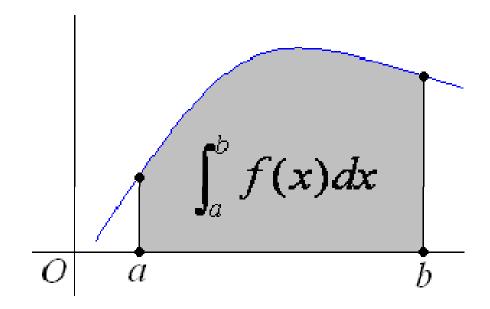
第5章 数值积分

求函数 f(x) 在区间 [a, b] 上的定积分

$$I(f) = \int_{a}^{b} f(x) dx$$



定积分的定义

$$\int_{a}^{b} f(x) dx = \lim_{n \to \infty} \sum_{k=0}^{n-1} f(x_k) \Delta x_k$$

Newton-Leibniz 公式

$$\int_{a}^{b} f(x)dx = F(b) - F(a), \quad F'(x) = f(x)$$

近似计算公式

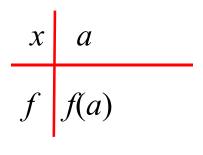
$$\int_{a}^{b} f(x) dx \approx \sum_{k=0}^{n-1} f(x_k) \Delta x_k$$

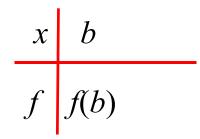
数值积分的必要性:

- f(x)的原函数不能用初等函数表示
- f(x)及其原函数的表达式很复杂
- f(x)是以表格形式给出

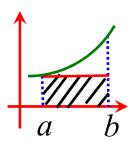
单个公式

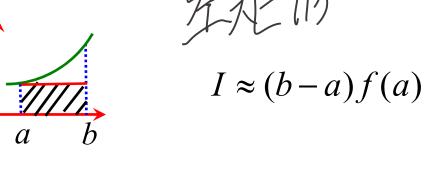
1. 矩形公式

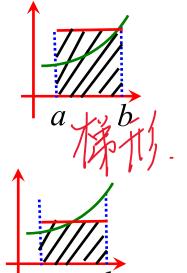




$$\begin{array}{c|c} x & (b+a)/2 \\ \hline f & f((b+a)/2) \end{array}$$







$$I \approx (b-a)f(b)$$

$$I \approx (b-a)f(b)$$

$$I \approx (b-a)f(b)$$

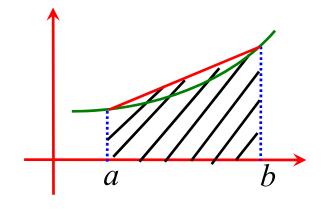
$$I \approx (b+a)$$

$$I \approx (b-a)f\left(\frac{b+a}{2}\right)$$

单个公式

2. 梯形公式

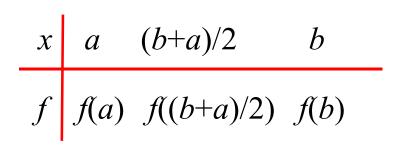
$$\begin{array}{c|ccc}
x & a & b \\
\hline
f & f(a) & f(b)
\end{array}$$



$$I \approx \frac{(b-a)}{2} [f(a) + f(b)]$$

单个公式

3. Simpson公式



$$\frac{1}{a}$$
 $\frac{(b+a)}{2}$ $\frac{b}{b}$

$$I \approx \frac{(b-a)}{6} \left[f(a) + 4f\left(\frac{a+b}{2}\right) + f(b) \right]$$

机械求积公式

$$I_n \approx \sum_{k=0}^n A_k f_k$$
 对节点 x_k 处的值进行加权平均

 x_{k} -- 积分节点; A_{k} -- 求积系数

 A_k 仅与节点值及区间 [a, b] 有关,而与被积函数 f(x)无关

$$R_n(f) = I(f) - I_n = \int_a^b f(x) dx - \sum_{k=0}^n A_k f(x_k)$$

代数精度(名称与测试逐步型有关)。

• 如果求积公式对于任何次数不高于m的多项式都精确成立,而对某个m+1次多项式不能精确成立,则称求积公式具有m次代数精度。

当
$$f(x) = 1, x, x^2, \dots x^m$$
 时

$$\int_{a}^{b} f(x) dx = \sum_{k=0}^{n} A_{k} f(x_{k})$$
 精确成立

$$\int_{a}^{b} x^{m+1} dx \neq \sum_{k=0}^{n} A_{k} x_{k}^{m+1}$$

几个常用的求积公式的代数精度

1. 梯形公式的代数精度

形 公 氏 的 代 剱 楠 段
当
$$f(x) = x$$
时
$$\int_{a}^{b} f(x)dx = \int_{a}^{b} xdx = \frac{1}{2}x^{2} \Big|_{a}^{b} = \frac{1}{2}(b^{2} - a^{2}) \qquad f(\chi) = 1$$

$$T[f] = \frac{b - a}{2}(f(a) + f(b)) = \frac{b - a}{2}(a + b) = \int_{a}^{b} f(x)dx \qquad \int_{a}^{b} f(x)dx \qquad$$

