

图

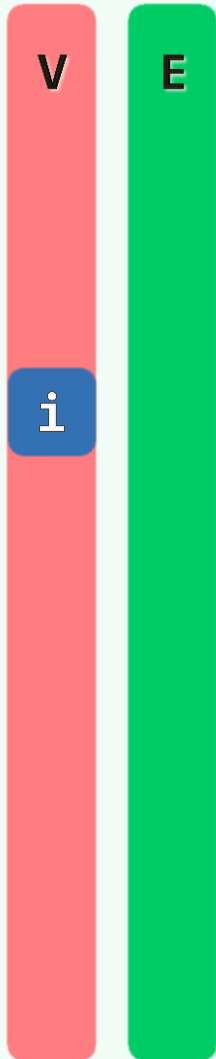
邻接矩阵：简单接口

06-B3

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# 顶点的读写



❖ Tv & vertex(int i) { return V[i].data; } //数据

int inDegree(int i) { return V[i].inDegree; } //入度

int outDegree(int i) { return V[i].outDegree; } //出度

VStatus & status(int i) { return V[i].status; } //状态

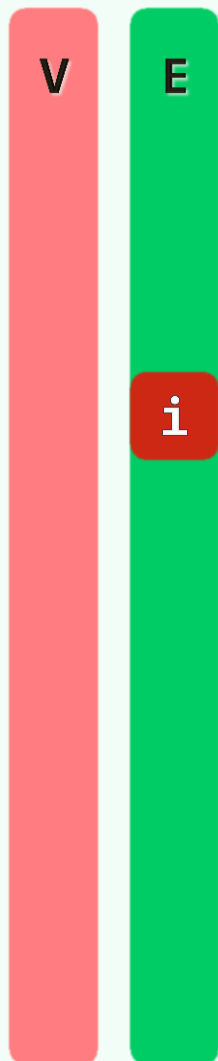
int & dTime(int i) { return V[i].dTime; } //时间标签dTime

int & fTime(int i) { return V[i].fTime; } //时间标签fTime

int & parent(int i) { return V[i].parent; } //在遍历树中的父亲

int & priority(int i) { return V[i].priority; } //优先级数

# 边的读写



```
❖ bool exists( int i, int j ) { //判断边(i, j)是否存在 ( 短路求值 )  
    return ( 0 <= i ) && ( i < n ) && ( 0 <= j ) && ( j < n )  
        && E[i][j] != NULL;  
} //以下假定exists(i, j) = true
```



```
❖ Te & edge( int i, int j ) //边的数据 , 0(1)  
    { return E[i][j]->data; }  
  
❖ EType & type( int i, int j ) //边的类型 , 0(1)  
    { return E[i][j]->type; }  
  
❖ int & weight( int i, int j ) //边的权重 , 0(1)  
    { return E[i][j]->weight; }
```

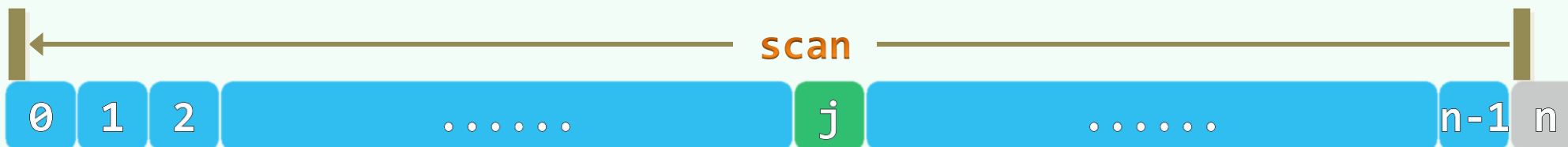
# 邻点的枚举

V

E

❖ 对于任意顶点  $i$  , 如何枚举其所有的邻接顶点 ( neighbor ) ?

❖ `int firstNbr( int i ) { return nextNbr( i, n ); } //假想哨兵`



❖ `int nextNbr( int i, int j ) { //若已枚举至邻居  $j$  , 则转向下一邻居`

`while ( (  $-1 < j$  ) && ! exists( i,  $--j$  ) ); //逆向顺序查找`

`return j;`

`} //O(n)——改用邻接表 , 可提高至  $O(1 + \text{outDegree}(i))$`