现代通信系统与网络仿真技术 大作业

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Problem1

题干略

After sampling, we have the energy spectral density

$$G_{MSKsampling}(f) = f_s^2 \sum_{n=-\infty}^{\infty} G_{MSK}(f - nf_s).$$

Hence we get the energy of signal and aliasing noise

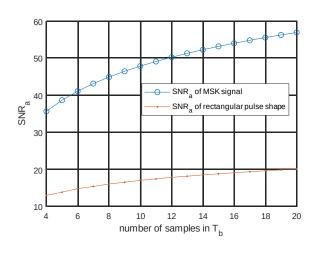
$$\begin{split} E_{signal} &= \int_{-\frac{f_s}{2}}^{\frac{f_s}{2}} f_s^2 G_{MSK}(f) df \,, \\ E_{noise} &= \int_{-\frac{f_s}{2}}^{\frac{f_s}{2}} f_s^2 \sum_{\substack{n=-\infty \\ n\neq 0}}^{\infty} G_{MSK}(f-nf_s) \, df \,, \end{split}$$

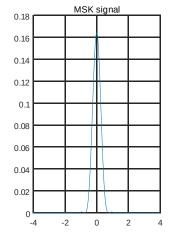
separately. By observing the representation of E_{nosie} and the propriety of $G_{MSK}(f)$, we derive

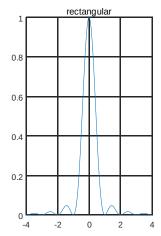
$$E_{noise} = \int_{\underline{f_s}}^{\infty} 2f_s^2 G_{MSK}(f) df ,$$

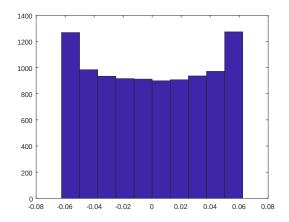
Then the resultant $(SNR)_a = \frac{E_{signal}}{E_{noise}}$. However, the integral is hard to solve, so we derive an approximation of $(SNR)_a$ by numerical integration:

$$(SNR)_a \approx \frac{\sum_{j=0}^{\frac{kf_s}{2}} G_{MSK}\left(\frac{j}{k}\right)}{\sum_{j=\frac{kf_s}{2}}^{\infty} G_{MSK}\left(\frac{j}{k}\right)}.$$









E: ---- Do -

The comparison of SNR is shown in Figure P1.a, which shows the signal to aliasing noise ratio of MSK signal is much higher than that of rectangular shape pulse signal, under same sampling frequency.

By plotting the energy spectral density in Figure P1.b, we can easily explain the gap: Energy spectral density of the rectangular shape pulse has a relative much higher sidelobe compared with that of MSK signal, whose sidelobe is very close to zero.

Problem 2

题干略

- (1) To avoid the vector being periodic, which will reduce the valid data.
- (3) Quantizing level $\Delta = \frac{2}{2^4} = 0.125$. Theoretical values:

$$E\{e[k]\} = 0.$$

$$E\{e^{2}[k]\} = \frac{\Delta^{2}}{12} = 0.0013.$$

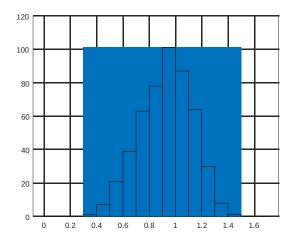
Experimental values:

$$\widehat{E}\{e[k]\} = -5.5155 \times 10^{-5} \approx 0$$

 $\widehat{E}\{e^2[k]\} = 0.0015.$

The theoretical values and experimental values are close.

(4) Histogram is shown as Figure P2.a. We conclude that it's a proper model to view quantizing error variant $e[k] = x_q[k] - x[k]$ as a uniformly distributed variant. But in this case of sin function, the probability of large error is greater than that of small error.



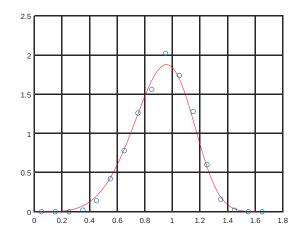


Figure P3.a

Problem 3

题干略

Solution:

$$f_X(x) = ax^{a-1}e^{-x^a}u(x)$$

$$F_X(x) = \int f_X(x)dx$$

$$= \int_0^\infty ax^{a-1}e^{-x^a}dx$$

$$= 1 - e^{-x^a}(x > 0)$$

Figure P3.b

Let $U = F_X(X)$, then $X = F_X^{-1}(U)$

$$X = \sqrt[a]{-\ln\left(1 - U\right)}$$

where U is uniformly distributed in (0,1).

