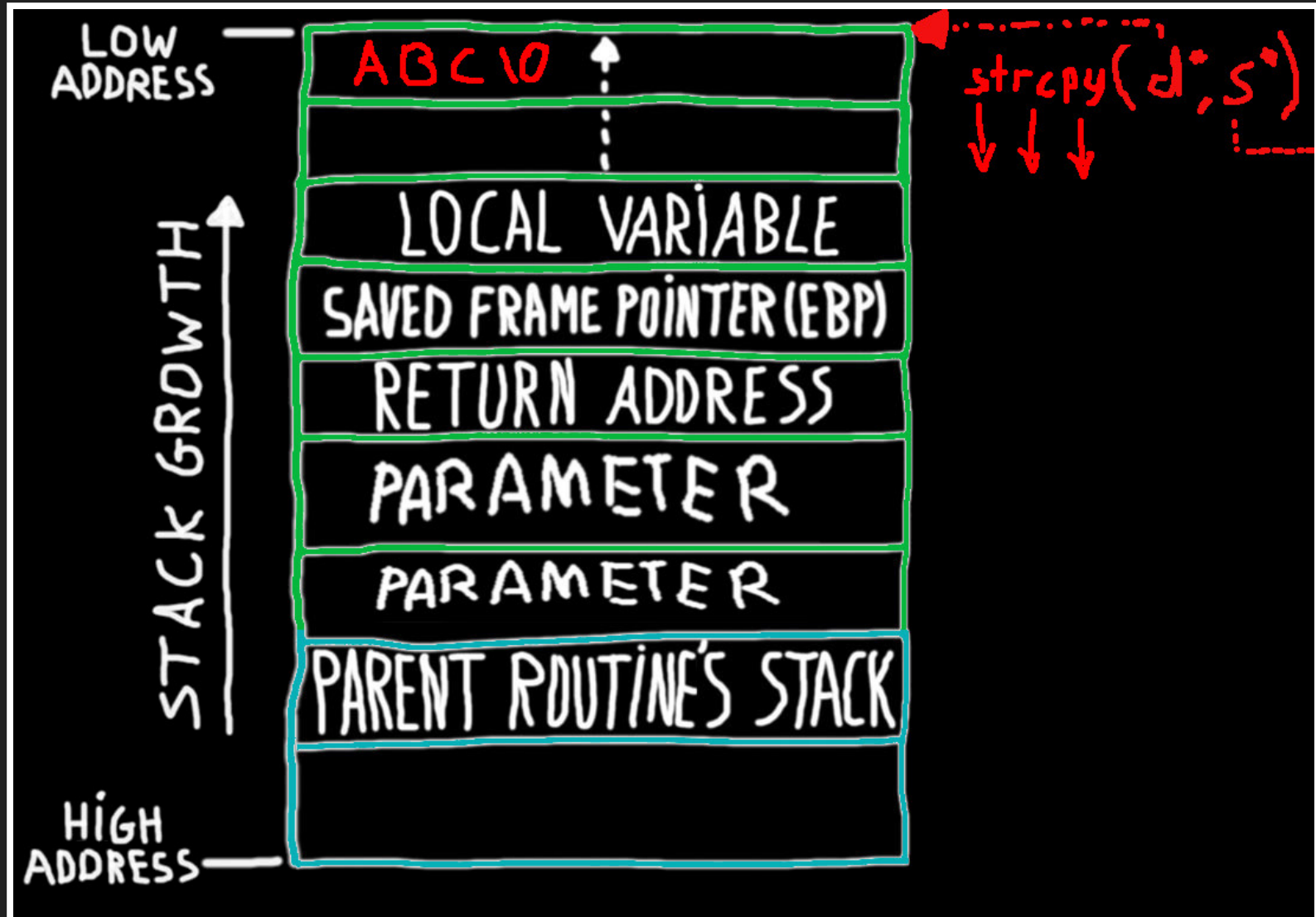
Nun soll mit `strcpy` das erste Programm Argument in den Buffer geschrieben werden.

```
char* strcpy(char *dest, char *src);
```

Kopiert solange bis ein null-byte kopiert wurde (String-Terminator)

# LOKALE VARIABLEN

Platz für lokale Variable geschaffen.



Strcpy schreibt entgegen der Stack growth...  
...solange bis ein NullTerminator kopiert wurde.  
Enthält \*src mehr bytes als \*dest reserviert hat  
werden Bereiche überschrieben  
Es gibt keine built-in safe-guards die sicherstellen,  
dass der reservierte Platz ausreicht

# DEMO



Beliebiges Ausführen von Memory (Kontrolle EIP)  
Beliebiges Schreiben und Lesen im Memory (StrCpy)



# POI - EIP

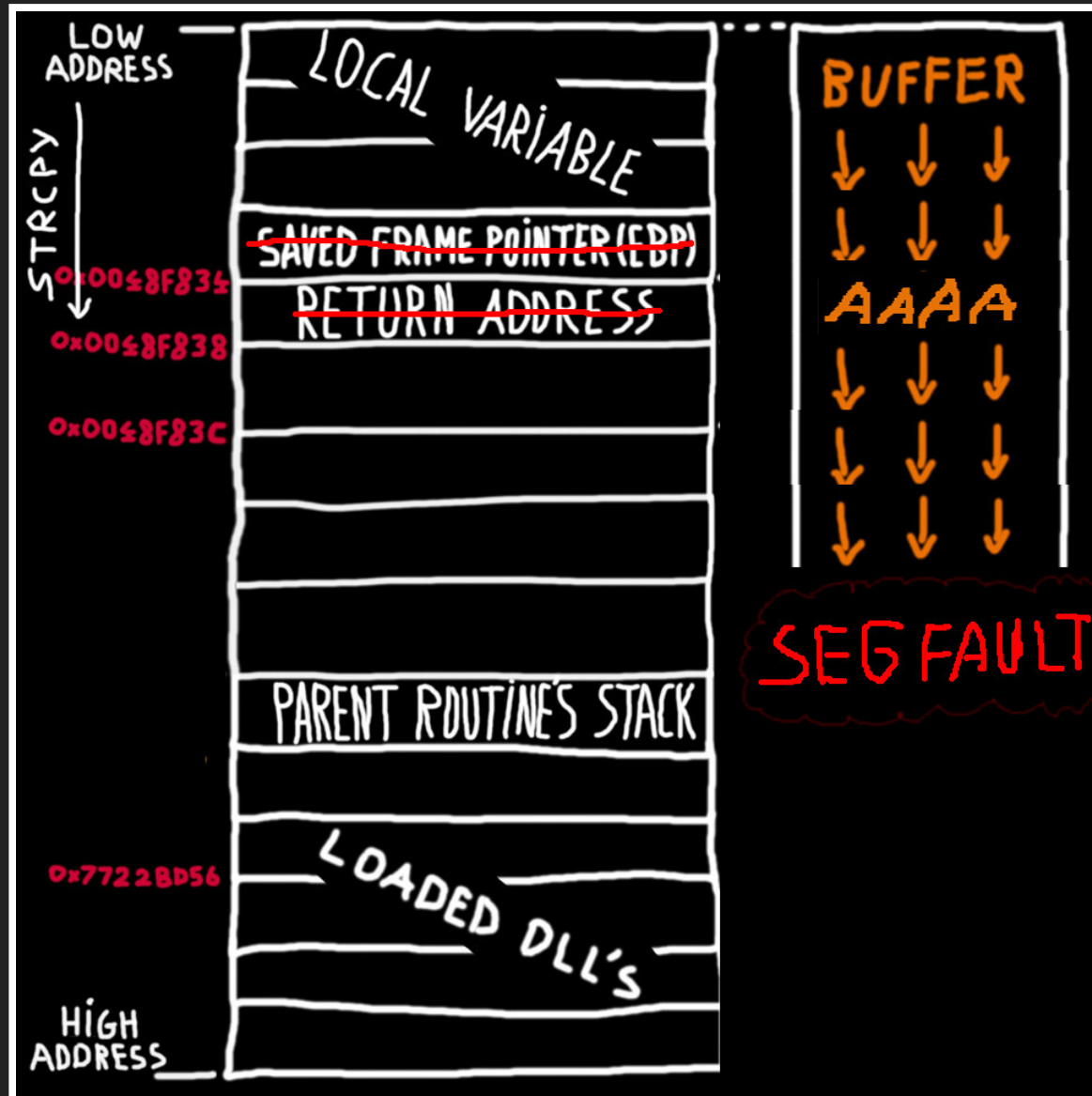
- EIP ist der Zeiger auf den auszuführenden Code
- Nachdem eine Funktion verlassen wird, wird der Return EIP von dem Stack gelesen...
- ...und in das EIP Register geschrieben.
- CPU Cycle führt den OpCode der sich an der Memory Adresse, worauf der EIP zeigt, aus
- Der Return EIP steht unterhalb der lokalen Variablen im Stackframe

