

图

邻接矩阵：简单接口

θ6-B3

邓俊辉

deng@tsinghua.edu.cn

顶点的读写



- ❖ `Tv & vertex(int i) { return V[i].data; } //数据`
- `int inDegree(int i) { return V[i].inDegree; } //入度`
- `int outDegree(int i) { return V[i].outDegree; } //出度`
- `vStatus & status(int i) { return V[i].status; } //状态`
- `int & dTime(int i) { return V[i].dTime; } //时间标签dTime`
- `int & fTime(int i) { return V[i].fTime; } //时间标签fTime`
- `int & parent(int i) { return V[i].parent; } //在遍历树中的父亲`
- `int & priority(int i) { return V[i].priority; } //优先级数`

边的读写

V

E

i

0 1 2

.....

j

.....

n-1 n

❖ bool exists(int i, int j) { //判断边(i, j)是否存在 (短路求值)
 return (0 <= i) && (i < n) && (0 <= j) && (j < n)
 && E[i][j] != NULL;

} //以下假定exists(i, j) = true

❖ TEdge & edge(int i, int j) //边的数据 , O(1)

{ return E[i][j]->data; }

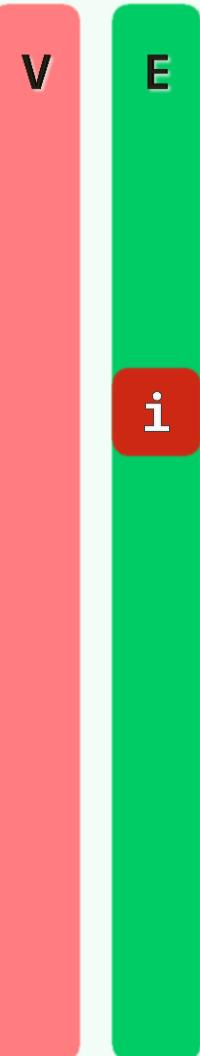
❖ EType & type(int i, int j) //边的类型 , O(1)

{ return E[i][j]->type; }

❖ int & weight(int i, int j) //边的权重 , O(1)

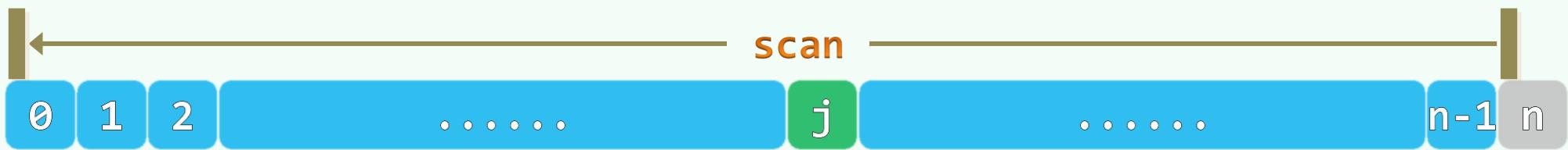
{ return E[i][j]->weight; }

邻点的枚举



❖ 对于任意顶点*i*，如何枚举其所有的邻接顶点（neighbor）？

❖ int firstNbr(int i) { return nextNbr(i, n); } //假想哨兵



❖ int nextNbr(int i, int j) { //若已枚举至邻居j，则转向下一邻居

```
while ( ( -1 < j ) && ! exists( i, --j ) ); //逆向顺序查找  
return j;
```

} // $O(n)$ ——改用邻接表，可提高至 $O(1 + \text{outDegree}(i))$