Histogram is shown in Figure P3.a and the PDF is shown in P3.b. To compare them, we use points in Figure P3.b to represent the relative value of histogram bins. Generating only 500 samples makes the gap between PDF and experimental values significant, but it still shows the correspondence between them, indicating that the algorithm is correct.

Problem 4

题目内容略

总服务时间见图P4.a,平均吞吐量见P4.b,平均接入时延见P4.c

可以看出在总服务时间随MTCD数量增加浮动较大,但总体趋势为MTCD数量越多,总服务时间越长。 在设备数量并不很多时,随着设备数量增多,平均吞吐量和和平均接入时延都较快增加,但到当设备数量

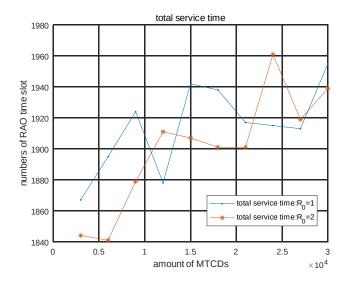


Figure P4.a

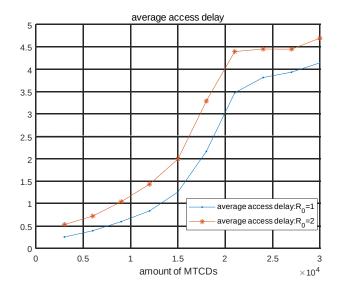


Figure P4 c

源码:

Problem1.m

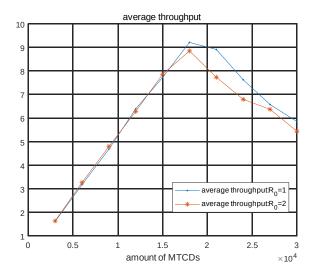


Figure P4.b

进一步增加时,平均吞吐量反而下降,而平均接入时延增速放缓。

对比不同最小传输速率约束 R_0 下的系统性能,发现其对总服务时间的影响规律性不强;而对于平均吞吐量,在设备数量较小时, R_0 对平均吞吐量基本没有影响,而在设备数量较大时, R_0 = 2 相较 R_0 = 1 时的吞吐量更小; R_0 对于平均接入时延的影响的规律性最强, R_0 = 2 时的接入时延显著高于 R_0 = 1 时的平均接入时延。

```
E_signal_MSK = sum(G_MSK(approx_points,T_b));
    E_signal_rec = sum(G_rec(approx_points,T_b));
    approx_points = f_s(idx)/2 : 1/approx_sample_n : infty_n;
    E_noise_MSK = sum(G_MSK(approx_points,T_b));
    E_noise_rec = sum(G_rec(approx_points,T_b));
    SNR_db_MSK(idx) = 10 * log10(E_signal_MSK / E_noise_MSK);
    SNR_db_rec(idx) = 10 * log10(E_signal_rec / E_noise_rec);
end
plot(m,SNR_db_MSK,'o-','DisplayName','SNR_a of MSK signal');
hold on
plot(m,SNR_db_rec,'.-','DisplayName','SNR_a of rectangular pulse shape');
xlabel('number of samples in T_b');
ylabel('SNR_a');
legend
grid on
figure(2)
f = -4:0.02:4;
y = G_MSK(f,1);
subplot(1,2,1)
plot(f,y,'DisplayName','MSK signal');
title('MSK signal');
grid on
subplot(1,2,2);
y = G_{rec}(f,1);
plot(f,y);
title('rectangular');
grid on
function y = G MSK(f,T b)
   y = 16 \cdot T_b \cdot (\cos(2 \cdot pi \cdot T_b \cdot f)) \cdot 2 \cdot / (pi^2 \cdot (1-(4 \cdot T_b \cdot f) \cdot 2)) \cdot 2;
function y = G_{rec}(f,T_b)
    y = (sinc(f.*T_b)).^2;
Problem2.m
clc,clear,clf
n_sample = 1e4;
t = 0 : 1/20 : n_sample / 20 - 1e-8;
signal = sin(6 .* t);
bias = 1/16;
q = quantizer('fixed','Round','saturate',[4,3]);
signal_q = quantize(q,signal - bias) + bias;
error = signal_q - signal;
% [signal_q_custom, error_custom] = quantize_costum(signal,1,3,bias);
mean_error = mean(error)
mean_error_square = mean(error .^ 2)
hist(error, 30)
% 下面这个函数只是检验一下自己对quantizer函数的理解。
% function [y,error] = quantize_costum(x,intLength,fracLength,bias)
      error = zeros(size(x));
%
      y = zeros(size(x));
```

```
%
       step = 2^{-fracLength};
%
       levels = -2^{intLength - 1}: step : 2^{intLength - 1} - 1e-8;
%
       levels = levels + bias;
%
       levels = levels(:);
%
       % levels is a column vector
       % suppose x is a row vector
%
%
       x_rep = repmat(x_length(levels), 1);
%
       err = x_rep - levels;
%
       [\sim, \min_i dx] = \min(abs(err));
%
       for idx = 1:length(x)
%
           error(idx) = err(min idx(idx),idx);
%
           y(idx) = levels(min_idx(idx));
%
       end
% end
```

Problem3.m

```
clc,clear,clf
num = 500;
a = 5;
figure(1)
interval_len = 0.1;
intervals = [0:interval_len:1.7];
U = rand([1,num]);
X = F_{weibull_inverse}(U,a);
h = histogram(X,intervals);
grid on
figure(2)
x = 0:0.02:1.6;
f = f_weibull(x,a);
plot(x,f,'r')
hold on
grid on
scatter(intervals(1:end-1) + interval_len /2 , h.Values ./ num ./ interval_len);
function f = f_weibull(x,a)
    f = a .* x.^{(a-1)} .* exp(-x.^a);
end
function X = F_weibull_inverse(U,a)
   X = (-(log(1-U))).^(1/a);
```

Problem4.m

clear,clc,clf

```
N_D_array = 3e3:3e3:3e4;

N_p = 54;

T = 10; % 10s

T_RAO = 5e-3; % 5ms, totally 2000 RAO time slot beta_a = 3; % paramaters of Beta distribution beta_b = 4;

J = 9;

target_SNR = 10;

R_O = 1;

filename_prefix = 'P4_R_O_1_';

total_service_time = zeros(size(N_D_array));
```

```
average throughput = zeros(size(N D array));
average access delay = zeros(size(N D array));
for N D idx = 1:length(N D array)
   \overline{N}D = ND_array(ND_idx);
   % get the first activation time
   activate time = betarnd(beta a,beta b,1,N D) * T;
   activate_time = floor(activate_time ./T_RAO);
   first activate time = activate time;
   time slot idx = min(activate time);
   is_done = 0;
   J[ocal = ones(1,N_D) * J; % when <math>J[ocal(idx) == -1, MTCD of idx won't access]
   sending_preamble_time = zeros(1,N_D);
   accessed_time = zeros(1,N_D);
   while is_done ~= 1
       % activate MTCD according to activate time
       activate_request_MTCD_idxs = find(activate_time == time_slot_idx);
       N_a = length(activate_request_MTCD_idxs);
       % ACB test
       ACB_factor = min(1, N_p / N_a);
       rand_{theta} = rand(1,N_a);
       bad_request_mask = rand_theta > ACB_factor;
       % bad request should request later
       delays = randi([1,4],size(activate_request_MTCD_idxs));
       delays = delays .* bad_request_mask;
       activate time(activate request MTCD idxs) = activate time(activate request MTCD idxs) + delays;
       % delete bad request ones and remains are activated successfully.
       activated MTCD idxs = find(activate time == time slot idx);
       sending preamble time(activated MTCD idxs) = time slot idx;
       sending preamble MTCD idxs = find(sending preamble time == time slot idx);
       % J validity check
       J_local(sending_preamble_MTCD_idxs) = J_local(sending_preamble_MTCD_idxs) - 1;
       sending preamble MTCD idxs = J validity check(J local, sending preamble MTCD idxs);
       % activated MTCDs choose preamble code
       [succeed_to_send_MTCD_idxs, fail_to_send_MTCD_idxs] =
preamble_code_uniqueness_check(sending_preamble_MTCD_idxs,N_p);
       % those who fail to send unique preamble code should send codes later.
       delays = randi([1,4],size(fail_to_send_MTCD_idxs));
       sending_preamble_time(fail_to_send_MTCD_idxs) = time_slot_idx + delays;
       % those succeed to send unique preamble code will face the SNR check
       SNR_{threshold} = 2^R_0 - 1;
       [succeed_to_access_MTCD_idxs, fail_to_access_MTCD_idxs] =
SNR_threshold_check(succeed_to_send_MTCD_idxs,target_SNR,SNR_threshold);
       delays = randi([1,4],size(fail_to_access_MTCD_idxs));
       sending_preamble_time(fail_to_access_MTCD_idxs) = time_slot_idx + delays;
       accessed time(succeed to access MTCD idxs) = time slot idx;
       accessed_number = sum(accessed_time \sim = 0);
       if (accessed_number + sum(J_local < 0)) == N_D
           is done = 1;
           % evalute
           total_service_time(N_D_idx) = time_slot_idx;
           average_throughput(N_D_idx) = accessed_number / time_slot_idx;
           access delay = accessed time - first activate time;
           access_delay(access_delay < 0) = []; % delete those failed
           average_access_delay(N_D_idx) = mean(access_delay);
       end
```

```
time_slot_idx = time_slot_idx + 1;
   end
   N D idx
end
figure(1)
plot(N D array,total service time, 'DisplayName', 'total service time: R 0=1');
beautify_figure();
figure(2)
plot(N D array, average throughput, 'DisplayName', 'average throughput: R 0=1');
beautify figure();
figure(3)
plot(N_D_array,average_access_delay,'DisplayName','average access_delay:R_0=1');
beautify_figure();
csvwrite([filename_prefix,'TST.csv'], total_service_time);
csvwrite([filename_prefix,'AT.csv'], average_throughput);
csvwrite([filename_prefix,'AAD.csv'], average_access_delay);
function beautify figure()
    grid on
   legend
end
function [success MTCD idxs,failure MTCD idxs] = SNR threshold check(total MTCD idxs, target SNR,
SNR_threshold)
   mu = 2;
   channel gains = exprnd(mu,size(total MTCD idxs));
   channel_gains = channel_gains * target_SNR;
   success MTCD idxs = total MTCD idxs(channel gains >= SNR threshold);
   failure MTCD idxs = total MTCD idxs(channel gains < SNR threshold);
function [success MTCD idxs, failure MTCD idxs] =
preamble_code_uniqueness_check(total_MTCD_idxs,N_p)
    codes_chosen = randi([1,N_p],size(total_MTCD_idxs));
   [codes_unique,success_MTCD_idxs,~] = unique(codes_chosen);
   codes_chosen = codes_chosen';
   failure mask = sum(codes unique == codes chosen) - 1; % it's a row vector
   success_MTCD_idxs = success_MTCD_idxs' .* (~failure_mask);
   success_MTCD_idxs(success_MTCD_idxs == 0) = [];
   success_MTCD_idxs = total_MTCD_idxs(success_MTCD_idxs);
   failure_MTCD_idxs = setdiff(total_MTCD_idxs,success_MTCD_idxs);
end
function idx = J_validity_check(J_local,pre_idx)
   J_exceed_mask = J_local(pre_idx) < 0;</pre>
   idx = pre_idx * (~J_exceed_mask);
   idx(idx == 0) = [];
end
Problem4_plot.m
clear,clc
N D array = 3e3:3e3:3e4;
filename prefix = 'P4 R 0 1 ';
tst_1 = csvread([filename_prefix,'TST.csv']);
at_1 = csvread([filename_prefix,'AT.csv']);
aad_1 = csvread([filename_prefix,'AAD.csv']);
filename prefix = 'P4 R 0 2 ';
tst_2 = csvread([filename_prefix, 'TST.csv']);
at 2 = csvread([filename prefix,'AT.csv']);
```

```
aad_2 = csvread([filename_prefix,'AAD.csv']);
figure(1);clf;
plot(N_D_array,tst_1,'.-','DisplayName','total service time:R_0=1');
hold on
plot(N_D_array,tst_2,'*-','DisplayName','total service time:R_0=2');
title('total service time');
ylabel('numbers of RAO time slot');
beautify_figure();
figure(2);clf;
plot(N_D_array,at_1,'.-','DisplayName','average throughput:R_0=1');
hold on
plot(N_D_array,at_2,'*-','DisplayName','average throughput:R_0=2');
title('average throughput');
beautify_figure();
figure(3);clf;
plot(N_D_array,aad_1,'.-','DisplayName','average access delay:R_0=1');
hold on
plot(N_D_array,aad_2,'*-','DisplayName','average access delay:R_0=2');
title('average access delay');
beautify_figure();
function beautify_figure()
   xlabel('amount of MTCDs');
   grid on
    legend
end
```

Problem4 流程图:

