由于在task在创建或者task从idle会wakeup的流程中涉及到sched domain和sched group的知识,所以现在提前看下这部分.

根据实际的物理属性.CPU domain分成下面三种:

CPU 分类	Linux内核分类	描述
SMT(多线程)	CONFIG_SHCED_SMT	一个物理核心可以有两个以及之上的执行线程,一个物理核心里面的线程共享相同的CPU资源和L1 cache,task的迁移不会影响cache利用率
MC(多核)	CONFIG_SHCED_MC	每个物理核心有独立的L1 cache,多个物理核心可以组成一个cluster,cluster共享L2 cache
DIE(SoC)	简称为DIE	SoC级别

# 一 cpu topology的获取

arm64架构的cpu拓扑结构存储在cpu topology[NR CPUS]变量当中:

```
struct cpu_topology {
    int thread_id;
    int core_id;
    int cluster_id;
    cpumask_t thread_sibling;
    cpumask_t core_sibling;
};

extern struct cpu_topology cpu_topology[NR_CPUS];

#define topology_physical_package_id(cpu) (cpu_topology[cpu].cluster_id)

#define topology_core_id(cpu) (cpu_topology[cpu].core_id)

#define topology_core_cpumask(cpu) (&cpu_topology[cpu].core_sibling)

#define topology_sibling_cpumask(cpu) (&cpu_topology[cpu].thread_sibling)
```

#### cpu topology[]数组获取的路径如下(都是从dts里面获取的):

```
kernel_init() -> kernel_init_freeable() -> smp_prepare_cpus() ->
init_cpu_topology() -> parse_dt_topology()
--->
static int __init parse_dt_topology(void)

{
    struct device_node *cn, *map;
    int ret = 0;
    int cpu;
    /*Mdts里面获取cpu node的跟节点*/
    cn = of_find_node_by_path("/cpus");
    if (!cn) {
        pr_err("No CPU information found in DT\n");
        return 0;
    }

    /*
    * When topology is provided cpu-map is essentially a root
    * cluster with restricted subnodes.
```

```
map = of_get_child_by_name(cn, "cpu-map");
       if (!map)
           goto out;
       ret = parse_cluster(map, 0);
       if (ret != 0)
           goto out_map;
        * Check that all cores are in the topology; the SMP code will
        * only mark cores described in the DT as possible.
        * /
       for_each_possible_cpu(cpu)
           if (cpu topology[cpu].cluster id == -1)
               ret = -EINVAL;
   out_map:
      of node put(map);
   out:
       of node_put(cn);
      return ret;
   /*看下解析cluster怎么做的*/
   static int __init parse_cluster(struct device_node *cluster, int depth)
       char name[10];
      bool leaf = true;
      bool has cores = false;
       struct device_node *c;
       static int cluster_id __initdata;
       int core id = 0;
•
       int i, ret;
        * First check for child clusters; we currently ignore any
        * information about the nesting of clusters and present the
•
        * scheduler with a flat list of them.
       i = 0;
•
       do { /*多个cluster, 迭代遍历*/
•
           snprintf(name, sizeof(name), "cluster%d", i);
           c = of_get_child_by_name(cluster, name);
           if (c) {
               leaf = false;
               ret = parse cluster(c, depth + 1);
               of_node_put(c);
               if (ret != 0)
                   return ret;
           i++;
•
       } while (c);
       /* Now check for cores */
       i = 0;
•
       do { /*遍历cluster下面core的节点*/
          snprintf(name, sizeof(name), "core%d", i);
           c = of_get_child_by_name(cluster, name);
           if (c) {
               has_cores = true;
```

```
if (depth == 0) {
                   pr_err("%s: cpu-map children should be clusters\n",
                          c->full name);
                   of node put(c);
•
                   return -EINVAL;
•
               /*如果是叶子cluster节点,继续遍历core中的cpu节点*/
•
               if (leaf) {
                   ret = parse_core(c, cluster_id, core_id++);
•
               } else {
                   pr_err("%s: Non-leaf cluster with core %s\n",
                         cluster->full_name, name);
•
                   ret = -EINVAL;
               }
•
               of_node_put(c);
•
               if (ret != 0)
                   return ret;
•
           i++;
       } while (c);
       if (leaf && !has cores)
•
           pr warn("%s: empty cluster\n", cluster->full name);
       if (leaf)
           cluster_id++;
       return 0;
•
   static int __init parse_core(struct device_node *core, int cluster id,
•
                    int core id)
•
•
       char name[10];
       bool leaf = true;
       int i = 0;
       int cpu;
       struct device node *t;
       do { /*如果存在thread层级,解析thread和cpu层级*/
           snprintf(name, sizeof(name), "thread%d", i);
           t = of_get_child_by_name(core, name);
           if (t) -{
•
               leaf = false;
               cpu = get_cpu_for_node(t);
               if (cpu >= 0) {
                   cpu_topology[cpu].cluster_id = cluster_id;
                   cpu_topology[cpu].core_id = core_id;
                   cpu_topology[cpu].thread_id = i;
•
               } else {
                   pr_err("%s: Can't get CPU for thread\n",
•
•
                          t->full name);
                   of node put(t);
                   return -EINVAL;
               of node put(t);
•
•
           i++;
       } while (t);
•
       /*否则直接解析cpu层级*/
       cpu = get_cpu_for_node(core);
       if (cpu >= 0) {
           if (!leaf) {
               pr err("%s: Core has both threads and CPU\n",
                      core->full_name);
              return -EINVAL;
```

```
/*得到了cpu的cluster id/core id*/
           cpu topology[cpu].cluster id = cluster id;
           cpu topology[cpu].core id = core id;
       } else if (leaf) {
           pr_err("%s: Can't get CPU for leaf core\n", core->full name);
           return -EINVAL;
•
       return 0;
   static int __init get_cpu_for_node(struct device node *node)
•
•
       struct device node *cpu node;
       int cpu;
       /*解析cpu层级*/
       cpu_node = of_parse_phandle(node, "cpu", 0);
       if (!cpu_node)
           return -1;
       for each possible cpu(cpu) {
           if (of get cpu node(cpu, NULL) == cpu node) {
               of_node_put(cpu_node);
               return cpu;
       }
       pr crit("Unable to find CPU node for %s\n", cpu node->full name);
       of_node_put(cpu_node);
       return -1;
```

cpu同一层次的关系cpu\_topology[cpu].core\_sibling/thread\_sibling会在 update siblings masks()中更新

```
kernel_init() -> kernel_init_freeable() -> smp_prepare_cpus()
-> store cpu topology() -> update siblings masks()
static void update siblings masks(unsigned int cpuid)
    struct cpu topology *cpu topo, *cpuid topo =
&cpu topology[cpuid];
    int cpu;
    /* update core and thread sibling masks */
    for_each_possible_cpu(cpu) {
        cpu topo = &cpu topology[cpu];
        if (cpuid topo->cluster id != cpu topo->cluster id)
            continue;
        cpumask set cpu(cpuid, &cpu topo->core sibling);
        if (cpu != cpuid)
            cpumask_set_cpu(cpu, &cpuid_topo->core sibling);
        if (cpuid topo->core id != cpu topo->core id)
            continue;
```

### 以我现在的手机平台为例,cpu相关的dts如下:

```
#address-cells = <2>;
#size-cells = <0>;
cpu-map {
   cluster0 {
       core0 {
         cpu = <&CPU0>;
        };
        core1 {
         cpu = < &CPU1>;
        } ;
        core2 {
           cpu = <&CPU2>;
        } ;
        core3 {
            cpu = <&CPU3>;
        } ;
    };
    cluster1 {
        core0 {
         cpu = <&CPU4>;
        } ;
        core1 {
            cpu = <&CPU5>;
        } ;
        core2 {
            cpu = <&CPU6>;
        };
        core3 {
         cpu = <&CPU7>;
        } ;
    };
};
CPU0: cpu@0 {
    device_type = "cpu";
    compatible = "arm, cortex-a55", "arm, armv8";
    reg = <0x0 0x0>;
    enable-method = "psci";
    cpu-supply = <&fan53555_dcdc>;
    cpufreq-data-v1 = <&cpufreq_clus0>;
    cpu-idle-states = <&CORE_PD>;
    sched-energy-costs = <&CPU_COST_0 &CLUSTER_COST_0>;
} ;
CPU1: cpu@100 {
    device_type = "cpu";
```

```
compatible = "arm, cortex-a55", "arm, armv8";
    reg = <0x0 0x100>;
    enable-method = "psci";
    cpu-supply = <&fan53555_dcdc>;
    cpufreq-data-v1 = <&cpufreq clus0>;
    cpu-idle-states = <&CORE PD>;
    sched-energy-costs = <&CPU COST 0 &CLUSTER COST 0>;
};
CPU2: cpu@200 {
    device_type = "cpu";
    compatible = "arm, cortex-a55", "arm, armv8";
    reg = <0x0 0x200>;
    enable-method = "psci";
    cpu-supply = <&fan53555_dcdc>;
    cpufreq-data-v1 = <&cpufreq clus0>;
    cpu-idle-states = <&CORE_PD>;
    sched-energy-costs = <&CPU COST 0 &CLUSTER COST 0>;
} ;
CPU3: cpu@300 {
    device_type = "cpu";
    compatible = "arm, cortex-a55", "arm, armv8";
   reg = <0x0 0x300>;
    enable-method = "psci";
    cpu-supply = <&fan53555 dcdc>;
    cpufreq-data-v1 = <&cpufreq clus0>;
    cpu-idle-states = <&CORE PD>;
   sched-energy-costs = <&CPU COST 0 &CLUSTER COST 0>;
} ;
CPU4: cpu@400 {
    device type = "cpu";
    compatible = "arm, cortex-a55", "arm, armv8";
    reg = <0x0 0x400>;
    enable-method = "psci";
    cpu-supply = <&vddcpu>;
    cpufreq-data-v1 = <&cpufreq clus1>;
    cpu-idle-states = <&CORE PD>;
    sched-energy-costs = <&CPU COST 1 &CLUSTER COST 1>;
};
CPU5: cpu@500 {
    device type = "cpu";
    compatible = "arm, cortex-a55", "arm, armv8";
   reg = <0x0 0x500>;
    enable-method = "psci";
    cpu-supply = <&vddcpu>;
    cpufreq-data-v1 = <&cpufreq_clus1>;
    cpu-idle-states = <&CORE PD>;
    sched-energy-costs = <&CPU_COST_1 &CLUSTER_COST 1>;
CPU6: cpu@600 {
   device type = "cpu";
    compatible = "arm, cortex-a55", "arm, armv8";
    reg = <0x0 0x600>;
    enable-method = "psci";
    cpu-supply = <&vddcpu>;
```

经过parse\_dt\_topology()、update\_siblings\_masks()解析后得到cpu\_topology[}的值为:

```
cpu 0 cluster id = 0, core id = 0, core sibling = 0xf
cpu 1 cluster id = 0, core id = 1, core sibling = 0xf
cpu 2 cluster id = 0, core id = 2, core sibling = 0xf
cpu 3 cluster id = 0, core id = 3, core sibling = 0xf
cpu 4 cluster id = 1, core id = 0, core sibling = 0xf0
cpu 5 cluster id = 1, core id = 1, core sibling = 0xf0
cpu 6 cluster id = 1, core id = 2, core sibling = 0xf0
cpu 7 cluster id = 1, core id = 3, core sibling = 0xf0
//adb 查看cpu topology信息
meizu1000:/sys/devices/system/cpu/cpu0/topology # ls -l
total 0
-r--r-- 1 root root 4096 2018-09-05 16:10 core id
-r--r-- 1 root root 4096 2018-09-05 16:10 core siblings //0f
-r--r-- 1 root root 4096 2018-09-05 16:10 core siblings list//0-3
-r--r-- 1 root root 4096 2018-09-05 16:10 physical package id //0
-r--r-- 1 root root 4096 2018-09-05 16:10 thread siblings //01
-r--r-- 1 root root 4096 2018-09-05 16:10 thread siblings list//0
```

### 二 调度域,调度组以及调度组capacity的初始化

内核中有预先定义的调度域的结构体struct sched\_domain\_topology\_level:

```
typedef const struct cpumask * (*sched domain mask f) (int cpu);
typedef int (*sched domain flags f) (void);
typedef
const struct sched group energy * const(*sched domain energy f) (int cpu);
struct sched domain topology level {
    sched domain mask f mask;
    sched_domain_flags_f sd_flags;
    sched_domain_energy_f energy;
               flags;
    int
               numa level;
    struct sd data data;
#ifdef CONFIG SCHED DEBUG
   char
                       *name;
#endif
```

```
struct sd data {
    struct sched_domain **__percpu sd;
    struct sched group ** percpu sq;
     struct sched group capacity ** percpu sgc;
 };
 * Topology list, bottom-up.
static struct sched domain topology level default topology[] = {
#ifdef CONFIG SCHED SMT
     { cpu smt mask, cpu smt flags, SD INIT NAME(SMT) },
#endif
#ifdef CONFIG SCHED MC
    { cpu_coregroup_mask, cpu_core_flags, SD_INIT_NAME(MC) },
#endif
    { cpu_cpu_mask, SD_INIT_NAME(DIE) },
    { NULL, },
};
static struct sched domain topology level *sched domain topology =
     default topology;
```

我们根据default\_topology[]这个结构体数组来解析sched\_domain\_topology\_level(简称SDTL) 结构体成员变量的含义

- sched\_domain\_mask\_f mask:函数指针,用于指定某个SDTL层级的cpumask位图
- sched\_domain\_flags\_f sd\_flags:函数指针,用于指定某个SDTL层级的标志位
- sched\_domain\_energy\_f energy:函数指针,用于指定某个SDTL层级的energy信息,包括 在active和idle状态下的capacity和power信息,具体获取是从dts(kernel/sched/energy.c) 获取的(energy是怎么获取的)
- struct sd\_data:表示SDTL层级关系,包括domain/group/group capacity 从default\_topology结构体数组看,DIE是标配,MC和SMT是可选择的.目前ARM不支持SMT超线程技术,目的只有两种topology,MC和DIE.

由于在建立topology的时候,涉及到cpumask的使用,下面几个是比较常用的cpumask:

```
static DECLARE_BITMAP(cpu_active_bits, CONFIG_NR_CPUS) __read_mostly;
const struct cpumask *const cpu_active_mask = to_cpumask(cpu_active_bits);
EXPORT_SYMBOL(cpu_active_mask);
```

上面的信息通过他们的名字就可以理解,需要注意的一点就是.online mask与active mask区别是:

- cpu\_online\_mask:表示当前系统online的cpu,包括idle+active cpu
- cpu\_active\_mask:表示当前系统online并且active的cpu

下面看看cpu topology怎么建立起来并经过初始化建立起sched domain和sched group的关系的.当boot cpu启动丛cpu的时候,cpu topology就开始建立:

start\_kernel-->rest\_init-->kernel\_init-->kernel\_init\_freeable-->sched\_init\_smp-->init\_sched\_d omains:

```
* Set up scheduler domains and groups. Callers must hold the hotplug lock.
   * For now this just excludes isolated cpus, but could be used to
   * exclude other special cases in the future.
  static int init sched domains (const struct cpumask *cpu map)
       int err;
       /*ARM没有定义*/
       arch update cpu topology();
       ndoms cur = 1;
       /*分配一个cpumask结构体变量*/
       doms_cur = alloc_sched_domains(ndoms_cur);
       if (!doms cur)
           doms_cur = &fallback_doms;
       /*doms curr[0] = cpu map & (~cpu isolated map), cpumask的与非的计算方式
       cpu isolated map表示有独立的domain的cpu,当前在建立的cpu topology过程中需要
      剔除cpu isolated map cpu*/
       cpumask andnot(doms cur[0], cpu map, cpu isolated map);
       /*开始建立cpu topology*/
       err = build sched domains(doms cur[0], NULL);
       /*注册sched domain的系统控制接口*/
       register_sched_domain_sysctl();
       return err;
   }
    * arch update cpu topology lets virtualized architectures update the
    * cpu core maps. It is supposed to return 1 if the topology changed
    \star or 0 if it stayed the same.
   int weak arch update cpu topology(void)
•
       return 0;
   cpumask_var_t *alloc_sched_domains(unsigned int ndoms)
•
       int i:
       cpumask var t *doms;
       doms = kmalloc(sizeof(*doms) * ndoms, GFP_KERNEL);
       if (!doms)
        return NULL;
```

```
for (i = 0; i < ndoms; i++) {</pre>
           if (!alloc cpumask var(&doms[i], GFP KERNEL)) {
               free sched domains(doms, i);
               return NULL;
•
•
       return doms;
•
   /*上面的cpu map是active cpu mask,那么active cpumask在哪里设置的呢?*/
   /* Called by boot processor to activate the rest. */
   /*arm_dt_init_cpu_maps() 函数会初始化若干个cpumask变量比如possible和present
   cpumask这个函数会将所有的cpu online起来,同时会等到online cpumask和active cpumask*/
   void init smp init(void)
   {
       unsigned int cpu;
       idle_threads_init();
       /* FIXME: This should be done in userspace --RR */
•
       for_each_present_cpu(cpu) {
           if (num_online_cpus() >= setup_max_cpus)
              break;
           if (!cpu online(cpu))
              cpu_up(cpu);
•
       }
       /* Any cleanup work */
       smp announce();
       smp cpus done(setup max cpus);
```

### 接下来解析真正构建cpu topology结构的函数:build\_sched\_domians

```
* Build sched domains for a given set of cpus and attach the sched domains
* to the individual cpus
static int build sched domains (const struct cpumask *cpu map,
                   struct sched domain attr *attr)
   enum s alloc alloc state;
    struct sched_domain *sd;
    struct s data d;
   int i, ret = -ENOMEM;
   alloc state = visit domain allocation hell(&d, cpu map);
    if (alloc state != sa rootdomain)
        goto error;
    /* Set up domains for cpus specified by the cpu map. */
    for each cpu(i, cpu map) {
        struct sched domain topology level *tl;
       sd = NULL;
       for_each_sd_topology(tl) {
            sd = build sched domain(tl, cpu map, attr, sd, i);
            if (tl == sched domain topology)
                *per cpu ptr(d.sd, i) = sd;
```

```
if (tl->flags & SDTL OVERLAP || sched feat(FORCE SD OVERLAP))
            sd->flags |= SD OVERLAP;
   }
}
/* Build the groups for the domains */
for each cpu(i, cpu map) {
    for (sd = *per_cpu_ptr(d.sd, i); sd; sd = sd->parent) {
        sd->span_weight = cpumask_weight(sched_domain_span(sd));
        if (sd->flags & SD_OVERLAP) {
            if (build overlap sched groups(sd, i))
                goto error;
        } else {
            if (build_sched_groups(sd, i))
                goto error;
        }
    }
}
/* Calculate CPU capacity for physical packages and nodes */
for (i = nr cpumask bits-1; i >= 0; i--) {
    struct sched_domain_topology_level *tl = sched_domain_topology;
    if (!cpumask_test_cpu(i, cpu_map))
        continue;
    for (sd = *per cpu ptr(d.sd, i); sd; sd = sd->parent, tl++) {
        init sched energy(i, sd, tl->energy);
        claim allocations(i, sd);
        init_sched_groups_capacity(i, sd);
    }
/* Attach the domains */
rcu read lock();
for_each_cpu(i, cpu_map) {
    int max_cpu = READ_ONCE(d.rd->max_cap_orig_cpu);
    int min_cpu = READ_ONCE(d.rd->min_cap_orig_cpu);
    if ((max cpu < 0) || (cpu rq(i)->cpu capacity orig >
        cpu rq(max cpu) ->cpu capacity orig))
        WRITE ONCE (d.rd->max cap orig cpu, i);
    if ((min cpu < 0) || (cpu rq(i)->cpu capacity orig <</pre>
        cpu_rq(min_cpu) ->cpu_capacity_orig))
        WRITE_ONCE(d.rd->min_cap_orig_cpu, i);
    sd = *per cpu ptr(d.sd, i);
    cpu attach domain(sd, d.rd, i);
rcu read unlock();
ret = 0;
```

```
error:
    __free_domain_allocs(&d, alloc_state, cpu_map);
    return ret;
}
```

对构建cpu topology函数分如下几个部分来分析:

1.分配空间:在build\_sched\_domains函数里面调用\_\_visit\_domain\_allocation\_hell函数

```
static enum s_alloc __visit_domain_allocation_hell(struct s_data *d,
                         const struct cpumask *cpu map)
{
   memset(d, 0, sizeof(*d));
    /*核心函数:创建调度域等数据结构,详细看下面的解析*/
   if ( sdt alloc(cpu map))
       return sa sd storage;
   d->sd = alloc percpu(struct sched domain *);
   if (!d->sd)
       return sa sd storage;
   d->rd = alloc rootdomain();
   if (!d->rd)
       return sa sd;
   return sa rootdomain;
static int      sdt alloc(const struct cpumask *cpu map)
   struct sched domain topology level *tl;
   int j;
static struct sched_domain_topology_level *sched_domain_topology =
     default_topology;
#define for_each_sd_topology(tl)
      for (tl = sched domain topology; tl->mask; tl++)
*/ /*对default topology结构体数据进行遍历*/
   for each sd topology(tl) {
        /*获取每个SDTL层级的sd data数据结构指针*/
       struct sd data *sdd = &tl->data;
       /*下面三个alloc_percpu为每一个SDTL层级的实体数据结构sd_data分配空间
        1.sched domain
        2. sched group
        sched group capacity*/
       sdd->sd = alloc percpu(struct sched domain *);
       if (!sdd->sd)
           return -ENOMEM;
       sdd->sg = alloc percpu(struct sched group *);
       if (!sdd->sg)
           return -ENOMEM;
       sdd->sgc = alloc percpu(struct sched group capacity *);
       if (!sdd->sgc)
           return -ENOMEM;
       /*针对每个active_cpu_mask进行遍历,每个cpu上都需要为sd/sg/sgc分配空间,并存放
       在percpu变量中.注意到sg->next指针,有两个不同的用途:
       1.对于MC 层级,不同cluster的每个cpu core sched group会组成一个链表,比如两个
       cluster,每个cluster 四个core,则每个cluster的sched group链表长度都为4,每个
       core的sched domain都有一个sched group.
```

```
2. DIE层级,不同cluster的sched group会链接起来.一个cluster所有core
         的sched domain都共享一个sched group,后面会以图表说明*/
        for_each_cpu(j, cpu_map) {
            struct sched_domain *sd;
            struct sched group *sg;
            struct sched_group_capacity *sgc;
            sd = kzalloc_node(sizeof(struct sched_domain) + cpumask_size(),
                    GFP KERNEL, cpu to node(j));
            if (!sd)
               return -ENOMEM;
            *per cpu ptr(sdd->sd, j) = sd;
            sg = kzalloc node(sizeof(struct sched group) + cpumask size(),
                   GFP_KERNEL, cpu_to_node(j));
            if (!sg)
               return -ENOMEM;
            sg->next = sg;
            *per cpu ptr(sdd->sg, j) = sg;
            sgc = kzalloc_node(sizeof(struct sched_group_capacity) +
cpumask_size(),
                    GFP KERNEL, cpu to node(j));
            if (!sqc)
               return -ENOMEM;
            *per_cpu_ptr(sdd->sgc, j) = sgc;
       }
    }
    return 0;
} static int __sdt_alloc(const struct cpumask *cpu_map)
    struct sched_domain_topology_level *tl;
    int j;
    for each sd topology(tl) {
        struct sd data *sdd = &tl->data;
        sdd->sd = alloc percpu(struct sched domain *);
        if (!sdd->sd)
            return -ENOMEM;
        sdd->sg = alloc_percpu(struct sched_group *);
        if (!sdd->sg)
            return -ENOMEM;
        sdd->sgc = alloc_percpu(struct sched_group_capacity *);
        if (!sdd->sgc)
            return -ENOMEM;
```

```
for each cpu(j, cpu map) {
            struct sched domain *sd;
            struct sched_group *sg;
            struct sched_group_capacity *sgc;
            sd = kzalloc node(sizeof(struct sched domain) + cpumask size(),
                    GFP KERNEL, cpu to node(j));
            if (!sd)
                return -ENOMEM;
            *per cpu ptr(sdd->sd, j) = sd;
            sq = kzalloc node(sizeof(struct sched group) + cpumask size(),
                   GFP_KERNEL, cpu_to_node(j));
            if (!sg)
               return -ENOMEM;
            sg->next = sg;
            *per cpu ptr(sdd->sg, j) = sg;
            sgc = kzalloc node(sizeof(struct sched group capacity) +
cpumask_size(),
                    GFP KERNEL, cpu to node(j));
            if (!sgc)
               return -ENOMEM;
            *per_cpu_ptr(sdd->sgc, j) = sgc;
        }
   }
   return 0;
```

#### 说明如下:

- 每个SDTL层级都有一个struct sched\_domain\_topology\_level数据结构来描述,并且内嵌了一个struct sd\_data数据结构,包含 sched domain,sched group,sched group capacity的二级指针
- 每个SDTL层级都分配一个percpu变量 sched\_domain,sched\_group,sched\_group\_capacity数据结构
- 在每个SDTL层级中为每个cpu都分配 sched\_domain,sched\_group,sched\_group\_capacity数据结构,即每个cpu在没给SDTL 层级中都有对应的调度域和调度组.
- 2. 初始化sched\_domain:build\_sched\_domain函数,\_\_visit\_domain\_allocation\_hell函数已经为每个SDTL层级分配了sched\_domain,sched\_group/sched\_group\_capacity空间了,有看空间之后,就开始为cpu\_map建立调度域关系了:

```
    build_sched_domains--->build_sched_domain
    /* Set up domains for cpus specified by the cpu_map. */
    /*对每个cpu_map进行遍历,并对每个CPU的SDTL层级进行遍历,相当于每个CPU都有一套SDTL对应的调度域,为每个CPU都初始化一整套SDTL对应的调度域和调度组,每个
    CPU的每个SDTL都调用build_sched_domain函数来建立调度域和调度组*/
    for_each_cpu(i, cpu_map) {
```

```
struct sched domain topology level *tl;
        sd = NULL;
        for_each_sd_topology(tl) {
            sd = build sched domain(tl, cpu map, attr, sd, i);
            if (tl == sched_domain_topology)
               *per cpu ptr(d.sd, i) = sd;
            if (tl->flags & SDTL_OVERLAP || sched_feat(FORCE_SD_OVERLAP))
                sd->flags |= SD OVERLAP;
        }
|->
static struct sched domain *
sd_init(struct sched_domain_topology_level *tl,
   struct sched domain *child, int cpu)
{ /*获取cpu的sched_domain实体*/
    struct sched_domain *sd = *per_cpu_ptr(tl->data.sd, cpu);
    int sd_weight, sd_flags = 0;
#ifdef CONFIG NUMA
     * Ugly hack to pass state to sd numa mask()...
    sched_domains_curr_level = tl->numa_level;
    /*通过default topology[]结构体数据填充对应的参数,涉及到一些cpumask,这个上面已经
    sd_weight = cpumask_weight(tl->mask(cpu));
    if (tl->sd flags)
       sd flags = (*tl->sd flags)();
    if (WARN_ONCE(sd_flags & ~TOPOLOGY_SD_FLAGS,
            "wrong sd flags in topology description\n"))
        sd flags &= ~TOPOLOGY SD FLAGS;
     /*填充sched domain一些成员变量*/
    *sd = (struct sched_domain) {
        .min_interval = sd_weight,
        .max_interval
                          = 2*sd_weight,
                          = 32,
        .busy_factor
        .imbalance_pct
                          = 125,
        .cache nice tries = 0,
        .busy_idx = 0,
.idle_idx = 0,
                       = 0,
        .newidle_idx
        .wake idx
                       = 0,
        .forkexec_idx
                          = 0,
                       = 1*SD LOAD BALANCE
        .flags
                   | 1*SD BALANCE NEWIDLE
                   | 1*SD_BALANCE_EXEC
                   | 1*SD BALANCE FORK
                    | 0*SD_BALANCE_WAKE
                    | 1*SD WAKE AFFINE
```

```
| 0*SD SHARE CPUCAPACITY
                   | 0*SD SHARE PKG RESOURCES
                   | 0*SD SERIALIZE
                   | 0*SD_PREFER_SIBLING
                   | 0*SD NUMA
#ifdef CONFIG_INTEL_DWS
                   | 0*SD INTEL DWS
#endif
                   | sd flags
       .last balance = jiffies,
       .balance interval = sd weight,
       .smt_gain = 0,
       .max_newidle_lb_cost = 0,
       .next_decay_max_lb_cost = jiffies,
       .child = child,
#ifdef CONFIG_INTEL_DWS
 .dws tf = 0,
#endif
#ifdef CONFIG SCHED DEBUG
 .name = tl->name,
#endif
   };
    * Convert topological properties into behaviour.
    if (sd->flags & SD_ASYM_CPUCAPACITY) {
       struct sched_domain *t = sd;
       for each lower domain(t)
          t->flags |= SD BALANCE WAKE;
    }
    if (sd->flags & SD SHARE CPUCAPACITY) {
       sd->flags |= SD_PREFER_SIBLING;
       sd->imbalance_pct = 110;
       sd->smt gain = 1178; /* ~15% */
    } else if (sd->flags & SD SHARE PKG RESOURCES) {
       sd->imbalance_pct = 117;
       sd->cache nice tries = 1;
       sd->busy_idx = 2;
#ifdef CONFIG NUMA
    } else if (sd->flags & SD NUMA) {
       sd->cache_nice_tries = 2;
       sd->busy idx = 3;
       sd->idle_idx = 2;
    sd->flags |= SD_SERIALIZE;
```

```
if (sched domains numa distance[t1->numa level] > RECLAIM DISTANCE)
            sd->flags &= ~(SD BALANCE EXEC |
                       SD BALANCE FORK |
                       SD WAKE AFFINE);
#endif
    } else {
       sd->flags |= SD_PREFER_SIBLING;
       sd->cache nice tries = 1;
       sd->busy idx = 2;
        sd->idle idx = 1;
    }
#ifdef CONFIG INTEL DWS
    if (sd->flags & SD INTEL DWS)
       sd->dws tf = 120;
    if (sd->flags & SD INTEL DWS && sd->flags & SD SHARE PKG RESOURCES)
        sd->dws tf = 160;
#endif
    sd->private = &tl->data;
    return sd;
}
struct sched domain *build sched domain(struct sched domain topology level
*tl,
        const struct cpumask *cpu map, struct sched domain attr *attr,
        struct sched domain *child, int cpu)
    /*上面解释了sd init函数*/
    struct sched domain *sd = sd init(tl, child, cpu);
    cpumask_and(sched_domain_span(sd), cpu_map, tl->mask(cpu));
    if (child) {
        sd->level = child->level + 1;
        sched_domain_level_max = max(sched_domain_level_max, sd->level);
        child->parent = sd;
        if (!cpumask subset(sched domain span(child),
                    sched domain span(sd))) {
            pr err("BUG: arch topology borken\n");
#ifdef CONFIG SCHED DEBUG
            pr err(" the %s domain not a subset of the %s domain\n",
                    child->name, sd->name);
#endif
            /* Fixup, ensure @sd has at least @child cpus. */
            cpumask or (sched domain span(sd),
                   sched_domain_span(sd),
                   sched domain span(child));
        }
```

#### 说明如下:

- sd\_init函数通过default\_topology结构体填充sched\_domain结构体并填充相关的一些参数
- cpumask\_and(sched\_domain\_span(sd), cpu\_map, tl->mask(cpu)); tl->mask(cpu)返回 该cpu某个SDTL层级下兄弟cpu的bitmap,cpumask\_and则将tl->mask(cpu)的bitmap复 制到span数组中.
- 我们知道遍历的顺序是SMT-->MC-->DIE,可以看成父子关系或者上下级关系,struct sched\_domain数据结构中有parent和child成员用于描述此关系.
- 后面完成了对调度域的初始化.
- 3.初始化sched\_group.

```
build_sched_domains--->build_sched_groups
/* Build the groups for the domains */
    for each cpu(i, cpu map) {
        /*per cpu ptr(d.sd, i) 获取每个cpu的最低层级的SDTL,并且通过指针sd->parent遍
历所有层级*/
        for (sd = *per cpu ptr(d.sd, i); sd; sd = sd->parent) {
            /*cpumask and(sched domain span(sd), cpu map, tl->mask(cpu))
     某个SDTL层级cpu兄弟 cpu bitmap,*/
            sd->span_weight = cpumask_weight(sched_domain_span(sd));
            if (sd->flags & SD OVERLAP) {
                if (build overlap sched groups(sd, i))
                    goto error;
             } else {
                if (build sched groups(sd, i))
                    goto error;
        }
    }
-->
/*
 * build sched groups will build a circular linked list of the groups
 ^{\star} covered by the given span, and will set each group's ->cpumask correctly,
 * and ->cpu capacity to 0.
  * Assumes the sched domain tree is fully constructed
 * 核心函数
 */
static int
build_sched_groups(struct sched_domain *sd, int cpu)
    struct sched group *first = NULL, *last = NULL;
    struct sd data *sdd = sd->private;
    const struct cpumask *span = sched_domain_span(sd);
    struct cpumask *covered;
    int i;
```

```
get group(cpu, sdd, &sd->groups);
   atomic inc(&sd->groups->ref);
   if (cpu != cpumask_first(span))
       return 0;
   lockdep assert held(&sched domains mutex);
   covered = sched_domains_tmpmask;
   cpumask clear(covered);
   /*span是某个SDTL层级的cpu bitmap,实质某个sched domain包含多少个cpu*/
   /*两个for循环依次设置了该调度域sd中不同CPU对应的调度组的包含关系,这些调度组通过next
   指针串联起来*/
   for_each_cpu(i, span) {
       struct sched group *sg;
       int group, j;
       if (cpumask_test_cpu(i, covered))
           continue;
       group = get group(i, sdd, &sg);
       cpumask setall(sched group mask(sg));
       for_each_cpu(j, span) {
           if (get_group(j, sdd, NULL) != group)
               continue;
           cpumask_set_cpu(j, covered);
           cpumask_set_cpu(j, sched_group_cpus(sg));
       }
       if (!first)
           first = sg;
       if (last)
           last->next = sq;
       last = sg;
   /*将一个sched domain里面的所有sched group串成一组链表*/
   last->next = first;
   return 0;
static int get group(int cpu, struct sd data *sdd, struct sched group **sg)
   struct sched_domain *sd = *per_cpu_ptr(sdd->sd, cpu);
   struct sched_domain *child = sd->child;
    /*如果是SDTL=DIE,那么child为true*/
   if (child)
       cpu = cpumask_first(sched_domain_span(child));
   if (sg) {
       *sg = *per_cpu_ptr(sdd->sg, cpu);
       (*sg) ->sgc = *per_cpu_ptr(sdd->sgc, cpu);
       atomic_set(&(*sg)->sgc->ref, 1); /* for claim_allocations */
   return cpu;
```

```
• }
•
```

4 初始化sched group capacity:

```
/\!\!\!\!\!^* Calculate CPU capacity for physical packages and nodes ^*/\!\!\!\!
for (i = nr cpumask bits-1; i >= 0; i--) {
    struct sched_domain_topology_level *tl = sched_domain_topology;
    if (!cpumask_test_cpu(i, cpu_map))
        continue;
    for (sd = *per cpu ptr(d.sd, i); sd; sd = sd->parent, tl++) {
        init sched energy(i, sd, tl->energy);
        claim allocations(i, sd);
        init_sched_groups_capacity(i, sd);
/*上面的函数目的是填充下面的结构体每个cpu的每个SDTL层级的sched group capacity*/
struct sched group capacity {
    atomic t ref;
    /*
     * CPU capacity of this group, SCHED LOAD SCALE being max capacity
     * for a single CPU.
    unsigned long capacity;
    unsigned long max_capacity; /* Max per-cpu capacity in group */
    unsigned long min_capacity; /* Min per-CPU capacity in group */
    unsigned long next_update;
    int imbalance; /* XXX unrelated to capacity but shared group state */
     * Number of busy cpus in this group.
     * /
    atomic t nr busy cpus;
    unsigned long cpumask[0]; /* iteration mask */
};
```

#### 5. 将每个cpu的调度域与每个cpu的struct rg root domain成员关联起来

```
/* Attach the domains */
rcu_read_lock();
for_each_cpu(i, cpu_map) {
    int max_cpu = READ_ONCE(d.rd->max_cap_orig_cpu);
    int min_cpu = READ_ONCE(d.rd->min_cap_orig_cpu);

if ((max_cpu < 0) || (cpu_rq(i)->cpu_capacity_orig > cpu_rq(max_cpu)->cpu_capacity_orig))
    WRITE_ONCE(d.rd->max_cap_orig_cpu, i);

if ((min_cpu < 0) || (cpu_rq(i)->cpu_capacity_orig < cpu_rq(min_cpu)->cpu_capacity_orig))
    WRITE_ONCE(d.rd->min_cap_orig_cpu, i);

sd = *per cpu ptr(d.sd, i);
```

```
/*将每个cpu的调度域关联到每个cpu的rq的root domain上*/
        cpu_attach_domain(sd, d.rd, i);
    }
    rcu read unlock();
 * Attach the domain 'sd' to 'cpu' as its base domain.
Callers must
 * hold the hotplug lock.
 * /
static void
cpu attach domain(struct sched domain *sd, struct root domain
*rd, int cpu)
    struct rq *rq = cpu rq(cpu);
    struct sched domain *tmp;
    /* Remove the sched domains which do not contribute to
scheduling. */
    for (tmp = sd; tmp; ) {
        struct sched_domain *parent = tmp->parent;
        if (!parent)
            break;
        if (sd parent degenerate(tmp, parent)) {
             tmp->parent = parent->parent;
             if (parent->parent)
                parent->parent->child = tmp;
             /*
             * Transfer SD PREFER SIBLING down in case of a
             * degenerate parent; the spans match for this
             * so the property transfers.
             * /
             if (parent->flags & SD PREFER SIBLING)
                tmp->flags |= SD PREFER SIBLING;
            destroy sched domain(parent, cpu);
         } else
            tmp = tmp->parent;
    }
    if (sd && sd degenerate(sd)) {
        tmp = sd;
        sd = sd->parent;
        destroy sched domain(tmp, cpu);
        if (sd)
             sd->child = NULL;
    }
    sched_domain_debug(sd, cpu);
    rq attach root(rq, rd);
```

```
tmp = rq->sd;
rcu_assign_pointer(rq->sd, sd);
destroy_sched_domains(tmp, cpu);

update_top_cache_domain(cpu);

#ifdef CONFIG_INTEL_DWS
update_dws_domain(sd, cpu);
#endif
}
```

## 通过调度域,调度组的建立过程,可以通过下面的图来直观的说明SDTL不同层级对应关系:

