

14 - C4

排序
希尔排序 : Pratt序列

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Pratt's Sequence, 1971

$$\mathcal{H}_{pratt} = \{ 2^p \cdot 3^q \mid p, q \in \mathcal{N} \}$$

$$= \{ 1, 2, 3, 4, 6, 8, 9, 12, 16, 18, 24, 27, 32, 36, \dots \}$$

❖ Note that

- adjacent items are **NOT** always relatively prime and
- there are $\mathcal{O}(\log^2 n)$ items no greater than n

❖ With \mathcal{H}_{pratt} ,

Shellsort sorts a sequence of length n in $\mathcal{O}(n \cdot \log^2 n)$ time ...

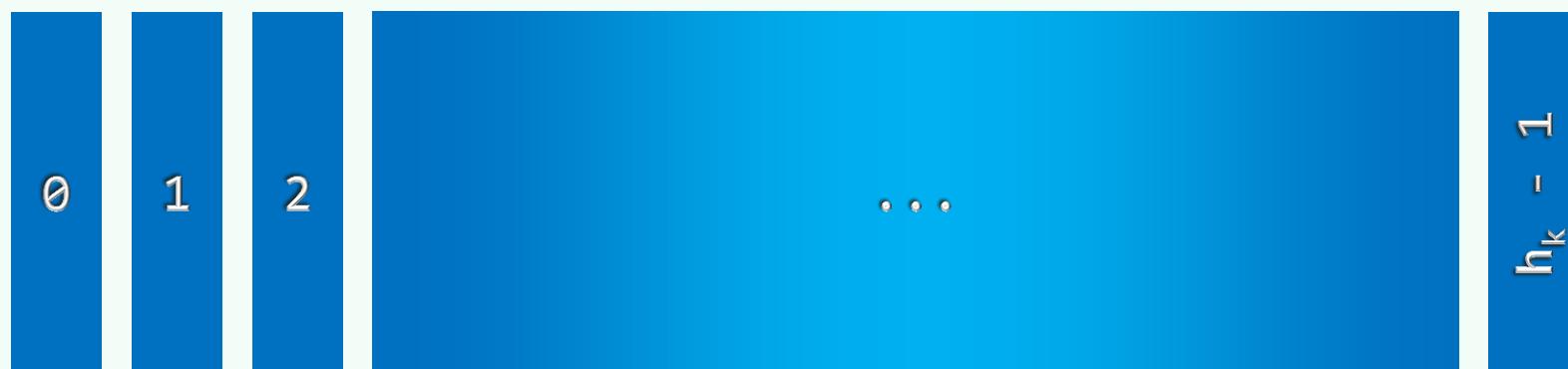
From $(2,3)$ -ordered to 1-ordered

$\because x(2,3) = 1$

\therefore To the right of each element in a $(2,3)$ -ordered sequence,

only the **NEXT** element can be smaller

\therefore It costs $\mathcal{O}(n)$ time to sort such a sequence



From $(2*h_k, 3*h_k)$ -ordered to h_k -ordered

❖ Divide S into h_k subsequences, each of which is $(2,3)$ -ordered

∴ it costs altogether $\mathcal{O}(n)$ time to sort them resp.

❖ ∵ there are altogether $\mathcal{O}(\log^2 n)$ iterations

∴ we need $\mathcal{O}(n \cdot \log^2 n)$ time

