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```

1 .PRECIOUS: ./p%
3 %: p%
  ulimit -s unlimited && ./%<
5 p%: p%.cpp
  g++ -o $@ $< -std=c++17 -Wall -Wextra -Wshadow \
7     -fsanitize=address,undefined

```

## 1.2. How Did We Get Here?

### 1.2.1. Macros

Use vectorizations and math optimizations at your own peril.  
For gcc $\geq$ 9, there are `[[likely]]` and `[[unlikely]]` attributes.

Call gcc with `-fopt-info-optimized-missed-optall` for optimization info.

```
1 #define _GLIBCXX_DEBUG 1 // for debug mode
2 #define _GLIBCXX_SANITIZE_VECTOR 1 // for asan on vectors
3 #pragma GCC optimize("O3", "unroll-loops")
4 #pragma GCC optimize("fast-math")
5 #pragma GCC target("avx,avx2,abm,bmi,bmi2") // tip: `lscpu`
6 // before a loop
7 #pragma GCC unroll 16 // 0 or 1 -> no unrolling
8 #pragma GCC ivdep
```

### 1.2.2. Fast I/O

```
1 struct scanner {
2     static constexpr size_t LEN = 32 << 20;
3     char *buf, *buf_ptr, *buf_end;
4     scanner()
5         : buf(new char[LEN]), buf_ptr(buf + LEN),
6           buf_end(buf + LEN) {}
7     ~scanner() { delete[] buf; }
8     char getc() {
9         if (buf_ptr == buf_end) [[unlikely]]
10            buf_end = buf + fread_unlocked(buf, 1, LEN, stdin),
11            buf_ptr = buf;
12        return *(buf_ptr++);
13    }
14    char seek(char del) {
15        char c;
16        while ((c = getc()) < del) {}
17        return c;
18    }
19    void read(int &t) {
20        bool neg = false;
21        char c = seek('-');
22        if (c == '-') neg = true, t = 0;
23        else t = c ^ '0';
24        while ((c = getc()) >= '0') t = t * 10 + (c ^ '0');
25        if (neg) t = -t;
26    }
27 };
28 struct printer {
29     static constexpr size_t CPI = 21, LEN = 32 << 20;
30     char *buf, *buf_ptr, *buf_end, *tbuf;
31     char *int_buf, *int_buf_end;
32     printer()
33         : buf(new char[LEN]), buf_ptr(buf),
34           buf_end(buf + LEN), int_buf(new char[CPI + 1]()),
35           int_buf_end(int_buf + CPI - 1) {}
36     ~printer() {
37         flush();
38         delete[] buf, delete[] int_buf;
39     }
40     void flush() {
41         fwrite_unlocked(buf, 1, buf_ptr - buf, stdout);
42         buf_ptr = buf;
43     }
44     void write(const char &c) {
45         *buf_ptr = c;
46         if (++buf_ptr == buf_end) [[unlikely]]
47             flush();
48     }
49     void write(const char *s) {
50         for (; *s != '\0'; ++s) write(*s);
51     }
52     void write(int x) {
53         if (x < 0) write('-'), x = -x;
54         if (x == 0) [[unlikely]]
55             return write('0');
56         for (tbuf = int_buf_end; x != 0; --tbuf, x /= 10)
57             *tbuf = '0' + char(x % 10);
58         write(++tbuf);
59     }
60 };
```

### Kotlin

```
1 import java.io.*
2 import java.util.*
3
4 @JvmField val cin = System.`in`.bufferedReader()
5 @JvmField val cout = PrintWriter(System.out, false)
6 @JvmField var tokenizer: StringTokenizer = StringTokenizer("")
7 fun nextLine() = cin.readLine()!!
8 fun read(): String {
9     while(!tokenizer.hasMoreTokens())
10         tokenizer = StringTokenizer(nextLine())
11     return tokenizer.nextToken()
12 }
```

```
// example
15 fun main() {
16     val n = read().toInt()
17     val a = DoubleArray(n) { read().toDouble() }
18     cout.println("omg hi")
19     cout.flush()
20 }
```

### 1.2.3. constexpr

Some default limits in gcc (7.x - trunk):

- constexpr recursion depth: 512
- constexpr loop iteration per function: 262144
- constexpr operation count per function: 33554432
- template recursion depth: 900 (gcc *might* segfault first)

```
1 constexpr array<int, 10> fibonacci{[] {
2     array<int, 10> a{};
3     a[0] = a[1] = 1;
4     for (int i = 2; i < 10; i++) a[i] = a[i - 1] + a[i - 2];
5     return a;
6 }}();
7 static_assert(fibonacci[9] == 55, "CE");
8
9 template <typename F, typename INT, INT... S>
10 constexpr void for_constexpr(integer_sequence<INT, S...>,
11                             F &&func) {
12     int _[] = {(func(integral_constant<INT, S>{}), 0)...};
13 }
14 // example
15 template <typename... T> void print_tuple(tuple<T...> t) {
16     for_constexpr(make_index_sequence<sizeof...(T)>{}),
17                 [&](auto i) { cout << get<i>(t) << '\n'; };
18 }
```

### 1.2.4. Bump Allocator

```
1 // global bump allocator
2 char mem[256 << 20]; // 256 MB
3 size_t rsp = sizeof mem;
4 void *operator new(size_t s) {
5     assert(s < rsp); // MLE
6     return (void *)&mem[rsp -= s];
7 }
8 void operator delete(void *) {}
9
10 // bump allocator for STL / pbds containers
11 char mem[256 << 20];
12 size_t rsp = sizeof mem;
13 template <typename T> struct bump {
14     typedef T value_type;
15     bump() {}
16     template <typename U> bump(U, ...) {}
17     T *allocate(size_t n) {
18         rsp -= n * sizeof(T);
19         rsp &= 0 - alignof(T);
20         return (T *)&(mem + rsp);
21     }
22     void deallocate(T *, size_t n) {}
23 };
```

### 1.3. Tools

#### 1.3.1. Floating Point Binary Search

```
1 union di {
2     double d;
3     ull i;
4 };
5 bool check(double);
6 // binary search in [L, R] with relative error 2^-eps
7 double binary_search(double L, double R, int eps) {
8     di l = {L}, r = {R}, m;
9     while (r.i - l.i > 1LL << (52 - eps)) {
10         m.i = (l.i + r.i) >> 1;
11         if (check(m.d)) r = m;
12         else l = m;
13     }
14     return l.d;
15 }
```

#### 1.3.2. SplitMix64

```
1 using ull = unsigned long long;
2 inline ull splitmix64(ull x) {
3     // change to `static ull x = SEED;` for DRBG
4     ull z = (x += 0x9E3779B97F4A7C15);
5     z = (z ^ (z >> 30)) * 0xBF58476D1CE4E5B9;
6     z = (z ^ (z >> 27)) * 0x94D049BB133111EB;
7     return z ^ (z >> 31);
8 }
```

### 1.3.3. <random>

```

1 #ifdef __unix__
  random_device rd;
3 mt19937_64 RNG(rd());
  #else
5 const auto SEED = chrono::high_resolution_clock::now()
                      .time_since_epoch()
                      .count();
7 mt19937_64 RNG(SEED);
9 #endif
// random uint_fast64_t: RNG();
11 // uniform random of type T (int, double, ...) in [l, r]:
// uniform_int_distribution<T> dist(l, r); dist(RNG);

```

### 1.3.4. x86 Stack Hack

```

1 constexpr size_t size = 200 << 20; // 200MiB
int main() {
3   register long rsp asm("rsp");
   char *buf = new char[size];
5   asm("movq %0, %%rsp\n" :: "r"(buf + size));
   // do stuff
7   asm("movq %0, %%rsp\n" :: "r"(rsp));
   delete[] buf;
9 }

```

## 1.4. Algorithms

### 1.4.1. Bit Hacks

```

1 // next permutation of x as a bit sequence
ull next_bits_permutation(ull x) {
3   ull c = __builtin_ctzll(x), r = x + (1ULL << c);
   return (r ^ x) >> (c + 2) | r;
5 }
// iterate over all (proper) subsets of bitset s
7 void subsets(ull s) {
   for (ull x = s; x;) { --x &= s; /* do stuff */ }
9 }

```

### 1.4.2. Aliens Trick

```

1 // min dp[i] value and its i (smallest one)
pll get_dp(int cost);
3 ll aliens(int k, int l, int r) {
   while (l != r) {
5     int m = (l + r) / 2;
     auto [f, s] = get_dp(m);
7     if (s == k) return f - m * k;
     if (s < k) r = m;
     else l = m + 1;
11  }
   return get_dp(l).first - l * k;
}

```

### 1.4.3. Hilbert Curve

```

1 ll hilbert(ll n, int x, int y) {
   ll res = 0;
3   for (ll s = n; s /= 2;) {
     int rx = !(x & s), ry = !(y & s);
5     res += s * s * ((3 * rx) ^ ry);
     if (ry == 0) {
7       if (rx == 1) x = s - 1 - x, y = s - 1 - y;
       swap(x, y);
9     }
   }
11  return res;
}

```

### 1.4.4. Infinite Grid Knight Distance

```

1 ll get_dist(ll dx, ll dy) {
   if (++(dx = abs(dx)) > ++(dy = abs(dy))) swap(dx, dy);
3   if (dx == 1 && dy == 2) return 3;
   if (dx == 3 && dy == 3) return 4;
5   ll lb = max(dy / 2, (dx + dy) / 3);
   return ((dx ^ dy ^ lb) & 1) ? ++lb : lb;
7 }

```

### 1.4.5. Poker Hand

```

1 using namespace std;

3 struct hand {
   static constexpr auto rk = [] {
5     array<int, 256> x{};
     auto s = "23456789TJQKACDHS";
7     for (int i = 0; i < 17; i++) x[s[i]] = i % 13;
     return x;
9   }();
   vector<pair<int, int>> v;
   vector<int> cnt, vf, vs;
11  int type;
   hand() : cnt(4), type(0) {}
13  void add_card(char suit, char rank) {
     ++cnt[rk[suit]];
15     for (auto &[f, s] : v)
       if (s == rk[rank]) return ++f, void();
     v.emplace_back(1, rk[rank]);
17   }
19  void process() {
     sort(v.rbegin(), v.rend());
21     for (auto [f, s] : v) vf.push_back(f), vs.push_back(s);
     bool str = 0, flu = find(all(cnt), 5) != cnt.end();
23     if ((str = v.size() == 5))
       for (int i = 1; i < 5; i++)
25         if (vs[i] != vs[i - 1] + 1) str = 0;
     if (vs == vector<int>{12, 3, 2, 1, 0})
27         str = 1, vs = {3, 2, 1, 0, -1};
     if (str && flu) type = 9;
29     else if (vf[0] == 4) type = 8;
     else if (vf[0] == 3 && vf[1] == 2) type = 7;
31     else if (str || flu) type = 5 + flu;
     else if (vf[0] == 3) type = 4;
33     else if (vf[0] == 2) type = 2 + (vf[1] == 2);
     else type = 1;
35   }
37  bool operator<(const hand &b) const {
     return make_tuple(type, vf, vs) <
39         make_tuple(b.type, b.vf, b.vs);
41  };
}

```

### 1.4.6. Longest Increasing Subsequence

```

1 template <class I> vi lis(const vector<I> &S) {
   if (S.empty()) return {};
3   vi prev(sz(S));
   typedef pair<I, int> p;
5   vector<p> res;
   rep(i, 0, sz(S)) {
7     // change 0 -> i for longest non-decreasing subsequence
     auto it = lower_bound(all(res), p{S[i], 0});
9     if (it == res.end())
       res.emplace_back(), it = res.end() - 1;
11    *it = {S[i], i};
     prev[i] = it == res.begin() ? 0 : (it - 1)->second;
13  }
   int L = sz(res), cur = res.back().second;
15  vi ans(L);
   while (L--) ans[L] = cur, cur = prev[cur];
17  return ans;
}

```

### 1.4.7. Mo's Algorithm on Tree

```

1 void MoAlgoOnTree() {
   Dfs(0, -1);
3   vector<int> euler(tk);
   for (int i = 0; i < n; ++i) {
5     euler[tin[i]] = i;
     euler[tout[i]] = i;
7   }
   vector<int> l(q), r(q), qr(q), sp(q, -1);
9   for (int i = 0; i < q; ++i) {
     if (tin[u[i]] > tin[v[i]]) swap(u[i], v[i]);
11    int z = GetLCA(u[i], v[i]);
     sp[i] = z[i];
13    if (z == u) l[i] = tin[u[i]], r[i] = tin[v[i]];
     else l[i] = tout[u[i]], r[i] = tin[v[i]];
15    qr[i] = i;
   }
17  sort(qr.begin(), qr.end(), [&](int i, int j) {
     if (l[i] / kB == l[j] / kB) return r[i] < r[j];
19    return l[i] / kB < l[j] / kB;
  });
   vector<bool> used(n);
21  // Add(v): add/remove v to/from the path based on used[v]
   for (int i = 0, tl = 0, tr = -1; i < q; ++i) {
23

```

```

25 while (tl < l[qr[i]]) Add(euler[tl++]);
   while (tl > l[qr[i]]) Add(euler[--tl]);
27 while (tr > r[qr[i]]) Add(euler[tr--]);
   while (tr < r[qr[i]]) Add(euler[++tr]);
   // add/remove LCA(u, v) if necessary
29 }
}

```

## 2. Data Structures

### 2.1. GNU PBDS

```

1 #include <ext/pb_ds/assoc_container.hpp>
  #include <ext/pb_ds/priority_queue.hpp>
3 #include <ext/pb_ds/tree_policy.hpp>
  using namespace __gnu_pbds;

   // most std::map + order_of_key, find_by_order, split, join
7 template <typename T, typename U = null_type>
   using ordered_map = tree<T, U, std::less<>, rb_tree_tag,
9   tree_order_statistics_node_update>;
   // useful tags: rb_tree_tag, splay_tree_tag

11 template <typename T> struct myhash {
13   size_t operator()(T x) const; // splitmix, bswap(x*R), ...
};
   // most of std::unordered_map, but faster (needs good hash)
15 template <typename T, typename U = null_type>
17 using hash_table = gp_hash_table<T, U, myhash<T>>;

19 // most std::priority_queue + modify, erase, split, join
   using heap = priority_queue<int, std::less<>>;
21 // useful tags: pairing_heap_tag, binary_heap_tag,
   // (rc)?binomial_heap_tag, thin_heap_tag

```

### 2.2. Segment Tree (ZKW)

```

1 struct segtree {
   using T = int;
3   T f(T a, T b) { return a + b; } // any monoid operation
   static constexpr T ID = 0; // identity element
5   int n;
   vector<T> v;
7   segtree(int n_) : n(n_), v(2 * n, ID) {}
   segtree(vector<T> &a) : n(a.size()), v(2 * n, ID) {
9     copy_n(a.begin(), n, v.begin() + n);
     for (int i = n - 1; i > 0; i--)
11       v[i] = f(v[i * 2], v[i * 2 + 1]);
   }
13   void update(int i, T x) {
     for (v[i += n] = x; i /= 2;)
15       v[i] = f(v[i * 2], v[i * 2 + 1]);
   }
17   T query(int l, int r) {
     T tl = ID, tr = ID;
19     for (l += n, r += n; l < r; l /= 2, r /= 2) {
       if (l & 1) tl = f(tl, v[l++]);
21       if (r & 1) tr = f(v[--r], tr);
     }
23     return f(tl, tr);
   }
25 };

```

### 2.3. Line Container

```

1 struct Line {
   mutable ll k, m, p;
3   bool operator<(const Line &o) const { return k < o.k; }
   bool operator<(ll x) const { return p < x; }
5 };
   // add: line y=kx+m, query: maximum y of given x
7 struct LineContainer : multiset<Line, less<>> {
   // (for doubles, use inf = 1/.0, div(a,b) = a/b)
9   static const ll inf = LLONG_MAX;
   ll div(ll a, ll b) { // floored division
11     return a / b - ((a ^ b) < 0 && a % b);
   }
13   bool isect(iterator x, iterator y) {
     if (y == end()) return x->p = inf, 0;
15     if (x->k == y->k) x->p = x->m > y->m ? inf : -inf;
     else x->p = div(y->m - x->m, x->k - y->k);
17     return x->p >= y->p;
   }
19   void add(ll k, ll m) {
     auto z = insert({k, m, 0}), y = z++, x = y;
21     while (isect(y, z)) z = erase(z);
     if (x != begin() && isect(--x, y))
23       isect(x, y = erase(y));
     while ((y = x) != begin() && (--x)->p >= y->p)

```

```

25   isect(x, erase(y));
   }
27   ll query(ll x) {
     assert(!empty());
29     auto l = *lower_bound(x);
     return l.k * x + l.m;
31   }
};

```

### 2.4. Li-Chao Tree

```

1 constexpr ll MAXN = 2e5, INF = 2e18;
   struct Line {
3     ll m, b;
     Line() : m(0), b(-INF) {}
5     Line(ll _m, ll _b) : m(_m), b(_b) {}
     ll operator()(ll x) const { return m * x + b; }
7   };
   struct Li_Chao {
9     Line a[MAXN * 4];
     void insert(Line seg, int l, int r, int v = 1) {
11       if (l == r) {
         if (seg(l) > a[v](l)) a[v] = seg;
13       }
         return;
15       int mid = (l + r) >> 1;
       if (a[v].m > seg.m) swap(a[v], seg);
17       if (a[v](mid) < seg(mid)) {
         swap(a[v], seg);
19       }
       insert(seg, l, mid, v << 1);
       insert(seg, mid + 1, r, v << 1 | 1);
21   }
   ll query(int x, int l, int r, int v = 1) {
23     if (l == r) return a[v](x);
     int mid = (l + r) >> 1;
25     if (x <= mid)
       return max(a[v](x), query(x, l, mid, v << 1));
27     else
       return max(a[v](x), query(x, mid + 1, r, v << 1 | 1));
29   }
};

```

### 2.5. Heavy-Light Decomposition

```

1 template <bool VALS_EDGES> struct HLD {
   int N, tim = 0;
3   vector<vi> adj;
   vi par, siz, depth, rt, pos;
5   Node *tree;
   HLD(vector<vi> adj_)
7     : N(sz(adj_)), adj(adj_), par(N, -1), siz(N, 1),
       depth(N), rt(N), pos(N), tree(new Node(0, N)) {}
9   dfsSz(0);
   dfsHld(0);
11 };
   void dfsSz(int v) {
13   if (par[v] != -1)
     adj[v].erase(find(all(adj[v]), par[v]));
15   for (int &u : adj[v]) {
     par[u] = v, depth[u] = depth[v] + 1;
17     dfsSz(u);
     siz[v] += siz[u];
19     if (siz[u] > siz[adj[v][0]]) swap(u, adj[v][0]);
   }
21 };
   void dfsHld(int v) {
23   pos[v] = tim++;
   for (int u : adj[v]) {
25     rt[u] = (u == adj[v][0] ? rt[v] : u);
     dfsHld(u);
27   }
   }
29 template <class B> void process(int u, int v, B op) {
   for (; rt[u] != rt[v]; v = par[rt[v]]) {
31     if (depth[rt[u]] > depth[rt[v]]) swap(u, v);
     op(pos[rt[v]], pos[v] + 1);
33   }
   if (depth[u] > depth[v]) swap(u, v);
35   op(pos[u] + VALS_EDGES, pos[v] + 1);
   }
37 void modifyPath(int u, int v, int val) {
   process(u, v,
39     [&](int l, int r) { tree->add(l, r, val); });
   }
41 int queryPath(int u,
   int v) { // Modify depending on problem
43   int res = -1e9;
   process(u, v, [&](int l, int r) {
45     res = max(res, tree->query(l, r));
   });
}

```

```

47     return res;
48 }
49 int querySubtree(int v) { // modifySubtree is similar
50     return tree->query(pos[v] + VALS_EDGES,
51                        pos[v] + siz[v]);
52 }
53 };

```

## 2.6. Wavelet Matrix

```

1  #pragma GCC target("popcnt,bmi2")
2  #include <immintrin.h>
3
4  // T is unsigned. You might want to compress values first
5  template <typename T> struct wavelet_matrix {
6      static_assert(is_unsigned_v<T>, "only unsigned T");
7      struct bit_vector {
8          static constexpr uint W = 64;
9          uint n, cnt0;
10         vector<ull> bits;
11         vector<uint> sum;
12         bit_vector(uint n_)
13             : n(n_), bits(n / W + 1), sum(n / W + 1) {}
14         void build() {
15             for (uint j = 0; j != n / W; ++j)
16                 sum[j + 1] = sum[j] + _mm_popcnt_u64(bits[j]);
17             cnt0 = rank0(n);
18         }
19         void set_bit(uint i) { bits[i / W] |= 1ULL << i % W; }
20         bool operator[](uint i) const {
21             return !(bits[i / W] & 1ULL << i % W);
22         }
23         uint rank1(uint i) const {
24             return sum[i / W] +
25                 _mm_popcnt_u64(_bzhi_u64(bits[i / W], i % W));
26         }
27         uint rank0(uint i) const { return i - rank1(i); }
28     };
29     uint n, lg;
30     vector<bit_vector> b;
31     wavelet_matrix(const vector<T> &a) : n(a.size()) {
32         lg =
33             __lg(max(*max_element(a.begin(), a.end()), T(1))) + 1;
34         b.assign(lg, n);
35         vector<T> cur = a, nxt(n);
36         for (int h = lg; h--;) {
37             for (uint i = 0; i < n; ++i)
38                 if (cur[i] & (T(1) << h)) b[h].set_bit(i);
39             b[h].build();
40             int il = 0, ir = b[h].cnt0;
41             for (uint i = 0; i < n; ++i)
42                 nxt[(b[h][i] ? ir : il)++] = cur[i];
43             swap(cur, nxt);
44         }
45     }
46     T operator[](uint i) const {
47         T res = 0;
48         for (int h = lg; h--;)
49             if (b[h][i])
50                 i += b[h].cnt0 - b[h].rank0(i), res |= T(1) << h;
51         else i = b[h].rank0(i);
52         return res;
53     }
54     // query k-th smallest (0-based) in a[l, r)
55     T kth(uint l, uint r, uint k) const {
56         T res = 0;
57         for (int h = lg; h--;) {
58             uint tl = b[h].rank0(l), tr = b[h].rank0(r);
59             if (k >= tr - tl) {
60                 k -= tr - tl;
61                 l += b[h].cnt0 - tl;
62                 r += b[h].cnt0 - tr;
63                 res |= T(1) << h;
64             } else l = tl, r = tr;
65         }
66         return res;
67     }
68     // count of i in [l, r) with a[i] < u
69     uint count(uint l, uint r, T u) const {
70         if (u >= T(1) << lg) return r - l;
71         uint res = 0;
72         for (int h = lg; h--;) {
73             uint tl = b[h].rank0(l), tr = b[h].rank0(r);
74             if (u & (T(1) << h)) {
75                 l += b[h].cnt0 - tl;
76                 r += b[h].cnt0 - tr;
77                 res += tr - tl;
78             } else l = tl, r = tr;
79         }
80         return res;
81     }

```

```

81 }
82 };

```

## 2.7. Link-Cut Tree

```

1  const int MXN = 100005;
2  const int MEM = 100005;
3
4  struct Splay {
5      static Splay nil, mem[MEM], *pmem;
6      Splay *ch[2], *f;
7      int val, rev, size;
8      Splay() : val(-1), rev(0), size(0) {
9          f = ch[0] = ch[1] = &nil;
10     }
11     Splay(int _val) : val(_val), rev(0), size(1) {
12         f = ch[0] = ch[1] = &nil;
13     }
14     bool isr() {
15         return f->ch[0] != this && f->ch[1] != this;
16     }
17     int dir() { return f->ch[0] == this ? 0 : 1; }
18     void setCh(Splay *c, int d) {
19         ch[d] = c;
20         if (c != &nil) c->f = this;
21         pull();
22     }
23     void push() {
24         if (rev) {
25             swap(ch[0], ch[1]);
26             if (ch[0] != &nil) ch[0]->rev ^= 1;
27             if (ch[1] != &nil) ch[1]->rev ^= 1;
28             rev = 0;
29         }
30     }
31     void pull() {
32         size = ch[0]->size + ch[1]->size + 1;
33         if (ch[0] != &nil) ch[0]->f = this;
34         if (ch[1] != &nil) ch[1]->f = this;
35     }
36 } Splay::nil, Splay::mem[MEM], *Splay::pmem = Splay::mem;
37 Splay *nil = &Splay::nil;
38
39 void rotate(Splay *x) {
40     Splay *p = x->f;
41     int d = x->dir();
42     if (!p->isr()) p->f->setCh(x, p->dir());
43     else x->f = p->f;
44     p->setCh(x->ch[!d], d);
45     x->setCh(p, !d);
46     p->pull();
47     x->pull();
48 }
49
50 vector<Splay*> splayVec;
51 void splay(Splay *x) {
52     splayVec.clear();
53     for (Splay *q = x; q; q = q->f) {
54         splayVec.push_back(q);
55         if (q->isr()) break;
56     }
57     reverse(begin(splayVec), end(splayVec));
58     for (auto it : splayVec) it->push();
59     while (!x->isr()) {
60         if (x->f->isr()) rotate(x);
61         else if (x->dir() == x->f->dir())
62             rotate(x->f), rotate(x);
63         else rotate(x), rotate(x);
64     }
65 }
66
67 Splay *access(Splay *x) {
68     Splay *q = nil;
69     for (; x != nil; x = x->f) {
70         splay(x);
71         x->setCh(q, 1);
72         q = x;
73     }
74     return q;
75 }
76
77 void evert(Splay *x) {
78     access(x);
79     splay(x);
80     x->rev ^= 1;
81     x->push();
82     x->pull();
83 }
84
85 void link(Splay *x, Splay *y) {
86     // evert(x);
87     access(x);

```



```

splay(x);
evert(y);
x->setCh(y, 1);
}
void cut(Splay *x, Splay *y) {
    // evert(x);
    access(y);
    splay(y);
    y->push();
    y->ch[0] = y->ch[0]->f = nil;
}

int N, Q;
Splay *vt[MXN];

int ask(Splay *x, Splay *y) {
    access(x);
    access(y);
    splay(x);
    int res = x->f->val;
    if (res == -1) res = x->val;
    return res;
}

int main(int argc, char **argv) {
    scanf("%d%d", &N, &Q);
    for (int i = 1; i <= N; i++)
        vt[i] = new (Splay::pmem++) Splay(i);
    while (Q--) {
        char cmd[105];
        int u, v;
        scanf("%s", cmd);
        if (cmd[1] == 'i') {
            scanf("%d%d", &u, &v);
            link(vt[u], vt[v]);
        } else if (cmd[0] == 'c') {
            scanf("%d", &v);
            cut(vt[1], vt[v]);
        } else {
            scanf("%d%d", &u, &v);
            int res = ask(vt[u], vt[v]);
            printf("%d\n", res);
        }
    }
}

```

### 3. Graph

#### 3.1. Modeling

- Maximum/Minimum flow with lower bound / Circulation problem
  - Construct super source  $S$  and sink  $T$ .
  - For each edge  $(x, y, l, u)$ , connect  $x \rightarrow y$  with capacity  $u - l$ .
  - For each vertex  $v$ , denote by  $in(v)$  the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
  - If  $in(v) > 0$ , connect  $S \rightarrow v$  with capacity  $in(v)$ , otherwise, connect  $v \rightarrow T$  with capacity  $-in(v)$ .
    - To maximize, connect  $t \rightarrow s$  with capacity  $\infty$  (skip this in circulation problem), and let  $f$  be the maximum flow from  $S$  to  $T$ . If  $f \neq \sum_{v \in V, in(v) > 0} in(v)$ , there's no solution. Otherwise, the maximum flow from  $s$  to  $t$  is the answer.
    - To minimize, let  $f$  be the maximum flow from  $S$  to  $T$ . Connect  $t \rightarrow s$  with capacity  $\infty$  and let the flow from  $S$  to  $T$  be  $f'$ . If  $f + f' \neq \sum_{v \in V, in(v) > 0} in(v)$ , there's no solution. Otherwise,  $f'$  is the answer.
  - The solution of each edge  $e$  is  $l_e + f_e$ , where  $f_e$  corresponds to the flow of edge  $e$  on the graph.
- Construct minimum vertex cover from maximum matching  $M$  on bipartite graph  $(X, Y)$ 
  - Redirect every edge:  $y \rightarrow x$  if  $(x, y) \in M$ ,  $x \rightarrow y$  otherwise.
  - DFS from unmatched vertices in  $X$ .
  - $x \in X$  is chosen iff  $x$  is unvisited.
  - $y \in Y$  is chosen iff  $y$  is visited.
- Minimum cost cyclic flow
  - Construct super source  $S$  and sink  $T$ .
  - For each edge  $(x, y, c)$ , connect  $x \rightarrow y$  with  $(cost, cap) = (c, 1)$  if  $c > 0$ , otherwise connect  $y \rightarrow x$  with  $(cost, cap) = (-c, 1)$ .
  - For each edge with  $c < 0$ , sum these cost as  $K$ , then increase  $d(y)$  by 1, decrease  $d(x)$  by 1.
  - For each vertex  $v$  with  $d(v) > 0$ , connect  $S \rightarrow v$  with  $(cost, cap) = (0, d(v))$ .
  - For each vertex  $v$  with  $d(v) < 0$ , connect  $v \rightarrow T$  with  $(cost, cap) = (0, -d(v))$ .
  - Flow from  $S$  to  $T$ , the answer is the cost of the flow  $C + K$ .
- Maximum density induced subgraph
  - Binary search on answer, suppose we're checking answer  $T$ .
  - Construct a max flow model, let  $K$  be the sum of all weights.
  - Connect source  $s \rightarrow v$ ,  $v \in G$  with capacity  $K$ .
  - For each edge  $(u, v, w)$  in  $G$ , connect  $u \rightarrow v$  and  $v \rightarrow u$  with capacity  $w$ .

- For  $v \in G$ , connect it with sink  $v \rightarrow t$  with capacity  $K + 2T - (\sum_{e \in E(v)} w(e)) - 2w(v)$ .
  - $T$  is a valid answer if the maximum flow  $f < K|V|$ .
- Minimum weight edge cover
    - For each  $v \in V$  create a copy  $v'$ , and connect  $u \rightarrow v'$  with weight  $w(u, v)$ .
    - Create edge  $v \rightarrow v'$  with weight  $2\mu(v)$ , where  $\mu(v)$  is the cost of the cheapest edge incident to  $v$ .
    - Find the minimum weight perfect matching on  $G'$ .
  - Project selection problem
    - If  $p_v > 0$ , create edge  $(s, v)$  with capacity  $p_v$ ; otherwise, create edge  $(v, t)$  with capacity  $-p_v$ .
    - Create edge  $(u, v)$  with capacity  $w$  with  $w$  being the cost of choosing  $u$  without choosing  $v$ .
    - The mincut is equivalent to the maximum profit of a subset of projects.
  - 0/1 quadratic programming

$$\sum_x c_x x + \sum_y c_y \bar{y} + \sum_{xy} c_{xy} x \bar{y} + \sum_{xyx'y'} c_{xyx'y'} (x \bar{y} + x' \bar{y}')$$

can be minimized by the mincut of the following graph:

- Create edge  $(x, t)$  with capacity  $c_x$  and create edge  $(s, y)$  with capacity  $c_y$ .
- Create edge  $(x, y)$  with capacity  $c_{xy}$ .
- Create edge  $(x, y)$  and edge  $(x', y')$  with capacity  $c_{xyx'y'}$ .

#### 3.2. Matching/Flows

##### 3.2.1. Dinic's Algorithm

```

1 struct Dinic {
2     struct edge {
3         int to, cap, flow, rev;
4     };
5     static constexpr int MAXN = 1000, MAXF = 1e9;
6     vector<edge> v[MAXN];
7     int top[MAXN], deep[MAXN], side[MAXN], s, t;
8     void make_edge(int s, int t, int cap) {
9         v[s].push_back({t, cap, 0, (int)v[t].size()});
10        v[t].push_back({s, 0, 0, (int)v[s].size() - 1});
11    }
12    int dfs(int a, int flow) {
13        if (a == t || !flow) return flow;
14        for (int &i = top[a]; i < v[a].size(); i++) {
15            edge &e = v[a][i];
16            if (deep[a] + 1 == deep[e.to] && e.cap - e.flow) {
17                int x = dfs(e.to, min(e.cap - e.flow, flow));
18                if (x) {
19                    e.flow += x, v[e.to][e.rev].flow -= x;
20                    return x;
21                }
22            }
23        }
24        deep[a] = -1;
25        return 0;
26    }
27    bool bfs() {
28        queue<int> q;
29        fill_n(deep, MAXN, 0);
30        q.push(s), deep[s] = 1;
31        int tmp;
32        while (!q.empty()) {
33            tmp = q.front(), q.pop();
34            for (edge &e : v[tmp])
35                if (!deep[e.to] && e.cap != e.flow)
36                    deep[e.to] = deep[tmp] + 1, q.push(e.to);
37        }
38        return deep[t];
39    }
40    int max_flow(int _s, int _t) {
41        s = _s, t = _t;
42        int flow = 0, tflow;
43        while (bfs()) {
44            fill_n(top, MAXN, 0);
45            while ((tflow = dfs(s, MAXF))) flow += tflow;
46        }
47        return flow;
48    }
49    void reset() {
50        fill_n(side, MAXN, 0);
51        for (auto &i : v) i.clear();
52    }
53 };

```

### 3.2.2. Minimum Cost Flow

```

1 struct MCF {
2     struct edge {
3         ll to, from, cap, flow, cost, rev;
4     } * fromE[MAXN];
5     vector<edge> v[MAXN];
6     ll n, s, t, flows[MAXN], dis[MAXN], pi[MAXN], flowlim;
7     void make_edge(int s, int t, ll cap, ll cost) {
8         if (!cap) return;
9         v[s].pb(edge{t, s, cap, 0LL, cost, v[t].size()});
10        v[t].pb(edge{s, t, 0LL, 0LL, -cost, v[s].size() - 1});
11    }
12    bitset<MAXN> vis;
13    void dijkstra() {
14        vis.reset();
15        __gnu_pbds::priority_queue<pair<ll, int>> q;
16        vector<decltype(q)::point_iterator> its(n);
17        q.push({0LL, s});
18        while (!q.empty()) {
19            int now = q.top().second;
20            q.pop();
21            if (vis[now]) continue;
22            vis[now] = 1;
23            ll ndis = dis[now] + pi[now];
24            for (edge &e : v[now]) {
25                if (e.flow == e.cap || vis[e.to]) continue;
26                if (dis[e.to] > ndis + e.cost - pi[e.to]) {
27                    dis[e.to] = ndis + e.cost - pi[e.to];
28                    flows[e.to] = min(flows[now], e.cap - e.flow);
29                    fromE[e.to] = &e;
30                    if (its[e.to] == q.end())
31                        its[e.to] = q.push({-dis[e.to], e.to});
32                    else q.modify(its[e.to], {-dis[e.to], e.to});
33                }
34            }
35        }
36    }
37    bool AP(ll &flow) {
38        fill_n(dis, n, INF);
39        fromE[s] = 0;
40        dis[s] = 0;
41        flows[s] = flowlim - flow;
42        dijkstra();
43        if (dis[t] == INF) return false;
44        flow += flows[t];
45        for (edge *e = fromE[t]; e; e = fromE[e->from]) {
46            e->flow += flows[t];
47            v[e->to][e->rev].flow -= flows[t];
48        }
49        for (int i = 0; i < n; i++)
50            pi[i] = min(pi[i] + dis[i], INF);
51        return true;
52    }
53    pll solve(int _s, int _t, ll _flowlim = INF) {
54        s = _s, t = _t, flowlim = _flowlim;
55        pll re;
56        while (re.F != flowlim && AP(re.F))
57            ;
58        for (int i = 0; i < n; i++)
59            for (edge &e : v[i])
60                if (e.flow != 0) re.S += e.flow * e.cost;
61        re.S /= 2;
62        return re;
63    }
64    void init(int _n) {
65        n = _n;
66        fill_n(pi, n, 0);
67        for (int i = 0; i < n; i++) v[i].clear();
68    }
69    void setpi(int s) {
70        fill_n(pi, n, INF);
71        pi[s] = 0;
72        for (ll it = 0, flag = 1, tdis; flag && it < n; it++) {
73            flag = 0;
74            for (int i = 0; i < n; i++)
75                if (pi[i] != INF)
76                    for (edge &e : v[i])
77                        if (e.cap && (tdis = pi[i] + e.cost) < pi[e.to])
78                            pi[e.to] = tdis, flag = 1;
79        }
80    }
81 };

```

### 3.2.3. Gomory-Hu Tree

Requires: Dinic's Algorithm

```

1 int e[MAXN][MAXN];
2 int p[MAXN];
3 Dinic D; // original graph

```

```

5 void gomory_hu() {
6     fill(p, p + n, 0);
7     fill(e[0], e[n], INF);
8     for (int s = 1; s < n; s++) {
9         int t = p[s];
10        Dinic F = D;
11        int tmp = F.max_flow(s, t);
12        for (int i = 1; i < s; i++)
13            e[s][i] = e[i][s] = min(tmp, e[t][i]);
14        for (int i = s + 1; i < n; i++)
15            if (p[i] == t && F.side[i]) p[i] = s;
16    }
17 }

```

### 3.2.4. Global Minimum Cut

```

1 // weights is an adjacency matrix, undirected
2 pair<int, vi> getMinCut(vector<vi> &weights) {
3     int N = sz(weights);
4     vi used(N), cut, best_cut;
5     int best_weight = -1;
6
7     for (int phase = N - 1; phase >= 0; phase--) {
8         vi w = weights[0], added = used;
9         int prev, k = 0;
10        rep(i, 0, phase) {
11            prev = k;
12            k = -1;
13            rep(j, 1, N) if (!added[j] &&
14                            (k == -1 || w[j] > w[k])) k = j;
15            if (i == phase - 1) {
16                rep(k, 0, N) weights[prev][j] += weights[k][j];
17                rep(j, 0, N) weights[j][prev] = weights[prev][j];
18                used[k] = true;
19                cut.push_back(k);
20                if (best_weight == -1 || w[k] < best_weight) {
21                    best_cut = cut;
22                    best_weight = w[k];
23                }
24            } else {
25                rep(j, 0, N) w[j] += weights[k][j];
26                added[k] = true;
27            }
28        }
29    }
30    return {best_weight, best_cut};
31 }

```

### 3.2.5. Bipartite Minimum Cover

Requires: Dinic's Algorithm

```

1 // maximum independent set = all vertices not covered
2 // x : [0, n), y : [0, m]
3 struct Bipartite_vertex_cover {
4     Dinic D;
5     int n, m, s, t, x[maxn], y[maxn];
6     void make_edge(int x, int y) { D.make_edge(x, y + n, 1); }
7     int matching() {
8         int re = D.max_flow(s, t);
9         for (int i = 0; i < n; i++)
10            for (Dinic::edge &e : D.v[i])
11                if (e.to != s && e.flow == 1) {
12                    x[i] = e.to - n, y[e.to - n] = i;
13                    break;
14                }
15        return re;
16    }
17    // init() and matching() before use
18    void solve(vector<int> &vx, vector<int> &vy) {
19        bitset<maxn * 2 + 10> vis;
20        queue<int> q;
21        for (int i = 0; i < n; i++)
22            if (x[i] == -1) q.push(i), vis[i] = 1;
23        while (!q.empty()) {
24            int now = q.front();
25            q.pop();
26            if (now < n) {
27                for (Dinic::edge &e : D.v[now])
28                    if (e.to != s && e.to - n != x[now] && !vis[e.to])
29                        vis[e.to] = 1, q.push(e.to);
30            } else {
31                if (!vis[y[now - n]])
32                    vis[y[now - n]] = 1, q.push(y[now - n]);
33            }
34        }
35        for (int i = 0; i < n; i++)
36            if (!vis[i]) vx.pb(i);
37        for (int i = 0; i < m; i++)
38            if (vis[i + n]) vy.pb(i);
39    }
40 }

```

```

39 }
40 void init(int _n, int _m) {
41     n = _n, m = _m, s = n + m, t = s + 1;
42     for (int i = 0; i < n; i++)
43         x[i] = -1, D.make_edge(s, i, 1);
44     for (int i = 0; i < m; i++)
45         y[i] = -1, D.make_edge(i + n, t, 1);
46 }
47 };

```

### 3.2.6. Edmonds' Algorithm

```

1 struct Edmonds {
2     int n, T;
3     vector<vector<int>>> g;
4     vector<int> pa, p, used, base;
5     Edmonds(int n)
6         : n(n), T(0), g(n), pa(n, -1), p(n), used(n),
7           base(n) {}
8     void add(int a, int b) {
9         g[a].push_back(b);
10        g[b].push_back(a);
11    }
12    int getBase(int i) {
13        while (i != base[i])
14            base[i] = base[base[i]], i = base[i];
15        return i;
16    }
17    vector<int> toJoin;
18    void mark_path(int v, int x, int b, vector<int> &path) {
19        for (; getBase(v) != b; v = p[x]) {
20            p[v] = x, x = pa[v];
21            toJoin.push_back(v);
22            toJoin.push_back(x);
23            if (!used[x]) used[x] = ++T, path.push_back(x);
24        }
25    }
26    bool go(int v) {
27        for (int x : g[v]) {
28            int b, bv = getBase(v), bx = getBase(x);
29            if (bv == bx) {
30                continue;
31            } else if (used[x]) {
32                vector<int> path;
33                toJoin.clear();
34                if (used[bx] < used[bv])
35                    mark_path(v, x, b = bx, path);
36                else mark_path(x, v, b = bv, path);
37                for (int z : toJoin) base[getBase(z)] = b;
38                for (int z : path)
39                    if (go(z)) return 1;
40            } else if (p[x] == -1) {
41                p[x] = v;
42                if (pa[x] == -1) {
43                    for (int y; x != -1; x = v)
44                        y = p[x], v = pa[y], pa[x] = y, pa[y] = x;
45                    return 1;
46                }
47                if (!used[pa[x]]) {
48                    used[pa[x]] = ++T;
49                    if (go(pa[x])) return 1;
50                }
51            }
52        }
53        return 0;
54    }
55    void init_dfs() {
56        for (int i = 0; i < n; i++)
57            used[i] = 0, p[i] = -1, base[i] = i;
58    }
59    bool dfs(int root) {
60        used[root] = ++T;
61        return go(root);
62    }
63    void match() {
64        int ans = 0;
65        for (int v = 0; v < n; v++)
66            for (int x : g[v])
67                if (pa[v] == -1 && pa[x] == -1) {
68                    pa[v] = x, pa[x] = v, ans++;
69                    break;
70                }
71        init_dfs();
72        for (int i = 0; i < n; i++)
73            if (pa[i] == -1 && dfs(i)) ans++, init_dfs();
74        cout << ans * 2 << "\n";
75        for (int i = 0; i < n; i++)
76            if (pa[i] > i)
77                cout << i + 1 << " " << pa[i] + 1 << "\n";
78    }
79 };

```

```

79 };

```

### 3.2.7. Minimum Weight Matching

```

1 struct Graph {
2     static const int MAXN = 105;
3     int n, e[MAXN][MAXN];
4     int match[MAXN], d[MAXN], onstk[MAXN];
5     vector<int> stk;
6     void init(int _n) {
7         n = _n;
8         for (int i = 0; i < n; i++)
9             for (int j = 0; j < n; j++)
10                // change to appropriate infinity
11                // if not complete graph
12                e[i][j] = 0;
13    }
14    void add_edge(int u, int v, int w) {
15        e[u][v] = e[v][u] = w;
16    }
17    bool SPFA(int u) {
18        if (onstk[u]) return true;
19        stk.push_back(u);
20        onstk[u] = 1;
21        for (int v = 0; v < n; v++) {
22            if (u != v && match[u] != v && !onstk[v]) {
23                int m = match[v];
24                if (d[m] > d[u] - e[v][m] + e[u][v]) {
25                    d[m] = d[u] - e[v][m] + e[u][v];
26                    onstk[v] = 1;
27                    stk.push_back(v);
28                    if (SPFA(m)) return true;
29                    stk.pop_back();
30                    onstk[v] = 0;
31                }
32            }
33        }
34        onstk[u] = 0;
35        stk.pop_back();
36        return false;
37    }
38    int solve() {
39        for (int i = 0; i < n; i += 2) {
40            match[i] = i + 1;
41            match[i + 1] = i;
42        }
43        while (true) {
44            int found = 0;
45            for (int i = 0; i < n; i++) onstk[i] = d[i] = 0;
46            for (int i = 0; i < n; i++) {
47                stk.clear();
48                if (!onstk[i] && SPFA(i)) {
49                    found = 1;
50                    while (stk.size() >= 2) {
51                        int u = stk.back();
52                        stk.pop_back();
53                        int v = stk.back();
54                        stk.pop_back();
55                        match[u] = v;
56                        match[v] = u;
57                    }
58                }
59            }
60            if (!found) break;
61        }
62        int ret = 0;
63        for (int i = 0; i < n; i++) ret += e[i][match[i]];
64        ret /= 2;
65        return ret;
66    }
67 } graph;

```

### 3.2.8. Stable Marriage

```

1 // normal stable marriage problem
2 /* input:
3 3
4 Albert Laura Nancy Marcy
5 Brad Marcy Nancy Laura
6 Chuck Laura Marcy Nancy
7 Laura Chuck Albert Brad
8 Marcy Albert Chuck Brad
9 Nancy Brad Albert Chuck
10 */
11
12 using namespace std;
13 const int MAXN = 505;
14
15 int n;

```



```

17 int favor[MAXN][MAXN]; // favor[boy_id][rank] = girl_id;
18 int order[MAXN][MAXN]; // order[girl_id][boy_id] = rank;
19 int current[MAXN]; // current[boy_id] = rank;
20 // boy_id will pursue current[boy_id] girl.
21 int girl_current[MAXN]; // girl[girl_id] = boy_id;

23 void initialize() {
24     for (int i = 0; i < n; i++) {
25         current[i] = 0;
26         girl_current[i] = n;
27         order[i][n] = n;
28     }
29 }

31 map<string, int> male, female;
32 string bname[MAXN], gname[MAXN];
33 int fit = 0;

35 void stable_marriage() {
36     queue<int> que;
37     for (int i = 0; i < n; i++) que.push(i);
38     while (!que.empty()) {
39         int boy_id = que.front();
40         que.pop();

41         int girl_id = favor[boy_id][current[boy_id]];
42         current[boy_id]++;

43         if (order[girl_id][boy_id] <
44             order[girl_id][girl_current[girl_id]]) {
45             if (girl_current[girl_id] < n)
46                 que.push(girl_current[girl_id]);
47             girl_current[girl_id] = boy_id;
48         } else {
49             que.push(boy_id);
50         }
51     }
52 }

54 int main() {
55     cin >> n;

56     for (int i = 0; i < n; i++) {
57         string p, t;
58         cin >> p;
59         male[p] = i;
60         bname[i] = p;
61         for (int j = 0; j < n; j++) {
62             cin >> t;
63             if (!female.count(t)) {
64                 gname[fit] = t;
65                 female[t] = fit++;
66             }
67             favor[i][j] = female[t];
68         }
69     }

71     for (int i = 0; i < n; i++) {
72         string p, t;
73         cin >> p;
74         for (int j = 0; j < n; j++) {
75             cin >> t;
76             order[female[p]][male[t]] = j;
77         }
78     }

79     initialize();
80     stable_marriage();

81     for (int i = 0; i < n; i++) {
82         cout << bname[i] << " "
83              << gname[favor[i][current[i] - 1]] << endl;
84     }
85 }

```

### 3.2.9. Kuhn-Munkres algorithm

```

1 // Maximum Weight Perfect Bipartite Matching
2 // Detect non-perfect-matching:
3 // 1. set all edge[i][j] as INF
4 // 2. if solve() >= INF, it is not perfect matching.
5
6 typedef long long ll;
7 struct KM {
8     static const int MAXN = 1050;
9     static const ll INF = 1LL << 60;
10    int n, match[MAXN], vx[MAXN], vy[MAXN];
11    ll edge[MAXN][MAXN], lx[MAXN], ly[MAXN], slack[MAXN];
12    void init(int _n) {
13        n = _n;

```

```

14        for (int i = 0; i < n; i++)
15            for (int j = 0; j < n; j++) edge[i][j] = 0;
16    }
17    void add_edge(int x, int y, ll w) { edge[x][y] = w; }
18    bool DFS(int x) {
19        vx[x] = 1;
20        for (int y = 0; y < n; y++) {
21            if (vy[y]) continue;
22            if (lx[x] + ly[y] > edge[x][y]) {
23                slack[y] =
24                    min(slack[y], lx[x] + ly[y] - edge[x][y]);
25            } else {
26                vy[y] = 1;
27                if (match[y] == -1 || DFS(match[y])) {
28                    match[y] = x;
29                    return true;
30                }
31            }
32        }
33        return false;
34    }
35    ll solve() {
36        fill(match, match + n, -1);
37        fill(lx, lx + n, -INF);
38        fill(ly, ly + n, 0);
39        for (int i = 0; i < n; i++)
40            for (int j = 0; j < n; j++)
41                lx[i] = max(lx[i], edge[i][j]);
42        for (int i = 0; i < n; i++) {
43            fill(slack, slack + n, INF);
44            while (true) {
45                fill(vx, vx + n, 0);
46                fill(vy, vy + n, 0);
47                if (DFS(i)) break;
48                ll d = INF;
49                for (int j = 0; j < n; j++)
50                    if (!vy[j]) d = min(d, slack[j]);
51                for (int j = 0; j < n; j++) {
52                    if (vx[j]) lx[j] -= d;
53                    if (vy[j]) ly[j] += d;
54                    else slack[j] -= d;
55                }
56            }
57        }
58        ll res = 0;
59        for (int i = 0; i < n; i++) {
60            res += edge[match[i]][i];
61        }
62        return res;
63    }
64 } graph;

```

### 3.3. Shortest Path Faster Algorithm

```

1 struct SPFA {
2     static const int maxn = 1010, INF = 1e9;
3     int dis[maxn];
4     bitset<maxn> inq, inneg;
5     queue<int> q, tq;
6     vector<pii> v[maxn];
7     void make_edge(int s, int t, int w) {
8         v[s].emplace_back(t, w);
9     }
10    void dfs(int a) {
11        inneg[a] = 1;
12        for (pii i : v[a])
13            if (!inneg[i.F]) dfs(i.F);
14    }
15    bool solve(int n, int s) { // true if have neg-cycle
16        for (int i = 0; i <= n; i++) dis[i] = INF;
17        dis[s] = 0, q.push(s);
18        for (int i = 0; i < n; i++) {
19            inq.reset();
20            int now;
21            while (!q.empty()) {
22                now = q.front(), q.pop();
23                for (pii &i : v[now]) {
24                    if (dis[i.F] > dis[now] + i.S) {
25                        dis[i.F] = dis[now] + i.S;
26                        if (!inq[i.F]) tq.push(i.F), inq[i.F] = 1;
27                    }
28                }
29            }
30            q.swap(tq);
31        }
32        bool re = !q.empty();
33        inneg.reset();
34        while (!q.empty()) {
35            if (!inneg[q.front()]) dfs(q.front());
36            q.pop();
37        }
38        return re;
39    }
40 }

```

```

37     }
38     return re;
39 }
40 void reset(int n) {
41     for (int i = 0; i <= n; i++) v[i].clear();
42 }
43 };

```

### 3.4. Strongly Connected Components

```

1 struct TarjanScc {
2     int n, step;
3     vector<int> time, low, instk, stk;
4     vector<vector<int>> e, scc;
5     TarjanScc(int n_)
6         : n(n_), step(0), time(n), low(n), instk(n), e(n) {}
7     void add_edge(int u, int v) { e[u].push_back(v); }
8     void dfs(int x) {
9         time[x] = low[x] = ++step;
10        stk.push_back(x);
11        instk[x] = 1;
12        for (int y : e[x])
13            if (!time[y]) {
14                dfs(y);
15                low[x] = min(low[x], low[y]);
16            } else if (instk[y]) {
17                low[x] = min(low[x], time[y]);
18            }
19        if (time[x] == low[x]) {
20            scc.emplace_back();
21            for (int y = -1; y != x; ) {
22                y = stk.back();
23                stk.pop_back();
24                instk[y] = 0;
25                scc.back().push_back(y);
26            }
27        }
28    }
29    void solve() {
30        for (int i = 0; i < n; i++)
31            if (!time[i]) dfs(i);
32        reverse(scc.begin(), scc.end());
33        // scc in topological order
34    }
35 };

```

#### 3.4.1. 2-Satisfiability

Requires: Strongly Connected Components

```

1 // 1 based, vertex in SCC = MAXN * 2
2 // (not i) is i + n
3 struct two_SAT {
4     int n, ans[MAXN];
5     SCC S;
6     void imply(int a, int b) { S.make_edge(a, b); }
7     bool solve(int _n) {
8         n = _n;
9         S.solve(n * 2);
10        for (int i = 1; i <= n; i++) {
11            if (S.scc[i] == S.scc[i + n]) return false;
12            ans[i] = (S.scc[i] < S.scc[i + n]);
13        }
14        return true;
15    }
16    void init(int _n) {
17        n = _n;
18        fill_n(ans, n + 1, 0);
19        S.init(n * 2);
20    }
21 } SAT;

```

### 3.5. Biconnected Components

#### 3.5.1. Articulation Points

```

1 void dfs(int x, int p) {
2     tin[x] = low[x] = ++t;
3     int ch = 0;
4     for (auto u : g[x])
5         if (u.first != p) {
6             if (!ins[u.second])
7                 st.push(u.second), ins[u.second] = true;
8             if (tin[u.first])
9                 low[x] = min(low[x], tin[u.first]);
10            continue;
11        }
12        ++ch;
13        dfs(u.first, x);
14        low[x] = min(low[x], low[u.first]);

```

```

15     if (low[u.first] >= tin[x]) {
16         cut[x] = true;
17         ++sz;
18         while (true) {
19             int e = st.top();
20             st.pop();
21             bcc[e] = sz;
22             if (e == u.second) break;
23         }
24     }
25     if (ch == 1 && p == -1) cut[x] = false;
26 }

```

#### 3.5.2. Bridges

```

1 // if there are multi-edges, then they are not bridges
2 void dfs(int x, int p) {
3     tin[x] = low[x] = ++t;
4     st.push(x);
5     for (auto u : g[x])
6         if (u.first != p) {
7             if (tin[u.first])
8                 low[x] = min(low[x], tin[u.first]);
9             continue;
10        }
11        dfs(u.first, x);
12        low[x] = min(low[x], low[u.first]);
13        if (low[u.first] == tin[u.first]) br[u.second] = true;
14    }
15    if (tin[x] == low[x]) {
16        ++sz;
17        while (st.size()) {
18            int u = st.top();
19            st.pop();
20            bcc[u] = sz;
21            if (u == x) break;
22        }
23    }
24 }

```

### 3.6. Triconnected Components

```

1 // requires a union-find data structure
2 struct ThreeEdgeCC {
3     int V, ind;
4     vector<int> id, pre, post, low, deg, path;
5     vector<vector<int>> components;
6     UnionFind uf;
7     template <class Graph>
8     void dfs(const Graph &G, int v, int prev) {
9         pre[v] = ++ind;
10        for (int w : G[v])
11            if (w != v) {
12                if (w == prev) {
13                    prev = -1;
14                    continue;
15                }
16                if (pre[w] != -1) {
17                    if (pre[w] < pre[v]) {
18                        deg[v]++;
19                        low[v] = min(low[v], pre[w]);
20                    } else {
21                        deg[v]--;
22                        int &u = path[v];
23                        for (; u != -1 && pre[u] <= pre[w] &&
24                             pre[w] <= post[u];) {
25                            uf.join(v, u);
26                            deg[v] += deg[u];
27                            u = path[u];
28                        }
29                    }
30                    continue;
31                }
32                dfs(G, w, v);
33                if (path[w] == -1 && deg[w] <= 1) {
34                    deg[v] += deg[w];
35                    low[v] = min(low[v], low[w]);
36                    continue;
37                }
38                if (deg[w] == 0) w = path[w];
39                if (low[v] > low[w]) {
40                    low[v] = min(low[v], low[w]);
41                    swap(w, path[v]);
42                }
43                for (; w != -1; w = path[w]) {
44                    uf.join(v, w);
45                    deg[v] += deg[w];
46                }
47            }

```

```

    post[v] = ind;
}
template <class Graph>
ThreeEdgeCC(const Graph &G)
: V(G.size()), ind(-1), id(V, -1), pre(V, -1),
  post(V), low(V, INT_MAX), deg(V, 0), path(V, -1),
  uf(V) {
    for (int v = 0; v < V; v++)
        if (pre[v] == -1) dfs(G, v, -1);
    components.reserve(uf.cnt);
    for (int v = 0; v < V; v++)
        if (uf.find(v) == v) {
            id[v] = components.size();
            components.emplace_back(1, v);
            components.back().reserve(uf.getSize(v));
        }
    for (int v = 0; v < V; v++)
        if (id[v] == -1)
            components[id[v] = id[uf.find(v)]] .push_back(v);
}
};

```

### 3.7. Centroid Decomposition

```

1 void get_center(int now) {
    v[now] = true;
    vtx.push_back(now);
    sz[now] = 1;
    mx[now] = 0;
    for (int u : G[now])
        if (!v[u]) {
            get_center(u);
            mx[now] = max(mx[now], sz[u]);
            sz[now] += sz[u];
        }
}
13 void get_dis(int now, int d, int len) {
    dis[d][now] = cnt;
    v[now] = true;
    for (auto u : G[now])
        if (!v[u.first]) { get_dis(u, d, len + u.second); }
}
19 void dfs(int now, int fa, int d) {
    get_center(now);
    int c = -1;
    for (int i : vtx) {
        if (max(mx[i], (int)vtx.size() - sz[i]) <=
            (int)vtx.size() / 2)
            c = i;
        v[i] = false;
    }
    get_dis(c, d, 0);
    for (int i : vtx) v[i] = false;
    v[c] = true;
    vtx.clear();
    dep[c] = d;
    p[c] = fa;
    for (auto u : G[c])
        if (u.first != fa && !v[u.first]) {
            dfs(u.first, c, d + 1);
        }
}

```

### 3.8. Minimum Mean Cycle

```

1 // d[i][j] == 0 if {i,j} !in E
long long d[1003][1003], dp[1003][1003];
3 pair<long long, long long> MMWC() {
    memset(dp, 0x3f, sizeof(dp));
    for (int i = 1; i <= n; ++i) dp[0][i] = 0;
    for (int i = 1; i <= n; ++i) {
        for (int j = 1; j <= n; ++j) {
            for (int k = 1; k <= n; ++k) {
                dp[i][k] = min(dp[i - 1][j] + d[j][k], dp[i][k]);
            }
        }
    }
    long long au = 1ll << 31, ad = 1;
    for (int i = 1; i <= n; ++i) {
        if (dp[n][i] == 0x3f3f3f3f3f3f3f3f) continue;
        long long u = 0, d = 1;
        for (int j = n - 1; j >= 0; --j) {
            if ((dp[n][i] - dp[j][i]) * d > u * (n - j)) {
                u = dp[n][i] - dp[j][i];
                d = n - j;
            }
        }
        if (u * ad < au * d) au = u, ad = d;
    }
}

```

```

long long g = __gcd(au, ad);
return make_pair(au / g, ad / g);
}

```

### 3.9. Directed MST

```

1 template <typename T> struct DMST {
    T g[maxn][maxn], fw[maxn];
    int n, fr[maxn];
    bool vis[maxn], inc[maxn];
    void clear() {
        for (int i = 0; i < maxn; ++i) {
            for (int j = 0; j < maxn; ++j) g[i][j] = inf;
            vis[i] = inc[i] = false;
        }
    }
    void addedge(int u, int v, T w) {
        g[u][v] = min(g[u][v], w);
    }
    T operator()(int root, int _n) {
        n = _n;
        if (dfs(root) != n) return -1;
        T ans = 0;
        while (true) {
            for (int i = 1; i <= n; ++i) fw[i] = inf, fr[i] = i;
            for (int i = 1; i <= n; ++i)
                if (!inc[i]) {
                    for (int j = 1; j <= n; ++j) {
                        if (!inc[j] && i != j && g[j][i] < fw[i]) {
                            fw[i] = g[j][i];
                            fr[i] = j;
                        }
                    }
                }
            int x = -1;
            for (int i = 1; i <= n; ++i)
                if (i != root && !inc[i]) {
                    int j = i, c = 0;
                    while (j != root && fr[j] != i && c <= n)
                        ++c, j = fr[j];
                    if (j == root || c > n) continue;
                    else {
                        x = i;
                        break;
                    }
                }
            if (!x) {
                for (int i = 1; i <= n; ++i)
                    if (i != root && !inc[i]) ans += fw[i];
                return ans;
            }
            int y = x;
            for (int i = 1; i <= n; ++i) vis[i] = false;
            do {
                ans += fw[y];
                y = fr[y];
                vis[y] = inc[y] = true;
            } while (y != x);
            inc[x] = false;
            for (int k = 1; k <= n; ++k)
                if (vis[k]) {
                    for (int j = 1; j <= n; ++j)
                        if (!vis[j]) {
                            if (g[x][j] > g[k][j]) g[x][j] = g[k][j];
                            if (g[j][k] < inf &&
                                g[j][k] - fw[k] < g[j][x])
                                g[j][x] = g[j][k] - fw[k];
                        }
                }
        }
    }
    int dfs(int now) {
        int r = 1;
        vis[now] = true;
        for (int i = 1; i <= n; ++i)
            if (g[now][i] < inf && !vis[i]) r += dfs(i);
        return r;
    }
};

```

### 3.10. Maximum Clique

```

1 // source: KACTL
3 typedef vector<bitset<200>> vb;
struct Maxclique {
    double limit = 0.025, pk = 0;
    struct Vertex {
        int i, d = 0;
    };
};

```

```

};
9 typedef vector<Vertex> vv;
  vb e;
11 vv V;
  vector<vi> C;
  vi qmax, q, S, old;
13 void init(vv &r) {
15     for (auto &v : r) v.d = 0;
     for (auto &v : r)
17         for (auto j : r) v.d += e[v.i][j.i];
     sort(all(r), [](auto a, auto b) { return a.d > b.d; });
19     int mxD = r[0].d;
     rep(i, 0, sz(r)) r[i].d = min(i, mxD) + 1;
21 }
  void expand(vv &R, int lev = 1) {
23     S[lev] += S[lev - 1] - old[lev];
     old[lev] = S[lev - 1];
25     while (sz(R)) {
         if (sz(q) + R.back().d <= sz(qmax)) return;
27         q.push_back(R.back().i);
         vv T;
29         for (auto v : R)
             if (e[R.back().i][v.i] T.push_back({v.i});
31         if (sz(T)) {
             if (S[lev]++ / ++pk < limit) init(T);
33             int j = 0, mxk = 1,
                 mnk = max(sz(qmax) - sz(q) + 1, 1);
             C[1].clear(), C[2].clear();
35             for (auto v : T) {
                 int k = 1;
37                 auto f = [&](int i) { return e[v.i][i]; };
                 while (any_of(all(C[k]), f)) k++;
39                 if (k > mxk) mxk = k, C[mxk + 1].clear();
                 if (k < mnk) T[j++].i = v.i;
                 C[k].push_back(v.i);
41             }
             if (j > 0) T[j - 1].d = 0;
43             rep(k, mnk, mxk + 1) for (int i : C[k]) T[j].i = i,
                                     T[j++].d = k;
45             expand(T, lev + 1);
47         } else if (sz(q) > sz(qmax)) qmax = q;
         q.pop_back(), R.pop_back();
51     }
  }
  vi maxClique() {
53     init(V), expand(V);
55     return qmax;
  }
  Maxclique(vb conn)
57     : e(conn), C(sz(e) + 1), S(sz(C)), old(S) {
59     rep(i, 0, sz(e)) V.push_back({i});
61 };

```

### 3.11. Dominator Tree

```

1 // idom[n] is the unique node that strictly dominates n but
  // does not strictly dominate any other node that strictly
  // dominates n. idom[n] = 0 if n is entry or the entry
  // cannot reach n.
5 struct DominatorTree {
7     static const int MAXN = 200010;
     int n, s;
     vector<int> g[MAXN], pred[MAXN];
     vector<int> cov[MAXN];
     int dfn[MAXN], nfd[MAXN], ts;
     int par[MAXN];
     int sdom[MAXN], idom[MAXN];
     int mom[MAXN], mn[MAXN];
13
15     inline bool cmp(int u, int v) { return dfn[u] < dfn[v]; }

17     int eval(int u) {
         if (mom[u] == u) return u;
19         int res = eval(mom[u]);
         if (cmp(sdom[mn[mom[u]]], sdom[mn[u]]))
21             mn[u] = mn[mom[u]];
         return mom[u] = res;
23     }

25     void init(int _n, int _s) {
         n = _n;
27         s = _s;
         REP1(i, 1, n) {
29             g[i].clear();
             pred[i].clear();
             idom[i] = 0;
31         }
33     }

```

```

void add_edge(int u, int v) {
35     g[u].push_back(v);
     pred[v].push_back(u);
37 }
  void DFS(int u) {
39     ts++;
     dfn[u] = ts;
41     nfd[ts] = u;
     for (int v : g[u])
43         if (dfn[v] == 0) {
             par[v] = u;
             DFS(v);
45         }
47 }
  void build() {
49     ts = 0;
     REP1(i, 1, n) {
51         dfn[i] = nfd[i] = 0;
         cov[i].clear();
53         mom[i] = mn[i] = sdom[i] = i;
     }
     DFS(s);
55     for (int i = ts; i >= 2; i--) {
         int u = nfd[i];
57         if (u == 0) continue;
         for (int v : pred[u])
59             if (dfn[v] != 0) {
                 eval(v);
                 if (cmp(sdom[mn[v]], sdom[u]))
61                     sdom[u] = sdom[mn[v]];
             }
63         cov[sdom[u]].push_back(u);
         mom[u] = par[u];
65         for (int w : cov[par[u]]) {
             eval(w);
67             if (cmp(sdom[mn[w]], par[u])) idom[w] = mn[w];
             else idom[w] = par[u];
69         }
         cov[par[u]].clear();
71     }
     REP1(i, 2, ts) {
73         int u = nfd[i];
         if (u == 0) continue;
75         if (idom[u] != sdom[u]) idom[u] = idom[idom[u]];
77     }
79 } dom;

```

### 3.12. Manhattan Distance MST

```

1 // returns [(dist, from, to), ...]
  // then do normal mst afterwards
3 typedef Point<int> P;
  vector<array<int, 3>> manhattanMST(vector<P> ps) {
5     vi id(sz(ps));
     iota(all(id), 0);
7     vector<array<int, 3>> edges;
     rep(k, 0, 4) {
9         sort(all(id), [&](int i, int j) {
             return (ps[i] - ps[j]).x < (ps[j] - ps[i]).y;
11         });
         map<int, int> sweep;
13         for (int i : id) {
             for (auto it = sweep.lower_bound(-ps[i].y);
15                 it != sweep.end(); sweep.erase(it++)) {
                 int j = it->second;
17                 P d = ps[i] - ps[j];
                 if (d.y > d.x) break;
                 edges.push_back({d.y + d.x, i, j});
19             }
             sweep[-ps[i].y] = i;
21         }
         for (P &p : ps)
23             if (k & 1) p.x = -p.x;
             else swap(p.x, p.y);
25     }
27     return edges;
  }

```

## 4. Math

### 4.1. Number Theory

#### 4.1.1. Mod Struct

A list of safe primes: 26003, 27767, 28319, 28979, 29243, 29759, 30467, 910927547, 919012223, 947326223, 990669467, 1007939579, 1019126699

929760389146037459, 975500632317046523, 989312547895528379

NTT prime $p$	$p - 1$	primitive root
65537	$1 \ll 16$	3
998244353	$119 \ll 23$	3
2748779069441	$5 \ll 39$	3
1945555039024054273	$27 \ll 56$	5

Requires: Extended GCD

```

1 template <typename T> struct M {
2     static T MOD; // change to constexpr if already known
3     T v;
4     M(T x = 0) {
5         v = (-MOD <= x && x < MOD) ? x : x % MOD;
6         if (v < 0) v += MOD;
7     }
8     explicit operator T() const { return v; }
9     bool operator==(const M &b) const { return v == b.v; }
10    bool operator!=(const M &b) const { return v != b.v; }
11    M operator-() { return M(-v); }
12    M operator+(M b) { return M(v + b.v); }
13    M operator-(M b) { return M(v - b.v); }
14    M operator*(M b) { return M((__int128)v * b.v % MOD); }
15    M operator/(M b) { return *this * (b ^ (MOD - 2)); }
16    // change above implementation to this if MOD is not prime
17    M inv() {
18        auto [p, _, g] = extgcd(v, MOD);
19        return assert(g == 1), p;
20    }
21    friend M operator^(M a, ll b) {
22        M ans(1);
23        for (; b >= 1, a *= a)
24            if (b & 1) ans *= a;
25        return ans;
26    }
27    friend M &operator+=(M &a, M b) { return a = a + b; }
28    friend M &operator-=(M &a, M b) { return a = a - b; }
29    friend M &operator*=(M &a, M b) { return a = a * b; }
30    friend M &operator/=(M &a, M b) { return a = a / b; }
31 };
32 using Mod = M<int>;
33 template <> int Mod::MOD = 1'000'000'007;
34 int &MOD = Mod::MOD;

```

## 4.1.2. Miller-Rabin

Requires: Mod Struct

```

1 // checks if Mod::MOD is prime
2 bool is_prime() {
3     if (MOD < 2 || MOD % 2 == 0) return MOD == 2;
4     Mod A[] = {2, 7, 61}; // for int values (< 2^31)
5     // ll: 2, 325, 9375, 28178, 450775, 9780504, 1795265022
6     int s = __builtin_ctzll(MOD - 1), i;
7     for (Mod a : A) {
8         Mod x = a ^ (MOD >> s);
9         for (i = 0; i < s && (x + 1).v > 2; i++) x *= x;
10        if (i && x != -1) return 0;
11    }
12    return 1;
13 }

```

## 4.1.3. Linear Sieve

```

1 constexpr ll MAXN = 1000000;
2 bitset<MAXN> is_prime;
3 vector<ll> primes;
4 ll mpf[MAXN], phi[MAXN], mu[MAXN];
5
6 void sieve() {
7     is_prime.set();
8     is_prime[1] = 0;
9     mu[1] = phi[1] = 1;
10    for (ll i = 2; i < MAXN; i++) {
11        if (is_prime[i]) {
12            mpf[i] = i;
13            primes.push_back(i);
14            phi[i] = i - 1;
15            mu[i] = -1;
16        }
17        for (ll p : primes) {
18            if (p > mpf[i] || i * p >= MAXN) break;
19            is_prime[i * p] = 0;
20            mpf[i * p] = p;
21            mu[i * p] = -mu[i];
22            if (i % p == 0)
23                phi[i * p] = phi[i] * p, mu[i * p] = 0;
24            else phi[i * p] = phi[i] * (p - 1);
25        }
26    }
27 }

```

## 4.1.4. Get Factors

Requires: Linear Sieve

```

1 vector<ll> all_factors(ll n) {
2     vector<ll> fac = {1};
3     while (n > 1) {
4         const ll p = mpf[n];
5         vector<ll> cur = {1};
6         while (n % p == 0) {
7             n /= p;
8             cur.push_back(cur.back() * p);
9         }
10        vector<ll> tmp;
11        for (auto x : fac)
12            for (auto y : cur) tmp.push_back(x * y);
13        tmp.swap(fac);
14    }
15    return fac;
16 }

```

## 4.1.5. Binary GCD

```

1 // returns the gcd of non-negative a, b
2 ull bin_gcd(ull a, ull b) {
3     if (!a || !b) return a + b;
4     int s = __builtin_ctzll(a | b);
5     a >>= __builtin_ctzll(a);
6     while (b) {
7         if ((b >>= __builtin_ctzll(b)) < a) swap(a, b);
8         b >>= 1;
9     }
10    return a << s;
11 }

```

## 4.1.6. Extended GCD

```

1 // returns (p, q, g): p * a + q * b == g == gcd(a, b)
2 // g is not guaranteed to be positive when a < 0 or b < 0
3 tuple<ll, ll, ll> extgcd(ll a, ll b) {
4     ll s = 1, t = 0, u = 0, v = 1;
5     while (b) {
6         ll q = a / b;
7         swap(a -= q * b, b);
8         swap(s -= q * t, t);
9         swap(u -= q * v, v);
10    }
11    return {s, u, a};
12 }

```

## 4.1.7. Chinese Remainder Theorem

Requires: Extended GCD

```

1 // for 0 <= a < m, 0 <= b < n, returns the smallest x >= 0
2 // such that x % m == a and x % n == b
3 ll crt(ll a, ll m, ll b, ll n) {
4     if (n > m) swap(a, b), swap(m, n);
5     auto [x, y, g] = extgcd(m, n);
6     assert((a - b) % g == 0); // no solution
7     x = ((b - a) / g * x) % (n / g) * m + a;
8     return x < 0 ? x + m / g * n : x;
9 }

```

## 4.1.8. Baby-Step Giant-Step

Requires: Mod Struct

```

1 // returns x such that a ^ x = b where x \in [l, r)
2 ll bsgs(Mod a, Mod b, ll l = 0, ll r = MOD - 1) {
3     int m = sqrt(r - l) + 1, i;
4     unordered_map<ll, ll> tb;
5     Mod d = (a ^ l) / b;
6     for (i = 0, d = (a ^ l) / b; i < m; i++, d *= a)
7         if (d == 1) return l + i;
8     else tb[(ll)d] = l + i;
9     Mod c = Mod(1) / (a ^ m);
10    for (i = 0, d = 1; i < m; i++, d *= c)
11        if (auto j = tb.find((ll)d); j != tb.end())
12            return j->second + i * m;
13    return assert(0), -1; // no solution
14 }

```



#### 4.1.9. Pollard's Rho

```

1 ll f(ll x, ll mod) { return (x * x + 1) % mod; }
  // n should be composite
3 ll pollard_rho(ll n) {
  if (!(n & 1)) return 2;
  while (1) {
    ll y = 2, x = RNG() % (n - 1) + 1, res = 1;
    for (int sz = 2; res == 1; sz *= 2) {
      for (int i = 0; i < sz && res <= 1; i++) {
        x = f(x, n);
        res = __gcd(abs(x - y), n);
      }
      y = x;
    }
    if (res != 0 && res != n) return res;
  }
}

```

#### 4.1.10. Tonelli-Shanks Algorithm

Requires: Mod Struct

```

1 int legendre(Mod a) {
  if (a == 0) return 0;
  return (a ^ ((MOD - 1) / 2)) == 1 ? 1 : -1;
}
5 Mod sqrt(Mod a) {
  assert(legendre(a) != -1); // no solution
  ll p = MOD, s = p - 1;
  if (a == 0) return 0;
  if (p == 2) return 1;
  if (p % 4 == 3) return a ^ ((p + 1) / 4);
  int r, m;
  for (r = 0; !(s & 1); r++) s >>= 1;
  Mod n = 2;
  while (legendre(n) != -1) n += 1;
  Mod x = a ^ ((s + 1) / 2), b = a ^ s, g = n ^ s;
  while (b != 1) {
    Mod t = b;
    for (m = 0; t != 1; m++) t *= t;
    Mod gs = g ^ (1LL << (r - m - 1));
    g = gs * gs, x *= gs, b *= g, r = m;
  }
  return x;
}
  // to get sqrt(X) modulo p^k, where p is an odd prime:
25 // c = x^2 (mod p), c = X^2 (mod p^k), q = p^(k-1)
  // X = x^q * c^((p^k-2q+1)/2) (mod p^k)

```

#### 4.1.11. Chinese Sieve

```

1 const ll N = 1000000;
  // f, g, h multiplicative, h = f (dirichlet convolution) g
3 ll pre_g(ll n);
  ll pre_h(ll n);
  // preprocessed prefix sum of f
  ll pre_f[N];
  // prefix sum of multiplicative function f
  ll solve_f(ll n) {
    static unordered_map<ll, ll> m;
    if (n < N) return pre_f[n];
    if (m.count(n)) return m[n];
    ll ans = pre_h(n);
    for (ll l = 2, r; l <= n; l = r + 1) {
      r = n / (n / l);
      ans -= (pre_g(r) - pre_g(l - 1)) * djs_f(n / l);
    }
    return m[n] = ans;
  }

```

#### 4.1.12. Rational Number Binary Search

```

1 struct QQ {
  ll p, q;
3 QQ go(QQ b, ll d) { return {p + b.p * d, q + b.q * d}; }
};
  bool pred(QQ);
  // returns smallest p/q in [lo, hi] such that
  // pred(p/q) is true, and 0 <= p,q <= N
  QQ frac_bs(ll N) {
    QQ lo{0, 1}, hi{1, 0};
    if (pred(lo)) return lo;
    assert(pred(hi));
    bool dir = 1, L = 1, H = 1;
    for (; L || H; dir = !dir) {
      ll len = 0, step = 1;
      for (int t = 0; t < 2 && (t ? step /= 2 : step *= 2);)
        if (QQ mid = hi.go(lo, len + step);
            mid.p > N || mid.q > N || dir ^ pred(mid))
          t++;
    }
  }

```

```

19 else len += step;
    swap(lo, hi = hi.go(lo, len));
    (dir ? L : H) = !len;
21 }
23 return dir ? hi : lo;
  }

```

#### 4.1.13. Farey Sequence

```

1 // returns (e/f), where (a/b, c/d, e/f) are
  // three consecutive terms in the order n farey sequence
3 // to start, call next_farey(n, 0, 1, 1, n)
  pll next_farey(ll n, ll a, ll b, ll c, ll d) {
5   ll p = (n + b) / d;
  return pll(p * c - a, p * d - b);
7 }

```

### 4.2. Combinatorics

#### 4.2.1. Matroid Intersection

This template assumes 2 weighted matroids of the same type, and that removing an element is much more expensive than checking if one can be added. **Remember to change the implementation details.**

The ground set is  $0, 1, \dots, n-1$ , where element  $i$  has weight  $w[i]$ . For the unweighted version, remove weights and change BF/SPFA to BFS.

```

1 constexpr int N = 100;
  constexpr int INF = 1e9;
3 struct Matroid {
  // represents an independent set
5   Matroid(bitset<N>); // initialize from an independent set
  bool can_add(int); // if adding will break independence
7   Matroid remove(int); // removing from the set
};
9 auto matroid_intersection(int n, const vector<int> &w) {
  bitset<N> S;
  for (int sz = 1; sz <= n; sz++) {
    Matroid M1(S), M2(S);
13
    vector<vector<pii>> e(n + 2);
    for (int j = 0; j < n; j++)
      if (!S[j]) {
17         if (M1.can_add(j)) e[n].emplace_back(j, -w[j]);
        if (M2.can_add(j)) e[j].emplace_back(n + 1, 0);
      }
    for (int i = 0; i < n; i++)
      if (S[i]) {
21         Matroid T1 = M1.remove(i), T2 = M2.remove(i);
        for (int j = 0; j < n; j++)
          if (!S[j]) {
25             if (T1.can_add(j)) e[i].emplace_back(j, -w[j]);
            if (T2.can_add(j)) e[j].emplace_back(i, w[i]);
          }
27
29   }
31   vector<pii> dis(n + 2, {INF, 0});
  vector<int> prev(n + 2, -1);
  dis[n] = {0, 0};
  // change to SPFA for more speed, if necessary
35   bool upd = 1;
  while (upd) {
    upd = 0;
    for (int u = 0; u < n + 2; u++)
      for (auto [v, c] : e[u]) {
39         pii x(dis[u].first + c, dis[u].second + 1);
        if (x < dis[v]) dis[v] = x, prev[v] = u, upd = 1;
      }
41
43   }
45   if (dis[n + 1].first < INF)
    for (int x = prev[n + 1]; x != n; x = prev[x])
      S.flip(x);
47   else break;
49
  // S is the max-weighted independent set with size sz
51 }
  return S;
53 }

```

#### 4.2.2. De Bruijn Sequence

```

1 int res[kN], aux[kN], a[kN], sz;
  void Rec(int t, int p, int n, int k) {
3   if (t > n) {
    if (n % p == 0)
      for (int i = 1; i <= p; ++i) res[sz++] = aux[i];
    else {
5

```

```

7   aux[t] = aux[t - p];
   Rec(t + 1, p, n, k);
9   for (aux[t] = aux[t - p] + 1; aux[t] < k; ++aux[t])
       Rec(t + 1, t, n, k);
11  }
13  int DeBruijn(int k, int n) {
   // return cyclic string of length k^n such that every
   // string of length n using k character appears as a
   // substring.
15  if (k == 1) return res[0] = 0, 1;
   fill(aux, aux + k * n, 0);
17  return sz = 0, Rec(1, 1, n, k), sz;
19  }

```

### 4.2.3. Multinomial

```

1  // ways to permute v[i]
ll multinomial(vi &v) {
3  ll c = 1, m = v.empty() ? 1 : v[0];
   for (int i = 1; i < v.size(); i++)
5     for (int j = 0; i < v[i]; j++) c = c * ++m / (j + 1);
   return c;
7  }

```

## 4.3. Algebra

### 4.3.1. Formal Power Series

```

1  template <typename mint>
   struct FormalPowerSeries : vector<mint> {
3     using vector<mint>::vector;
       using FPS = FormalPowerSeries;

5     FPS &operator+=(const FPS &r) {
7         if (r.size() > this->size()) this->resize(r.size());
         for (int i = 0; i < (int)r.size(); i++)
9             (*this)[i] += r[i];
         return *this;
11    }

13    FPS &operator+=(const mint &r) {
         if (this->empty()) this->resize(1);
         (*this)[0] += r;
         return *this;
15    }

17    FPS &operator-=(const FPS &r) {
         if (r.size() > this->size()) this->resize(r.size());
         for (int i = 0; i < (int)r.size(); i++)
21             (*this)[i] -= r[i];
         return *this;
23    }

25    FPS &operator-=(const mint &r) {
         if (this->empty()) this->resize(1);
         (*this)[0] -= r;
         return *this;
27    }

29    FPS &operator*=(const mint &v) {
         for (int k = 0; k < (int)this->size(); k++)
33             (*this)[k] *= v;
         return *this;
35    }

37    FPS &operator/=(const FPS &r) {
         if (this->size() < r.size()) {
39             this->clear();
             return *this;
41         }
         int n = this->size() - r.size() + 1;
         if ((int)r.size() <= 64) {
43             FPS f(*this), g(r);
             g.shrink();
             mint coeff = g.back().inverse();
             for (auto &x : g) x *= coeff;
45             int deg = (int)f.size() - (int)g.size() + 1;
             int gs = g.size();
             FPS quo(deg);
             for (int i = deg - 1; i >= 0; i--) {
51                 quo[i] = f[i + gs - 1];
                 for (int j = 0; j < gs; j++)
53                     f[i + j] -= quo[i] * g[j];
             }
             *this = quo * coeff;
             this->resize(n, mint(0));
             return *this;
57         }
59         return *this = ((*this).rev().pre(n) * r.rev().inv(n))
61     };

```

```

63         .pre(n)
        .rev();
65     }

67     FPS &operator%=(const FPS &r) {
         *this -= *this / r * r;
         shrink();
69         return *this;
71     }

73     FPS operator+(const FPS &r) const {
         return FPS(*this) += r;
75     }

77     FPS operator+(const mint &v) const {
         return FPS(*this) += v;
79     }

81     FPS operator-(const FPS &r) const {
         return FPS(*this) -= r;
83     }

85     FPS operator-(const mint &v) const {
         return FPS(*this) -= v;
87     }

89     FPS operator*(const FPS &r) const {
         return FPS(*this) *= r;
91     }

93     FPS operator*(const mint &v) const {
         return FPS(*this) *= v;
95     }

97     FPS operator/(const FPS &r) const {
         return FPS(*this) /= r;
99     }

101    FPS operator/(const FPS &r) const {
         return FPS(*this) /= r;
103    }

105    FPS operator%(const FPS &r) const {
         return FPS(*this) %= r;
107    }

109    FPS operator-() const {
         FPS ret(this->size());
         for (int i = 0; i < (int)this->size(); i++)
111             ret[i] = -(*this)[i];
         return ret;
113    }

115    void shrink() {
         while (this->size() && this->back() == mint(0))
117             this->pop_back();
119    }

121    FPS rev() const {
         FPS ret(*this);
         reverse(begin(ret), end(ret));
         return ret;
123    }

125    FPS dot(FPS r) const {
         FPS ret(min(this->size(), r.size()));
         for (int i = 0; i < (int)ret.size(); i++)
127             ret[i] = (*this)[i] * r[i];
         return ret;
129    }

131    FPS pre(int sz) const {
         return FPS(begin(*this),
133             begin(*this) + min((int)this->size(), sz));
135    }

137    FPS operator>>(int sz) const {
         if ((int)this->size() <= sz) return {};
         FPS ret(*this);
         ret.erase(ret.begin(), ret.begin() + sz);
         return ret;
139    }

141    FPS operator<<(int sz) const {
         FPS ret(*this);
         ret.insert(ret.begin(), sz, mint(0));
         return ret;
143    }

145    FPS diff() const {
         const int n = (int)this->size();
         FPS ret(max(0, n - 1));
         mint one(1), coeff(1);
         for (int i = 1; i < n; i++) {
147             ret[i - 1] = (*this)[i] * coeff;
             coeff += one;
149         }
         return ret;
151    }

    FPS integral() const {
         const int n = (int)this->size();

```

```

153 FPS ret(n + 1);
    ret[0] = mint(0);
    if (n > 0) ret[1] = mint(1);
155 auto mod = mint::get_mod();
    for (int i = 2; i <= n; i++)
157     ret[i] = (-ret[mod % i]) * (mod / i);
    for (int i = 0; i < n; i++) ret[i + 1] *= (*this)[i];
159 return ret;
}

161 mint eval(mint x) const {
    mint r = 0, w = 1;
    for (auto &v : *this) r += w * v, w *= x;
165 return r;
}

167 FPS log(int deg = -1) const {
    assert((*this)[0] == mint(1));
    if (deg == -1) deg = (int)this->size();
169 return (this->diff() * this->inv(deg))
    .pre(deg - 1)
    .integral();
173 }

175 FPS pow(int64_t k, int deg = -1) const {
    const int n = (int)this->size();
    if (deg == -1) deg = n;
177 for (int i = 0; i < n; i++) {
        if ((*this)[i] != mint(0)) {
181             if (i * k > deg) return FPS(deg, mint(0));
            mint rev = mint(1) / (*this)[i];
            FPS ret =
183             (((*this * rev) >> i).log(deg) * k).exp(deg) *
            ((*this)[i].pow(k));
            ret = (ret << (i * k)).pre(deg);
185             if ((int)ret.size() < deg) ret.resize(deg, mint(0));
            return ret;
187         }
    }
189 return FPS(deg, mint(0));
}

191 static void ntt_ptr;
193 static void set_fft();
195 FPS &operator*=(const FPS &r);
197 void ntt();
199 void intt();
201 FPS inv(int deg = -1) const;
203 FPS exp(int deg = -1) const;
205 template <typename mint>
void *FormalPowerSeries<mint>::ntt_ptr = nullptr;

```

## 4.4. Theorems

### 4.4.1. Kirchhoff's Theorem

Denote  $L$  be a  $n \times n$  matrix as the Laplacian matrix of graph  $G$ , where  $L_{ii} = d(i)$ ,  $L_{ij} = -c$  where  $c$  is the number of edge  $(i, j)$  in  $G$ .

- The number of undirected spanning in  $G$  is  $|\det(\tilde{L}_{11})|$ .
- The number of directed spanning tree rooted at  $r$  in  $G$  is  $|\det(\tilde{L}_{rr})|$ .

### 4.4.2. Tutte's Matrix

Let  $D$  be a  $n \times n$  matrix, where  $d_{ij} = x_{ij}$  ( $x_{ij}$  is chosen uniformly at random) if  $i < j$  and  $(i, j) \in E$ , otherwise  $d_{ij} = -d_{ji}$ .  $\frac{\text{rank}(D)}{2}$  is the maximum matching on  $G$ .

### 4.4.3. Cayley's Formula

- Given a degree sequence  $d_1, d_2, \dots, d_n$  for each *labeled* vertices, there are

$$\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$$

spanning trees.

- Let  $T_{n,k}$  be the number of *labeled* forests on  $n$  vertices with  $k$  components, such that vertex  $1, 2, \dots, k$  belong to different components. Then  $T_{n,k} = kn^{n-k-1}$ .

### 4.4.4. Erdős–Gallai Theorem

A sequence of non-negative integers  $d_1 \geq d_2 \geq \dots \geq d_n$  can be represented as the degree sequence of a finite simple graph on  $n$  vertices if and only if  $d_1 + d_2 + \dots + d_n$  is even and

$$\sum_{i=1}^k d_i \leq k(k-1) + \sum_{i=k+1}^n \min(d_i, k)$$

holds for all  $1 \leq k \leq n$ .

### 4.4.5. Burnside's Lemma

Let  $X$  be a set and  $G$  be a group that acts on  $X$ . For  $g \in G$ , denote by  $X^g$  the elements fixed by  $g$ :

$$X^g = \{x \in X \mid gx \in X\}$$

Then

$$|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|.$$

## 5. Numeric

### 5.1. Barrett Reduction

```

1 using ull = unsigned long long;
2 using ul = __uint128_t;
3 // very fast calculation of a % m
4 struct reduction {
5     const ull m, d;
6     explicit reduction(ull m) : m(m), d(((ul)1 << 64) / m) {}
7     inline ull operator()(ull a) const {
8         ull q = (ull)((ul)d * a) >> 64;
9         return (a - q * m) >= m ? a - m : a;
10    }
11 };

```

### 5.2. Long Long Multiplication

```

1 using ull = unsigned long long;
2 using ll = long long;
3 using ld = long double;
4 // returns a * b % M where a, b < M < 2**63
5 ull mult(ull a, ull b, ull M) {
6     ll ret = a * b - M * ull(ld(a) * ld(b) / ld(M));
7     return ret + M * (ret < 0) - M * (ret >= (ll)M);
8 }

```

### 5.3. Fast Fourier Transform

```

1 template <typename T>
2 void fft(int n, vector<T> &a, vector<T> &rt, bool inv) {
3     vector<int> br(n);
4     for (int i = 1; i < n; i++) {
5         br[i] = (i & 1) ? br[i - 1] + n / 2 : br[i / 2] / 2;
6         if (br[i] > i) swap(a[i], a[br[i]]);
7     }
8     for (int len = 2; len <= n; len *= 2)
9         for (int i = 0; i < n; i += len)
10             for (int j = 0; j < len / 2; j++) {
11                 int pos = n / len * (inv ? len - j : j);
12                 T u = a[i + j], v = a[i + j + len / 2] * rt[pos];
13                 a[i + j] = u + v, a[i + j + len / 2] = u - v;
14             }
15     if (T minv = T(1) / T(n); inv)
16         for (T &x : a) x *= minv;
17 }

```

Requires: Mod Struct

```

1 void ntt(vector<Mod> &a, bool inv, Mod primitive_root) {
2     int n = a.size();
3     Mod root = primitive_root ^ (MOD - 1) / n;
4     vector<Mod> rt(n + 1, 1);
5     for (int i = 0; i < n; i++) rt[i + 1] = rt[i] * root;
6     fft_(n, a, rt, inv);
7 }
8 void fft(vector<complex<double>> &a, bool inv) {
9     int n = a.size();
10    vector<complex<double>> rt(n + 1);
11    double arg = acos(-1) * 2 / n;
12    for (int i = 0; i <= n; i++)
13        rt[i] = {cos(arg * i), sin(arg * i)};
14    fft_(n, a, rt, inv);
15 }

```

### 5.4. Fast Walsh-Hadamard Transform

Requires: Mod Struct

```

1 void fwht(vector<Mod> &a, bool inv) {
2     int n = a.size();
3     for (int d = 1; d < n; d <= 1)
4         for (int m = 0; m < n; m++)
5             if (!(m & d)) {
6                 inv ? a[m | d] -= a[m] : a[m] += a[m | d]; // AND
7                 inv ? a[m | d] -= a[m] : a[m | d] += a[m]; // OR
8                 Mod x = a[m], y = a[m | d]; // XOR
9                 a[m] = x + y, a[m | d] = x - y; // XOR
10            }
11    if (Mod iv = Mod(1) / n; inv) // XOR
12        for (Mod &i : a) i *= iv; // XOR
13 }

```

## 5.5. Subset Convolution

Requires: Mod Struct

```

1 #pragma GCC target("popcnt")
2 #include <immintrin.h>
3
4 void fwht(int n, vector<vector<Mod>> &a, bool inv) {
5     for (int h = 0; h < n; h++)
6         for (int i = 0; i < (1 << n); i++)
7             if (!(i & (1 << h)))
8                 for (int k = 0; k <= n; k++)
9                     inv ? a[i | (1 << h)][k] -= a[i][k] :
10                        a[i | (1 << h)][k] += a[i][k];
11 }
12 // c[k] = sum(popcnt(i & j) == sz && i | j == k) a[i] * b[j]
13 vector<Mod> subset_convolution(int n, int sz,
14                               const vector<Mod> &a_,
15                               const vector<Mod> &b_) {
16     int len = n + sz + 1, N = 1 << n;
17     vector<vector<Mod>> a(1 << n, vector<Mod>(len, 0)), b = a;
18     for (int i = 0; i < N; i++)
19         a[i][_mm_popcnt_u64(i)] = a_[i],
20         b[i][_mm_popcnt_u64(i)] = b_[i];
21     fwht(n, a, 0), fwht(n, b, 0);
22     for (int i = 0; i < N; i++) {
23         vector<Mod> tmp(len);
24         for (int j = 0; j < len; j++)
25             for (int k = 0; k <= j; k++)
26                 tmp[j] += a[i][k] * b[i][j - k];
27         a[i] = tmp;
28     }
29     fwht(n, a, 1);
30     vector<Mod> c(N);
31     for (int i = 0; i < N; i++)
32         c[i] = a[i][_mm_popcnt_u64(i) + sz];
33     return c;
34 }

```

## 5.6. Linear Recurrences

### 5.6.1. Berlekamp-Massey Algorithm

```

1 template <typename T>
2 vector<T> berlekamp_massey(const vector<T> &s) {
3     int n = s.size(), l = 0, m = 1;
4     vector<T> r(n), p(n);
5     r[0] = p[0] = 1;
6     T b = 1, d = 0;
7     for (int i = 0; i < n; i++, m++, d = 0) {
8         for (int j = 0; j <= l; j++) d += r[j] * s[i - j];
9         if ((d /= b) == 0) continue; // change if T is float
10        auto t = r;
11        for (int j = m; j < n; j++) r[j] -= d * p[j - m];
12        if (l * 2 <= i) l = i + 1 - l, b = d, m = 0, p = t;
13    }
14    return r.resize(l + 1), reverse(r.begin(), r.end()), r;
15 }

```

### 5.6.2. Linear Recurrence Calculation

```

1 template <typename T> struct lin_rec {
2     using poly = vector<T>;
3     poly mul(poly a, poly b, poly m) {
4         int n = m.size();
5         poly r(n);
6         for (int i = n - 1; i >= 0; i--) {
7             r.insert(r.begin(), 0, r.pop_back());
8             T c = r[n - 1] + a[n - 1] * b[i];
9             // c /= m[n - 1]; if m is not monic
10            for (int j = 0; j < n; j++)
11                r[j] += a[j] * b[i] - c * m[j];
12        }
13        return r;
14    }
15    poly pow(poly p, ll k, poly m) {
16        poly r(m.size());
17        r[0] = 1;
18        for (; k >= 1, p = mul(p, p, m))
19            if (k & 1) r = mul(r, p, m);
20        return r;
21    }
22    T calc(poly t, poly r, ll k) {
23        int n = r.size();
24        poly p(n);
25        p[1] = 1;
26        poly q = pow(p, k, r);
27        T ans = 0;
28        for (int i = 0; i < n; i++) ans += t[i] * q[i];
29        return ans;
30    }
31 };

```

## 5.7. Matrices

### 5.7.1. Determinant

Requires: Mod Struct

```

1 Mod det(vector<vector<Mod>> a) {
2     int n = a.size();
3     Mod ans = 1;
4     for (int i = 0; i < n; i++) {
5         int b = i;
6         for (int j = i + 1; j < n; j++)
7             if (a[j][i] != 0) {
8                 b = j;
9                 break;
10            }
11        if (i != b) swap(a[i], a[b]), ans = -ans;
12        ans *= a[i][i];
13        if (ans == 0) return 0;
14        for (int j = i + 1; j < n; j++) {
15            Mod v = a[j][i] / a[i][i];
16            if (v != 0)
17                for (int k = i + 1; k < n; k++)
18                    a[j][k] -= v * a[i][k];
19        }
20    }
21    return ans;
22 }

```

```

1 double det(vector<vector<double>> a) {
2     int n = a.size();
3     double ans = 1;
4     for (int i = 0; i < n; i++) {
5         int b = i;
6         for (int j = i + 1; j < n; j++)
7             if (fabs(a[j][i]) > fabs(a[b][i])) b = j;
8         if (i != b) swap(a[i], a[b]), ans = -ans;
9         ans *= a[i][i];
10        if (ans == 0) return 0;
11        for (int j = i + 1; j < n; j++) {
12            double v = a[j][i] / a[i][i];
13            if (v != 0)
14                for (int k = i + 1; k < n; k++)
15                    a[j][k] -= v * a[i][k];
16        }
17    }
18    return ans;
19 }

```

### 5.7.2. Inverse

```

1 // Returns rank.
2 // Result is stored in A unless singular (rank < n).
3 // For prime powers, repeatedly set
4 // A^{-1} = A^{-1} (2I - A*A^{-1}) (mod p^k)
5 // where A^{-1} starts as the inverse of A mod p,
6 // and k is doubled in each step.
7
8 int matInv(vector<vector<double>> &A) {
9     int n = sz(A);
10    vi col(n);
11    vector<vector<double>> tmp(n, vector<double>(n));
12    rep(i, 0, n) tmp[i][i] = 1, col[i] = i;
13
14    rep(i, 0, n) {
15        int r = i, c = i;
16        rep(j, i, n)
17            rep(k, i, n) if (fabs(A[j][k]) > fabs(A[r][c])) r = j, c = k;
18
19        if (fabs(A[r][c]) < 1e-12) return i;
20        A[i].swap(A[r]);
21        tmp[i].swap(tmp[r]);
22        rep(j, 0, n) swap(A[j][i], A[j][c]),
23        swap(tmp[j][i], tmp[j][c]);
24        swap(col[i], col[c]);
25        double v = A[i][i];
26        rep(j, i + 1, n) {
27            double f = A[j][i] / v;
28            A[j][i] = 0;
29            rep(k, i + 1, n) A[j][k] -= f * A[i][k];
30            rep(k, 0, n) tmp[j][k] -= f * tmp[i][k];
31        }
32        rep(j, i + 1, n) A[i][j] /= v;
33        rep(j, 0, n) tmp[i][j] /= v;
34        A[i][i] = 1;
35    }
36
37    for (int i = n - 1; i > 0; --i) rep(j, 0, i) {
38        double v = A[j][i];
39    }

```

```

    rep(k, 0, n) tmp[j][k] -= v * tmp[i][k];
41 }
43 rep(i, 0, n) rep(j, 0, n) A[col[i]][col[j]] = tmp[i][j];
45 return n;
47 int matInv_mod(vector<vector<ll>> &A) {
48     int n = sz(A);
49     vi col(n);
50     vector<vector<ll>> tmp(n, vector<ll>(n));
51     rep(i, 0, n) tmp[i][i] = 1, col[i] = i;
53     rep(i, 0, n) {
54         int r = i, c = i;
55         rep(j, i, n) rep(k, i, n) if (A[j][k]) {
56             r = j;
57             c = k;
58             goto found;
59         }
60         return i;
61     found:
62         A[i].swap(A[r]);
63         tmp[i].swap(tmp[r]);
64         rep(j, 0, n) swap(A[j][i], A[j][c]),
65         swap(tmp[j][i], tmp[j][c]);
66         swap(col[i], col[c]);
67         ll v = modpow(A[i][i], mod - 2);
68         rep(j, i + 1, n) {
69             ll f = A[j][i] * v % mod;
70             A[j][i] = 0;
71             rep(k, i + 1, n) A[j][k] =
72             (A[j][k] - f * A[i][k]) % mod;
73             rep(k, 0, n) tmp[j][k] =
74             (tmp[j][k] - f * tmp[i][k]) % mod;
75         }
76         rep(j, i + 1, n) A[i][j] = A[i][j] * v % mod;
77         rep(j, 0, n) tmp[i][j] = tmp[i][j] * v % mod;
78         A[i][i] = 1;
79     }
81     for (int i = n - 1; i > 0; --i) rep(j, 0, i) {
82         ll v = A[j][i];
83         rep(k, 0, n) tmp[j][k] =
84         (tmp[j][k] - v * tmp[i][k]) % mod;
85     }
87     rep(i, 0, n) rep(j, 0, n) A[col[i]][col[j]] =
88     tmp[i][j] % mod + (tmp[i][j] < 0 ? mod : 0);
89     return n;
90 }

```

### 5.7.3. Characteristic Polynomial

```

1 // calculate det(a - xI)
2 template <typename T>
3 vector<T> CharacteristicPolynomial(vector<vector<T>> a) {
4     int N = a.size();
5     for (int j = 0; j < N - 2; j++) {
6         for (int i = j + 1; i < N; i++) {
7             if (a[i][j] != 0) {
8                 swap(a[j + 1], a[i]);
9                 for (int k = 0; k < N; k++)
10                     swap(a[k][j + 1], a[k][i]);
11                 break;
12             }
13         }
14         if (a[j + 1][j] != 0) {
15             T inv = T(1) / a[j + 1][j];
16             for (int i = j + 2; i < N; i++) {
17                 if (a[i][j] == 0) continue;
18                 T coe = inv * a[i][j];
19                 for (int l = j; l < N; l++)
20                     a[i][l] -= coe * a[j + 1][l];
21                 for (int k = 0; k < N; k++)
22                     a[k][j + 1] += coe * a[k][i];
23             }
24         }
25     }
26     vector<vector<T>> p(N + 1);
27     p[0] = {T(1)};
28     for (int i = 1; i <= N; i++) {
29         p[i].resize(i + 1);
30         for (int j = 0; j < i; j++) {
31             p[i][j + 1] = p[i - 1][j];
32             p[i][j] += p[i - 1][j] * a[i - 1][i - 1];
33         }
34         T x = 1;

```

```

37     for (int m = 1; m < i; m++) {
38         x *= -a[i - m][i - m - 1];
39         T coe = x * a[i - m - 1][i - 1];
40         for (int j = 0; j < i - m; j++)
41             p[i][j] += coe * p[i - m - 1][j];
42     }
43     return p[N];
44 }

```

### 5.7.4. Solve Linear Equation

```

1 typedef vector<double> vd;
2 const double eps = 1e-12;
3 // solves for x: A * x = b
4 int solveLinear(vector<vd> &A, vd &b, vd &x) {
5     int n = sz(A), m = sz(x), rank = 0, br, bc;
6     if (n) assert(sz(A[0]) == m);
7     vi col(m);
8     iota(all(col), 0);
9     rep(i, 0, n) {
10         double v, bv = 0;
11         rep(r, i, n) rep(c, i, m) if ((v = fabs(A[r][c])) > bv)
12             br = r, bc = c, bv = v;
13         if (bv <= eps) {
14             rep(j, i, n) if (fabs(b[j]) > eps) return -1;
15             break;
16         }
17         swap(A[i], A[br]);
18         swap(b[i], b[br]);
19         swap(col[i], col[bc]);
20         rep(j, 0, n) swap(A[j][i], A[j][bc]);
21         bv = 1 / A[i][i];
22         rep(j, i + 1, n) {
23             double fac = A[j][i] * bv;
24             b[j] -= fac * b[i];
25             rep(k, i + 1, m) A[j][k] -= fac * A[i][k];
26         }
27         rank++;
28     }
29     x.assign(m, 0);
30     for (int i = rank; i--;) {
31         b[i] /= A[i][i];
32         x[col[i]] = b[i];
33         rep(j, 0, i) b[j] -= A[j][i] * b[i];
34     }
35     return rank; // (multiple solutions if rank < m)
36 }

```

### 5.8. Polynomial Interpolation

```

1 // returns a, such that a[0]x^0 + a[1]x^1 + a[2]x^2 + ...
2 // passes through the given points
3 typedef vector<double> vd;
4 vd interpolate(vd x, vd y, int n) {
5     vd res(n), temp(n);
6     rep(k, 0, n - 1) rep(i, k + 1, n) y[i] =
7     (y[i] - y[k]) / (x[i] - x[k]);
8     double last = 0;
9     temp[0] = 1;
10     rep(k, 0, n) rep(i, 0, n) {
11         res[i] += y[k] * temp[i];
12         swap(last, temp[i]);
13         temp[i] -= last * x[k];
14     }
15     return res;
16 }

```

### 5.9. Simplex Algorithm

```

1 // Two-phase simplex algorithm for solving linear programs
2 // of the form
3 //
4 //      maximize      c^T x
5 //      subject to    Ax <= b
6 //                   x >= 0
7 //
8 // INPUT: A -- an m x n matrix
9 //         b -- an m-dimensional vector
10 //         c -- an n-dimensional vector
11 //         x -- a vector where the optimal solution will be
12 //             stored
13 //
14 // OUTPUT: value of the optimal solution (infinity if
15 //         unbounded
16 //         above, nan if infeasible)

```



```

17 //
18 // To use this code, create an LPSolver object with A, b,
19 // and c as arguments. Then, call Solve(x).
21 typedef long double ld;
22 typedef vector<ld> vd;
23 typedef vector<vd> vvd;
24 typedef vector<int> vi;
25
26 const ld EPS = 1e-9;
27
28 struct LPSolver {
29     int m, n;
30     vi B, N;
31     vvd D;
32
33     LPSolver(const vvd &A, const vd &b, const vd &c)
34         : m(b.size()), n(c.size()), N(n + 1), B(m),
35           D(m + 2, vd(n + 2)) {
36         for (int i = 0; i < m; i++)
37             for (int j = 0; j < n; j++) D[i][j] = A[i][j];
38         for (int i = 0; i < m; i++) {
39             B[i] = n + i;
40             D[i][n] = -1;
41             D[i][n + 1] = b[i];
42         }
43         for (int j = 0; j < n; j++) {
44             N[j] = j;
45             D[m][j] = -c[j];
46         }
47         N[n] = -1;
48         D[m + 1][n] = 1;
49     }
50
51     void Pivot(int r, int s) {
52         double inv = 1.0 / D[r][s];
53         for (int i = 0; i < m + 2; i++)
54             if (i != r)
55                 for (int j = 0; j < n + 2; j++)
56                     if (j != s) D[i][j] -= D[r][j] * D[i][s] * inv;
57         for (int j = 0; j < n + 2; j++)
58             if (j != s) D[r][j] *= inv;
59         for (int i = 0; i < m + 2; i++)
60             if (i != r) D[i][s] *= -inv;
61         D[r][s] = inv;
62         swap(B[r], N[s]);
63     }
64
65     bool Simplex(int phase) {
66         int x = phase == 1 ? m + 1 : m;
67         while (true) {
68             int s = -1;
69             for (int j = 0; j <= n; j++) {
70                 if (phase == 2 && N[j] == -1) continue;
71                 if (s == -1 || D[x][j] < D[x][s] ||
72                     D[x][j] == D[x][s] && N[j] < N[s])
73                     s = j;
74             }
75             if (D[x][s] > -EPS) return true;
76             int r = -1;
77             for (int i = 0; i < m; i++) {
78                 if (D[i][s] < EPS) continue;
79                 if (r == -1 ||
80                     D[i][n + 1] / D[i][s] < D[r][n + 1] / D[r][s] ||
81                     (D[i][n + 1] / D[i][s] ==
82                      D[r][n + 1] / D[r][s]) &&
83                     B[i] < B[r])
84                     r = i;
85             }
86             if (r == -1) return false;
87             Pivot(r, s);
88         }
89     }
90
91     ld Solve(vd &x) {
92         int r = 0;
93         for (int i = 1; i < m; i++)
94             if (D[i][n + 1] < D[r][n + 1]) r = i;
95         if (D[r][n + 1] < -EPS) {
96             Pivot(r, n);
97             if (!Simplex(1) || D[m + 1][n + 1] < -EPS)
98                 return -numeric_limits<ld>::infinity();
99             for (int i = 0; i < m; i++)
100                 if (B[i] == -1) {
101                     int s = -1;
102                     for (int j = 0; j <= n; j++)
103                         if (s == -1 || D[i][j] < D[i][s] ||
104                             D[i][j] == D[i][s] && N[j] < N[s])
105                             s = j;
106                     Pivot(i, s);

```

```

107     }
108 }
109 if (!Simplex(2)) return numeric_limits<ld>::infinity();
110 x = vd(n);
111 for (int i = 0; i < m; i++)
112     if (B[i] < n) x[B[i]] = D[i][n + 1];
113 return D[m][n + 1];
114 };
115
116 int main() {
117     const int m = 4;
118     const int n = 3;
119     ld _A[m][n] = {
120         {6, -1, 0}, {-1, -5, 0}, {1, 5, 1}, {-1, -5, -1}};
121     ld _b[m] = {10, -4, 5, -5};
122     ld _c[n] = {1, -1, 0};
123
124     vvd A(m);
125     vd b(_b, _b + m);
126     vd c(_c, _c + n);
127     for (int i = 0; i < m; i++) A[i] = vd(_A[i], _A[i] + n);
128
129     LPSolver solver(A, b, c);
130     vd x;
131     ld value = solver.Solve(x);
132
133     cerr << "VALUE: " << value << endl; // VALUE: 1.29032
134     cerr << "SOLUTION: "; // SOLUTION: 1.74194 0.451613 1
135     for (size_t i = 0; i < x.size(); i++) cerr << " " << x[i];
136     cerr << endl;
137     return 0;
138 }

```

## 6. Geometry

### 6.1. Point

```

1 template <typename T> struct P {
2     T x, y;
3     P(T x = 0, T y = 0) : x(x), y(y) {}
4     bool operator<(const P &p) const {
5         return tie(x, y) < tie(p.x, p.y);
6     }
7     bool operator==(const P &p) const {
8         return tie(x, y) == tie(p.x, p.y);
9     }
10    P operator-() const { return {-x, -y}; }
11    P operator+(P p) const { return {x + p.x, y + p.y}; }
12    P operator-(P p) const { return {x - p.x, y - p.y}; }
13    P operator*(T d) const { return {x * d, y * d}; }
14    P operator/(T d) const { return {x / d, y / d}; }
15    T dist2() const { return x * x + y * y; }
16    double len() const { return sqrt(dist2()); }
17    P unit() const { return *this / len(); }
18    friend T dot(P a, P b) { return a.x * b.x + a.y * b.y; }
19    friend T cross(P a, P b) { return a.x * b.y - a.y * b.x; }
20    friend T cross(P a, P b, P o) {
21        return cross(a - o, b - o);
22    }
23 };
24 using pt = P<ll>;

```

#### 6.1.1. Quaternion

```

1 constexpr double PI = 3.141592653589793;
2 constexpr double EPS = 1e-7;
3 struct Q {
4     using T = double;
5     T x, y, z, r;
6     Q(T r = 0) : x(0), y(0), z(0), r(r) {}
7     Q(T x, T y, T z, T r = 0) : x(x), y(y), z(z), r(r) {}
8     friend bool operator==(const Q &a, const Q &b) {
9         return (a - b).abs2() <= EPS;
10    }
11    friend bool operator!=(const Q &a, const Q &b) {
12        return !(a == b);
13    }
14    Q operator-() const { return {-x, -y, -z, -r}; }
15    Q operator+(const Q &b) const {
16        return Q(x + b.x, y + b.y, z + b.z, r + b.r);
17    }
18    Q operator-(const Q &b) const {
19        return Q(x - b.x, y - b.y, z - b.z, r - b.r);
20    }
21    Q operator*(const T &t) const {
22        return Q(x * t, y * t, z * t, r * t);
23    }

```

```

25 Q operator*(const Q &b) const {
    return Q(r * b.x + x * b.r + y * b.z - z * b.y,
            r * b.y - x * b.z + y * b.r + z * b.x,
            r * b.z + x * b.y - y * b.x + z * b.r,
            r * b.r - x * b.x - y * b.y - z * b.z);
29 }
31 Q operator/(const Q &b) const { return *this * b.inv(); }
33 T abs2() const { return r * r + x * x + y * y + z * z; }
35 T len() const { return sqrt(abs2()); }
37 Q conj() const { return Q(-x, -y, -z, r); }
39 Q unit() const { return *this * (1.0 / len()); }
41 Q inv() const { return conj() * (1.0 / abs2()); }
43 friend T dot(Q a, Q b) {
    return a.x * b.x + a.y * b.y + a.z * b.z;
45 }
47 friend Q cross(Q a, Q b) {
    return Q(a.y * b.z - a.z * b.y, a.z * b.x - a.x * b.z,
            a.x * b.y - a.y * b.x);
49 }
51 friend Q rotation_around(Q axis, T angle) {
    return axis.unit() * sin(angle / 2) + cos(angle / 2);
53 }
55 Q rotated_around(Q axis, T angle) {
    Q u = rotation_around(axis, angle);
    return u * *this / u;
57 }
59 friend Q rotation_between(Q a, Q b) {
    a = a.unit(), b = b.unit();
    if (a == -b) {
        // degenerate case
        Q ortho = abs(a.y) > EPS ? cross(a, Q(1, 0, 0))
                                : cross(a, Q(0, 1, 0));
        return rotation_around(ortho, PI);
    }
    return (a * (a + b)).conj();
}

```

### 6.1.2. Spherical Coordinates

```

1 struct car_p {
    double x, y, z;
2 };
3 struct sph_p {
    double r, theta, phi;
4 };
5 sph_p conv(car_p p) {
    double r = sqrt(p.x * p.x + p.y * p.y + p.z * p.z);
    double theta = asin(p.y / r);
    double phi = atan2(p.y, p.x);
    return {r, theta, phi};
6 }
7 car_p conv(sph_p p) {
    double x = p.r * cos(p.theta) * sin(p.phi);
    double y = p.r * cos(p.theta) * cos(p.phi);
    double z = p.r * sin(p.theta);
    return {x, y, z};
8 }
9 }

```

### 6.2. Segments

```

1 // for non-collinear ABCD, if segments AB and CD intersect
2 bool intersects(pt a, pt b, pt c, pt d) {
3     if (cross(b, c, a) * cross(b, d, a) > 0) return false;
4     if (cross(d, a, c) * cross(d, b, c) > 0) return false;
5     return true;
6 }
7 // the intersection point of lines AB and CD
8 pt intersect(pt a, pt b, pt c, pt d) {
9     auto x = cross(b, c, a), y = cross(b, d, a);
10    if (x == y) {
11        // if(abs(x, y) < 1e-8) {
12        // is parallel
13    } else {
14        return d * (x / (x - y)) - c * (y / (x - y));
15    }
16 }

```

### 6.3. Convex Hull

```

1 // returns a convex hull in counterclockwise order
2 // for a non-strict one, change cross >= to >
3 vector<pt> convex_hull(vector<pt> p) {
4     sort(ALL(p));
5     if (p[0] == p.back()) return {p[0]};
6     int n = p.size(), t = 0;
7     vector<pt> h(n + 1);
8     for (int _ = 2, s = 0; _--; s = --t, reverse(ALL(p)))
9         for (pt i : p) {

```

```

11         while (t > s + 1 && cross(i, h[t - 1], h[t - 2]) >= 0)
12             t--;
13         h[t++] = i;
14     }
15     return h.resize(t), h;

```

#### 6.3.1. 3D Hull

```

1 typedef Point3D<double> P3;
2
3 struct PR {
4     void ins(int x) { (a == -1 ? a : b) = x; }
5     void rem(int x) { (a == x ? a : b) = -1; }
6     int cnt() { return (a != -1) + (b != -1); }
7     int a, b;
8 };
9
10 struct F {
11     P3 q;
12     int a, b, c;
13 };
14
15 vector<F> hull3d(const vector<P3> &A) {
16     assert(sz(A) >= 4);
17     vector<vector<PR>> E(sz(A), vector<PR>(sz(A), {-1, -1}));
18     #define E(x, y) E[f.x][f.y]
19     vector<F> FS;
20     auto mf = [&](int i, int j, int k, int l) {
21         P3 q = (A[j] - A[i]).cross((A[k] - A[i]));
22         if (q.dot(A[l]) > q.dot(A[i])) q = q * -1;
23         F f{i, j, k};
24         E(a, b).ins(k);
25         E(a, c).ins(j);
26         E(b, c).ins(i);
27         FS.push_back(f);
28     };
29     rep(i, 0, 4) rep(j, i + 1, 4) rep(k, j + 1, 4)
30     mf(i, j, k, 6 - i - j - k);
31
32     rep(i, 4, sz(A)) {
33         rep(j, 0, sz(FS)) {
34             F f = FS[j];
35             if (f.q.dot(A[i]) > f.q.dot(A[f.a])) {
36                 E(a, b).rem(f.c);
37                 E(a, c).rem(f.b);
38                 E(b, c).rem(f.a);
39                 swap(FS[j--], FS.back());
40                 FS.pop_back();
41             }
42         }
43         int nw = sz(FS);
44         rep(j, 0, nw) {
45             F f = FS[j];
46             #define C(a, b, c)
47             if (E(a, b).cnt() != 2) mf(f.a, f.b, i, f.c);
48             C(a, b, c);
49             C(a, c, b);
50             C(b, c, a);
51         }
52     }
53     for (F &it : FS)
54         if ((A[it.b] - A[it.a])
55             .cross(A[it.c] - A[it.a])
56             .dot(it.q) <= 0)
57             swap(it.c, it.b);
58     return FS;
59 }

```

#### 6.4. Angular Sort

```

1 auto angle_cmp = [](const pt &a, const pt &b) {
2     auto btm = [](const pt &a) {
3         return a.y < 0 || (a.y == 0 && a.x < 0);
4     };
5     return make_tuple(btm(a), a.y * b.x, abs2(a)) <
6            make_tuple(btm(b), a.x * b.y, abs2(b));
7 };
8 void angular_sort(vector<pt> &p) {
9     sort(p.begin(), p.end(), angle_cmp);
10 }

```

#### 6.5. Convex Polygon Minkowski Sum

```

1 // O(n) convex polygon minkowski sum
2 // must be sorted and counterclockwise
3 vector<pt> minkowski_sum(vector<pt> p, vector<pt> q) {
4     auto diff = [](vector<pt> &c) {
5         auto rcmp = [](pt a, pt b) {
6             return pt{a.y, a.x} < pt{b.y, b.x};

```

```

7   };
8   rotate(c.begin(), min_element(ALL(c), rcmp), c.end());
9   c.push_back(c[0]);
10  vector<pt> ret;
11  for (int i = 1; i < c.size(); i++)
12      ret.push_back(c[i] - c[i - 1]);
13  return ret;
14  };
15  auto dp = diff(p), dq = diff(q);
16  pt cur = p[0] + q[0];
17  vector<pt> d(dp.size() + dq.size(), ret = {cur});
18  // include angle_cmp from angular-sort.cpp
19  merge(ALL(dp), ALL(dq), d.begin(), angle_cmp);
20  // optional: make ret strictly convex (UB if degenerate)
21  int now = 0;
22  for (int i = 1; i < d.size(); i++) {
23      if (cross(d[i], d[now]) == 0) d[now] = d[now] + d[i];
24      else d[++now] = d[i];
25  }
26  d.resize(now + 1);
27  // end optional part
28  for (pt v : d) ret.push_back(cur = cur + v);
29  return ret.pop_back(), ret;
30  }

```

## 6.6. Point In Polygon

```

1  bool on_segment(pt a, pt b, pt p) {
2      return cross(a, b, p) == 0 && dot((p - a), (p - b)) <= 0;
3  }
4  // p can be any polygon, but this is O(n)
5  bool inside(const vector<pt> &p, pt a) {
6      int cnt = 0, n = p.size();
7      for (int i = 0; i < n; i++) {
8          pt l = p[i], r = p[(i + 1) % n];
9          // change to return 0; for strict version
10         if (on_segment(l, r, a)) return 1;
11         cnt ^= ((a.y < l.y) - (a.y < r.y)) * cross(l, r, a) > 0;
12     }
13     return cnt;
14 }

```

### 6.6.1. Convex Version

```

1  // no preprocessing version
2  // p must be a strict convex hull, counterclockwise
3  // if point is inside or on border
4  bool is_inside(const vector<pt> &c, pt p) {
5      int n = c.size(), l = 1, r = n - 1;
6      if (cross(c[0], c[1], p) < 0) return false;
7      if (cross(c[n - 1], c[0], p) < 0) return false;
8      while (l < r - 1) {
9          int m = (l + r) / 2;
10         T a = cross(c[0], c[m], p);
11         if (a > 0) l = m;
12         else if (a < 0) r = m;
13         else return dot(c[0] - p, c[m] - p) <= 0;
14     }
15     if (l == r) return dot(c[0] - p, c[l] - p) <= 0;
16     else return cross(c[l], c[r], p) >= 0;
17 }
18
19 // with preprocessing version
20 vector<pt> vecs;
21 pt center;
22 // p must be a strict convex hull, counterclockwise
23 // BEWARE OF OVERFLOWS!!
24 void preprocess(vector<pt> p) {
25     for (auto &v : p) v = v * 3;
26     center = p[0] + p[1] + p[2];
27     center.x /= 3, center.y /= 3;
28     for (auto &v : p) v = v - center;
29     vecs = (angular_sort(p), p);
30 }
31 bool intersect_strict(pt a, pt b, pt c, pt d) {
32     if (cross(b, c, a) * cross(b, d, a) > 0) return false;
33     if (cross(d, a, c) * cross(d, b, c) >= 0) return false;
34     return true;
35 }
36 // if point is inside or on border
37 bool query(pt p) {
38     p = p * 3 - center;
39     auto pr = upper_bound(ALL(vecs), p, angle_cmp);
40     if (pr == vecs.end()) pr = vecs.begin();
41     auto pl = (pr == vecs.begin()) ? vecs.back() : *(pr - 1);
42     return !intersect_strict({0, 0}, p, pl, *pr);
43 }

```

## 6.6.2. Offline Multiple Points Version

Requires: Point, GNU PBDS

```

1  using Double = __float128;
2  using Point = pt<Double, Double>;
3
4  int n, m;
5  vector<Point> poly;
6  vector<Point> query;
7  vector<int> ans;
8
9  struct Segment {
10     Point a, b;
11     int id;
12 };
13 vector<Segment> segs;
14
15 Double Xnow;
16 inline Double get_y(const Segment &u, Double xnow = Xnow) {
17     const Point &a = u.a;
18     const Point &b = u.b;
19     return (a.y * (b.x - xnow) + b.y * (xnow - a.x)) /
20         (b.x - a.x);
21 }
22
23 bool operator<(Segment u, Segment v) {
24     Double yu = get_y(u);
25     Double yv = get_y(v);
26     if (yu != yv) return yu < yv;
27     return u.id < v.id;
28 }
29
30 ordered_map<Segment> st;
31
32 struct Event {
33     int type; // +1 insert seg, -1 remove seg, 0 query
34     Double x, y;
35     int id;
36 };
37
38 bool operator<(Event a, Event b) {
39     if (a.x != b.x) return a.x < b.x;
40     if (a.type != b.type) return a.type < b.type;
41     return a.y < b.y;
42 }
43
44 vector<Event> events;
45
46 void solve() {
47     set<Double> xs;
48     set<Point> ps;
49     for (int i = 0; i < n; i++) {
50         xs.insert(poly[i].x);
51         ps.insert(poly[i]);
52     }
53     for (int i = 0; i < n; i++) {
54         Segment s{poly[i], poly[(i + 1) % n], i};
55         if (s.a.x > s.b.x ||
56             (s.a.x == s.b.x && s.a.y > s.b.y)) {
57             swap(s.a, s.b);
58         }
59         segs.push_back(s);
60     }
61     for (int i = 0; i < m; i++) {
62         events.push_back({+1, s.a.x + 0.2, s.a.y, i});
63         events.push_back({-1, s.b.x - 0.2, s.b.y, i});
64     }
65     sort(events.begin(), events.end());
66     int cnt = 0;
67     for (Event e : events) {
68         int i = e.id;
69         Xnow = e.x;
70         if (e.type == 0) {
71             Double x = e.x;
72             Double y = e.y;
73             Segment tmp = {{x - 1, y}, {x + 1, y}, -1};
74             auto it = st.lower_bound(tmp);
75
76             if (ps.count(query[i]) > 0) {
77                 ans[i] = 0;
78             } else if (xs.count(x) > 0) {
79                 ans[i] = -2;
80             } else if (it != st.end() &&
81                 get_y(*it) == get_y(tmp)) {
82                 ans[i] = 0;
83             } else if (it != st.begin() &&
84                 get_y(*prev(it)) == get_y(tmp)) {
85                 ans[i] = 0;
86             } else {
87                 int rk = st.order_of_key(tmp);
88                 if (rk % 2 == 1) {

```

```

89     ans[i] = 1;
90   } else {
91     ans[i] = -1;
92   }
93 }
94 } else if (e.type == 1) {
95   st.insert(segs[i]);
96   assert((int)st.size() == ++cnt);
97 } else if (e.type == -1) {
98   st.erase(segs[i]);
99   assert((int)st.size() == --cnt);
100 }
101 }

```

### 6.7. Closest Pair

```

1 vector<pll> p; // sort by x first!
2 bool cmpy(const pll &a, const pll &b) const {
3   return a.y < b.y;
4 }
5 ll sq(ll x) { return x * x; }
6 // returns (minimum dist)^2 in [l, r]
7 ll solve(int l, int r) {
8   if (r - l <= 1) return 1e18;
9   int m = (l + r) / 2;
10  ll mid = p[m].x, d = min(solve(l, m), solve(m, r));
11  auto pb = p.begin();
12  inplace_merge(pb + l, pb + m, pb + r, cmpy);
13  vector<pll> s;
14  for (int i = l; i < r; i++)
15    if (sq(p[i].x - mid) < d) s.push_back(p[i]);
16  for (int i = 0; i < s.size(); i++)
17    for (int j = i + 1;
18         j < s.size() && sq(s[j].y - s[i].y) < d; j++)
19      d = min(d, dis(s[i], s[j]));
20  return d;
21 }

```

### 6.8. Minimum Enclosing Circle

```

1 typedef Point<double> P;
2 double ccRadius(const P &A, const P &B, const P &C) {
3   return (B - A).dist() * (C - B).dist() * (A - C).dist() /
4     abs((B - A).cross(C - A)) / 2;
5 }
6 P ccCenter(const P &A, const P &B, const P &C) {
7   P b = C - A, c = B - A;
8   return A + (b * c.dist2() - c * b.dist2()).perp() /
9     b.cross(c) / 2;
10 }
11 pair<P, double> mec(vector<P> ps) {
12   shuffle(all(ps), mt19937(time(0)));
13   P o = ps[0];
14   double r = 0, EPS = 1 + 1e-8;
15   rep(i, 0, sz(ps)) if ((o - ps[i]).dist() > r * EPS) {
16     o = ps[i], r = 0;
17     rep(j, 0, i) if ((o - ps[j]).dist() > r * EPS) {
18       o = (ps[i] + ps[j]) / 2;
19       r = (o - ps[i]).dist();
20       rep(k, 0, j) if ((o - ps[k]).dist() > r * EPS) {
21         o = ccCenter(ps[i], ps[j], ps[k]);
22         r = (o - ps[i]).dist();
23       }
24     }
25   }
26   return {o, r};
27 }

```

### 6.9. Delaunay Triangulation

```

1 typedef Point<ll> P;
2 typedef struct Quad *Q;
3 typedef __int128_t lll; // (can be ll if coords are < 2e4)
4 P arb(LLONG_MAX, LLONG_MAX); // not equal to any other point
5
6 struct Quad {
7   bool mark;
8   Q o, rot;
9   P p;
10  P F() { return r()->p; }
11  Q r() { return rot->rot; }
12  Q prev() { return rot->o->rot; }
13  Q next() { return r()->prev(); }
14 };
15
16 bool circ(P p, P a, P b, P c) { // is p in the circumcircle?
17   lll p2 = p.dist2(), A = a.dist2() - p2,
18   B = b.dist2() - p2, C = c.dist2() - p2;
19   return p.cross(a, b) * C + p.cross(b, c) * A +

```

```

    p.cross(c, a) * B >
    0;
21 }
22 Q makeEdge(P orig, P dest) {
23   Q q[] = {new Quad{0, 0, 0, orig}, new Quad{0, 0, 0, arb},
24     new Quad{0, 0, 0, dest}, new Quad{0, 0, 0, arb}};
25   rep(i, 0, 4) q[i]->o = q[(i + 1) % 3];
26   return *q;
27 }
28 void splice(Q a, Q b) {
29   swap(a->o->rot->o, b->o->rot->o);
30   swap(a->o, b->o);
31 }
32 Q connect(Q a, Q b) {
33   Q q = makeEdge(a->F(), b->p);
34   splice(q, a->next());
35   splice(q->r(), b);
36   return q;
37 }
38 pair<Q, Q> rec(const vector<P> &s) {
39   if (sz(s) <= 3) {
40     Q a = makeEdge(s[0], s[1]),
41     b = makeEdge(s[1], s.back());
42     if (sz(s) == 2) return {a, a->r()};
43     splice(a->r(), b);
44     auto side = s[0].cross(s[1], s[2]);
45     Q c = side > 0 ? connect(b, a) : 0;
46     return {side < 0 ? c->r() : a, side < 0 ? c : b->r()};
47   }
48   #define H(e) e->F(), e->p
49   #define valid(e) (e->F().cross(H(base)) > 0)
50   Q A, B, ra, rb;
51   int half = sz(s) / 2;
52   tie(ra, A) = rec({all(s) - half});
53   tie(B, rb) = rec({sz(s) - half + all(s)});
54   while ((B->p.cross(H(A)) < 0 && (A = A->next())) ||
55     (A->p.cross(H(B)) > 0 && (B = B->r()->o)))
56     ;
57   Q base = connect(B->r(), A);
58   if (A->p == ra->p) ra = base->r();
59   if (B->p == rb->p) rb = base;
60
61   #define DEL(e, init, dir)
62   Q e = init->dir;
63   if (valid(e))
64     while (circ(e->dir->F(), H(base), e->F())) {
65       Q t = e->dir;
66       splice(e, e->prev());
67       splice(e->r(), e->r()->prev());
68       e = t;
69     }
70   for (;) {
71     DEL(LC, base->r(), o);
72     DEL(RC, base, prev());
73     if (!valid(LC) && !valid(RC)) break;
74     if (!valid(LC) || (valid(RC) && circ(H(RC), H(LC))))
75       base = connect(RC, base->r());
76     else base = connect(base->r(), LC->r());
77   }
78   return {ra, rb};
79 }
80
81 // returns [A_0, B_0, C_0, A_1, B_1, ...]
82 // where A_i, B_i, C_i are counter-clockwise triangles
83 vector<P> triangulate(vector<P> pts) {
84   sort(all(pts));
85   assert(unique(all(pts)) == pts.end());
86   if (sz(pts) < 2) return {};
87   Q e = rec(pts).first;
88   vector<Q> q = {e};
89   int qi = 0;
90   while (e->o->F().cross(e->F(), e->p) < 0) e = e->o;
91   #define ADD
92   {
93     Q c = e;
94     do {
95       c->mark = 1;
96       pts.push_back(c->p);
97       q.push_back(c->r());
98       c = c->next();
99     } while (c != e);
100   }
101   ADD;
102   pts.clear();
103   while (qi < sz(q))
104     if (!(e = q[qi++])->mark) ADD;
105   return pts;
106 }

```

}

### 6.9.1. Slower Version

```

1 template <class P, class F>
2 void delaunay(vector<P> &ps, F trifun) {
3     if (sz(ps) == 3) {
4         int d = (ps[0].cross(ps[1], ps[2]) < 0);
5         trifun(0, 1 + d, 2 - d);
6     }
7     vector<P> p3;
8     for (P p : ps) p3.emplace_back(p.x, p.y, p.dist2());
9     if (sz(p3) > 3)
10         for (auto t : hull3d(p3))
11             if ((p3[t.b] - p3[t.a])
12                 .cross(p3[t.c] - p3[t.a])
13                 .dot(P3(0, 0, 1)) < 0)
14                 trifun(t.a, t.c, t.b);
15 }

```

### 6.10. Half Plane Intersection

```

1 struct Line {
2     Point P;
3     Vector v;
4     bool operator<(const Line &b) const {
5         return atan2(v.y, v.x) < atan2(b.v.y, b.v.x);
6     }
7 };
8 bool OnLeft(const Line &L, const Point &p) {
9     return Cross(L.v, p - L.P) > 0;
10 }
11 Point GetIntersection(Line a, Line b) {
12     Vector u = a.P - b.P;
13     Double t = Cross(b.v, u) / Cross(a.v, b.v);
14     return a.P + a.v * t;
15 }
16 int HalfplaneIntersection(Line *L, int n, Point *poly) {
17     sort(L, L + n);
18
19     int first, last;
20     Point *p = new Point[n];
21     Line *q = new Line[n];
22     q[first = last = 0] = L[0];
23     for (int i = 1; i < n; i++) {
24         while (first < last && !OnLeft(L[i], p[last - 1]))
25             last--;
26         while (first < last && !OnLeft(L[i], p[first])) first++;
27         q[++last] = L[i];
28         if (fabs(Cross(q[last].v, q[last - 1].v)) < EPS) {
29             last--;
30             if (OnLeft(q[last], L[i].P)) q[last] = L[i];
31         }
32         if (first < last)
33             p[last - 1] = GetIntersection(q[last - 1], q[last]);
34     }
35     while (first < last && !OnLeft(q[first], p[last - 1]))
36         last--;
37     if (last - first <= 1) return 0;
38     p[last] = GetIntersection(q[last], q[first]);
39
40     int m = 0;
41     for (int i = first; i <= last; i++) poly[m++] = p[i];
42     return m;
43 }

```

## 7. Strings

### 7.1. Knuth-Morris-Pratt Algorithm

```

1 vector<int> pi(const string &s) {
2     vector<int> p(s.size());
3     for (int i = 1; i < s.size(); i++) {
4         int g = p[i - 1];
5         while (g && s[i] != s[g]) g = p[g - 1];
6         p[i] = g + (s[i] == s[g]);
7     }
8     return p;
9 }
10 vector<int> match(const string &s, const string &pat) {
11     vector<int> p = pi(pat + '\0' + s), res;
12     for (int i = p.size() - s.size(); i < p.size(); i++)
13         if (p[i] == pat.size())
14             res.push_back(i - 2 * pat.size());
15     return res;
16 }

```

### 7.2. Aho-Corasick Automaton

```

1 struct Aho_Corasick {
2     static const int maxc = 26, maxn = 4e5;
3     struct NODES {
4         int Next[maxc], fail, ans;
5     };
6     NODES T[maxn];
7     int top, qtop, q[maxn];
8     int get_node(const int &fail) {
9         fill_n(T[top].Next, maxc, 0);
10        T[top].fail = fail;
11        T[top].ans = 0;
12        return top++;
13    }
14    int insert(const string &s) {
15        int ptr = 1;
16        for (char c : s) { // change char id
17            c -= 'a';
18            if (!T[ptr].Next[c]) T[ptr].Next[c] = get_node(ptr);
19            ptr = T[ptr].Next[c];
20        }
21        return ptr;
22    } // return ans_last_place
23    void build_fail(int ptr) {
24        int tmp;
25        for (int i = 0; i < maxc; i++)
26            if (T[ptr].Next[i]) {
27                tmp = T[ptr].fail;
28                while (tmp != 1 && !T[tmp].Next[i])
29                    tmp = T[tmp].fail;
30                if (T[tmp].Next[i] != T[ptr].Next[i])
31                    if (T[tmp].Next[i]) tmp = T[tmp].Next[i];
32                T[T[ptr].Next[i]].fail = tmp;
33                q[qtop++] = T[ptr].Next[i];
34            }
35    }
36    void AC_auto(const string &s) {
37        int ptr = 1;
38        for (char c : s) {
39            while (ptr != 1 && !T[ptr].Next[c]) ptr = T[ptr].fail;
40            if (T[ptr].Next[c]) {
41                ptr = T[ptr].Next[c];
42                T[ptr].ans++;
43            }
44        }
45    }
46    void Solve(string &s) {
47        for (char &c : s) // change char id
48            c -= 'a';
49        for (int i = 0; i < qtop; i++) build_fail(q[i]);
50        AC_auto(s);
51        for (int i = qtop - 1; i > -1; i--)
52            T[T[q[i]].fail].ans += T[q[i]].ans;
53    }
54    void reset() {
55        qtop = top = q[0] = 1;
56        get_node(1);
57    }
58 } AC;
59 // usage example
60 string s, S;
61 int n, t, ans_place[50000];
62 int main() {
63     Tie cin >> t;
64     while (t--) {
65         AC.reset();
66         cin >> S >> n;
67         for (int i = 0; i < n; i++) {
68             cin >> s;
69             ans_place[i] = AC.insert(s);
70         }
71         AC.Solve(S);
72         for (int i = 0; i < n; i++)
73             cout << AC.T[ans_place[i]].ans << '\n';
74     }
75 }

```

### 7.3. Suffix Array

```

1 // sa[i]: starting index of suffix at rank i
2 // 0-indexed, sa[0] = n (empty string)
3 // lcp[i]: lcp of sa[i] and sa[i - 1], lcp[0] = 0
4 struct SuffixArray {
5     vector<int> sa, lcp;
6     SuffixArray(string &s,
7         int lim = 256) { // or basic_string<int>
8         int n = sz(s) + 1, k = 0, a, b;
9         vector<int> x(all(s) + 1), y(n), ws(max(n, lim)),
10             rank(n);
11         sa = lcp = y, iota(all(sa), 0);

```



```

13   for (int j = 0, p = 0; p < n;
      j = max(1, j * 2), lim = p) {
15       p = j, iota(all(y), n - j);
      for (int i = 0; i < n; i++)
17         if (sa[i] >= j) y[p++] = sa[i] - j;
      fill(all(ws), 0);
      for (int i = 0; i < n; i++) ws[x[i]]++;
19       for (int i = 1; i < lim; i++) ws[i] += ws[i - 1];
      for (int i = n; i--;) sa[--ws[x[y[i]]]] = y[i];
21       swap(x, y), p = 1, x[sa[0]] = 0;
      for (int i = 1; i < n; i++)
23         a = sa[i - 1], b = sa[i],

        x[b] = (y[a] == y[b] && y[a + j] == y[b + j])
            ? p - 1 : p++;

25   }
27   for (int i = 1; i < n; i++) rank[sa[i]] = i;
   for (int i = 0, j; i < n - 1; lcp[rank[i+1]] = k)
31     for (k && k--, j = sa[rank[i] - 1];
          s[i + k] == s[j + k]; k++)
33       ;
35 };

```

#### 7.4. Suffix Tree

```

1 struct SAM {
   static const int maxc = 26; // char range
   static const int maxn = 10010; // string len
   struct Node {
       Node *green, *edge[maxc];
       int max_len, in, times;
   } * root, *last, reg[maxn * 2];
   int top;
   Node *get_node(int _max) {
       Node *re = &reg[top++];
       re->in = 0, re->times = 1;
       re->max_len = _max, re->green = 0;
       for (int i = 0; i < maxc; i++) re->edge[i] = 0;
       return re;
   }
   void insert(const char c) { // c in range [0, maxc)
       Node *p = last;
       last = get_node(p->max_len + 1);
       while (p && !p->edge[c])
           p->edge[c] = last, p = p->green;
       if (!p) last->green = root;
       else {
           Node *pot_green = p->edge[c];
           if ((pot_green->max_len) == (p->max_len + 1))
               last->green = pot_green;
           else {
               Node *wish = get_node(p->max_len + 1);
               wish->times = 0;
               while (p && p->edge[c] == pot_green)
                   p->edge[c] = wish, p = p->green;
               for (int i = 0; i < maxc; i++)
                   wish->edge[i] = pot_green->edge[i];
               wish->green = pot_green->green;
               pot_green->green = wish;
               last->green = wish;
           }
       }
   }
   Node *q[maxn * 2];
   int ql, qr;
   void get_times(Node *p) {
       ql = 0, qr = -1, reg[0].in = 1;
       for (int i = 1; i < top; i++) reg[i].green->in++;
       for (int i = 0; i < top; i++)
           if (!reg[i].in) q[++qr] = &reg[i];
       while (ql <= qr) {
           q[ql]->green->times += q[ql]->times;
           if (!(--q[ql]->green->in)) q[++qr] = q[ql]->green;
           ql++;
       }
   }
   void build(const string &s) {
       top = 0;
       root = last = get_node(0);
       for (char c : s) insert(c - 'a'); // change char id
       get_times(root);
   }
   // call build before solve
   int solve(const string &s) {
       Node *p = root;
       for (char c : s)
           if (!(p = p->edge[c - 'a'])) // change char id
               return 0;
   }

```

```

65   return p->times;
   };

```

#### 7.5. Cocke-Younger-Kasami Algorithm

```

1 struct rule {
   // s -> xy
   // if y == -1, then s -> x (unit rule)
   int s, x, y, cost;
2   };
   int state;
   // state (id) for each letter (variable)
   // lowercase letters are terminal symbols
   map<char, int> rules;
   vector<rule> cnf;
   void init() {
       state = 0;
       rules.clear();
       cnf.clear();
   }
   // convert a cfg rule to cnf (but with unit rules) and add
   // it
   void add_to_cnf(char s, const string &p, int cost) {
       if (!rules.count(s)) rules[s] = state++;
       for (char c : p)
           if (!rules.count(c)) rules[c] = state++;
       if (p.size() == 1) {
           cnf.push_back({rules[s], rules[p[0]], -1, cost});
       } else {
           // length >= 3 -> split
           int left = rules[s];
           int sz = p.size();
           for (int i = 0; i < sz - 2; i++) {
               cnf.push_back({left, rules[p[i]], state, 0});
               left = state++;
           }
           cnf.push_back(
               {left, rules[p[sz - 2]], rules[p[sz - 1]], cost});
       }
   }

   constexpr int MAXN = 55;
   vector<long long> dp[MAXN][MAXN];
   // unit rules with negative costs can cause negative cycles
   vector<bool> neg_INF[MAXN][MAXN];

   void relax(int l, int r, rule c, long long cost,
              bool neg_c = 0) {
       if (!neg_INF[l][r][c.s] &&
           (neg_INF[l][r][c.x] || cost < dp[l][r][c.s])) {
           if (neg_c || neg_INF[l][r][c.x]) {
               dp[l][r][c.s] = 0;
               neg_INF[l][r][c.s] = true;
           } else {
               dp[l][r][c.s] = cost;
           }
       }
   }

   void bellman(int l, int r, int n) {
       for (int k = 1; k <= state; k++)
           for (rule c : cnf)
               if (c.y == -1)
                   relax(l, r, c, dp[l][r][c.x] + c.cost, k == n);
   }

   void cyk(const string &s) {
       vector<int> tok;
       for (char c : s) tok.push_back(rules[c]);
       for (int i = 0; i < tok.size(); i++) {
           for (int j = 0; j < tok.size(); j++) {
               dp[i][j] = vector<long long>(state + 1, INT_MAX);
               neg_INF[i][j] = vector<bool>(state + 1, false);
           }
           dp[i][i][tok[i]] = 0;
           bellman(i, i, tok.size());
       }
       for (int r = 1; r < tok.size(); r++) {
           for (int l = r - 1; l >= 0; l--) {
               for (int k = l; k < r; k++)
                   for (rule c : cnf)
                       if (c.y != -1)
                           relax(l, r, c,
                               dp[l][k][c.x] + dp[k + 1][r][c.y] +
                               c.cost);
               bellman(l, r, tok.size());
           }
       }
   }
   // usage example

```

```

85 int main() {
    init();
    add_to_cnf('S', "aSc", 1);
    add_to_cnf('S', "BBB", 1);
    add_to_cnf('S', "SB", 1);
    add_to_cnf('B', "b", 1);
    cyk("abbbbc");
    // dp[0][s.size() - 1][rules[start]] = min cost to
    // generate s
    cout << dp[0][5][rules['S']] << '\n'; // 7
    cyk("acbc");
    cout << dp[0][3][rules['S']] << '\n'; // INT_MAX
    add_to_cnf('S', "S", -1);
    cyk("abbbbc");
    cout << neg_INF[0][5][rules['S']] << '\n'; // 1
}

```

## 7.6. Z Value

```

1 int z[n];
void zval(string s) {
    // z[i] => longest common prefix of s and s[i:], i > 0
    int n = s.size();
    z[0] = 0;
    for (int b = 0, i = 1; i < n; i++) {
        if (z[b] + b <= i) z[i] = 0;
        else z[i] = min(z[i - b], z[b] + b - i);
        while (s[i + z[i]] == s[z[i]]) z[i]++;
        if (i + z[i] > b + z[b]) b = i;
    }
}

```

## 7.7. Manacher's Algorithm

```

1 int z[n];
void manacher(string s) {
    // z[i] => longest odd palindrome centered at i is
    // s[i - z[i] ... i + z[i]]
    // to get all palindromes (including even length),
    // insert a '#' between each s[i] and s[i + 1]
    int n = s.size();
    z[0] = 0;
    for (int b = 0, i = 1; i < n; i++) {
        if (z[b] + b >= i)
            z[i] = min(z[2 * b - i], b + z[b] - i);
        else z[i] = 0;
        while (i + z[i] + 1 < n && i - z[i] - 1 >= 0 &&
            s[i + z[i] + 1] == s[i - z[i] - 1])
            z[i]++;
        if (z[i] + i > z[b] + b) b = i;
    }
}

```

## 7.8. Minimum Rotation

```

1 int min_rotation(string s) {
    int a = 0, n = s.size();
    s += s;
    for (int b = 0; b < n; b++) {
        for (int k = 0; k < n; k++) {
            if (a + k == b || s[a + k] < s[b + k]) {
                b += max(0, k - 1);
                break;
            }
            if (s[a + k] > s[b + k]) {
                a = b;
                break;
            }
        }
    }
    return a;
}

```

## 7.9. Palindromic Tree

```

1 struct palindromic_tree {
    struct node {
        int next[26], fail, len;
        int cnt,
        num; // cnt: appear times, num: number of pal. suf.
        node(int l = 0) : fail(0), len(l), cnt(0), num(0) {
            for (int i = 0; i < 26; ++i) next[i] = 0;
        }
    };
    vector<node> St;
    vector<char> s;
    int last, n;
    palindromic_tree() : St(2), last(1), n(0) {
        St[0].fail = 1, St[1].len = -1, s.pb(-1);
    }
}

```

```

inline void clear() {
    St.clear(), s.clear(), last = 1, n = 0;
    St.pb(0), St.pb(-1);
    St[0].fail = 1, s.pb(-1);
}

inline int get_fail(int x) {
    while (s[n - St[x].len - 1] != s[n]) x = St[x].fail;
    return x;
}

inline void add(int c) {
    s.push_back(c -= 'a'), ++n;
    int cur = get_fail(last);
    if (!St[cur].next[c]) {
        int now = SZ(St);
        St.pb(St[cur].len + 2);
        St[now].fail = St[get_fail(St[cur].fail)].next[c];
        St[cur].next[c] = now;
        St[now].num = St[St[now].fail].num + 1;
    }
    last = St[cur].next[c], ++St[last].cnt;
}

inline void count() { // counting cnt
    auto i = St.rbegin();
    for (; i != St.rend(); ++i) {
        St[i->fail].cnt += i->cnt;
    }
}

inline int size() { // The number of diff. pal.
    return SZ(St) - 2;
}
};

```

## 8. Debug List

- 1 - Pre-submit:
  - Did you make a typo when copying a template?
  - Test more cases if unsure.
  - Write a naive solution and check small cases.
  - Submit the correct file.
- 7 - General Debugging:
  - Read the whole problem again.
  - Have a teammate read the problem.
  - Have a teammate read your code.
  - Explain your solution to them (or a rubber duck).
  - Print the code and its output / debug output.
  - Go to the toilet.
- 15 - Wrong Answer:
  - Any possible overflows?
    - > ``_int128``?
    - Try ``-ftrapv`` or ``#pragma GCC optimize("trapv")``
  - Floating point errors?
    - > ``long double``?
    - turn off math optimizations
    - check for ``==``, ``>=``, ``acos(1.000000001)``, etc.
  - Did you forget to sort or unique?
  - Generate large and worst "corner" cases.
  - Check your ``m`` / ``n``, ``i`` / ``j`` and ``x`` / ``y``.
  - Are everything initialized or reset properly?
  - Are you sure about the STL thing you are using?
    - Read cppreference (should be available).
  - Print everything and run it on pen and paper.
- 31 - Time Limit Exceeded:
  - Calculate your time complexity again.
  - Does the program actually end?
    - Check for ``while(q.size())`` etc.
  - Test the largest cases locally.
  - Did you do unnecessary stuff?
    - e.g. pass vectors by value
    - e.g. ``memset`` for every test case
  - Is your constant factor reasonable?
- 41 - Runtime Error:
  - Check memory usage.
  - Forget to clear or destroy stuff?
    - > ``vector::shrink_to_fit()``
  - Stack overflow?
  - Bad pointer / array access?
    - Try ``-fsanitize=address``
  - Division by zero? NaN's?