



EDA 软件设计 - 调度算法

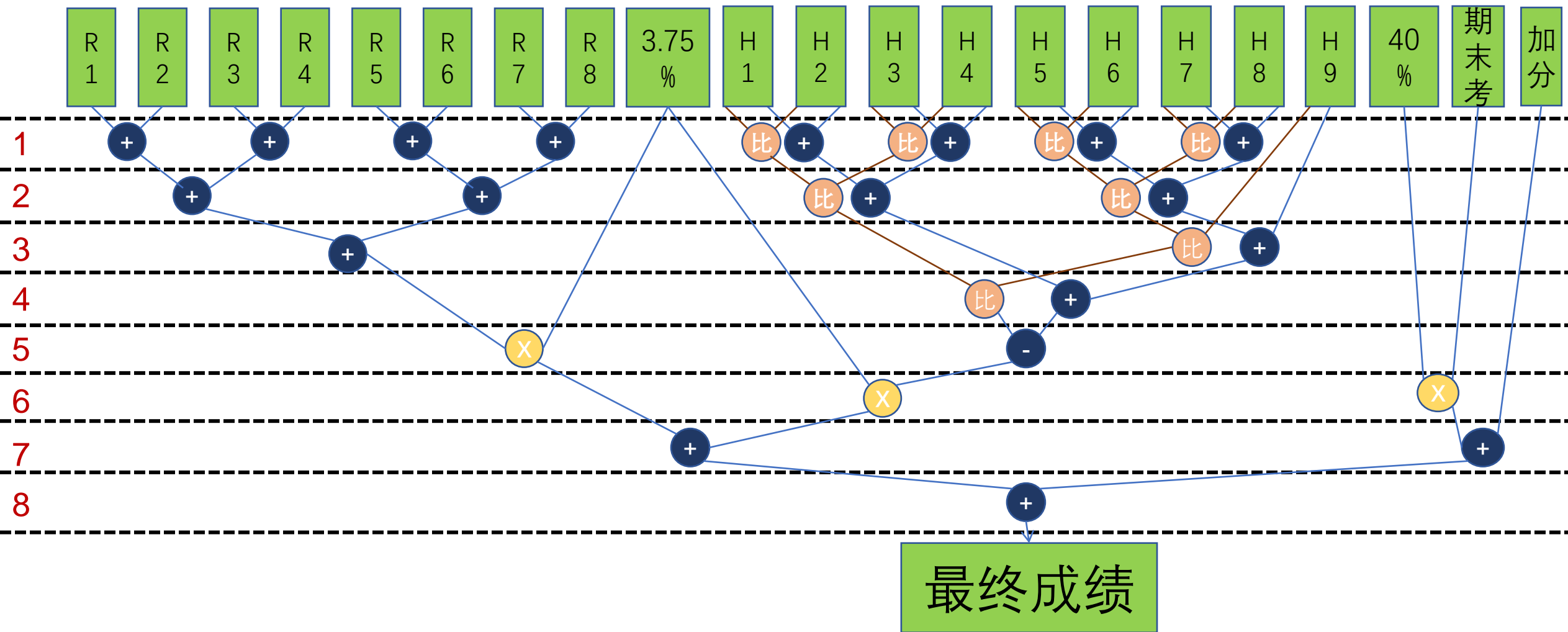
在正式开始之前



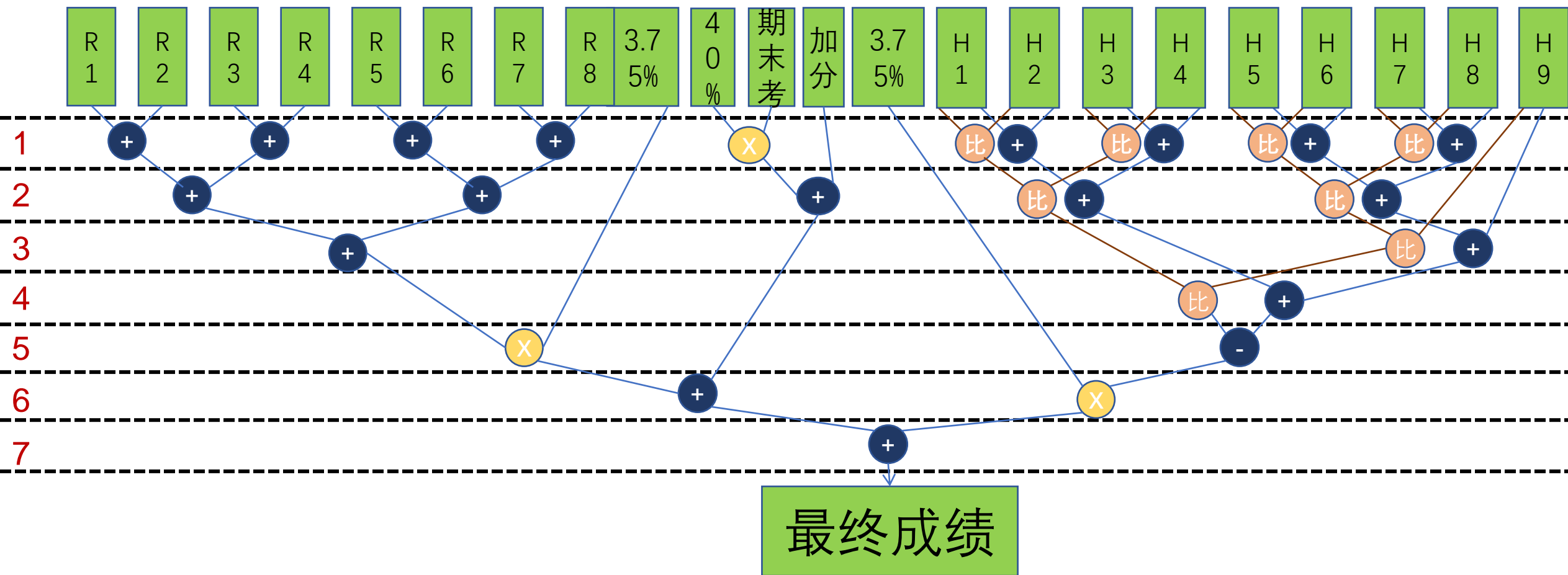
软件开发者

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顺序图



顺序图





软件开发者

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在这门课程中我们主要关注

高级
综合

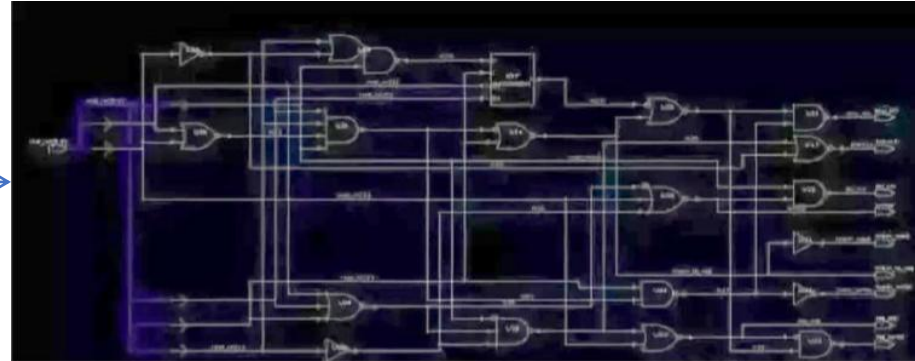
```
LIBRARY IEEE;
USE IEEE.STD_LOGIC_1164.ALL;

ENTITY MUX41A IS
    PORT (a, b, c, d: IN STD_LOGIC;
          s0, s1: IN STD_LOGIC;
          y: OUT STD_LOGIC);
END ENTITY MUX41A;

ARCHITECTURE BHV OF MUX41A IS
    SIGNAL S: STD_LOGIC_VECTOR(1 DOWNTO 0);
BEGIN
    S <= s1 & s0;
    PROCESS (s)
    BEGIN
        CASE (S) IS
            WHEN "00" => y <= a;
            WHEN "01" => y <= b;
            WHEN "10" => y <= c;
            WHEN "11" => y <= d;
            WHEN OTHERS => NULL;
        END CASE;
    END PROCESS;
END ARCHITECTURE BHV;
```

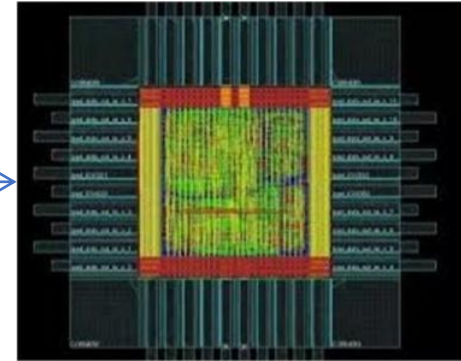
HDL代码

逻辑
综合



门级网表

布局
布线



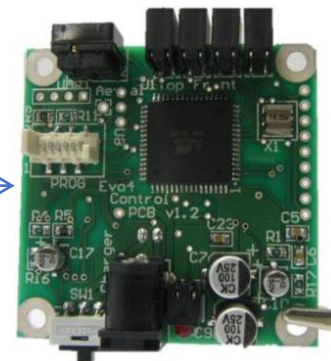
物理版图

工厂



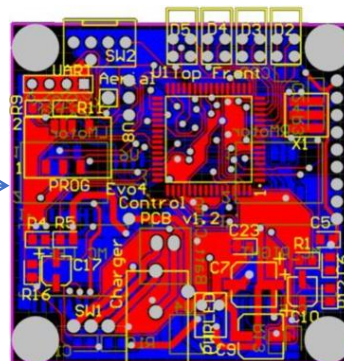
芯片

工厂



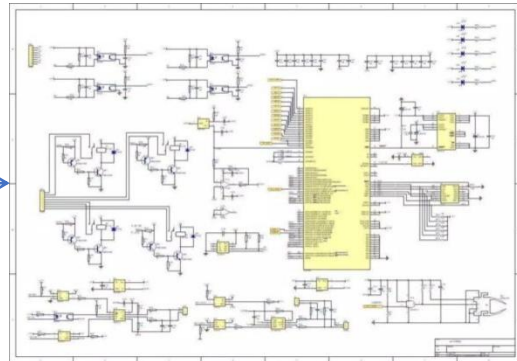
PCB实物板

工厂



PCB设计文件

PCB
设计

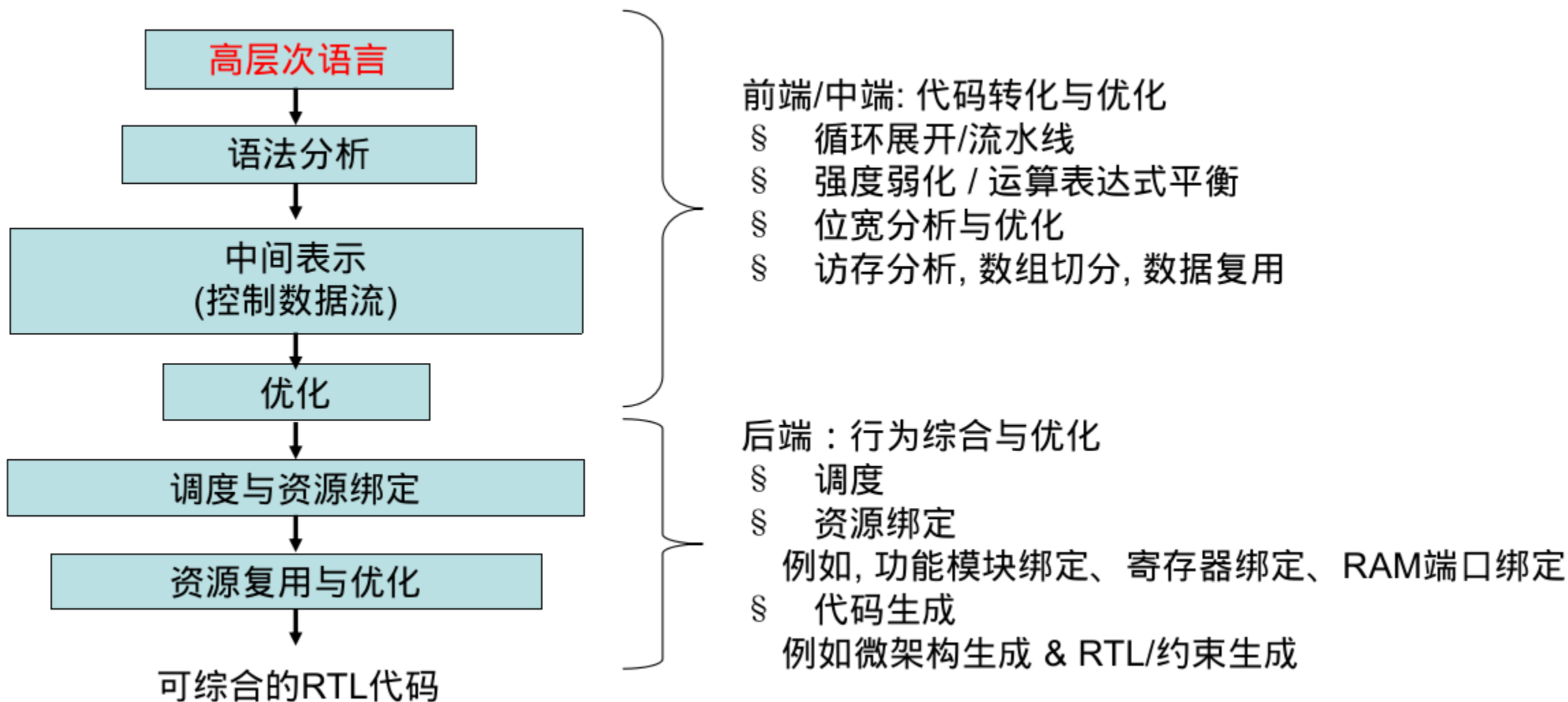


电路原理图

原理图
设计

高层次综合编译器的基本流程

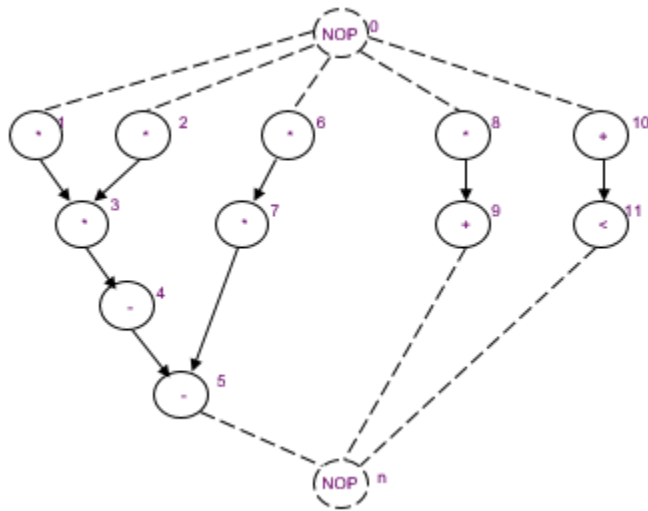
The basic flow of a high-level synthesis compiler



高层次综合

High-level synthesis

- 以顺序图（即依赖图或数据流图）为行为参考



- 操作时序的约束
- 硬件资源可用性的约束

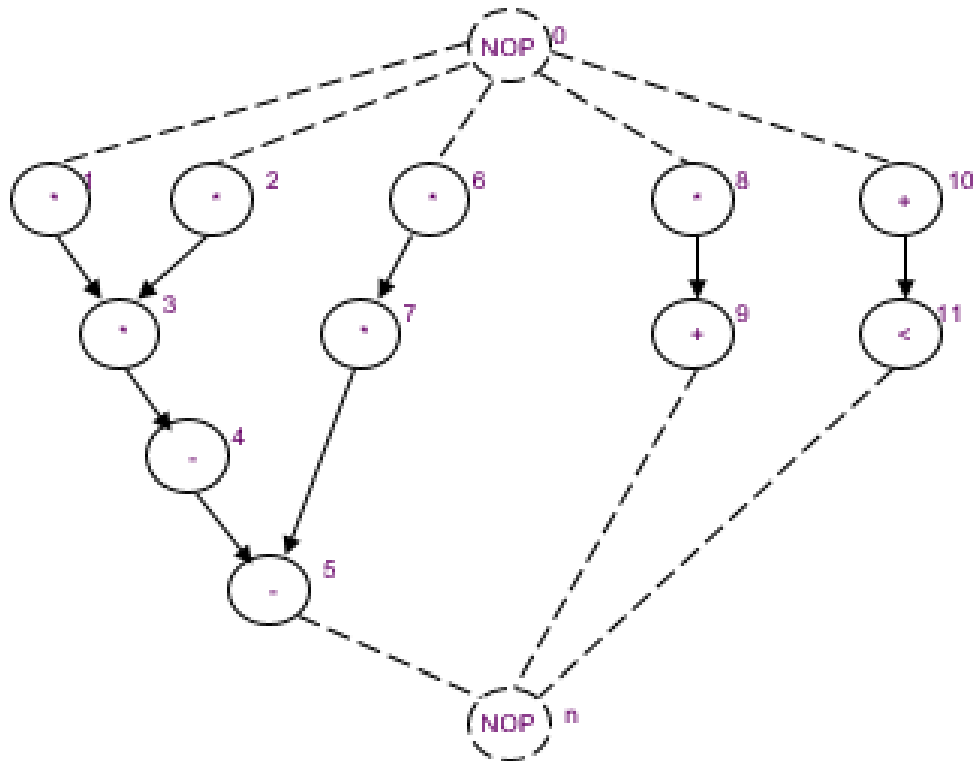
目标:

- 生成带有开始时间和资源注释的顺序图
- 可能有多个可行的解决方案（探索权衡）
- 满足约束，最小化目标：
 - 在面积约束的基础上最大化效率
 - 在效率约束的基础上最小化面积

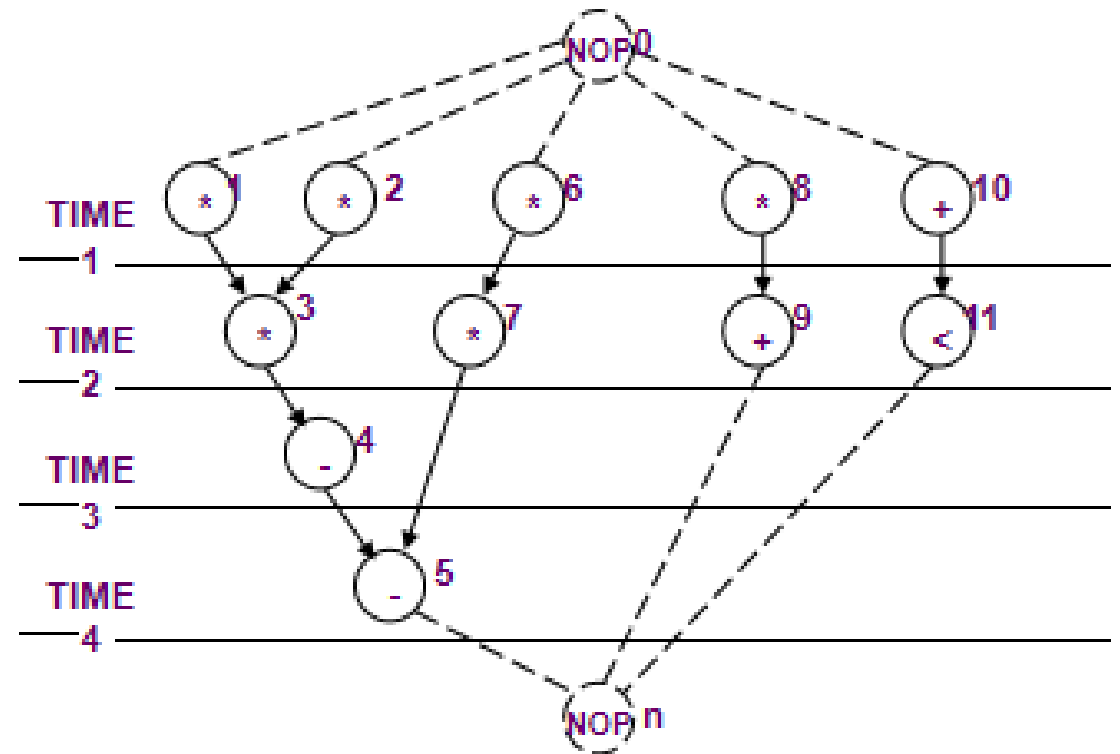
在时间域中的综合

Synthesis in the time domain

- 调度：
- 为每个操作关联一个开始时间



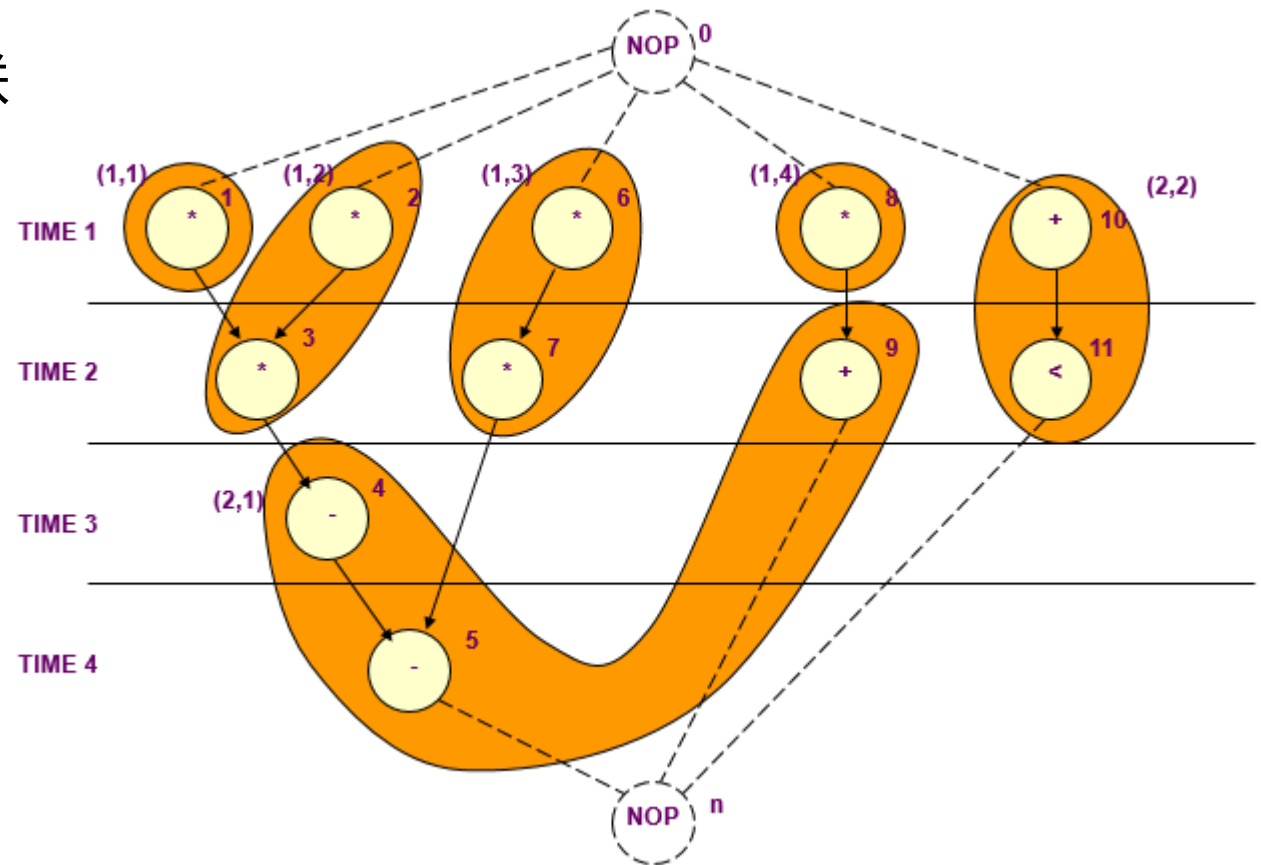
- 已调度的顺序图：
- 带有开始时间注释的顺序图



在空间域中的综合

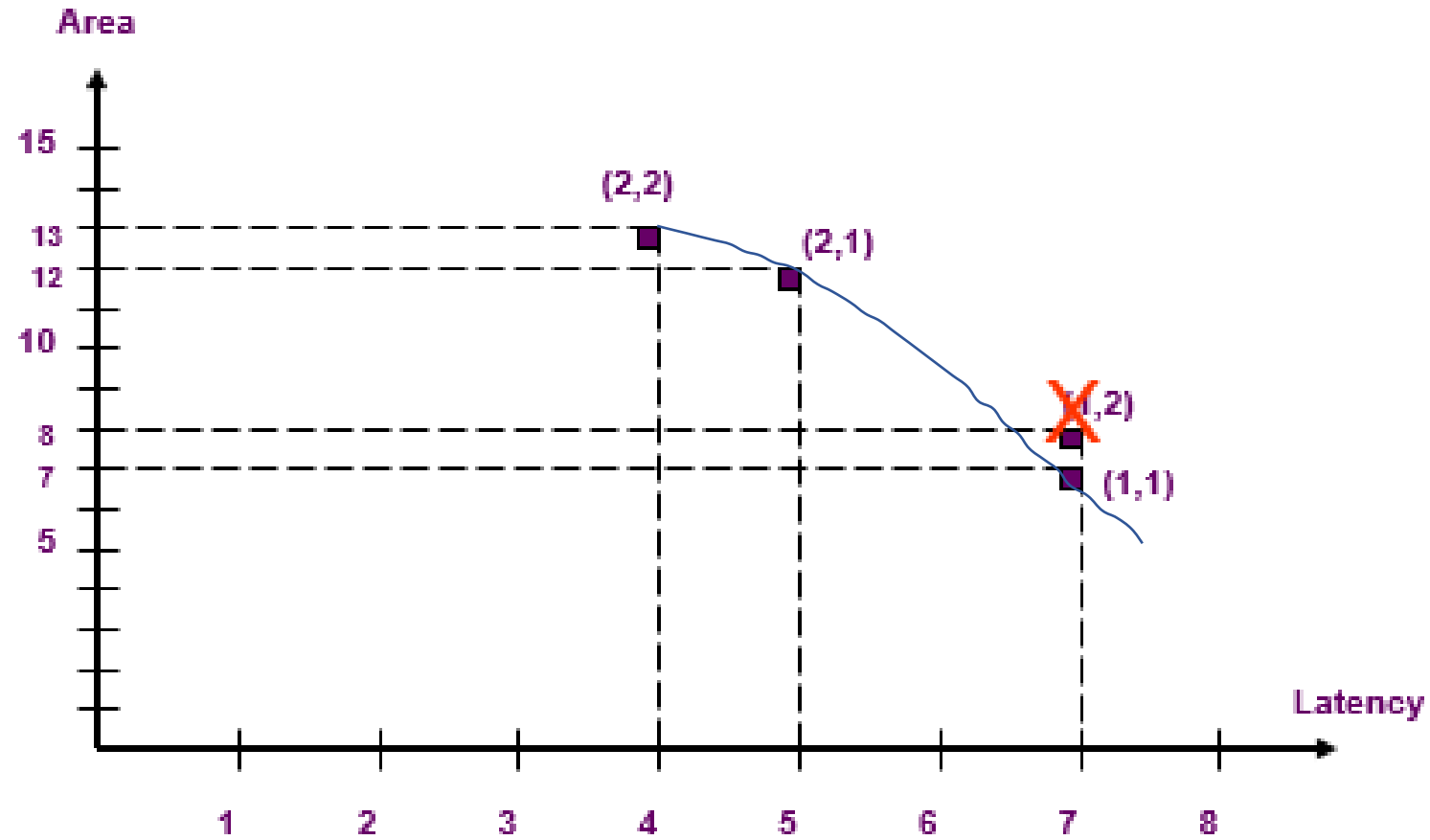
Synthesis in the spatial domain

- 绑定：
 - 将资源与具有相同类型的每个操作关联
 - 确定实现的面积
- 共享：
 - 将资源绑定到多个操作
 - 操作不能同时执行
- 绑定的顺序图：
 - 带有资源注释的顺序图



架构优化

Architecture optimization





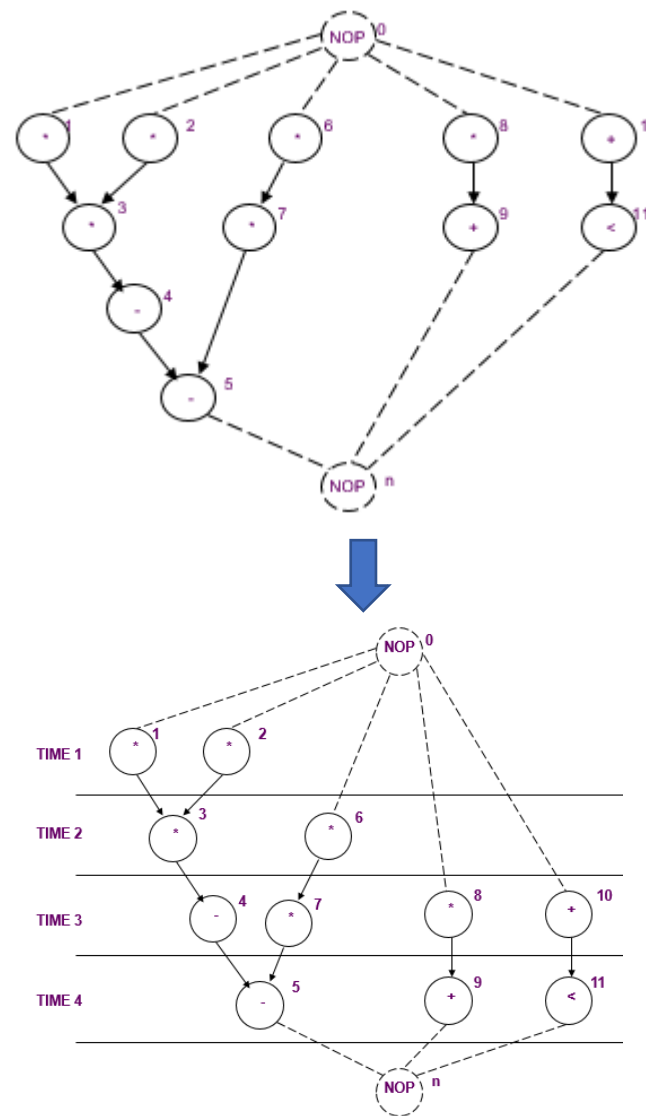
调度算法

Scheduling Algorithms

操作调度

Operation Scheduling

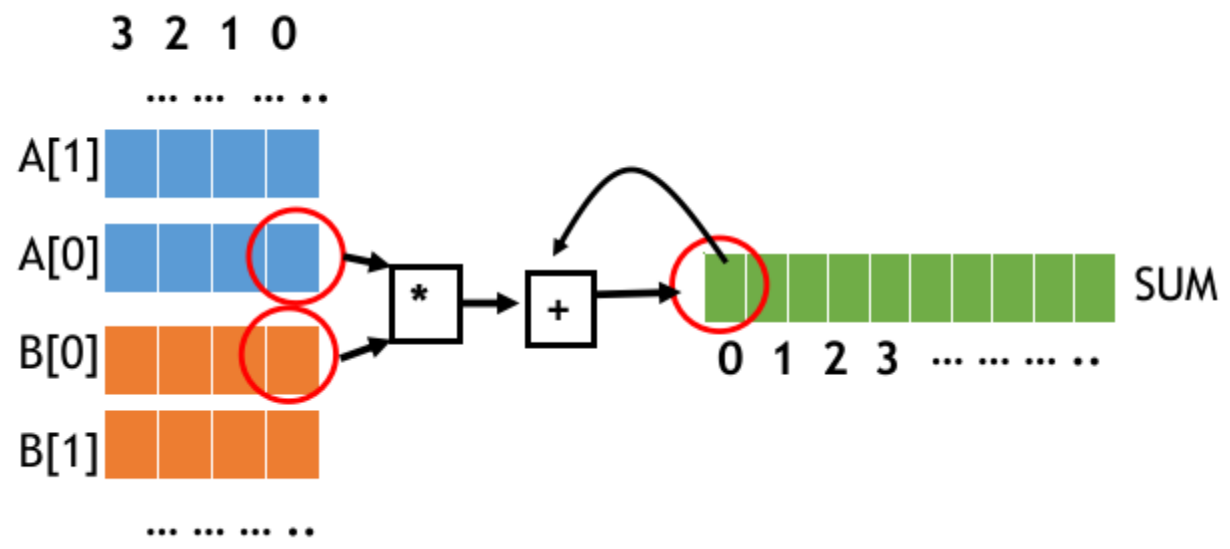
- 输入：
 - 顺序图 $G(V, E)$, 含 n 个顶点
 - 操作延迟 $D = \{d_i; i=0..n\}$
- 输出：
 - 调度 φ : 确定操作 v_i 的开始时间 t_i
 - 延迟 $\lambda = t_n - t_0$
- 目标：权衡资源和时间

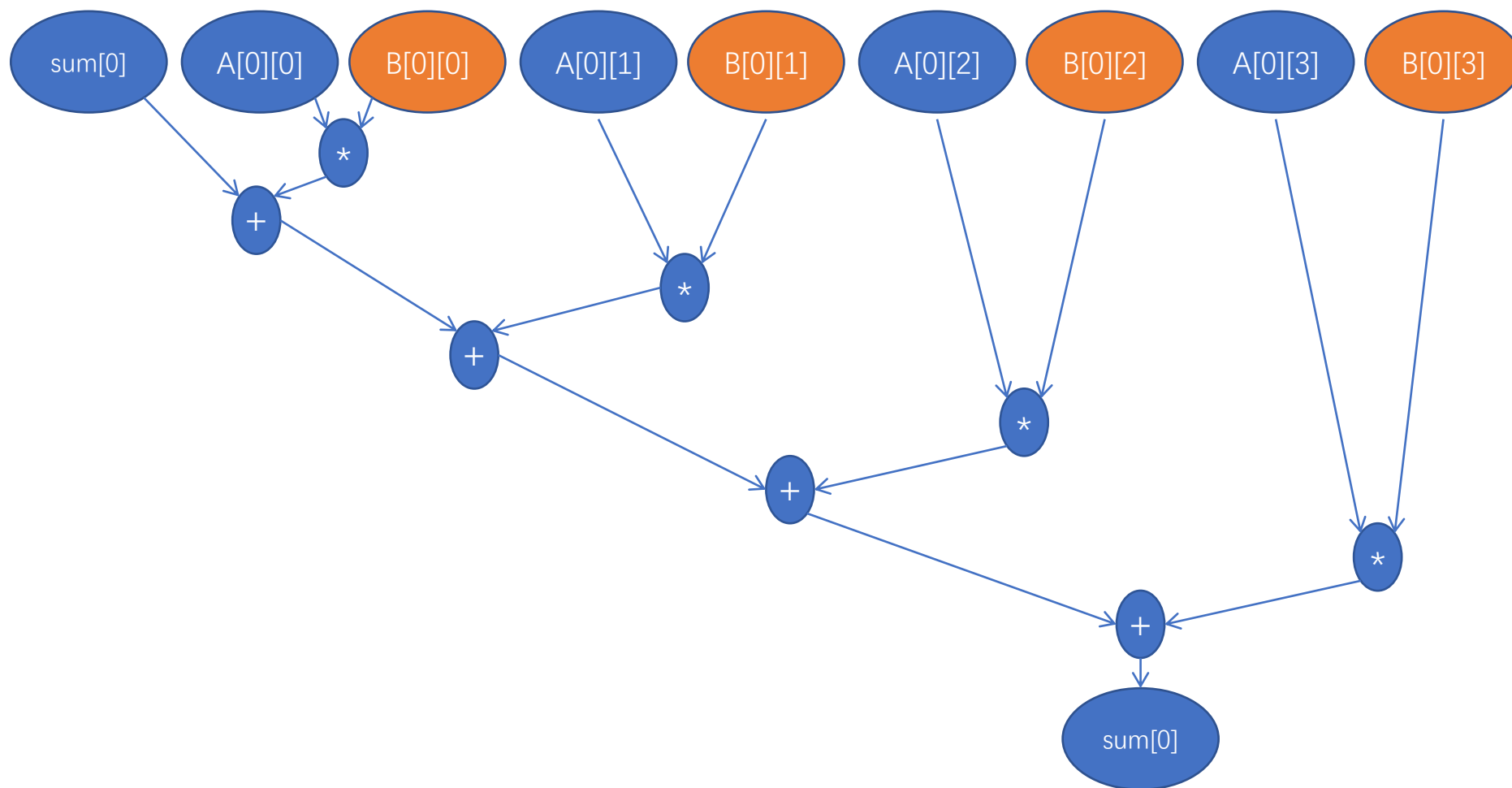


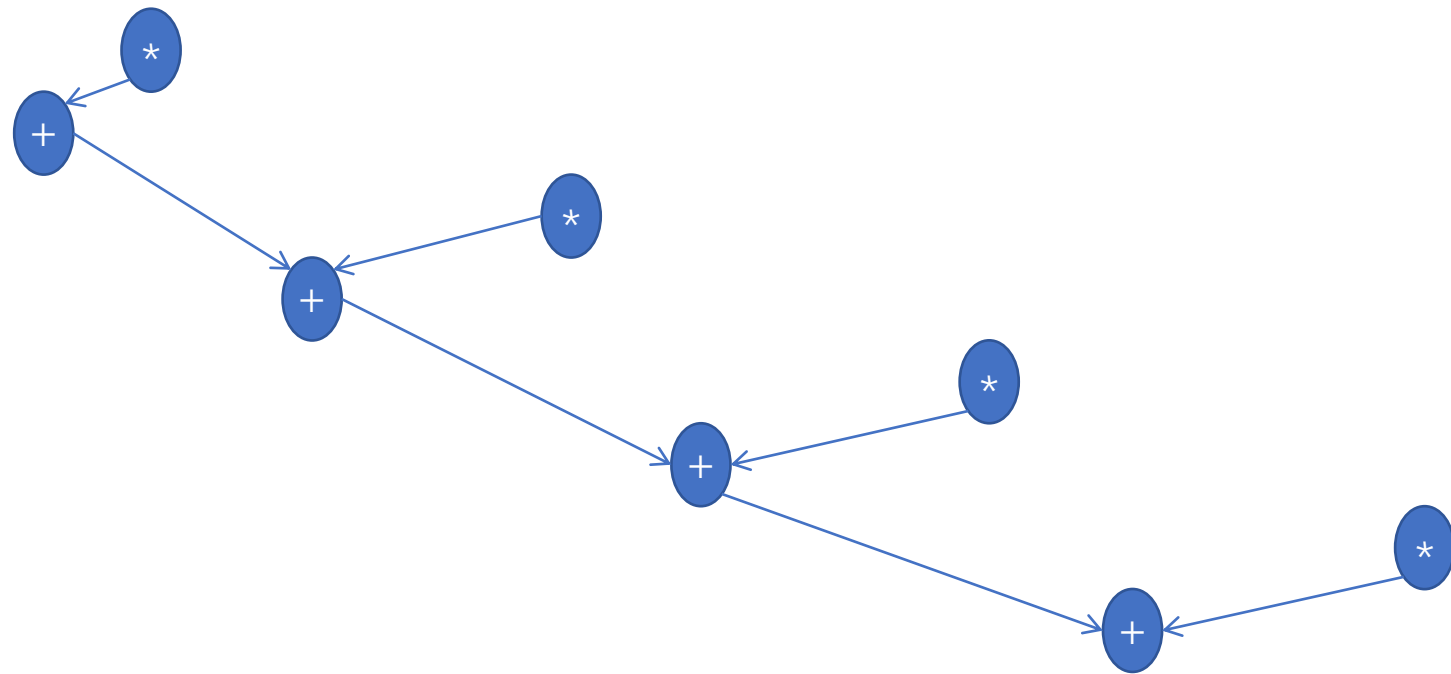
```

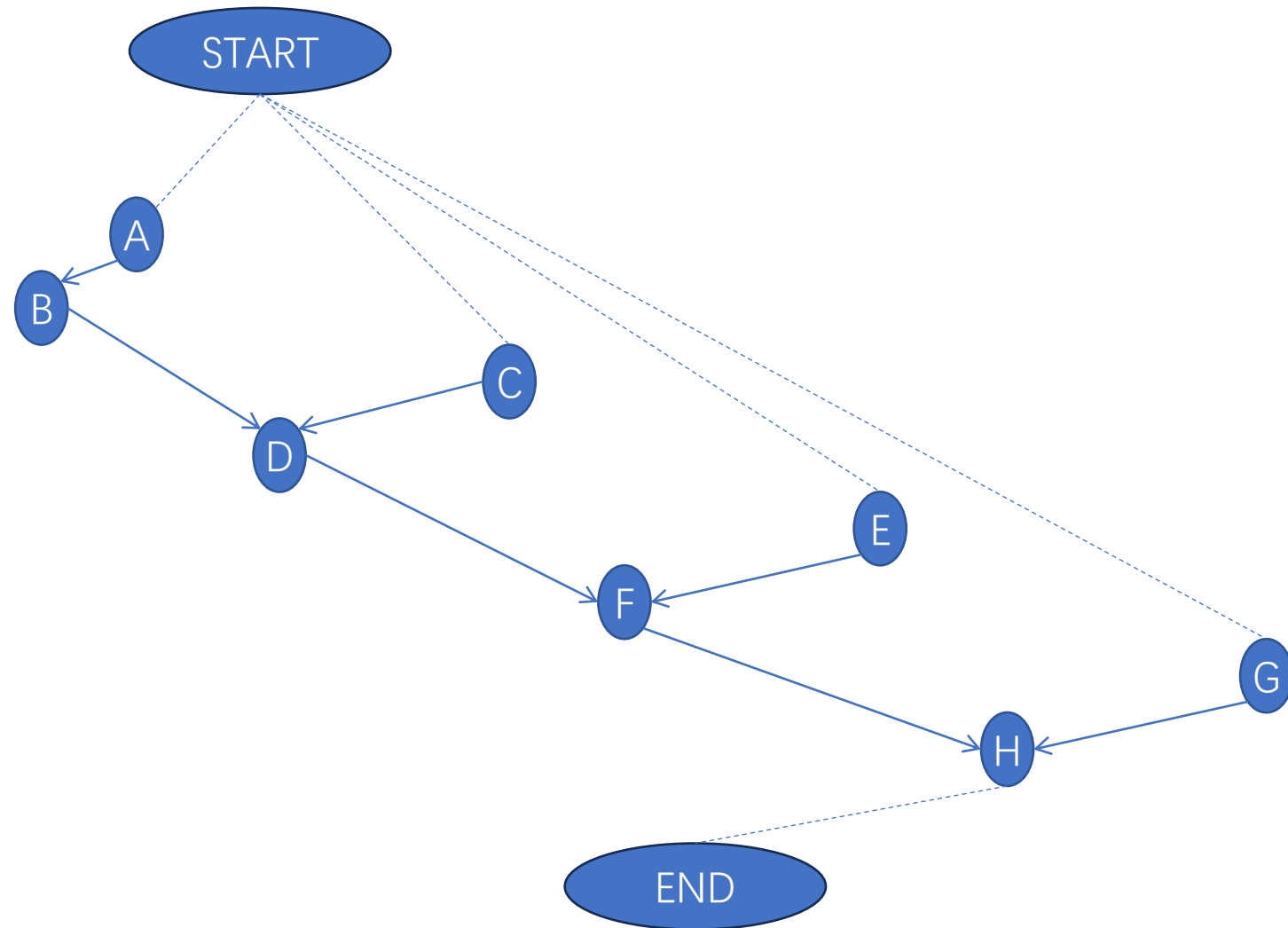
void vectorDot(float A[10][4], float B[10][4], float SUM[10])
{
    for (int i=0;i<10;i++)
    {
        SUM[i] = 0;
        for (int j=0;<4;j++)
        {
            SUM[i] += A[i][j] * B[i][j];
        }
    }
}

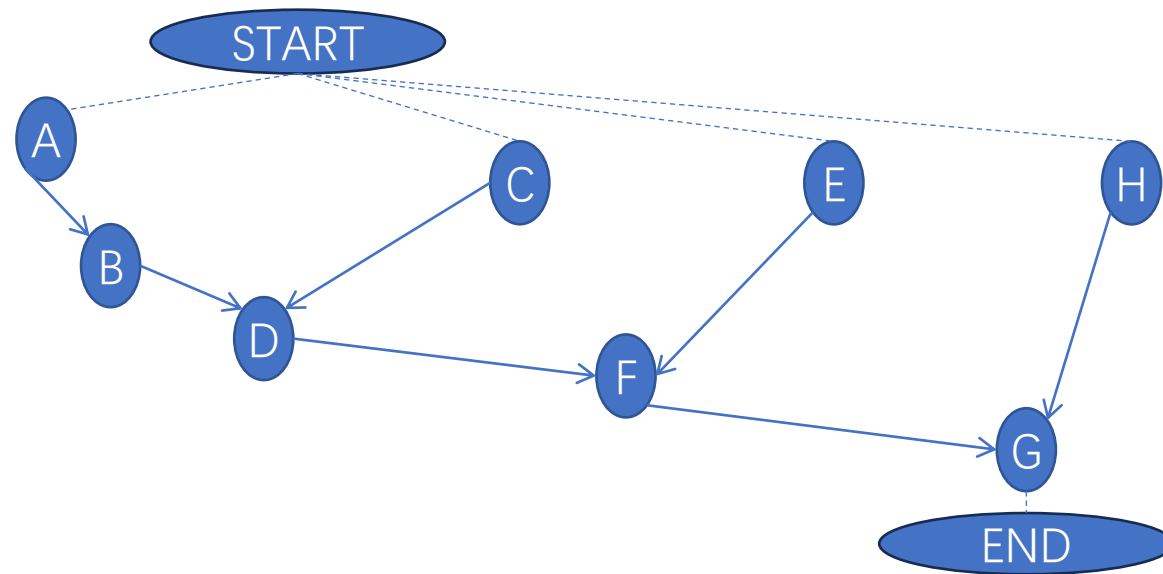
```











操作调度

Operation Scheduling

- 输入：
 - 顺序图 $G(V, E)$, 含 n 个顶点
 - 操作延迟 $D = \{d_i: i=0..n\}$
- 输出：
 - 调度 φ : 确定操作 v_i 的开始时间 t_i
 - 延迟 $\lambda = t_n - t_0$
- 目标: 权衡资源和时间

输入文件内容:

7

START 1

A 2

B 1

C 3

D 1

E 2

END 1

A B

C B

A E

输出文件内容:

START 0

A 1

C 1

B 4

E 3

D 1

END 5

最小延迟无约束调度问题



最小延迟无约束调度

Minimum-latency Unconstrained Scheduling

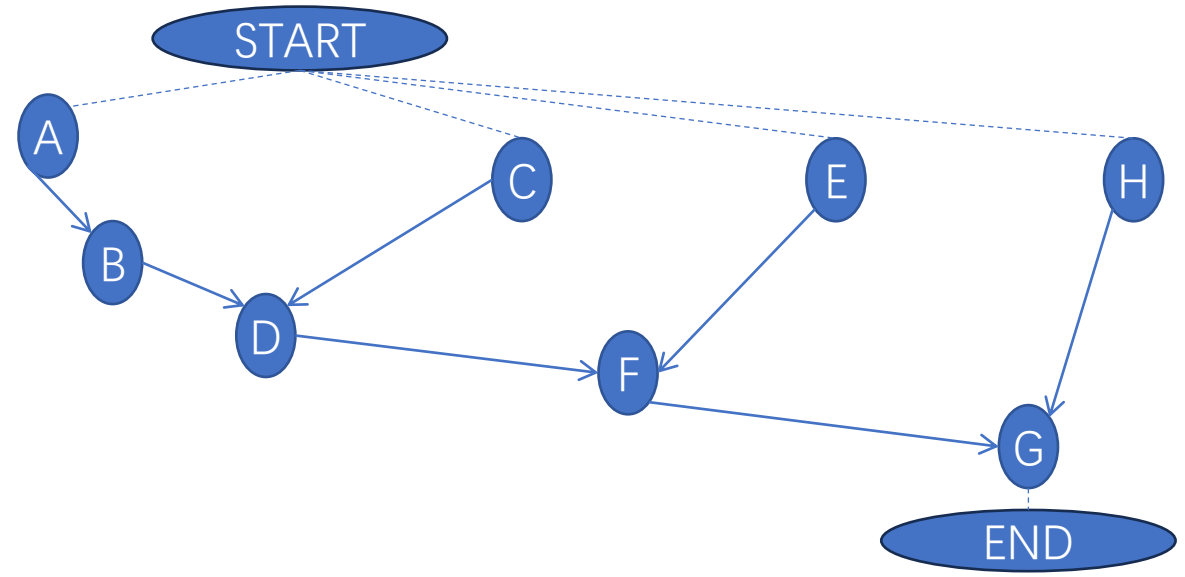
- 最简单的情况：没有约束，找到最小的延迟
- 输入：顶点集 V ，延迟 D 和偏序关系集 E ,
- 输出：对每个操作进行标记 $\varphi: V \rightarrow \mathbb{Z}^+$ ，使得：
 - $t_i = \varphi(v_i)$
 - $t_i \geq t_j + d_j$ 对于所有 $(v_i, v_j) \in E$
 - $\lambda = t_n - t_0$ 是最小的
- 可在多项式时间内解决
- 使用的 ASAP 算法：拓扑顺序

ASAP调度算法

ASAP调度算法

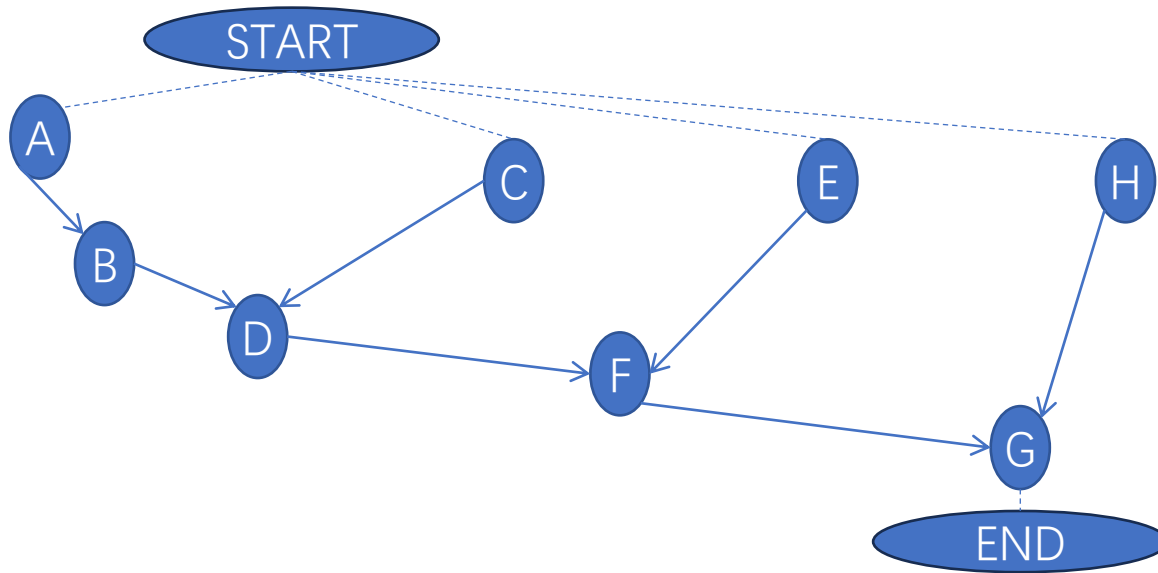
ASAP scheduling algorithm

```
1 ASAP (  $G_s(V,E)$  ) {  
2     Schedule  $v_0$  by setting  $t_0 = 0$ ;  
3     repeat {  
4         Select a vertex  $v_i$  whose predecessors are all scheduled;  
5         Schedule  $v_i$  by setting  $t_i = \max t_j + d_j$ ;  
6     }  
7     until ( $v_n$  is scheduled);  
8     return (  $G$  );  
9 }
```

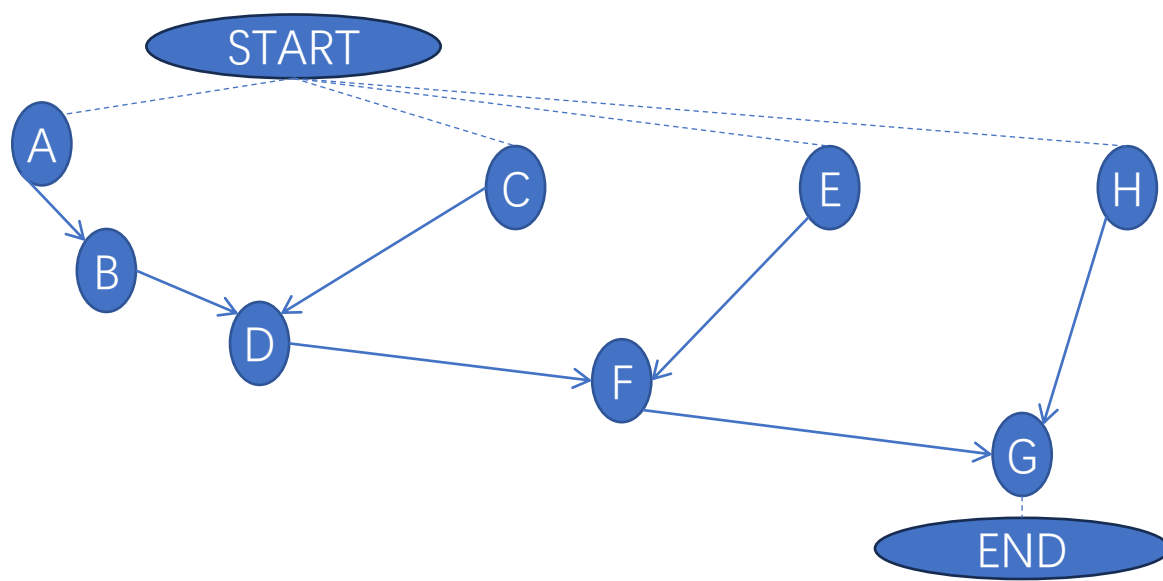


ASAP调度算法

ASAP scheduling algorithm



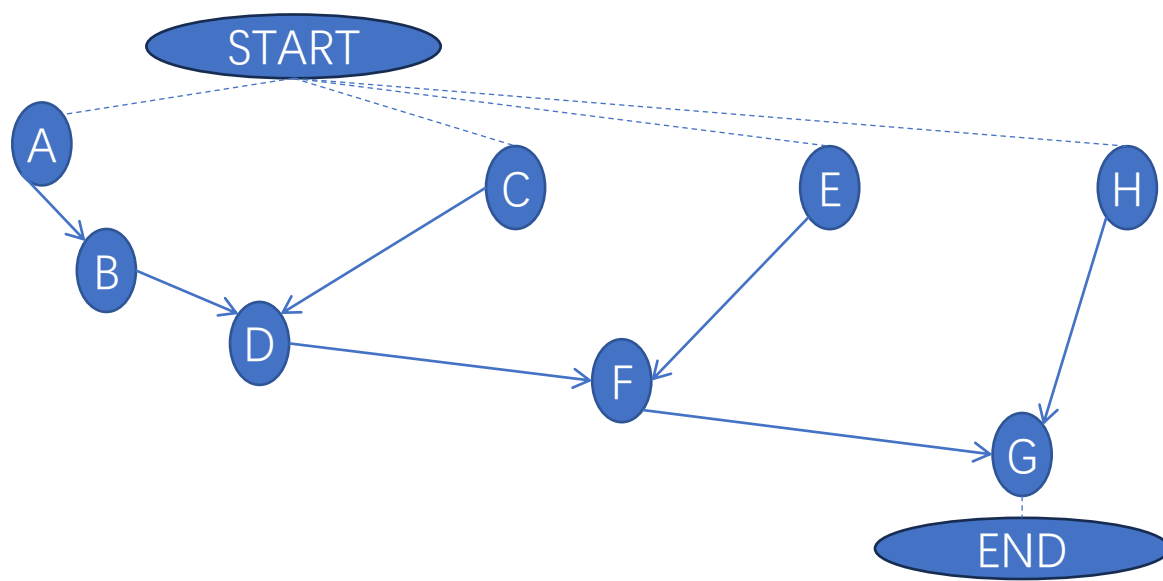
| | |
|---------|---------|
| 10 | START A |
| START 1 | START C |
| A 1 | START E |
| B 1 | START G |
| C 1 | A B |
| D 1 | C D |
| E 1 | B D |
| F 1 | E F |
| G 1 | D F |
| H 1 | H G |
| END 1 | F G |
| | G END |



START 0

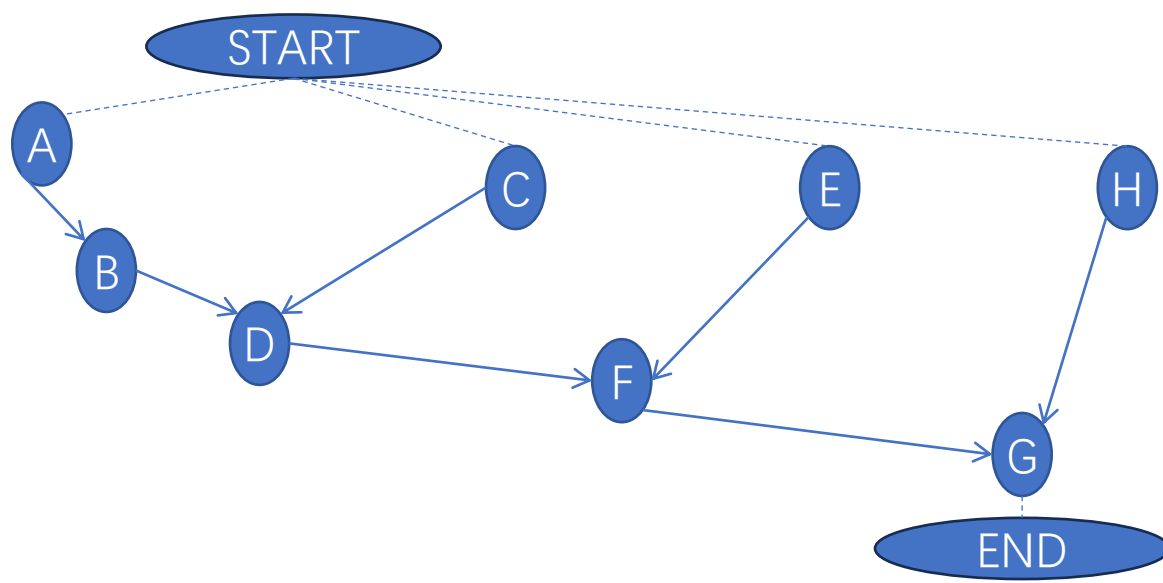
```

1 ASAP (  $G_s(V,E)$  ) {
2   Schedule  $v_0$  by setting  $t_0 = 0$ ;
3   repeat {
4     Select a vertex  $v_i$  whose predecessors are all scheduled;
5     Schedule  $v_i$  by setting  $t_i = \max t_j + d_j$ ;
7   until ( $v_n$  is scheduled);
8   return (  $G$  );}
  
```



START 0

```
1 ASAP (  $G_s(V,E)$  ) {  
2     Schedule  $v_0$  by setting  $t_0 = 0$ ;  
3     repeat {  
4         Select a vertex  $v_i$  whose predecessors are all scheduled;  
5         Schedule  $v_i$  by setting  $t_i = \max t_j + d_j$  ;}  
7     until ( $v_n$  is scheduled);  
8     return (  $G$  );}
```

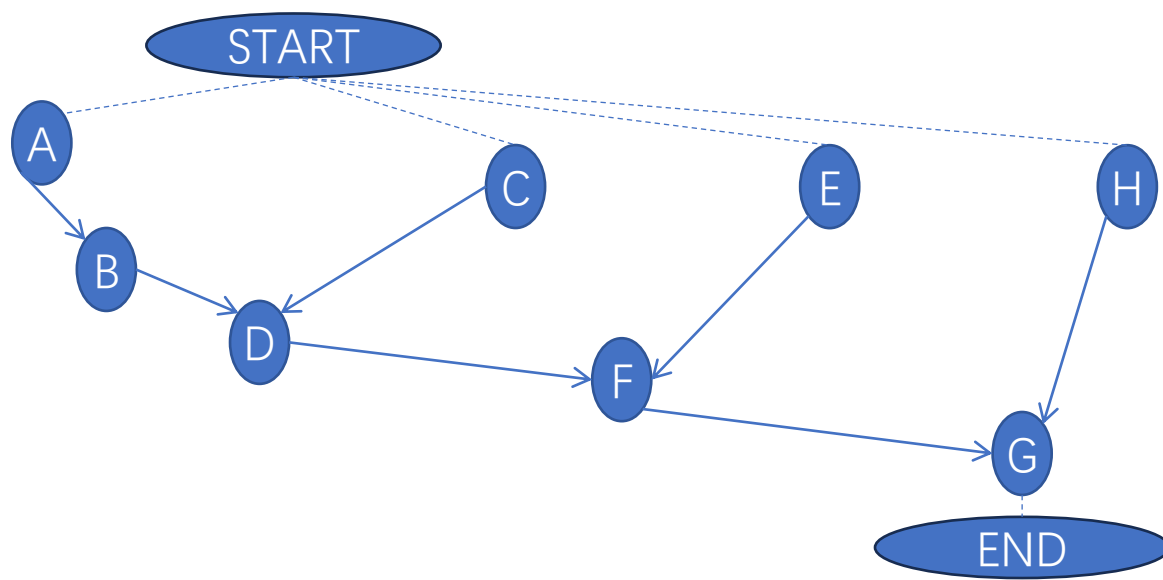


START 0
A



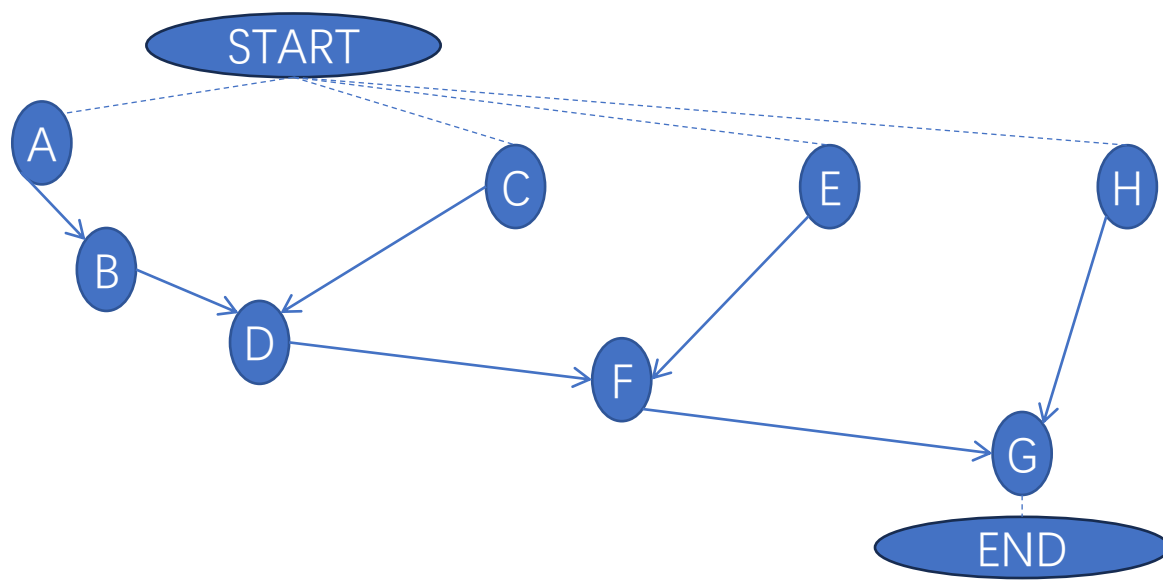
```

1 ASAP (  $G_s(V,E)$  ) {
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```



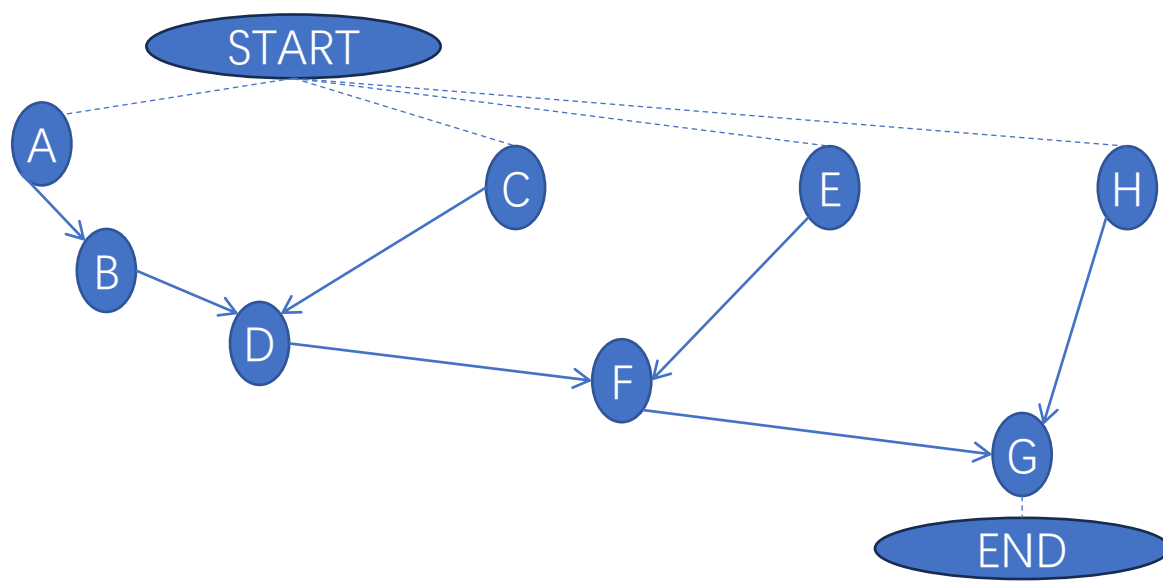
START 0
A 1

```
1 ASAP (  $G_s(V,E)$  ) {  
2     Schedule  $v_0$  by setting  $t_0 = 0$ ;  
3     repeat {  
4         Select a vertex  $v_i$  whose predecessors are all scheduled;  
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7     until ( $v_n$  is scheduled);  
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```



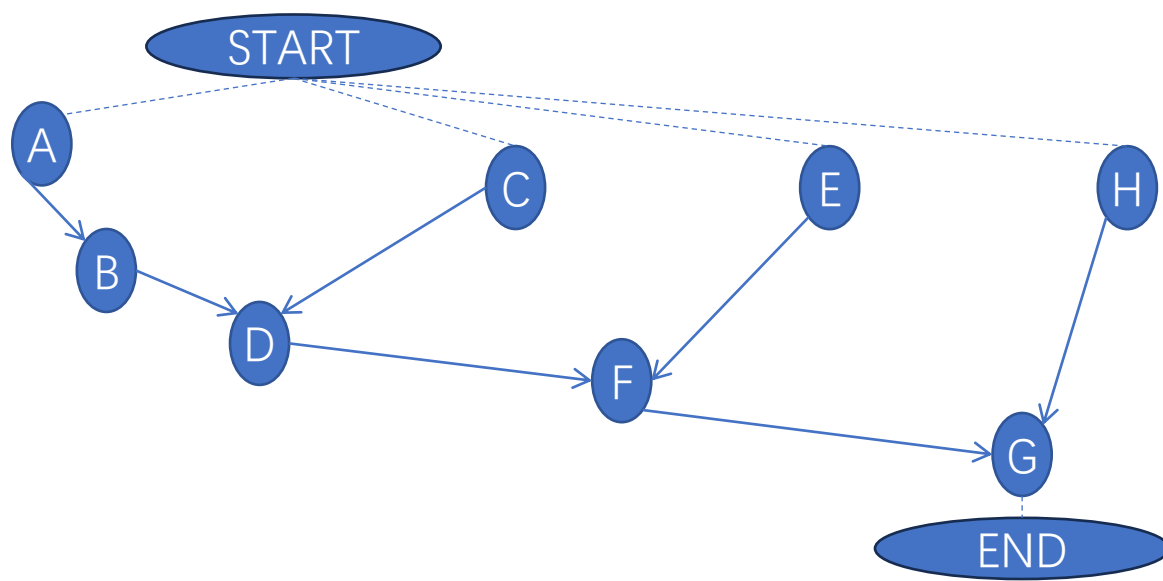
START 0
A 1

```
1 ASAP (  $G_s(V,E)$  ) {  
2     Schedule  $v_0$  by setting  $t_0 = 0$ ;  
3     repeat {  
4         Select a vertex  $v_i$  whose predecessors are all scheduled;  
5         Schedule  $v_i$  by setting  $t_i = \max t_j + d_j$  ;}  
7     until ( $v_n$  is scheduled);  
8     return (  $G$  );}
```



START 0
A 1
B

```
1 ASAP (  $G_s(V,E)$  ) {  
2     Schedule  $v_0$  by setting  $t_0 = 0$ ;  
3     repeat {  
4         Select a vertex  $v_i$  whose predecessors are all scheduled;  
5         Schedule  $v_i$  by setting  $t_i = \max t_j + d_j$ ;}  
7     until ( $v_n$  is scheduled);  
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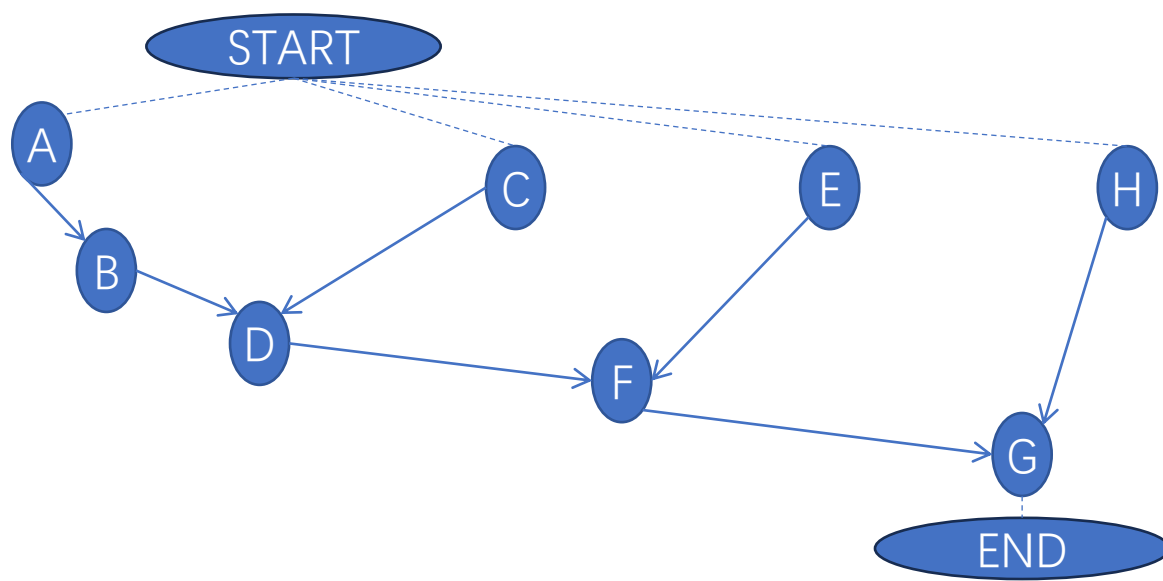


START 0

A 1

B 2

```
1 ASAP (  $G_s(V,E)$  ) {  
2     Schedule  $v_0$  by setting  $t_0 = 0$ ;  
3     repeat {  
4         Select a vertex  $v_i$  whose predecessors are all scheduled;  
5         Schedule  $v_i$  by setting  $t_i = \max t_j + d_j$  ;}  
7     until ( $v_n$  is scheduled);  
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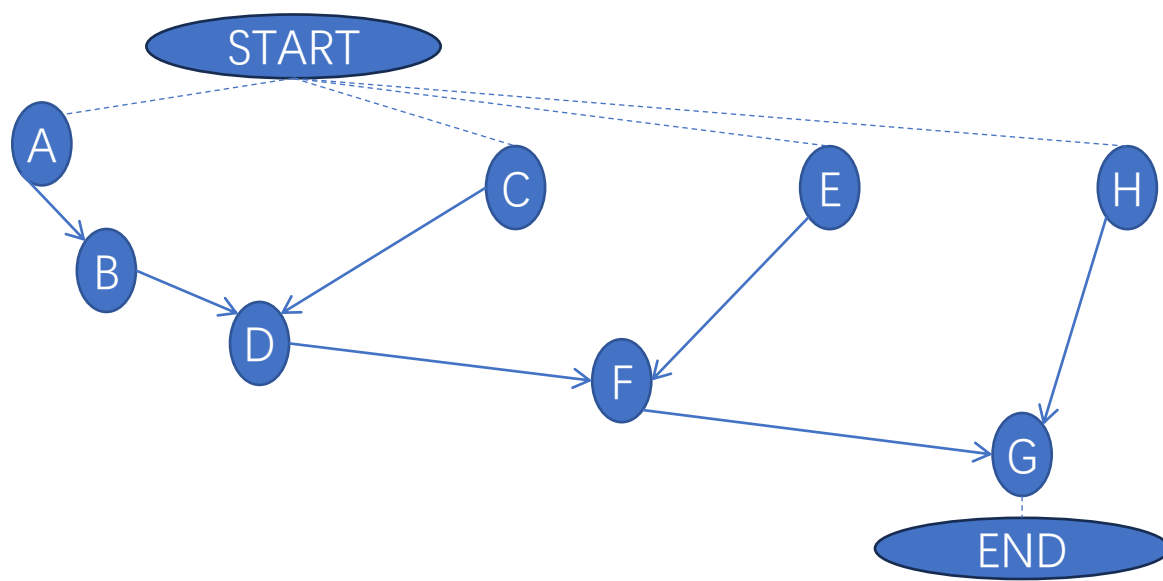



START 0

A 1

B 2

```
1 ASAP (  $G_s(V,E)$  ) {  
2     Schedule  $v_0$  by setting  $t_0 = 0$ ;  
3     repeat {  
4         Select a vertex  $v_i$  whose predecessors are all scheduled;  
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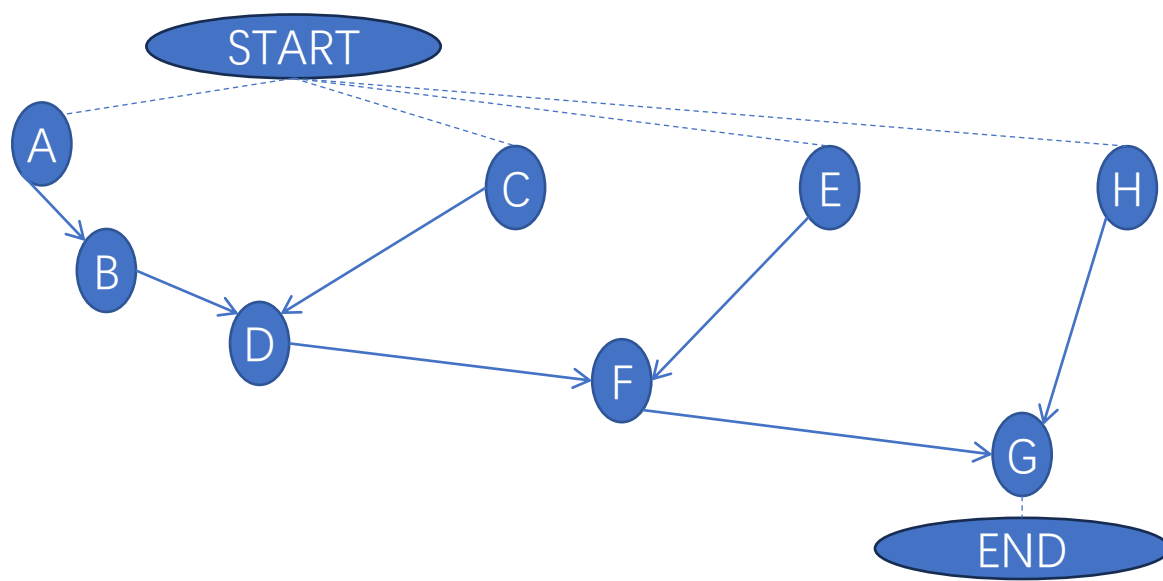
START 0

A 1

B 2

C

```
1 ASAP (  $G_s(V,E)$  ) {  
2     Schedule  $v_0$  by setting  $t_0 = 0$ ;  
3     repeat {  
4         Select a vertex  $v_i$  whose predecessors are all scheduled;  
5         Schedule  $v_i$  by setting  $t_i = \max t_j + d_j$ ;}  
7     until ( $v_n$  is scheduled);  
8     return (  $G$  );}
```



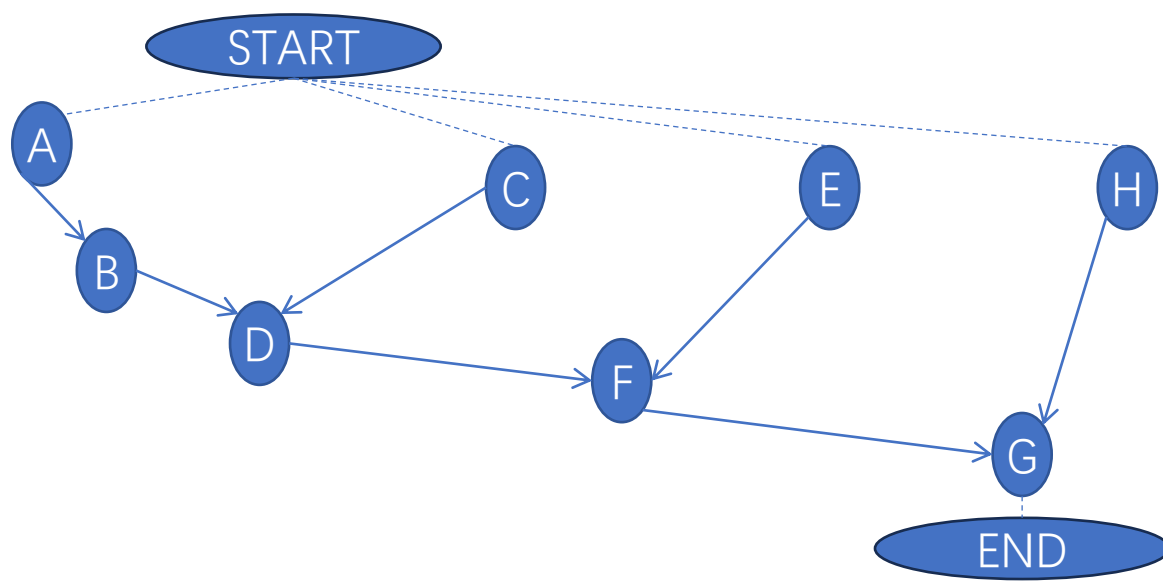
START 0

A 1

B 2

C 1

```
1 ASAP (  $G_s(V,E)$  ) {  
2     Schedule  $v_0$  by setting  $t_0 = 0$ ;  
3     repeat {  
4         Select a vertex  $v_i$  whose predecessors are all scheduled;  
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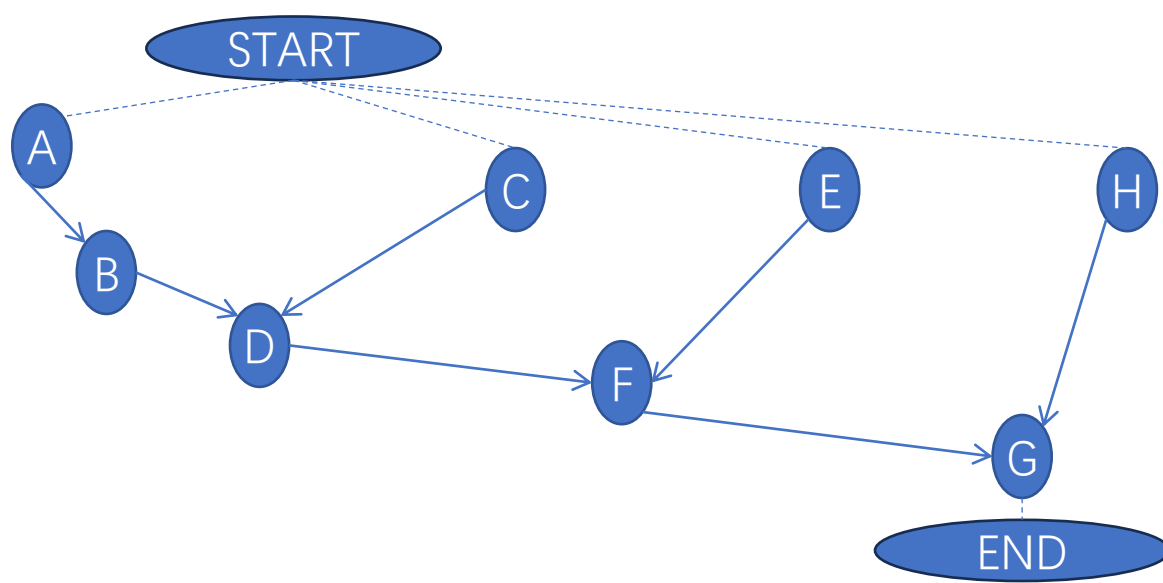
START 0

A 1

B 2

C 1

```
1 ASAP (  $G_s(V,E)$  ) {  
2     Schedule  $v_0$  by setting  $t_0 = 0$ ;  
3     repeat {  
4         Select a vertex  $v_i$  whose predecessors are all scheduled;  
5         Schedule  $v_i$  by setting  $t_i = \max t_j + d_j$ ;}  
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START 0

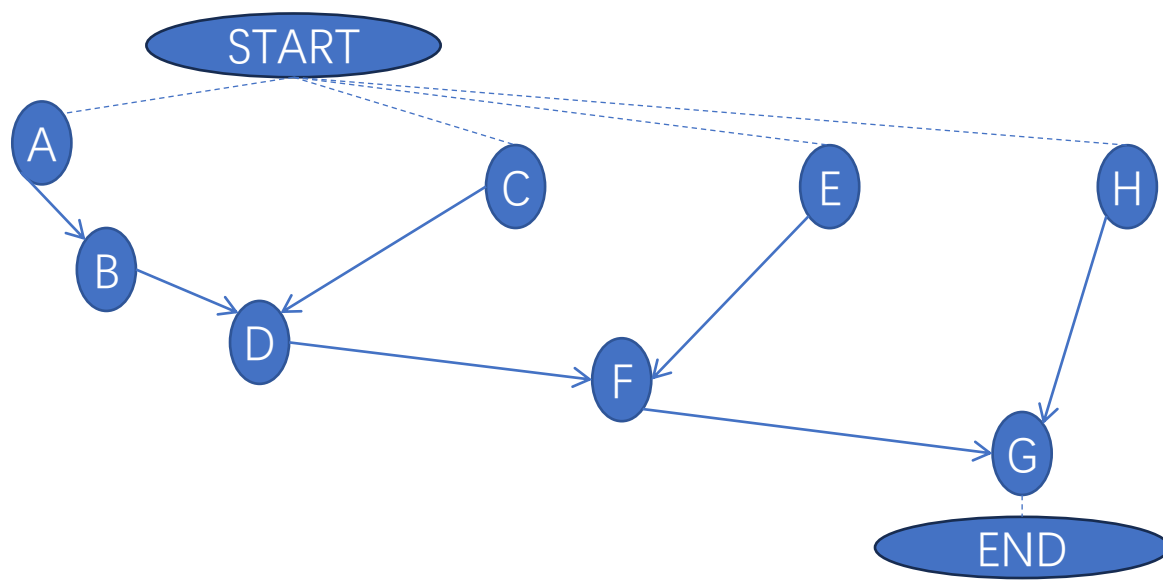
A 1

B 2

C 1

D

```
1 ASAP (  $G_s(V,E)$  ) {  
2     Schedule  $v_0$  by setting  $t_0 = 0$ ;  
3     repeat {  
4         Select a vertex  $v_i$  whose predecessors are all scheduled;  
5         Schedule  $v_i$  by setting  $t_i = \max t_j + d_j$ ;}  
7     until ( $v_n$  is scheduled);  
8     return (  $G$  );}
```



START 0

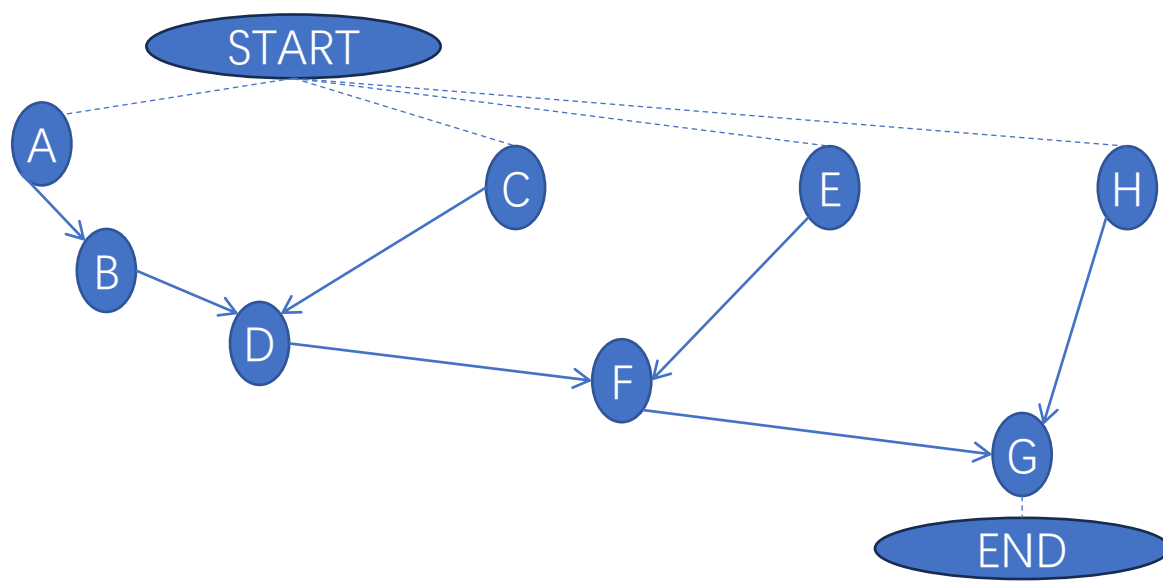
A 1

B 2

C 1

D 3

```
1 ASAP (  $G_s(V,E)$  ) {  
2     Schedule  $v_0$  by setting  $t_0 = 0$ ;  
3     repeat {  
4         Select a vertex  $v_i$  whose predecessors are all scheduled;  
5         Schedule  $v_i$  by setting  $t_i = \max t_j + d_j$  ;}  
7     until ( $v_n$  is scheduled);  
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START 0

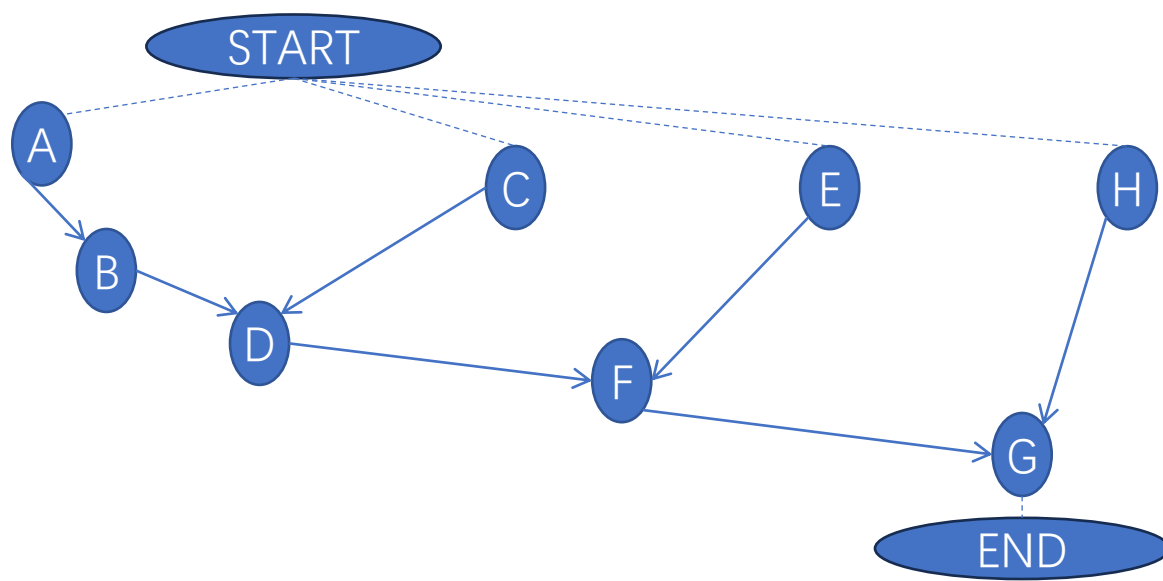
A 1

B 2

C 1

D 3

```
1 ASAP (  $G_s(V,E)$  ) {  
2     Schedule  $v_0$  by setting  $t_0 = 0$ ;  
3     repeat {  
4         Select a vertex  $v_i$  whose predecessors are all scheduled;  
5         Schedule  $v_i$  by setting  $t_i = \max t_j + d_j$ ;}  
7     until ( $v_n$  is scheduled);  
8     return (  $G$  );}
```



START 0

A 1

B 2

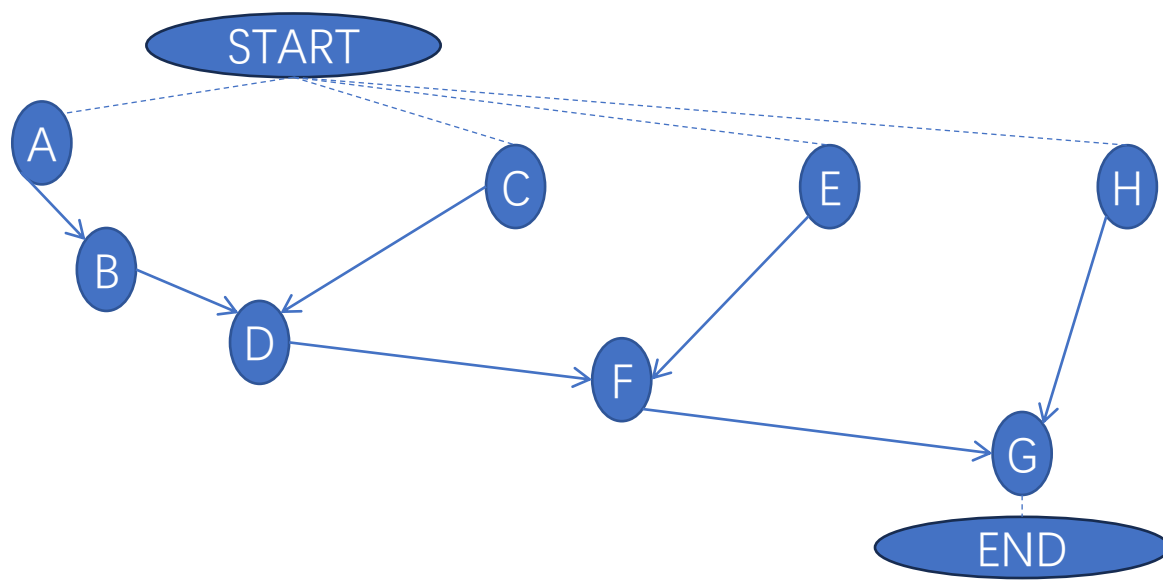
C 1

D 3

E

```

1 ASAP (  $G_s(V,E)$  ) {
2   Schedule  $v_0$  by setting  $t_0 = 0$ ;
3   repeat {
4     Select a vertex  $v_i$  whose predecessors are all scheduled;
5     Schedule  $v_i$  by setting  $t_i = \max t_j + d_j$ ; }
7   until ( $v_n$  is scheduled);
8   return (  $G$  ); }
  
```

START 0

A 1

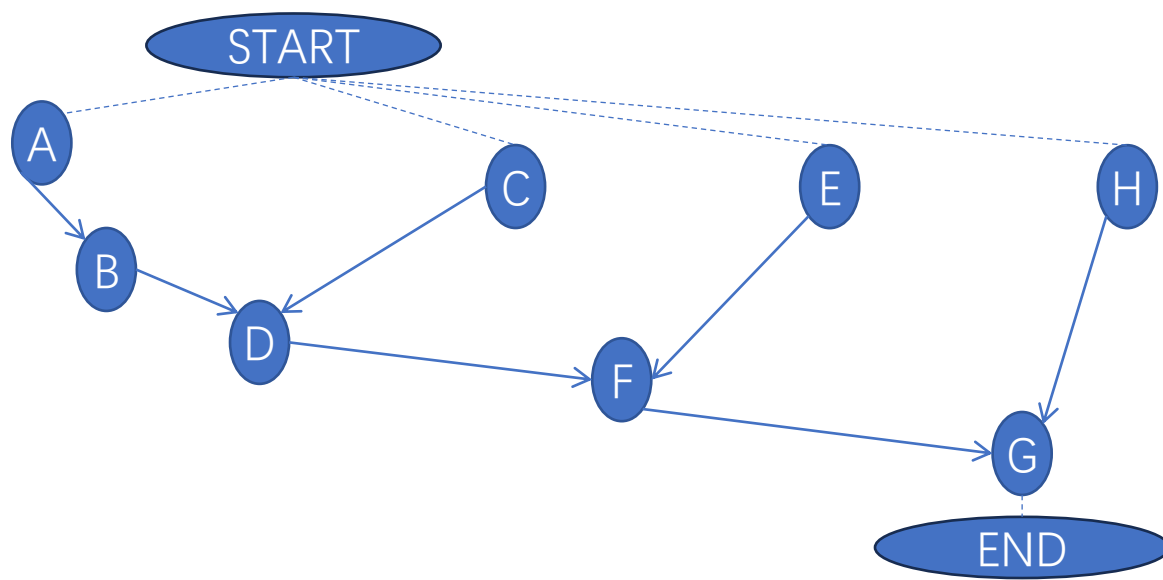
B 2

C 1

D 3

E 1

```
1 ASAP (  $G_s(V,E)$  ) {  
2     Schedule  $v_0$  by setting  $t_0 = 0$ ;  
3     repeat {  
4         Select a vertex  $v_i$  whose predecessors are all scheduled;  
5         Schedule  $v_i$  by setting  $t_i = \max t_j + d_j$ ;}  
7     until ( $v_n$  is scheduled);  
8     return (  $G$  );}
```



START 0

A 1

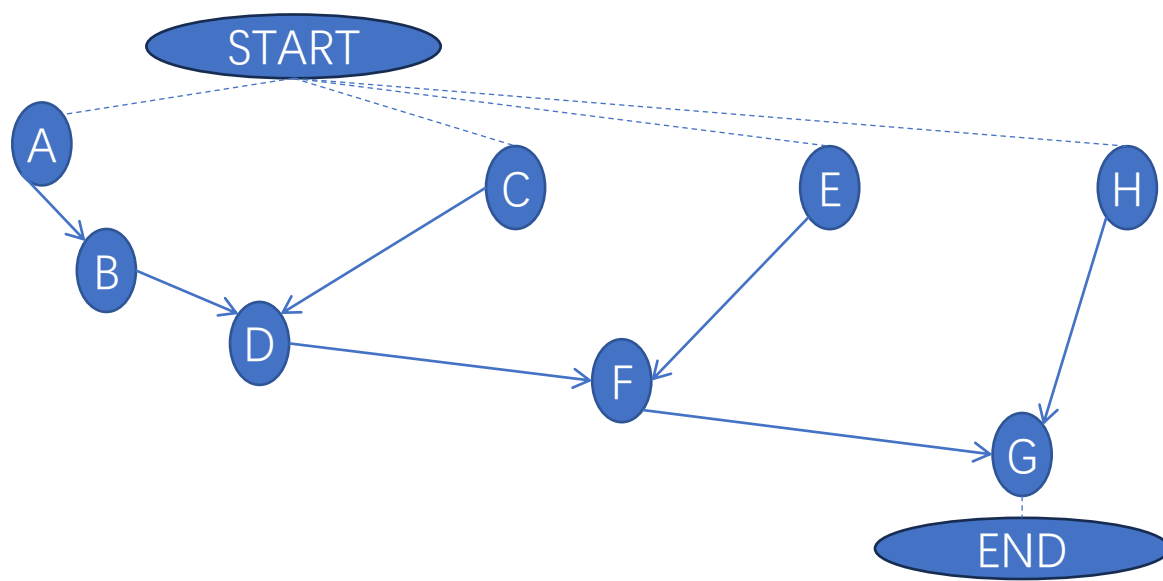
B 2

C 1

D 3

E 1

```
1 ASAP (  $G_s(V,E)$  ) {  
2     Schedule  $v_0$  by setting  $t_0 = 0$ ;  
3     repeat {  
4         Select a vertex  $v_i$  whose predecessors are all scheduled;  
5         Schedule  $v_i$  by setting  $t_i = \max t_j + d_j$ ;}  
7     until ( $v_n$  is scheduled);  
8     return (  $G$  );}
```



START 0

A 1

B 2

C 1

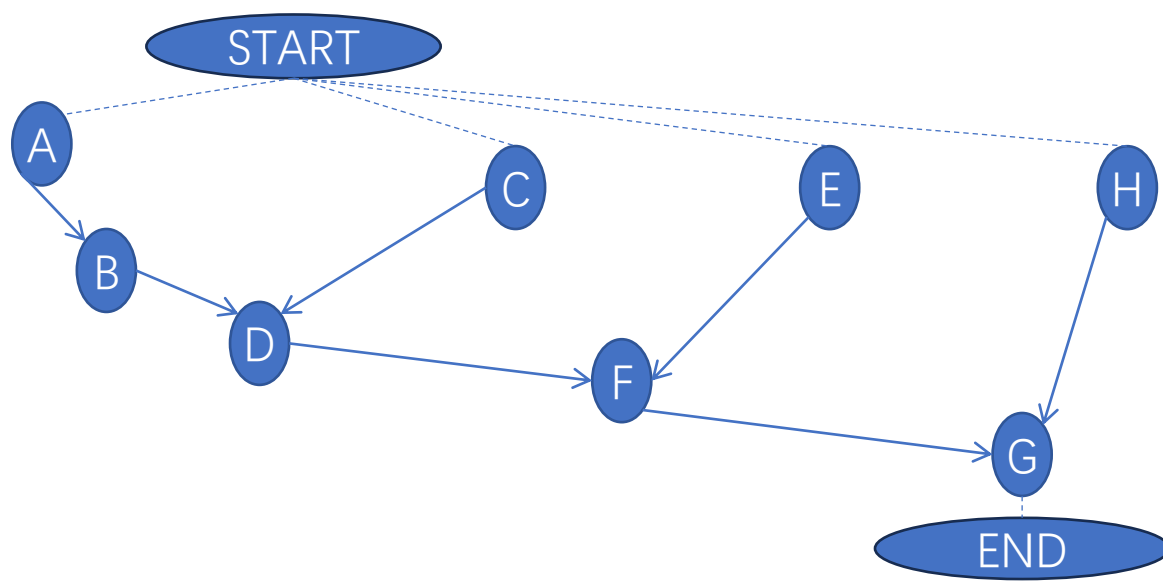
D 3

E 1

F

```

1 ASAP (  $G_s(V,E)$  ) {
2   Schedule  $v_0$  by setting  $t_0 = 0$ ;
3   repeat {
4     Select a vertex  $v_i$  whose predecessors are all scheduled;
5     Schedule  $v_i$  by setting  $t_i = \max t_j + d_j$ ; }
7   until ( $v_n$  is scheduled);
8   return (  $G$  ); }
  
```



START 0

A 1

B 2

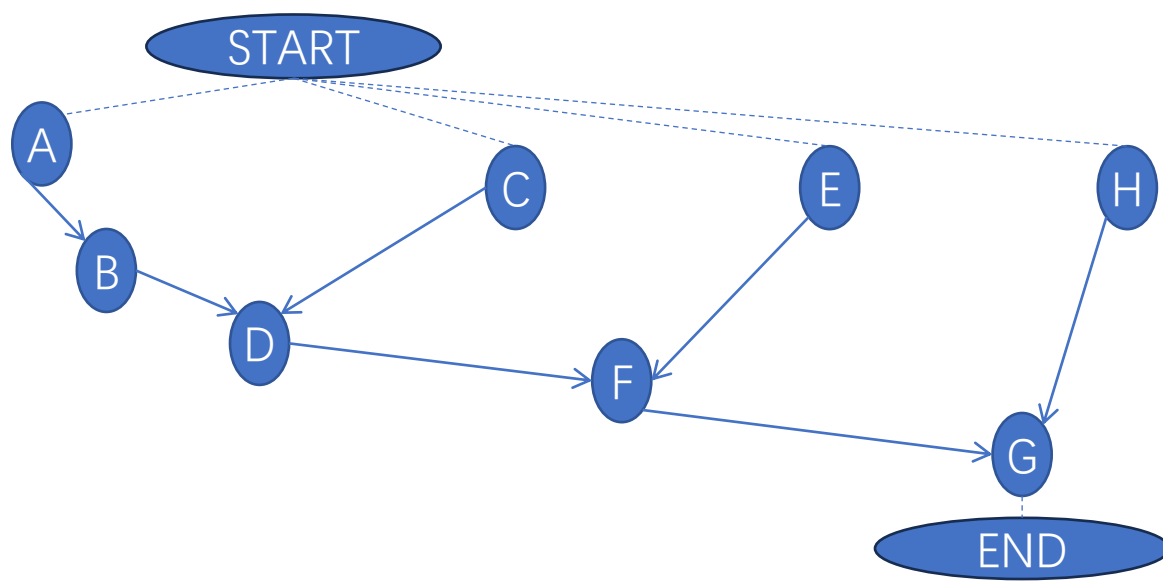
C 1

D 3

E 1

F 4

```
1 ASAP (  $G_s(V,E)$  ) {  
2     Schedule  $v_0$  by setting  $t_0 = 0$ ;  
3     repeat {  
4         Select a vertex  $v_i$  whose predecessors are all scheduled;  
5         Schedule  $v_i$  by setting  $t_i = \max t_j + d_j$ ;}  
7     until ( $v_n$  is scheduled);  
8     return (  $G$  );}
```



START 0

A 1

B 2

C 1

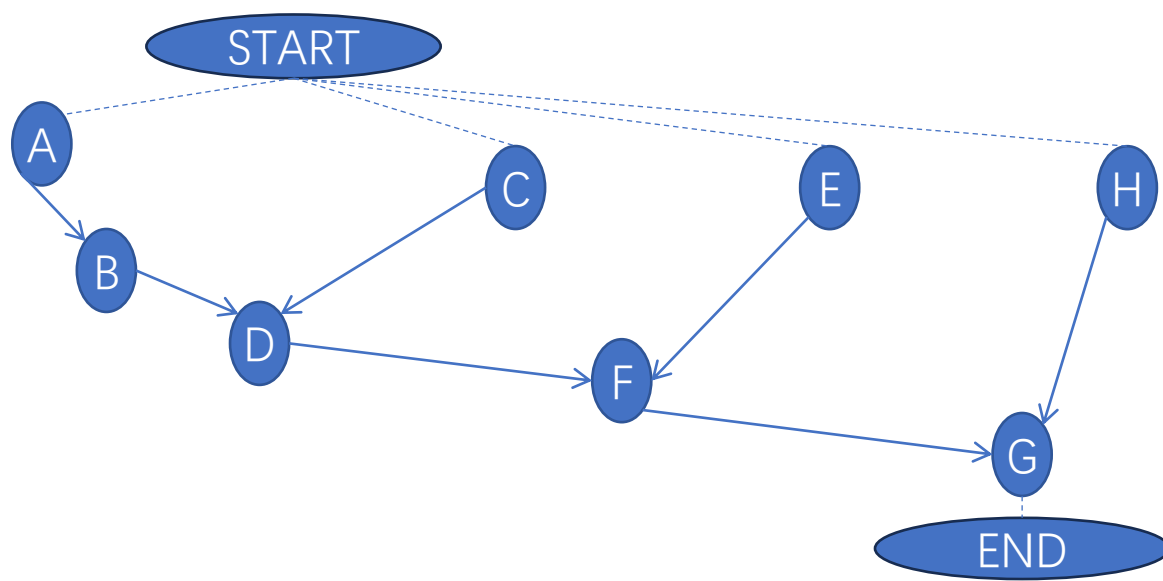
D 3

E 1

F 4

```

1 ASAP (  $G_s(V,E)$  ) {
2   Schedule  $v_0$  by setting  $t_0 = 0$ ;
3   repeat {
4     Select a vertex  $v_i$  whose predecessors are all scheduled;
5     Schedule  $v_i$  by setting  $t_i = \max t_j + d_j$ ;
7   until ( $v_n$  is scheduled);
8   return (  $G$  );}
  
```



START 0

A 1

B 2

C 1

D 3

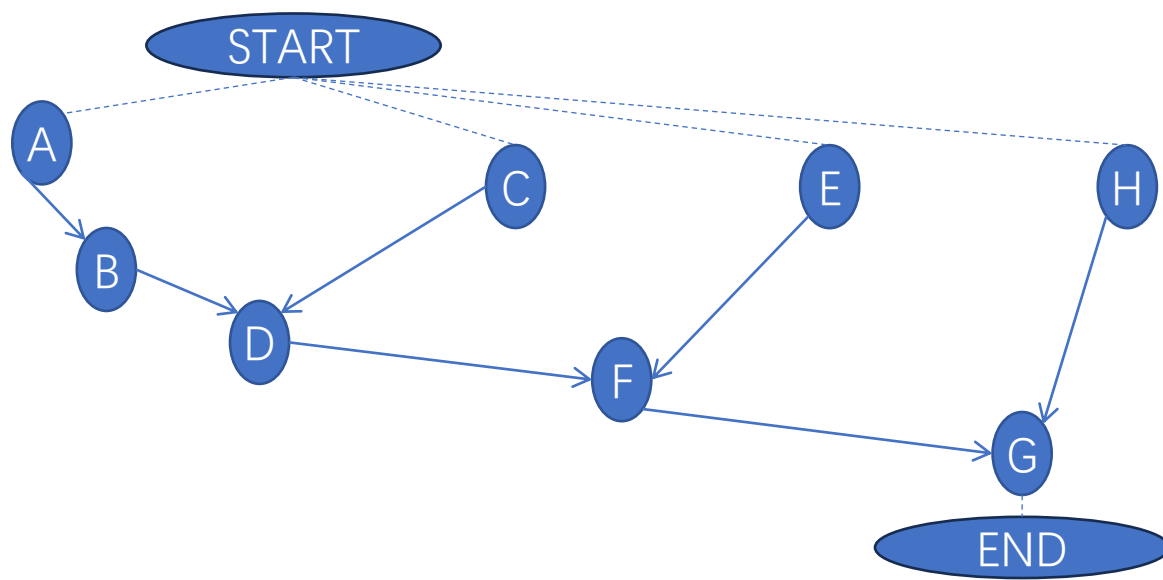
E 1

F 4

H

```

1 ASAP (  $G_s(V,E)$  ) {
2     Schedule  $v_0$  by setting  $t_0 = 0$ ;
3     repeat {
4         Select a vertex  $v_i$  whose predecessors are all scheduled;
5         Schedule  $v_i$  by setting  $t_i = \max t_j + d_j$ ; }
7     until ( $v_n$  is scheduled);
8     return (  $G$  ); }
  
```



START 0

A 1

B 2

C 1

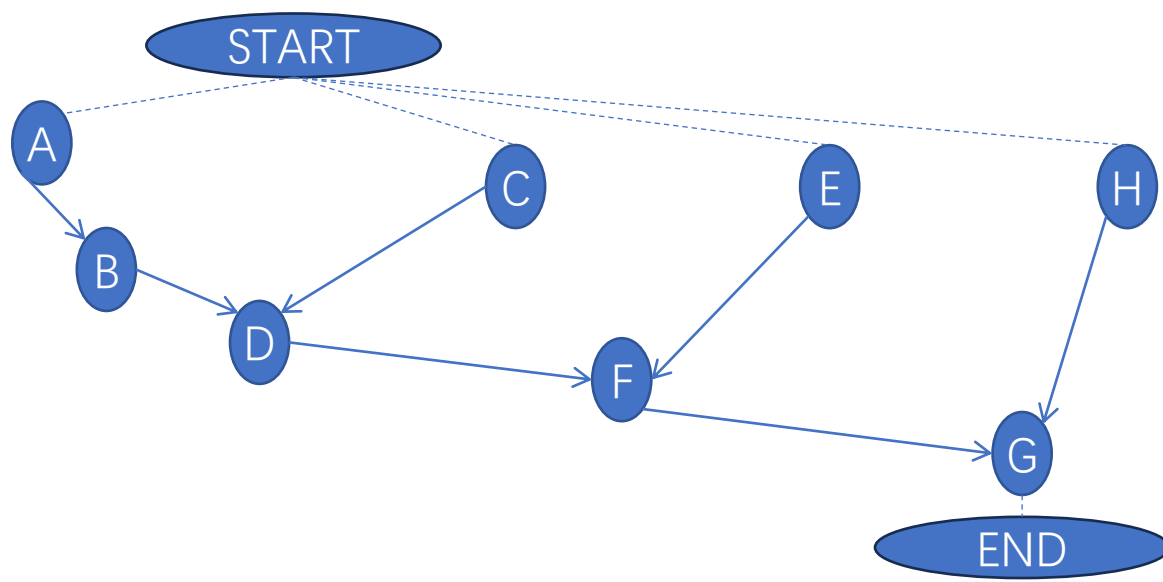
D 3

E 1

F 4

H 1

```
1 ASAP (  $G_s(V,E)$  ) {  
2     Schedule  $v_0$  by setting  $t_0 = 0$ ;  
3     repeat {  
4         Select a vertex  $v_i$  whose predecessors are all scheduled;  
5         Schedule  $v_i$  by setting  $t_i = \max t_j + d_j$ ;}  
7     until ( $v_n$  is scheduled);  
8     return (  $G$  );}
```



START 0

A 1

B 2

C 1

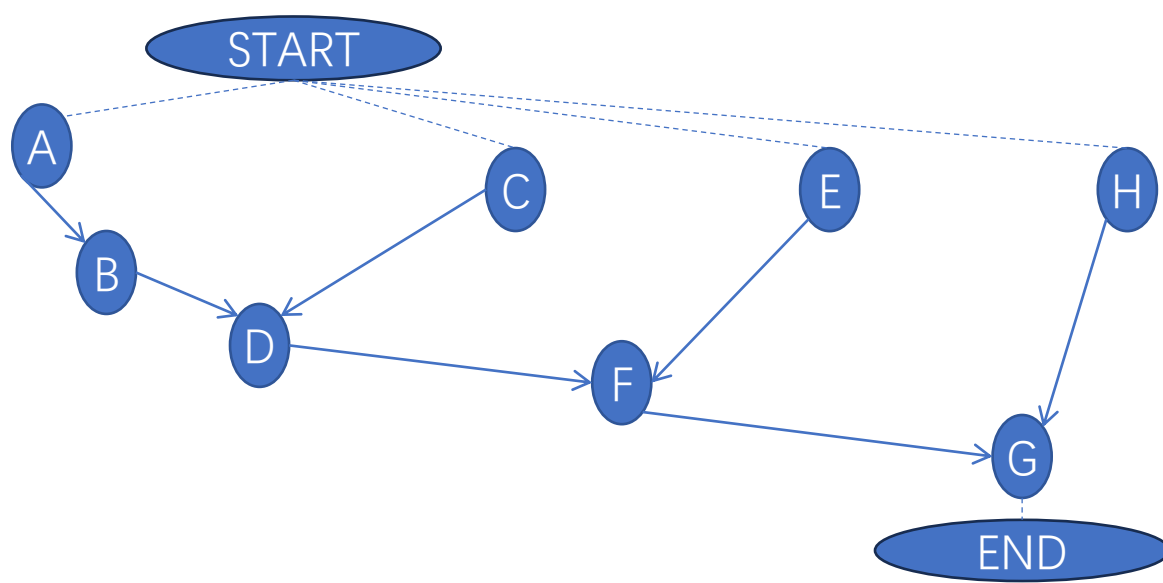
D 3

E 1

F 4

H 1

```
1 ASAP (  $G_s(V,E)$  ) {  
2     Schedule  $v_0$  by setting  $t_0 = 0$ ;  
3     repeat {  
4         Select a vertex  $v_i$  whose predecessors are all scheduled;  
5         Schedule  $v_i$  by setting  $t_i = \max t_j + d_j$ ;}  
7     until ( $v_n$  is scheduled);  
8     return (  $G$  );}
```

START 0

A 1

B 2

C 1

D 3

E 1

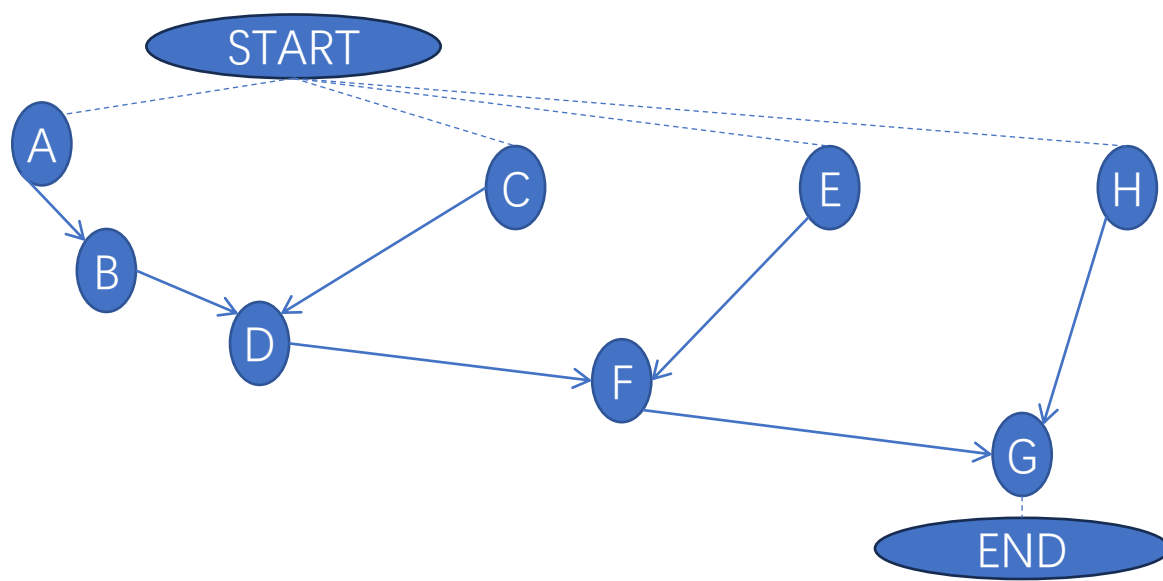
F 4

H 1

G

```

1 ASAP (  $G_s(V,E)$  ) {
2     Schedule  $v_0$  by setting  $t_0 = 0$ ;
3     repeat {
4         Select a vertex  $v_i$  whose predecessors are all scheduled;
5         Schedule  $v_i$  by setting  $t_i = \max t_j + d_j$ ; }
7     until ( $v_n$  is scheduled);
8     return (  $G$  ); }
  
```



START 0

A 1

B 2

C 1

D 3

E 1

F 4

H 1

G 5

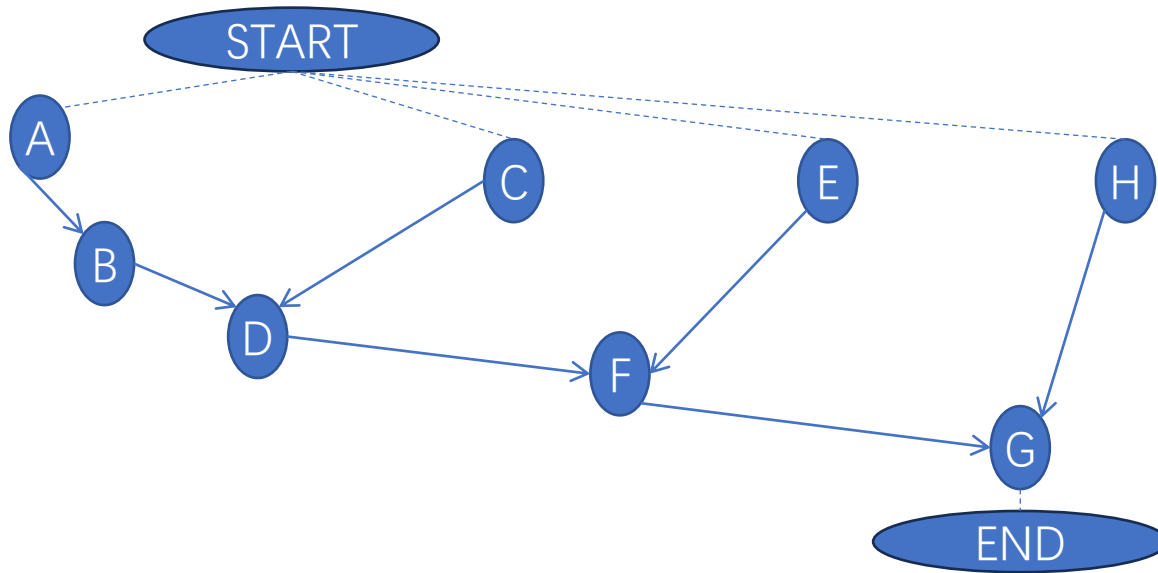
END 6

```

1 ASAP (  $G_s(V,E)$  ) {
2   Schedule  $v_0$  by setting  $t_0 = 0$ ;
3   repeat {
4     Select a vertex  $v_i$  whose predecessors are all scheduled;
5     Schedule  $v_i$  by setting  $t_i = \max t_j + d_j$ ;
7   until ( $v_n$  is scheduled);
8   return (  $G$  );}
  
```

ASAP调度算法

ASAP scheduling algorithm



| | |
|---------|---------|
| 10 | START A |
| START 1 | START C |
| A 1 | START E |
| B 1 | START G |
| C 1 | A B |
| D 1 | C D |
| E 1 | B D |
| F 1 | E F |
| G 1 | D F |
| H 1 | H G |
| END 1 | F G |
| | G END |



ASAP调度算法

ASAP scheduling algorithm

```
class Node:
    def __init__(self, name, duration):
        self.name = name
        self.duration = duration
        self.is_scheduled = False
        self.start_time = -1
        self.predecessors = []
```

```
node=Node(START, 1)
```

| | |
|--------------|-------|
| name | START |
| duration | 1 |
| Is_scheduled | False |
| start_time | -1 |
| predecessors | [] |

node

| | |
|---------|---------|
| 10 | START A |
| START 1 | START C |
| A 1 | START E |
| B 1 | START G |
| C 1 | A B |
| D 1 | C D |
| E 1 | B D |
| F 1 | E F |
| G 1 | D F |
| H 1 | H G |
| END 1 | F G |
| | G END |



ASAP调度算法

ASAP scheduling algorithm

```
def read_graph_from_file(input_file):  
    with open(input_file, "r", encoding="utf-8") as f:  
        N = int(f.readline().strip()) # 读取节点数  
        nodes = []  
        # 读取节点信息  
        for _ in range(N):  
            parts = f.readline().strip().split()  
            node_name = parts[0]  
            duration = int(parts[1])  
            nodes.append(Node(node_name, duration))
```

| | |
|---------|---------|
| 10 | START A |
| START 1 | START C |
| A 1 | START E |
| B 1 | START G |
| C 1 | A B |
| D 1 | C D |
| E 1 | B D |
| F 1 | E F |
| G 1 | D F |
| H 1 | H G |
| END 1 | F G |
| | G END |



ASAP调度算法

ASAP scheduling algorithm

```
def read_graph_from_file(input_file):
    with open(input_file, "r", encoding="utf-8") as f:
        N = int(f.readline().strip()) # 读取节点数
        nodes = []
        # 读取节点信息
        for _ in range(N):
            parts = f.readline().strip().split()
            node_name = parts[0]
            duration = int(parts[1])
            nodes.append(Node(node_name, duration))
```

| | | |
|-----|--------------|-------|
| 0 | name | START |
| | duration | 1 |
| | Is_scheduled | False |
| | start_time | -1 |
| | predecessors | [] |
| 1 | name | A |
| | duration | 1 |
| | Is_scheduled | False |
| | start_time | -1 |
| | predecessors | [] |
| 2 | ... | |
| ... | | |

nodes



ASAP调度算法

```
def read_graph_from_file(input_file):  
    with open(input_file, "r", encoding="utf-8") as f:  
        N = int(f.readline().strip()) # 读取节点数  
        nodes = []  
        edges = []  
        # 读取节点信息 ...  
        # 读取边信息并建立前驱关系  
        for line in f:  
            src, dst = line.strip().split()  
            src_node = None  
            dst_node = None  
            for node in nodes:  
                if node.name == src:  
                    src_node = node  
                if node.name == dst:  
                    dst_node = node  
            if dst_node and src_node:  
                dst_node.predecessors.append(src_node)
```

| | |
|---------|---------|
| 10 | START A |
| START 1 | START C |
| A 1 | START E |
| B 1 | START G |
| C 1 | A B |
| D 1 | C D |
| E 1 | B D |
| F 1 | E F |
| G 1 | D F |
| H 1 | H G |
| END 1 | F G |
| | G END |

ASAP调度算法

```
def read_graph_from_file(input_file):
    with open(input_file, "r", encoding="utf-8") as f:
        N = int(f.readline().strip()) # 读取节点数
        nodes = []
        # 读取节点信息 ...
        # 读取边信息并建立前驱关系
        for line in f:
            src, dst = line.strip().split()
            src_node = None
            dst_node = None
            for node in nodes:
                if node.name == src:
                    src_node = node
                if node.name == dst:
                    dst_node = node
            if dst_node and src_node:
                dst_node.predecessors.append(src_node)
```

| | | |
|-----|--------------|------------|
| 0 | name | START |
| | duration | 1 |
| | Is_scheduled | False |
| | start_time | -1 |
| | predecessors | [] |
| 1 | name | A |
| | duration | 1 |
| | Is_scheduled | False |
| | start_time | -1 |
| | predecessors | [nodes[0]] |
| 2 | ... | |
| ... | | |

nodes



ASAP调度算法

ASAP scheduling algorithm

```
1 ASAP (  $G_s(V,E)$  ) {  
2     Schedule  $v_0$  by setting  $t_0 = 0$ ;  
3     repeat {  
4         Select a vertex  $v_i$  whose predecessors are all scheduled;  
5         Schedule  $v_i$  by setting  $t_i = \max t_j + d_j$ ;  
6     }  
7     until ( $v_n$  is scheduled);  
8     return (  $G$  );  
9 }
```



ASAP调度算法

ASAP scheduling algorithm

```
def asap_algorithm(nodes):  
    # 1. 令 `START` 先调度  
    start_node = next(node for node in nodes if node.name == "START")  
    start_node.is_scheduled = True  
    start_node.start_time = 0
```



ASAP调度算法

ASAP scheduling algorithm

```
1 ASAP (  $G_s(V,E)$  ) {  
2     Schedule  $v_0$  by setting  $t_0 = 0$ ;  
3     repeat {  
4         Select a vertex  $v_i$  whose predecessors are all scheduled;  
5         Schedule  $v_i$  by setting  $t_i = \max t_j + d_j$ ;  
6     }  
7     until ( $v_n$  is scheduled);  
8     return (  $G$  );  
9 }
```

```
def asap_algorithm(nodes):  
    # 1. 令 `START` 先调度  
    start_node = next(node for node in nodes if node.name == "START")  
    start_node.is_scheduled = True  
    start_node.start_time = 0  
    # 2. 直到 `END` 也被调度  
    end_node = next(node for node in nodes if node.name == "END")  
    while not end_node.is_scheduled:  
        # 找到一个满足条件的节点  
        node_to_schedule = find_suitable_node(nodes)  
        # 决定该节点的开始时间  
        schedule_node(node_to_schedule)
```



ASAP调度算法

ASAP scheduling algorithm

```
1 ASAP (  $G_s(V,E)$  ) {  
2     Schedule  $v_0$  by setting  $t_0 = 0$ ;  
3     repeat {  
4         Select a vertex  $v_i$  whose predecessors are all scheduled;  
5         Schedule  $v_i$  by setting  $t_i = \max t_j + d_j$ ;  
6     }  
7     until ( $v_n$  is scheduled);  
8     return (  $G$  );  
9 }
```

```
def find_suitable_node(nodes):  
    for node in nodes:  
        # 只对当前还没被调度的结点进行检查  
        if not node.is_scheduled:  
            # 检查所有前置结点是否已调度  
            all_predecessors_scheduled = True  
            for p in node.predecessors:  
                if not p.is_scheduled:  
                    all_predecessors_scheduled = False  
                    break  
            if all_predecessors_scheduled:  
                return node
```



ASAP调度算法

ASAP scheduling algorithm

```
1 ASAP (  $G_s(V,E)$  ) {  
2     Schedule  $v_0$  by setting  $t_0 = 0$ ;  
3     repeat {  
4         Select a vertex  $v_i$  whose predecessors are all scheduled;  
5         Schedule  $v_i$  by setting  $t_i = \max t_j + d_j$ ;  
6     }  
7     until ( $v_n$  is scheduled);  
8     return (  $G$  );  
9 }
```

```
def schedule_node(node):  
    # 统计所有前置结点的结束时间  
    pred_finish_times = []  
    for pred in node.predecessors:  
        finish_time = pred.start_time + pred.duration  
        pred_finish_times.append(finish_time)  
    # 设定当前节点的开始时间  
    node.start_time = max(pred_finish_times)  
    node.is_scheduled = True
```



```
1 ASAP (  $G_s(V,E)$  ) {  
2     Schedule  $v_0$  by setting  $t_0 = 0$ ;  
3     repeat {  
4         Select a vertex  $v_i$  whose predecessors are all scheduled;  
5         Schedule  $v_i$  by setting  $t_i = \max t_j + d_j$ ;  
6     }  
7     until ( $v_n$  is scheduled);  
8     return (  $G$  );  
9 }
```

```
def asap_algorithm(nodes):  
    start_node = next(node for node in nodes if node.name == "START")  
    start_node.is_scheduled = True  
    start_node.start_time = 0  
    end_node = next(node for node in nodes if node.name == "END")  
    while not end_node.is_scheduled:  
        node_to_schedule = find_suitable_node(nodes)  
        schedule_node(node_to_schedule, nodes)  
    return nodes
```

```
def write_schedule_to_file(nodes, output_file):  
    with open(output_file, 'w', encoding="utf-8") as f:  
        for node in nodes:  
            f.write(f"{node.name} {node.start_time}\n")
```

```
input_file = "input.txt"  
output_file = "output.txt"  
nodes = read_graph_from_file(input_file)  
nodes = asap_algorithm(nodes)  
write_schedule_to_file(nodes, output_file)
```

START 0

A 1

B 2

C 1

D 3

E 1

F 4

H 1

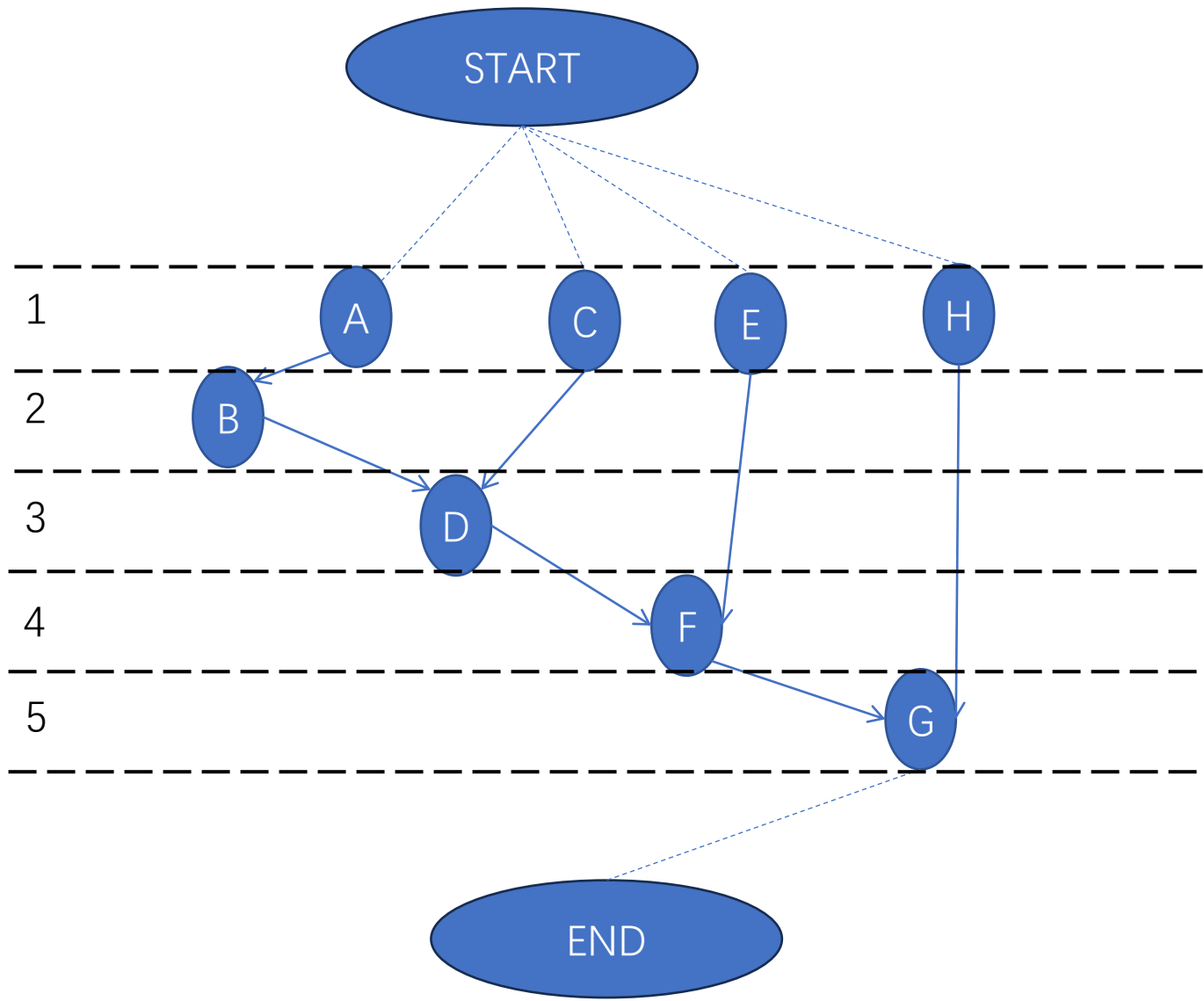
G 5

END 6

输入文件内容:

10
START 1
A 1
B 1
C 1
D 1
E 1
F 1
G 1
H 1
END 0
START B
START A
START C
START E
START G
A B
C D
B D
E F
D F
H G
F G
B END
D END
F END
G END

ASAP的结果



为顶点决定开始周期的顺序:
START->
A->B->C->D->E->F->H->G
->END

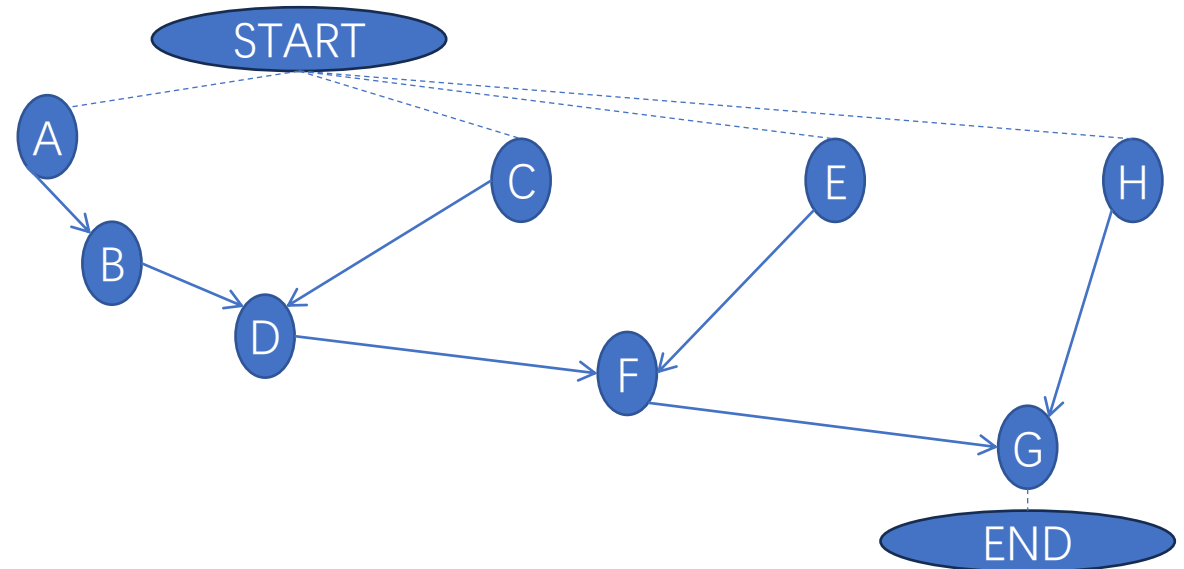
输出文件内容:
START 0
A 1
B 2
C 1
D 3
E 1
F 4
H 1
G 5
END 6

ALAP调度算法

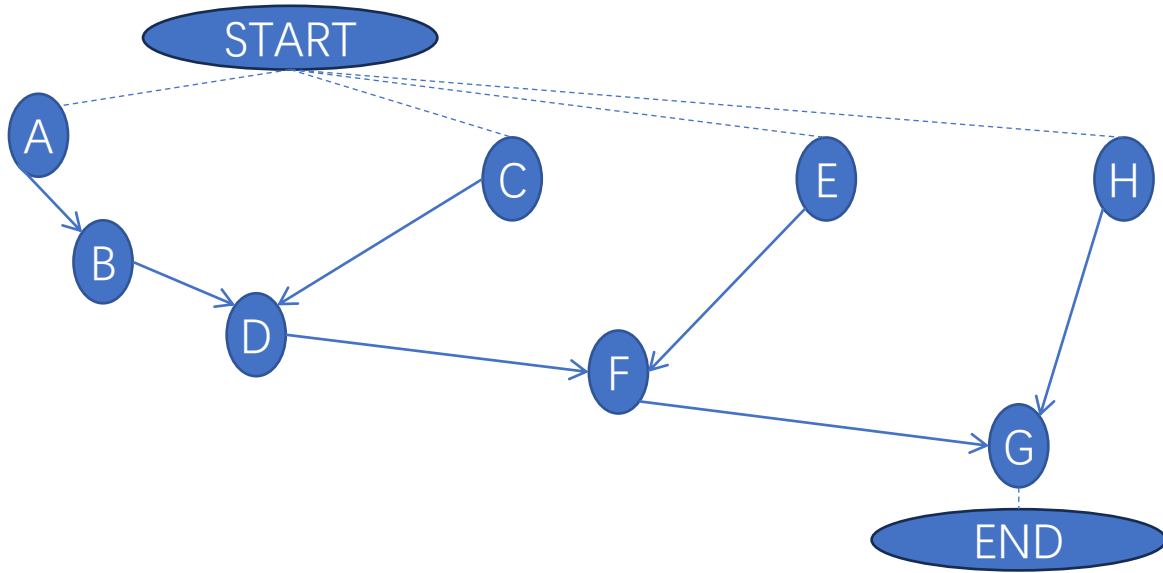
ALAP调度算法（最晚时间约束）

ALAP scheduling algorithm (Latency Constrained)

```
1 ALAP (  $G_s(V,E)$ ,  $x$  ) {  
2     Schedule  $v_n$  by setting  $t_n = x+1$  ;  
3     repeat {  
4         Select a vertex  $v_i$  whose successors are all scheduled;  
5         Schedule  $v_i$  by setting  $t_i = \min t_j - d_i$ ;  
6     }  
7     until ( $v_0$  is scheduled);  
8     return ( $t$ );  
9 }
```

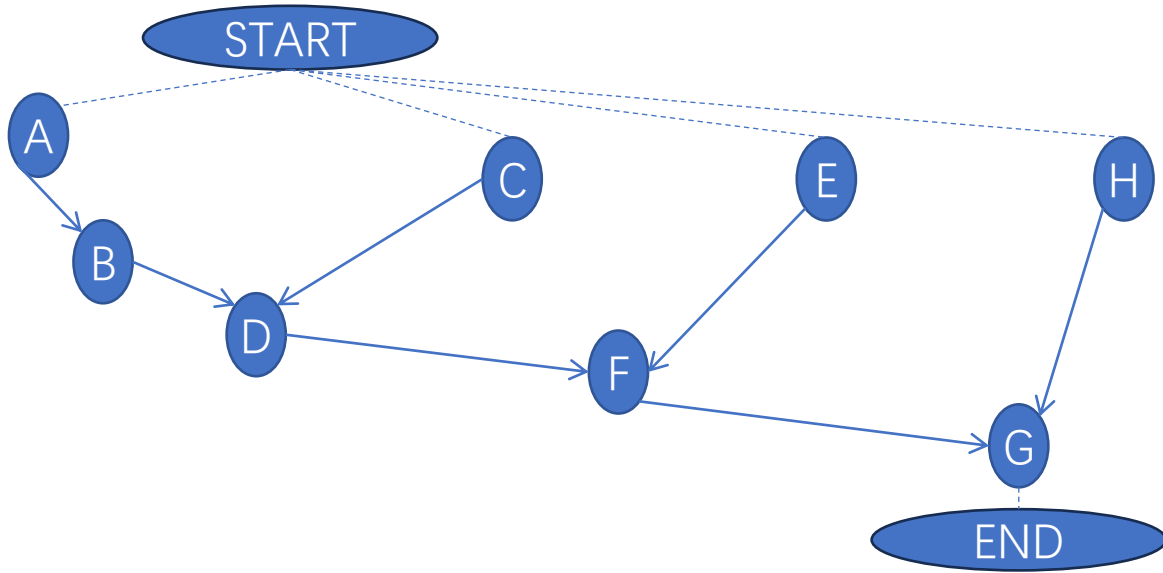


END 6

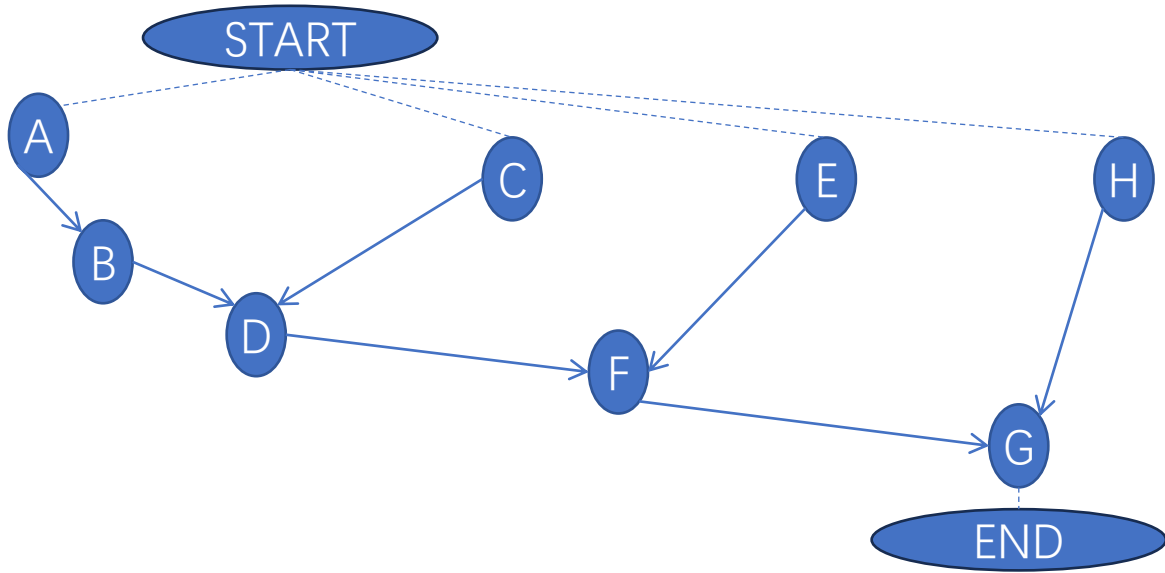


```
1 ALAP (  $G_s(V,E)$ ,  $x$  ) {  
2     Schedule  $v_n$  by setting  $t_n = x + 1$ ;  
3     repeat {  
4         Select a vertex  $v_i$  whose successors are all scheduled;  
5         Schedule  $v_i$  by setting  $t_i = \min t_j - d_i$ ;  
7     until ( $v_0$  is scheduled);  
8     return ( $t$ );}
```

END 6

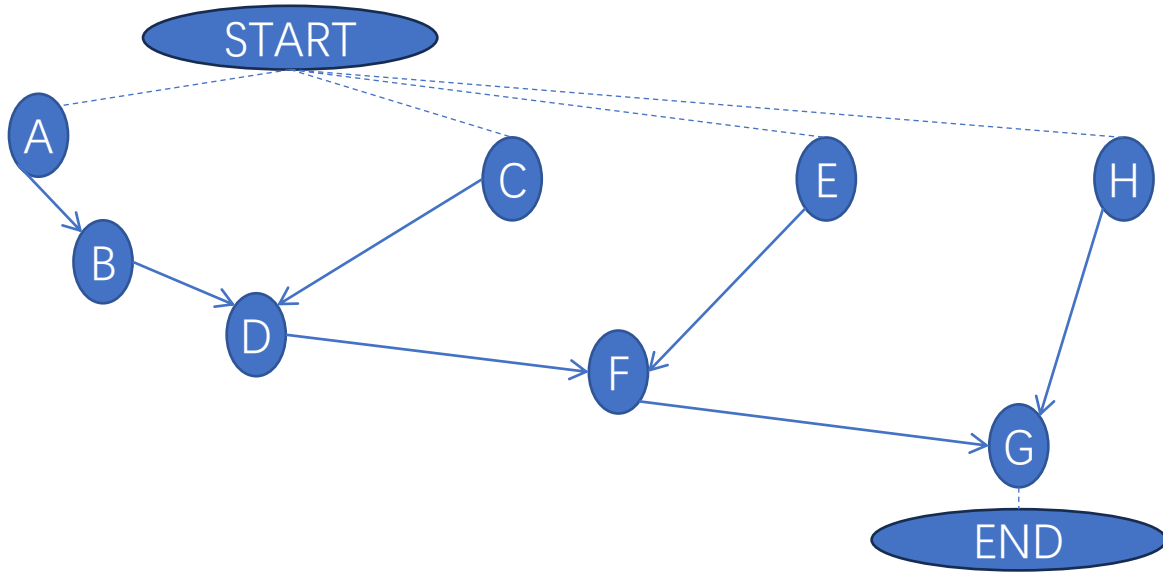


```
1 ALAP (  $G_s(V,E)$ ,  $x$  ) {  
2     Schedule  $v_n$  by setting  $t_n = x + 1$ ;  
3     repeat {  
4         Select a vertex  $v_i$  whose successors are all scheduled;  
5         Schedule  $v_i$  by setting  $t_i = \min t_j - d_i$ ;  
7     until ( $v_0$  is scheduled);  
8     return ( $t$ );}
```

END 6
G

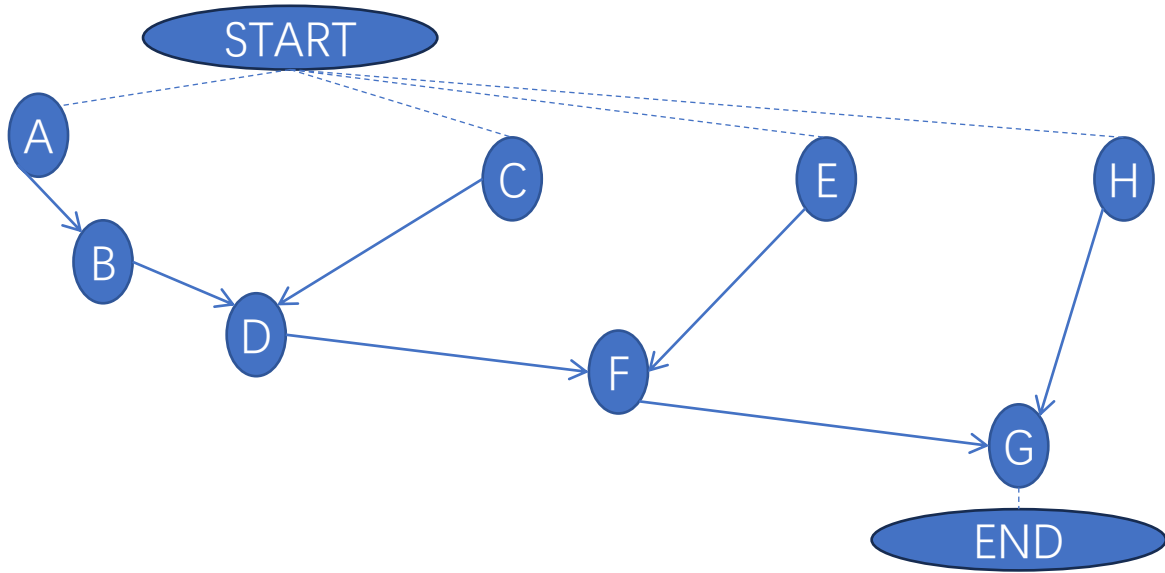
```
1 ALAP (  $G_s(V,E)$ ,  $x$  ) {  
2     Schedule  $v_n$  by setting  $t_n = x + 1$ ;  
3     repeat {  
4         Select a vertex  $v_i$  whose successors are all scheduled;  
5         Schedule  $v_i$  by setting  $t_i = \min t_j - d_i$ ;}  
7     until ( $v_0$  is scheduled);  
8     return ( $t$ );}
```



END 6
G 5

```

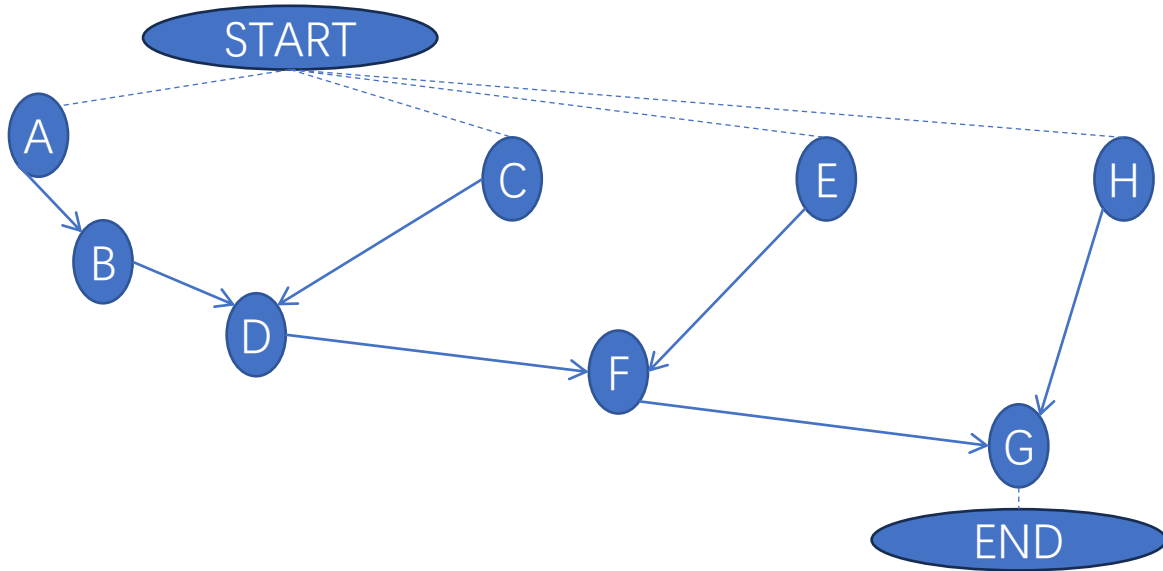
1 ALAP (  $G_s(V,E), x$  ) {
2     Schedule  $v_n$  by setting  $t_n = x + 1$ ;
3     repeat {
4         Select a vertex  $v_i$  whose successors are all scheduled;
5         Schedule  $v_i$  by setting  $t_i = \min t_j - d_i$ ;
7     until ( $v_0$  is scheduled);
8     return ( $t$ );}
  
```



END 6
G 5
F

```

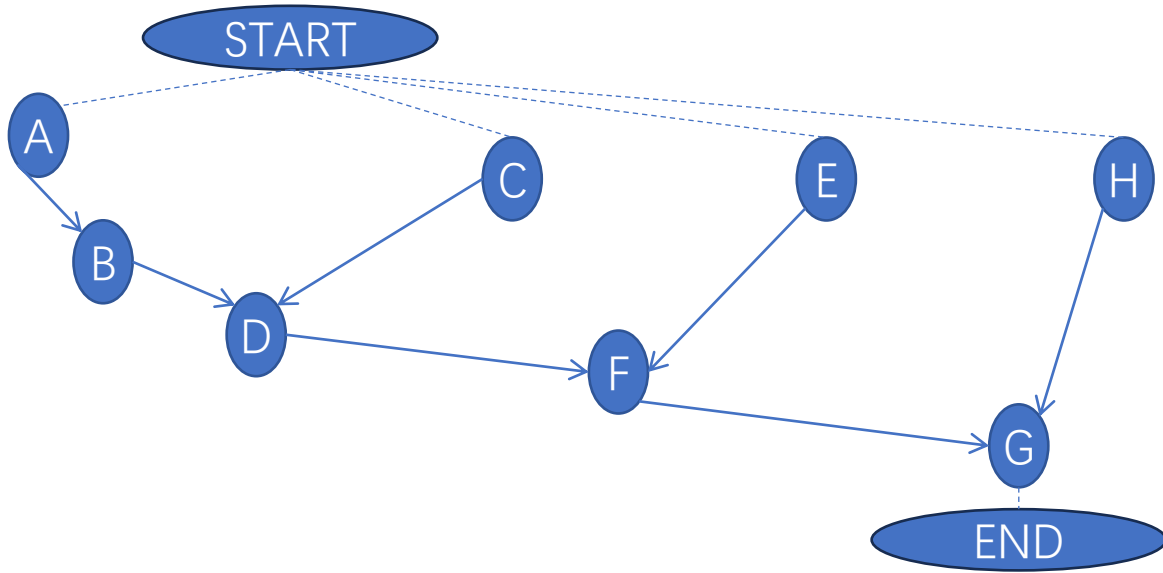
1 ALAP (  $G_s(V,E)$ ,  $x$  ) {
2     Schedule  $v_n$  by setting  $t_n = x + 1$ ;
3     repeat {
4         Select a vertex  $v_i$  whose successors are all scheduled;
5         Schedule  $v_i$  by setting  $t_i = \min t_j - d_i$ ;
7     until ( $v_0$  is scheduled);
8     return ( $t$ );}
  
```



END 6
G 5
F 4

```

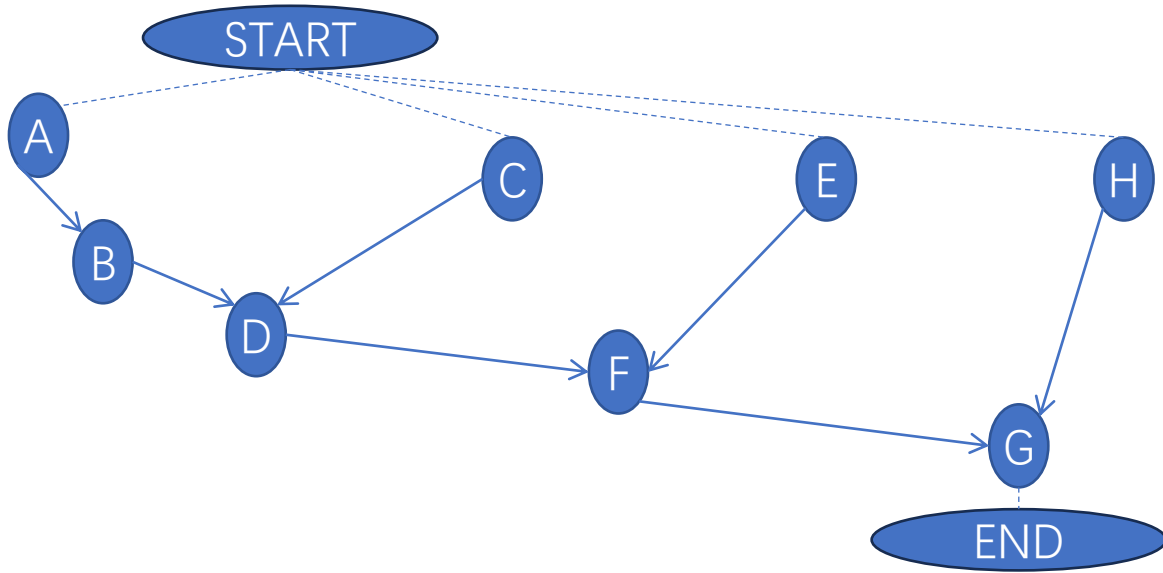
1 ALAP (  $G_s(V,E)$ ,  $x$  ) {
2     Schedule  $v_n$  by setting  $t_n = x + 1$ ;
3     repeat {
4         Select a vertex  $v_i$  whose successors are all scheduled;
5         Schedule  $v_i$  by setting  $t_i = \min t_j - d_i$ ;
7     until ( $v_0$  is scheduled);
8     return ( $t$ );}
  
```



END 6
G 5
F 4
D

```

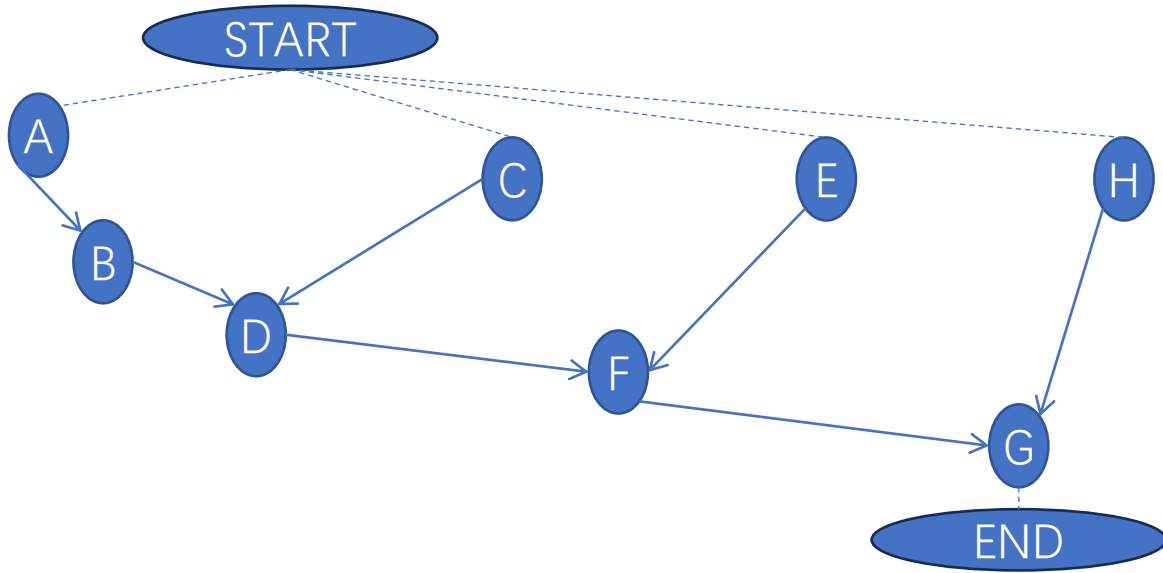
1 ALAP (  $G_s(V,E)$ ,  $x$  ) {
2     Schedule  $v_n$  by setting  $t_n = x + 1$ ;
3     repeat {
4         Select a vertex  $v_i$  whose successors are all scheduled;
5         Schedule  $v_i$  by setting  $t_i = \min t_j - d_i$ ;
7     until ( $v_0$  is scheduled);
8     return ( $t$ );}
  
```



END 6
G 5
F 4
D 3

```

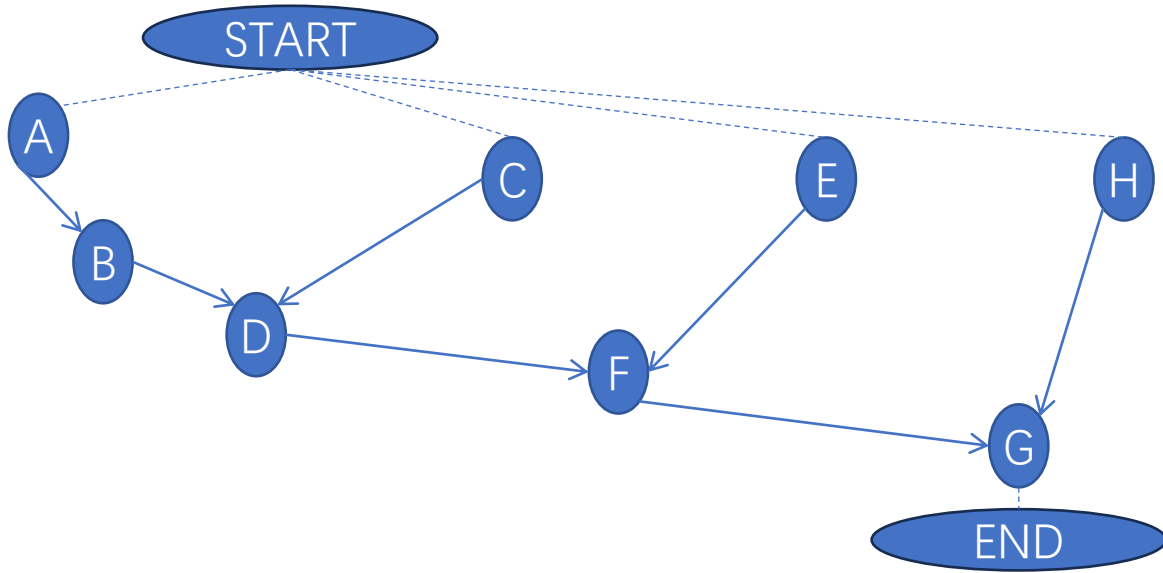
1 ALAP (  $G_s(V,E), x$  ) {
2     Schedule  $v_n$  by setting  $t_n = x + 1$ ;
3     repeat {
4         Select a vertex  $v_i$  whose successors are all scheduled;
5         Schedule  $v_i$  by setting  $t_i = \min t_j - d_i$ ;
7     until ( $v_0$  is scheduled);
8     return ( $t$ );}
  
```



END 6
G 5
F 4
D 3
B

```

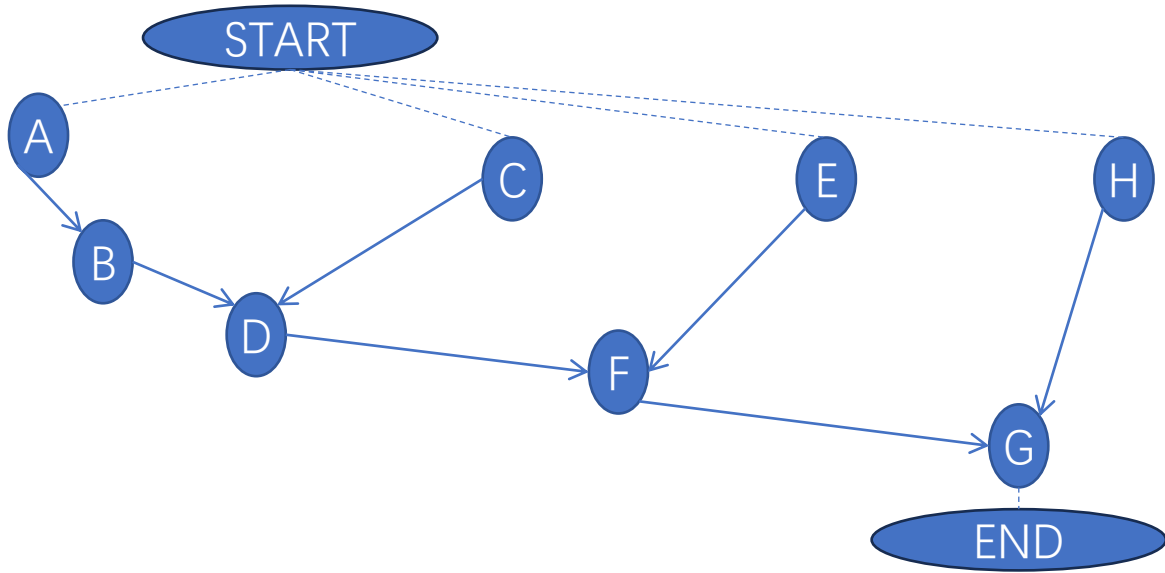
1 ALAP (  $G_s(V,E)$ ,  $x$  ) {
2     Schedule  $v_n$  by setting  $t_n = x + 1$ ;
3     repeat {
4         Select a vertex  $v_i$  whose successors are all scheduled;
5         Schedule  $v_i$  by setting  $t_i = \min t_j - d_i$ ;
7     until ( $v_0$  is scheduled);
8     return ( $t$ );}
  
```



END 6
G 5
F 4
D 3
B 2

```

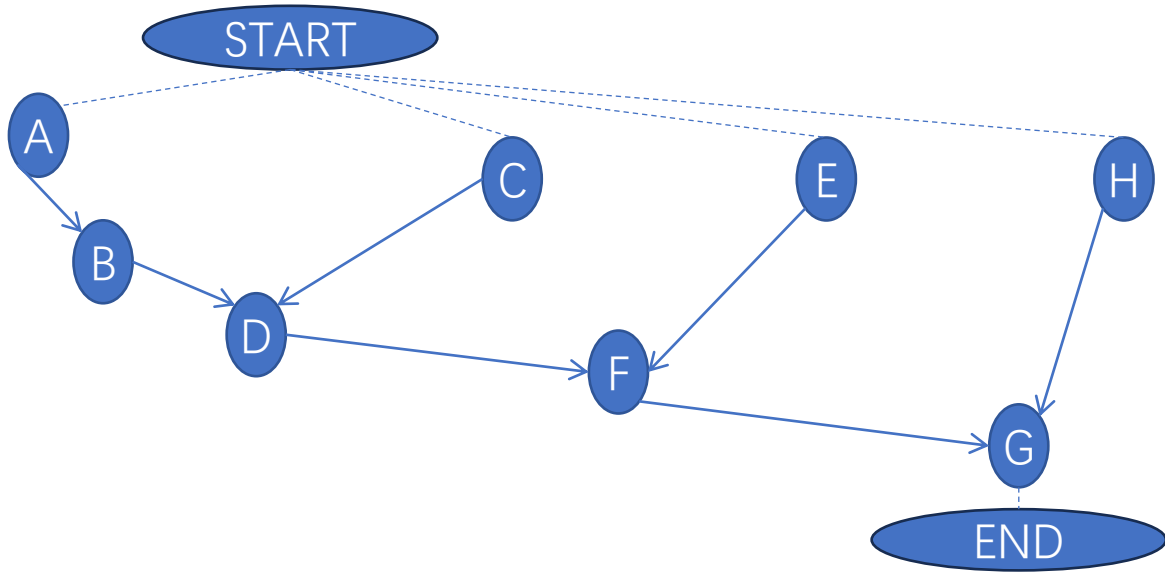
1 ALAP (  $G_s(V,E)$ ,  $x$  ) {
2     Schedule  $v_n$  by setting  $t_n = x + 1$ ;
3     repeat {
4         Select a vertex  $v_i$  whose successors are all scheduled;
5         Schedule  $v_i$  by setting  $t_i = \min t_j - d_i$ ;
7     until ( $v_0$  is scheduled);
8     return ( $t$ );}
  
```

END 6
 G 5
 F 4
 D 3
 B 2
 A

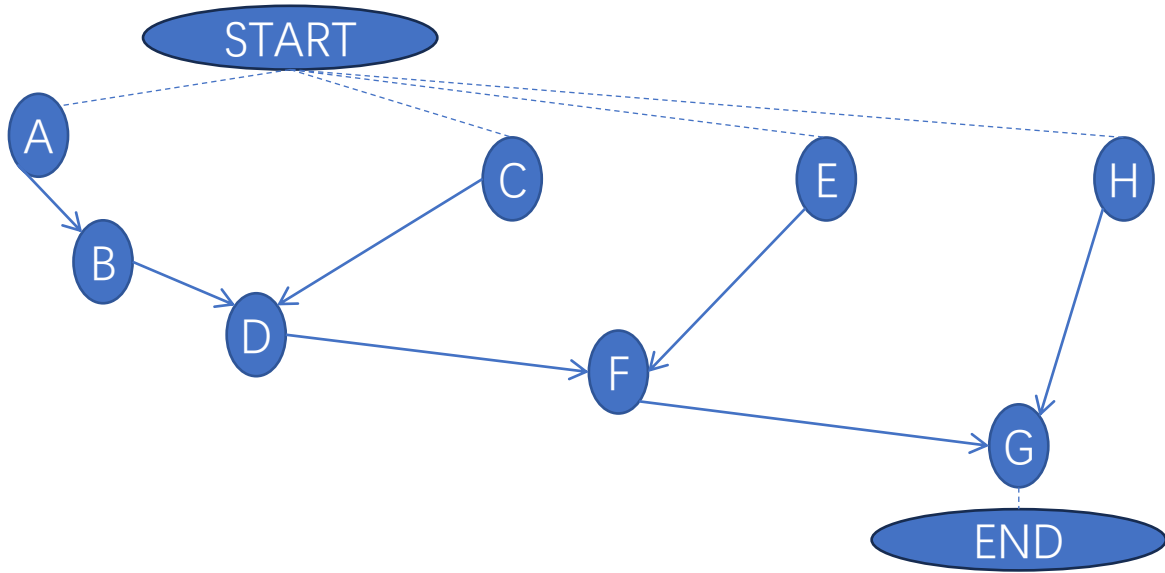
```

1 ALAP (  $G_s(V,E)$ ,  $x$  ) {
2     Schedule  $v_n$  by setting  $t_n = x + 1$ ;
3     repeat {
4         Select a vertex  $v_i$  whose successors are all scheduled;
5         Schedule  $v_i$  by setting  $t_i = \min t_j - d_i$ ;
6     }
7     until ( $v_0$  is scheduled);
8     return ( $t$ );}
  
```



END 6
G 5
F 4
D 3
B 2
A 1

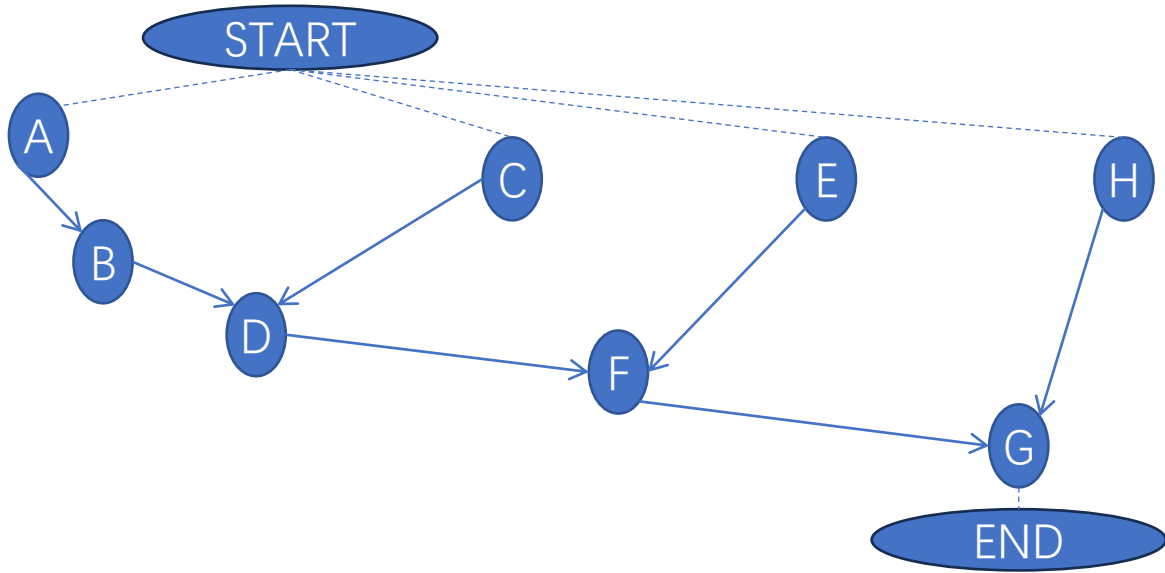
```
1 ALAP (  $G_s(V,E)$ ,  $x$  ) {  
2     Schedule  $v_n$  by setting  $t_n = x + 1$ ;  
3     repeat {  
4         Select a vertex  $v_i$  whose successors are all scheduled;  
5         Schedule  $v_i$  by setting  $t_i = \min t_j - d_i$ ;  
7     until ( $v_0$  is scheduled);  
8     return ( $t$ );}
```



END 6
 G 5
 F 4
 D 3
 B 2
 A 1
 C

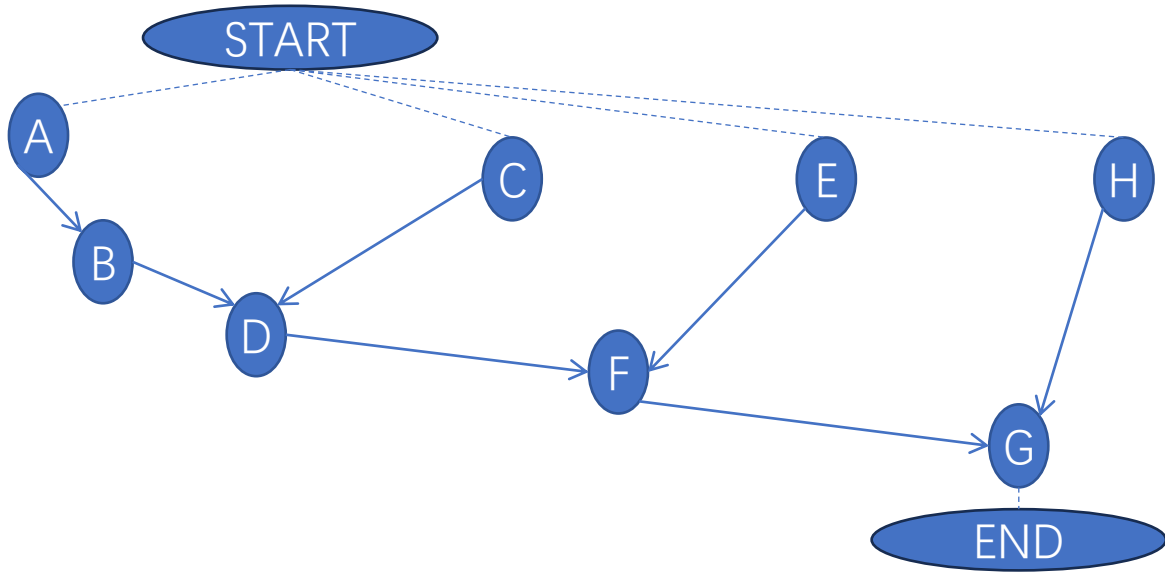
```

1 ALAP (  $G_s(V,E)$ ,  $x$  ) {
2     Schedule  $v_n$  by setting  $t_n = x + 1$ ;
3     repeat {
4         Select a vertex  $v_i$  whose successors are all scheduled;
5         Schedule  $v_i$  by setting  $t_i = \min t_j - d_i$ ;
7     until ( $v_0$  is scheduled);
8     return ( $t$ );}
  
```



END 6
G 5
F 4
D 3
B 2
A 1
C 2

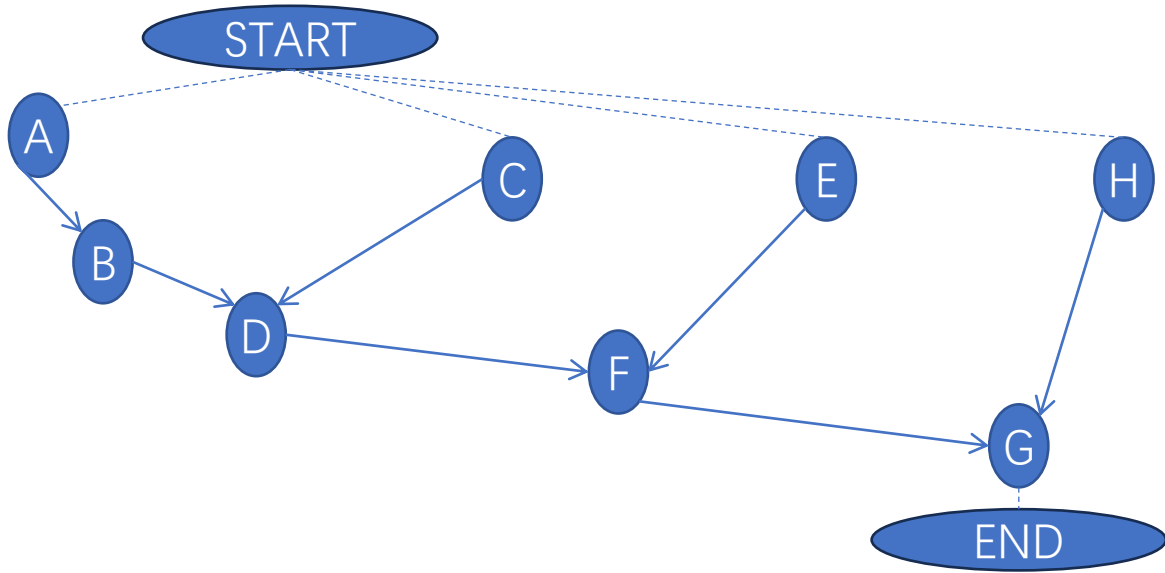
```
1 ALAP (  $G_s(V,E)$ ,  $x$  ) {  
2     Schedule  $v_n$  by setting  $t_n = x + 1$ ;  
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5         Schedule  $v_i$  by setting  $t_i = \min t_j - d_i$ ;  
7     until ( $v_0$  is scheduled);  
8     return ( $t$ );}
```



END 6
 G 5
 F 4
 D 3
 B 2
 A 1
 C 2
 E

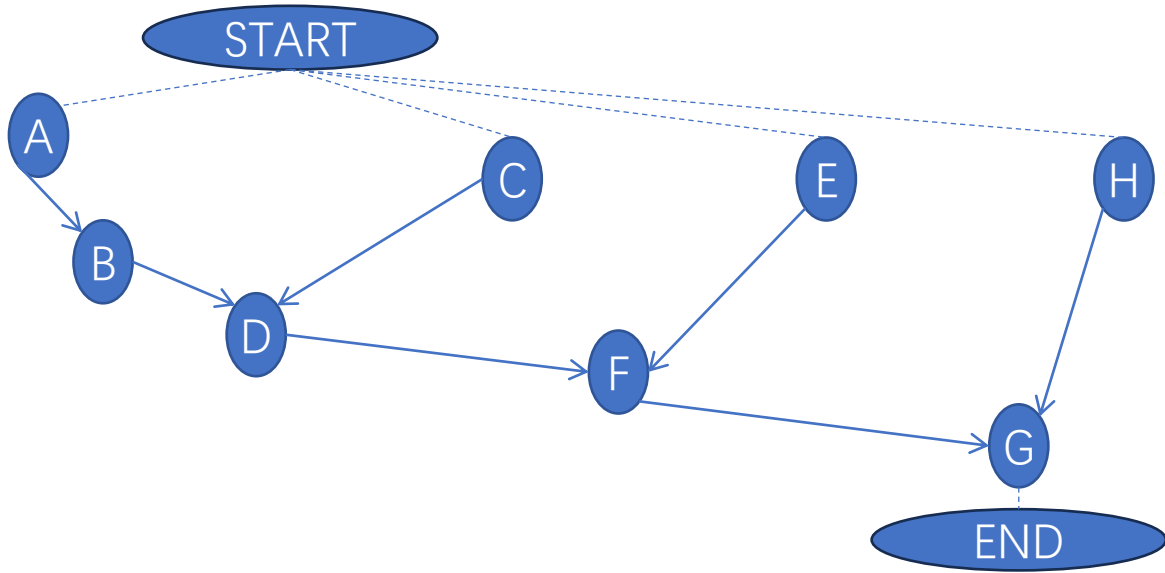
```

1 ALAP (  $G_s(V,E)$ ,  $x$  ) {
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3     repeat {
4         Select a vertex  $v_i$  whose successors are all scheduled;
5         Schedule  $v_i$  by setting  $t_i = \min t_j - d_i$ ;
7     until ( $v_0$  is scheduled);
8     return ( $t$ );}
  
```



END 6
G 5
F 4
D 3
B 2
A 1
C 2
E 3

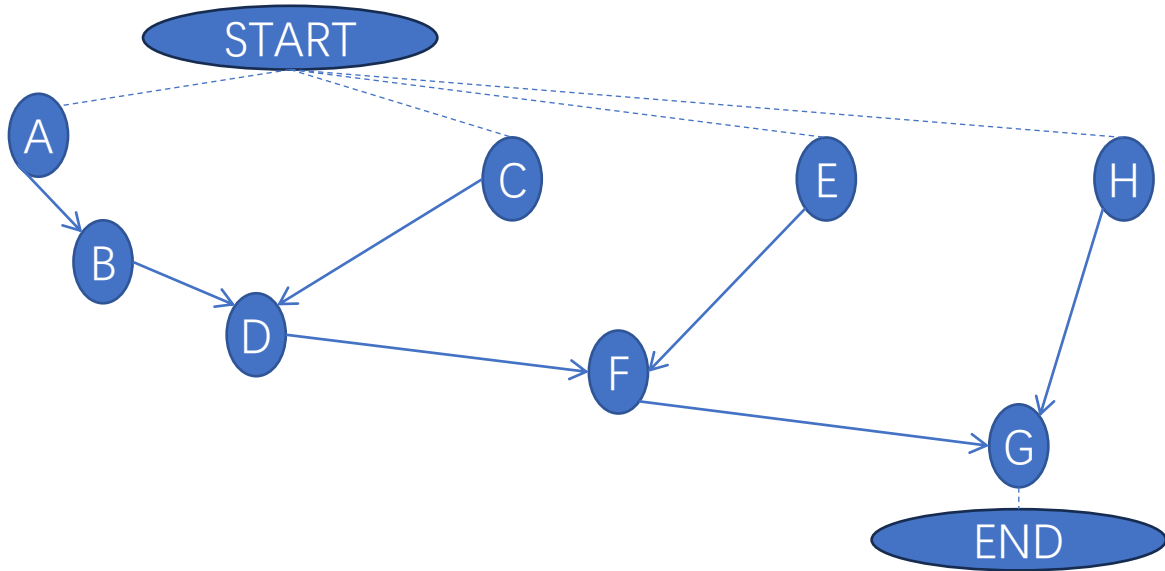
```
1 ALAP (  $G_s(V,E)$ ,  $x$  ) {  
2     Schedule  $v_n$  by setting  $t_n = x + 1$ ;  
3     repeat {  
4         Select a vertex  $v_i$  whose successors are all scheduled;  
5         Schedule  $v_i$  by setting  $t_i = \min t_j - d_i$ ;  
7     until ( $v_0$  is scheduled);  
8     return ( $t$ );}
```



END 6
 G 5
 F 4
 D 3
 B 2
 A 1
 C 2
 E 3
 H

```

1 ALAP (  $G_s(V,E)$ ,  $x$  ) {
2     Schedule  $v_n$  by setting  $t_n = x + 1$ ;
3     repeat {
4         Select a vertex  $v_i$  whose successors are all scheduled;
5         Schedule  $v_i$  by setting  $t_i = \min t_j - d_i$ ;
7     until ( $v_0$  is scheduled);
8     return ( $t$ );}
  
```



END 6

G 5

F 4

D 3

B 2

A 1

C 2

E 3

H 4

1 ALAP ($G_s(V,E), x$) {

2 Schedule v_n by setting $t_n = x + 1$;

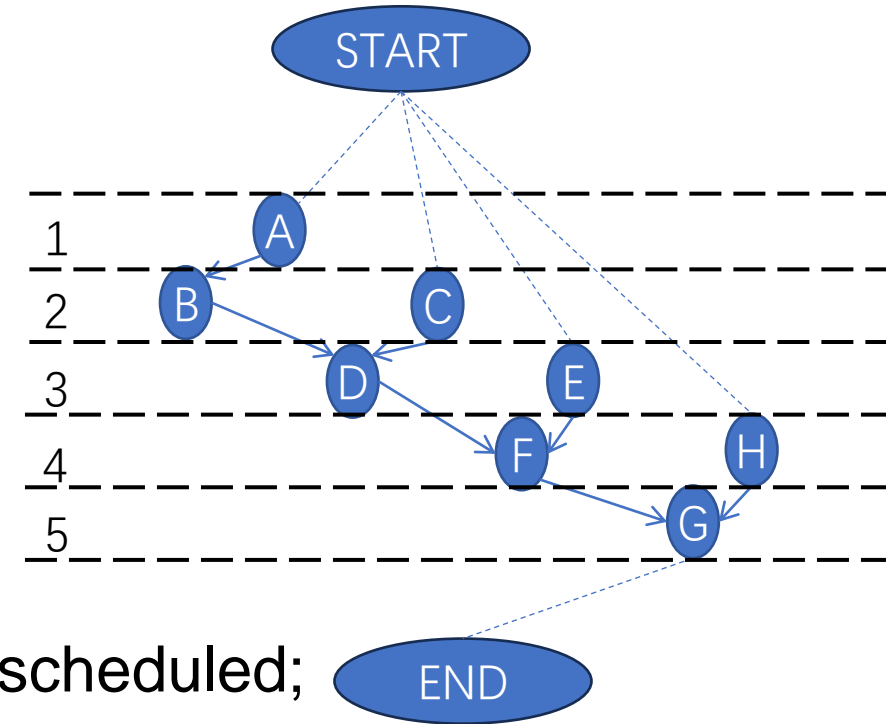
3 repeat {

4 Select a vertex v_i whose successors are all scheduled;

5 Schedule v_i by setting $t_i = \min t_j - d_i$;

7 until (v_0 is scheduled);

8 return (t);}



休息一下

随堂作业

in-class assignment

已知输入文件的格式如下：

第1行：一个数字 n ，代表顶点个数

第2到 $2+n$ 行：每一行由一个字符串和一个数字组成，以空格隔开，分别代表顶点的名字和所需的延迟时间

首尾两个顶点是虚拟顶点，代表起始点和结束点

第 $2+n+1$ 行到结束：每一行由两个顶点名组成，以空格隔开，代表第一个顶点需要在第二个顶点之前完成

请分别写出以数组方式依次存储和遍历顶点时，**ASAP**和**ALAP**调度下的程序**为顶点决定开始周期的顺序**和**输出文件的内容**，输出文件格式如下：

每一行由一个字符串和一个数字组成，以空格隔开，分别代表顶点名和开始的调度周期

输入文件内容：

7

START 1

A 2

B 1

C 3

D 2

E 1

END 1

START A

START C

A B

C B

A E

B D

E END

D END



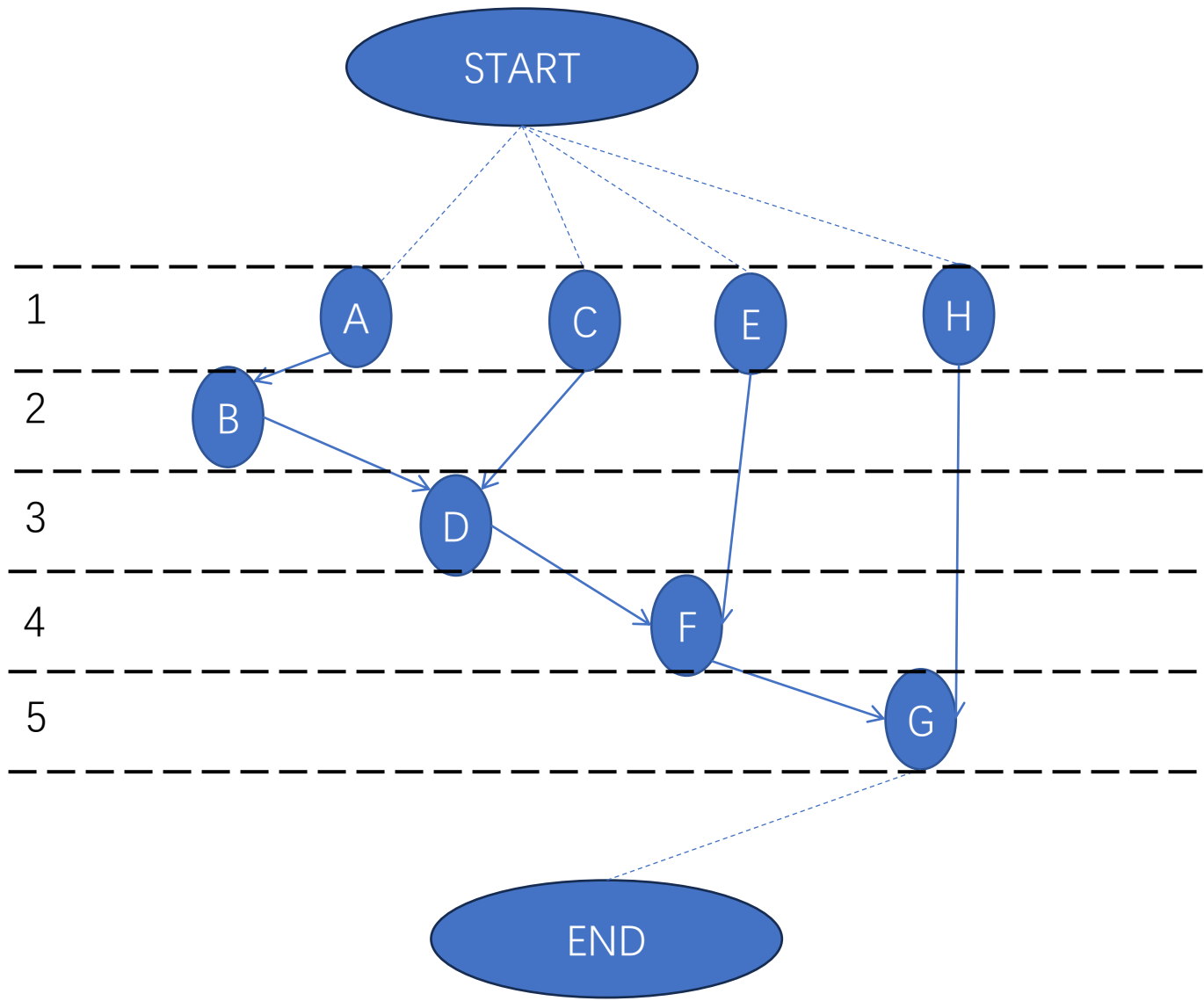
你觉得ASAP
和ALAP算法？

- A. 太难
- B. 适中
- C. 太简单

输入文件内容:

10
START 1
A 1
B 1
C 1
D 1
E 1
F 1
G 1
H 1
END 0
START B
START A
START C
START E
START G
A B
C D
B D
E F
D F
H G
F G
B END
D END
F END
G END

输入输出的例子



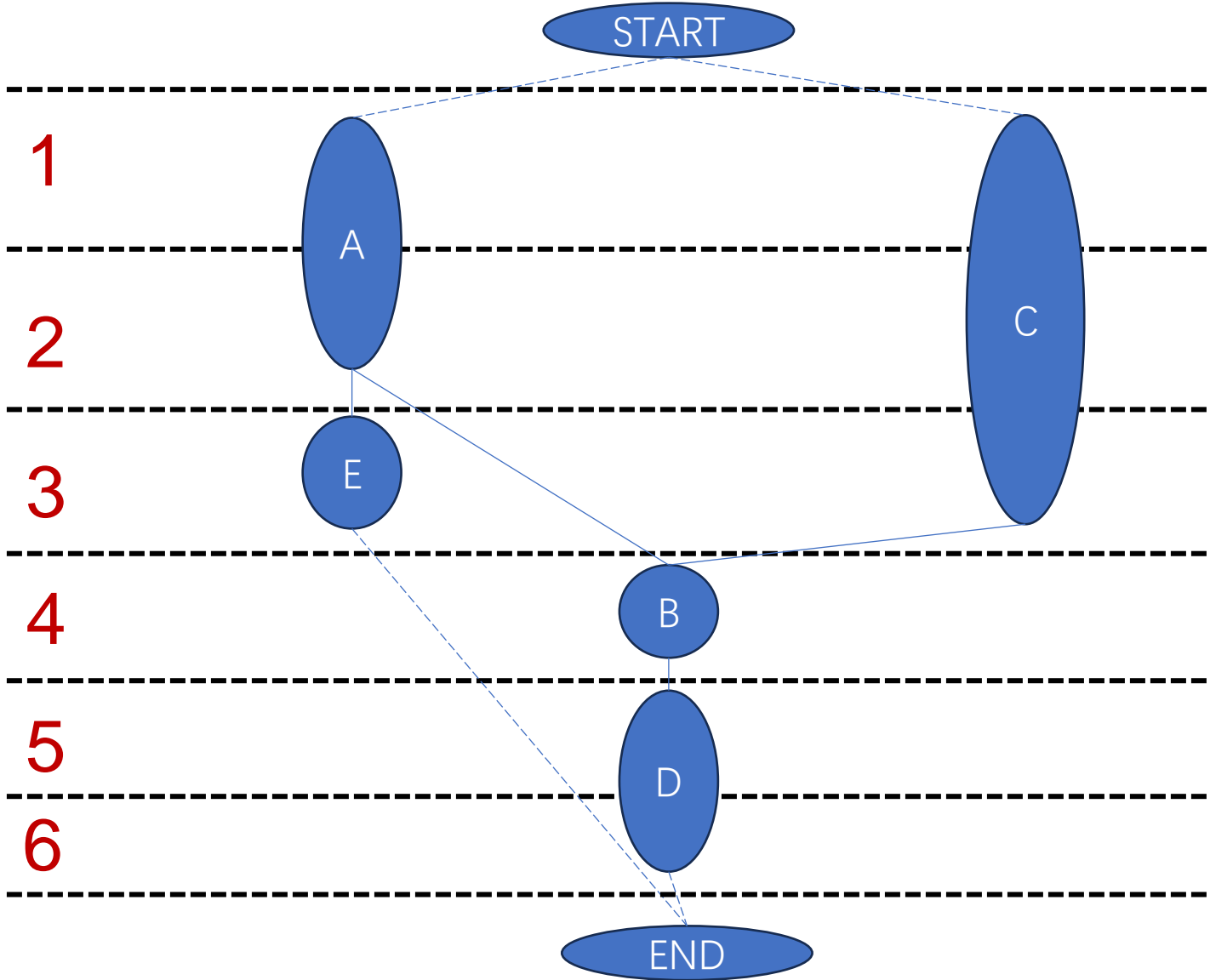
为顶点决定开始周期的顺序:
START->
A->B->C->D->E->F->H->G
->END

输出文件内容:
START 0
A 1
B 2
C 1
D 3
E 1
F 4
H 1
G 5
END 6

输入文件内容:

7
START 1
A 2
B 1
C 3
D 2
E 1
END 0
START A
START C
A B
C B
A E
B D
E END
D END

ASAP的结果



为顶点决定开始周期的顺序:
START->
A->C->B->D->E
->END

输出文件内容:
START 0
A 1
C 1
B 4
D 5
E 3
END 7

输入文件内容:

7

START 1

A 2

B 1

C 3

D 2

E 1

END 0

START A

START C

A B

C B

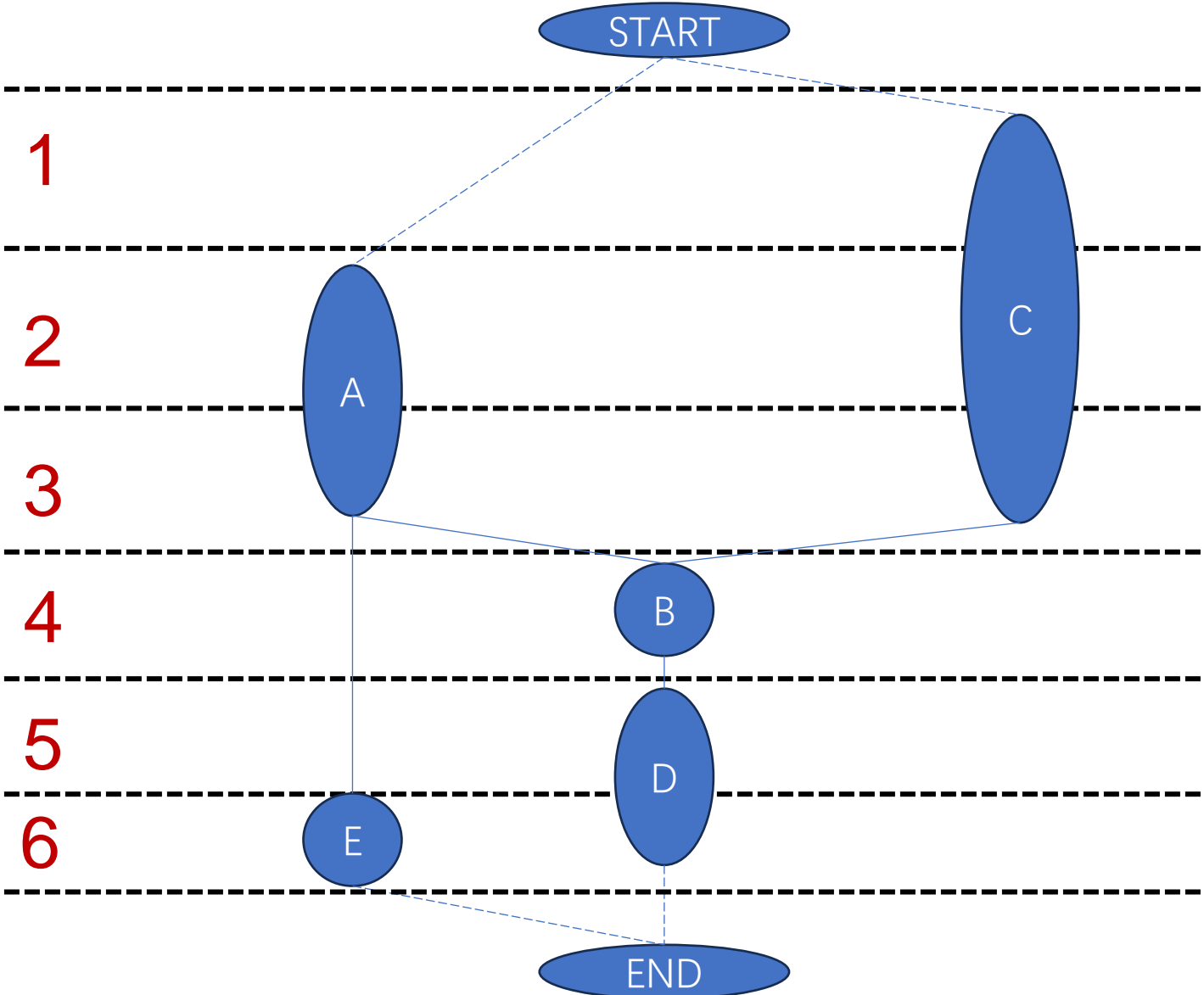
A E

B D

E END

D END

ALAP的结果



为顶点决定开始周期的顺序:

END->

D -> B -> C -> E -> A

->START

输出文件内容:

START 0

D 5

B 4

C 1

E 6

A 2

END 7

有约束的调度问题



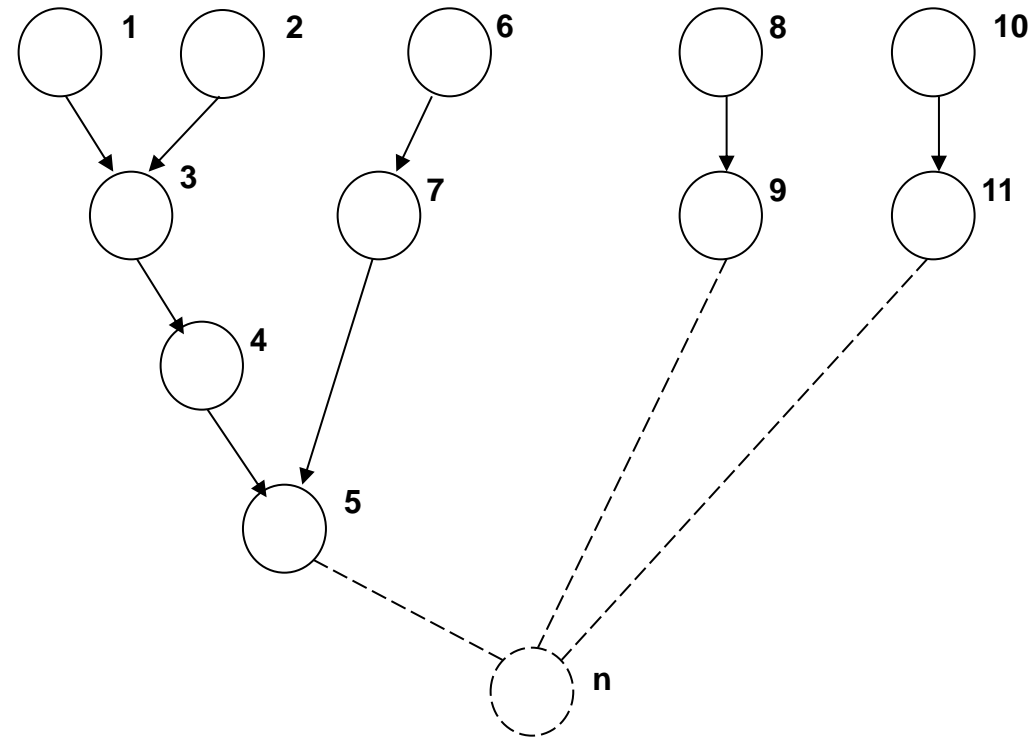
有约束的调度

Constrained Scheduling

- 约束调度
 - 一般情况下是NP完全问题
 - 在面积或资源的约束下最小化延迟 (ML-RCS)
 - 使受到延迟约束的资源最小化 (MR-LCS)
- 确切解决方法
 - ILP: 整数线性规划 (Integer linear program)
 - Hu算法: 适用于只有一种资源类型的问题
- 启发式算法
 - 列表调度 (List scheduling)
 - 力导向调度 (Force-directed scheduling)

优先级约束的多处理器调度

Precedence-constrained Multiprocessor Scheduling

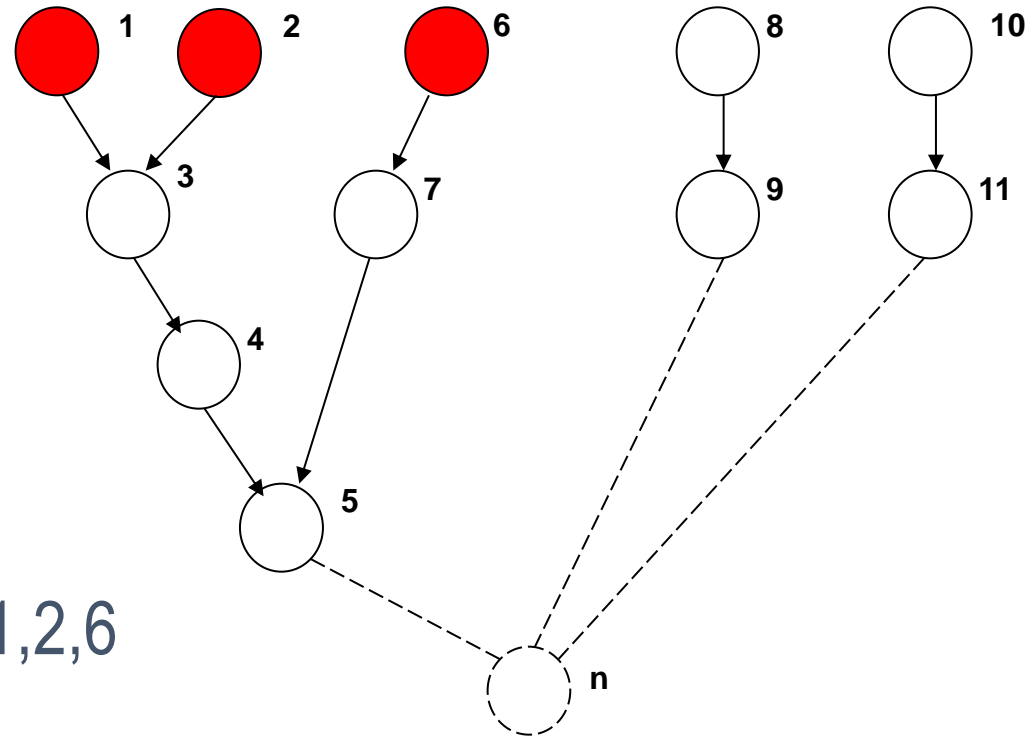


优先级约束的多处理器调度

Precedence-constrained Multiprocessor Scheduling

$\bar{a} = 3$

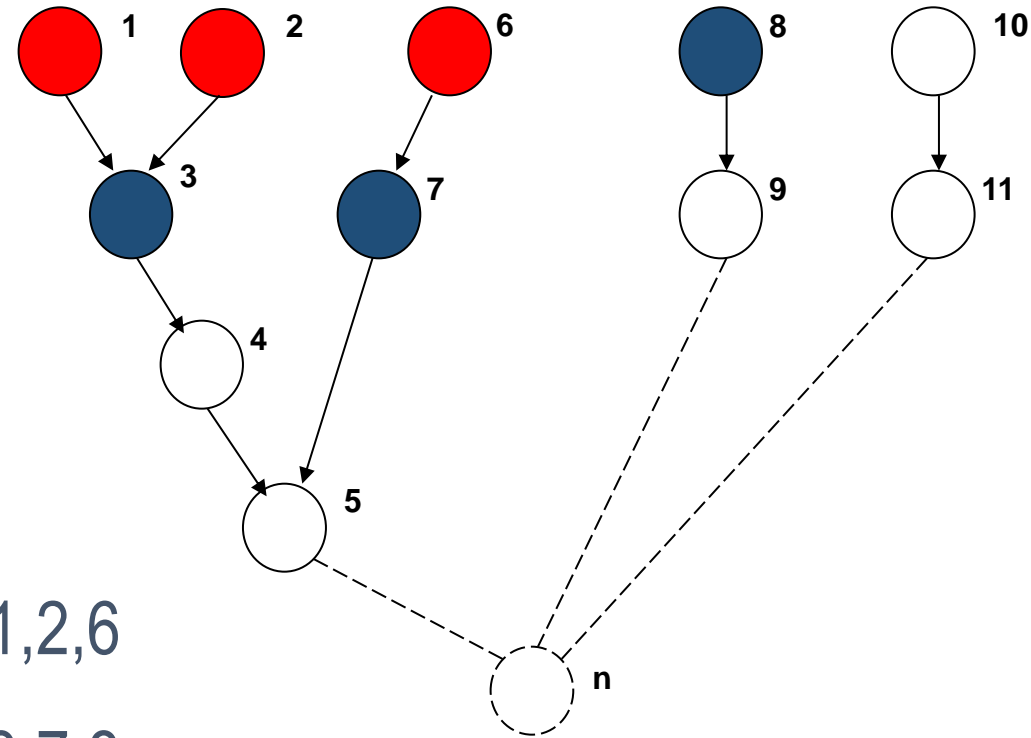
Step 1: Op 1,2,6



优先级约束的多处理器调度

Precedence-constrained Multiprocessor Scheduling

$\bar{a} = 3$



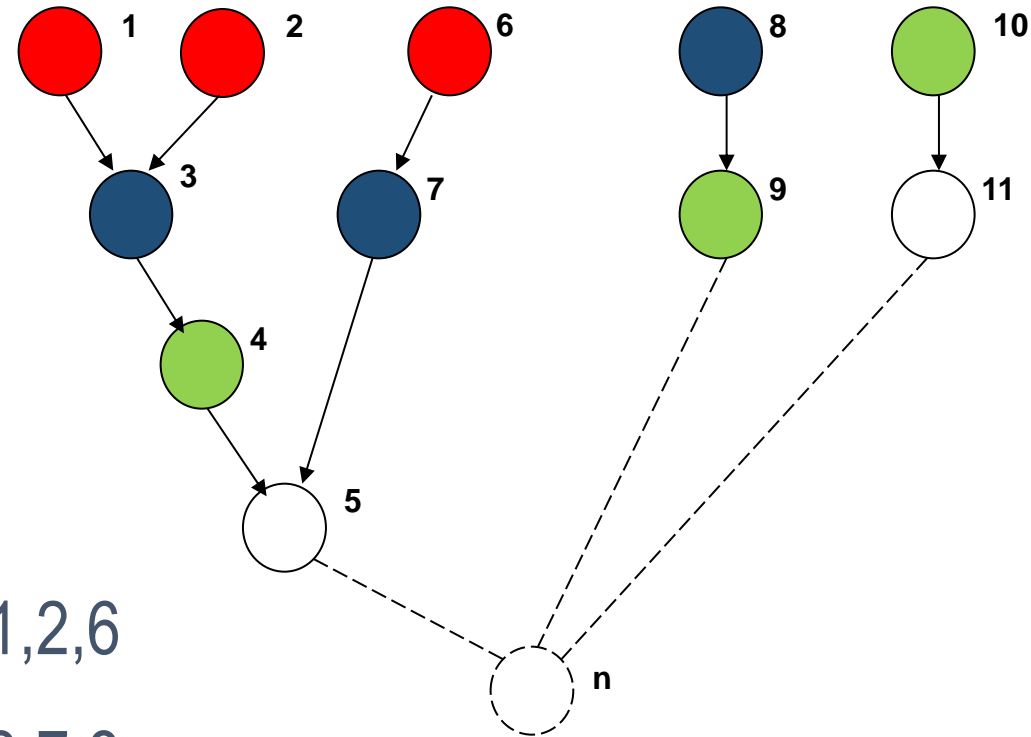
Step 1: Op 1,2,6

Step 2: Op 3,7,8

优先级约束的多处理器调度

Precedence-constrained Multiprocessor Scheduling

$\bar{a} = 3$



Step 1: Op 1,2,6

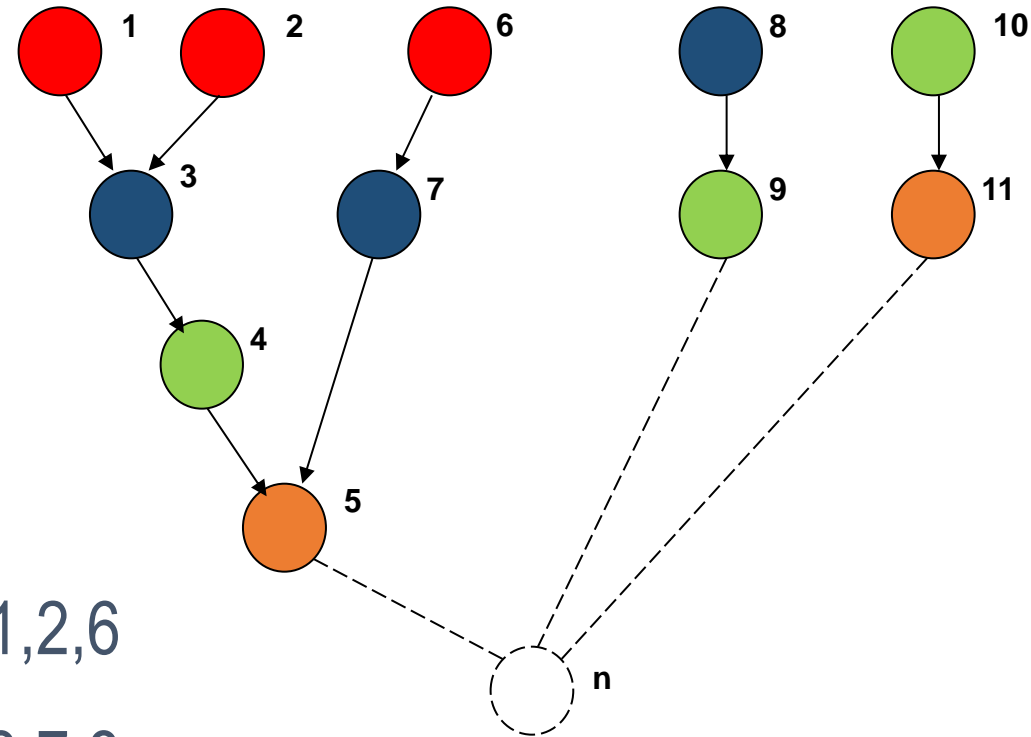
Step 2: Op 3,7,8

Step 3: Op 4,9,10

优先级约束的多处理器调度

Precedence-constrained Multiprocessor Scheduling

$\bar{a} = 3$



Step 1: Op 1,2,6

Step 2: Op 3,7,8

Step 3: Op 4,9,10

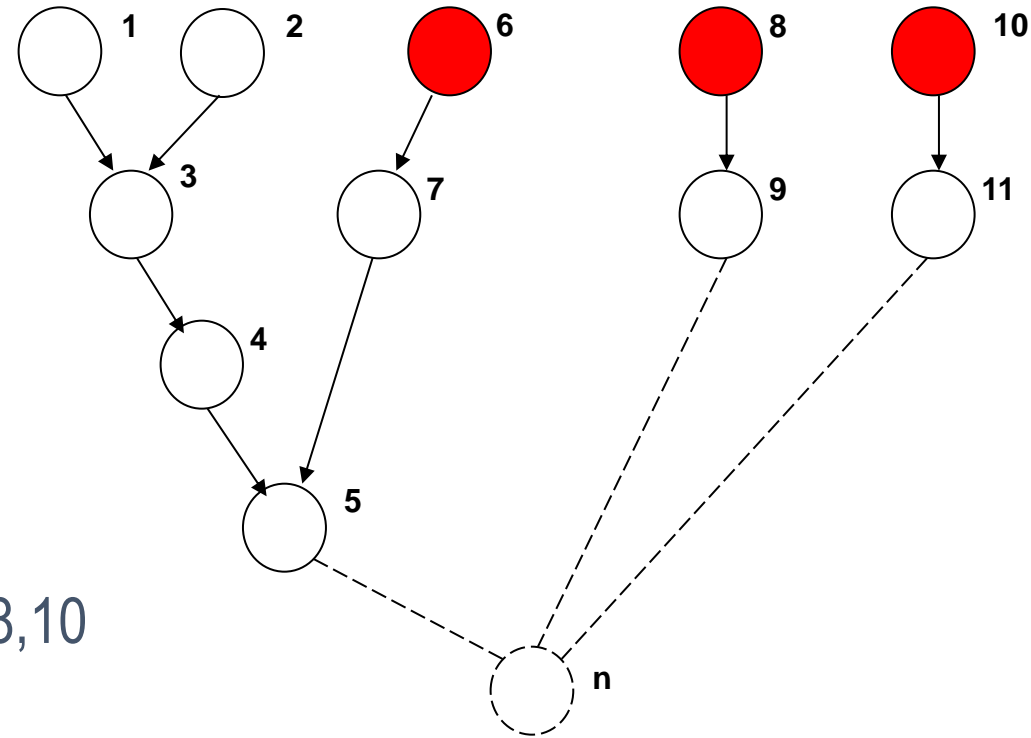
Step 4: Op 5,11

优先级约束的多处理器调度

Precedence-constrained Multiprocessor Scheduling

$\bar{a} = 3$

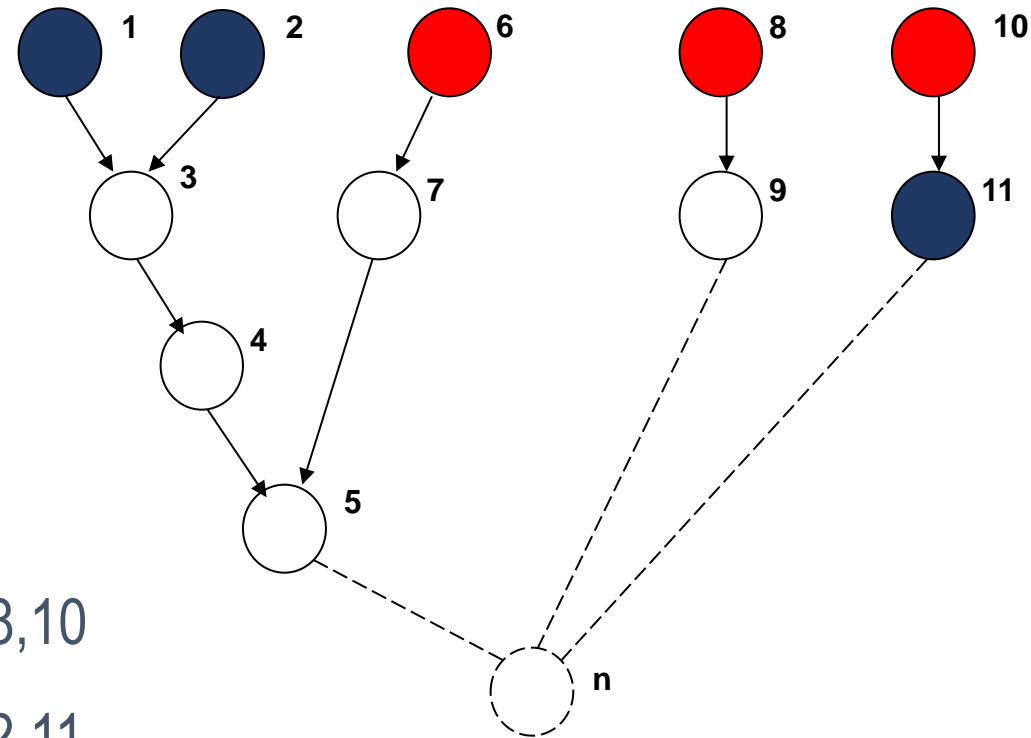
Step 1: Op 6,8,10



优先级约束的多处理器调度

Precedence-constrained Multiprocessor Scheduling

$\bar{a} = 3$



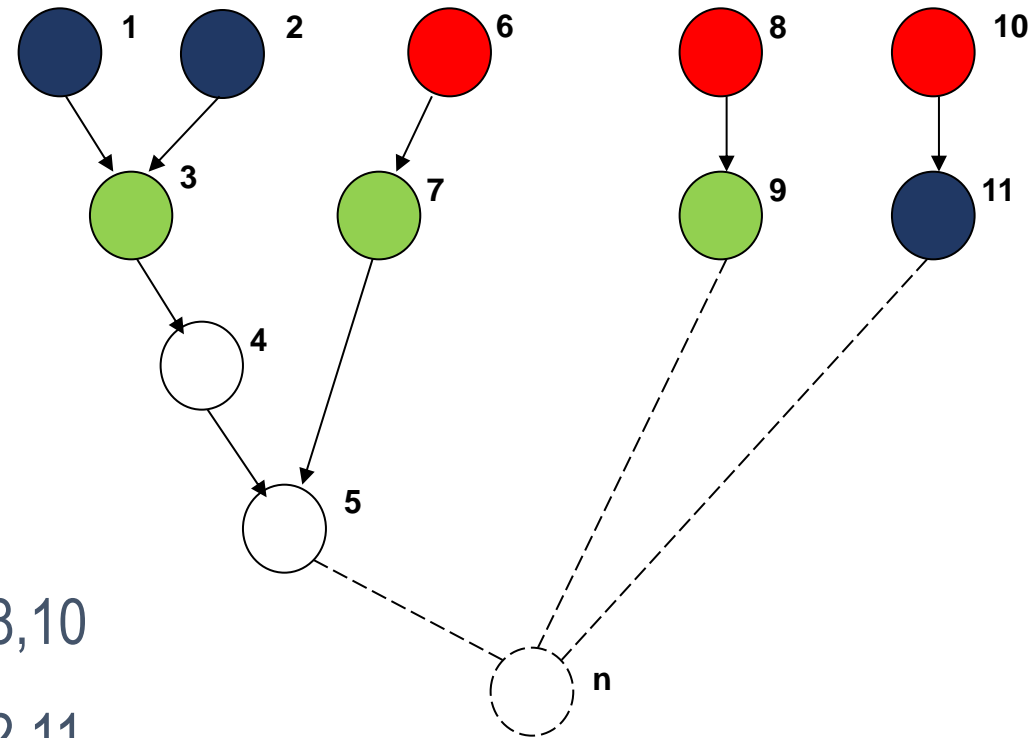
Step 1: Op 6,8,10

Step 2: Op 1,2,11

优先级约束的多处理器调度

Precedence-constrained Multiprocessor Scheduling

$\bar{a} = 3$



Step 1: Op 6,8,10

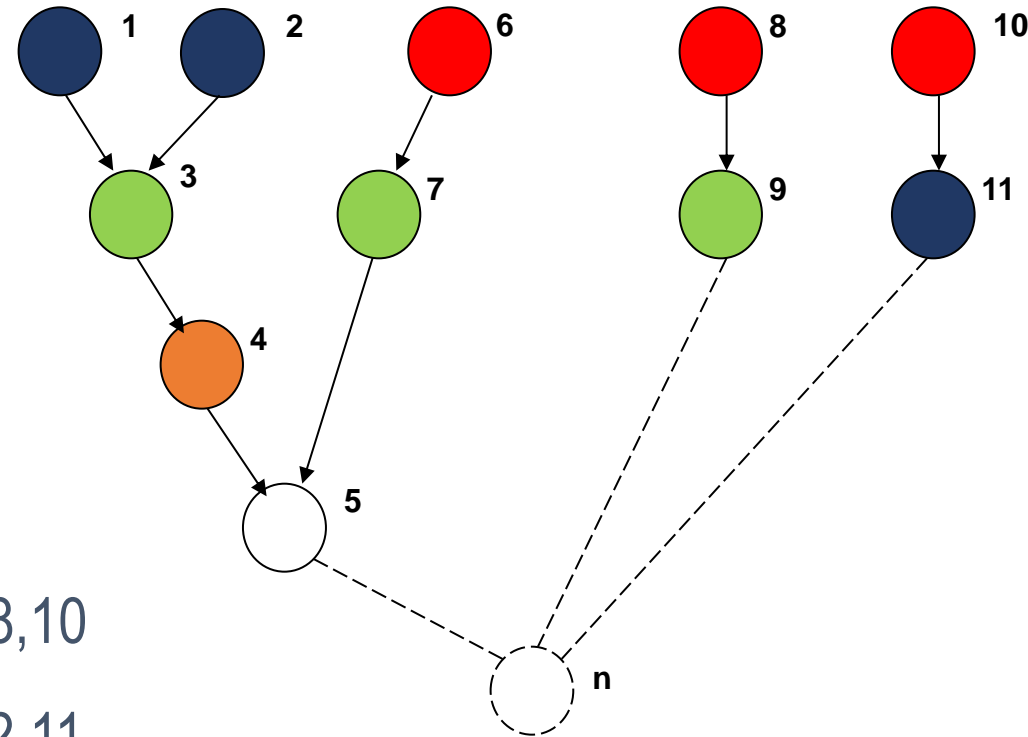
Step 2: Op 1,2,11

Step 3: Op 3,7,9

优先级约束的多处理器调度

Precedence-constrained Multiprocessor Scheduling

$\bar{a} = 3$



Step 1: Op 6,8,10

Step 2: Op 1,2,11

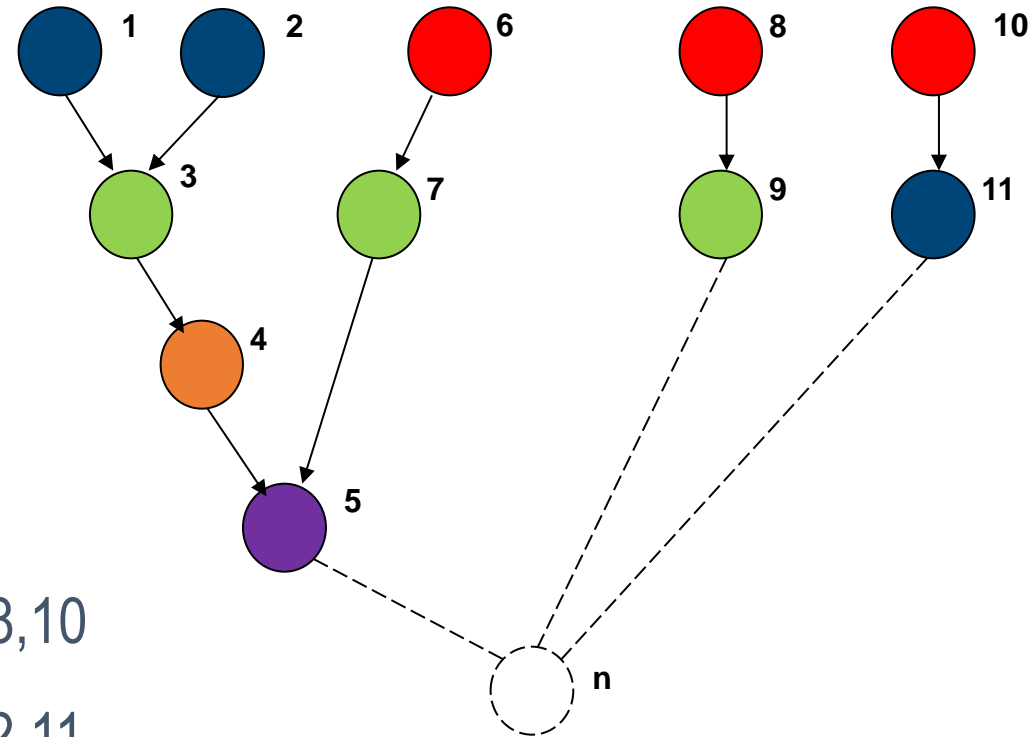
Step 3: Op 3,7,9

Step 4: Op 4

优先级约束的多处理器调度

Precedence-constrained Multiprocessor Scheduling

$\bar{a} = 3$



Step 1: Op 6,8,10

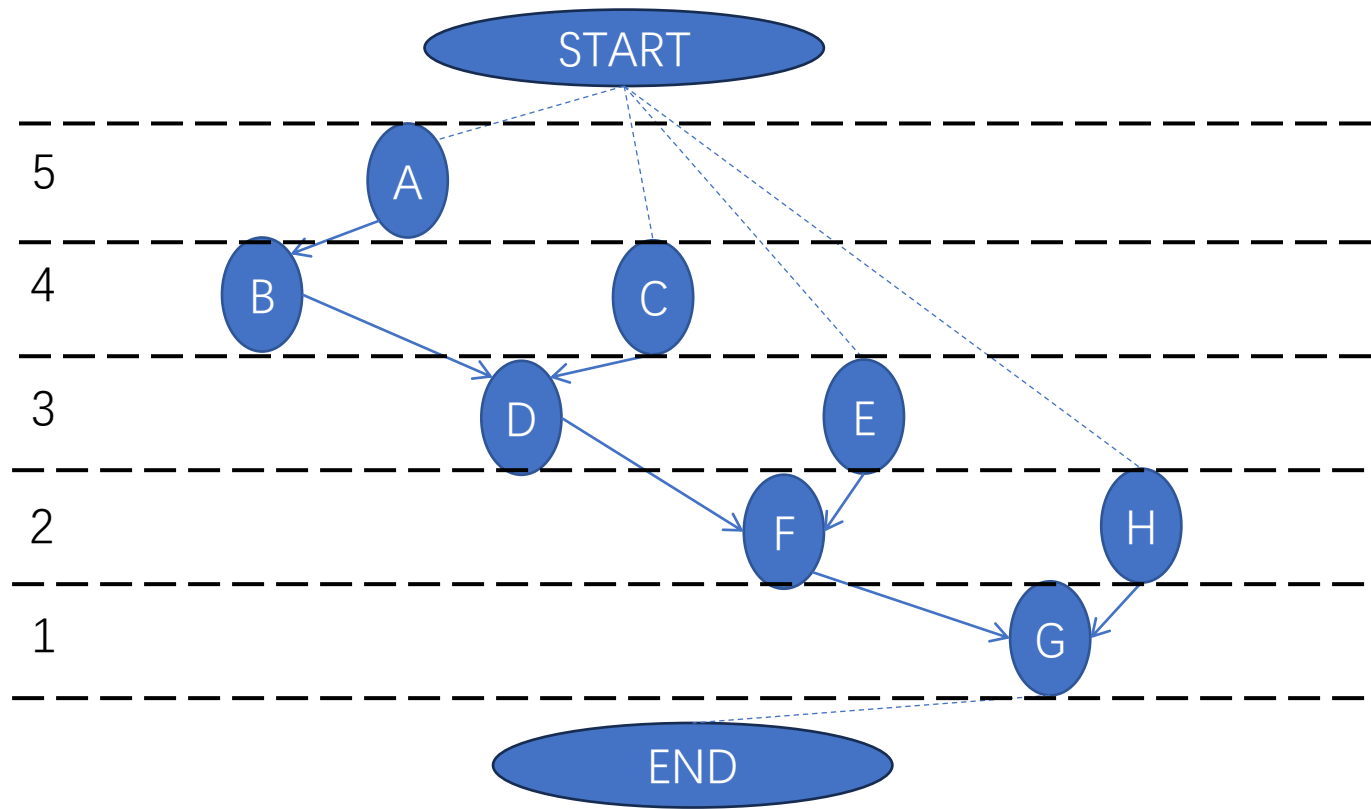
Step 2: Op 1,2,11

Step 3: Op 3,7,9

Step 4: Op 4

Step 5: Op 5

距离终点的距离



END 6
G 5
F 4
D 3
B 2
A 1
C 2
E 3
H 4



HU调度算法

HU scheduling algorithm

```
HU (G(V, E), a) {  
    Label the vertices;  
    t0 = 0; l = 1;  
    repeat {  
        U = unscheduled vertices in V without predecessors or whose predecessors have been scheduled;  
        Select  $S \subseteq U$  vertices, such that  $|S| \leq a$  and labels in S are maximal;  
        Schedule the S operations at step t by setting  $t_i = l \ \forall v_i \in S$ ;  
        l = l + 1;  
    }  
    until ( $v_n$  is scheduled);  
}
```



HU调度算法

HU scheduling algorithm

HU ($G(V, E), a$) {

对顶点进行标号;

$t_0 = 0; l = 1;$

重复执行以下步骤:

U = 仍未调度的顶点集合, 这些顶点要么没有前驱, 要么其所有前驱都已被调度;

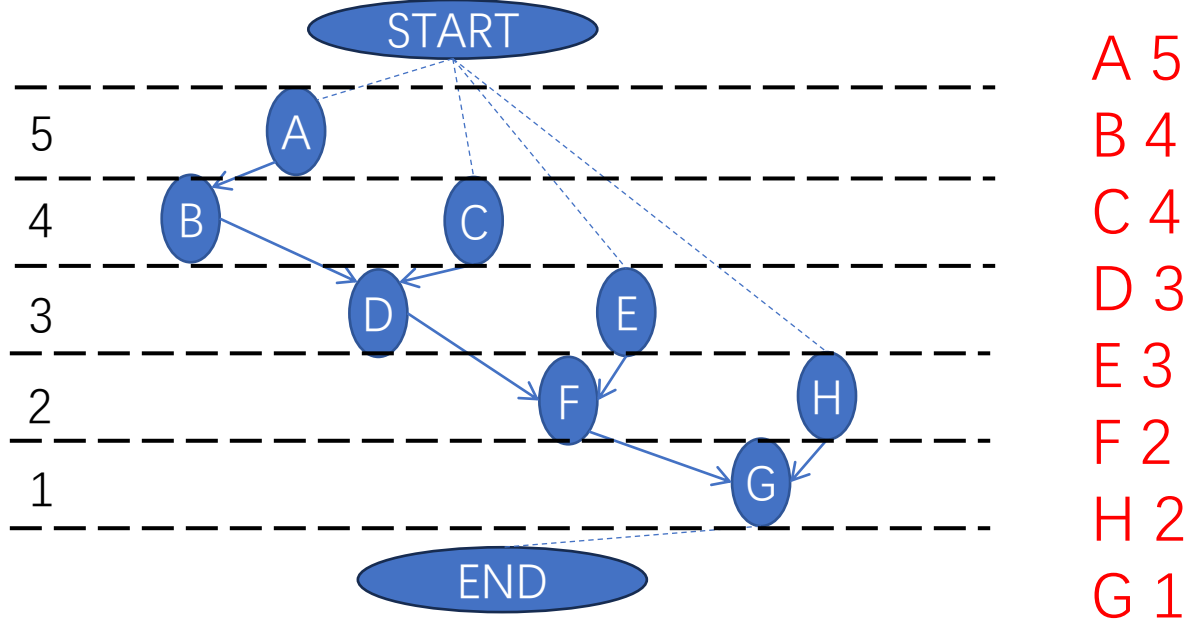
选择一个子集 $S \subseteq U$, 使得 $|S| \leq a$, 且 S 中的顶点标签值最大;

在步骤 t 处调度 S 中的所有操作, 即对所有 $v_i \in S$ 设定调度时间 $t_i = l$;

$l = l + 1;$

直到 v_n 被调度;

}



HU (G(V, E), a) {

对顶点进行标号;

$t_0 = 0$; $l = 1$;

重复执行以下步骤:

U = 仍未调度的顶点集合, 这些顶点要么没有前驱, 要么其所有前驱都已被调度;

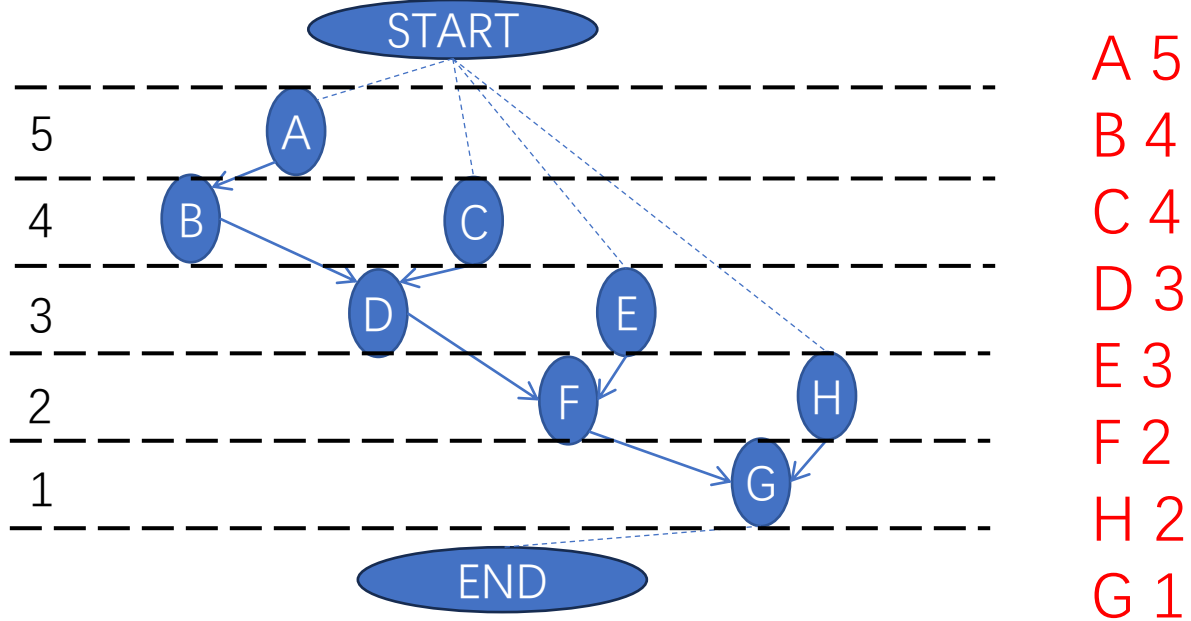
选择一个子集 $S \subseteq U$, 使得 $|S| \leq a$, 且 S 中的顶点标签值最大;

在步骤 t 处调度 S 中的所有操作, 即对所有 $v_i \in S$ 设定调度时间 $t_i = l$;

$l = l + 1$;

直到 v_n 被调度;

}



START 0

HU (G(V, E), a) {

对顶点进行标号;

$t_0 = 0; l = 1;$

重复执行以下步骤:

U = 仍未调度的顶点集合, 这些顶点要么没有前驱, 要么其所有前驱都已被调度;

选择一个子集 $S \subseteq U$, 使得 $|S| \leq a$, 且 S 中的顶点标签值最大;

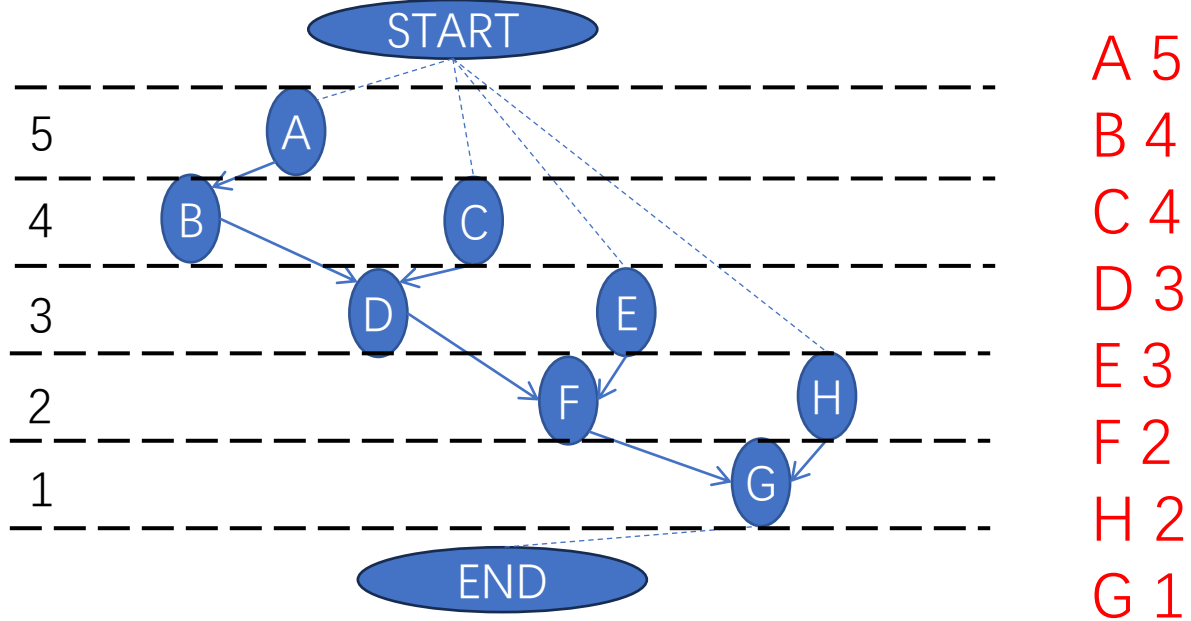
在步骤 t 处调度 S 中的所有操作, 即对所有 $v_i \in S$ 设定调度时间 $t_i = l$;

$l = l + 1;$

直到 v_n 被调度;

}

| 变量 | 当前值 |
|----|-----|
| l | 1 |
| U | |
| S | |



START 0

HU (G(V, E), a) {

对顶点进行标号;

t0 = 0; l = 1;

重复执行以下步骤:

U = 仍未调度的顶点集合, 这些顶点要么没有前驱, 要么其所有前驱都已被调度;

选择一个子集 $S \subseteq U$, 使得 $|S| \leq a$, 且 S 中的顶点标签值最大;

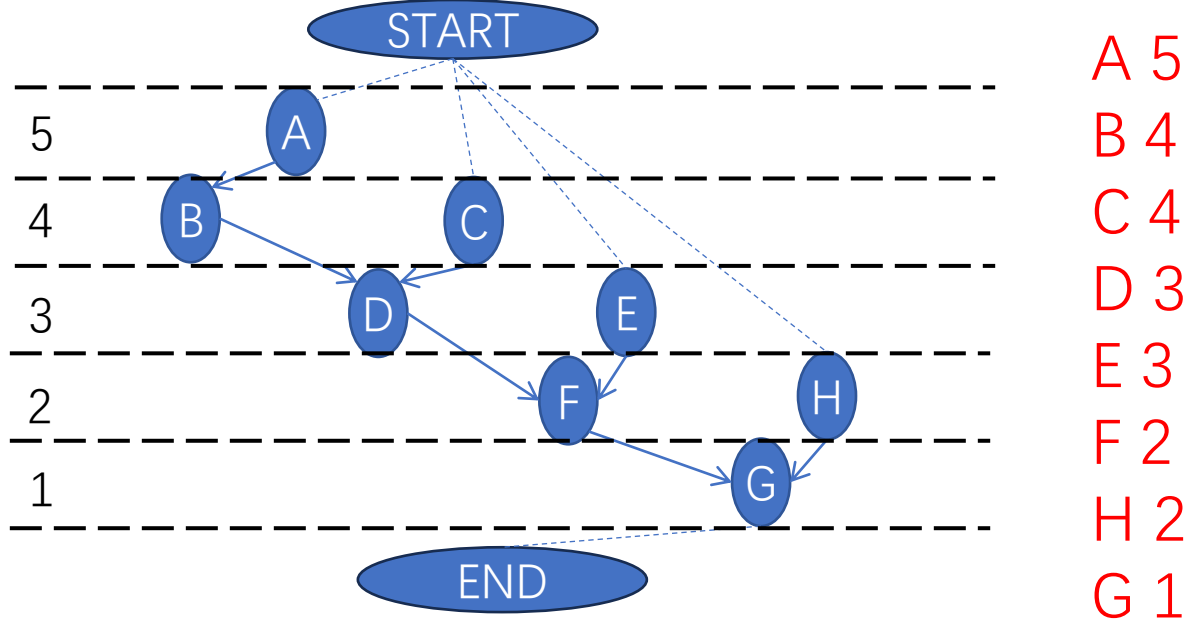
在步骤 t 处调度 S 中的所有操作, 即对所有 $v_i \in S$ 设定调度时间 $t_i = l$;

$l = l + 1$;

直到 vn 被调度;

}

| 变量 | 当前值 |
|----|-----|
| l | 1 |
| U | |
| S | |



HU (G(V, E), a) {

对顶点进行标号;

t0 = 0; l = 1;

重复执行以下步骤:

U = 仍未调度的顶点集合, 这些顶点要么没有前驱, 要么其所有前驱都已被调度;

选择一个子集 $S \subseteq U$, 使得 $|S| \leq a$, 且 S 中的顶点标签值最大;

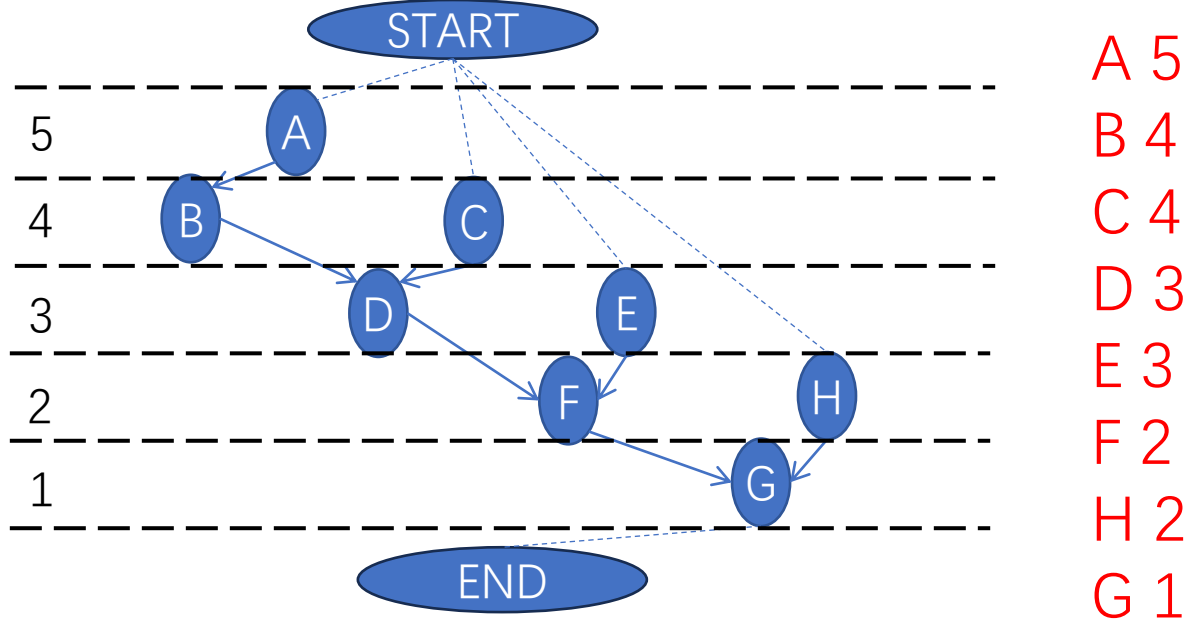
在步骤 t 处调度 S 中的所有操作, 即对所有 $v_i \in S$ 设定调度时间 $t_i = l$;

$l = l + 1$;

直到 v_n 被调度;

}

| 变量 | 当前值 |
|----|--------------|
| l | 1 |
| U | {A, C, E, H} |
| S | |



HU (G(V, E), a) {

对顶点进行标号;

t0 = 0; l = 1;

重复执行以下步骤:

U = 仍未调度的顶点集合, 这些顶点要么没有前驱, 要么其所有前驱都已被调度;

选择一个子集 $S \subseteq U$, 使得 $|S| \leq a$, 且 S 中的顶点标签值最大;

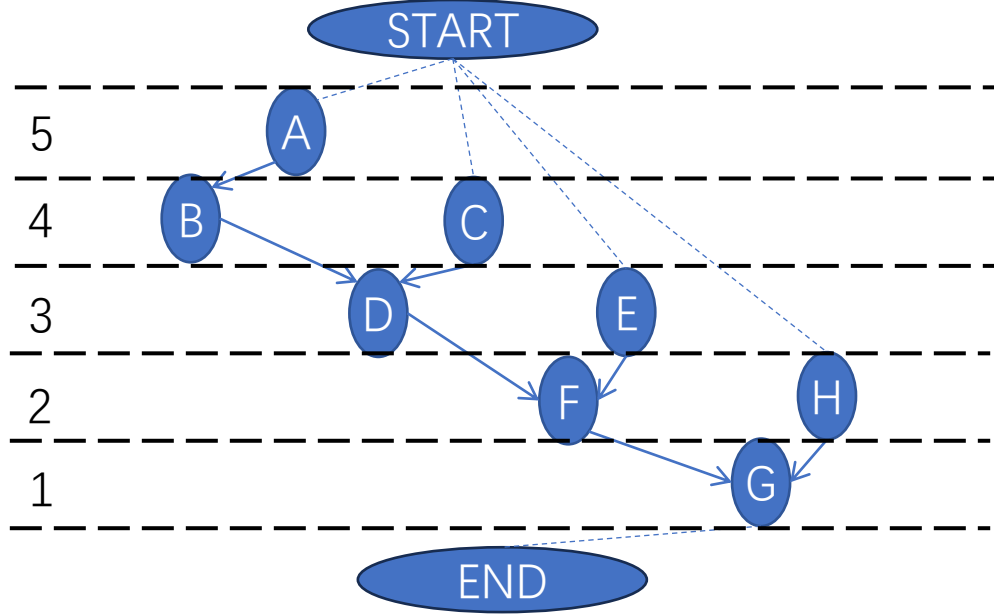
在步骤 t 处调度 S 中的所有操作, 即对所有 $v_i \in S$ 设定调度时间 $t_i = l$;

$l = l + 1$;

直到 vn 被调度;

}

| 变量 | 当前值 |
|----|--------------|
| l | 1 |
| U | {A, C, E, H} |
| S | {A,C,E} |



A 5
B 4
C 4
D 3
E 3
F 2
H 2
G 1

START 0
A 1
C 1
E 1

HU (G(V, E), a) {

对顶点进行标号;

t0 = 0; l = 1;

重复执行以下步骤:

U = 仍未调度的顶点集合, 这些顶点要么没有前驱, 要么其所有前驱都已被调度;

选择一个子集 $S \subseteq U$, 使得 $|S| \leq a$, 且 S 中的顶点标签值最大;

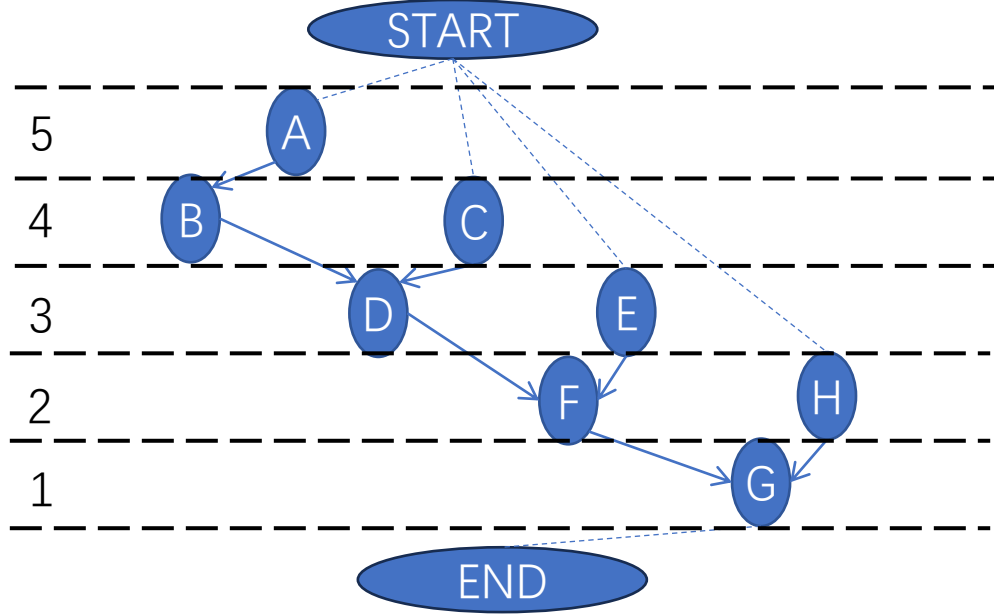
在步骤 t 处调度 S 中的所有操作, 即对所有 $v_i \in S$ 设定调度时间 $t_i = l$;

$l = l + 1$;

直到 vn 被调度;

}

| 变量 | 当前值 |
|----|--------------|
| l | 1 |
| U | {A, C, E, H} |
| S | {A,C,E} |



A 5
B 4
C 4
D 3
E 3
F 2
H 2
G 1

START 0
A 1
C 1
E 1

HU (G(V, E), a) {

对顶点进行标号;

t0 = 0; l = 1;

重复执行以下步骤:

U = 仍未调度的顶点集合, 这些顶点要么没有前驱, 要么其所有前驱都已被调度;

选择一个子集 $S \subseteq U$, 使得 $|S| \leq a$, 且 S 中的顶点标签值最大;

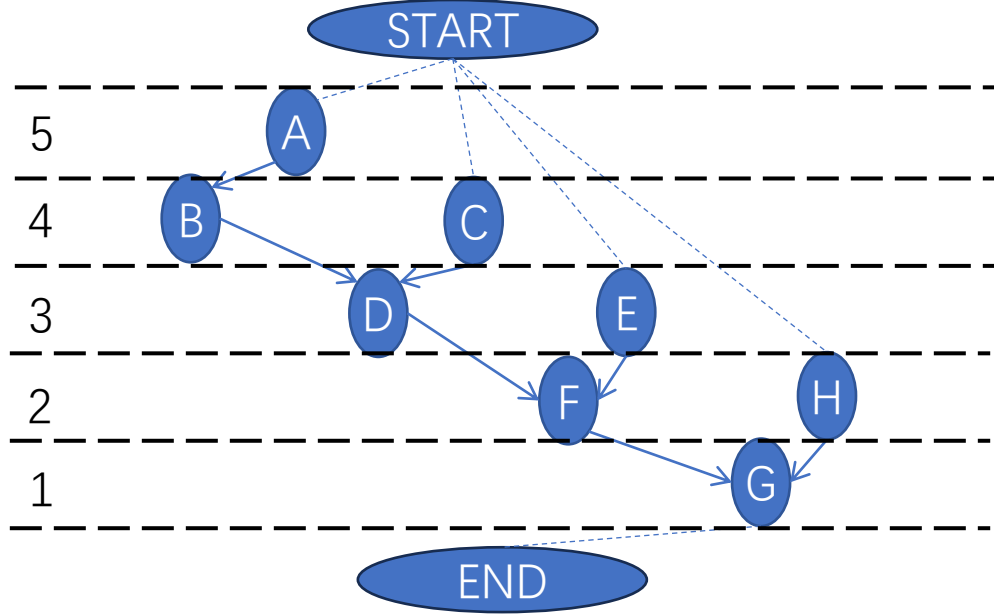
在步骤 t 处调度 S 中的所有操作, 即对所有 $v_i \in S$ 设定调度时间 $t_i = l$;

$l = l + 1$;

直到 vn 被调度;

}

| 变量 | 当前值 |
|----|--------------|
| l | 2 |
| U | {A, C, E, H} |
| S | {A,C,E} |



A 5
B 4
C 4
D 3
E 3
F 2
H 2
G 1

START 0
A 1
C 1
E 1

HU (G(V, E), a) {

对顶点进行标号;

t0 = 0; l = 1;

重复执行以下步骤:

U = 仍未调度的顶点集合, 这些顶点要么没有前驱, 要么其所有前驱都已被调度;

选择一个子集 $S \subseteq U$, 使得 $|S| \leq a$, 且 S 中的顶点标签值最大;

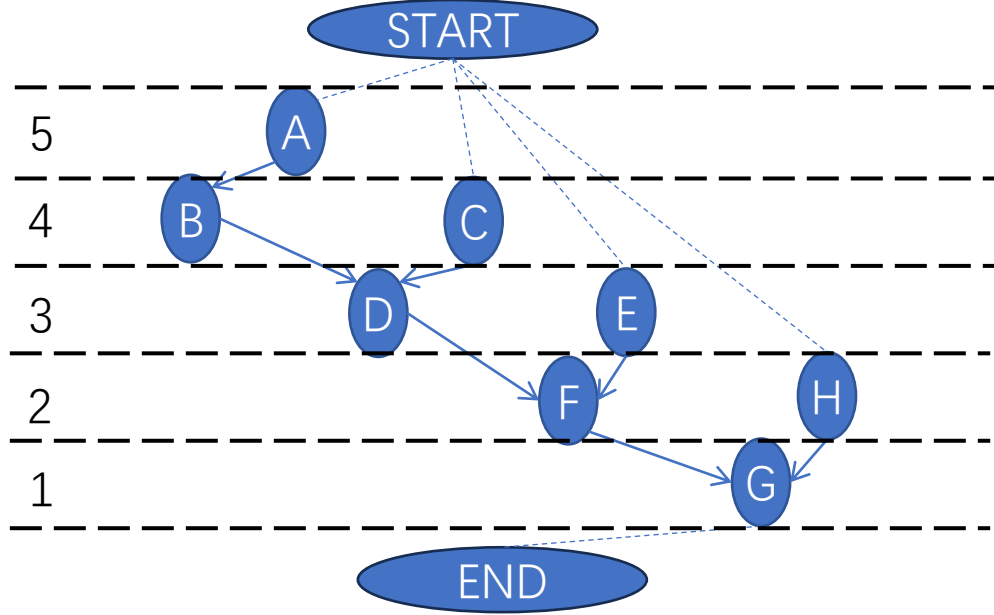
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A 5
B 4
C 4
D 3
E 3
F 2
H 2
G 1

START 0
A 1
C 1
E 1

HU (G(V, E), a) {

对顶点进行标号;

t0 = 0; l = 1;

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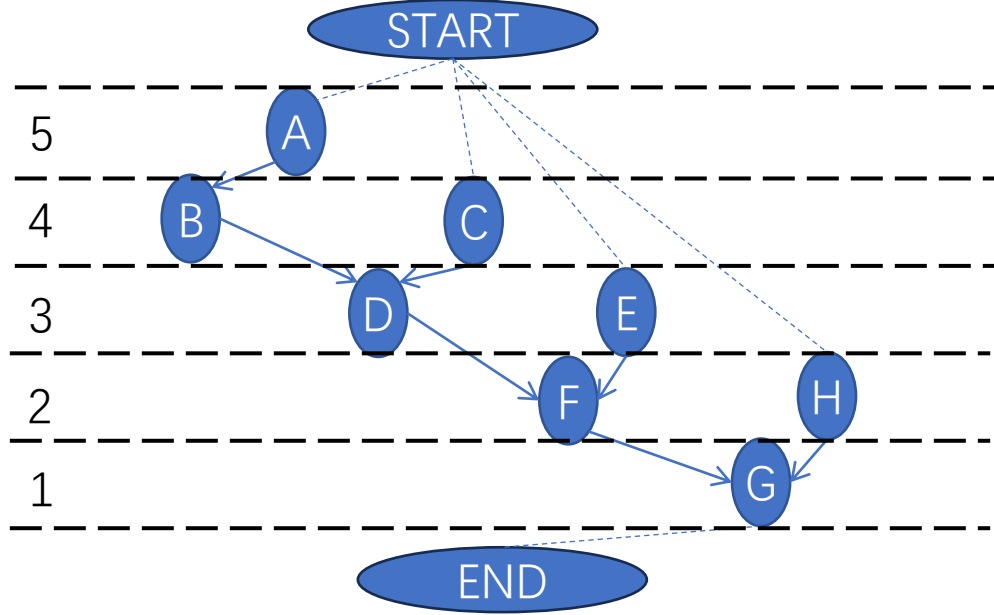
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$l = l + 1$;

直到 vn 被调度;

}

| 变量 | 当前值 |
|----|-----------|
| l | 2 |
| U | {H, B} |
| S | {A, C, E} |



A 5
B 4
C 4
D 3
E 3
F 2
H 2
G 1

START 0
A 1
C 1
E 1

HU (G(V, E), a) {

对顶点进行标号;

t0 = 0; l = 1;

重复执行以下步骤:

U = 仍未调度的顶点集合, 这些顶点要么没有前驱, 要么其所有前驱都已被调度;

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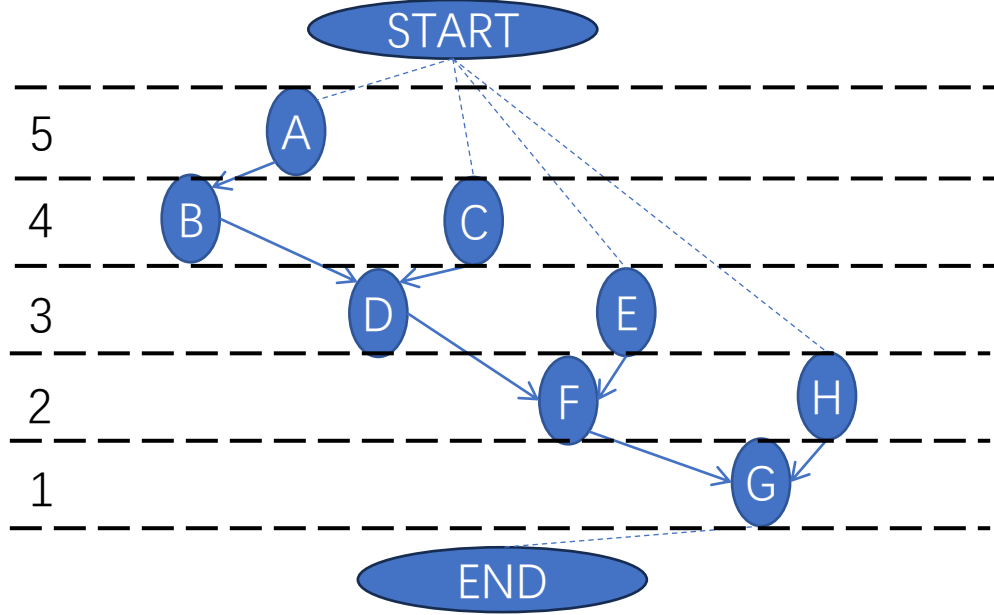
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$l = l + 1$;

直到 vn 被调度;

}

| 变量 | 当前值 |
|----|--------|
| l | 2 |
| U | {H, B} |
| S | {H, B} |



A 5
B 4
C 4
D 3
E 3
F 2
H 2
G 1

START 0
A 1
C 1
E 1
H 2
B 2

HU (G(V, E), a) {

对顶点进行标号;

t0 = 0; l = 1;

重复执行以下步骤:

U = 仍未调度的顶点集合, 这些顶点要么没有前驱, 要么其所有前驱都已被调度;

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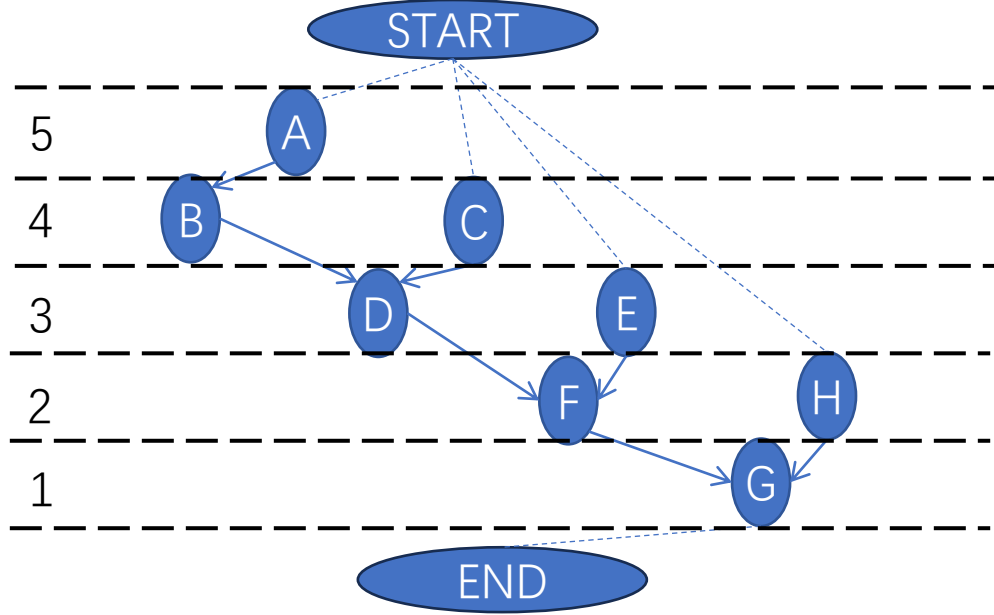
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$l = l + 1$;

直到 vn 被调度;

}

| 变量 | 当前值 |
|----|--------|
| l | 2 |
| U | {H, B} |
| S | {H, B} |



A 5
B 4
C 4
D 3
E 3
F 2
H 2
G 1

START 0
A 1
C 1
E 1
H 2
B 2

HU (G(V, E), a) {

对顶点进行标号;

t0 = 0; l = 1;

重复执行以下步骤:

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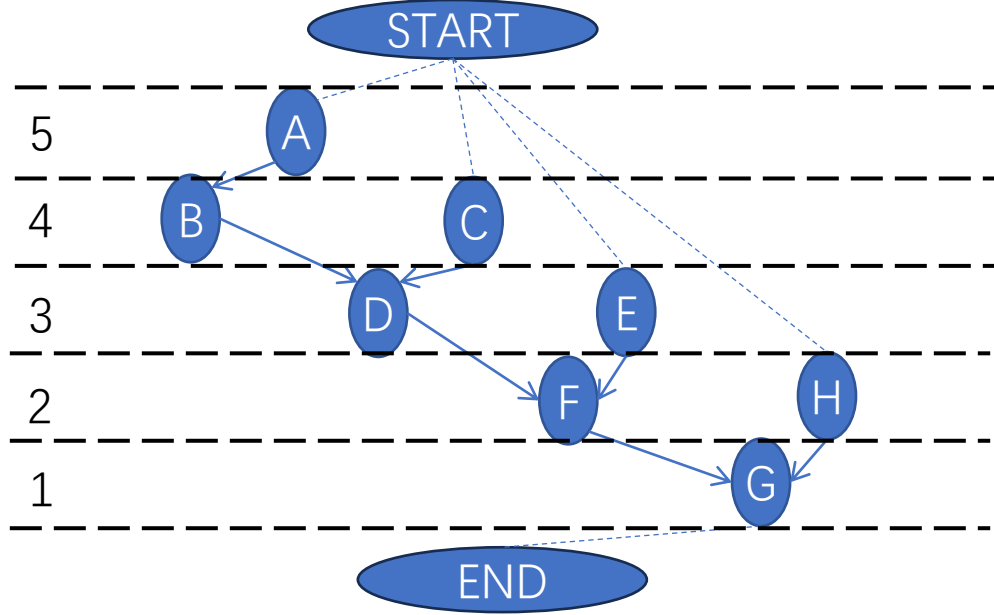
在步骤 t 处调度 S 中的所有操作, 即对所有 $v_i \in S$ 设定调度时间 $t_i = l$;

$l = l + 1$;

直到 vn 被调度;

}

| 变量 | 当前值 |
|----|--------|
| l | 3 |
| U | {H, B} |
| S | {H, B} |



A 5
B 4
C 4
D 3
E 3
F 2
H 2
G 1

START 0
A 1
C 1
E 1
H 2
B 2

HU (G(V, E), a) {

对顶点进行标号;

t0 = 0; l = 1;

重复执行以下步骤:

U = 仍未调度的顶点集合, 这些顶点要么没有前驱, 要么其所有前驱都已被调度;

选择一个子集 S ⊆ U, 使得 |S| ≤ a, 且 S 中的顶点标签值最大;

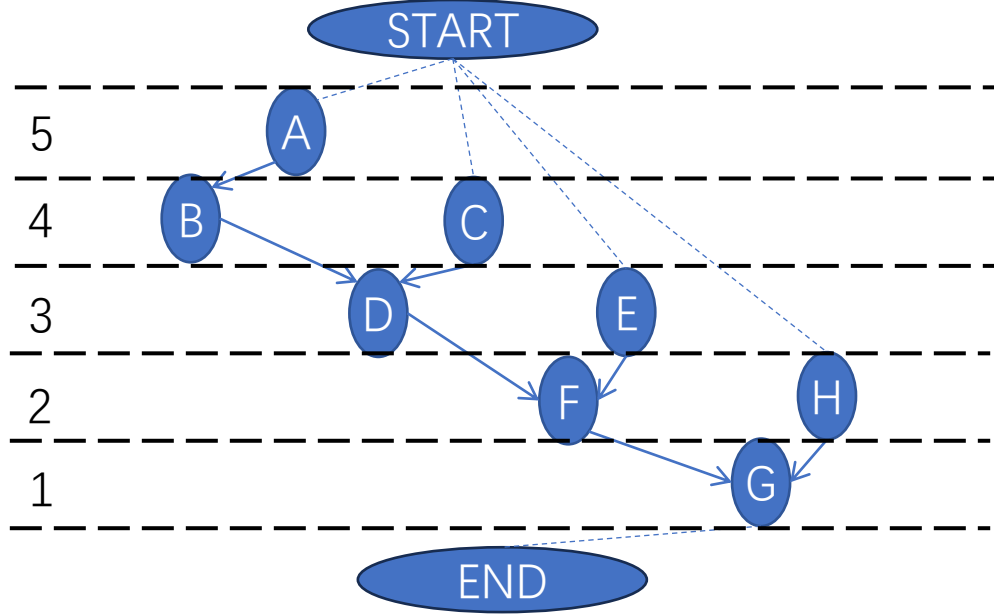
在步骤 t 处调度 S 中的所有操作, 即对所有 vi ∈ S 设定调度时间 ti = l;

l = l + 1;

直到 vn 被调度;

| 变量 | 当前值 |
|----|--------|
| l | 3 |
| U | {D} |
| S | {H, B} |

}



A 5
B 4
C 4
D 3
E 3
F 2
H 2
G 1

START 0
A 1
C 1
E 1
H 2
B 2
D 3

HU (G(V, E), a) {

对顶点进行标号;

t0 = 0; l = 1;

重复执行以下步骤:

U = 仍未调度的顶点集合, 这些顶点要么没有前驱, 要么其所有前驱都已被调度;

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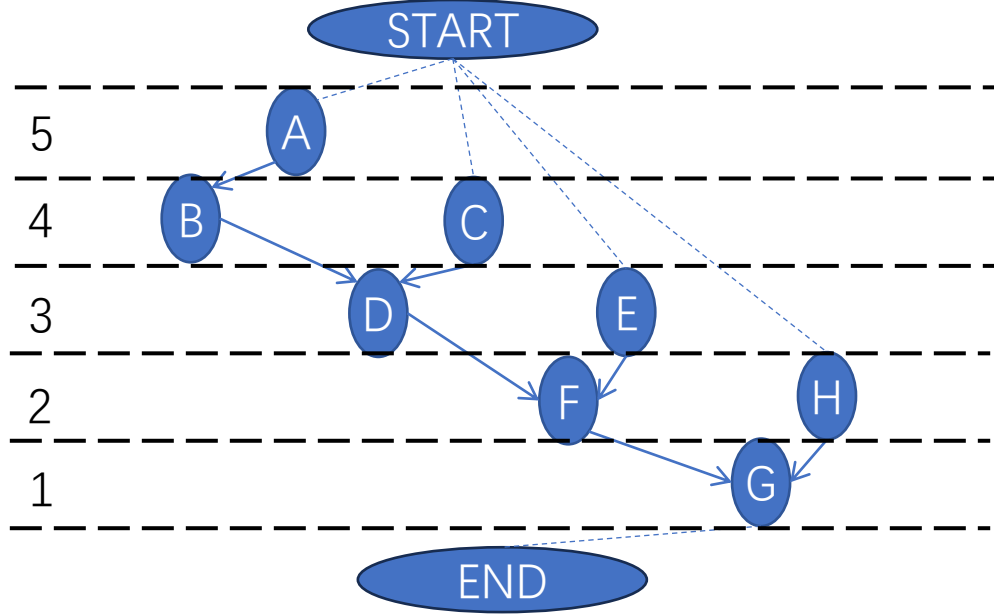
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$l = l + 1$;

直到 vn 被调度;

}

| 变量 | 当前值 |
|----|-----|
| l | 3 |
| U | {D} |
| S | {D} |



A 5
B 4
C 4
D 3
E 3
F 2
H 2
G 1

START 0
A 1
C 1
E 1
H 2
B 2
D 3

HU (G(V, E), a) {

对顶点进行标号;

t0 = 0; l = 1;

重复执行以下步骤:

U = 仍未调度的顶点集合, 这些顶点要么没有前驱, 要么其所有前驱都已被调度;

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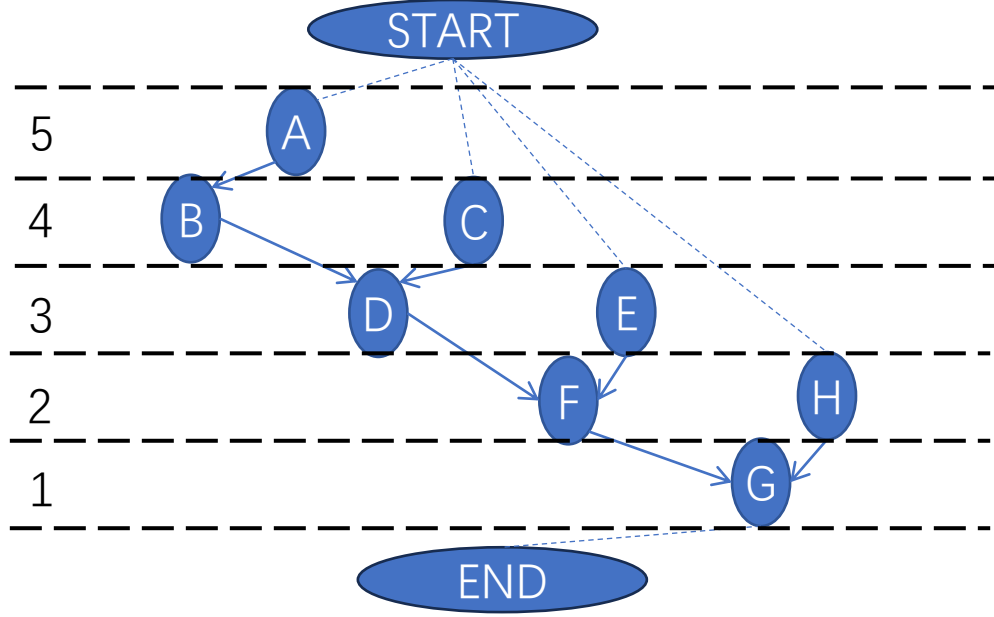
在步骤 t 处调度 S 中的所有操作, 即对所有 $v_i \in S$ 设定调度时间 $t_i = l$;

$l = l + 1$;

直到 vn 被调度;

}

| 变量 | 当前值 |
|----|-----|
| l | 4 |
| U | {D} |
| S | {D} |



A 5
B 4
C 4
D 3
E 3
F 2
H 2
G 1

START 0
A 1
C 1
E 1
H 2
B 2
D 3

HU (G(V, E), a) {

对顶点进行标号;

t0 = 0; l = 1;

重复执行以下步骤:

U = 仍未调度的顶点集合, 这些顶点要么没有前驱, 要么其所有前驱都已被调度;

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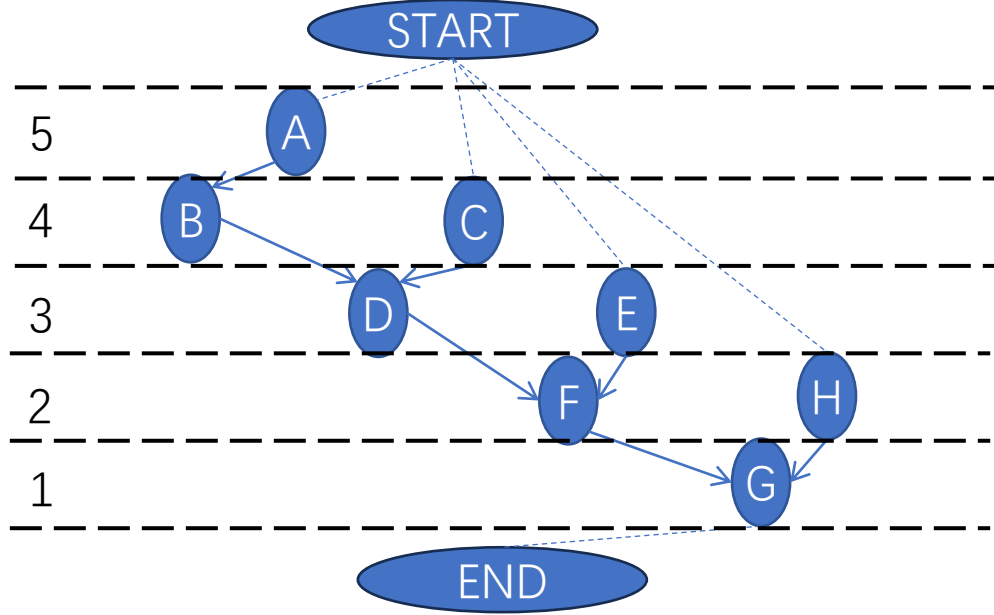
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$l = l + 1$;

直到 vn 被调度;

}

| 变量 | 当前值 |
|----|-----|
| l | 4 |
| U | {F} |
| S | {D} |



A 5
B 4
C 4
D 3
E 3
F 2
H 2
G 1

START 0
A 1
C 1
E 1
H 2
B 2
D 3
F 4

HU (G(V, E), a) {

对顶点进行标号;

t0 = 0; l = 1;

重复执行以下步骤:

U = 仍未调度的顶点集合, 这些顶点要么没有前驱, 要么其所有前驱都已被调度;

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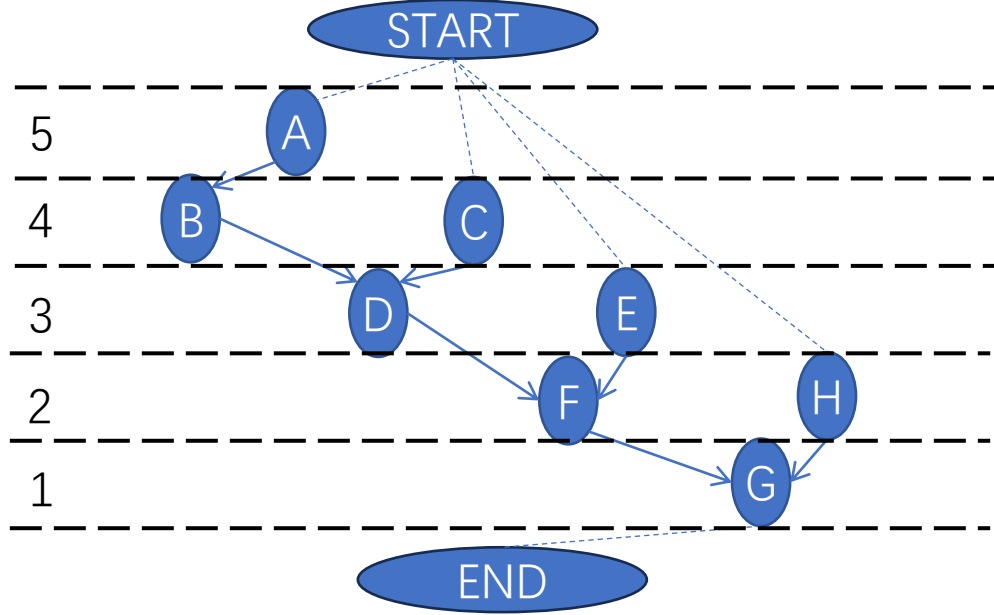
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$l = l + 1$;

直到 vn 被调度;

}

| 变量 | 当前值 |
|----|-----|
| l | 4 |
| U | {F} |
| S | {F} |



A 5
B 4
C 4
D 3
E 3
F 2
H 2
G 1

START 0
A 1
C 1
E 1
H 2
B 2
D 3
F 4

HU (G(V, E), a) {

对顶点进行标号;

t0 = 0; l = 1;

重复执行以下步骤:

U = 仍未调度的顶点集合, 这些顶点要么没有前驱, 要么其所有前驱都已被调度;

选择一个子集 $S \subseteq U$, 使得 $|S| \leq a$, 且 S 中的顶点标签值最大;

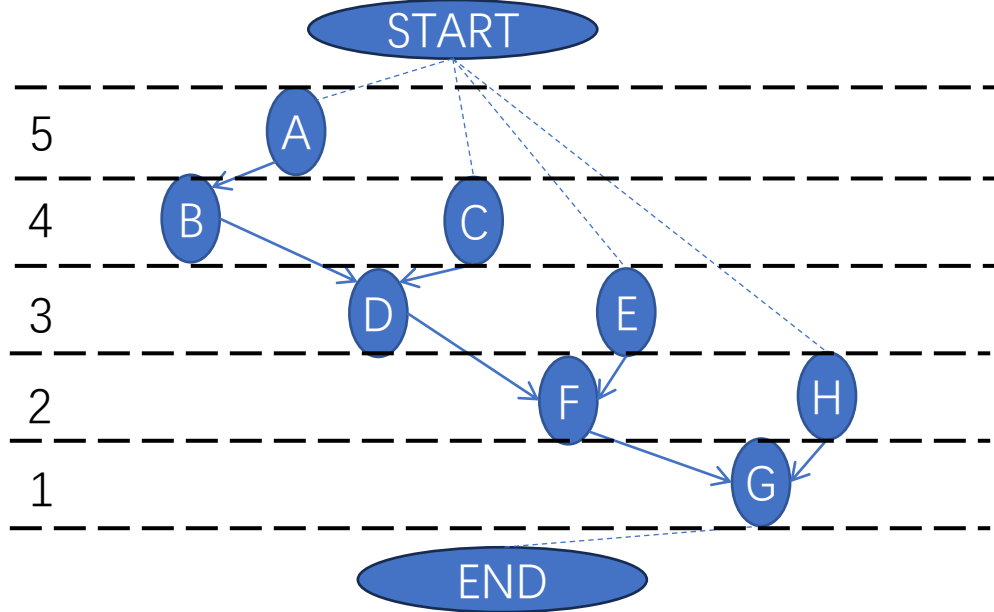
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$l = l + 1$;

直到 vn 被调度;

}

| 变量 | 当前值 |
|----|-----|
| l | 5 |
| U | {F} |
| S | {F} |



A 5
B 4
C 4
D 3
E 3
F 2
H 2
G 1

START 0
A 1
C 1
E 1
H 2
B 2
D 3
F 4

HU (G(V, E), a) {

对顶点进行标号;

t0 = 0; l = 1;

重复执行以下步骤:

U = 仍未调度的顶点集合, 这些顶点要么没有前驱, 要么其所有前驱都已被调度;

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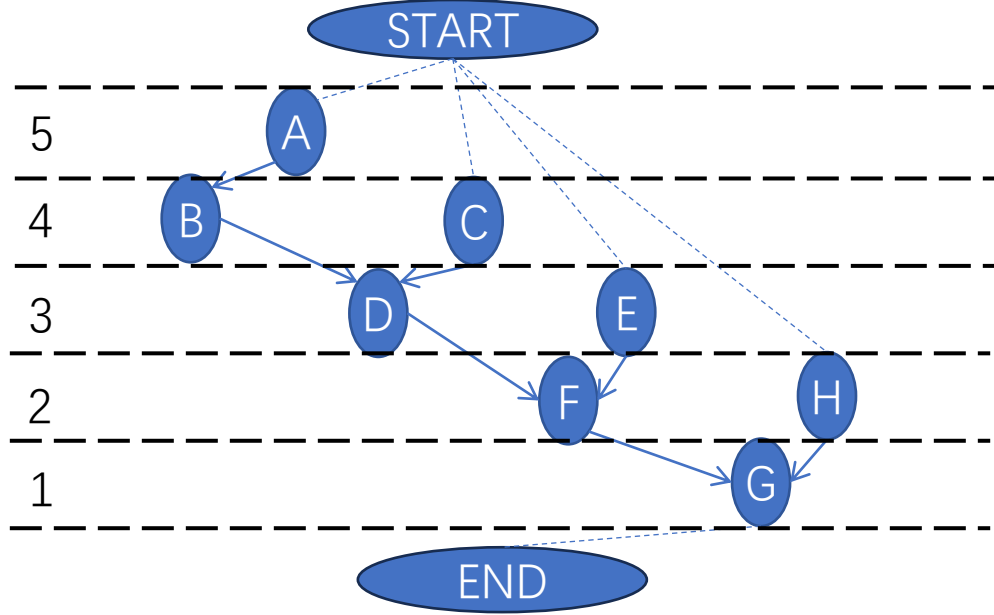
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l = l + 1;

直到 vn 被调度;

}

| 变量 | 当前值 |
|----|-----|
| l | 5 |
| U | {G} |
| S | {F} |



A 5
B 4
C 4
D 3
E 3
F 2
H 2
G 1

START 0
A 1
C 1
E 1
H 2
B 2
D 3
F 4
G 5

HU (G(V, E), a) {

对顶点进行标号;

t0 = 0; l = 1;

重复执行以下步骤:

U = 仍未调度的顶点集合, 这些顶点要么没有前驱, 要么其所有前驱都已被调度;

选择一个子集 $S \subseteq U$, 使得 $|S| \leq a$, 且 S 中的顶点标签值最大;

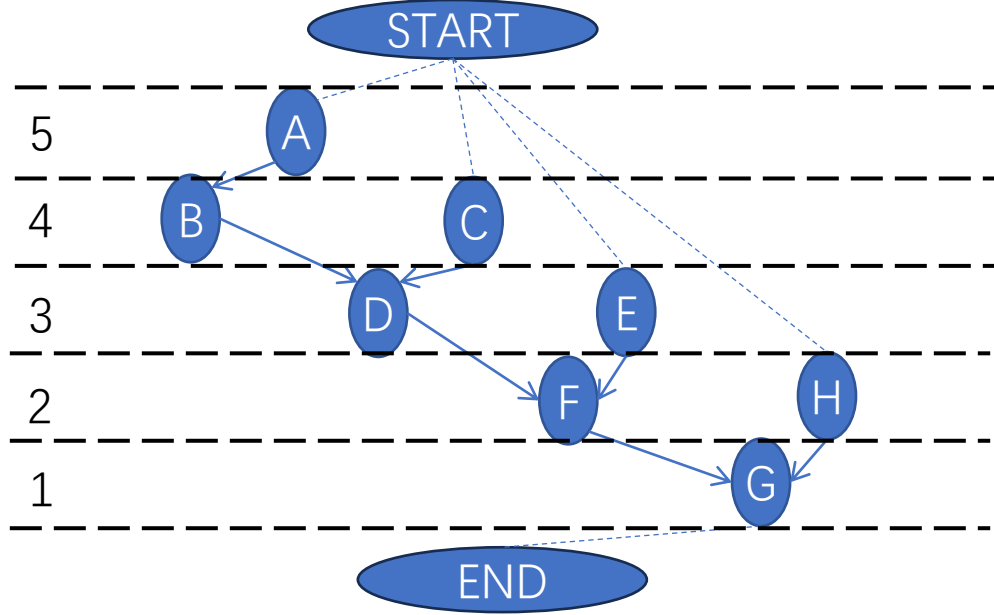
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直到 vn 被调度;

}

| 变量 | 当前值 |
|----|-----|
| l | 5 |
| U | {G} |
| S | {G} |



A 5
B 4
C 4
D 3
E 3
F 2
H 2
G 1

START 0
A 1
C 1
E 1
H 2
B 2
D 3
F 4
G 5

HU (G(V, E), a) {

对顶点进行标号;

t0 = 0; l = 1;

重复执行以下步骤:

U = 仍未调度的顶点集合, 这些顶点要么没有前驱, 要么其所有前驱都已被调度;

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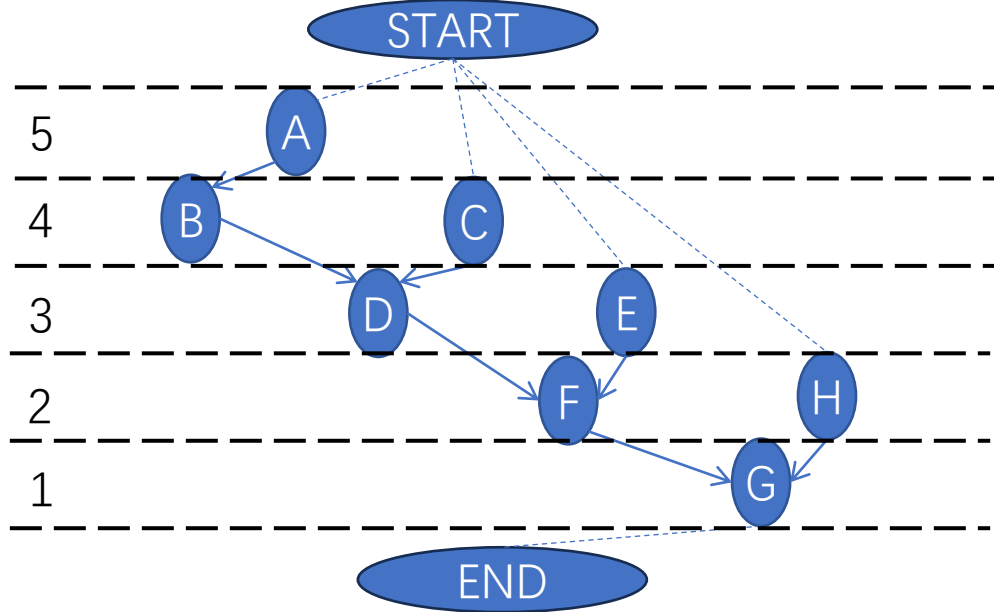
在步骤 t 处调度 S 中的所有操作, 即对所有 $v_i \in S$ 设定调度时间 $t_i = l$;

$l = l + 1$;

直到 vn 被调度;

}

| 变量 | 当前值 |
|----|-----|
| l | 6 |
| U | {G} |
| S | {G} |



A 5
B 4
C 4
D 3
E 3
F 2
H 2
G 1

START 0
A 1
C 1
E 1
H 2
B 2
D 3
F 4
G 5

HU (G(V, E), a) {

对顶点进行标号;

t0 = 0; l = 1;

重复执行以下步骤:

U = 仍未调度的顶点集合, 这些顶点要么没有前驱, 要么其所有前驱都已被调度;

选择一个子集 $S \subseteq U$, 使得 $|S| \leq a$, 且 S 中的顶点标签值最大;

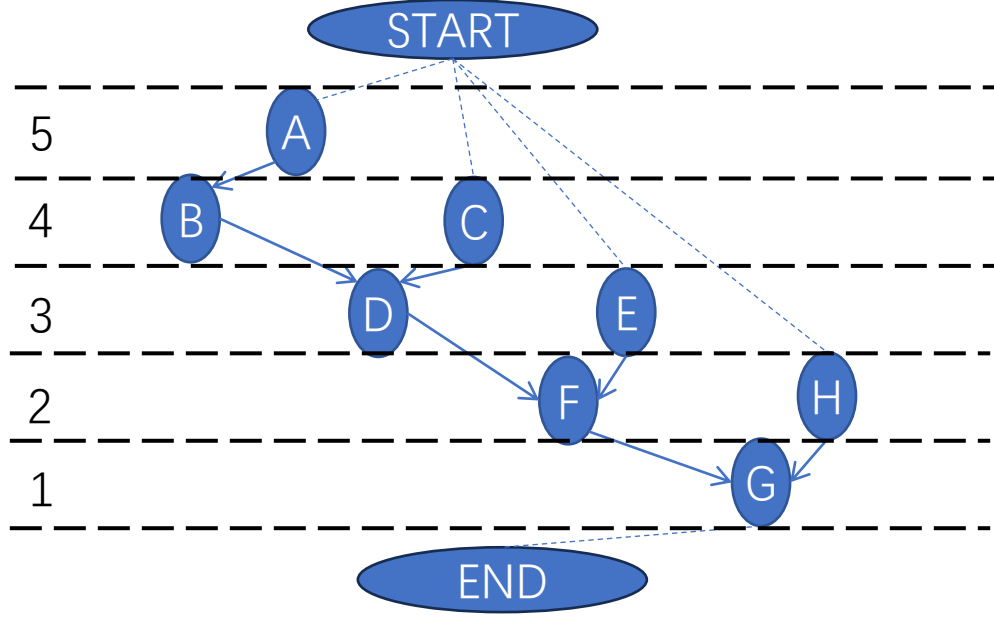
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$l = l + 1$;

直到 vn 被调度;

}

| 变量 | 当前值 |
|----|-------|
| l | 6 |
| U | {END} |
| S | {G} |



A 5
B 4
C 4
D 3
E 3
F 2
H 2
G 1

START 0
A 1
C 1
E 1
H 2
B 2
D 3
F 4
G 5
END 6

HU (G(V, E), a) {

对顶点进行标号;

t0 = 0; l = 1;

重复执行以下步骤:

U = 仍未调度的顶点集合, 这些顶点要么没有前驱, 要么其所有前驱都已被调度;

选择一个子集 $S \subseteq U$, 使得 $|S| \leq a$, 且 S 中的顶点标签值最大;

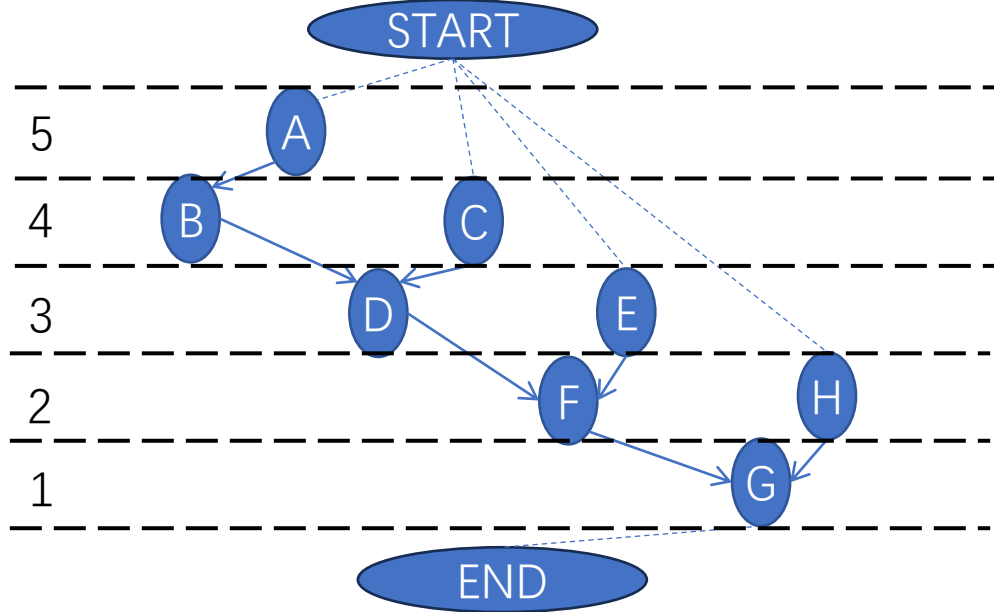
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$l = l + 1$;

直到 vn 被调度;

}

| 变量 | 当前值 |
|----|-------|
| l | 6 |
| U | {END} |
| S | {END} |



A 5
B 4
C 4
D 3
E 3
F 2
H 2
G 1

START 0
A 1
C 1
E 1
H 2
B 2
D 3
F 4
G 5
END 6

HU (G(V, E), a) {

对顶点进行标号;

t0 = 0; l = 1;

重复执行以下步骤:

U = 仍未调度的顶点集合, 这些顶点要么没有前驱, 要么其所有前驱都已被调度;

选择一个子集 $S \subseteq U$, 使得 $|S| \leq a$, 且 S 中的顶点标签值最大;

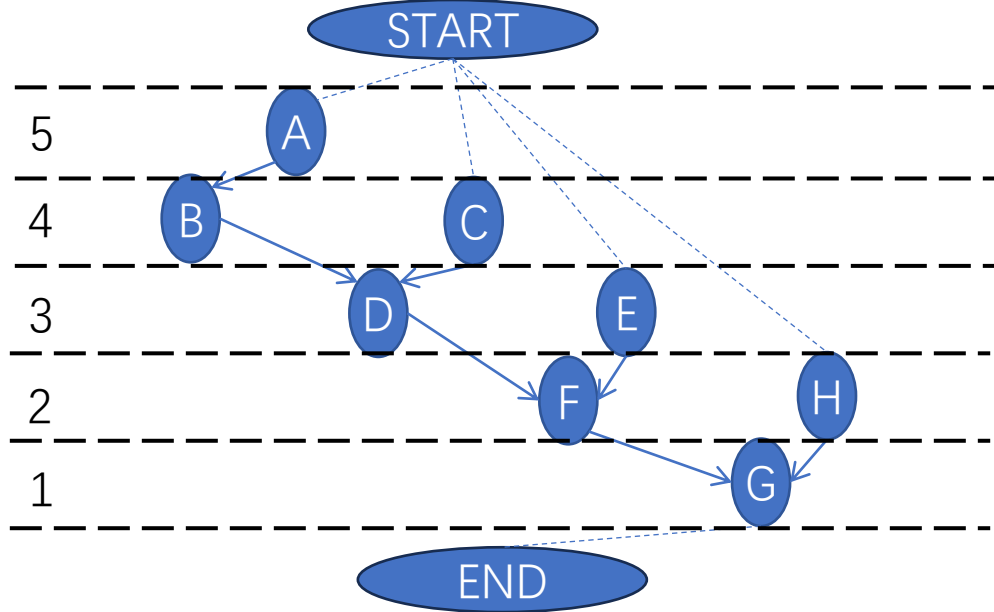
在步骤 t 处调度 S 中的所有操作, 即对所有 $v_i \in S$ 设定调度时间 $t_i = l$;

$l = l + 1$;

直到 vn 被调度;

}

| 变量 | 当前值 |
|----|-------|
| l | 7 |
| U | {END} |
| S | {END} |



A 5
B 4
C 4
D 3
E 3
F 2
H 2
G 1

START 0
A 1
C 1
E 1
H 2
B 2
D 3
F 4
G 5
END 6

HU (G(V, E), a) {

对顶点进行标号;

t0 = 0; l = 1;

重复执行以下步骤:

U = 仍未调度的顶点集合, 这些顶点要么没有前驱, 要么其所有前驱都已被调度;

选择一个子集 $S \subseteq U$, 使得 $|S| \leq a$, 且 S 中的顶点标签值最大;

在步骤 t 处调度 S 中的所有操作, 即对所有 $v_i \in S$ 设定调度时间 $t_i = l$;

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直到 vn 被调度;

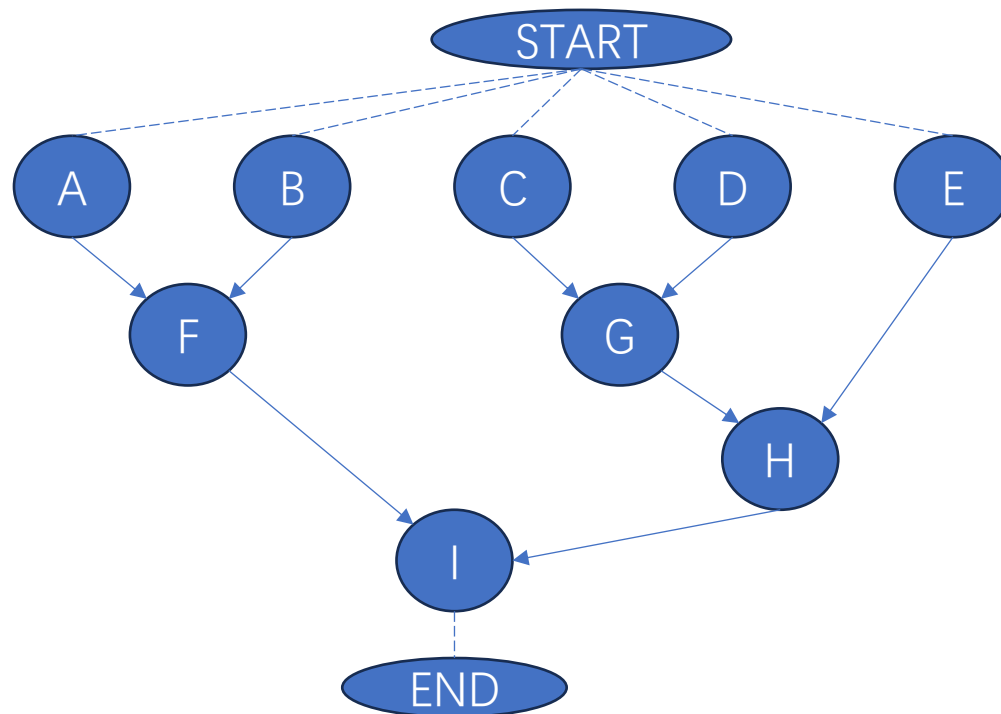
}

| 变量 | 当前值 |
|----|-------|
| l | 7 |
| U | {END} |
| S | {END} |

作业2

in-class assignment

已知输入的图如下（每个都是加法操作）：



假设我们有两个加法器，请写出HU调度下的程序的输出文件的内容，输出文件格式如下：

每一行由一个字符串和一个数字组成，以空格隔开，分别代表顶点名和开始的调度周期