

优先级队列

左式堆：沿藤合并

12-XB1

邓俊辉

deng@tsinghua.edu.cn

God's right hand is gentle
But terrible is his left hand

堆合并

❖ $H = \text{merge}(A, B)$: 将堆A和B合二为一 //不妨设 $|A| = n \geq m = |B|$

❖ 方法一 : $A.\text{insert}(B.\text{delMax}())$

$$O(m * (\log m + \log(n + m)))$$

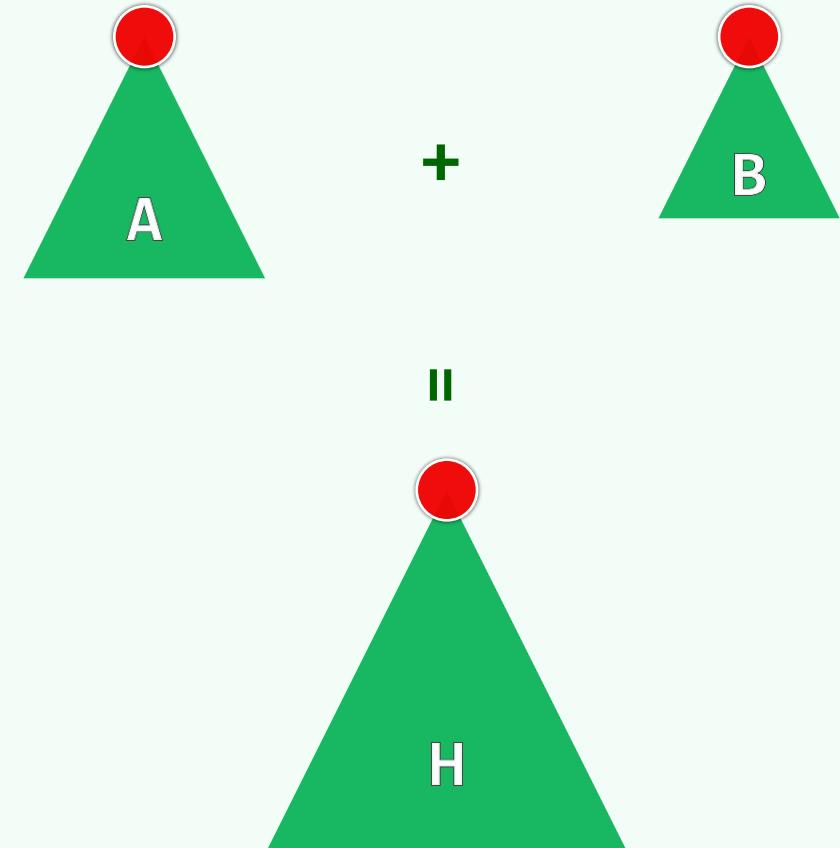
$$= O(m * \log(n + m))$$

❖ 方法二 : $\text{union}(A, B).\text{heapify}(n + m)$

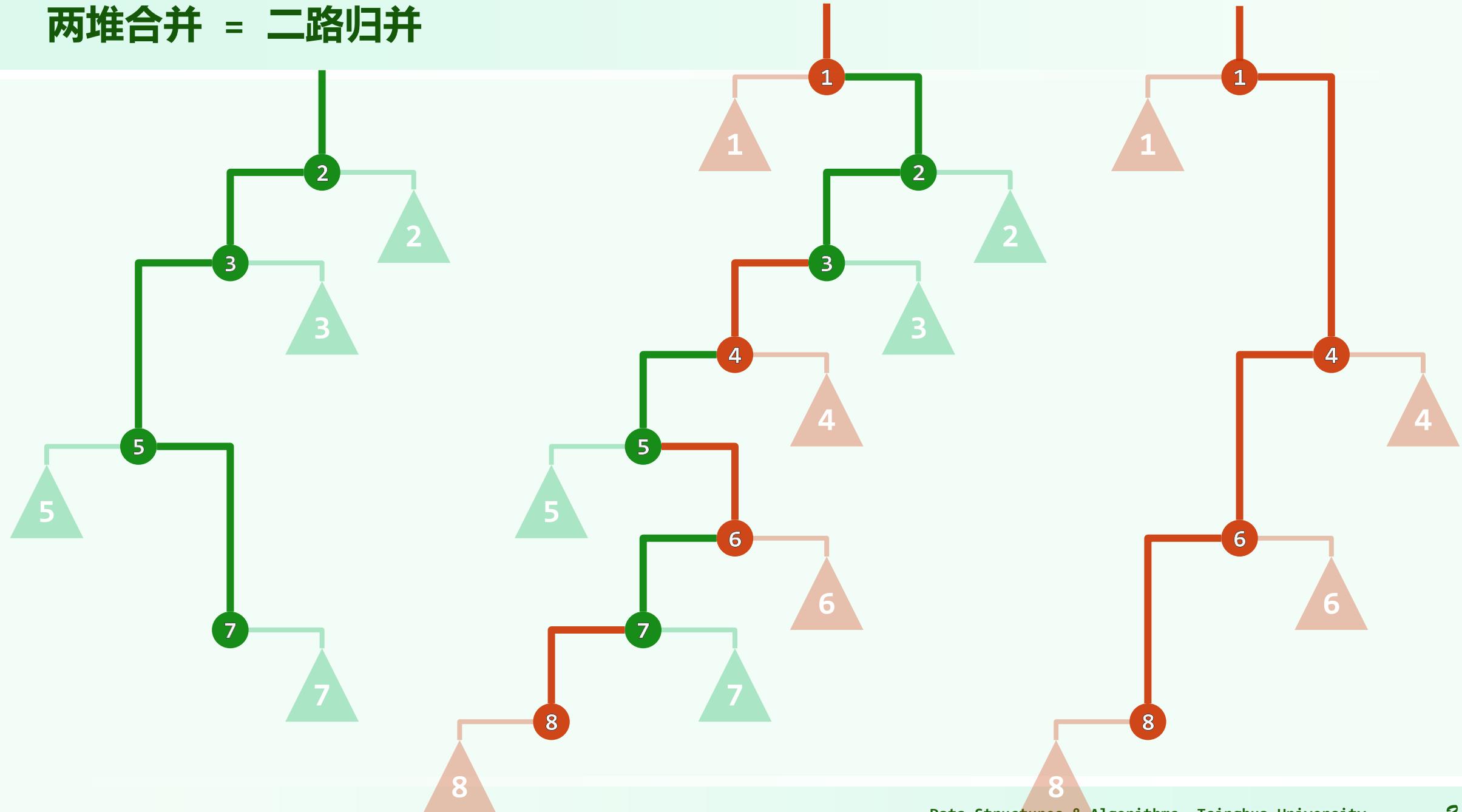
$$O(m + n)$$

❖ 有没有更好的办法 ? 比如...

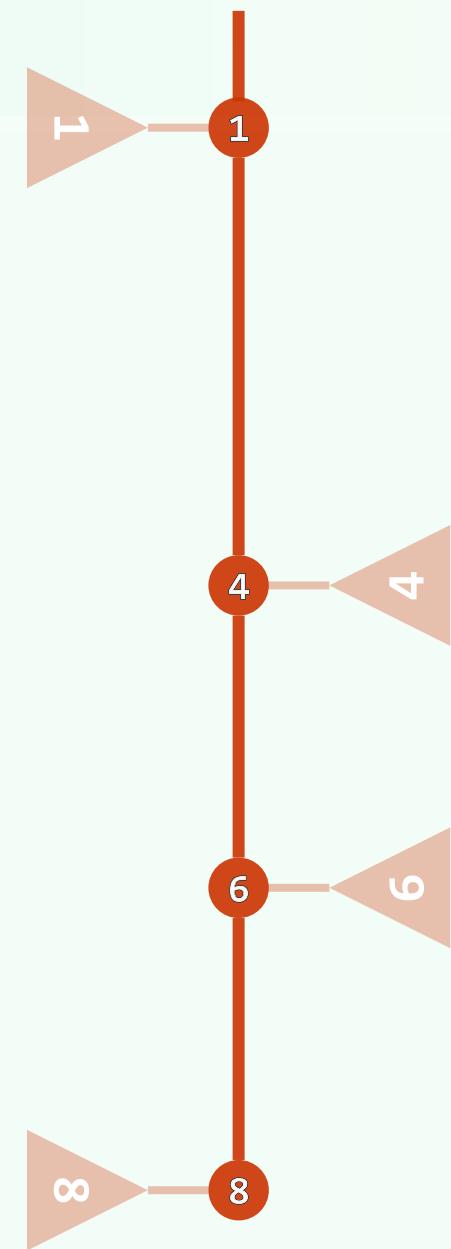
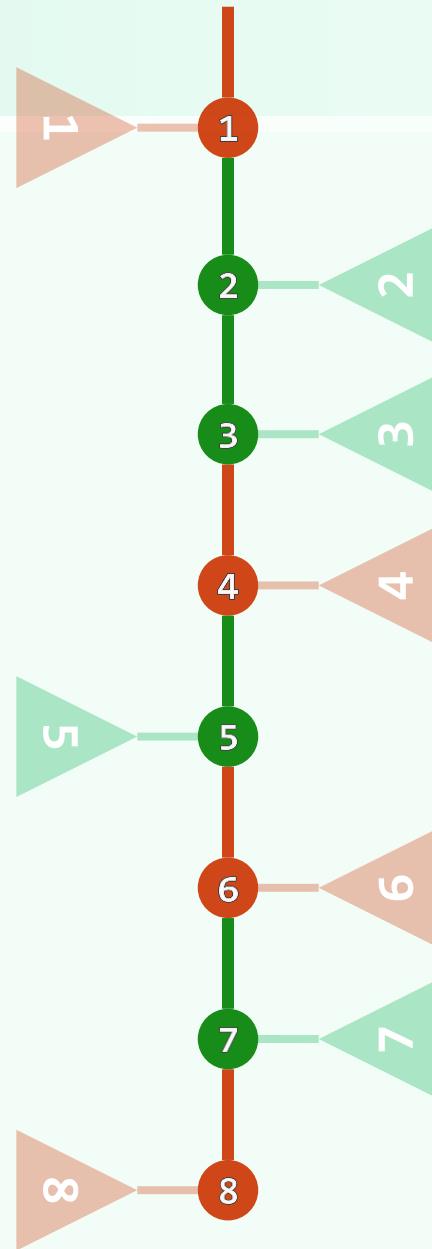
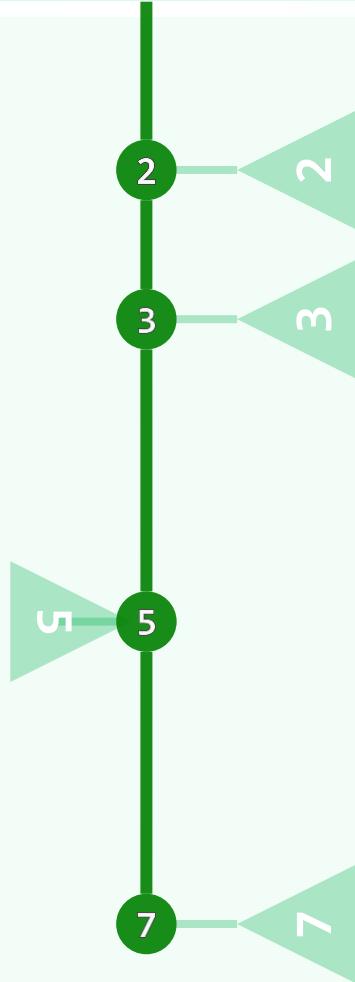
❖ 可否奢望在... $O(\log n)$...时间内实现merge() ?



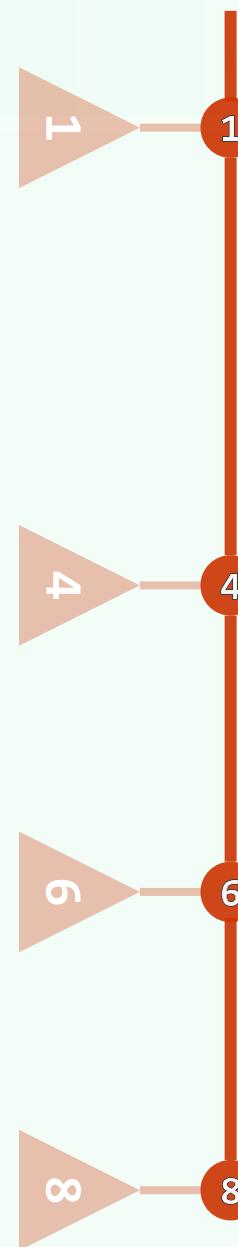
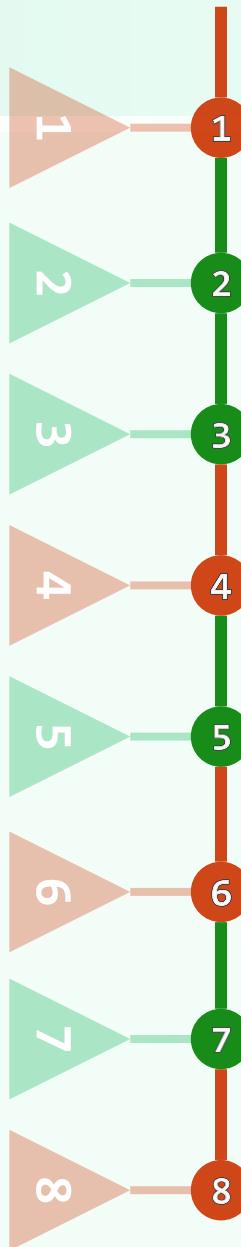
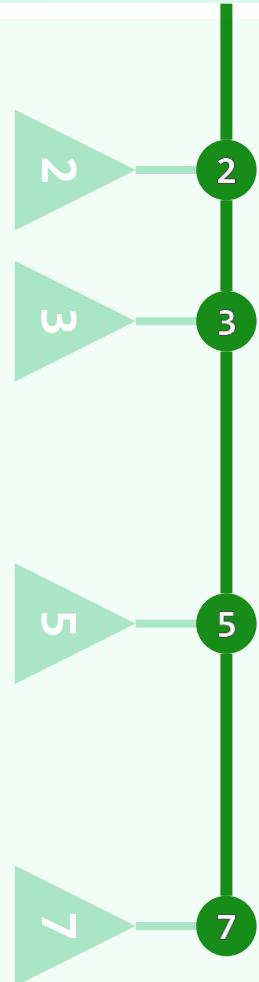
两堆合并 = 二路归并



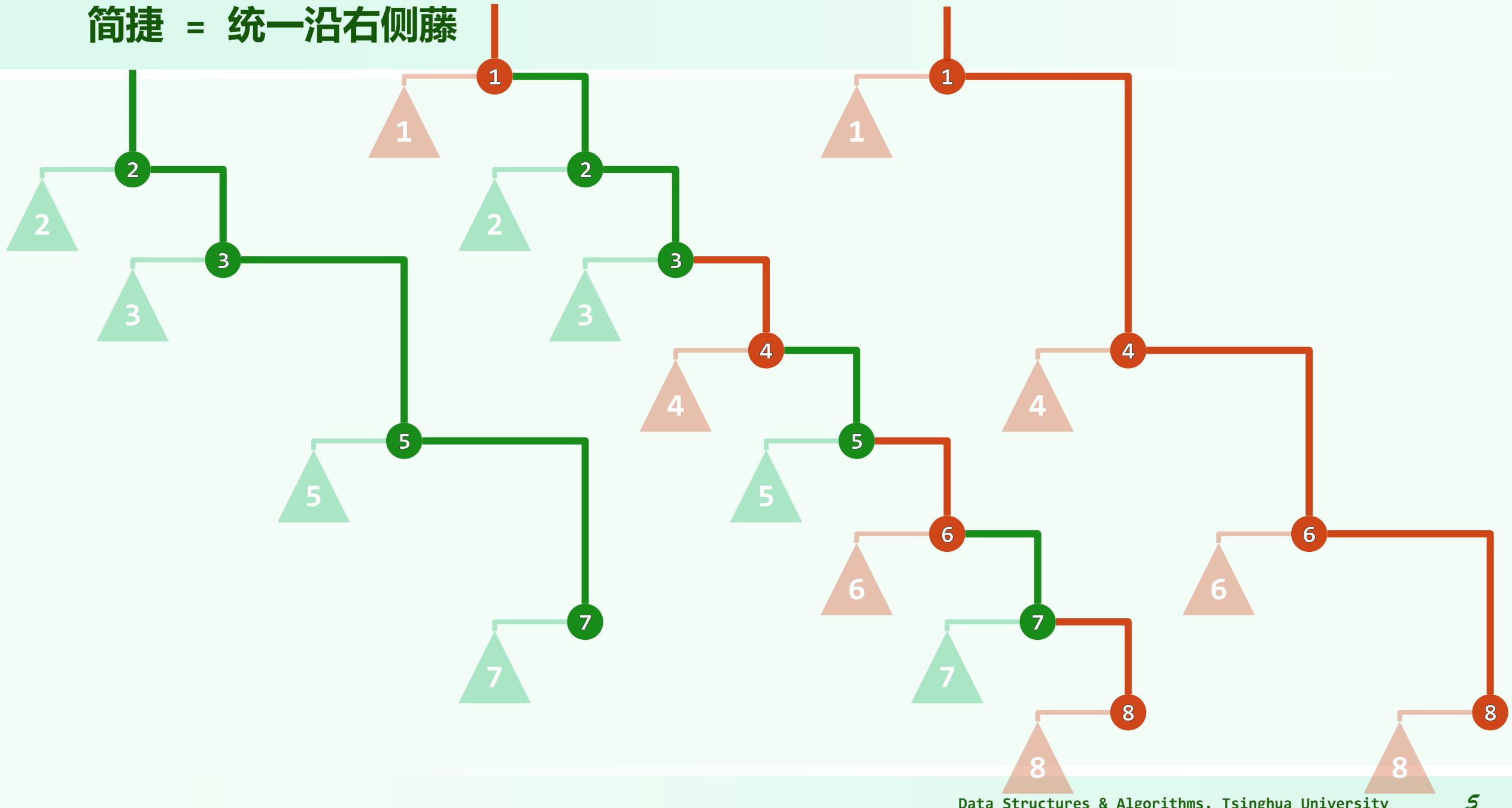
两堆合并 = 二路归并



简捷 = 统一右侧藤



简捷 = 统一沿右侧藤

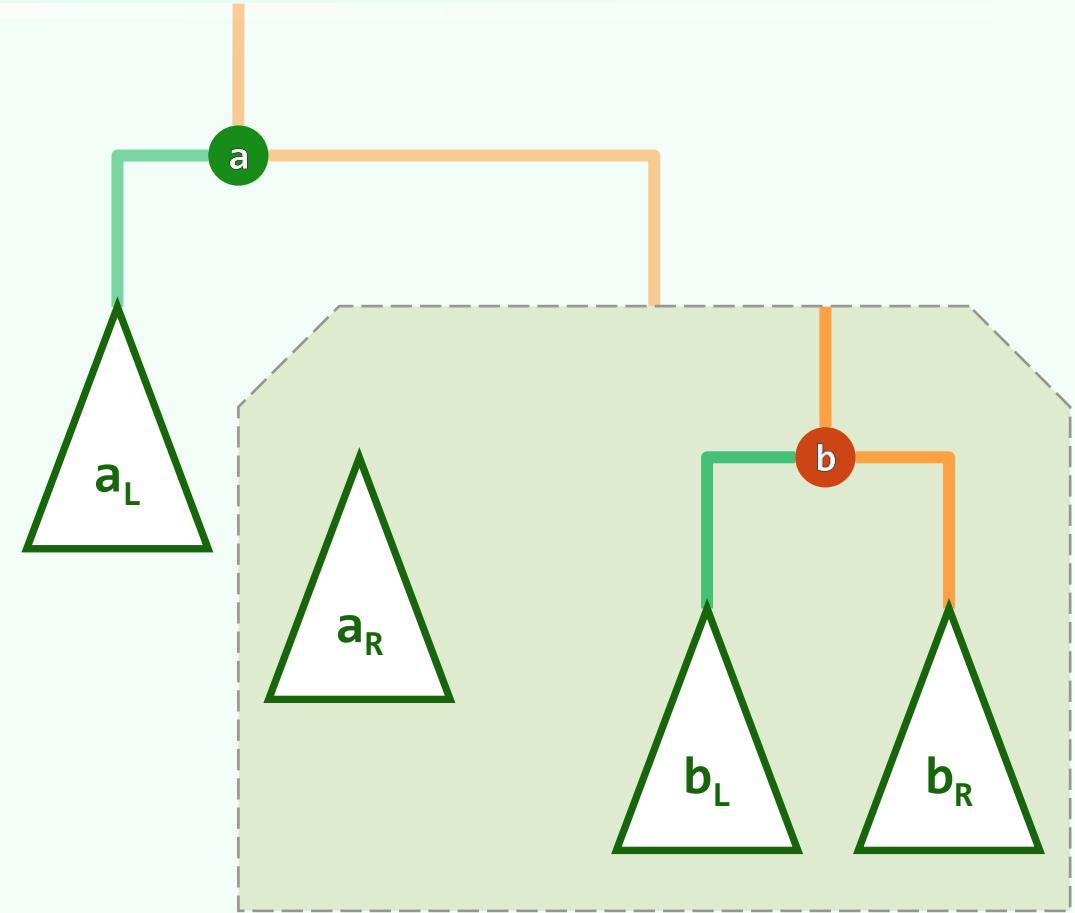
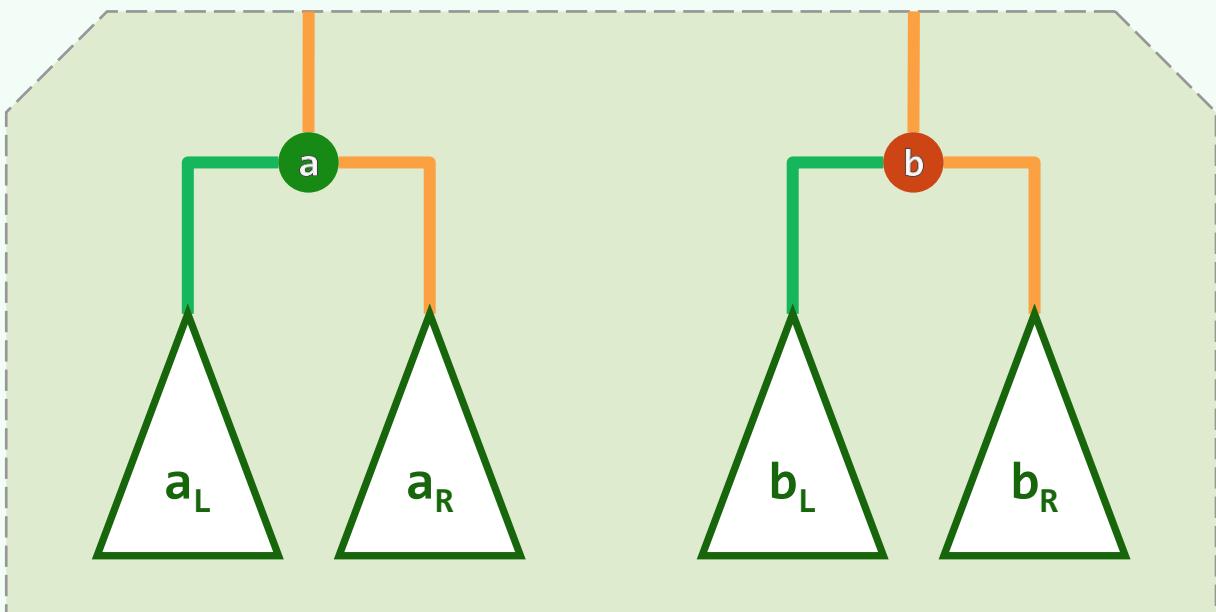


递归实现

所需时间 \propto 右侧藤总长

$\mathcal{O}(n)$

$\mathcal{O}(\log n)$



如何才能...控制藤长以...持续合并？