当
$$f(x) = x^2$$
时

$$\int_{a}^{b} f(x) dx = \int_{a}^{b} x^{2} dx = \frac{1}{3}x^{3} \Big|_{a}^{b} = \frac{1}{3}(b^{3} - a^{3})$$

$$\int_{a}^{b} f(x) dx = \int_{a}^{b} x^{2} dx = \frac{1}{3}x^{3} \Big|_{a}^{b} = \frac{1}{3}(b^{3} - a^{3})$$

$$\int_{a}^{b} f(x) dx = \int_{a}^{b} x^{2} dx = \frac{1}{3}x^{3} \Big|_{a}^{b} = \frac{1}{3}(b^{3} - a^{3})$$

$$T[f] = \frac{b-a}{2}(f(a)+f(b)) = \frac{b-a}{2}(a^2+b^2) \neq \int_a^b f(x)dx \int_a^b f(x)dx$$

$$\frac{b-a}{2} \cdot \frac{1}{2} (f(a)+f(c)) + \frac{b-a}{2} \cdot \frac{1}{2} (f(c)+f(b)) \\
= \frac{b-a}{4} (f(a)+2f(c)+f(b)) \\
f(x)=| I=b-a. \\
I_2=b-a. \\
f(x)= X I=\frac{1}{2} (b^2-a^2) \\
I_2=\frac{b-a}{4} (a+b+a+b) \\
= \frac{1}{2} (b^2-a^2)$$

$$f(x) = x^{2} \qquad I = \frac{1}{3}(b^{3} - a^{3}).$$

$$I_{2} = \frac{b-a}{4}(a^{2}+b^{2}+\frac{(a+b)^{2}}{2})$$

$$= \frac{b-a}{4}(a^{2}+b^{2}+\frac{1}{2}(a^{2}+b^{2}+2ab))$$

$$= \frac{b-a}{4}(\frac{3}{2}a^{2}+\frac{3}{2}b^{3}+ab)$$

$$I_{2} \neq I \qquad M=1$$

2. Simpson-公式的代数精度

当
$$f(x) = x$$
时
$$\int_{a}^{b} f(x)dx = \int_{a}^{b} xdx = \frac{1}{2}(b^{2} - a^{2})$$

$$S[f] = \frac{b-a}{6}(f(a) + 4f(\frac{b+a}{2}) + f(b))$$

$$= \frac{b-a}{6}(a+4\frac{a+b}{2}+b) = \frac{1}{2}(b^{2} - a^{2})$$
所以 $\int_{a}^{b} f(x)dx = S[f]$ 成立

当
$$f(x) = x^2$$
时
$$\int_a^b f(x)dx = \int_a^b x^2 dx = \frac{1}{3}(b^3 - a^3)$$

$$S[f] = \frac{b-a}{6}(f(a) + 4f(\frac{b+a}{2}) + f(b))$$

$$= \frac{b-a}{6}(a^2 + 4(\frac{a+b}{2})^2 + b^2)$$

$$= \frac{b-a}{6}(2a^2 + 4ab + 2b^2) = \frac{1}{3}(b^3 - a^3)$$
即 $\int_a^b f(x)dx = S[f]$ 精确成立

当
$$f(x) = x^3$$
时
$$\int_a^b f(x)dx = \int_a^b x^3 dx = \frac{1}{4}(b^4 - a^4)$$

$$S[f] = \frac{b-a}{6}(f(a) + 4f(\frac{b+a}{2}) + f(b))$$

$$= \frac{b-a}{6}(a^3 + 4(\frac{a+b}{2})^3 + b^3)$$

$$= \frac{b-a}{6}(a^3 + \frac{1}{2}(a^3 + 3a^2b + 3ab^2 + b^3) + b^3)$$

$$= \frac{b-a}{6}\frac{3}{2}(a^3 + a^2b + ab^2 + b^3) = \frac{1}{4}(b^4 - a^4)$$
即 $\int_a^b f(x)dx = S[f]$ 精确成立

In=Asto tAfit ... + Anfin

构造n阶代数精度的待定系数法

$$f(x) = 1: \int_{a}^{b} dx = (b-a) = A_{0} + A_{1} + \dots + A_{n}$$

$$x: \int_{a}^{b} x dx = \frac{b^{2} - a^{2}}{2} = A_{0}x_{0} + A_{1}x_{1} + \dots + A_{n}x_{n}$$

$$x^{2}: \int_{a}^{b} x^{2} dx = \frac{b^{3} - a^{3}}{3} = A_{0}x_{0}^{2} + A_{1}x_{1}^{2} + \dots + A_{n}x_{n}^{2}$$

$$\therefore \mathcal{N}: \int_{a}^{b} \mathcal{N} d\chi = \frac{b^{n+1} - a^{n+1}}{2}$$

$$= \left[\begin{array}{ccccc} 1 & 1 & \dots & 1 \\ x_{0} & x_{1} & \dots & x_{n} \\ \dots & \dots & \dots & \dots \\ x_{0}^{n} & x_{1}^{n} & \dots & x_{n}^{n} \end{array} \right] \begin{bmatrix} A_{0} \\ A_{1} \\ \dots \\ A_{n} \end{bmatrix} = \begin{bmatrix} b - a \\ (b^{2} - a^{2})/2 \\ \dots \\ (b^{n+1} - a^{n+1})/(n+1) \end{bmatrix}$$

届了一个3/2代数楼度积分线 $\int_{0}^{5} f(x) dx = A_{5}f(0) + A_{1}f(1) + A_{2}f(2) + A_{3}f(3)$ 3 = Abt At Azt Az + Az. 0 1 2 3 +(x)=| $f(x)=x^{3^{2}}=A_{1}+2A_{2}+3A_{3}$ $f(x)=x^2$ = A₁ + 4A₂+9A₃ f(X)=x = A,· otA, + 8A2+27A3 由处标(性得) A。=A3 A1=A= $|3A|+3A_3 = \frac{9}{2}$ $|5A|+9A_3 = 9$ $|9A|+27A_3 = \frac{81}{4}$ A3= 3 A1= 3

C,X浸偿还的校童确定/f(x)dx= C[f(X,)+f(X,)+f(X)]使得代数精影量高. 取f(X)=1, X, X, X, X, 使得公式准确就立 f(x)=1 3C=2 f(x)=x C[x₀+x₁+x₂]=0 $f(x) = x^2 \left(\frac{1}{12} \left(\frac{1}{12} x_0 + \frac{1}{12} x_1 + \frac{1}{12} x_2 \right) = \frac{3}{5}$ $f(x) = x^3 ((x_0^3 + x_1^3 + x_2^3) = 0.$ C=3 · 12-15 x02 X15 X15 X15 $\chi_0 + \chi_1 + \chi_2 = 0$ X3+X1+X2=1 入3+入13+入23二〇、 $\chi_0 = -\chi_1$ $\chi_1 = 0$ 270^{2} 70^{2} 70^{2} X2=5

f(X)=X4 = (x,+x,+x,+)+= 求Ao,A,Az使信XXX=Af(w)+Af(cc)构加 例数指度是言 $+(x) = 1, X, X^{2}$ $A_0 + A_1 + A_2 = b - a$ A. a+A, C+A2h=b-a2 $A_0 a^2 + A_1 c^2 + A_2 b^2 = \frac{1}{37} b^3 - a^3$ 又批外 人。二人 $\frac{1}{2}A_{b}+A_{1}=b-a$. $\frac{1}{2}(a+b)A_{b}+A_{1}(c=\frac{1}{2}b^{2}-a^{2})$ $\frac{1}{2}(a+b^{2})A_{b}+A_{1}(c^{2}=\frac{1}{2}(b^{3}-a^{3})$.

$$\begin{cases}
2A_{0} + A_{1} = b - \alpha \\
A_{0} + \frac{1}{2}A_{1} = \frac{1}{2}(b - \alpha)
\end{cases}$$

$$A_{1} = b - \alpha - 2A_{3}$$

$$(a^{2} + b^{2}) A_{0} + (b - \alpha - 2A_{0})(\frac{\alpha + b^{2}}{2}) = \frac{1}{3}(b^{3} - \alpha^{3})$$

$$A_{3} = \frac{1}{6}(b - \alpha)$$

$$A_{1} = \frac{1}{6}(b - \alpha)$$

• 多项式插值

插值函数类是多种多样的,一般根据问题的特征与研究的要求来选择。最常用到的是代多项式函数插值,多项式函数形式简单,便于计算。

插值函数是 n 次多项式

$$p_n(x) = a_0 + a_1 x + \dots + a_n x^n$$

 $a_0, a_1, \dots, a_n \longrightarrow$ 待定系数

由插值条件得

$$\begin{cases} a_0 + a_1 x_0 + a_2 x_0^2 + \dots + a_n x_0^n = y_0 \\ a_0 + a_1 x_1 + a_2 x_1^2 + \dots + a_n x_1^n = y_1 \\ \dots \\ a_0 + a_1 x_n + a_2 x_n^2 + \dots + a_n x_n^n = y_n \end{cases}$$

系数矩阵 (Vandermonde-- 范德蒙德 行列式)

$$|D| = \begin{vmatrix} 1 & x_0 & x_0^2 & \cdots & x_0^n \\ 1 & x_1 & x_1^2 & \cdots & x_1^n \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 1 & x_n & x_n^2 & \cdots & x_n^n \end{vmatrix} = \prod_{0 \le i < j \le n} (x_j - x_i)$$

