

排序

希尔排序：Shell序列 + 输入敏感性

14-C2

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# Shell's Sequence

❖ Shell 1959 :  $\mathcal{H}_{shell} = \{ 1, 2, 4, 8, \dots, 2^k, \dots \}$

❖ 实际上, 采用  $\mathcal{H}_{shell}$ , 在最坏情况下需要运行  $\Omega(n^2)$  时间...

❖ 考查由子序列  $A = \text{unsort}[0, 2^{N-1})$  和  $B = \text{unsort}[2^{N-1}, 2^N)$  交错而成的序列



❖ 在做2-sorting时, A、B各成一行; 故此必然各自有序



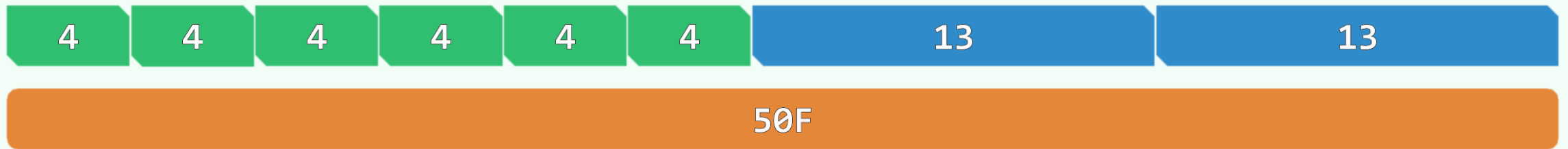
❖ 然而其中的逆序对依然很多, 最后的1-sorting仍需  $1 + 2 + 3 + \dots + 2^{N-1} = \Omega(n^2/4)$  时间

❖ 根源在于,  $\mathcal{H}_{shell}$  中各项并不互素, 甚至相邻项也非互素

# Postage Problem

❖ The postage for a letter is 50F, and a postcard 35F

But there are only stamps of 4F and 13F available



❖ Possible to stamp



the letter and



the postcard



EXACTLY?

❖ Given a postage  $P$ , determine whether  $P \in \{ n \cdot 4 + m \cdot 13 \mid n, m \in \mathcal{N} \}$

# Linear Combination

❖ Let  $g, h \in \mathcal{N}$

❖ For any  $n, m \in \mathcal{N}$ ,  $n \cdot g + m \cdot h$  is called a **linear combination** of  $g$  and  $h$

❖ Denote  $\mathbf{C}(g, h) = \{ ng + mh \mid n, m \in \mathcal{N} \}$

$\mathbf{N}(g, h) = \mathcal{N} \setminus \mathbf{C}(g, h)$  //numbers that are **NOT** combinations of  $g$  and  $h$

$\mathbf{x}(g, h) = \max\{ \mathbf{N}(g, h) \}$  //always exists?

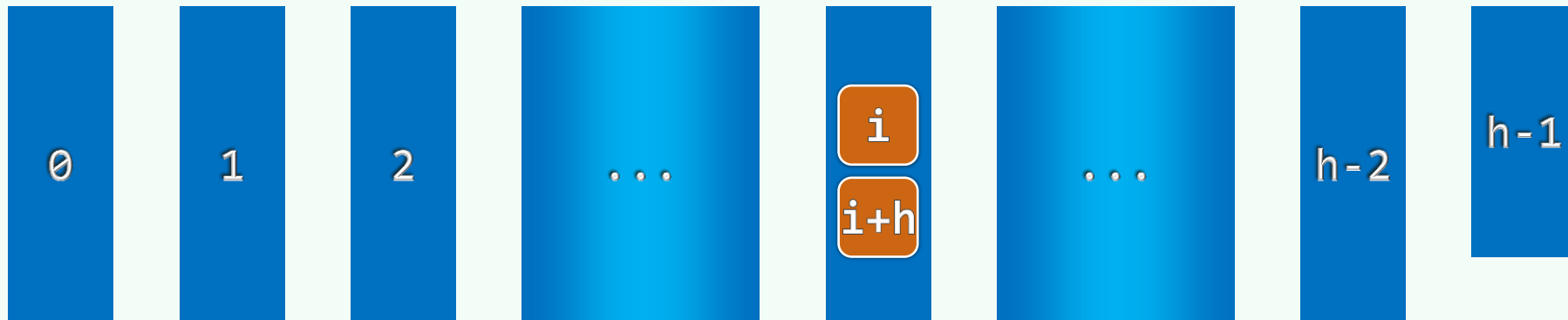
❖ Theorem: when  $g$  and  $h$  are **relatively prime**, we have

$$\mathbf{x}(g, h) = (g - 1) \cdot (h - 1) - 1 = gh - g - h$$

e.g.  $\mathbf{x}(3, 7) = 11$ ,  $\mathbf{x}(4, 9) = 23$ ,  $\mathbf{x}(\boxed{4}, \boxed{13}) = \boxed{35}$ ,  $\mathbf{x}(5, 14) = 51$