# POI - STRCPY

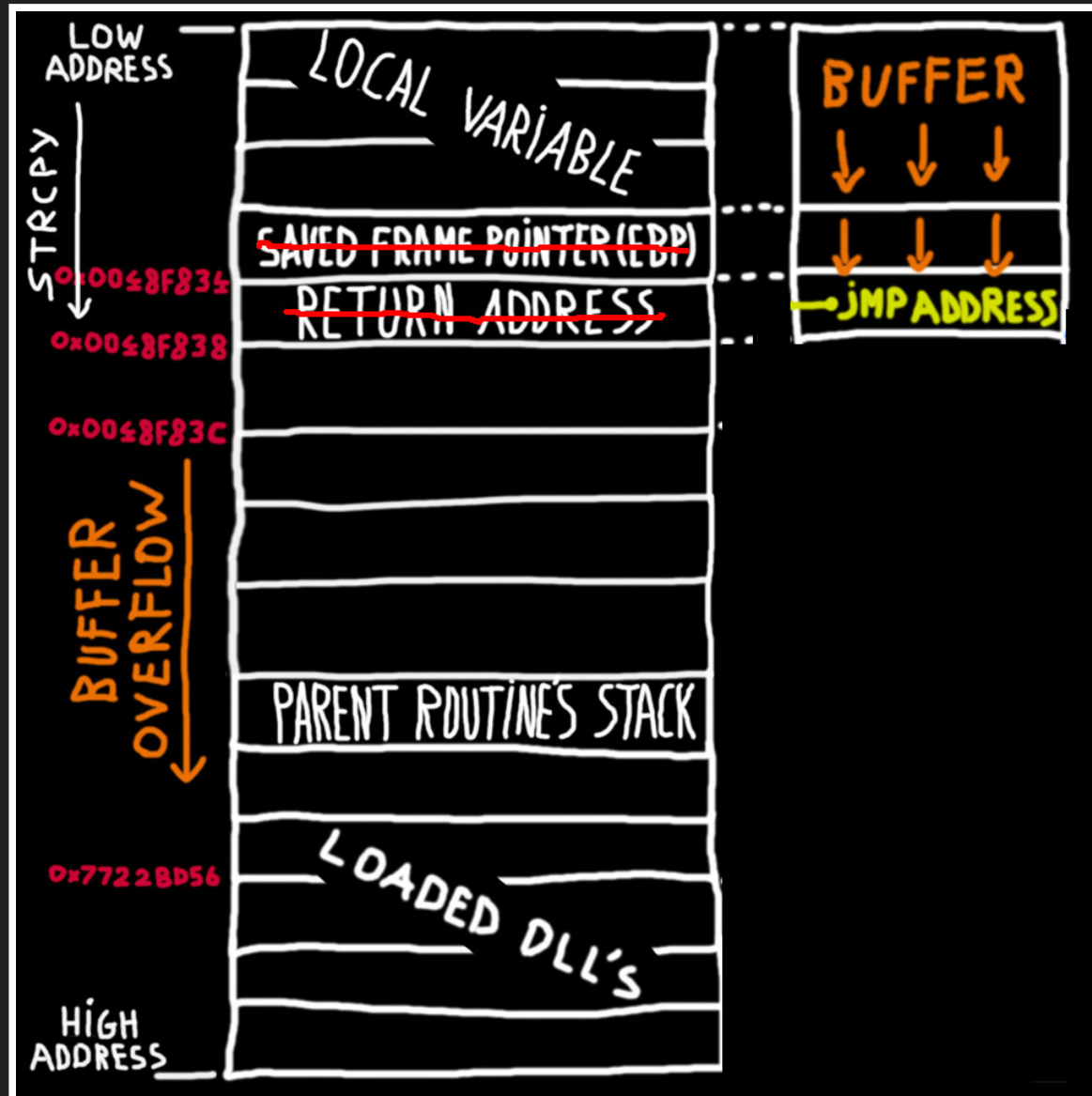
- `strcpy` schreibt entgegen der Stack growth
- und unabhängig von der Größe des Zielbuffers bis ein Null-byte kopiert wurde

# BUFFEROVERFLOW

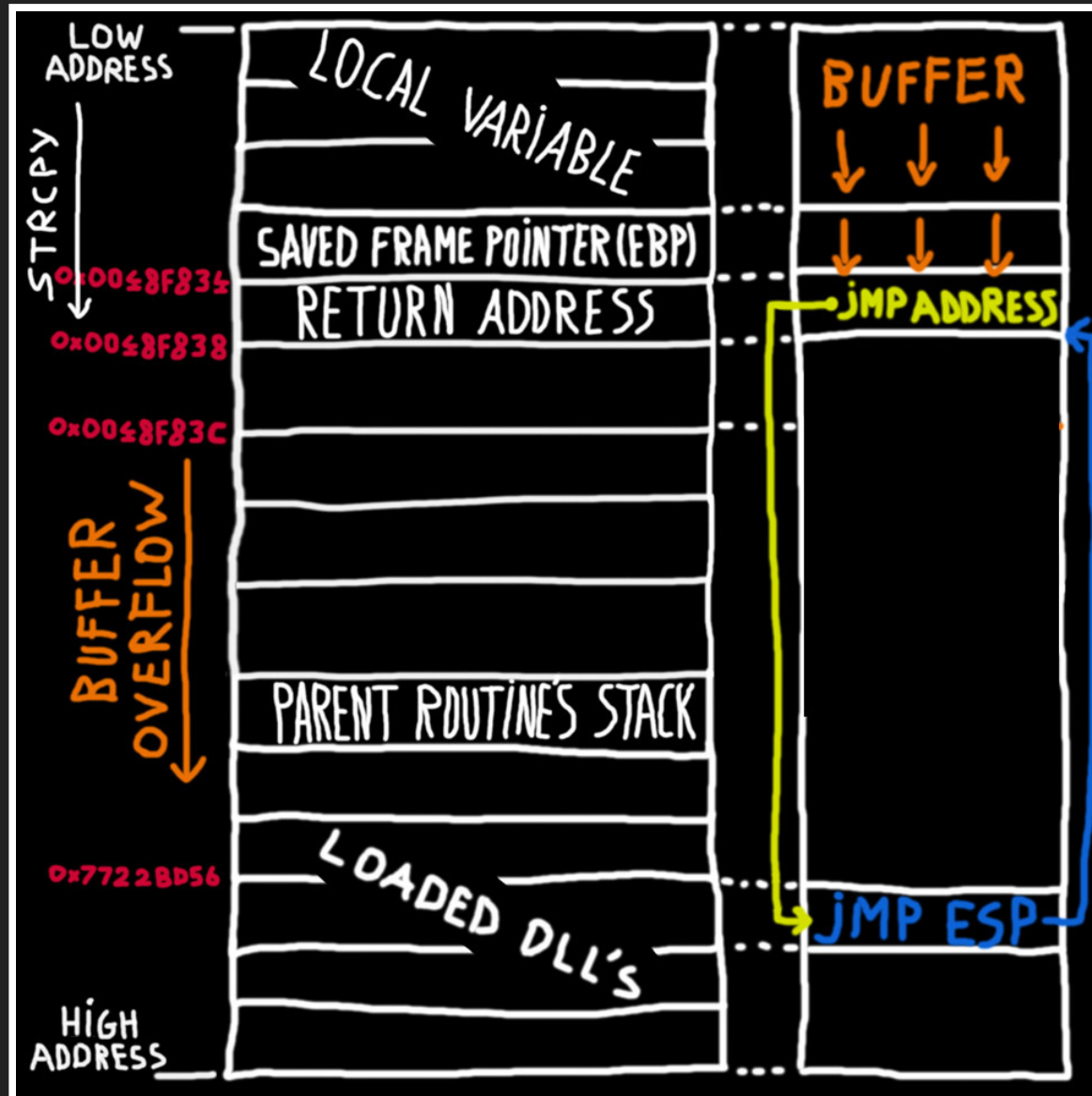
# BUFFEROVERFLOW



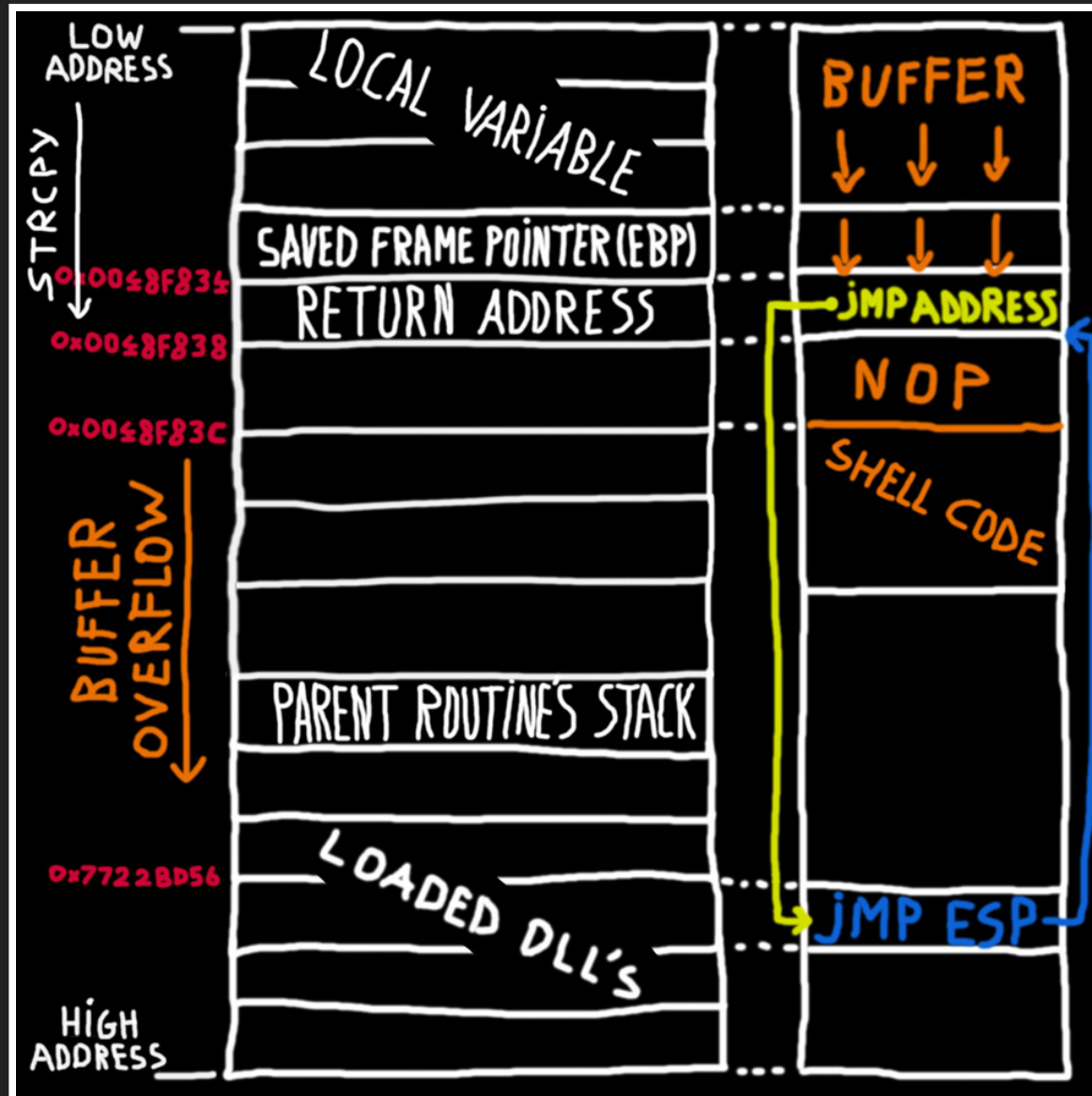
# BUFFEROVERFLOW



# BUFFEROVERFLOW



# BUFFEROVERFLOW



# TODO 4 EXPLOIT

- Pattern erstellen und EIP Offset finden
- BadChars identifizieren
- Addr für EIP finden (Shellcode ausführen z.B. `Jump ESP`)
- Shellcode generieren (ggf. mit Encoder wg. BadChars)



# PREREQUISITES LINUX

...for 32bit on 64bit

```
apt-get install g++-multilib libc6-  
dev-i386
```

...GDB Extension

Install PEDA

```
#include <stdio.h>
#include <string.h>

int countInputLength(char *input){
    char buffer[50];
    strcpy(buffer, input);
    int len = strlen(buffer);
    return len;
}

int main(int argCount, char *args[]){
    if(argCount < 2){
        return 0;
    }

    char input[200];
    strncpy(input, args[1], 200);

    int len = countInputLength(input);
    printf("Your input is %d characters long!\n", len);
    return 0;
}
```



# COMPILE

```
gcc vulnByDesign.c -m32 -g -fno-builtin -fno-stack-protector  
-z execstack -no-pie -o vulnByDesign
```

-m32	32 Bit
-g	debug information
-fno-builtin	Do not optimize library Functions never copied to the code
-fno-stack-protector	Disable CANARY
-z execstack	Disable DEP (data execution prevention)
-no-pie	Disable PIE (Position Independent Executable)

# DISABLE ASLR

```
echo 0 > /proc/sys/kernel/randomize_va_space
```

oder

```
gdb-peda$ aslr [on|off]
```

# START GDB

```
gdb --args ./vulnByDesign abcdef
```

```
gdb --args ./vulnByDesign `for i in {1..200}; do echo -n  
a;done`
```

```
gdb --args ./vulnByDesign `python -c 'print 200 * 'A''`
```

A man with dark hair, wearing a black t-shirt, is performing a yoga pose on a blue patterned mat. He is in a low lunge position with his right leg bent and his left leg extended back. His hands are raised, with his fingers spread, and he is looking directly at the camera. The background shows a room with bookshelves and a desk.

# HANDS ON

## WARMUP

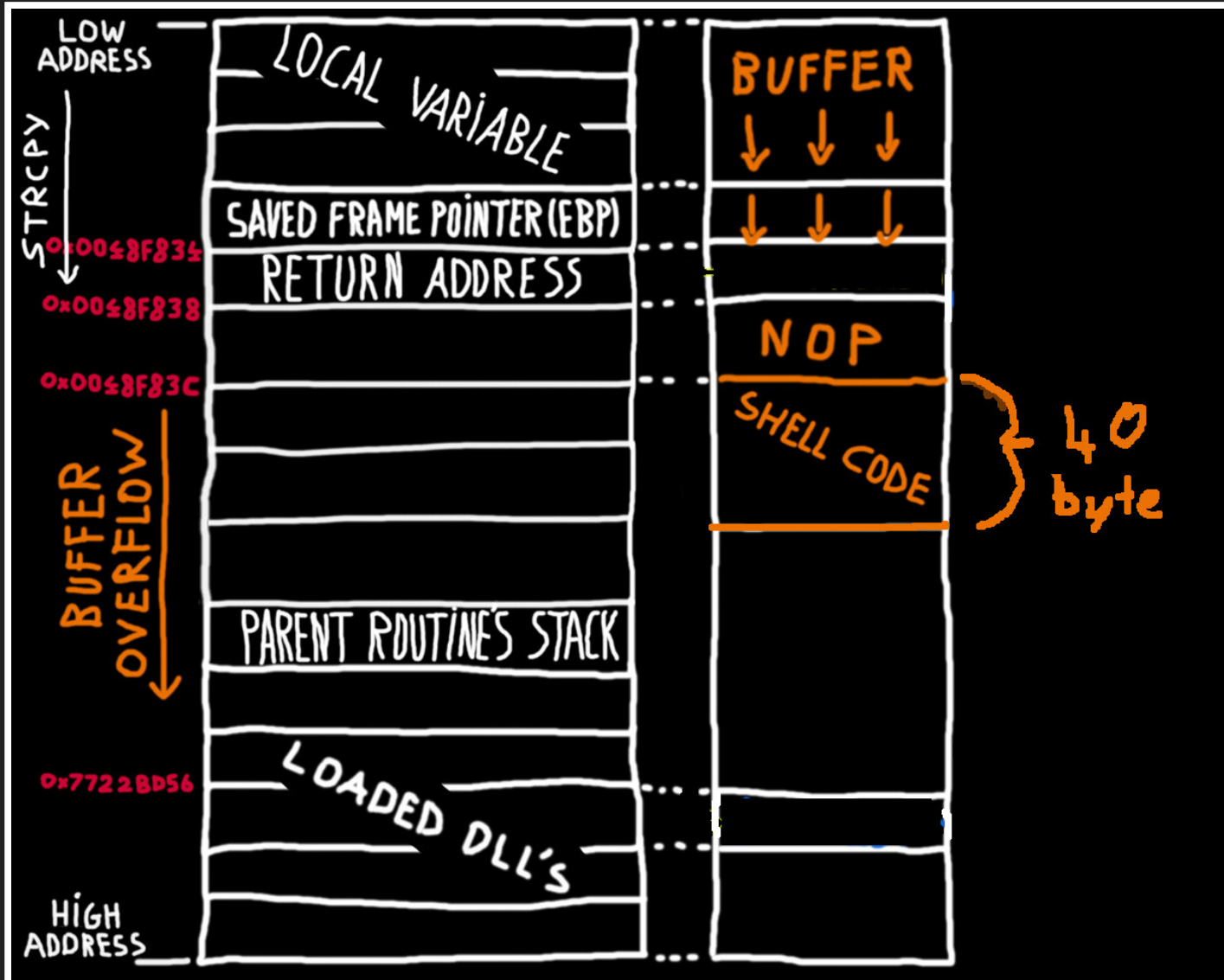
# EGGHUNTING





# WHY EGGHUNTING?

## Zu wenig Platz auf dem Stack für den Shellcode - Use staged Shellcode



# FUNKTIONSWEISE

- Shellcode auf den Heap/Stack
- Shellcode ist mit einem Egg markiert
- Egghunter über BufferOverflow auf den Stack
- Egghunter durchsucht den RAM nach dem Egg
- Mit Egg markierten Shellcode ausführen

# FUNKTIONSWEISE



# INPUTVECTORS

- Zusätzliche Eingabemöglichkeiten finden (Manuell bzw. Fuzzing)
- Nicht alle Eingaben sind dauerhaft im RAM zu finden (free()/methode wird verlassen/anderer thread/etc.)
- Eingaben müssen zum Zeitpunkt des BufferOverflows zur Verfügung stehen

# PSEUDOEXPLOIT

```
egg = "{unique 4byte hex}";
```

```
shellcode = "{opens a reverse shell}";
```

```
buffer4Heap = egg + egg + shellcode;
```

```
egghunter = "{searches for eggs in RAM and jmp to it}";
```

```
buffer4Stack = offset + eip + nopsled + egghunter;
```

# EGG

- unique
- eight byte egg when doing the searching  
stored once in the searching code and accidentally run into the egg hunter itself vice running into the expected buffer
- 0x50905090 repeated twice  
(opcode for nop; push eax)
- some egg hunter implementations require that the egg itself be executable assembly

# EGGHUNTER - ANFORDERUNGEN

- Robust (inaccessible mem)
- Small
- Fast

# EGGHUNTER - IMPLEMENTIERUNG

- most obvious approach would be to register a **SIGSEGV handler** to catch invalid memory address  
(but egghunter size is too large)
- system call for checking if memory address is valid

Safely Searching Process Virtual Address Space





# HANDS ON

## EGGHUNTING - LINUX

# EGGHUNTER BUGFIX

Der 35byte große Egghunter terminiert mit einem SIGSEGV.

# SMALL EGGHUNTER

- Mit 18 byte sehr klein. Aber nicht robust!
- Wurde "optimiert" indem der AccessCheck entfernt wurde
- Außerdem können 4byte eggs im shellcode verwendet werden

Um nicht versehentlich im egghunter zu landen wird das (egg -1) als Hex gespeichert und dann im egghunter vor der Suche hochgezählt

A photograph of four brown eggs standing in a row behind a plate of scrambled eggs. Each egg has a sad face drawn on it with black marker. The egg on the far left has a tear running down its cheek. The egg in the second position from the left has a wide-open mouth as if crying. The egg in the third position has a frown and heavy-lidded eyes. The egg on the far right has a surprised or shocked expression with wide eyes and an open mouth. The plate of scrambled eggs is in the foreground, and the background shows a kitchen setting with a stove and a washing machine.

# SCRAMBLED EGG(HUNTER)

# HOW TO IDENTIFY VULNERABILITIES AND INPUTVECTORS

- Manual testing
- Code Review
- Static Code Analysis
- Fuzzing
- ...

# FUZZING

# FUZZING

Fuzzing is a process of  
**sending arbitrary/malicious/random data**  
to a program in order to  
**generate failures, or errors in the application**

# FUZZING

typically **local** analysis technique and not something we would typically use in a remote penetration test.



# FUZZING

- Identify Inputs
- Generate Fuzzed Data (Define Data Schema)
- Send Fuzzed Data
- Monitor for Exception, Log and Restart
- Determine Exploitability

# IDENTIFY INPUTS

Prior to engaging in a fuzzing test

- It is important to understand the protocol you will be evaluating
- Identify all input that the application accepts and processes

# MONITOR FOR EXCEPTION, LOG AND RESTART

we typically need to attach a debugger to the target process to capture stack and register dump information.

# CHOOSE TOOL

- Manually
- Code yourself (Python,...)
- Fuzzer-Tool with config (sFuzz, Peach, etc.)
- Frameworks (sully, kitty, SPIKE, etc.)

Config vs. Coding

[fuzzing resources](#)

# SIMPLE FUZZER

Fuzzing mit sfuzz

```
sfuzz -S 127.0.0.1 -p 9999 -f fuzz.vulnserver.config -T0ev
```

# SIMPLE FUZZER CONFIG

```
#generates sequences incremented by and max
    seqstep=1000
    maxseqlen=10000
    literal=abcdefg
    sequence=A
    sequence=1
#Overwrite Line Terminator
    !CRLF=0d 0a
    lineterm=CRLF
# wait between requests (default 100ms)
    reqwait=200
    endcfg
    STATS FUZZ

    - -
    GMON FUZZ

    - -
```

# BOOFUZZ

Fork from Sully

Install boofuzz

# FUZZING FOR BUFFEROVERFLOW

- Bufferlänge verändern
- Bedingungen (IF) im Code berücksichtigen  
(Zeichen, Länge, Reihenfolge der Aufrufe etc.)
- Ist die Länge ein Parameter?



# FUZZING FOR EGGSHELLCODE SPACE

- Bufferlänge verändern
- Bedingungen (IF) im Code berücksichtigen  
(Zeichen, Länge, Reihenfolge der Aufrufe etc.)
- Nicht alle Eingaben sind dauerhaft im RAM zu finden  
(free())/methode wird verlassen/anderer thread/etc.)

# FUZZING FOR EGGSHELLCODE SPACE

```
sendPatternData()
```

```
sendBufferOverflowExploit()
```

```
memDump = DumpMemory()
```

```
searchForPatternData(memDump)
```

# PREPARE 4 WINDOWS

- Lunch vulnserver.exe on Desktop
- Start Immunity Debugger with Mona
- File > Attach
- Or use debugVulnserver.bat

# PREPARE 4 WINDOWS

- !mona findmsp
- !mona jmp -r {register}
- !mona pattern\_create
- !mona pattern\_offset
- !mona asm -s {assembler instruction}

A man with dark hair, wearing a dark blue or black t-shirt, is in a room. He has both hands raised in the air, palms facing forward, in a gesture of surprise or surrender. He is looking directly at the camera with a neutral expression. The room has warm, yellowish lighting. In the background, there are bookshelves filled with books. The floor appears to be covered with a patterned rug or blanket. The overall mood is casual and slightly humorous.

# HANDS ON

## FUZZING & EGGHUNTING - WINDOWS