插值方法

方法
$$L_n(x) = \sum_{k=0}^n f(x_k) \cdot l_k(x)$$

$$= \int_0^n \int_{\mathbb{R}^{2}} f(x_k) \left(\frac{1}{k} \right) \left(\frac{1}{k} \right) \left(\frac{1}{k} \right) \left(\frac{1}{k} \right) dx + \mathcal{L}_n(x) = \sum_{k=0}^n f(x_k) \cdot l_k(x) + \frac{f^{(n+1)}(\xi)}{(n+1)!} \omega_{n+1}(x) \qquad (a < \xi < b) \right)$$

$$A_k = \int_a^b l_k(x) dx = \int_a^b \frac{1}{\omega'_{n+1}(x_k)} \cdot \frac{\omega_{n+1}(x)}{(x - x_k)} dx$$

$$R_n(f) = \int_a^b \frac{f^{(n+1)}(\xi)}{(n+1)!} \omega_{n+1}(x) dx \qquad (a < \xi < b)$$

$$\begin{array}{c|cccc}
x & x_{0} = a \\
\hline
f & f(a) \\
L_{0}(x) = f(x_{0}) = f(a) \\
L_{0}(x) = f(x_{0}) = f(a) \\
L_{0}(x) = f(a) \\
\hline
f & f(b) \\
\hline
f & f(a) \\
L_{1}(x) = f(a) (a, cx) + f(b) (a, cx) + f(b) (a, cx) \\
= f(a) & x - b \\
A_{0} = \int_{a}^{b} L_{0}(x) dx - \frac{1}{a - b} \int_{a}^{b} (x - b) dx \\
= \frac{1}{a - b} \left[\frac{b^{2} - a^{2}}{2} - b(b - a) \right]
\end{array}$$

$$= -\frac{\alpha + b}{2} + b = \frac{b-\alpha}{2}$$

$$A_1 = \frac{b-\alpha}{2}$$

$$\chi_0 = \alpha$$
 $\chi_1 = \zeta$ $\chi_2 = b$

$$X_1 = C$$
 $X_2 = k$

$$\begin{cases}
6(x) = \frac{(x-c)(x-b)}{(a-c)(a-b)} \\
A_0 = \int_{a}^{b} \frac{b-a}{b} \\
A_1 = \frac{4(b-a)}{b} \\
A_2 = \frac{b-a}{b}
\end{cases}$$

Newton Cotes 积分公式(等距节点)

$$[a,b]$$
区间等分n等分,取 $h = \frac{b-a}{n}, x_j = a+kh$

$$(j = 0,1,2...,n)$$

$$f(x) = L_n(x) + R_n(x)$$

$$\int_a^b f(x)dx = \int_a^b L_n(x)dx + \int_a^b R_n(x)dx$$

$$= (b-a)\sum_{j=0}^n \binom{n}{j} f_j + R[f]$$

$$Cotes 系数$$

$$A_j = \int_a^b \int_a (x)dx$$

$$A_j = \int_a^b \int_a (x)dx$$

$$\begin{array}{ll}
-(b-a) \stackrel{>}{\underset{j=0}{\sum}} \stackrel{>}{A} f(x_j) \\
-(b-a) \stackrel{>}{\underset{j=0}{\sum}} \stackrel{>}{A} f(x_j) \\
-(a+i) \stackrel{>}{\underset{j=0}{\sum}} \stackrel{>}{\underset{j=0}{\sum}$$

$$C_j^{(n)} = \frac{1}{b-a} \int_a^b l_j(x) dx = \frac{1}{b-a} \int_a^b \prod_{\substack{i=0\\i\neq j}}^n \frac{x - x_i}{x_j - x_i} dx$$

$$x = a + th$$

$$C_{j}^{(n)} = \frac{1}{b-a} \int_{a}^{b} \prod_{i=0}^{n} \frac{x-x_{i}}{x_{j}-x_{i}} dx = \frac{1}{nh} \int_{0}^{n} \prod_{i=0}^{n} \frac{t-i}{j-i} h dt$$

$$= \frac{1}{n} \prod_{i=0}^{n} \frac{1}{j-i} \int_{0}^{n} \prod_{i=0}^{n} (t-i) dt$$

$$N_{j} = \frac{(-1)^{n-j}}{nj!(n-j)!} \int_{0}^{n} \prod_{i=0}^{n} (t-i) dt$$

 $\mathbf{n} = 1$ 时,仅有两个节点:

$$C_0^{(1)} = \frac{(-1)^{1-0}}{1 \times 0! \times (1-0)!} \int_0^1 (t-1)dt = \frac{-1}{1} \frac{(t-1)^2}{2} \Big|_0^1 = \frac{1}{2}$$

$$C_1^{(1)} = \frac{(-1)^{1-1}}{1 \times 1! \times (1-1)!} \int_0^1 (t-0)dt = \frac{1}{1} \frac{t^2}{2} \Big|_0^1 = \frac{1}{2}$$