# h-sorting & h-ordered

- ❖ A sequence  $S[0,n)$  is called **h-ordered** if  $S[i] \leq S[i+h], \forall 0 \leq i < n-h$
- ❖ A 1-ordered sequence is sorted
- ❖ **h-sorting**: an h-ordered sequence is obtained by
  - arranging  $S$  into a 2D matrix with **h** columns and
  - sorting each column respectively

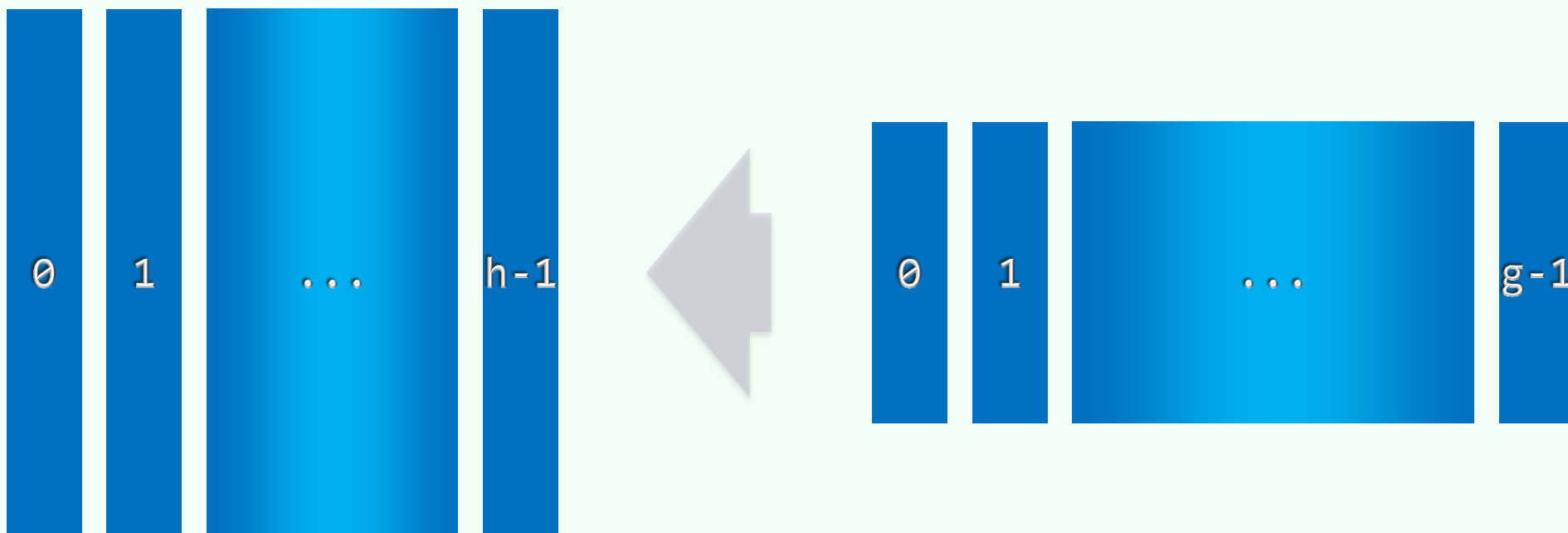


# Theorem K

❖ [Knuth, ACP Vol.3 p.90]

//习题解析[12-12, 12-13]

A **g**-ordered sequence REMAINS **g**-ordered after being **h**-sorted.

