

# Chapter 9: Traversable functors and contrafunctors

Sergei Winitzki

Academy by the Bay

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# Motivation for the `traverse` operation

- Consider data of type  $\text{List}^A$  and processing  $f : A \Rightarrow \text{Future}^B$
- Typically, we want to wait until the entire data set is processed
- What we need is  $\text{List}^A \Rightarrow (A \Rightarrow \text{Future}^B) \Rightarrow \text{Future}^{\text{List}^B}$
- Generalize:  $L^A \Rightarrow (A \Rightarrow F^B) \Rightarrow F^{L^B}$  for some type constructors  $F, L$
- This operation is called `traverse`
  - ▶ How to implement it: for example, a 3-element list is  $A \times A \times A$
  - ▶ Consider  $L^A \equiv A \times A \times A$ , apply map  $f$  and get  $F^B \times F^B \times F^B$
  - ▶ We will get  $F^{L^B} \equiv F^{B \times B \times B}$  if we can apply `zip` as  $F^B \times F^B \Rightarrow F^{B \times B}$
- So we need to assume that  $F$  is applicative
- In Scala, we have `Future.traverse()` that assumes  $L$  to be a sequence
  - ▶ This is the iconic example that fixes the requirements
- Questions:
  - ▶ Which functors  $L$  can have this operation?
  - ▶ Can we express `traverse` through a simpler operation?
  - ▶ What are the required laws for `traverse`?
  - ▶ What about contrafunctors or profunctors?

# Deriving the `sequence` operation

- The type signature of `traverse` is a complicated “lifting”
  - ▶ A “lifting” is always equivalent to a simpler natural transformation
- To derive it, ask: what is missing from `fmap` to do the job of `traverse`?

$$\text{fmap} : (A \Rightarrow F^B) \Rightarrow L^A \Rightarrow L^{F^B}$$

- We need  $F^{L^B}$ , but the `traverse` operation gives us  $L^{F^B}$  instead
  - ▶ What’s missing is a natural transformation `sequence` :  $L^{F^B} \Rightarrow F^{L^B}$
- The functions `traverse` and `sequence` are computationally equivalent:

$$\text{trav } f^{A \Rightarrow F^B} = \text{fmap } f \circ \text{seq}$$

A commutative diagram illustrating the relationship between `fmap`, `seq`, and `trav`. The diagram has three nodes:  $L^A$  on the left,  $L^{F^B}$  at the top center, and  $F^{L^B}$  on the right. An arrow labeled `fmap` with  $f^{A \Rightarrow F^B}$  above it points from  $L^A$  to  $L^{F^B}$ . An arrow labeled `seq` points from  $L^{F^B}$  to  $F^{L^B}$ . A long arrow labeled `trav` with  $f^{A \Rightarrow F^B}$  below it points directly from  $L^A$  to  $F^{L^B}$ .

Here  $F$  is an arbitrary applicative functor

- ▶ Keep in mind the example `Future.sequence` :  $\text{List}^{\text{Future}^X} \Rightarrow \text{Future}^{\text{List}^X}$
- ▶ Examples: `List`, all “finite” polynomial functors (see Bird et al., 2013)
- ▶ Non-traversable:  $L^A \equiv R \Rightarrow A$ ; lazy lists (“infinite streams”)

★ Note: We *cannot* have the opposite transformation  $F^{L^B} \Rightarrow L^{F^B}$

# Motivation for the laws of the `traverse` operation

- The “**law of traversals**” paper (2012) argues that `traverse` should “visit each element” of the container  $L^A$  exactly once, and evaluate each corresponding “effect”  $F^B$  exactly once; then they formulate the laws
- To derive the laws, use the “lifting” intuition for `traverse`,

$$\text{trav} : (A \Rightarrow F^B) \Rightarrow L^A \Rightarrow F^{L^B}$$

Look for “identity” and “composition” laws:

- ① “Identity” as `pure` :  $A \Rightarrow F^A$  must be lifted to `pure` :  $L^A \Rightarrow F^{L^A}$
- ② “Identity” as  $\text{id}^{A \Rightarrow A}$  with  $F^A \equiv A$  (identity functor) lifted to  $\text{id}^{L^A \Rightarrow L^A}$
- ③ “Compose”  $f : A \Rightarrow F^B$  and  $g : B \Rightarrow G^C$  to get  $h : A \Rightarrow F^{G^C}$ , where  $F, G$  are applicative; a traversal with  $h$  maps  $L^A$  to  $F^{G^{L^C}}$  and must be somehow equal to the composition of traversals with  $f$  and then with  $g$

Questions:

- Are the laws for the `sequence` operation simpler?
- Are all these laws independent?
- What functors  $L$  satisfy these laws *for all* applicative functors  $F$ ?

# Formulation of the laws for `traverse`

The

# Derivation of the laws for `sequence`

The

# Constructions of traversable functors

The

# Traversable profunctors

The



# Traversability with respect to profunctors

The

# “Folding” a traversable functor into a monoid

The

- ① Show that any traversable functor  $L$  admits a method

$$\text{consume} : (L^A \Rightarrow B) \Rightarrow L^{F^A} \Rightarrow F^B$$

for any applicative functor  $F$ . Show that `traverse` and `consume` are equivalent.