当
$$n=2$$
时

$$C_0^{(2)} = \frac{(-1)^{2-0}}{2 \times 0! \times (2-0)!} \int_0^2 (t-1)(t-2) dt$$

$$= \frac{1}{4} \int_0^2 [(t-2)^2 + (t-2)] dt$$

$$= \frac{1}{4} \left[\frac{1}{3} (t-2)^3 + \frac{1}{2} (t-2)^2 \right] \Big|_0^2 = \frac{1}{6}$$
同理可得 $C_1^{(2)} = \frac{4}{6}$, $C_2^{(2)} = \frac{1}{6}$

• 以此类推得Cotes系数表:

\overline{n}	$C_k^{(n)}$
1	$\frac{1}{2}\{1,1\}$
2	$\frac{1}{6}\{1,4,1\}$
3	$\frac{1}{8}\{1,3,3,1\}$
4	$\frac{1}{90} \{7,32,12,32,7\}$
5	$\frac{1}{288} \{19, 75, 50, 50, 75, 19\}$
6	$\frac{1}{840} \{41, 216, 27, 272, 27, 216, 41\}$
7	$\frac{1}{17280} \{ 751, 3577, 1323, 2989, 2989, 1323, 3577, 751 \}$
8	$\frac{1}{28350} \{ 989, 5888, -928, 10496, -4540, 10496, -928, 5888, 989 \}$

常用的几个积分公式

• 梯形公式(n=1)

• Simpson 公式 (n=2)

因为
$$C_0^{(2)} = \frac{1}{6}, C_1^{(2)} = \frac{4}{6}, C_2^{(2)} = \frac{1}{6}$$
所以 $\int_a^b f(x)dx = S[f] + R_S[f]$
且 $S[f] = \frac{b-a}{6}(f(a) + 4f(\frac{a+b}{2}) + f(b))$
 $R_S[f] = -\frac{(b-a)^5}{2880}f^{(4)}(\xi) \quad \xi \in (a,b)$

• Newton 公式 (n=3)

因为
$$C_0^{(3)} = \frac{1}{8}, C_1^{(3)} = \frac{3}{8}, C_2^{(3)} = \frac{3}{8}, C_3^{(3)} = \frac{1}{8}$$
所以 $\int_a^b f(x)dx = N[f] + R_N[f]$
且 $N[f] = \frac{b-a}{8}(f(a) + 3f(a+h) + 3f(a+2h) + f(b))$
其中 $h = \frac{b-a}{3}$ 。

• Cotes 公式 (n=4)

因为
$$C_0^{(4)} = \frac{7}{90}$$
, $C_1^{(4)} = \frac{32}{90}$, $C_2^{(4)} = \frac{12}{90}$, $C_3^{(4)} = \frac{32}{90}$, $C_4^{(4)} = \frac{7}{90}$ 所以 $\int_a^b f(x)dx = C[f] + R_C[f]$ 且 $C[f] = \frac{b-a}{90} (7f(a) + 32f(a+h) + 12f(a+2h) + 32f(a+2h) + 7f(b))$ 其中 $h = \frac{b-a}{4}$ 。

例

计算
$$I = \int_1^2 \frac{1}{r} dx$$
 。

解: 由Newton-Leibniz公式得

$$I = \int_{1}^{2} \frac{1}{x} dx = \ln 2 = 0.69314718$$

由梯形公式
$$I = \frac{1}{2}(\frac{1}{2} + \frac{1}{1}) = 0.75;$$

曲
$$Simpson$$
公式 $I = \frac{1}{6}(\frac{1}{1} + 4\frac{1}{3} + \frac{1}{2}) = 0.6944;$

曲 Newton 公式
$$I = \frac{1}{8}(1+3\frac{1}{4}+3\frac{1}{5}+\frac{1}{2}) = 0.69375$$

由Cotes公式得I=0.693175

$$\int_{0}^{0.5} \frac{1}{(x+1)} \frac{1}$$

$$f(t)=0.84548 \quad f(t)=0.7[15]$$

$$I = \frac{0.5}{8}[1+3\times0.84548+3\times0.7115] + 0.496285].$$

$$\int_{0.5}^{1} \sin(x^{2}) dx$$

$$Q=0, b=1$$

$$I = \frac{b-a}{90}(7f(0)+32f(0.25)+72f(0.5) + 32f(0.75)+72f(1)]$$

$$= 0.3[0.26]...$$

$$Rn[H]=h^{1+3}f^{(n+2)}(1)f^{(H-\frac{1}{2})(H-1)}...$$

$$h=H^{1}f^{(n+2)}(1-1)(1-2)(1-3)(1-4)dt$$

$$h=H^{1}f^{(H-\frac{1}{2})(1-1)(1-2)(1-3)(1-4)dt}$$

$$|R| \leq |H22|X|0^{-3}.$$

Newton-Cotes公式截断误差及代数精度 代数精度:f(X)=1;X,X,···X 拉格础划的是以积分起间

Newton – Cotes 公式的截断误差为
$$\{(x)=1, x, x, y, y\}$$
 $\{(x)=1, x, y\}$ $\{(x)=1, x, y\}$

其中,
$$h = \frac{b-a}{n}, \eta \in [a,b]$$

令n = 4 得Cotes 求积公式的截断误差

$$R_C[f] = -\frac{(b-a)^7}{483840} f^{(6)}(\eta) \quad \eta \in [a,b]$$

定理: 当n 为偶数时,Newton-Cotes 公式具有n+1 次代数精度。证明 只需证明当 $f(x)=x^{n+1}$ 时,余式R[f]=0。

注意到此时 $f^{n+1}(\xi) = (n+1)!$ 则

$$R[f] = \int_{a}^{b} \omega(x) dx = h^{n+2} \int_{0}^{n} t(t-1)...(t-n) dt$$

因为n为偶数,所以令n = 2m, u = t - m

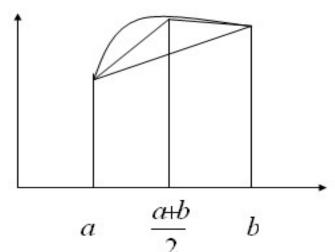
$$\mathbb{M} \quad R[f] = h^{2m+2} \int_{-m}^{m} u(u^2 - 1)...(u^2 - m^2) du = 0$$

复化求积公式

• 将区间[a,b]适当分割成若干个字区间,对每个子区间使用求积公式,构成所谓的复化求积公式,这是提高积分精度的一个常用的方法。

定步长复化求积公式

1.复化梯形求积公式



$$T(h) = \frac{a-b}{2}(f(a) + f(b))$$

$$T(\frac{h}{2}) = \frac{\frac{a-b}{2}}{2}(f(a) + f(\frac{a+b}{2}))$$

$$\frac{a-b}{2}(f(\frac{a+b}{2}) + f(b))$$

$$= \frac{b-a}{4}(f(a) + 2f(\frac{a+b}{2}) + f(b))$$

• 一般地将[a,b]区间n等分,则

$$h = \frac{1}{n} \cdot x_j = a + jh$$
 $(j = 0,1,2,...n)$ 对每个子区间 $[x_{j-1},x_j]$ $(j = 1,2,...n)$ 使用 T 公式有 $S_j = \frac{h}{2}(f(x_{j-1}) + f(x_j))$

• 斯以
$$\int_{a}^{b} f(x) dx = \sum_{j=1}^{n} S_{j} + R_{T}[f, h]$$

$$\overline{m} \sum_{j=1}^{n} S_{j} = \frac{h}{2} \sum_{j=1}^{n} (f(x_{j-1}) + f(x_{j}))$$

$$= \frac{h}{2} (f(a) + f(x_{1}) + f(x_{1}) + f(x_{2}) + \dots + f(x_{n-1}) + f(b) + f(b) - f(b))$$

$$= \frac{h}{2} (f(a) - f(b) + 2 \sum_{j=1}^{n} f(x_{j}))$$

$$= \frac{h}{2} (f(a) - f(b) + 2 \sum_{j=1}^{n} f(a + jh)) = T_{n}(h)$$

様形 等分

$$\chi_{j-1}$$
 $\chi_{j} R_{T}[f,h] = -\frac{h^{3}}{12} \sum_{j=1}^{n} f''(\xi_{j})$ を 変数 数 $\xi_{j} \in (x_{j-1},x_{j})$

若f(x)是[a,b]区间上的二阶连续函数则必存在一点 $\xi \in (a,b)$ 使得

$$f''(\xi) = \frac{1}{n} \sum_{j=1}^{n} f''(\xi_{j}) \qquad \qquad |f''(\xi)| \max = M$$

$$ign times the equation of the$$

2. 复化Simpson公式



在每个子区间[
$$x_{j-1}, x_j$$
] $j = 1, 2,n$ 上有

$$S_{j} = \frac{h}{6} (f(x_{j-1}) + 4f(x_{j-\frac{1}{2}}) + f(x_{j}))$$

则
$$\int_{a}^{b} f(x)dx = \sum_{j=1}^{n} S_{j} + R_{S}[f,h]$$

$$\overrightarrow{m} \sum_{j=1}^{n} S_{j} = \frac{h}{6} \sum_{j=1}^{n} (f(x_{j-1}) + 4f(x_{j-\frac{1}{2}}) + f(x_{j}))$$

$$= \frac{h}{6} (f_{0} + 4f_{\frac{1}{2}} + f_{1} + f_{1} + 4f_{\frac{3}{2}} + f_{2}$$

$$+ f_{n-1} + 4 f_{n-\frac{1}{2}} + f_n + f_n - f_n$$

$$= \frac{h}{6} (f(a) + f(b) + 4 \sum_{j=1}^{n} f(x_{j-\frac{1}{2}}) + 2 \sum_{j=1}^{n} f(x_{j}))$$

$$= \frac{h}{3} (\frac{f(a) - f(b)}{2} + 2 \sum_{j=1}^{n} f(x_{j-\frac{1}{2}}) + \sum_{j=1}^{n} f(x_{j}))$$

$$= \frac{h}{3} (\frac{f(a) + f(b)}{2} + 2 \sum_{j=1}^{n} f(x_{j-\frac{1}{2}}) + \sum_{j=1}^{n} f(x_{j}))$$

$$= \frac{h}{3} \left(\frac{f(a) - f(b)}{2} + 2 \sum_{j=1}^{n} f(a + (j - \frac{1}{2})h) + \sum_{j=1}^{n} f(a + jh) \right)$$

$$= \frac{h}{3} \left(\frac{f(a) - f(b)}{2} + \left(\sum_{j=1}^{n} f(a + jh) + 2 f(a + (j - \frac{1}{2})h) \right) \right)$$

$$= S_{n}(h)$$

$$R_{S}[f, h] = -\frac{h^{5}}{2880} \sum_{j=1}^{n} f^{(4)}(\xi_{i}) \quad \xi_{i} \in (x_{j-1}, x_{j})$$

$$= -\frac{(b - a)h^{4}}{2880} f^{(4)}(\xi) \quad \xi \in (a, b)$$

$$\begin{cases} \sum_{j=1}^{n} f(a + jh) + 2 f(a + (j - \frac{1}{2})h) \\ \sum_{j=1}^{n} f(a + jh) + 2 f(a + (j - \frac{1}{2})h) \\ \sum_{j=1}^{n} f(a + jh) + 2 f(a + (j - \frac{1}{2})h) \end{cases}$$

例题

$$I = \int_0^{\frac{\pi}{2}} \sin x dx$$

如果用复化梯形公式和用复化Simpson

公式计算,要使得截断误差不超过 $\frac{1}{2} \times 10^{-5}$

试问划分数n至少取多少?

解: 由截断误差有

$$R_{T}[f,h] = -\frac{(b-a)h^{2}}{12}f''(\xi) = -\frac{\frac{\pi}{2} - 0}{12}(\frac{\pi}{2})^{2}(\sin^{"}\xi)$$

$$\mathbb{E}[R_{T}[f,h]] = \left|\frac{\pi}{24}\frac{\pi}{2n}\sin\xi\right| \le \left|\frac{\pi^{2}}{48n}\right| \le \frac{1}{2} \times 10^{-5}$$

解得 n > 254

由

$$R_{S}[f,h] = -\frac{(b-a)h^{4}}{2880} f^{(4)}(\xi) = -\frac{\frac{\pi}{2} - 0}{2880} (h)^{4} (\sin^{(4)} \xi)$$

$$\mathbb{P}[R_{S}[f,h]] = \left| \frac{\pi}{2880 \times 2} h^{4} \sin \xi \right| \le \frac{1}{2} \times 10^{-5}$$

解得
$$h < 0.31$$
,故 $n > \frac{\frac{n}{2}}{0.31} \approx 5.1$,取 $n = 6$ 。

$$f(x) = x^{4} \qquad I = \int_{0}^{4} f(x) dx = \frac{45}{5} = 2.04.8$$

$$h = 4 \qquad I_{1} = \frac{1}{2} [f(a) + f(b)]$$

$$= 2 \times (0 + 4^{4}) = 512$$

$$h = 2 \qquad I_{2} = \frac{1}{2} (T_{1} + h_{1} + 0.2 + f(a) + h_{2} + f(a) + h_{3} + h_{4} + h_{4} + h_{5} + h_$$

TI TZ 14 T8 T16 T32

$$T_{2k} = T_{k} + W(T_{2k} - T_{k}).$$

变步长求积公式

 $\frac{1}{a} = \frac{b^{-\alpha}}{2} (f(\alpha) + f(b))$ $\frac{1}{1} = \frac{b^{-\alpha}}{4} (f(\alpha) + f(b)) + f(b)$ $\frac{1}{1} = \frac{b^{-\alpha}}{4} (f(\alpha) + f(b)) + f(b)$

定步长复化求积公式的一个明显缺点是:事先 很难估计分划数 n 使结果达到预期精度。由于适当 加密分点,精度会有所改善,为此采用自动加密分 点的方法,并利用事后估计来控制加密次数,以判 断是否达到预期精度,从而停止计算。

首先我们讨论变步长梯形求积公式。

变步长梯形求积公式

设区间[a,b]划分为 n 等分,即步长 $h = \frac{b-a}{n}$, 计算 $T_n(h)$: 然后将区间[a,b]分点加密一倍,即步长 缩小一半为 $\frac{h}{2}$,再计算出 $T_{2n}(\frac{h}{2})$ 。 如果 $\left|T_{2n}(\frac{h}{2}) - T_n(h)\right| \leq \varepsilon$ 则取 $S=T_{2n}(\frac{h}{2})$ 作为定积分的近似值。已知 $T_n(h)$,如何计算 $T_{2n}(\frac{h}{2})$ 且计算量小? $T_{2n}(\frac{h}{2}) = T_{2n}(\frac{h}{2}) = T_{2n}(\frac$ +2= f(x=1) $S_{j} = \frac{\frac{h}{2}}{2} (f_{j-1} + 2f_{j-\frac{1}{2}} + f_{j}) \qquad -\frac{1}{2} \int_{h}^{R^{2}} (h) + \frac{h}{2} \int_{k=1}^{R^{2}} f(k-\frac{1}{2})$

所以
$$T_{2n} = \sum_{j=1}^{n} S_j = \frac{h}{4} \{ \sum_{j=1}^{n} (f_{j-1} + f_j) + 2 \sum_{j=1}^{n} f_{j-\frac{1}{2}} \}$$

$$= \frac{1}{2} \left\{ \frac{h}{2} \sum_{j=1}^{n} (f_{j-1} + f_j) + h \sum_{j=1}^{n} f_{j-\frac{1}{2}} \right\}$$
$$= \frac{1}{2} (T_n(h) + H_n(h))$$

其中
$$H_n(h) = h \sum_{j=1}^n f(a + (j - \frac{1}{2})h)$$

变 步 长 Simpson 求 积 公 式

由复化Simpson公式

$$S_n(h) = \frac{h}{3} \left(\frac{f(a) - f(b)}{2} + \left(\sum_{j=1}^n f(a+jh) + 2f(a+(j-\frac{1}{2})h) \right) \right)$$

$$= \frac{1}{3} \left\{ h \left[\frac{f(a) - f(b)}{2} + \sum_{j=1}^n f(a+jh) \right] + 2h \sum_{j=1}^n f(a+(j-\frac{1}{2})h) \right\}$$

$$= \frac{1}{3} \left(T_n(h) + 2H_n(h) \right)$$
所以 $S_{2n}(\frac{h}{2}) = \frac{1}{3} \left(T_{2n}(\frac{h}{2}) + 2H_{2n}(\frac{h}{2}) \right)$

程序实现的基本思想:

定积分的近似值,否则 将分点加密一倍, 重复上述过程。 复化 Simpson 公式与复化梯形公式 有如下关系

$$S_n(h) = \frac{4T_{2n}(\frac{h}{2}) - T_n(h)}{4 - 1}$$

同理也可以推出复化 Cotes 公式

$$C_n(h) = \frac{4^2 S_{2n}(\frac{h}{2}) - S_n(h)}{4^2 - 1}$$

例题

计算
$$I = \int_0^1 \frac{dx}{1+x}$$

- (1)用Simpson公式;
- (2)用n = 5的复化Simpson计算,并估计误差;
- (3)用变步长Simpson公式计算,使其误差小于10⁻⁵。

(1)
$$I = \frac{1}{6} (f(0) + 4f(\frac{1}{2}) + f(1))$$
$$= \frac{1}{6} (1 + 4\frac{1}{3} + \frac{1}{2}) = \frac{25}{36} \approx 0.69444$$
$$(2) \cdots b = \frac{b - a}{3} - \frac{1}{3} = 0.2$$

(2) :
$$h = \frac{b-a}{n} = \frac{1}{5} = 0.2$$

则节点
$$x_i = 0 + ih$$
 $(i = 0,1,2,3,4,5)$

所以有:

$$I = S_5(h) = \frac{1}{6} \frac{1}{5} \left[\left(\frac{1}{1+0} - \frac{1}{1+1} \right) + 2 \times \left(\frac{1}{1+0.2} \right) \right]$$

$$+ \frac{1}{1+0.4} + \frac{1}{1+0.6} + \frac{1}{1+0.8} + \frac{1}{1+1} \right]$$

$$= \frac{1}{30} \left(\frac{1}{2} + 2 \times \left(\frac{1}{1.2} + \frac{1}{1.4} + \frac{1}{1.6} + \frac{1}{1.8} + \frac{1}{2} \right) \right)$$

$$\approx 0.69315$$

因为
$$\left| f^{(4)}(x) \right| = \frac{24}{\left| 1 + x \right|^5} \le 24 \quad x \in [0,1]$$
所以 $R_S[f,h] \le \frac{h^4}{2880} \times 24 = \frac{(0.2)^4}{120}$
 $= 1.3333 \times 10^{-5}$