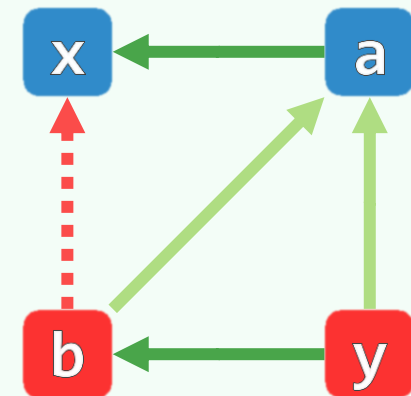
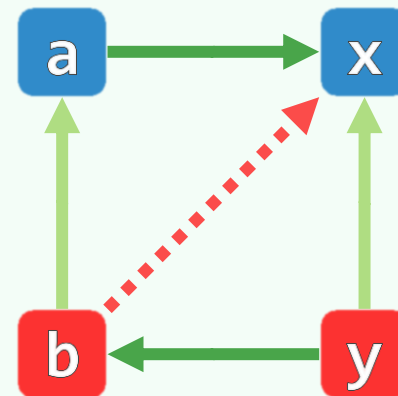
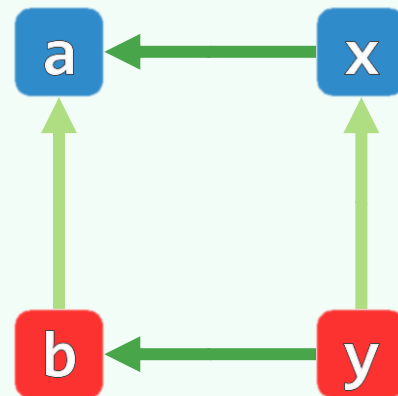
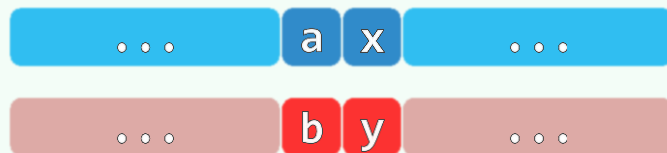


# Order Preservation

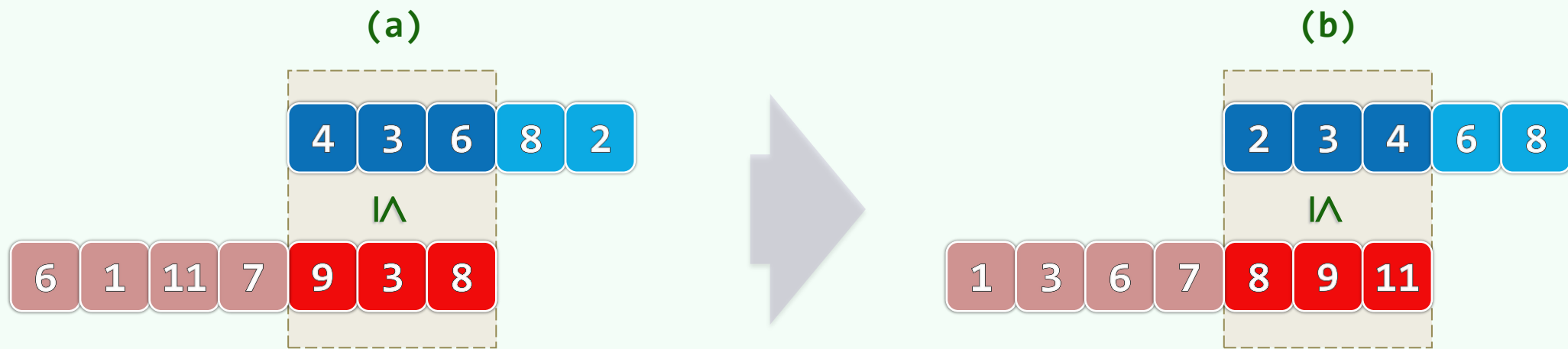
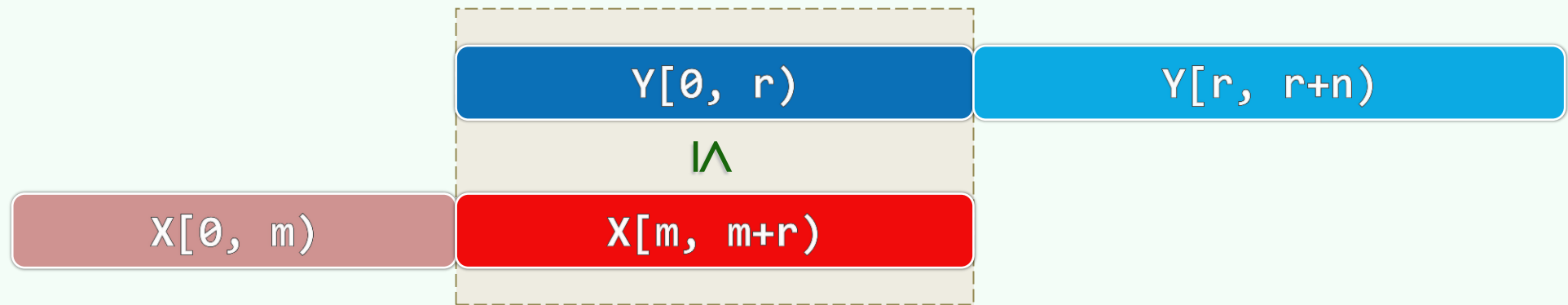
5	8	7	3	5
1	5	2	8	8
0	9	4	6	2
6	3	1	4	7

0	3	1	3	2
1	5	2	4	5
5	8	4	6	7
6	9	7	8	8

0	1	2	3	3
1	2	4	5	5
4	5	6	7	8
6	7	8	8	9



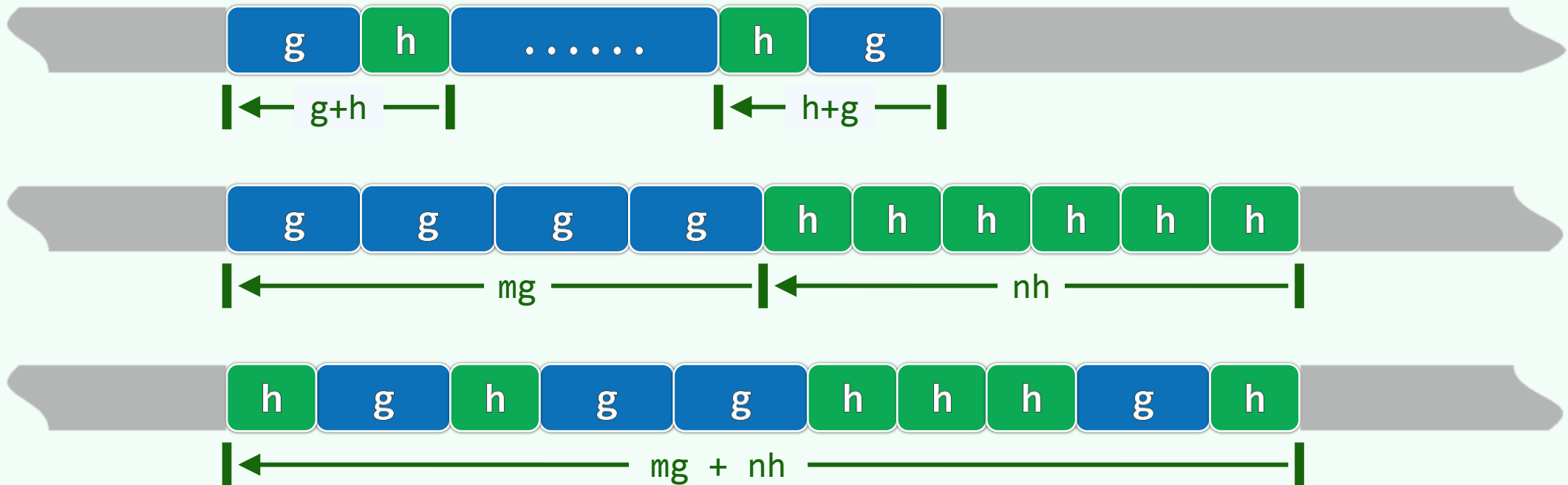
## Lemma L





# Linear Combination

- ❖ A sequence that is both **g**-ordered and **h**-ordered  
is called **(g,h)**-ordered, which must be both  
**(g+h)**-ordered and **(mg+nh)**-ordered for any  $m, n \in \mathbb{N}$



# Inversion

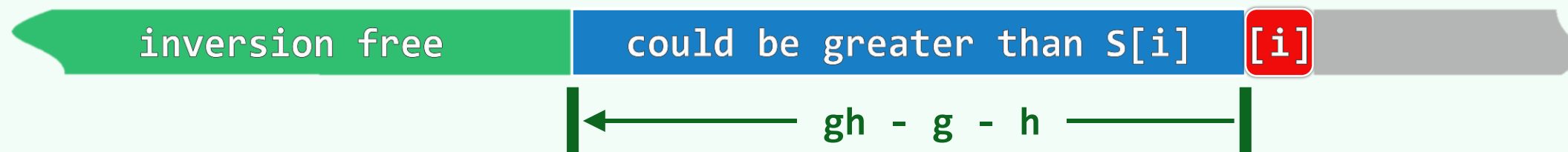
❖ Let  $S[0,n)$  be a  $(g,h)$ -ordered sequence, where  $g$  and  $h$  are **relatively prime**

❖ Then for all elements  $S[j]$  and  $S[i]$ , we have

$$i - j > x(g, h) \quad \text{only if} \quad S[j] \leq S[i]$$

❖ This implies that to the **LEFT** of each element,

only the previous  $x(g, h)$  elements could be **GREATER**



❖ There would be no more than  $n \cdot x(g, h)$  **INVERSIONS** altogether