前有schedutil governor来调节cpu频率,现有schedfreq governor来调节频率,现在来讲解它的来龙去脉。

源代码参考AOSP kernel的这个分支:remotes/origin/android-msm-wahoo-4.4-pie, kernel version:4.4.116。源代码下载地址:<u>https://aosp.tuna.tsinghua.edu.cn/kernel/msm.git</u> 使用我的分析的kernel code方式如下:

- git clone https://aosp.tuna.tsinghua.edu.cn/kernel/msm.git
- git branch -a 查看有什么分支

- git checkout -b remotes/origin/android-msm-wahoo-4.4-pie
- 正确切换到我所看的kernel 分支, Android 9.0

闲话少扯。。。。

按照老样子,先看governor相关的结构体成员变量:

```
    /*降频和升频的最小间隔,可以修改*/
    #define THROTTLE_DOWN_NSEC 50000000 /* 50ms default */
    #define THROTTLE_UP_NSEC 500000 /* 500us default */
    static DEFINE_PER_CPU(unsigned long, enabled);
    DEFINE_PER_CPU(struct sched_capacity_reqs, cpu_sched_capacity_reqs);
```

```
/*tunable的参数,对于升频和降频的时间限制,attr set设置sys接口,可供userspace调节*/
struct gov tunables {
    struct gov_attr_set attr_set;
    unsigned int up_throttle_nsec;
    unsigned int down throttle nsec;
};
 * gov data - per-policy data internal to the governor
 * @up throttle: next throttling period expiry if increasing OPP
 * @down throttle: next throttling period expiry if decreasing OPP
* @up throttle nsec: throttle period length in nanoseconds if increasing
* @down throttle nsec: throttle period length in nanoseconds if decreasing
OPP
* @task: worker thread for dvfs transition that may block/sleep
 * @irq work: callback used to wake up worker thread
 * @requested_freq: last frequency requested by the sched governor
 * struct gov data is the per-policy cpufreq sched-specific data structure.
* per-policy instance of it is created when the cpufreq sched governor
* the CPUFREQ GOV START condition and a pointer to it exists in the
gov data
 * member of struct cpufreq policy.
 * Readers of this data must call down read(policy->rwsem). Writers must
 * call down_write(policy->rwsem).
 */
struct gov data {
    ktime t up throttle; /*升频的时间节点*/
    /*降频的时间节点,是当前时间+门限数值,即为下次降频的最近时间节点*/
    ktime t down throttle;
    /*tunable参数*/
    struct gov_tunables *tunables;
    struct list head tunables hook;
    /*频率修改的进程*/
    struct task_struct *task;
    /*slow adjust freq的worker*/
    struct irq work irq work;
    /*更加capacity数值设定的请求频率,在通过频率table挑选idx并获取对于
    freq table的freq value,最后通过dvfs进行频率的调整*/
    unsigned int requested freq;
};
```

接着看cpufreq governor结构体的填充。

static int cpufreq sched setup(struct cpufreq policy *policy,

```
unsigned int event)

{
    switch (event) {
    case CPUFREQ_GOV_POLICY_INIT:
        return cpufreq_sched_policy_init(policy);
    case CPUFREQ_GOV_POLICY_EXIT:
    return cpufreq_sched_policy_exit(policy);
```

```
case CPUFREQ GOV START:
         return cpufreq sched start(policy);
     case CPUFREQ_GOV_STOP:
        return cpufreq_sched_stop(policy);
     case CPUFREQ GOV LIMITS:
         cpufreq sched limits(policy);
         break;
    return 0;
#ifndef CONFIG CPU FREQ DEFAULT GOV SCHED
static
#endif
struct cpufreq_governor cpufreq_gov_sched = {
    .name = "sched",
.governor = cpufreq_sched_setup,
.owner = THIS_MODULE,
};
static int    init cpufreq sched init(void)
    int cpu;
    for each cpu(cpu, cpu possible mask)
        per cpu(enabled, cpu) = 0;
    return cpufreq_register_governor(&cpufreq_gov_sched);
}
/* Try to make this the default governor */
fs initcall(cpufreq sched init);
```

可以看到governor名字为"sched",顾明思议就是更加调度器的某些变量来调节cpu频率的。我们看起governor callback函数的init:

```
static int cpufreq_sched_policy_init(struct cpufreq_policy *policy)
   struct gov data *gd;
   int cpu;
   int rc;
   /*对每个cpu上的cpu_sched_capacity结构体进行初始化为0*/
   for each cpu(cpu, policy->cpus)
       memset(&per_cpu(cpu_sched_capacity_reqs, cpu), 0,
              sizeof(struct sched capacity reqs));
   /*为sched governor data分配空间*/
   gd = kzalloc(sizeof(*gd), GFP KERNEL);
   if (!gd)
       return -ENOMEM;
   /*将sched governor data挂载到cpu policy governor data上,即关联上*/
   policy->governor data = gd;
   if (!global tunables) {
       /*对tunable结构体变量分配空间*/
       gd->tunables = kzalloc(sizeof(*gd->tunables), GFP KERNEL);
       if (!gd->tunables)
```

```
goto free qd;
    /*设置频率升高的时间限制,也就是升频率间隔不能小于这个间隔*/
   gd->tunables->up_throttle_nsec =
       policy->cpuinfo.transition_latency ?
       policy->cpuinfo.transition latency :
       THROTTLE UP NSEC;
   /*设置频率降低的时间限制*/
    gd->tunables->down throttle nsec =
       THROTTLE DOWN NSEC;
   /*初始化tunable结构体成员变量的kobject, 并产生sys fs*/
   rc = kobject init and add(&gd->tunables->attr set.kobj,
                 &tunables ktype,
                 get governor parent kobj (policy),
                 "%s", cpufreq_gov_sched.name);
   if (rc)
       goto free_tunables;
   /*属性设置*/
   gov_attr_set_init(&gd->tunables->attr_set,
             &gd->tunables hook);
   pr debug("%s: throttle threshold = %u [ns]\n",
        __func__, gd->tunables->up_throttle_nsec);
   if (!have_governor_per_policy())
       global tunables = gd->tunables;
} else {
   gd->tunables = global tunables;
   gov_attr_set_get(&global_tunables->attr set,
            &gd->tunables hook);
/*再次update, 上面那个是否有点多余哈??*/
policy->governor data = gd;
if (cpufreq driver is slow()) {
   cpufreq driver slow = true;
    /*cpufreq driver slow这个参数有点意思。下面创建thread, wakeup函数为:
   cpufreq_sched_thread, 最后会创建kschedfreq:0和kschedfreq:4。
   对于两个cluster的cpu架构*/
   gd->task = kthread_create(cpufreq_sched_thread, policy,
                 "kschedfreq:%d",
                 cpumask first(policy->related cpus));
   if (IS ERR OR NULL(gd->task)) {
       pr err("%s: failed to create kschedfreq thread\n",
              __func__);
       goto free tunables;
   get task struct(gd->task);
    /*绑定相关联的cpu*/
   kthread bind mask(gd->task, policy->related cpus);
   wake_up_process(gd->task);
    /*初始化irq work*/
   init_irq_work(&gd->irq_work, cpufreq_sched_irq_work);
set sched freq();
```

```
return 0;

free_tunables:
    kfree(gd->tunables);

free_gd:
    policy->governor_data = NULL;
    kfree(gd);
    return -ENOMEM;
}
```

我们能够看到,上面最重要的信息如下:

- 频率升高的时间限制
- 频率下降的时间限制
- 创建的thread, cpufreq_sched_thread, 频率调节的进程
- 初始化一个irq_work, callback函数为cpufreq_sched_irq_work, 最后还是 wakeupgd->task, callback cpufreq_sched_thread这个函数。

接下来看一下, cpufreq_sched_thread这个函数的实现过程:

```
* we pass in struct cpufreq policy. This is safe because changing out the
* policy requires a call to cpufreq governor(policy, CPUFREQ GOV STOP),
* which tears down all of the data structures and
__cpufreq_governor(policy,
* CPUFREQ GOV START) will do a full rebuild, including this kthread with
* new policy pointer
static int cpufreq_sched_thread(void *data)
    struct sched_param param;
    struct cpufreq policy *policy;
    struct gov data *gd;
    unsigned int new request = 0;
    unsigned int last request = 0;
    int ret;
    /*获取当前cpufreq_policy*/
    policy = (struct cpufreq_policy *) data;
    /*获取sched governor data*/
    gd = policy->governor data;
    param.sched priority = 50;
    ret = sched setscheduler nocheck(gd->task, SCHED FIFO, ¶m);
    if (ret) {
        pr warn("%s: failed to set SCHED FIFO\n", func );
        do_exit(-EINVAL);
    } else {
        pr_debug("%s: kthread (%d) set to SCHED_FIFO\n",
                __func__, gd->task->pid);
    }
    do {
        /*governor请求的频率*/
        new request = gd->requested freq;
        /*如果频率一致,则进程进入TASK INTERRUPTIBLE状态并sleep*/
```

```
if (new request == last request) {
       set current state(TASK INTERRUPTIBLE);
       if (kthread_should_stop())
           break;
       schedule(); /*放弃cpu的运行, sleep*/
    } else {
        * if the frequency thread sleeps while waiting to be
        * unthrottled, start over to check for a newer request
       /*是否最后启动频率请求, */
       if (finish last request(gd, policy->cur))
           continue;
       last request = new request;
       /*update 升和降频率的时间=当前时间+升/降频率时间间隔,同时调用DVFS
       进行频率的调节*/
       cpufreq_sched_try_driver_target(policy, new_request);
} while (!kthread should stop());
return 0;
```

我们来看一下finish_last_request函数的实现:

```
static bool finish last request(struct gov data *gd, unsigned int cur freq)
   ktime t now = ktime get();
   ktime t throttle = gd->requested freq < cur freq ?</pre>
       gd->down throttle : gd->up throttle;
   /*当前时间与下个周期频率变化的时间,如果现在时间已经超过本应该频率调节的时间
   节点, 那么就必须进行频率调节了*/
   if (ktime after(now, throttle))
      return false;
   while (1) {
       /*由于throtle > now*/
       int usec left = ktime to ns(ktime sub(throttle, now));
        /*将差值转换为us*/
       usec left /= NSEC PER USEC;
       trace cpufreq sched throttled(usec left);
       /*休眠[usec left,usec left+100]这个时间间隔*/
       usleep range(usec left, usec left + 100);
       now = ktime get();
                         /*当前时间*/
       /*再次比较当前时间与throttle的差值,如果now + sleep 时间 > throttle,
        则在thread里面重新计算是否需要频率调整。即休眠[usec left,usec left+100]
        这个时间段之后,判断finish last update为false。。
       if (ktime after(now, throttle))
          return true;
```

如果没有外部触发这个thread,最后last_request = new_request一直会相等,导致频率不会update,那么肯定还有调度算法来trigger governor来调节频率:

```
void update_cpu_capacity_request(int cpu, bool request)
```

```
unsigned long new capacity;
struct sched_capacity_reqs *scr;
/* The rq lock serializes access to the CPU's sched capacity reqs. */
lockdep assert held(&cpu rq(cpu)->lock);
/*获取当前cpu的sche capacity reqs数据结果数据*/
scr = &per_cpu(cpu_sched_capacity_reqs, cpu);
/*capacity受cfs和rt task的影响*/
new_capacity = scr->cfs + scr->rt;
/*将capacity增加10%*/
new capacity = new capacity * capacity margin
   / SCHED CAPACITY SCALE;
new capacity += scr->dl;
/*如果当前cpu的需要调整的capacity与源total一致,没有必要频率调节了*/
if (new_capacity == scr->total)
   return;
trace cpufreq sched update capacity(cpu, request, scr, new capacity);
/*更新源total数据为最新数据*/
scr->total = new capacity;
/*如果需要频率调整,则调用下面函数进行调整*/
if (request)
   update_fdomain_capacity_request(cpu);
```

看看update_fdomain_capacity_request函数的实现原理(比较简单):

```
static void update fdomain capacity request(int cpu)
   unsigned int freq new, index new, cpu tmp;
   struct cpufreq_policy *policy;
   struct gov data *gd;
   unsigned long capacity = 0;
   /*
    * Avoid grabbing the policy if possible. A test is still
    * required after locking the CPU's policy to avoid racing
    * with the governor changing.
    *//*是否启用sched governor*/
   if (!per cpu(enabled, cpu))
       return;
   /*获取当前cpu的cpufreg policy*/
   policy = cpufreq cpu get(cpu);
   if (IS ERR OR NULL(policy))
       return;
   if (policy->governor != &cpufreq gov sched ||
       !policy->governor data)
       goto out;
    /*获取当前sched governor data结构体*/
   gd = policy->governor data;
    /*对同一个policy的cpu找出最大的capacity数值*/
   /* find max capacity requested by cpus in this policy */
   for each cpu(cpu tmp, policy->cpus) {
        struct sched capacity reqs *scr;
```

```
scr = &per_cpu(cpu_sched_capacity_reqs, cpu_tmp);
        capacity = max(capacity, scr->total);
    /*使用capacity容量归一化最高频率*/
    /* Convert the new maximum capacity request into a cpu frequency */
    freq new = capacity * policy->max >> SCHED CAPACITY SHIFT;
    /*找出符合freq new频率的index new索引, 即频率table索引号*/
    if (cpufreq frequency table target(policy, policy->freq table,
                      freq_new, CPUFREQ_RELATION_L,
                      &index new))
        goto out;
    /*索引号在table里面对应的频率*/
    freq_new = policy->freq_table[index_new].frequency;
    /*频率校正*/
    if (freq_new > policy->max)
        freq new = policy->max;
    if (freq new < policy->min)
        freq new = policy->min;
    trace_cpufreq_sched_request_opp(cpu, capacity, freq_new,
                   gd->requested freq);
    if (freq_new == gd->requested_freq)
        goto out;
    /*更新governor data结构体请求的频率参数*/
    gd->requested freq = freq new;
    * Throttling is not yet supported on platforms with fast cpufreq
    * drivers.
    * /
   if (cpufreq driver slow)
        /*触发频率调节*/
       irq_work_queue_on(&gd->irq_work, cpu);
    else
       cpufreq_sched_try_driver_target(policy, freq_new);
out:
   cpufreq cpu put (policy);
```

最后都是执行到如下函数中:

```
gd->tunables->up_throttle_nsec);
gd->down_throttle = ktime_add_ns(ktime_get(),
gd->tunables->down_throttle_nsec);
up_write(&policy->rwsem);
}
```

至此完毕,比较简单。关键还是这个函数update_cpu_capacity_request在哪里调用的。调用关系方框图如下:

