动态规划 Dynamic Programming



本节内容

- 1. 递推! 而递归+记忆化是少数情况
- 2. 状态的定义 fib[n]
- 3. 最优子结构 opt[n] = best_of(opt[n-1], opt[n-2], ...)
- 4. 状态转移方程(DP方程)



要点

- 1. 递推! 而递归+记忆化是少数情况
- 2. 状态的定义 fib[n]
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Fibonacci array

• 0, 1, 1, 2, 3, 5, 8, 13, 21, ...

• F[n] = F[n-1] + F[n-2]

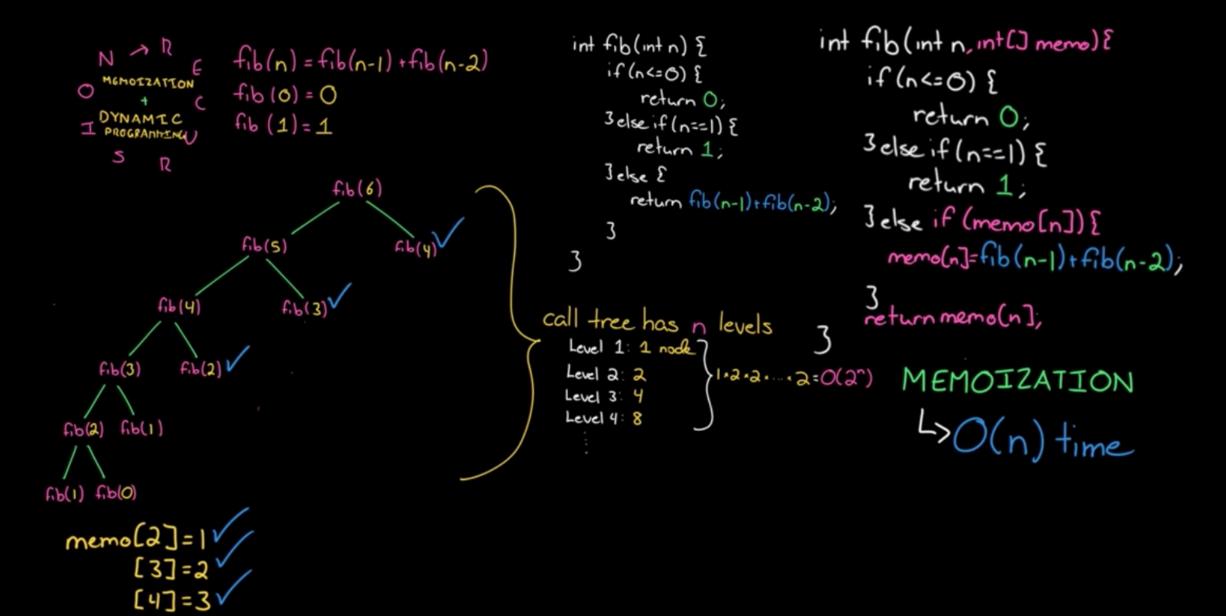


```
int fib (int n) {
                   fib(n) = fib(n-1) + fib(n-2)
                                                                  if (n<=0) {
   MEMOTZATION
0
                   f_1b(0) = 0
                                                                      return 0;
  DYNAMIC
                                                                  3 else if (n==1) {
                   fib (1) = 1
I PROGRAMMING
                                                                      return 1;
    S
                                                                  Jelse &
                                f.p(6)
                                                                     return fib(n-1)+fib(n-2);
                    fib(s)
                                       F.P(A)
                                              f.b(2)
         fib (4)
                                      fib(3)
                         f(P(3))
                                                          call tree has n levels
                                                              Level 1: 1 node 7
                       fib(2) fib(1) fib(2)
  f.b(3)
            f.b(2)
                                                                               >1*2 *2 *... * 2 = O(2^)
                                                             Level 2: 2
                                 t.P(1) t.P(0)
                                                             Level 3: 4
                      £₽(1) £₽(0)
                                                             Level 4: 8
    4P(1) 4P(1) 4P(0)
```

- int fib(int n) {
- return n <= 1 ? n : fib(n 1) + fib(n 2);
- }



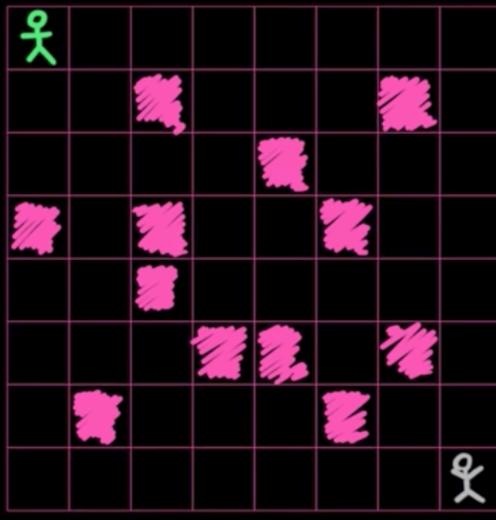
```
int fib (int n) {
                   fib(n) = fib(n-1) + fib(n-2)
                                                                  if (n<=0) {
   MEMOTZATTON
0
                   fib(0) = 0
                                                                      return 0;
  DYNAMIC
                                                                 3 else if (n==1) {
                   fib (1)=1
I PROGRAMMENCY
                                                                     return 1;
    S
          R
                                                                 Jelse &
                               f.p(6)
                                                                     return fib(n-1)+fib(n-2);
                    fib(5)
                                       F.P(A)
                                              f.b(2)
         fib (4)
                                      f.b(3)
                         f.b(3)
                                                          call tree has n levels
                                                             Level 1: 1 node 7
                                        trp(1) trp(1) trp(0)
                        fib(2) fib(1)
  f.b(3)
            f.b(2)
                                                                              >1,2,2, =0(2")
                                                             Level a: 2
                                 £P(1) €P(0)
                                                             Level 3: 4
                      t.P(1) t.P(0)
                                                             Level 4: 8
    tip(1) tip(1) tip(0)
```



[5]=5

COUNT THE PATHS

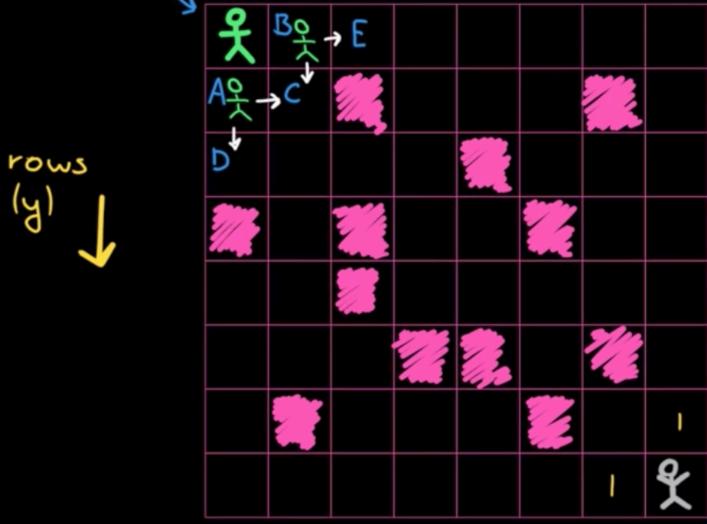
start





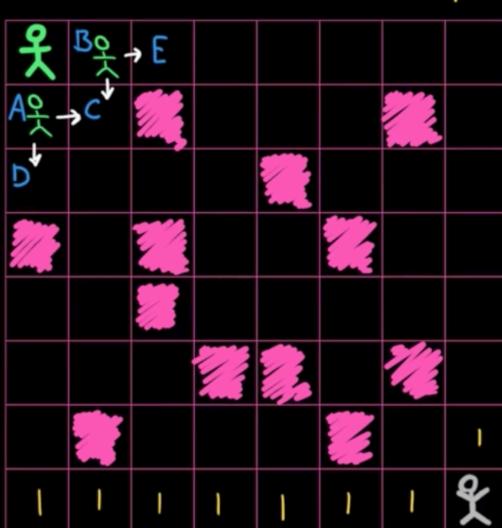
```
COUNT THE PATHS
               paths(start, end) =
                     + paths(B,enQ)
   paths(A,en@)
paths(D,end)+paths(C,end) paths(C,end)+paths(E,end)
         int countPaths(boolean[]] grid, int row, int col) {
           if (!validSquare(grid, row, col)) return 0;
          if (isAtEnd(grid, row, col)) return 1;
           return countPaths(grid, row+1, col) + countPaths(grid, row, col+1);
```

COUNT THE PATHS columns (x)





COUNT THE PATHS columns (x)







$$opt[i, j] = opt[i + 1, j] + opt[i, j + 1]$$

```
if a[i, j] = '空地':
    opt[i , j] = opt[i + 1, j] + opt[i, j + 1]
    else: // 石头
    opt[i , j] = 0
```



COUNT THE PATHS

columns (x) start

	_	옷	Bo.	• E					
		A옷-	, C	獲					
25		D				獲			
				X			製		
				**					
						意		1/4	
			青	3	a	ı		2	1
		ı	1	ı	1	1	1	1	¥

row



COUNT THE PATHS

٨	Ļ			L
0	Ľ	×	r	T
		۰		
			`	٠

	,d/						
옷	17	12	12	7	4	1	1
10	5	灩	5	3	3	黨	1
5	S	2	2	覆	3	3	١
	3	業	2	1	製	2	1
フ	3	藩		1	1	1	1
4	3	3	7	黎	0	1/4	ı
1	膏	3	a	ı		2	1
١	1	1	1	1	1	1	4



动态规划 Dynamic Programming

- 1. 递推! (recursion)
- 2. 状态的定义 fib[n]
- 3. 最优子结构 optimized = fib[n]
- 4. 状态转移方程(DP方程: fib[n] = fib[n-1] + fib[n-2])



动态规划 Dynamic Programming

1. 打破自己的思维惯性,形成机器思维

2. 理解复杂逻辑的关键

3. 也是职业进阶的要点要领



01 背包问题

- •假设山洞里共有a,b,c,d,e这5件宝物(不是5种宝物),它们的重量分别是 2,2,6,5,4,它们的价值分别是6,3,5,4,6,现在给你个承重为10的背包, 怎么装背 包,可以才能带走最多的财富。
- 01背包的状态转换方程 f[i,j] = Max{ f[i-1,j-Wi]+Pi(j >= Wi), f[i-1,j] }
- f[i,j]表示在前i件物品中选择若干件放在承重为 j 的背包中,可以取得的最大价值。
- Pi表示第i件物品的价值。
- 决策:为了背包中物品总价值最大化,第 i件物品应该放入背包中吗?



预习题目(上课考察)

- 1. https://leetcode.com/problems/best-time-to-buy-and-sell-stock/#/description
- 2. https://leetcode.com/problems/climbing-stairs/description/
- 3. https://leetcode.com/problems/triangle/description/
- 4. https://leetcode.com/problems/maximum-productsubarray/description/
- 5. https://leetcode.com/problems/coin-change/description/



预习题目(上课考察)

- 1. https://leetcode.com/problems/house-robber/
- 2. https://leetcode.com/problems/house-robber-ii/description/
- 3. https://leetcode.com/problems/best-time-to-buy-and-sell-stock/#/description
- 4. https://leetcode.com/problems/best-time-to-buy-and-sell-stock-ii/
- 5. https://leetcode.com/problems/best-time-to-buy-and-sell-stock-iii/
- 6. https://leetcode.com/problems/best-time-to-buy-and-sell-stock-with-cooldown/
- 7. https://leetcode.com/problems/best-time-to-buy-and-sell-stock-iv/
- 8. https://leetcode.com/problems/best-time-to-buy-and-sell-stock-with-transaction-fee/



实战题目

- 1. https://leetcode-cn.com/problems/perfect-squares/
- 2. https://leetcode-cn.com/problems/burst-balloons/
- 3. https://leetcode-cn.com/problems/jump-game/
- 4. https://leetcode-cn.com/problems/jump-game-ii/
- 5. https://leetcode-cn.com/problems/unique-paths/
- 6. https://leetcode-cn.com/problems/unique-paths-ii/
- 7. https://leetcode-cn.com/problems/unique-paths-iii/
- 8. https://leetcode-cn.com/problems/coin-change/
- 9. https://leetcode-cn.com/problems/coin-change-2/



实战题目

- 1. https://leetcode-cn.com/problems/longest-valid-parentheses/
- 2. https://leetcode-cn.com/problems/minimum-path-sum/
- 3. https://leetcode-cn.com/problems/edit-distance/
- 4. https://leetcode-cn.com/problems/decode-ways
- 5. https://leetcode-cn.com/problems/maximal-square/
- 6. https://leetcode-cn.com/problems/max-sum-of-rectangle-no-larger-than-k/
- 7. https://leetcode-cn.com/problems/frog-jump/
- 8. https://leetcode.com/problems/split-array-largest-sum
- 9. https://leetcode-cn.com/problems/student-attendance-record-ii/
- 10.https://leetcode-cn.com/problems/task-scheduler/
- 11.https://leetcode-cn.com/problems/palindromic-substrings/
- 12.https://leetcode-cn.com/problems/minimum-window-substring/



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