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```

15 - Wrong Answer:
    - Any possible overflows?
    - > `__int128`?
    - Try `-ftrapv` or `#pragma GCC optimize("trapv")`
19 - Floating point errors?
    - > `long double`?
    - turn off math optimizations
    - check for `==`, `>=`, `acos(1.000000001)`, etc.
23 - Did you forget to sort or unique?
    - Generate large and worst "corner" cases.
25 - Check your `m` / `n`, `i` / `j` and `x` / `y`.
    - Are everything initialized or reset properly?
27 - Are you sure about the STL thing you are using?
    - Read cppreference (should be available).
29 - Print everything and run it on pen and paper.

31 - Time Limit Exceeded:
    - Calculate your time complexity again.
33 - Does the program actually end?
    - Check for `while(q.size())` etc.
35 - Test the largest cases locally.
    - Did you do unnecessary stuff?
    - e.g. pass vectors by value
    - e.g. `memset` for every test case
39 - Is your constant factor reasonable?

41 - Runtime Error:
    - Check memory usage.
    - Forget to clear or destroy stuff?
    - > `vector::shrink_to_fit()`
43 - Stack overflow?
45 - Bad pointer / array access?
    - Try `-fsanitize=address`
47 - Division by zero? NaN's?

```

1.2. How Did We Get Here?

1.2.1. Macros

Use vectorizations and math optimizations at your own peril.
 For gcc \geq 9, there are `[[likely]]` and `[[unlikely]]` attributes.
 Call gcc with `-fopt-info-optimized-missed-optall` for optimization info.

```

1 #define _GLIBCXX_DEBUG 1 // for debug mode
2 #define _GLIBCXX_SANITIZE_VECTOR 1 // for asan on vectors
3 #pragma GCC optimize("O3", "unroll-loops")
4 #pragma GCC optimize("fast-math")
5 #pragma GCC target("avx,avx2,abm,bmi,bmi2") // tip: `lscpu`
6 // before a loop
7 #pragma GCC unroll 16 // 0 or 1 -> no unrolling
8 #pragma GCC ivdep

```

1.2.2. Fast I/O

```

1 struct scanner {
2     static constexpr size_t LEN = 32 << 20;
3     char *buf, *buf_ptr, *buf_end;
4     scanner() {
5         buf(new char[LEN]), buf_ptr(buf + LEN),
6         buf_end(buf + LEN) {}
7     }
8     ~scanner() { delete[] buf; }
9     char getc() {
10         if (buf_ptr == buf_end) [[unlikely]]
11             buf_end = buf + fread_unlocked(buf, 1, LEN, stdin),
12             buf_ptr = buf;
13         return *(buf_ptr++);
14     }
15     char seek(char del) {
16         char c;
17         while ((c = getc()) < del) {}
18         return c;
19     }
20     void read(int &t) {
21         bool neg = false;
22         char c = seek('-');
23         if (c == '-') neg = true, t = 0;
24         else t = c ^ '0';
25         while ((c = getc()) >= '0') t = t * 10 + (c ^ '0');
26         if (neg) t = -t;
27     }
28 };
29 struct printer {
30     static constexpr size_t CPI = 21, LEN = 32 << 20;
31     char *buf, *buf_ptr, *buf_end, *tbuf;
32     char *int_buf, *int_buf_end;
33     printer() {
34         buf(new char[LEN]), buf_ptr(buf),
35         buf_end(buf + LEN), int_buf(new char[CPI + 1]()),
36         int_buf_end(int_buf + CPI - 1) {}
37     }
38     ~printer() {
39         flush();
40         delete[] buf, delete[] int_buf;
41     }
42     void flush() {
43         fwrite_unlocked(buf, 1, buf_ptr - buf, stdout);
44         buf_ptr = buf;
45     }
46     void write(const char &c) {
47         *buf_ptr = c;
48         if (++buf_ptr == buf_end) [[unlikely]]
49             flush();
50     }
51     void write(const char *s) {
52         for (; *s != '\0'; ++s) write(*s);
53     }
54     void write(int x) {
55         if (x < 0) write('-'), x = -x;
56         if (x == 0) [[unlikely]]
57             return write('0');
58         for (tbuf = int_buf_end; x != 0; --tbuf, x /= 10)
59             *tbuf = '0' + char(x % 10);
60         write(++tbuf);
61     }
62 };

```

```

37     flush();
38     delete[] buf, delete[] int_buf;
39 }
40 void flush() {
41     fwrite_unlocked(buf, 1, buf_ptr - buf, stdout);
42     buf_ptr = buf;
43 }
44 void write(const char &c) {
45     *buf_ptr = c;
46     if (++buf_ptr == buf_end) [[unlikely]]
47         flush();
48 }
49 void write(const char *s) {
50     for (; *s != '\0'; ++s) write(*s);
51 }
52 void write(int x) {
53     if (x < 0) write('-'), x = -x;
54     if (x == 0) [[unlikely]]
55         return write('0');
56     for (tbuf = int_buf_end; x != 0; --tbuf, x /= 10)
57         *tbuf = '0' + char(x % 10);
58     write(++tbuf);
59 }
60 };

```

Kotlin

```

1 import java.io.*
2 import java.util.*
3
4 @JvmField val cin = System.`in`.bufferedReader()
5 @JvmField val cout = PrintWriter(System.out, false)
6 @JvmField var tokenizer: StringTokenizer = StringTokenizer("")
7 fun nextLine() = cin.readLine()!!
8 fun read(): String {
9     while(!tokenizer.hasMoreTokens())
10         tokenizer = StringTokenizer(nextLine())
11     return tokenizer.nextToken()
12 }
13
14 // example
15 fun main() {
16     val n = read().toInt()
17     val a = DoubleArray(n) { read().toDouble() }
18     cout.println("omg hi")
19     cout.flush()
20 }

```

1.2.3. constexpr

Some default limits in gcc (7.x - trunk):

- constexpr recursion depth: 512
- constexpr loop iteration per function: 262144
- constexpr operation count per function: 33554432
- template recursion depth: 900 (gcc *might* segfault first)

```

1 constexpr array<int, 10> fibonacci{[] {
2     array<int, 10> a{};
3     a[0] = a[1] = 1;
4     for (int i = 2; i < 10; i++) a[i] = a[i - 1] + a[i - 2];
5     return a;
6 }}();
7 static_assert(fibonacci[9] == 55, "CE");
8
9 template <typename F, typename INT, INT... S>
10 constexpr void for_constexpr(integer_sequence<INT, S...>,
11                             F &&func) {
12     int _[] = {(func(integral_constant<INT, S>{}), 0)...};
13 }
14 // example
15 template <typename... T> void print_tuple(tuple<T...> t) {
16     for_constexpr(make_index_sequence<sizeof...(T)>{}),
17     [&](auto i) { cout << get<i>(t) << '\n'; };
18 }

```

1.2.4. Bump Allocator

```

1 // global bump allocator
2 char mem[256 << 20]; // 256 MB
3 size_t rsp = sizeof mem;
4 void *operator new(size_t s) {
5     assert(s < rsp); // MLE
6     return (void *)&mem[rsp -= s];
7 }
8 void operator delete(void *) {}
9
10 // bump allocator for STL / pbds containers

```

```

11 char mem[256 << 20];
12 size_t rsp = sizeof mem;
13 template <typename T> struct bump {
14     typedef T value_type;
15     bump() {}
16     template <typename U> bump(U, ...) {}
17     T *allocate(size_t n) {
18         rsp -= n * sizeof(T);
19         rsp &= 0 - alignof(T);
20         return (T *) (mem + rsp);
21     }
22     void deallocate(T *, size_t n) {}
23 };

```

1.3. Tools

1.3.1. Floating Point Binary Search

```

1 union di {
2     double d;
3     ull i;
4 };
5 bool check(double);
6 // binary search in [L, R) with relative error 2^-eps
7 double binary_search(double L, double R, int eps) {
8     di l = {L}, r = {R}, m;
9     while (r.i - l.i > 1LL << (52 - eps)) {
10         m.i = (l.i + r.i) >> 1;
11         if (check(m.d)) r = m;
12         else l = m;
13     }
14     return l.d;
15 }

```

1.3.2. SplitMix64

```

1 using ull = unsigned long long;
2 inline ull splitmix64(ull x) {
3     // change to `static ull x = SEED;` for DRBG
4     ull z = (x += 0x9E3779B97F4A7C15);
5     z = (z ^ (z >> 30)) * 0xBF58476D1CE4E5B9;
6     z = (z ^ (z >> 27)) * 0x94D049BB133111EB;
7     return z ^ (z >> 31);
8 }

```

1.3.3. <random>

```

1 #ifdef __unix__
2     random_device rd;
3     mt19937_64 RNG(rd());
4 #else
5     const auto SEED = chrono::high_resolution_clock::now()
6         .time_since_epoch()
7         .count();
8     mt19937_64 RNG(SEED);
9 #endif
10 // random uint_fast64_t: RNG();
11 // uniform random of type T (int, double, ...) in [l, r]:
12 // uniform_int_distribution<T> dist(l, r); dist(RNG);

```

1.3.4. x86 Stack Hack

```

1 constexpr size_t size = 200 << 20; // 200MiB
2 int main() {
3     register long rsp asm("rsp");
4     char *buf = new char[size];
5     asm("movq %0, %%rsp\n" :: "r" (buf + size));
6     // do stuff
7     asm("movq %0, %%rsp\n" :: "r" (rsp));
8     delete[] buf;
9 }

```

1.4. Algorithms

1.4.1. Bit Hacks

```

1 // next permutation of x as a bit sequence
2 ull next_bits_permutation(ull x) {
3     ull c = __builtin_ctzll(x), r = x + (1 << c);
4     return (r ^ x) >> (c + 2) | r;
5 }
6 // iterate over all (proper) subsets of bitset s
7 void subsets(ull s) {
8     for (ull x = s; x;) { --x &= s; /* do stuff */ }
9 }

```

1.4.2. Aliens Trick

```

1 // min dp[i] value and its i (smallest one)
2 pll get_dp(int cost);
3 ll aliens(int k, int l, int r) {
4     while (l != r) {
5         int m = (l + r) / 2;
6         auto [f, s] = get_dp(m);
7         if (s == k) return f - m * k;
8         if (s < k) r = m;
9         else l = m + 1;
10     }
11     return get_dp(l).first - l * k;
12 }

```

1.4.3. Hilbert Curve

```

1 ll hilbert(ll n, int x, int y) {
2     ll res = 0;
3     for (ll s = n; s /= 2;) {
4         int rx = !(x & s), ry = !(y & s);
5         res += s * s * ((3 * rx) ^ ry);
6         if (ry == 0) {
7             if (rx == 1) x = s - 1 - x, y = s - 1 - y;
8             swap(x, y);
9         }
10     }
11     return res;
12 }

```

1.4.4. Infinite Grid Knight Distance

```

1 ll get_dist(ll dx, ll dy) {
2     if (++(dx = abs(dx)) > ++(dy = abs(dy))) swap(dx, dy);
3     if (dx == 1 && dy == 2) return 3;
4     if (dx == 3 && dy == 3) return 4;
5     ll lb = max(dy / 2, (dx + dy) / 3);
6     return ((dx ^ dy ^ lb) & 1) ? ++lb : lb;
7 }

```

1.4.5. Poker Hand

```

1 using namespace std;
2
3 struct hand {
4     static constexpr auto rk = [] {
5         array<int, 256> x{};
6         auto s = "23456789TJQKACDHS";
7         for (int i = 0; i < 17; i++) x[s[i]] = i % 13;
8         return x;
9     }();
10     vector<pair<int, int>> v;
11     vector<int> cnt, vf, vs;
12     int type;
13     hand() : cnt(4), type(0) {}
14     void add_card(char suit, char rank) {
15         ++cnt[rk[suit]];
16         for (auto &[f, s] : v)
17             if (s == rk[rank]) return ++f, void();
18         v.emplace_back(1, rk[rank]);
19     }
20     void process() {
21         sort(v.rbegin(), v.rend());
22         for (auto [f, s] : v) vf.push_back(f), vs.push_back(s);
23         bool str = 0, flu = find(all(cnt), 5) != cnt.end();
24         if ((str = v.size() == 5))
25             for (int i = 1; i < 5; i++)
26                 if (vs[i] != vs[i - 1] + 1) str = 0;
27         if (vs == vector<int>{12, 3, 2, 1, 0})
28             str = 1, vs = {3, 2, 1, 0, -1};
29         if (str && flu) type = 9;
30         else if (vf[0] == 4) type = 8;
31         else if (vf[0] == 3 && vf[1] == 2) type = 7;
32         else if (str || flu) type = 5 + flu;
33         else if (vf[0] == 3) type = 4;
34         else if (vf[0] == 2) type = 2 + (vf[1] == 2);
35         else type = 1;
36     }
37     bool operator<(const hand &b) const {
38         return make_tuple(type, vf, vs) <
39             make_tuple(b.type, b.vf, b.vs);
40     }
41 };

```

1.4.6. Longest Increasing Subsequence

```

1 template <class I> vi lis(const vector<I> &S) {
2     if (S.empty()) return {};
3     vi prev(sz(S));
4     typedef pair<I, int> p;

```

```

5 vector<p> res;
6 rep(i, 0, sz(S)) {
7     // change 0 -> i for longest non-decreasing subsequence
8     auto it = lower_bound(all(res), p[S[i], 0]);
9     if (it == res.end())
10         res.emplace_back(), it = res.end() - 1;
11     *it = {S[i], i};
12     prev[i] = it == res.begin() ? 0 : (it - 1)->second;
13 }
14 int L = sz(res), cur = res.back().second;
15 vi ans(L);
16 while (L--) ans[L] = cur, cur = prev[cur];
17 return ans;
18 }

```

1.4.7. Mo's Algorithm on Tree

```

1 void MoAlgoOnTree() {
2     Dfs(0, -1);
3     vector<int> euler(tk);
4     for (int i = 0; i < n; ++i) {
5         euler[tin[i]] = i;
6         euler[tout[i]] = i;
7     }
8     vector<int> l(q), r(q), qr(q), sp(q, -1);
9     for (int i = 0; i < q; ++i) {
10         if (tin[u[i]] > tin[v[i]]) swap(u[i], v[i]);
11         int z = GetLCA(u[i], v[i]);
12         sp[i] = z[i];
13         if (z == u) l[i] = tin[u[i]], r[i] = tin[v[i]];
14         else l[i] = tout[u[i]], r[i] = tin[v[i]];
15         qr[i] = i;
16     }
17     sort(qr.begin(), qr.end(), [&](int i, int j) {
18         if (l[i] / KB == l[j] / KB) return r[i] < r[j];
19         return l[i] / KB < l[j] / KB;
20     });
21     vector<bool> used(n);
22     // Add(v): add/remove v to/from the path based on used[v]
23     for (int i = 0, tl = 0, tr = -1; i < q; ++i) {
24         while (tl < l[qr[i]]) Add(euler[tl++]);
25         while (tl > l[qr[i]]) Add(euler[--tl]);
26         while (tr > r[qr[i]]) Add(euler[tr--]);
27         while (tr < r[qr[i]]) Add(euler[++tr]);
28         // add/remove LCA(u, v) if necessary
29     }
30 }

```

```

13 void update(int i, T x) {
14     for (v[i += n] = x; i /= 2;)
15         v[i] = f(v[i * 2], v[i * 2 + 1]);
16 }
17 T query(int l, int r) {
18     T tl = ID, tr = ID;
19     for (l += n, r += n; l < r; l /= 2, r /= 2) {
20         if (l & 1) tl = f(tl, v[l++]);
21         if (r & 1) tr = f(v[--r], tr);
22     }
23     return f(tl, tr);
24 }
25 };

```

2.3. Line Container

```

1 struct Line {
2     mutable ll k, m, p;
3     bool operator<(const Line &o) const { return k < o.k; }
4     bool operator<(ll x) const { return p < x; }
5 };
6 // add: line y=kx+m, query: maximum y of given x
7 struct LineContainer : multiset<Line, less<>> {
8     // (for doubles, use inf = 1/.0, div(a,b) = a/b)
9     static const ll inf = LLONG_MAX;
10     ll div(ll a, ll b) { // floored division
11         return a / b - ((a ^ b) < 0 && a % b);
12     }
13     bool isect(iterator x, iterator y) {
14         if (y == end()) return x->p = inf, 0;
15         if (x->k == y->k) x->p = x->m > y->m ? inf : -inf;
16         else x->p = div(y->m - x->m, x->k - y->k);
17         return x->p >= y->p;
18     }
19     void add(ll k, ll m) {
20         auto z = insert({k, m, 0}), y = z++, x = y;
21         while (isect(y, z)) z = erase(z);
22         if (x != begin() && isect(--x, y))
23             isect(x, y = erase(y));
24         while ((y = x) != begin() && (--x)->p >= y->p)
25             isect(x, erase(y));
26     }
27     ll query(ll x) {
28         assert(!empty());
29         auto l = *lower_bound(x);
30         return l.k * x + l.m;
31     }
32 };

```

2. Data Structures

2.1. GNU PBDS

```

1 #include <ext/pb_ds/assoc_container.hpp>
2 #include <ext/pb_ds/priority_queue.hpp>
3 #include <ext/pb_ds/tree_policy.hpp>
4 using namespace __gnu_pbds;
5
6 // most std::map + order_of_key, find_by_order, split, join
7 template <typename T, typename U = null_type>
8 using ordered_map = tree<T, U, std::less<>, rb_tree_tag,
9     tree_order_statistics_node_update>;
10 // useful tags: rb_tree_tag, splay_tree_tag
11
12 template <typename T> struct myhash {
13     size_t operator()(T x) const; // splitmix, bswap(x*R), ...
14 };
15 // most of std::unordered_map, but faster (needs good hash)
16 template <typename T, typename U = null_type>
17 using hash_table = gp_hash_table<T, U, myhash<T>>;
18
19 // most std::priority_queue + modify, erase, split, join
20 using heap = priority_queue<int, std::less<>>;
21 // useful tags: pairing_heap_tag, binary_heap_tag,
22 // (rc_)?binomial_heap_tag, thin_heap_tag

```

2.2. Segment Tree (ZKW)

```

1 struct segtree {
2     using T = int;
3     T f(T a, T b) { return a + b; } // any monoid operation
4     static constexpr T ID = 0; // identity element
5     int n;
6     vector<T> v;
7     segtree(int n_) : n(n_), v(2 * n, ID) {}
8     segtree(vector<T> &a) : n(a.size()), v(2 * n, ID) {
9         copy_n(a.begin(), n, v.begin() + n);
10         for (int i = n - 1; i > 0; i--)
11             v[i] = f(v[i * 2], v[i * 2 + 1]);
12     }
13 }

```

2.4. Heavy-Light Decomposition

```

1 template <bool VALS_EDGES> struct HLD {
2     int N, tim = 0;
3     vector<vi> adj;
4     vi par, siz, depth, rt, pos;
5     Node *tree;
6     HLD(vector<vi> adj_)
7         : N(sz(adj_)), adj(adj_), par(N, -1), siz(N, 1),
8           depth(N), rt(N), pos(N), tree(new Node(0, N)) {}
9     dfsSz(0);
10     dfsHld(0);
11 }
12 void dfsSz(int v) {
13     if (par[v] != -1)
14         adj[v].erase(find(all(adj[v]), par[v]));
15     for (int &u : adj[v]) {
16         par[u] = v, depth[u] = depth[v] + 1;
17         dfsSz(u);
18         siz[v] += siz[u];
19         if (siz[u] > siz[adj[v][0]]) swap(u, adj[v][0]);
20     }
21 }
22 void dfsHld(int v) {
23     pos[v] = tim++;
24     for (int u : adj[v]) {
25         rt[u] = (u == adj[v][0] ? rt[v] : u);
26         dfsHld(u);
27     }
28 }
29 template <class B> void process(int u, int v, B op) {
30     for (; rt[u] != rt[v]; v = par[rt[v]]) {
31         if (depth[rt[u]] > depth[rt[v]]) swap(u, v);
32         op(pos[rt[v]], pos[v] + 1);
33     }
34     if (depth[u] > depth[v]) swap(u, v);
35     op(pos[u] + VALS_EDGES, pos[v] + 1);
36 }
37 void modifyPath(int u, int v, int val) {
38     process(u, v,
39         [&](int l, int r) { tree->add(l, r, val); });
40 }

```

```

}
41 int queryPath(int u,
               int v) { // Modify depending on problem
43     int res = -1e9;
    process(u, v, [&](int l, int r) {
45         res = max(res, tree->query(l, r));
    });
47     return res;
}
49 int querySubtree(int v) { // modifySubtree is similar
    return tree->query(pos[v] + VALS_EDGES,
51                     pos[v] + siz[v]);
}
53 };

```

2.5. Wavelet Matrix

```

1 #pragma GCC target("popcnt,bmi2")
#include <immintrin.h>
3 // T is unsigned. You might want to compress values first
5 template <typename T> struct wavelet_matrix {
    static_assert(is_unsigned_v<T>, "only unsigned T");
7     struct bit_vector {
        static constexpr uint W = 64;
        uint n, cnt0;
        vector<ull> bits;
        vector<uint> sum;
        bit_vector(uint n_)
            : n(n_), bits(n / W + 1), sum(n / W + 1) {}
13         void build() {
            for (uint j = 0; j != n / W; ++j)
15                 sum[j + 1] = sum[j] + _mm_popcnt_u64(bits[j]);
            cnt0 = rank0(n);
17         }
        void set_bit(uint i) { bits[i / W] |= 1ULL << i % W; }
        bool operator[](uint i) const {
21             return !(bits[i / W] & 1ULL << i % W);
        }
        uint rank1(uint i) const {
23             return sum[i / W] +
                _mm_popcnt_u64(_bzh_u64(bits[i / W], i % W));
25         }
        uint rank0(uint i) const { return i - rank1(i); }
    };
    uint n, lg;
    vector<bit_vector> b;
    wavelet_matrix(const vector<T> &a) : n(a.size()) {
33         lg =
            __lg(max(*max_element(a.begin(), a.end()), T(1))) + 1;
        b.assign(lg, n);
        vector<T> cur = a, nxt(n);
        for (int h = lg; h--;) {
37             for (uint i = 0; i < n; ++i)
                if (cur[i] & (T(1) << h)) b[h].set_bit(i);
            b[h].build();
            int il = 0, ir = b[h].cnt0;
            for (uint i = 0; i < n; ++i)
41                 nxt[(b[h][i] ? ir : il)++] = cur[i];
            swap(cur, nxt);
43         }
    }
    T operator[](uint i) const {
47         T res = 0;
        for (int h = lg; h--;)
49             if (b[h][i])
                i += b[h].cnt0 - b[h].rank0(i), res |= T(1) << h;
51         else i = b[h].rank0(i);
        return res;
53     }
    // query k-th smallest (0-based) in a[l, r]
55     T kth(uint l, uint r, uint k) const {
        T res = 0;
57         for (int h = lg; h--;) {
            uint tl = b[h].rank0(l), tr = b[h].rank0(r);
59             if (k >= tr - tl) {
                k -= tr - tl;
                l += b[h].cnt0 - tl;
                r += b[h].cnt0 - tr;
                res |= T(1) << h;
                } else l = tl, r = tr;
65         }
        return res;
67     }
    // count of i in [l, r] with a[i] < u
69     uint count(uint l, uint r, T u) const {
        if (u >= T(1) << lg) return r - l;
        uint res = 0;
71         for (int h = lg; h--;) {
            uint tl = b[h].rank0(l), tr = b[h].rank0(r);
73

```

```

        if (u & (T(1) << h)) {
            l += b[h].cnt0 - tl;
            r += b[h].cnt0 - tr;
            res += tr - tl;
        } else l = tl, r = tr;
79     }
    return res;
81 }
};

```

2.6. Link-Cut Tree

```

1 const int MXN = 100005;
const int MEM = 100005;
3
struct Splay {
5     static Splay nil, mem[MEM], *pmem;
    Splay *ch[2], *f;
    int val, rev, size;
    Splay() : val(-1), rev(0), size(0) {
        f = ch[0] = ch[1] = &nil;
    }
11     Splay(int _val) : val(_val), rev(0), size(1) {
        f = ch[0] = ch[1] = &nil;
    }
13     bool isr() {
15         return f->ch[0] != this && f->ch[1] != this;
    }
    int dir() { return f->ch[0] == this ? 0 : 1; }
    void setCh(Splay *c, int d) {
19         ch[d] = c;
        if (c != &nil) c->f = this;
        pull();
21     }
    void push() {
23         if (rev) {
            swap(ch[0], ch[1]);
            if (ch[0] != &nil) ch[0]->rev ^= 1;
            if (ch[1] != &nil) ch[1]->rev ^= 1;
            rev = 0;
25         }
    }
    void pull() {
31         size = ch[0]->size + ch[1]->size + 1;
        if (ch[0] != &nil) ch[0]->f = this;
        if (ch[1] != &nil) ch[1]->f = this;
33     }
} Splay::nil, Splay::mem[MEM], *Splay::pmem = Splay::mem;
Splay *nil = &Splay::nil;
37
void rotate(Splay *x) {
    Splay *p = x->f;
    int d = x->dir();
    if (!p->isr()) p->f->setCh(x, p->dir());
    else x->f = p->f;
    p->setCh(x->ch[!d], d);
    x->setCh(p, !d);
    p->pull();
    x->pull();
47 }
49
vector<Splay *> splayVec;
51 void splay(Splay *x) {
    splayVec.clear();
53     for (Splay *q = x; q = q->f) {
        splayVec.push_back(q);
        if (q->isr()) break;
55     }
    reverse(begin(splayVec), end(splayVec));
    for (auto it : splayVec) it->push();
57     while (!x->isr()) {
        if (x->f->isr()) rotate(x);
        else if (x->dir() == x->f->dir())
61             rotate(x->f), rotate(x);
        else rotate(x), rotate(x);
63     }
65 }
67 Splay *access(Splay *x) {
    Splay *q = nil;
    for (; x != nil; x = x->f) {
        splay(x);
        x->setCh(q, 1);
        q = x;
71     }
    return q;
73 }
75 void evert(Splay *x) {
    access(x);
    splay(x);
77

```



```

79  x->rev ^= 1;
80  x->push();
81  x->pull();
82  }
83  void link(Splay *x, Splay *y) {
84      // evert(x);
85      access(x);
86      splay(x);
87      evert(y);
88      x->setCh(y, 1);
89  }
90  void cut(Splay *x, Splay *y) {
91      // evert(x);
92      access(y);
93      splay(y);
94      y->push();
95      y->ch[0] = y->ch[0]->f = nil;
96  }
97
98  int N, Q;
99  Splay *vt[MXN];
100
101  int ask(Splay *x, Splay *y) {
102      access(x);
103      access(y);
104      splay(x);
105      int res = x->f->val;
106      if (res == -1) res = x->val;
107      return res;
108  }
109
110  int main(int argc, char **argv) {
111      scanf("%d%d", &N, &Q);
112      for (int i = 1; i <= N; i++)
113          vt[i] = new (Splay::pmem++) Splay(i);
114      while (Q--) {
115          char cmd[105];
116          int u, v;
117          scanf("%s", cmd);
118          if (cmd[1] == 'i') {
119              scanf("%d%d", &u, &v);
120              link(vt[u], vt[v]);
121          } else if (cmd[0] == 'c') {
122              scanf("%d", &v);
123              cut(vt[1], vt[v]);
124          } else {
125              scanf("%d%d", &u, &v);
126              int res = ask(vt[u], vt[v]);
127              printf("%d\n", res);
128          }
129      }
130  }

```

3. Graph

3.1. Modeling

- Maximum/Minimum flow with lower bound / Circulation problem
 - Construct super source S and sink T .
 - For each edge (x, y, l, u) , connect $x \rightarrow y$ with capacity $u - l$.
 - For each vertex v , denote by $in(v)$ the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
 - If $in(v) > 0$, connect $S \rightarrow v$ with capacity $in(v)$, otherwise, connect $v \rightarrow T$ with capacity $-in(v)$.
 - To maximize, connect $t \rightarrow s$ with capacity ∞ (skip this in circulation problem), and let f be the maximum flow from S to T . If $f \neq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, the maximum flow from s to t is the answer.
 - To minimize, let f be the maximum flow from S to T . Connect $t \rightarrow s$ with capacity ∞ and let the flow from S to T be f' . If $f + f' \neq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, f' is the answer.
 - The solution of each edge e is $l_e + f_e$, where f_e corresponds to the flow of edge e on the graph.
- Construct minimum vertex cover from maximum matching M on bipartite graph (X, Y)
 - Redirect every edge: $y \rightarrow x$ if $(x, y) \in M$, $x \rightarrow y$ otherwise.
 - DFS from unmatched vertices in X .
 - $x \in X$ is chosen iff x is unvisited.
 - $y \in Y$ is chosen iff y is visited.
- Minimum cost cyclic flow
 - Construct super source S and sink T
 - For each edge (x, y, c) , connect $x \rightarrow y$ with $(cost, cap) = (c, 1)$ if $c > 0$, otherwise connect $y \rightarrow x$ with $(cost, cap) = (-c, 1)$
 - For each edge with $c < 0$, sum these cost as K , then increase $d(y)$ by 1, decrease $d(x)$ by 1
 - For each vertex v with $d(v) > 0$, connect $S \rightarrow v$ with $(cost, cap) = (0, d(v))$

- For each vertex v with $d(v) < 0$, connect $v \rightarrow T$ with $(cost, cap) = (0, -d(v))$
- Flow from S to T , the answer is the cost of the flow $C + K$
- Maximum density induced subgraph
 - Binary search on answer, suppose we're checking answer T
 - Construct a max flow model, let K be the sum of all weights
 - Connect source $s \rightarrow v$, $v \in G$ with capacity K
 - For each edge (u, v, w) in G , connect $u \rightarrow v$ and $v \rightarrow u$ with capacity w
 - For $v \in G$, connect it with sink $v \rightarrow t$ with capacity $K + 2T - (\sum_{e \in E(v)} w(e)) - 2w(v)$
 - T is a valid answer if the maximum flow $f < K|V|$
- Minimum weight edge cover
 - For each $v \in V$ create a copy v' , and connect $u' \rightarrow v'$ with weight $w(u, v)$.
 - Connect $v \rightarrow v'$ with weight $2\mu(v)$, where $\mu(v)$ is the cost of the cheapest edge incident to v .
 - Find the minimum weight perfect matching on G' .
- Project selection problem
 - If $p_v > 0$, create edge (s, v) with capacity p_v ; otherwise, create edge (v, t) with capacity $-p_v$.
 - Create edge (u, v) with capacity w with w being the cost of choosing u without choosing v .
 - The mincut is equivalent to the maximum profit of a subset of projects.
- 0/1 quadratic programming

$$\sum_x c_x x + \sum_y c_y \bar{y} + \sum_{xy} c_{xy} x \bar{y} + \sum_{xyx'y'} c_{xyx'y'} (x \bar{y} + x' \bar{y}')$$

can be minimized by the mincut of the following graph:

- Create edge (x, t) with capacity c_x and create edge (s, y) with capacity c_y .
- Create edge (x, y) with capacity c_{xy} .
- Create edge (x, y) and edge (x', y') with capacity $c_{xyx'y'}$.

3.2. Matching/Flows

3.2.1. Dinic's Algorithm

```

1  struct Dinic {
2      struct edge {
3          int to, cap, flow, rev;
4      };
5      static constexpr int MAXN = 1000, MAXF = 1e9;
6      vector<edge> v[MAXN];
7      int top[MAXN], deep[MAXN], side[MAXN], s, t;
8      void make_edge(int s, int t, int cap) {
9          v[s].push_back({t, cap, 0, (int)v[t].size()});
10         v[t].push_back({s, 0, 0, (int)v[s].size() - 1});
11     }
12     int dfs(int a, int flow) {
13         if (a == t || !flow) return flow;
14         for (int &i = top[a]; i < v[a].size(); i++) {
15             edge &e = v[a][i];
16             if (deep[a] + 1 == deep[e.to] && e.cap - e.flow) {
17                 int x = dfs(e.to, min(e.cap - e.flow, flow));
18                 if (x) {
19                     e.flow += x, v[e.to][e.rev].flow -= x;
20                     return x;
21                 }
22             }
23         }
24         deep[a] = -1;
25         return 0;
26     }
27     bool bfs() {
28         queue<int> q;
29         fill_n(deep, MAXN, 0);
30         q.push(s), deep[s] = 1;
31         int tmp;
32         while (!q.empty()) {
33             tmp = q.front(), q.pop();
34             for (edge &e : v[tmp])
35                 if (!deep[e.to] && e.cap != e.flow)
36                     deep[e.to] = deep[tmp] + 1, q.push(e.to);
37         }
38         return deep[t];
39     }
40     int max_flow(int _s, int _t) {
41         s = _s, t = _t;
42         int flow = 0, tflow;
43         while (bfs()) {
44             fill_n(top, MAXN, 0);
45             while ((tflow = dfs(s, MAXF))) flow += tflow;
46         }
47         return flow;
48     }
49     void reset() {
50         fill_n(side, MAXN, 0);
51     }
52 }

```

```

51     for (auto &i : v) i.clear();
52 }
53 };

```

3.2.2. Minimum Cost Flow

```

1 struct MCF {
2     struct edge {
3         ll to, from, cap, flow, cost, rev;
4     } * fromE[MAXN];
5     vector<edge> v[MAXN];
6     ll n, s, t, flows[MAXN], dis[MAXN], pi[MAXN], flowlim;
7     void make_edge(int s, int t, ll cap, ll cost) {
8         if (!cap) return;
9         v[s].pb(edge{t, s, cap, 0LL, cost, v[t].size()});
10        v[t].pb(edge{s, t, 0LL, 0LL, -cost, v[s].size() - 1});
11    }
12    bitset<MAXN> vis;
13    void dijkstra() {
14        vis.reset();
15        __gnu_pbds::priority_queue<pair<ll, int>> q;
16        vector<decltype(q)::point_iterator> its(n);
17        q.push({0LL, s});
18        while (!q.empty()) {
19            int now = q.top().second;
20            q.pop();
21            if (vis[now]) continue;
22            vis[now] = 1;
23            ll ndis = dis[now] + pi[now];
24            for (edge &e : v[now]) {
25                if (e.flow == e.cap || vis[e.to]) continue;
26                if (dis[e.to] > ndis + e.cost - pi[e.to]) {
27                    dis[e.to] = ndis + e.cost - pi[e.to];
28                    flows[e.to] = min(flows[now], e.cap - e.flow);
29                    fromE[e.to] = &e;
30                    if (its[e.to] == q.end())
31                        its[e.to] = q.push({-dis[e.to], e.to});
32                    else q.modify(its[e.to], {-dis[e.to], e.to});
33                }
34            }
35        }
36    }
37    bool AP(ll &flow) {
38        fill_n(dis, n, INF);
39        fromE[s] = 0;
40        dis[s] = 0;
41        flows[s] = flowlim - flow;
42        dijkstra();
43        if (dis[t] == INF) return false;
44        flow += flows[t];
45        for (edge *e = fromE[t]; e; e = fromE[e->from]) {
46            e->flow += flows[t];
47            v[e->to][e->rev].flow -= flows[t];
48        }
49        for (int i = 0; i < n; i++)
50            pi[i] = min(pi[i] + dis[i], INF);
51        return true;
52    }
53    pll solve(int _s, int _t, ll _flowlim = INF) {
54        s = _s, t = _t, flowlim = _flowlim;
55        pll re;
56        while (re.F != flowlim && AP(re.F))
57            ;
58        for (int i = 0; i < n; i++)
59            for (edge &e : v[i])
60                if (e.flow != 0) re.S += e.flow * e.cost;
61        re.S /= 2;
62        return re;
63    }
64    void init(int _n) {
65        n = _n;
66        fill_n(pi, n, 0);
67        for (int i = 0; i < n; i++) v[i].clear();
68    }
69    void setpi(int s) {
70        fill_n(pi, n, INF);
71        pi[s] = 0;
72        for (ll it = 0, flag = 1, tdis; flag && it < n; it++) {
73            flag = 0;
74            for (int i = 0; i < n; i++)
75                if (pi[i] != INF)
76                    for (edge &e : v[i])
77                        if (e.cap && (tdis = pi[i] + e.cost) < pi[e.to])
78                            pi[e.to] = tdis, flag = 1;
79        }
80    }
81 };

```

3.2.3. Gomory-Hu Tree

Requires: Dinic's Algorithm

```

1 int e[MAXN][MAXN];
2 int p[MAXN];
3 Dinic D; // original graph
4 void gomory_hu() {
5     fill(p, p + n, 0);
6     fill(e[0], e[n], INF);
7     for (int s = 1; s < n; s++) {
8         int t = p[s];
9         Dinic F = D;
10        int tmp = F.max_flow(s, t);
11        for (int i = 1; i < s; i++)
12            e[s][i] = e[i][s] = min(tmp, e[t][i]);
13        for (int i = s + 1; i <= n; i++)
14            if (p[i] == t && F.side[i]) p[i] = s;
15    }
16 }

```

3.2.4. Global Minimum Cut

```

1 // weights is an adjacency matrix, undirected
2 pair<int, vi> getMinCut(vector<vi> &weights) {
3     int N = sz(weights);
4     vi used(N), cut, best_cut;
5     int best_weight = -1;
6
7     for (int phase = N - 1; phase >= 0; phase--) {
8         vi w = weights[0], added = used;
9         int prev, k = 0;
10        rep(i, 0, phase) {
11            prev = k;
12            k = -1;
13            rep(j, 1, N) if (!added[j] &&
14                            (k == -1 || w[j] > w[k])) k = j;
15            if (i == phase - 1) {
16                rep(j, 0, N) weights[prev][j] += weights[k][j];
17                rep(j, 0, N) weights[j][prev] = weights[prev][j];
18                used[k] = true;
19                cut.push_back(k);
20                if (best_weight == -1 || w[k] < best_weight) {
21                    best_cut = cut;
22                    best_weight = w[k];
23                }
24            } else {
25                rep(j, 0, N) w[j] += weights[k][j];
26                added[k] = true;
27            }
28        }
29    }
30    return {best_weight, best_cut};
31 }

```

3.2.5. Bipartite Minimum Cover

Requires: Dinic's Algorithm

```

1 // maximum independent set = all vertices not covered
2 // x : [0, n), y : [0, m]
3 struct Bipartite_vertex_cover {
4     Dinic D;
5     int n, m, s, t, x[maxn], y[maxn];
6     void make_edge(int x, int y) { D.make_edge(x, y + n, 1); }
7     int matching() {
8         int re = D.max_flow(s, t);
9         for (int i = 0; i < n; i++)
10            for (Dinic::edge &e : D.v[i])
11                if (e.to != s && e.flow == 1) {
12                    x[i] = e.to - n, y[e.to - n] = i;
13                    break;
14                }
15        return re;
16    }
17    // init() and matching() before use
18    void solve(vector<int> &vx, vector<int> &vy) {
19        bitset<maxn * 2 + 10> vis;
20        queue<int> q;
21        for (int i = 0; i < n; i++)
22            if (x[i] == -1) q.push(i), vis[i] = 1;
23        while (!q.empty()) {
24            int now = q.front();
25            q.pop();
26            if (now < n) {
27                for (Dinic::edge &e : D.v[now])
28                    if (e.to != s && e.to - n != x[now] && !vis[e.to])
29                        vis[e.to] = 1, q.push(e.to);
30            } else {
31                if (!vis[y[now - n]])
32                    vis[y[now - n]] = 1, q.push(y[now - n]);
33            }
34        }
35    }
36 }

```

```

35     for (int i = 0; i < n; i++)
36         if (!vis[i]) vx.pb(i);
37     for (int i = 0; i < m; i++)
38         if (vis[i + n]) vy.pb(i);
39 }
40 void init(int _n, int _m) {
41     n = _n, m = _m, s = n + m, t = s + 1;
42     for (int i = 0; i < n; i++)
43         x[i] = -1, D.make_edge(s, i, 1);
44     for (int i = 0; i < m; i++)
45         y[i] = -1, D.make_edge(i + n, t, 1);
46 }
47 };

```

3.2.6. Edmonds' Algorithm

```

1 struct Edmonds {
2     int n, T;
3     vector<vector<int>>> g;
4     vector<int> pa, p, used, base;
5     Edmonds(int n)
6         : n(n), T(0), g(n), pa(n, -1), p(n), used(n),
7           base(n) {}
8     void add(int a, int b) {
9         g[a].push_back(b);
10        g[b].push_back(a);
11    }
12    int getBase(int i) {
13        while (i != base[i])
14            base[i] = base[base[i]], i = base[i];
15        return i;
16    }
17    vector<int> toJoin;
18    void mark_path(int v, int x, int b, vector<int> &path) {
19        for (; getBase(v) != b; v = p[x]) {
20            p[v] = x, x = pa[v];
21            toJoin.push_back(v);
22            toJoin.push_back(x);
23            if (!used[x]) used[x] = ++T, path.push_back(x);
24        }
25    }
26    bool go(int v) {
27        for (int x : g[v]) {
28            int b, bv = getBase(v), bx = getBase(x);
29            if (bv == bx) continue;
30            else if (used[x]) {
31                vector<int> path;
32                toJoin.clear();
33                if (used[bx] < used[bv])
34                    mark_path(v, x, b = bx, path);
35                else mark_path(x, v, b = bv, path);
36                for (int z : toJoin) base[getBase(z)] = b;
37                for (int z : path)
38                    if (go(z)) return 1;
39            } else if (p[x] == -1) {
40                p[x] = v;
41                if (pa[x] == -1) {
42                    for (int y; x != -1; x = v)
43                        y = p[x], v = pa[y], pa[x] = y, pa[y] = x;
44                    return 1;
45                }
46                if (!used[pa[x]]) {
47                    used[pa[x]] = ++T;
48                    if (go(pa[x])) return 1;
49                }
50            }
51        }
52        return 0;
53    }
54    void init_dfs() {
55        for (int i = 0; i < n; i++)
56            used[i] = 0, p[i] = -1, base[i] = i;
57    }
58    bool dfs(int root) {
59        used[root] = ++T;
60        return go(root);
61    }
62    void match() {
63        int ans = 0;
64        for (int v = 0; v < n; v++)
65            for (int x : g[v])
66                if (pa[v] == -1 && pa[x] == -1) {
67                    pa[v] = x, pa[x] = v, ans++;
68                    break;
69                }
70        init_dfs();
71        for (int i = 0; i < n; i++)
72            if (pa[i] == -1 && dfs(i)) ans++, init_dfs();
73        cout << ans * 2 << "\n";

```

```

75     for (int i = 0; i < n; i++)
76         if (pa[i] > i)
77             cout << i + 1 << " " << pa[i] + 1 << "\n";
78 }
79 };

```

3.2.7. Minimum Weight Matching

```

1 struct Graph {
2     static const int MAXN = 105;
3     int n, e[MAXN][MAXN];
4     int match[MAXN], d[MAXN], onstk[MAXN];
5     vector<int> stk;
6     void init(int _n) {
7         n = _n;
8         for (int i = 0; i < n; i++)
9             for (int j = 0; j < n; j++)
10                // change to appropriate infinity
11                // if not complete graph
12                e[i][j] = 0;
13    }
14    void add_edge(int u, int v, int w) {
15        e[u][v] = e[v][u] = w;
16    }
17    bool SPFA(int u) {
18        if (onstk[u]) return true;
19        stk.push_back(u);
20        onstk[u] = 1;
21        for (int v = 0; v < n; v++) {
22            if (u != v && match[u] != v && !onstk[v]) {
23                int m = match[v];
24                if (d[m] > d[u] - e[v][m] + e[u][v]) {
25                    d[m] = d[u] - e[v][m] + e[u][v];
26                    onstk[v] = 1;
27                    stk.push_back(v);
28                    if (SPFA(m)) return true;
29                    stk.pop_back();
30                    onstk[v] = 0;
31                }
32            }
33        }
34        onstk[u] = 0;
35        stk.pop_back();
36        return false;
37    }
38    int solve() {
39        for (int i = 0; i < n; i += 2) {
40            match[i] = i + 1;
41            match[i + 1] = i;
42        }
43        while (true) {
44            int found = 0;
45            for (int i = 0; i < n; i++) onstk[i] = d[i] = 0;
46            for (int i = 0; i < n; i++) {
47                stk.clear();
48                if (!onstk[i] && SPFA(i)) {
49                    found = 1;
50                    while (stk.size() >= 2) {
51                        int u = stk.back();
52                        stk.pop_back();
53                        int v = stk.back();
54                        stk.pop_back();
55                        match[u] = v;
56                        match[v] = u;
57                    }
58                }
59            }
60            if (!found) break;
61        }
62        int ret = 0;
63        for (int i = 0; i < n; i++) ret += e[i][match[i]];
64        ret /= 2;
65        return ret;
66    }
67 } graph;

```

3.2.8. Stable Marriage

```

1 // normal stable marriage problem
2 /* input:
3 3
4 Albert Laura Nancy Marcy
5 Brad Marcy Nancy Laura
6 Chuck Laura Marcy Nancy
7 Laura Chuck Albert Brad
8 Marcy Albert Chuck Brad
9 Nancy Brad Albert Chuck
10 */
11

```



```

13 using namespace std;
14 const int MAXN = 505;
15
16 int n;
17 int favor[MAXN][MAXN]; // favor[boy_id][rank] = girl_id;
18 int order[MAXN][MAXN]; // order[girl_id][boy_id] = rank;
19 int current[MAXN]; // current[boy_id] = rank;
20 // boy_id will pursue current[boy_id] girl.
21 int girl_current[MAXN]; // girl[girl_id] = boy_id;
22
23 void initialize() {
24     for (int i = 0; i < n; i++) {
25         current[i] = 0;
26         girl_current[i] = n;
27         order[i][n] = n;
28     }
29 }
30
31 map<string, int> male, female;
32 string bname[MAXN], gname[MAXN];
33 int fit = 0;
34
35 void stable_marriage() {
36     queue<int> que;
37     for (int i = 0; i < n; i++) que.push(i);
38     while (!que.empty()) {
39         int boy_id = que.front();
40         que.pop();
41
42         int girl_id = favor[boy_id][current[boy_id]];
43         current[boy_id]++;
44
45         if (order[girl_id][boy_id] <
46             order[girl_id][girl_current[girl_id]]) {
47             if (girl_current[girl_id] < n)
48                 que.push(girl_current[girl_id]);
49             girl_current[girl_id] = boy_id;
50         } else {
51             que.push(boy_id);
52         }
53     }
54 }
55
56 int main() {
57     cin >> n;
58
59     for (int i = 0; i < n; i++) {
60         string p, t;
61         cin >> p;
62         male[p] = i;
63         bname[i] = p;
64         for (int j = 0; j < n; j++) {
65             cin >> t;
66             if (!female.count(t)) {
67                 gname[fit] = t;
68                 female[t] = fit++;
69             }
70             favor[i][j] = female[t];
71         }
72     }
73
74     for (int i = 0; i < n; i++) {
75         string p, t;
76         cin >> p;
77         for (int j = 0; j < n; j++) {
78             cin >> t;
79             order[female[p]][male[t]] = j;
80         }
81     }
82
83     initialize();
84     stable_marriage();
85
86     for (int i = 0; i < n; i++) {
87         cout << bname[i] << " "
88              << gname[favor[i][current[i] - 1]] << endl;
89     }
90 }

```

3.2.9. Kuhn-Munkres algorithm

```

1 // Maximum Weight Perfect Bipartite Matching
2 // Detect non-perfect-matching:
3 // 1. set all edge[i][j] as INF
4 // 2. if solve() >= INF, it is not perfect matching.
5
6 typedef long long ll;
7 struct KM {
8     static const int MAXN = 1050;
9     static const ll INF = 1LL << 60;

```

```

11 int n, match[MAXN], vx[MAXN], vy[MAXN];
12 ll edge[MAXN][MAXN], lx[MAXN], ly[MAXN], slack[MAXN];
13 void init(int _n) {
14     n = _n;
15     for (int i = 0; i < n; i++)
16         for (int j = 0; j < n; j++) edge[i][j] = 0;
17
18 void add_edge(int x, int y, ll w) { edge[x][y] = w; }
19 bool DFS(int x) {
20     vx[x] = 1;
21     for (int y = 0; y < n; y++) {
22         if (vy[y] continue;
23         if (lx[x] + ly[y] > edge[x][y]) {
24             slack[y] =
25                 min(slack[y], lx[x] + ly[y] - edge[x][y]);
26         } else {
27             vy[y] = 1;
28             if (match[y] == -1 || DFS(match[y])) {
29                 match[y] = x;
30                 return true;
31             }
32         }
33     }
34     return false;
35 }
36 ll solve() {
37     fill(match, match + n, -1);
38     fill(lx, lx + n, -INF);
39     fill(ly, ly + n, 0);
40     for (int i = 0; i < n; i++)
41         for (int j = 0; j < n; j++)
42             lx[i] = max(lx[i], edge[i][j]);
43     for (int i = 0; i < n; i++) {
44         fill(slack, slack + n, INF);
45         while (true) {
46             fill(vx, vx + n, 0);
47             fill(vy, vy + n, 0);
48             if (DFS(i)) break;
49             ll d = INF;
50             for (int j = 0; j < n; j++)
51                 if (!vy[j]) d = min(d, slack[j]);
52             for (int j = 0; j < n; j++) {
53                 if (vx[j]) lx[j] -= d;
54                 if (vy[j]) ly[j] += d;
55                 else slack[j] -= d;
56             }
57         }
58     }
59     ll res = 0;
60     for (int i = 0; i < n; i++) {
61         res += edge[match[i]][i];
62     }
63     return res;
64 }
65 } graph;

```

3.2.10. Network Simplex

```

1 #pragma once
2
3 using namespace std;
4
5 struct linked_lists {
6     int L, N;
7     vector<int> next, prev;
8
9     // L: lists are [0...L], N: elements are [0...N]
10    explicit linked_lists(int L = 0, int N = 0) {
11        assign(L, N);
12    }
13
14    int rep(int l) const {
15        return l + N;
16    } // "representative" of list l
17    int head(int l) const { return next[rep(l)]; }
18    int tail(int l) const { return prev[rep(l)]; }
19    bool empty(int l) const { return next[rep(l)] == rep(l); }
20
21    void push_front(int l, int n) {
22        link(rep(l), n, head(l));
23    }
24
25    void push_back(int l, int n) { link(tail(l), n, rep(l)); }
26    void insert_before(int i, int n) { link(prev[i], n, i); }
27    void insert_after(int i, int n) { link(i, n, next[i]); }
28    void erase(int n) { link(prev[n], next[n]); }
29    void pop_front(int l) { link(rep(l), next[head(l)]); }
30    void pop_back(int l) { link(prev[tail(l)], rep(l)); }
31
32    void clear() {

```

```

33     iota(begin(next) + N, end(next),
34           N); // sets next[rep(l)] = rep(l)
35     iota(begin(prev) + N, end(prev),
36           N); // sets prev[rep(l)] = rep(l)
37 }
38 void assign(int L, int N) {
39     this->L = L, this->N = N;
40     next.resize(N + L), prev.resize(N + L), clear();
41 }
42
43 private:
44     inline void link(int u, int v) {
45         next[u] = v, prev[v] = u;
46     }
47     inline void link(int u, int v, int w) {
48         link(u, v), link(v, w);
49     }
50 };
51
52 // Iterate through elements (call them i) of the list #l in
53 // "lists"
54 #define FOR_EACH_IN_LINKED_LIST(i, l, lists) \
55     for (int z##i = l, i = lists.head(z##i); \
56          i != lists.rep(z##i); i = lists.next[i])
57
58 /**
59  * Network simplex for minimum cost circulation with fixed
60  * supply/demand at nodes. Supports negative costs, negative
61  * cost cycles, self-loops and multiple edges fine.
62  *
63  * Flow should be large enough to hold sum of supplies and
64  * edge flows Cost should be large enough to hold sum of
65  * absolute costs * 2V (usually >=64 bits)
66  *
67  * Complexity: O(V) expected pivot, O(E) worst-case
68  *
69  * Usage:
70  *     network_simplex<int, long> mcc(V);
71  *     for (edges...) {
72  *         mcc.add(u, v, cap, cost);
73  *     }
74  *     for (nodes...) {
75  *         mcc.set_supply(u, supply);
76  *         mcc.set_demand(v, demand);
77  *     }
78  *     bool feasible = mcc.mincost_circulation();
79  *     auto mincost = mcc.circulation_cost();
80  *
81  * References:
82  *     LEMON network_simplex.h
83  *     OCW MIT MIT15_082JF10_av16.pdf
84  *     OCW MIT MIT15_082JF10_lec16.pdf
85  */
86 template <typename Flow, typename Cost>
87 struct network_simplex {
88     // we number the vertices 0,...,V-1, R is given number V
89
90     explicit network_simplex(int V) : V(V), node(V + 1) {}
91
92     void add(int u, int v, Flow cap, Cost cost) {
93         assert(0 <= u && u < V && 0 <= v && v < V);
94         edge.push_back({u, v}, cap, cost), E++;
95     }
96
97     void set_supply(int u, Flow supply) {
98         node[u].supply = supply;
99     }
100    void set_demand(int u, Flow demand) {
101        node[u].supply = -demand;
102    }
103    auto get_supply(int u) const { return node[u].supply; }
104
105    auto get_potential(int u) const { return node[u].pi; }
106    auto get_flow(int e) const { return edge[e].flow; }
107
108    auto reduced_cost(int e) const {
109        auto [u, v] = edge[e].node;
110        return edge[e].cost + node[u].pi - node[v].pi;
111    }
112
113    template <typename CostSum = Cost>
114    auto circulation_cost() const {
115        CostSum sum = 0;
116        for (int e = 0; e < E; e++) {
117            sum += edge[e].flow * CostSum(edge[e].cost);
118        }
119        return sum;
120    }
121
122    void verify_spanning_tree() const {
123        for (int e = 0; e < E; e++) {
124            assert(0 <= edge[e].flow &&
125                  edge[e].flow <= edge[e].cap);
126            assert(edge[e].flow == 0 || reduced_cost(e) <= 0);
127            assert(edge[e].flow == edge[e].cap ||
128                  reduced_cost(e) >= 0);
129        }
130    }
131
132    bool mincost_circulation() {
133        static constexpr bool INFEASIBLE = false,
134                               OPTIMAL = true;
135
136        // Check trivialities: positive cap[e] and sum of
137        // supplies is 0
138        Flow sum_supply = 0;
139        for (int u = 0; u < V; u++) {
140            sum_supply += node[u].supply;
141        }
142        if (sum_supply != 0) { return INFEASIBLE; }
143        for (int e = 0; e < E; e++) {
144            if (edge[e].cap < 0) { return INFEASIBLE; }
145        }
146
147        // Compute inf_cost as sum of all costs + 1, and reset
148        // the flow network
149        Cost inf_cost = 1;
150        for (int e = 0; e < E; e++) {
151            edge[e].flow = 0;
152            edge[e].state = STATE_LOWER;
153            inf_cost += abs(edge[e].cost);
154        }
155
156        edge.resize(E + V); // make space for V artificial edges
157        bfs.resize(V + 1);
158        children.assign(V + 1, V + 1);
159
160        // Add V artificial edges with infinite cost and initial
161        // supply for feasible flow
162        int root = V;
163        node[root] = {-1, -1, 0, 0};
164
165        for (int u = 0, e = E; u < V; u++, e++) {
166            // spanning tree links
167            node[u].parent = root, node[u].pred = e;
168            children.push_back(root, u);
169
170            auto supply = node[u].supply;
171            if (supply >= 0) {
172                node[u].pi = -inf_cost;
173                edge[e] = {
174                    u, root, supply, inf_cost, supply, STATE_TREE;
175                }
176            } else {
177                node[u].pi = inf_cost;
178                edge[e] = {
179                    root, u, -supply, inf_cost, -supply, STATE_TREE;
180                }
181            }
182
183            // We want to, hopefully, find a pivot edge in
184            // O(sqrt(E)). This can be tuned.
185            block_size =
186                max(int(ceil(sqrt(E + V))), min(10, V + 1));
187            next_arc = 0;
188
189            // Pivot pivot pivot
190            int e_in = select_pivot_edge();
191            while (e_in != -1) {
192                pivot(e_in);
193                e_in = select_pivot_edge();
194            }
195
196            // If there is >0 flow through an artificial edge, the
197            // problem is infeasible.
198            for (int e = E; e < E + V; e++) {
199                if (edge[e].flow > 0) {
200                    edge.resize(E);
201                    return INFEASIBLE;
202                }
203            }
204            edge.resize(E);
205            return OPTIMAL;
206        }
207    private:
208        enum ArcState : int8_t {
209            STATE_UPPER = -1,
210            STATE_TREE = 0,
211            STATE_LOWER = 1
212        };

```

```

213 struct Node {
215     int parent, pred;
        Flow supply;
        Cost pi;
    };
219 struct Edge {
        array<int, 2> node; // [0]->[1]
        Flow cap;
        Cost cost;
        Flow flow = 0;
        ArcState state = STATE_LOWER;
225 };

227 int V, E = 0, next_arc = 0, block_size = 0;
        vector<Node> node;
        vector<Edge> edge;
        linked_lists children;
231 vector<int> bfs; // scratchpad for downwards bfs and evert

233 int select_pivot_edge() {
        // block search: check block_size edges looping, and
        // pick the lowest reduced cost
        Cost minimum = 0;
237     int e_in = -1;
        int count = block_size, seen_edges = E + V;
239     for (int &e = next_arc; seen_edges-- > 0;
            e = e + 1 == E + V ? 0 : e + 1) {
241         if (minimum > edge[e].state * reduced_cost(e)) {
            minimum = edge[e].state * reduced_cost(e);
            e_in = e;
243         }
245         if (--count == 0 && minimum < 0) {
            break;
247         } else if (count == 0) {
            count = block_size;
249         }
251     }
    return e_in;
253 }

255 void pivot(int e_in) {
        // Find lca of u_in and v_in with two pointers technique
        auto [u_in, v_in] = edge[e_in].node;
257     int a = u_in, b = v_in;
        while (a != b) {
259         a = node[a].parent == -1 ? v_in : node[a].parent;
            b = node[b].parent == -1 ? u_in : node[b].parent;
261     }
        int lca = a;

263     // Orient the edge so that we add flow along u_in->v_in
        if (edge[e_in].state == STATE_UPPER) {
265         swap(u_in, v_in);
267     }

269     // Let's find the saturating flow and exiting arc
        enum OutArcSide { SAME_EDGE, U_IN_SIDE, V_IN_SIDE };
        OutArcSide side = SAME_EDGE;
        Flow flow_delta = edge[e_in].cap;
273     int u_out = -1;

275     // Go up from u_in to lca, break ties by preferring lower
        // vertices
277     for (int u = u_in; u != lca && flow_delta > 0;
            u = node[u].parent) {
279         int e = node[u].pred;
            bool edge_down = u == edge[e].node[1];
281         Flow to_saturate =
            edge_down ? edge[e].cap - edge[e].flow : edge[e].flow;

283         if (flow_delta > to_saturate) {
            flow_delta = to_saturate;
285             u_out = u;
            side = U_IN_SIDE;
287         }
289     }

291     // Go up from v_in to lca, break ties by preferring
        // higher vertices
293     for (int u = v_in; u != lca; u = node[u].parent) {
        int e = node[u].pred;
295         bool edge_up = u == edge[e].node[0];
            Flow to_saturate =
            edge_up ? edge[e].cap - edge[e].flow : edge[e].flow;

297         if (flow_delta >= to_saturate) {
            flow_delta = to_saturate;
299             u_out = u;
            side = V_IN_SIDE;
301         }
303     }

305     // Augment along the cycle if we can push anything
        if (flow_delta > 0) {
307         auto delta = edge[e_in].state * flow_delta;
            edge[e_in].flow += delta;

309         for (int u = edge[e_in].node[0]; u != lca;
            u = node[u].parent) {
311             int e = node[u].pred;
            edge[e].flow +=
313                 u == edge[e].node[0] ? -delta : delta;
        }
315         for (int u = edge[e_in].node[1]; u != lca;
            u = node[u].parent) {
317             int e = node[u].pred;
            edge[e].flow +=
319                 u == edge[e].node[0] ? delta : -delta;
        }
321     }

323     // Return now if we didn't change the spanning tree. The
        // state of e_in flipped.
325     if (side == SAME_EDGE) {
        edge[e_in].state = ArcState(-edge[e_in].state);
327         return;
329     }

331     // Basis exchange: Replace out_arc with e_in in the
        // spanning tree
333     int out_arc = node[u_out].pred;
        edge[e_in].state = STATE_TREE;
335     edge[out_arc].state =
        edge[out_arc].flow ? STATE_UPPER : STATE_LOWER;
337

339     // Put u_in on the same side as u_out
        if (side == V_IN_SIDE) { swap(u_in, v_in); }

341     // Evert: Walk up from u_in to u_out, and fix
        // parent/pred/child pointers downwards
343     int i = 0, S = 0;
        for (int u = u_in; u != u_out; u = node[u].parent) {
345         bfs[S++] = u;
347     }
        for (i = S - 1; i >= 0; i--) {
349             int u = bfs[i], p = node[u].parent;
            children.erase(p); // remove p from its children list
            // and add it to u's
351             children.push_back(u, p);
            node[p].parent = u;
353             node[p].pred = node[u].pred;
355         }
        children.erase(u_in); // remove u_in from its children
        // list and add it to v_in's
357         children.push_back(v_in, u_in);
        node[u_in].parent = v_in;
359         node[u_in].pred = e_in;

361     // Fix potentials: Visit the subtree of u_in (pi_delta
        // is not 0).
        Cost current_pi = reduced_cost(e_in);
363         Cost pi_delta =
        u_in == edge[e_in].node[0] ? -current_pi : current_pi;

365         bfs[0] = u_in;
        for (i = 0, S = 1; i < S; i++) {
367             int u = bfs[i];
            node[u].pi += pi_delta;
            FOR_EACH_IN_LINKED_LIST(v, u, children) {
371                 bfs[S++] = v;
373             }
375         }
377     };
}

```

3.3. Shortest Path Faster Algorithm

```

1 struct SPFA {
        static const int maxn = 1010, INF = 1e9;
        int dis[maxn];
        bitset<maxn> inq, inneg;
        queue<int> q, tq;
        vector<pii> v[maxn];
        void make_edge(int s, int t, int w) {
            v[s].emplace_back(t, w);
        }
        void dfs(int a) {
            inneg[a] = 1;
            for (pii i : v[a])
                if (!inneg[i.F]) dfs(i.F);
13     }
}

```

```

15 bool solve(int n, int s) { // true if have neg-cycle
    for (int i = 0; i <= n; i++) dis[i] = INF;
    dis[s] = 0, q.push(s);
    for (int i = 0; i < n; i++) {
        inq.reset();
        int now;
        while (!q.empty()) {
            now = q.front(), q.pop();
            for (pii &i : v[now]) {
                if (dis[i.F] > dis[now] + i.S) {
                    dis[i.F] = dis[now] + i.S;
                    if (!inq[i.F]) tq.push(i.F), inq[i.F] = 1;
                }
            }
            q.swap(tq);
        }
        bool re = !q.empty();
        inneg.reset();
        while (!q.empty()) {
            if (!inneg[q.front()]) dfs(q.front());
            q.pop();
        }
        return re;
    }
    void reset(int n) {
        for (int i = 0; i <= n; i++) v[i].clear();
    }
};

```

3.4. Strongly Connected Components

```

1 struct TarjanScc {
    int n, step;
    vector<int> time, low, instk, stk;
    vector<vector<int>> e, scc;
    TarjanScc(int n_)
        : n(n_), step(0), time(n), low(n), instk(n), e(n) {}
    void add_edge(int u, int v) { e[u].push_back(v); }
    void dfs(int x) {
        time[x] = low[x] = ++step;
        stk.push_back(x);
        instk[x] = 1;
        for (int y : e[x])
            if (!time[y]) {
                dfs(y);
                low[x] = min(low[x], low[y]);
            } else if (instk[y]) {
                low[x] = min(low[x], time[y]);
            }
        if (time[x] == low[x]) {
            scc.emplace_back();
            for (int y = -1; y != x; ) {
                y = stk.back();
                stk.pop_back();
                instk[y] = 0;
                scc.back().push_back(y);
            }
        }
    }
    void solve() {
        for (int i = 0; i < n; i++)
            if (!time[i]) dfs(i);
        reverse(scc.begin(), scc.end());
        // scc in topological order
    }
};

```

3.4.1. 2-Satisfiability

Requires: Strongly Connected Components

```

1 // 1 based, vertex in SCC = MAXN * 2
// (not i) is i + n
2 struct two_SAT {
    int n, ans[MAXN];
    SCC S;
    void imply(int a, int b) { S.make_edge(a, b); }
    bool solve(int _n) {
        n = _n;
        S.solve(n * 2);
        for (int i = 1; i <= n; i++) {
            if (S.scc[i] == S.scc[i + n]) return false;
            ans[i] = (S.scc[i] < S.scc[i + n]);
        }
        return true;
    }
    void init(int _n) {
        n = _n;
        fill_n(ans, n + 1, 0);
    }
};

```

```

19 S.init(n * 2);
20 }
21 } SAT;

```

3.5. Biconnected Components

3.5.1. Articulation Points

```

1 void dfs(int x, int p) {
    tin[x] = low[x] = ++t;
    int ch = 0;
    for (auto u : g[x])
        if (u.first != p) {
            if (!ins[u.second])
                st.push(u.second), ins[u.second] = true;
            if (tin[u.first]) {
                low[x] = min(low[x], tin[u.first]);
                continue;
            }
            ++ch;
            dfs(u.first, x);
            low[x] = min(low[x], low[u.first]);
            if (low[u.first] >= tin[x]) {
                cut[x] = true;
                ++sz;
                while (true) {
                    int e = st.top();
                    st.pop();
                    bcc[e] = sz;
                    if (e == u.second) break;
                }
            }
        }
    if (ch == 1 && p == -1) cut[x] = false;
}

```

3.5.2. Bridges

```

1 // if there are multi-edges, then they are not bridges
2 void dfs(int x, int p) {
    tin[x] = low[x] = ++t;
    st.push(x);
    for (auto u : g[x])
        if (u.first != p) {
            if (tin[u.first]) {
                low[x] = min(low[x], tin[u.first]);
                continue;
            }
            dfs(u.first, x);
            low[x] = min(low[x], low[u.first]);
            if (low[u.first] == tin[u.first]) br[u.second] = true;
        }
    if (tin[x] == low[x]) {
        ++sz;
        while (st.size()) {
            int u = st.top();
            st.pop();
            bcc[u] = sz;
            if (u == x) break;
        }
    }
}

```

3.6. Triconnected Components

```

1 // requires a union-find data structure
2 struct ThreeEdgeCC {
    int V, ind;
    vector<int> id, pre, post, low, deg, path;
    vector<vector<int>> components;
    UnionFind uf;
    template <class Graph>
    void dfs(const Graph &G, int v, int prev) {
        pre[v] = ++ind;
        for (int w : G[v])
            if (w != v) {
                if (w == prev) {
                    prev = -1;
                    continue;
                }
                if (pre[w] != -1) {
                    if (pre[w] < pre[v]) {
                        deg[v]++;
                        low[v] = min(low[v], pre[w]);
                    } else {
                        deg[v]--;
                        int &u = path[v];
                        for (; u != -1 && pre[u] <= pre[w] &&
                             pre[w] <= post[u];) {
                            uf.join(v, u);
                        }
                    }
                }
                dfs(G, w, v);
            }
    }
};

```

```

    deg[v] += deg[u];
    u = path[u];
}
}
continue;
}
dfs(G, w, v);
if (path[w] == -1 && deg[w] <= 1) {
    deg[v] += deg[w];
    low[v] = min(low[v], low[w]);
    continue;
}
if (deg[w] == 0) w = path[w];
if (low[v] > low[w]) {
    low[v] = min(low[v], low[w]);
    swap(w, path[v]);
}
for (; w != -1; w = path[w]) {
    uf.join(v, w);
    deg[v] += deg[w];
}
}
post[v] = ind;
}
template <class Graph>
ThreeEdgeCC(const Graph &G)
: V(G.size()), ind(-1), id(V, -1), pre(V, -1),
  post(V), low(V, INT_MAX), deg(V, 0), path(V, -1),
  uf(V) {
    for (int v = 0; v < V; v++)
        if (pre[v] == -1) dfs(G, v, -1);
    components.reserve(uf.cnt);
    for (int v = 0; v < V; v++)
        if (uf.find(v) == v) {
            id[v] = components.size();
            components.emplace_back(1, v);
            components.back().reserve(uf.getSize(v));
        }
    for (int v = 0; v < V; v++)
        if (id[v] == -1)
            components[id[v] = id[uf.find(v)]] .push_back(v);
};

```

3.7. Centroid Decomposition

```

1 void get_center(int now) {
    v[now] = true;
    vtx.push_back(now);
    sz[now] = 1;
    mx[now] = 0;
    for (int u : G[now])
        if (!v[u]) {
            get_center(u);
            mx[now] = max(mx[now], sz[u]);
            sz[now] += sz[u];
        }
}
13 void get_dis(int now, int d, int len) {
    dis[d][now] = cnt;
    v[now] = true;
    for (auto u : G[now])
        if (!v[u.first]) { get_dis(u, d, len + u.second); }
}
19 void dfs(int now, int fa, int d) {
    get_center(now);
    int c = -1;
    for (int i : vtx) {
        if (max(mx[i], (int)vtx.size() - sz[i]) <=
            (int)vtx.size() / 2)
            c = i;
        v[i] = false;
    }
    get_dis(c, d, 0);
    for (int i : vtx) v[i] = false;
    v[c] = true;
    vtx.clear();
    dep[c] = d;
    p[c] = fa;
    for (auto u : G[c])
        if (u.first != fa && !v[u.first]) {
            dfs(u.first, c, d + 1);
        }
}

```

3.8. Minimum Mean Cycle

```

1 // source: waynedisonitau123
3 // d[i][j] == 0 if {i,j} !in E

```

```

long long d[1003][1003], dp[1003][1003];
5 pair<long long, long long> MMWC() {
    memset(dp, 0x3f, sizeof(dp));
    for (int i = 1; i <= n; ++i) dp[0][i] = 0;
    for (int i = 1; i <= n; ++i) {
        for (int j = 1; j <= n; ++j) {
            for (int k = 1; k <= n; ++k) {
                dp[i][k] = min(dp[i - 1][j] + d[j][k], dp[i][k]);
            }
        }
    }
    long long au = 1ll << 31, ad = 1;
    for (int i = 1; i <= n; ++i) {
        if (dp[n][i] == 0x3f3f3f3f3f3f3f3f) continue;
        long long u = 0, d = 1;
        for (int j = n - 1; j >= 0; --j) {
            if ((dp[n][i] - dp[j][i]) * d > u * (n - j)) {
                u = dp[n][i] - dp[j][i];
                d = n - j;
            }
        }
        if (u * ad < au * d) au = u, ad = d;
    }
    long long g = __gcd(au, ad);
    return make_pair(au / g, ad / g);
}

```

3.9. Directed MST

```

1 template <typename T> struct DMST {
    T g[maxn][maxn], fw[maxn];
    int n, fr[maxn];
    bool vis[maxn], inc[maxn];
    void clear() {
        for (int i = 0; i < maxn; ++i) {
            for (int j = 0; j < maxn; ++j) g[i][j] = inf;
            vis[i] = inc[i] = false;
        }
    }
    void addedge(int u, int v, T w) {
        g[u][v] = min(g[u][v], w);
    }
    T operator()(int root, int _n) {
        n = _n;
        if (dfs(root) != n) return -1;
        T ans = 0;
        while (true) {
            for (int i = 1; i <= n; ++i) fw[i] = inf, fr[i] = i;
            for (int i = 1; i <= n; ++i)
                if (!inc[i]) {
                    for (int j = 1; j <= n; ++j) {
                        if (!inc[j] && i != j && g[j][i] < fw[i]) {
                            fw[i] = g[j][i];
                            fr[i] = j;
                        }
                    }
                }
            int x = -1;
            for (int i = 1; i <= n; ++i)
                if (i != root && !inc[i]) {
                    int j = i, c = 0;
                    while (j != root && fr[j] != i && c <= n)
                        ++c, j = fr[j];
                    if (j == root || c > n) continue;
                    else {
                        x = i;
                        break;
                    }
                }
            if (!x) break;
            for (int i = 1; i <= n; ++i)
                if (i != root && !inc[i]) ans += fw[i];
            return ans;
        }
        int y = x;
        for (int i = 1; i <= n; ++i) vis[i] = false;
        do {
            ans += fw[y];
            y = fr[y];
            vis[y] = inc[y] = true;
        } while (y != x);
        inc[x] = false;
        for (int k = 1; k <= n; ++k)
            if (vis[k]) {
                for (int j = 1; j <= n; ++j)
                    if (!vis[j]) {
                        if (g[x][j] > g[k][j]) g[x][j] = g[k][j];
                        if (g[j][k] < inf &&
                            g[j][k] - fw[k] < g[j][x])

```



```

61         g[j][x] = g[j][k] - fw[k];
62     }
63 }
64 }
65 return ans;
66 }
67 int dfs(int now) {
68     int r = 1;
69     vis[now] = true;
70     for (int i = 1; i <= n; ++i)
71         if (g[now][i] < inf && !vis[i]) r += dfs(i);
72     return r;
73 };

```

3.10. Maximum Clique

```

1 // source: KACTL
3 typedef vector<bitset<200>> vb;
4 struct Maxclique {
5     double limit = 0.025, pk = 0;
6     struct Vertex {
7         int i, d = 0;
8     };
9     typedef vector<Vertex> vv;
10    vb e;
11    vv V;
12    vector<vi> C;
13    vi qmax, q, S, old;
14    void init(vv &r) {
15        for (auto &v : r) v.d = 0;
16        for (auto &v : r)
17            for (auto j : r) v.d += e[v.i][j.i];
18        sort(all(r), [](auto a, auto b) { return a.d > b.d; });
19        int mxD = r[0].d;
20        rep(i, 0, sz(r)) r[i].d = min(i, mxD) + 1;
21    }
22    void expand(vv &r, int lev = 1) {
23        S[lev] += S[lev - 1] - old[lev];
24        old[lev] = S[lev - 1];
25        while (sz(R)) {
26            if (sz(q) + R.back().d <= sz(qmax)) return;
27            q.push_back(R.back().i);
28            vv T;
29            for (auto v : R)
30                if (e[R.back().i][v.i]) T.push_back({v.i});
31            if (sz(T)) {
32                if (S[lev]++ / ++pk < limit) init(T);
33                int j = 0, mxk = 1,
34                    mnk = max(sz(qmax) - sz(q) + 1, 1);
35                C[1].clear(), C[2].clear();
36                for (auto v : T) {
37                    int k = 1;
38                    auto f = [&](int i) { return e[v.i][i]; };
39                    while (any_of(all(C[k]), f)) k++;
40                    if (k > mxk) mxk = k, C[mxk + 1].clear();
41                    if (k < mnk) T[j++].i = v.i;
42                    C[k].push_back(v.i);
43                }
44                if (j > 0) T[j - 1].d = 0;
45                rep(k, mnk, mxk + 1) for (int i : C[k]) T[j].i = i,
46                    T[j++].d = k;
47                expand(T, lev + 1);
48            } else if (sz(q) > sz(qmax)) qmax = q;
49            q.pop_back(), R.pop_back();
50        }
51    }
52    vi maxClique() {
53        init(V), expand(V);
54        return qmax;
55    }
56    Maxclique(vb conn)
57        : e(conn), C(sz(e) + 1), S(sz(C)), old(S) {
58        rep(i, 0, sz(e)) V.push_back({i});
59    }
60 };

```

3.11. Dominator Tree

```

1 // idom[n] is the unique node that strictly dominates n but
2 // does not strictly dominate any other node that strictly
3 // dominates n. idom[n] = 0 if n is entry or the entry
4 // cannot reach n.
5 struct DominatorTree {
6     static const int MAXN = 200010;
7     int n, s;
8     vector<int> g[MAXN], pred[MAXN];
9     vector<int> cov[MAXN];

```

```

11 int dfn[MAXN], nfd[MAXN], ts;
12 int par[MAXN];
13 int sdom[MAXN], idom[MAXN];
14 int mom[MAXN], mn[MAXN];
15 inline bool cmp(int u, int v) { return dfn[u] < dfn[v]; }
16
17 int eval(int u) {
18     if (mom[u] == u) return u;
19     int res = eval(mom[u]);
20     if (cmp(sdom[mn[mom[u]]], sdom[mn[u]]))
21         mn[u] = mn[mom[u]];
22     return mom[u] = res;
23 }
24
25 void init(int _n, int _s) {
26     n = _n;
27     s = _s;
28     REP1(i, 1, n) {
29         g[i].clear();
30         pred[i].clear();
31         idom[i] = 0;
32     }
33 }
34 void add_edge(int u, int v) {
35     g[u].push_back(v);
36     pred[v].push_back(u);
37 }
38 void DFS(int u) {
39     ts++;
40     dfn[u] = ts;
41     nfd[ts] = u;
42     for (int v : g[u])
43         if (dfn[v] == 0) {
44             par[v] = u;
45             DFS(v);
46         }
47 }
48 void build() {
49     ts = 0;
50     REP1(i, 1, n) {
51         dfn[i] = nfd[i] = 0;
52         cov[i].clear();
53         mom[i] = mn[i] = sdom[i] = i;
54     }
55     DFS(s);
56     for (int i = ts; i >= 2; i--) {
57         int u = nfd[i];
58         if (u == 0) continue;
59         for (int v : pred[u])
60             if (dfn[v]) {
61                 eval(v);
62                 if (cmp(sdom[mn[v]], sdom[u]))
63                     sdom[u] = sdom[mn[v]];
64             }
65         cov[sdom[u]].push_back(u);
66         mom[u] = par[u];
67         for (int w : cov[par[u]]) {
68             eval(w);
69             if (cmp(sdom[mn[w]], par[u])) idom[w] = mn[w];
70             else idom[w] = par[u];
71         }
72         cov[par[u]].clear();
73     }
74     REP1(i, 2, ts) {
75         int u = nfd[i];
76         if (u == 0) continue;
77         if (idom[u] != sdom[u]) idom[u] = idom[idom[u]];
78     }
79 } dom;

```

3.12. Manhattan Distance MST

```

1 // source: KACTL
2
3 // returns [(dist, from, to), ...]
4 // then do normal mst afterwards
5 typedef Point<int> P;
6 vector<array<int, 3>> manhattanMST(vector<P> ps) {
7     vi id(sz(ps));
8     iota(all(id), 0);
9     vector<array<int, 3>> edges;
10    rep(k, 0, 4) {
11        sort(all(id), [&](int i, int j) {
12            return (ps[i] - ps[j]).x < (ps[j] - ps[i]).y;
13        });
14        map<int, int> sweep;
15        for (int i : id) {
16            for (auto it = sweep.lower_bound(-ps[i].y);

```

```

17     it != sweep.end(); sweep.erase(it++)) {
18         int j = it->second;
19         P d = ps[i] - ps[j];
20         if (d.y > d.x) break;
21         edges.push_back({d.y + d.x, i, j});
22     }
23     sweep[-ps[i].y] = i;
24 }
25 for (P &p : ps)
26     if (k & 1) p.x = -p.x;
27     else swap(p.x, p.y);
28 }
29 return edges;

```

4. Math

4.1. Number Theory

4.1.1. Mod Struct

A list of safe primes: 26003, 27767, 28319, 28979, 29243, 29759, 30467, 910927547, 919012223, 947326223, 990669467, 1007939579, 1019126699, 929760389146037459, 975500632317046523, 989312547895528379

NTT prime p	$p - 1$	primitive root
65537	$1 \ll 16$	3
998244353	$119 \ll 23$	3
2748779069441	$5 \ll 39$	3
1945555039024054273	$27 \ll 56$	5

Requires: Extended GCD

```

1 template <typename T> struct M {
2     static T MOD; // change to constexpr if already known
3     T v;
4     M(T x = 0) {
5         v = (-MOD <= x && x < MOD) ? x : x % MOD;
6         if (v < 0) v += MOD;
7     }
8     explicit operator T() const { return v; }
9     bool operator==(const M &b) const { return v == b.v; }
10    bool operator!=(const M &b) const { return v != b.v; }
11    M operator-() { return M(-v); }
12    M operator+(M b) { return M(v + b.v); }
13    M operator-(M b) { return M(v - b.v); }
14    M operator*(M b) { return M((__int128)v * b.v % MOD); }
15    M operator/(M b) { return *this * (b ^ (MOD - 2)); }
16    // change above implementation to this if MOD is not prime
17    M inv() {
18        auto [p, _, g] = extgcd(v, MOD);
19        return assert(g == 1), p;
20    }
21    friend M operator^(M a, ll b) {
22        M ans(1);
23        for (; b >= 1, a *= a)
24            if (b & 1) ans *= a;
25        return ans;
26    }
27    friend M operator+=(M &a, M b) { return a = a + b; }
28    friend M operator-=(M &a, M b) { return a = a - b; }
29    friend M operator*=(M &a, M b) { return a = a * b; }
30    friend M operator/=(M &a, M b) { return a = a / b; }
31 };
32 using Mod = M<int>;
33 template <> int Mod::MOD = 1'000'000'007;
34 int &MOD = Mod::MOD;

```

4.1.2. Miller-Rabin

Requires: Mod Struct

```

1 // checks if Mod::MOD is prime
2 bool is_prime() {
3     if (MOD < 2 || MOD % 2 == 0) return MOD == 2;
4     Mod A[] = {2, 7, 61}; // for int values (< 2^31)
5     // ll: 2, 325, 9375, 28178, 450775, 9780504, 1795265022
6     int s = __builtin_ctzll(MOD - 1), i;
7     for (Mod a : A) {
8         Mod x = a ^ (MOD >> s);
9         for (i = 0; i < s && (x + 1).v > 2; i++) x *= x;
10        if (i && x != -1) return 0;
11    }
12    return 1;
13 }

```

4.1.3. Linear Sieve

```

1 constexpr ll MAXN = 1000000;
2 bitset<MAXN> is_prime;
3 vector<ll> primes;
4 ll mpf[MAXN], phi[MAXN], mu[MAXN];
5
6 void sieve() {
7     is_prime.set();
8     is_prime[1] = 0;
9     mu[1] = phi[1] = 1;
10    for (ll i = 2; i < MAXN; i++) {
11        if (is_prime[i]) {
12            mpf[i] = i;
13            primes.push_back(i);
14            phi[i] = i - 1;
15            mu[i] = -1;
16        }
17        for (ll p : primes) {
18            if (p > mpf[i] || i * p >= MAXN) break;
19            is_prime[i * p] = 0;
20            mpf[i * p] = p;
21            mu[i * p] = -mu[i];
22            if (i % p == 0)
23                phi[i * p] = phi[i] * p, mu[i * p] = 0;
24            else phi[i * p] = phi[i] * (p - 1);
25        }
26    }
27 }

```

4.1.4. Get Factors

Requires: Linear Sieve

```

1 vector<ll> all_factors(ll n) {
2     vector<ll> fac = {1};
3     while (n > 1) {
4         const ll p = mpf[n];
5         vector<ll> cur = {1};
6         while (n % p == 0) {
7             n /= p;
8             cur.push_back(cur.back() * p);
9         }
10        vector<ll> tmp;
11        for (auto x : fac)
12            for (auto y : cur) tmp.push_back(x * y);
13        tmp.swap(fac);
14    }
15    return fac;
16 }

```

4.1.5. Binary GCD

```

1 // returns the gcd of non-negative a, b
2 ull bin_gcd(ull a, ull b) {
3     if (!a || !b) return a + b;
4     int s = __builtin_ctzll(a | b);
5     a >>= __builtin_ctzll(a);
6     while (b) {
7         if ((b >>= __builtin_ctzll(b)) < a) swap(a, b);
8         b >>= a;
9     }
10    return a << s;
11 }

```

4.1.6. Extended GCD

```

1 // returns (p, q, g): p * a + q * b == g == gcd(a, b)
2 // g is not guaranteed to be positive when a < 0 or b < 0
3 tuple<ll, ll, ll> extgcd(ll a, ll b) {
4     ll s = 1, t = 0, u = 0, v = 1;
5     while (b) {
6         ll q = a / b;
7         swap(a -= q * b, b);
8         swap(s -= q * t, t);
9         swap(u -= q * v, v);
10    }
11    return {s, u, a};
12 }

```

4.1.7. Chinese Remainder Theorem

Requires: Extended GCD

```

1 // for 0 <= a < m, 0 <= b < n, returns the smallest x >= 0
2 // such that x % m == a and x % n == b
3 ll crt(ll a, ll m, ll b, ll n) {
4     if (n > m) swap(a, b), swap(m, n);
5     auto [x, y, g] = extgcd(m, n);
6     assert((a - b) % g == 0); // no solution
7     x = ((b - a) / g * x) % (n / g) * m + a;
8     return x < 0 ? x + m / g * n : x;
9 }

```

4.1.8. Baby-Step Giant-Step

Requires: Mod Struct

```

1 // returns x such that a ^ x = b where x \in [l, r)
2 ll bsgs(Mod a, Mod b, ll l = 0, ll r = MOD - 1) {
3     int m = sqrt(r - l) + 1, i;
4     unordered_map<ll, ll> tb;
5     Mod d = (a ^ l) / b;
6     for (i = 0, d = (a ^ l) / b; i < m; i++, d *= a)
7         if (d == 1) return l + i;
8     else tb[(ll)d] = l + i;
9     Mod c = Mod(1) / (a ^ m);
10    for (i = 0, d = 1; i < m; i++, d *= c)
11        if (auto j = tb.find((ll)d); j != tb.end())
12            return j->second + i * m;
13    return assert(0), -1; // no solution
14 }

```

4.1.9. Pollard's Rho

```

1 ll f(ll x, ll mod) { return (x * x + 1) % mod; }
2 // n should be composite
3 ll pollard_rho(ll n) {
4     if (!(n & 1)) return 2;
5     while (1) {
6         ll y = 2, x = RNG() % (n - 1) + 1, res = 1;
7         for (int sz = 2; res == 1; sz *= 2) {
8             for (int i = 0; i < sz && res == 1; i++) {
9                 x = f(x, n);
10                res = __gcd(abs(x - y), n);
11            }
12            y = x;
13        }
14        if (res != 0 && res != n) return res;
15    }
16 }

```

4.1.10. Tonelli-Shanks Algorithm

Requires: Mod Struct

```

1 int legendre(Mod a) {
2     if (a == 0) return 0;
3     return (a ^ ((MOD - 1) / 2)) == 1 ? 1 : -1;
4 }
5 Mod sqrt(Mod a) {
6     assert(legendre(a) != -1); // no solution
7     ll p = MOD, s = p - 1;
8     if (a == 0) return 0;
9     if (p == 2) return 1;
10    if (p % 4 == 3) return a ^ ((p + 1) / 4);
11    int r, m;
12    for (r = 0; !(s & 1); r++) s >>= 1;
13    Mod n = 2;
14    while (legendre(n) != -1) n += 1;
15    Mod x = a ^ ((s + 1) / 2), b = a ^ s, g = n ^ s;
16    while (b != 1) {
17        Mod t = b;
18        for (m = 0; t != 1; m++) t *= t;
19        Mod gs = g ^ (1LL << (r - m - 1));
20        g = gs * gs, x *= gs, b *= g, r = m;
21    }
22    return x;
23 }
24 // to get sqrt(x) modulo p^k, where p is an odd prime:
25 // c = x^2 (mod p), c = x^2 (mod p^k), q = p^(k-1)
26 // X = x^q * c^((p^k-2q+1)/2) (mod p^k)

```

4.1.11. Chinese Sieve

```

1 const ll N = 1000000;
2 // f, g, h multiplicative, h = f (dirichlet convolution) g
3 ll pre_g(ll n);
4 ll pre_h(ll n);
5 // preprocessed prefix sum of f
6 ll pre_f[N];
7 // prefix sum of multiplicative function f
8 ll solve_f(ll n) {
9     static unordered_map<ll, ll> m;
10    if (n < N) return pre_f[n];
11    if (m.count(n)) return m[n];
12    ll ans = pre_h(n);
13    for (ll l = 2, r; l <= n; l = r + 1) {
14        r = n / (n / l);
15        ans -= (pre_g(r) - pre_g(l - 1)) * djs_f(n / l);
16    }
17    return m[n] = ans;
18 }

```

4.1.12. Rational Number Binary Search

```

1 struct QQ {
2     ll p, q;
3     QQ go(QQ b, ll d) { return {p + b.p * d, q + b.q * d}; }
4 };
5 bool pred(QQ);
6 // returns smallest p/q in [lo, hi] such that
7 // pred(p/q) is true, and 0 <= p, q <= N
8 QQ frac_bs(ll N) {
9     QQ lo{0, 1}, hi{1, 0};
10    if (pred(lo)) return lo;
11    assert(pred(hi));
12    bool dir = 1, L = 1, H = 1;
13    for (; L || H; dir = !dir) {
14        ll len = 0, step = 1;
15        for (int t = 0; t < 2 && (t ? step /= 2 : step *= 2);)
16            if (QQ mid = hi.go(lo, len + step);
17                mid.p > N || mid.q > N || dir ^ pred(mid))
18                t++;
19        else len += step;
20        swap(lo, hi = hi.go(lo, len));
21        (dir ? L : H) = !len;
22    }
23    return dir ? hi : lo;
24 }

```

4.1.13. Farey Sequence

```

1 // returns (e/f), where (a/b, c/d, e/f) are
2 // three consecutive terms in the order n farey sequence
3 // to start, call next_farey(n, 0, 1, 1, n)
4 pll next_farey(ll n, ll a, ll b, ll c, ll d) {
5     ll p = (n + b) / d;
6     return pll(p * c - a, p * d - b);
7 }

```

4.2. Combinatorics

4.2.1. Matroid Intersection

This template assumes 2 weighted matroids of the same type, and that removing an element is much more expensive than checking if one can be added. **Remember to change the implementation details.**

The ground set is $0, 1, \dots, n-1$, where element i has weight $w[i]$. For the unweighted version, remove weights and change BF/SPFA to BFS.

```

1 constexpr int N = 100;
2 constexpr int INF = 1e9;
3
4 struct Matroid { // represents an independent set
5     Matroid(bitset<N>); // initialize from an independent set
6     bool can_add(int); // if adding will break independence
7     Matroid remove(int); // removing from the set
8 };
9
10 auto matroid_intersection(int n, const vector<int> &w) {
11     bitset<N> S;
12     for (int sz = 1; sz <= n; sz++) {
13         Matroid M1(S), M2(S);
14
15         vector<vector<pii>> e(n + 2);
16         for (int j = 0; j < n; j++)
17             if (!S[j]) {
18                 if (M1.can_add(j)) e[n].emplace_back(j, -w[j]);
19                 if (M2.can_add(j)) e[j].emplace_back(n + 1, 0);
20             }
21         for (int i = 0; i < n; i++)
22             if (S[i]) {
23                 Matroid T1 = M1.remove(i), T2 = M2.remove(i);
24                 for (int j = 0; j < n; j++)
25                     if (!S[j]) {
26                         if (T1.can_add(j)) e[i].emplace_back(j, -w[j]);
27                         if (T2.can_add(j)) e[j].emplace_back(i, w[i]);
28                     }
29             }
30
31         vector<pii> dis(n + 2, {INF, 0});
32         vector<int> prev(n + 2, -1);
33         dis[n] = {0, 0};
34         // change to SPFA for more speed, if necessary
35         bool upd = 1;
36         while (upd) {
37             upd = 0;
38             for (int u = 0; u < n + 2; u++)
39                 for (auto [v, c] : e[u]) {
40                     pii x(dis[u].first + c, dis[u].second + 1);
41                     if (x < dis[v]) dis[v] = x, prev[v] = u, upd = 1;
42                 }
43         }
44     }
45 }

```

```

43 }
45 if (dis[n + 1].first < INF)
46     for (int x = prev[n + 1]; x != n; x = prev[x])
47         S.flip(x);
48 else break;
49 // S is the max-weighted independent set with size sz
51 }
52 return S;
53 }

```

4.2.2. De Bruijn Sequence

```

1 int res[kN], aux[kN], a[kN], sz;
2 void Rec(int t, int p, int n, int k) {
3     if (t > n) {
4         if (n % p == 0)
5             for (int i = 1; i <= p; ++i) res[sz++] = aux[i];
6     } else {
7         aux[t] = aux[t - p];
8         Rec(t + 1, p, n, k);
9         for (aux[t] = aux[t - p] + 1; aux[t] < k; ++aux[t])
10             Rec(t + 1, t, n, k);
11     }
12 }
13 int DeBruijn(int k, int n) {
14     // return cyclic string of length k^n such that every
15     // string of length n using k character appears as a
16     // substring.
17     if (k == 1) return res[0] = 0, 1;
18     fill(aux, aux + k * n, 0);
19     return sz = 0, Rec(1, 1, n, k), sz;
20 }

```

4.2.3. Multinomial

```

1 // ways to permute v[i]
2 ll multinomial(vi &v) {
3     ll c = 1, m = v.empty() ? 1 : v[0];
4     for (int i = 1; i < v.size(); i++)
5         for (int j = 0; i < v[i]; j++) c = c * ++m / (j + 1);
6     return c;
7 }

```

4.3. Theorems

4.3.1. Kirchhoff's Theorem

Denote L be a $n \times n$ matrix as the Laplacian matrix of graph G , where $L_{ii} = d(i)$, $L_{ij} = -c$ where c is the number of edge (i, j) in G .

- The number of undirected spanning in G is $|\det(\tilde{L}_{11})|$.
- The number of directed spanning tree rooted at r in G is $|\det(\tilde{L}_{rr})|$.

4.3.2. Tutte's Matrix

Let D be a $n \times n$ matrix, where $d_{ij} = x_{ij}$ (x_{ij} is chosen uniformly at random) if $i < j$ and $(i, j) \in E$, otherwise $d_{ij} = -d_{ji}$. $\frac{\text{rank}(D)}{2}$ is the maximum matching on G .

4.3.3. Cayley's Formula

- Given a degree sequence d_1, d_2, \dots, d_n for each *labeled* vertices, there are

$$\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$$

spanning trees.

- Let $T_{n,k}$ be the number of *labeled* forests on n vertices with k components, such that vertex $1, 2, \dots, k$ belong to different components. Then $T_{n,k} = kn^{n-k-1}$.

4.3.4. Erdős–Gallai Theorem

A sequence of non-negative integers $d_1 \geq d_2 \geq \dots \geq d_n$ can be represented as the degree sequence of a finite simple graph on n vertices if and only if $d_1 + d_2 + \dots + d_n$ is even and

$$\sum_{i=1}^k d_i \leq k(k-1) + \sum_{i=k+1}^n \min(d_i, k)$$

holds for all $1 \leq k \leq n$.

4.3.5. Burnside's Lemma

Let X be a set and G be a group that acts on X . For $g \in G$, denote by X^g the elements fixed by g :

$$X^g = \{x \in X \mid gx \in X\}$$

Then

$$|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|.$$

5. Numeric

5.1. Barrett Reduction

```

1 using ull = unsigned long long;
2 using ul = __uint128_t;
3 // very fast calculation of a % m
4 struct reduction {
5     const ull m, d;
6     explicit reduction(ull m) : m(m), d(((ul)1 << 64) / m) {}
7     inline ull operator()(ull a) const {
8         ull q = (ull)((ul)d * a) >> 64;
9         return (a - q * m) >= m ? a - m : a;
10    }
11 };

```

5.2. Long Long Multiplication

```

1 using ull = unsigned long long;
2 using ll = long long;
3 using ld = long double;
4 // returns a * b % M where a, b < M < 2**63
5 ull mult(ull a, ull b, ull M) {
6     ll ret = a * b - M * ull(ld(a) * ld(b) / ld(M));
7     return ret + M * (ret < 0) - M * (ret >= (ll)M);
8 }

```

5.3. Fast Fourier Transform

```

1 template <typename T>
2 void fft_(int n, vector<T> &a, vector<T> &rt, bool inv) {
3     vector<int> br(n);
4     for (int i = 1; i < n; i++) {
5         br[i] = (i & 1) ? br[i - 1] + n / 2 : br[i / 2] / 2;
6         if (br[i] > i) swap(a[i], a[br[i]]);
7     }
8     for (int len = 2; len <= n; len *= 2)
9         for (int i = 0; i < n; i += len)
10             for (int j = 0; j < len / 2; j++) {
11                 int pos = n / len * (inv ? len - j : j);
12                 T u = a[i + j], v = a[i + j + len / 2] * rt[pos];
13                 a[i + j] = u + v, a[i + j + len / 2] = u - v;
14             }
15     if (T minv = T(1) / T(n); inv)
16         for (T &x : a) x *= minv;
17 }

```

Requires: Mod Struct

```

1 void ntt(vector<Mod> &a, bool inv, Mod primitive_root) {
2     int n = a.size();
3     Mod root = primitive_root ^ (MOD - 1) / n;
4     vector<Mod> rt(n + 1, 1);
5     for (int i = 0; i < n; i++) rt[i + 1] = rt[i] * root;
6     fft_(n, a, rt, inv);
7 }
8 void fft(vector<complex<double>> &a, bool inv) {
9     int n = a.size();
10    vector<complex<double>> rt(n + 1);
11    double arg = acos(-1) * 2 / n;
12    for (int i = 0; i <= n; i++)
13        rt[i] = {cos(arg * i), sin(arg * i)};
14    fft_(n, a, rt, inv);
15 }

```

5.4. Fast Walsh-Hadamard Transform

Requires: Mod Struct

```

1 void fwht(vector<Mod> &a, bool inv) {
2     int n = a.size();
3     for (int d = 1; d < n; d <= 1)
4         for (int m = 0; m < n; m++)
5             if (!(m & d)) {
6                 inv ? a[m] -= a[m | d] : a[m] += a[m | d]; // AND
7                 inv ? a[m | d] -= a[m] : a[m | d] += a[m]; // OR
8                 Mod x = a[m], y = a[m | d]; // XOR
9                 a[m] = x + y, a[m | d] = x - y; // XOR
10            }
11    if (Mod iv = Mod(1) / n; inv) // XOR
12        for (Mod &i : a) i *= iv; // XOR
13 }

```

5.5. Linear Recurrences

5.5.1. Berlekamp-Massey Algorithm

```

1 template <typename T>
  vector<T> berlekamp_massey(const vector<T> &s) {
3   int n = s.size(), l = 0, m = 1;
  vector<T> r(n), p(n);
5   r[0] = p[0] = 1;
  T b = 1, d = 0;
7   for (int i = 0; i < n; i++, m++, d = 0) {
      for (int j = 0; j <= l; j++) d += r[j] * s[i - j];
9      if ((d /= b) == 0) continue; // change if T is float
      auto t = r;
11     for (int j = m; j < n; j++) r[j] -= d * p[j - m];
      if (l * 2 <= i) l = i + 1 - l, b = d, m = 0, p = t;
13  }
15  return r.resize(l + 1), reverse(r.begin(), r.end()), r;
}

```

5.5.2. Linear Recurrence Calculation

```

1 template <typename T> struct lin_rec {
  using poly = vector<T>;
3   poly mul(poly a, poly b, poly m) {
      int n = m.size();
5      poly r(n);
      for (int i = n - 1; i >= 0; i--) {
9          r.insert(r.begin(), 0), r.pop_back();
          T c = r[n - 1] + a[n - 1] * b[i];
          // c /= m[n - 1]; if m is not monic
11         for (int j = 0; j < n; j++)
            r[j] += a[j] * b[i] - c * m[j];
13     }
    return r;
15  }
  poly pow(poly p, ll k, poly m) {
      poly r(m.size());
17      r[0] = 1;
      for (; k >= 1; p = mul(p, p, m))
19          if (k & 1) r = mul(r, p, m);
      return r;
21  }
  T calc(poly t, poly r, ll k) {
      int n = r.size();
23      poly p(n);
      p[1] = 1;
      poly q = pow(p, k, r);
25      T ans = 0;
      for (int i = 0; i < n; i++) ans += t[i] * q[i];
27      return ans;
29  }
31 };

```

5.6. Matrices

5.6.1. Determinant

Requires: Mod Struct

```

1 Mod det(vector<vector<Mod>> a) {
  int n = a.size();
3   Mod ans = 1;
  for (int i = 0; i < n; i++) {
5       int b = i;
       for (int j = i + 1; j < n; j++)
9           if (a[j][i] != 0) {
               b = j;
               break;
11          }
       if (i != b) swap(a[i], a[b]), ans = -ans;
       ans *= a[i][i];
13      if (ans == 0) return 0;
       for (int j = i + 1; j < n; j++) {
15          Mod v = a[j][i] / a[i][i];
          if (v != 0)
17              for (int k = i + 1; k < n; k++)
                  a[j][k] -= v * a[i][k];
19      }
21  }
  return ans;
}

```

```

1 double det(vector<vector<double>> a) {
  int n = a.size();
3   double ans = 1;
  for (int i = 0; i < n; i++) {
5       int b = i;
       for (int j = i + 1; j < n; j++)
9           if (fabs(a[j][i]) > fabs(a[b][i])) b = j;

```

```

9       if (i != b) swap(a[i], a[b]), ans = -ans;
       ans *= a[i][i];
       if (ans == 0) return 0;
11      for (int j = i + 1; j < n; j++) {
          double v = a[j][i] / a[i][i];
13          if (v != 0)
              for (int k = i + 1; k < n; k++)
15                  a[j][k] -= v * a[i][k];
17      }
19  }
  return ans;
}

```

5.6.2. Inverse

```

1 // Returns rank.
  // Result is stored in A unless singular (rank < n).
3 // For prime powers, repeatedly set
  // A^{-1} = A^{-1} (2I - A * A^{-1}) (mod p^k)
5 // where A^{-1} starts as the inverse of A mod p,
  // and k is doubled in each step.
7
9 int matInv(vector<vector<double>> &A) {
  int n = sz(A);
11  vi col(n);
  vector<vector<double>> tmp(n, vector<double>(n));
  rep(i, 0, n) tmp[i][i] = 1, col[i] = i;
13
15  rep(i, 0, n) {
      int r = i, c = i;
      rep(j, i, n)
17          rep(k, i, n) if (fabs(A[j][k]) > fabs(A[r][c])) r = j, c = k;
19      if (fabs(A[r][c]) < 1e-12) return i;
      A[i].swap(A[r]);
      tmp[i].swap(tmp[r]);
21      rep(j, 0, n) swap(A[j][i], A[j][c]),
          swap(tmp[j][i], tmp[j][c]);
      swap(col[i], col[c]);
23      double v = A[i][i];
      rep(j, i + 1, n) {
25          double f = A[j][i] / v;
          A[j][i] = 0;
          rep(k, i + 1, n) A[j][k] -= f * A[i][k];
          rep(k, 0, n) tmp[j][k] -= f * tmp[i][k];
31      }
      rep(j, i + 1, n) A[i][j] /= v;
      rep(j, 0, n) tmp[i][j] /= v;
33      A[i][i] = 1;
35  }
37
39  for (int i = n - 1; i > 0; --i) rep(j, 0, i) {
      double v = A[j][i];
      rep(k, 0, n) tmp[j][k] -= v * tmp[i][k];
41  }
43  rep(i, 0, n) rep(j, 0, n) A[col[i]][col[j]] = tmp[i][j];
  return n;
45 }
47
49 int matInv_mod(vector<vector<ll>> &A) {
  int n = sz(A);
  vi col(n);
  vector<vector<ll>> tmp(n, vector<ll>(n));
  rep(i, 0, n) tmp[i][i] = 1, col[i] = i;
51
53  rep(i, 0, n) {
      int r = i, c = i;
      rep(j, i, n) rep(k, i, n) if (A[j][k]) {
55          r = j;
          c = k;
          goto found;
57      }
      return i;
59  }
  found:
61  A[i].swap(A[r]);
  tmp[i].swap(tmp[r]);
63  rep(j, 0, n) swap(A[j][i], A[j][c]),
      swap(tmp[j][i], tmp[j][c]);
  swap(col[i], col[c]);
65  ll v = modpow(A[i][i], mod - 2);
  rep(j, i + 1, n) {
67      ll f = A[j][i] * v % mod;
      A[j][i] = 0;
      rep(k, i + 1, n) A[j][k] =
71          (A[j][k] - f * A[i][k]) % mod;
      rep(k, 0, n) tmp[j][k] =
73          (tmp[j][k] - f * tmp[i][k]) % mod;
75  }

```



```

77     rep(j, i + 1, n) A[i][j] = A[i][j] * v % mod;
78     rep(j, 0, n) tmp[i][j] = tmp[i][j] * v % mod;
79     A[i][i] = 1;
80 }
81 for (int i = n - 1; i > 0; --i) rep(j, 0, i) {
82     ll v = A[j][i];
83     rep(k, 0, n) tmp[j][k] =
84         (tmp[j][k] - v * tmp[i][k]) % mod;
85 }
86 rep(i, 0, n) rep(j, 0, n) A[col[i]][col[j]] =
87     tmp[i][j] % mod + (tmp[i][j] < 0 ? mod : 0);
88 return n;
89 }

```

5.6.3. Solve Linear Equation

```

1  typedef vector<double> vd;
2  const double eps = 1e-12;
3
4  // solves for x: A * x = b
5  int solveLinear(vector<vd> &A, vd &b, vd &x) {
6      int n = sz(A), m = sz(x), rank = 0, br, bc;
7      if (n) assert(sz(A[0]) == m);
8      vi col(m);
9      iota(all(col), 0);
10
11     rep(i, 0, n) {
12         double v, bv = 0;
13         rep(r, i, n) rep(c, i, m) if ((v = fabs(A[r][c])) > bv)
14             br = r, bc = c, bv = v;
15         if (bv <= eps) {
16             rep(j, i, n) if (fabs(b[j]) > eps) return -1;
17             break;
18         }
19         swap(A[i], A[br]);
20         swap(b[i], b[br]);
21         swap(col[i], col[bc]);
22         rep(j, 0, n) swap(A[j][i], A[j][bc]);
23         bv = 1 / A[i][i];
24         rep(j, i + 1, n) {
25             double fac = A[j][i] * bv;
26             b[j] -= fac * b[i];
27             rep(k, i + 1, m) A[j][k] -= fac * A[i][k];
28         }
29         rank++;
30     }
31
32     x.assign(m, 0);
33     for (int i = rank; i--;) {
34         b[i] /= A[i][i];
35         x[col[i]] = b[i];
36         rep(j, 0, i) b[j] -= A[j][i] * b[i];
37     }
38     return rank; // (multiple solutions if rank < m)
39 }

```

5.7. Polynomial Interpolation

```

1  // returns a, such that a[0]x^0 + a[1]x^1 + a[2]x^2 + ...
2  // passes through the given points
3  typedef vector<double> vd;
4  vd interpolate(vd x, vd y, int n) {
5      vd res(n), temp(n);
6      rep(k, 0, n - 1) rep(i, k + 1, n) y[i] =
7          (y[i] - y[k]) / (x[i] - x[k]);
8      double last = 0;
9      temp[0] = 1;
10     rep(k, 0, n) rep(i, 0, n) {
11         res[i] += y[k] * temp[i];
12         swap(last, temp[i]);
13         temp[i] -= last * x[k];
14     }
15     return res;
16 }

```

5.8. Simplex Algorithm

```

1  // Two-phase simplex algorithm for solving linear programs
2  // of the form
3  //
4  //      maximize      c^T x
5  //      subject to    Ax <= b
6  //                   x >= 0
7  //
8  // INPUT: A -- an m x n matrix
9  //         b -- an m-dimensional vector
10 //         c -- an n-dimensional vector

```

```

11 //      x -- a vector where the optimal solution will be
12 //      stored
13 //
14 // OUTPUT: value of the optimal solution (infinity if
15 // unbounded
16 //         above, nan if infeasible)
17 //
18 // To use this code, create an LPSolver object with A, b,
19 // and c as arguments. Then, call Solve(x).

```

```

21 typedef long double ld;
22 typedef vector<ld> vd;
23 typedef vector<vd> vvd;
24 typedef vector<int> vi;
25
26 const ld EPS = 1e-9;
27
28 struct LPSolver {
29     int m, n;
30     vi B, N;
31     vvd D;
32
33     LPSolver(const vvd &A, const vd &b, const vd &c)
34         : m(b.size()), n(c.size()), N(n + 1), B(m),
35           D(m + 2, vd(n + 2)) {
36         for (int i = 0; i < m; i++)
37             for (int j = 0; j < n; j++) D[i][j] = A[i][j];
38         for (int i = 0; i < m; i++) {
39             B[i] = n + i;
40             D[i][n] = -1;
41             D[i][n + 1] = b[i];
42         }
43         for (int j = 0; j < n; j++) {
44             N[j] = j;
45             D[m][j] = -c[j];
46         }
47         N[n] = -1;
48         D[m + 1][n] = 1;
49     }
50
51     void Pivot(int r, int s) {
52         double inv = 1.0 / D[r][s];
53         for (int i = 0; i < m + 2; i++)
54             if (i != r)
55                 for (int j = 0; j < n + 2; j++)
56                     if (j != s) D[i][j] -= D[r][j] * D[i][s] * inv;
57         for (int j = 0; j < n + 2; j++)
58             if (j != s) D[r][j] *= inv;
59         for (int i = 0; i < m + 2; i++)
60             if (i != r) D[i][s] *= -inv;
61         D[r][s] = inv;
62         swap(B[r], N[s]);
63     }
64
65     bool Simplex(int phase) {
66         int x = phase == 1 ? m + 1 : m;
67         while (true) {
68             int s = -1;
69             for (int j = 0; j <= n; j++) {
70                 if (phase == 2 && N[j] == -1) continue;
71                 if (s == -1 || D[x][j] < D[x][s] ||
72                     D[x][j] == D[x][s] && N[j] < N[s])
73                     s = j;
74             }
75             if (D[x][s] > -EPS) return true;
76             int r = -1;
77             for (int i = 0; i < m; i++) {
78                 if (D[i][s] < EPS) continue;
79                 if (r == -1 ||
80                     D[i][n + 1] / D[i][s] < D[r][n + 1] / D[r][s] ||
81                     (D[i][n + 1] / D[i][s]) ==
82                     (D[r][n + 1] / D[r][s]) &&
83                     B[i] < B[r])
84                     r = i;
85             }
86             if (r == -1) return false;
87             Pivot(r, s);
88         }
89     }
90 }

```

```

91 ld Solve(vd &x) {
92     int r = 0;
93     for (int i = 1; i < m; i++)
94         if (D[i][n + 1] < D[r][n + 1]) r = i;
95     if (D[r][n + 1] < -EPS) {
96         Pivot(r, n);
97         if (!Simplex(1) || D[m + 1][n + 1] < -EPS)
98             return -numeric_limits<ld>::infinity();
99         for (int i = 0; i < m; i++)
100             if (B[i] == -1) {

```

```

101     int s = -1;
102     for (int j = 0; j <= n; j++)
103         if (s == -1 || D[i][j] < D[i][s] ||
104             D[i][j] == D[i][s] && N[j] < N[s])
105             s = j;
106     Pivot(i, s);
107 }
108 if (!Simplex(2)) return numeric_limits<ld>::infinity();
109 x = vd(n);
110 for (int i = 0; i < m; i++)
111     if (B[i] < n) x[B[i]] = D[i][n + 1];
112 return D[m][n + 1];
113 }
114 };
115
116 int main() {
117     const int m = 4;
118     const int n = 3;
119     ld _A[m][n] = {
120         {6, -1, 0}, {-1, -5, 0}, {1, 5, 1}, {-1, -5, -1}};
121     ld _b[m] = {10, -4, 5, -5};
122     ld _c[n] = {1, -1, 0};
123
124     vvd A(m);
125     vd b(_b, _b + m);
126     vd c(_c, _c + n);
127     for (int i = 0; i < m; i++) A[i] = vd(_A[i], _A[i] + n);
128
129     LPSolver solver(A, b, c);
130     vd x;
131     ld value = solver.Solve(x);
132
133     cerr << "VALUE: " << value << endl; // VALUE: 1.29032
134     cerr << "SOLUTION:"; // SOLUTION: 1.74194 0.451613 1
135     for (size_t i = 0; i < x.size(); i++) cerr << " " << x[i];
136     cerr << endl;
137     return 0;
138 }

```

6. Geometry

6.1. Point

```

1 template <typename T> struct P {
2     T x, y;
3     P(T x = 0, T y = 0) : x(x), y(y) {}
4     bool operator<(const P &p) const {
5         return tie(x, y) < tie(p.x, p.y);
6     }
7     bool operator==(const P &p) const {
8         return tie(x, y) == tie(p.x, p.y);
9     }
10    P operator-() const { return {-x, -y}; }
11    P operator+(P p) const { return {x + p.x, y + p.y}; }
12    P operator-(P p) const { return {x - p.x, y - p.y}; }
13    P operator*(T d) const { return {x * d, y * d}; }
14    P operator/(T d) const { return {x / d, y / d}; }
15    T dist2() const { return x * x + y * y; }
16    double len() const { return sqrt(dist2()); }
17    P unit() const { return *this / len(); }
18    friend T dot(P a, P b) { return a.x * b.x + a.y * b.y; }
19    friend T cross(P a, P b) { return a.x * b.y - a.y * b.x; }
20    friend T cross(P a, P b, P o) {
21        return cross(a - o, b - o);
22    }
23 };
24 using pt = P<ll>;

```

6.1.1. Quarternion

```

1 constexpr double PI = 3.141592653589793;
2 constexpr double EPS = 1e-7;
3 struct Q {
4     using T = double;
5     T x, y, z, r;
6     Q(T r = 0) : x(0), y(0), z(0), r(r) {}
7     Q(T x, T y, T z, T r = 0) : x(x), y(y), z(z), r(r) {}
8     friend bool operator==(const Q &a, const Q &b) {
9         return (a - b).abs2() <= EPS;
10    }
11    friend bool operator!=(const Q &a, const Q &b) {
12        return !(a == b);
13    }
14    Q operator-() { return Q(-x, -y, -z, -r); }
15    Q operator+(const Q &b) const {
16        return Q(x + b.x, y + b.y, z + b.z, r + b.r);
17    }

```

```

19    Q operator-(const Q &b) const {
20        return Q(x - b.x, y - b.y, z - b.z, r - b.r);
21    }
22    Q operator*(const T &t) const {
23        return Q(x * t, y * t, z * t, r * t);
24    }
25    Q operator*(const Q &b) const {
26        return Q(r * b.x + x * b.r + y * b.z - z * b.y,
27                r * b.y - x * b.z + y * b.r + z * b.x,
28                r * b.z + x * b.y - y * b.x + z * b.r,
29                r * b.r - x * b.x - y * b.y - z * b.z);
30    }
31    Q operator/(const Q &b) const { return *this * b.inv(); }
32    T abs2() const { return r * r + x * x + y * y + z * z; }
33    T len() const { return sqrt(abs2()); }
34    Q conj() const { return Q(-x, -y, -z, r); }
35    Q unit() const { return *this * (1.0 / len()); }
36    Q inv() const { return conj() * (1.0 / abs2()); }
37    friend T dot(Q a, Q b) {
38        return a.x * b.x + a.y * b.y + a.z * b.z;
39    }
40    friend Q cross(Q a, Q b) {
41        return Q(a.y * b.z - a.z * b.y, a.z * b.x - a.x * b.z,
42                a.x * b.y - a.y * b.x);
43    }
44    friend Q rotation_around(Q axis, T angle) {
45        return axis.unit() * sin(angle / 2) + cos(angle / 2);
46    }
47    Q rotated_around(Q axis, T angle) {
48        Q u = rotation_around(axis, angle);
49        return u * *this / u;
50    }
51    friend Q rotation_between(Q a, Q b) {
52        a = a.unit(), b = b.unit();
53        if (a == -b) {
54            // degenerate case
55            Q ortho = abs(a.y) > EPS ? cross(a, Q(1, 0, 0))
56                                     : cross(a, Q(0, 1, 0));
57            return rotation_around(ortho, PI);
58        }
59        return (a * (a + b)).conj();
60    }
61 };

```

6.1.2. Spherical Coordinates

```

1 struct car_p {
2     double x, y, z;
3 };
4 struct sph_p {
5     double r, theta, phi;
6 };
7 sph_p conv(car_p p) {
8     double r = sqrt(p.x * p.x + p.y * p.y + p.z * p.z);
9     double theta = asin(p.y / r);
10    double phi = atan2(p.y, p.x);
11    return {r, theta, phi};
12 }
13 car_p conv(sph_p p) {
14     double x = p.r * cos(p.theta) * sin(p.phi);
15     double y = p.r * cos(p.theta) * cos(p.phi);
16     double z = p.r * sin(p.theta);
17     return {x, y, z};
18 }

```

6.2. Convex Hull

```

1 // returns a convex hull in counterclockwise order
2 // for a non-strict one, change cross >= to >
3 vector<pt> convex_hull(vector<pt> p) {
4     sort(ALL(p));
5     if (p[0] == p.back()) return {p[0]};
6     int n = p.size(), t = 0;
7     vector<pt> h(n + 1);
8     for (int _ = 2, s = 0; _--; s = --t, reverse(ALL(p)))
9         for (pt i : p) {
10             while (t > s + 1 && cross(i, h[t - 1], h[t - 2]) >= 0)
11                 t--;
12             h[t++] = i;
13         }
14     return h.resize(t), h;
15 }

```

6.2.1. 3D Hull

```

1 // source: KACTL
2
3 typedef Point3D<double> P3;

```

```

5 struct PR {
6     void ins(int x) { (a == -1 ? a : b) = x; }
7     void rem(int x) { (a == x ? a : b) = -1; }
8     int cnt() { return (a != -1) + (b != -1); }
9     int a, b;
10 };
11
12 struct F {
13     P3 q;
14     int a, b, c;
15 };
16
17 vector<F> hull3d(const vector<P3> &A) {
18     assert(sz(A) >= 4);
19     vector<vector<PR>> E(sz(A), vector<PR>(sz(A), {-1, -1}));
20     #define E(x, y) E[f.x][f.y]
21     vector<F> FS;
22     auto mf = [&](int i, int j, int k, int l) {
23         P3 q = (A[j] - A[i]).cross((A[k] - A[i]));
24         if (q.dot(A[l]) > q.dot(A[i])) q = q * -1;
25         F f{q, i, j, k};
26         E(a, b).ins(k);
27         E(a, c).ins(j);
28         E(b, c).ins(i);
29         FS.push_back(f);
30     };
31     rep(i, 0, 4) rep(j, i + 1, 4) rep(k, j + 1, 4)
32     mf(i, j, k, 6 - i - j - k);
33
34     rep(i, 4, sz(A)) {
35         rep(j, 0, sz(FS)) {
36             F f = FS[j];
37             if (f.q.dot(A[i]) > f.q.dot(A[f.a])) {
38                 E(a, b).rem(f.c);
39                 E(a, c).rem(f.b);
40                 E(b, c).rem(f.a);
41                 swap(FS[j--], FS.back());
42                 FS.pop_back();
43             }
44         }
45         int nw = sz(FS);
46         rep(j, 0, nw) {
47             F f = FS[j];
48             #define C(a, b, c)
49             if (E(a, b).cnt() != 2) mf(f.a, f.b, i, f.c);
50             C(a, b, c);
51             C(a, c, b);
52             C(b, c, a);
53         }
54         for (F &it : FS)
55             if ((A[it.b] - A[it.a])
56                 .cross(A[it.c] - A[it.a])
57                 .dot(it.q) <= 0)
58                 swap(it.c, it.b);
59         return FS;
60     };
61 };

```

6.3. Angular Sort

```

1 auto angle_cmp = [](const pt &a, const pt &b) {
2     auto btm = [](const pt &a) {
3         return a.y < 0 || (a.y == 0 && a.x < 0);
4     };
5     return make_tuple(btm(a), a.y * b.x, abs2(a)) <
6            make_tuple(btm(b), a.x * b.y, abs2(b));
7 };
8 void angular_sort(vector<pt> &p) {
9     sort(p.begin(), p.end(), angle_cmp);
10 }

```

6.4. Convex Polygon Minkowski Sum

```

1 // O(n) convex polygon minkowski sum
2 // must be sorted and counterclockwise
3 vector<pt> minkowski_sum(vector<pt> p, vector<pt> q) {
4     auto diff = [](vector<pt> &c) {
5         auto rcmp = [](pt a, pt b) {
6             return pt{a.y, a.x} < pt{b.y, b.x};
7         };
8         rotate(c.begin(), min_element(ALL(c), rcmp), c.end());
9         c.push_back(c[0]);
10        vector<pt> ret;
11        for (int i = 1; i < c.size(); i++)
12            ret.push_back(c[i] - c[i - 1]);
13        return ret;
14    };
15    auto dp = diff(p), dq = diff(q);
16    pt cur = p[0] + q[0];
17    vector<pt> d(dp.size() + dq.size(), ret = {cur};

```

```

19 // include angle_cmp from angular-sort.cpp
20 merge(ALL(dp), ALL(dq), d.begin(), angle_cmp);
21 // optional: make ret strictly convex (UB if degenerate)
22 int now = 0;
23 for (int i = 1; i < d.size(); i++) {
24     if (cross(d[i], d[now]) == 0) d[now] = d[now] + d[i];
25     else d[++now] = d[i];
26 }
27 d.resize(now + 1);
28 // end optional part
29 for (pt v : d) ret.push_back(cur = cur + v);
30 return ret.pop_back(), ret;
31 }

```

6.5. Point In Polygon

```

1 bool on_segment(pt a, pt b, pt p) {
2     return cross(a, b, p) == 0 && dot((p - a), (p - b)) <= 0;
3 }
4 // p can be any polygon, but this is O(n)
5 bool inside(const vector<pt> &p, pt a) {
6     int cnt = 0, n = p.size();
7     for (int i = 0; i < n; i++) {
8         pt l = p[i], r = p[(i + 1) % n];
9         // change to return 0; for strict version
10        if (on_segment(l, r, a)) return 1;
11        cnt ^= ((a.y < l.y) - (a.y < r.y)) * cross(l, r, a) > 0;
12    }
13    return cnt;
14 }

```

6.5.1. Convex Version

```

1 // no preprocessing version
2 // p must be a strict convex hull, counterclockwise
3 // if point is inside or on border
4 bool is_inside(const vector<pt> &c, pt p) {
5     int n = c.size(), l = 1, r = n - 1;
6     if (cross(c[0], c[1], p) < 0) return false;
7     if (cross(c[n - 1], c[0], p) < 0) return false;
8     while (l < r - 1) {
9         int m = (l + r) / 2;
10        T a = cross(c[0], c[m], p);
11        if (a > 0) l = m;
12        else if (a < 0) r = m;
13        else return dot(c[0] - p, c[m] - p) <= 0;
14    }
15    if (l == r) return dot(c[0] - p, c[l] - p) <= 0;
16    else return cross(c[l], c[r], p) >= 0;
17 }
18
19 // with preprocessing version
20 vector<pt> vecs;
21 pt center;
22 // p must be a strict convex hull, counterclockwise
23 // BEWARE OF OVERFLOWS!!
24 void preprocess(vector<pt> p) {
25     for (auto &v : p) v = v * 3;
26     center = p[0] + p[1] + p[2];
27     center.x /= 3, center.y /= 3;
28     for (auto &v : p) v = v - center;
29     vecs = (angular_sort(p), p);
30 }
31 bool intersect_strict(pt a, pt b, pt c, pt d) {
32     if (cross(b, c, a) * cross(b, d, a) > 0) return false;
33     if (cross(d, a, c) * cross(d, b, c) >= 0) return false;
34     return true;
35 }
36 // if point is inside or on border
37 bool query(pt p) {
38     p = p * 3 - center;
39     auto pr = upper_bound(ALL(vecs), p, angle_cmp);
40     if (pr == vecs.end()) pr = vecs.begin();
41     auto pl = (pr == vecs.begin()) ? vecs.back() : *(pr - 1);
42     return !intersect_strict({0, 0}, p, pl, *pr);
43 }

```

6.5.2. Offline Multiple Points Version

Requires: GNU PBDS, Point

```

1 using Double = __float128;
2 using Point = pt<Double, Double>;
3
4 int n, m;
5 vector<Point> poly;
6 vector<Point> query;
7 vector<int> ans;
8
9 struct Segment {

```

```

11 Point a, b;
12 int id;
13 };
14 vector<Segment> segs;
15
16 Double Xnow;
17 inline Double get_y(const Segment &u, Double xnow = Xnow) {
18     const Point &a = u.a;
19     const Point &b = u.b;
20     return (a.y * (b.x - xnow) + b.y * (xnow - a.x)) /
21           (b.x - a.x);
22 }
23 bool operator<(Segment u, Segment v) {
24     Double yu = get_y(u);
25     Double yv = get_y(v);
26     if (yu != yv) return yu < yv;
27     return u.id < v.id;
28 }
29 ordered_map<Segment> st;
30
31 struct Event {
32     int type; // +1 insert seg, -1 remove seg, 0 query
33     Double x, y;
34     int id;
35 };
36 bool operator<(Event a, Event b) {
37     if (a.x != b.x) return a.x < b.x;
38     if (a.type != b.type) return a.type < b.type;
39     return a.y < b.y;
40 }
41 vector<Event> events;
42
43 void solve() {
44     set<Double> xs;
45     set<Point> ps;
46     for (int i = 0; i < n; i++) {
47         xs.insert(poly[i].x);
48         ps.insert(poly[i]);
49     }
50     for (int i = 0; i < n; i++) {
51         Segment s{poly[i], poly[(i + 1) % n], i};
52         if (s.a.x > s.b.x ||
53             (s.a.x == s.b.x && s.a.y > s.b.y)) {
54             swap(s.a, s.b);
55         }
56         segs.push_back(s);
57
58         if (s.a.x != s.b.x) {
59             events.push_back({+1, s.a.x + 0.2, s.a.y, i});
60             events.push_back({-1, s.b.x - 0.2, s.b.y, i});
61         }
62     }
63     for (int i = 0; i < m; i++) {
64         events.push_back({0, query[i].x, query[i].y, i});
65     }
66     sort(events.begin(), events.end());
67     int cnt = 0;
68     for (Event e : events) {
69         int i = e.id;
70         Xnow = e.x;
71         if (e.type == 0) {
72             Double x = e.x;
73             Double y = e.y;
74             Segment tmp = {{x - 1, y}, {x + 1, y}, -1};
75             auto it = st.lower_bound(tmp);
76
77             if (ps.count(query[i]) > 0) {
78                 ans[i] = 0;
79             } else if (xs.count(x) > 0) {
80                 ans[i] = -2;
81             } else if (it != st.end() &&
82                 get_y(*it) == get_y(tmp)) {
83                 ans[i] = 0;
84             } else if (it != st.begin() &&
85                 get_y(*prev(it)) == get_y(tmp)) {
86                 ans[i] = 0;
87             } else {
88                 int rk = st.order_of_key(tmp);
89                 if (rk % 2 == 1) {
90                     ans[i] = 1;
91                 } else {
92                     ans[i] = -1;
93                 }
94             }
95         } else if (e.type == 1) {
96             st.insert(segs[i]);
97             assert((int)st.size() == ++cnt);
98         } else if (e.type == -1) {
99             st.erase(segs[i]);
100             assert((int)st.size() == --cnt);

```

```

101     }
102 }

```

6.6. Closest Pair

```

1 vector<pll> p; // sort by x first!
2 bool cmpy(const pll &a, const pll &b) const {
3     return a.y < b.y;
4 }
5 ll sq(ll x) { return x * x; }
6 // returns (minimum dist)^2 in [l, r)
7 ll solve(int l, int r) {
8     if (r - l <= 1) return 1e18;
9     int m = (l + r) / 2;
10    ll mid = p[m].x, d = min(solve(l, m), solve(m, r));
11    auto pb = p.begin();
12    inplace_merge(pb + l, pb + m, pb + r, cmpy);
13    vector<pll> s;
14    for (int i = l; i < r; i++)
15        if (sq(p[i].x - mid) < d) s.push_back(p[i]);
16    for (int i = 0; i < s.size(); i++)
17        for (int j = i + 1;
18            j < s.size() && sq(s[j].y - s[i].y) < d; j++)
19            d = min(d, dis(s[i], s[j]));
20    return d;
21 }

```

6.7. Minimum Enclosing Circle

```

1 typedef Point<double> P;
2 double ccRadius(const P &A, const P &B, const P &C) {
3     return (B - A).dist() * (C - B).dist() * (A - C).dist() /
4           abs((B - A).cross(C - A)) / 2;
5 }
6 P ccCenter(const P &A, const P &B, const P &C) {
7     P b = C - A, c = B - A;
8     return A + (b * c.dist2() - c * b.dist2()).perp() /
9           b.cross(c) / 2;
10 }
11 pair<P, double> mec(vector<P> ps) {
12     shuffle(all(ps), mt19937(time(0)));
13     P o = ps[0];
14     double r = 0, EPS = 1 + 1e-8;
15     rep(i, 0, sz(ps)) if ((o - ps[i]).dist() > r * EPS) {
16         o = ps[i], r = 0;
17     }
18     rep(j, 0, i) if ((o - ps[j]).dist() > r * EPS) {
19         o = (ps[i] + ps[j]) / 2;
20         r = (o - ps[i]).dist();
21     }
22     rep(k, 0, j) if ((o - ps[k]).dist() > r * EPS) {
23         o = ccCenter(ps[i], ps[j], ps[k]);
24         r = (o - ps[i]).dist();
25     }
26     return {o, r};
27 }

```

6.8. Delaunay Triangulation

```

1 // source: KACTL
2
3 typedef Point<ll> P;
4 typedef struct Quad *Q;
5 typedef __int128_t lll; // (can be ll if coords are < 2e4)
6 P arb(LLONG_MAX, LLONG_MAX); // not equal to any other point
7
8 struct Quad {
9     bool mark;
10    Q o, rot;
11    P p;
12    P F() { return r()->p; }
13    Q r() { return rot->rot; }
14    Q prev() { return rot->o->rot; }
15    Q next() { return r()->prev(); }
16 };
17
18 bool circ(P p, P a, P b, P c) { // is p in the circumcircle?
19     lll p2 = p.dist2(), A = a.dist2() - p2,
20         B = b.dist2() - p2, C = c.dist2() - p2;
21     return p.cross(a, b) * C + p.cross(b, c) * A +
22           p.cross(c, a) * B >
23           0;
24 }
25 Q makeEdge(P orig, P dest) {
26     Q q[] = {new Quad{0, 0, 0, orig}, new Quad{0, 0, 0, arb},
27             new Quad{0, 0, 0, dest}, new Quad{0, 0, 0, arb}};
28     rep(i, 0, 4) q[i]->o = q[(i & 3)],
29             q[i]->rot = q[(i + 1) & 3];
30     return *q;

```

```

31 }
32 void splice(Q a, Q b) {
33     swap(a->o->rot->o, b->o->rot->o);
34     swap(a->o, b->o);
35 }
36 Q connect(Q a, Q b) {
37     Q q = makeEdge(a->F(), b->p);
38     splice(q, a->next());
39     splice(q->r(), b);
40     return q;
41 }
42
43 pair<Q, Q> rec(const vector<P> &s) {
44     if (sz(s) <= 3) {
45         Q a = makeEdge(s[0], s[1]),
46           b = makeEdge(s[1], s.back());
47         if (sz(s) == 2) return {a, a->r()};
48         splice(a->r(), b);
49         auto side = s[0].cross(s[1], s[2]);
50         Q c = side ? connect(b, a) : 0;
51         return {side < 0 ? c->r() : a, side < 0 ? c : b->r()};
52     }
53
54     #define H(e) e->F(), e->p
55     #define valid(e) (e->F().cross(H(base)) > 0)
56     Q A, B, ra, rb;
57     int half = sz(s) / 2;
58     tie(ra, A) = rec({all(s) - half});
59     tie(B, rb) = rec({sz(s) - half + all(s)});
60     while ((B->p.cross(H(A)) < 0 && (A = A->next())) ||
61           (A->p.cross(H(B)) > 0 && (B = B->r()->o)))
62         ;
63     Q base = connect(B->r(), A);
64     if (A->p == ra->p) ra = base->r();
65     if (B->p == rb->p) rb = base;
66
67     #define DEL(e, init, dir)
68     Q e = init->dir;
69     if (valid(e))
70         while (circ(e->dir->F(), H(base), e->F())) {
71             Q t = e->dir;
72             splice(e, e->prev());
73             splice(e->r(), e->r()->prev());
74             e = t;
75         }
76     for (;;) {
77         DEL(LC, base->r(), o);
78         DEL(RC, base, prev());
79         if (!valid(LC) && !valid(RC)) break;
80         if (!valid(LC) || (valid(RC) && circ(H(RC), H(LC))))
81             base = connect(RC, base->r());
82         else base = connect(base->r(), LC->r());
83     }
84     return {ra, rb};
85 }
86
87 // returns [A_0, B_0, C_0, A_1, B_1, ...]
88 // where A_i, B_i, C_i are counter-clockwise triangles
89 vector<P> triangulate(vector<P> pts) {
90     sort(all(pts));
91     assert(unique(all(pts)) == pts.end());
92     if (sz(pts) < 2) return {};
93     Q e = rec(pts).first;
94     vector<Q> q = {e};
95     int qi = 0;
96     while (e->o->F().cross(e->F(), e->p) < 0) e = e->o;
97     #define ADD
98     {
99         Q c = e;
100         do {
101             c->mark = 1;
102             pts.push_back(c->p);
103             q.push_back(c->r());
104             c = c->next();
105         } while (c != e);
106     }
107     ADD;
108     pts.clear();
109     while (qi < sz(q))
110         if (!(e = q[qi++])->mark) ADD;
111     return pts;
112 }

```

6.9. Half Plane Intersection

```

1 struct Line {
2     Point P;
3     Vector v;
4     bool operator<(const Line &b) const {
5         return atan2(v.y, v.x) < atan2(b.v.y, b.v.x);

```

```

6     };
7 };
8 bool OnLeft(const Line &L, const Point &p) {
9     return Cross(L.v, p - L.P) > 0;
10 }
11 Point GetIntersection(Line a, Line b) {
12     Vector u = a.P - b.P;
13     Double t = Cross(b.v, u) / Cross(a.v, b.v);
14     return a.P + a.v * t;
15 }
16 int HalfplaneIntersection(Line *L, int n, Point *poly) {
17     sort(L, L + n);
18
19     int first, last;
20     Point *p = new Point[n];
21     Line *q = new Line[n];
22     q[first = last = 0] = L[0];
23     for (int i = 1; i < n; i++) {
24         while (first < last && !OnLeft(L[i], p[last - 1]))
25             last--;
26         while (first < last && !OnLeft(L[i], p[first])) first++;
27         q[++last] = L[i];
28         if (fabs(Cross(q[last].v, q[last - 1].v)) < EPS) {
29             last--;
30             if (OnLeft(q[last], L[i].P)) q[last] = L[i];
31         }
32         if (first < last)
33             p[last - 1] = GetIntersection(q[last - 1], q[last]);
34     }
35     while (first < last && !OnLeft(q[first], p[last - 1]))
36         last--;
37     if (last - first <= 1) return 0;
38     p[last] = GetIntersection(q[last], q[first]);
39
40     int m = 0;
41     for (int i = first; i <= last; i++) poly[m++] = p[i];
42     return m;
43 }

```

7. Strings

7.1. Aho-Corasick Automaton

```

1 struct Aho_Corasick {
2     static const int maxc = 26, maxn = 4e5;
3     struct NODES {
4         int Next[maxc], fail, ans;
5     };
6     NODES T[maxn];
7     int top, qtop, q[maxn];
8     int get_node(const int &fail) {
9         fill_n(T[top].Next, maxc, 0);
10        T[top].fail = fail;
11        T[top].ans = 0;
12        return top++;
13    }
14    int insert(const string &s) {
15        int ptr = 1;
16        for (char c : s) { // change char id
17            c -= 'a';
18            if (!T[ptr].Next[c]) T[ptr].Next[c] = get_node(ptr);
19            ptr = T[ptr].Next[c];
20        }
21        return ptr;
22    } // return ans_last_place
23    void build_fail(int ptr) {
24        int tmp;
25        for (int i = 0; i < maxc; i++)
26            if (T[ptr].Next[i]) {
27                tmp = T[ptr].fail;
28                while (tmp != 1 && !T[tmp].Next[i])
29                    tmp = T[tmp].fail;
30                if (T[tmp].Next[i] != T[ptr].Next[i])
31                    if (T[tmp].Next[i]) tmp = T[tmp].Next[i];
32                T[T[ptr].Next[i]].fail = tmp;
33                q[qtop++] = T[ptr].Next[i];
34            }
35    }
36    void AC_auto(const string &s) {
37        int ptr = 1;
38        for (char c : s) {
39            while (ptr != 1 && !T[ptr].Next[c]) ptr = T[ptr].fail;
40            if (T[ptr].Next[c]) {
41                ptr = T[ptr].Next[c];
42                T[ptr].ans++;
43            }
44        }
45    }
46    void Solve(string &s) {
47        for (char &c : s) // change char id

```



```

    c -= 'a';
49   for (int i = 0; i < qtop; i++) build_fail(q[i]);
    AC_auto(s);
51   for (int i = qtop - 1; i > -1; i--)
        T[T[q[i]].fail].ans += T[q[i]].ans;
53   }
    void reset() {
55       qtop = top = q[0] = 1;
        get_node(1);
57   } AC;
59   // usage example
    string s, S;
61   int n, t, ans_place[50000];
    int main() {
63       Tie cin >> t;
        while (t--) {
65           AC.reset();
            cin >> S >> n;
67           for (int i = 0; i < n; i++) {
                cin >> s;
69                 ans_place[i] = AC.insert(s);
            }
71         AC.Solve(S);
            for (int i = 0; i < n; i++)
73                 cout << AC.T[ans_place[i]].ans << '\n';
        }
75   }

```

7.2. Suffix Array

```

1 // sa[i]: starting index of suffix at rank i
  // 0-indexed, sa[0] = n (empty string)
3 // lcp[i]: lcp of sa[i] and sa[i - 1], lcp[0] = 0
    struct SuffixArray {
5         vector<int> sa, lcp;
        SuffixArray(string &s,
7             int lim = 256) { // or basic_string<int>
            int n = sz(s) + 1, k = 0, a, b;
9             vector<int> x(all(s) + 1), y(n), ws(max(n, lim)),
                rank(n);
11            sa = lcp = y, iota(all(sa), 0);
            for (int j = 0, p = 0; p < n;
13                j = max(1, j * 2), lim = p) {
                p = j, iota(all(y), n - j);
15                for (int i = 0; i < n; i++)
                    if (sa[i] >= j) y[p++] = sa[i] - j;
17                fill(all(ws), 0);
                for (int i = 0; i < n; i++) ws[x[i]]++;
19                for (int i = 1; i < lim; i++) ws[i] += ws[i - 1];
                for (int i = n; i--;) sa[--ws[x[i]]] = i;
21                swap(x, y), p = 1, x[sa[0]] = 0;
                for (int i = 1; i < n; i++)
23                    a = sa[i - 1], b = sa[i],
                    x[b] = (y[a] == y[b] && y[a + j] == y[b + j])
                        ? p - 1 : p++;
25            }
27            for (int i = 1; i < n; i++) rank[sa[i]] = i;
            for (int i = 0, j; i < n - 1; lcp[rank[i+]] = k)
31                for (k && k--; j = sa[rank[i] - 1];
                    s[i + k] == s[j + k]; k++)
33                    ;
35        }
    };

```

7.3. Suffix Tree

```

1 struct SAM {
    static const int maxc = 26; // char range
3     static const int maxn = 10010; // string len
    struct Node {
5         Node *green, *edge[maxc];
            int max_len, in, times;
7     } * root, *last, reg[maxn * 2];
    int top;
9     Node *get_node(int _max) {
        Node *re = &reg[top++];
11        re->in = 0, re->times = 1;
        re->max_len = _max, re->green = 0;
13        for (int i = 0; i < maxc; i++) re->edge[i] = 0;
        return re;
15    }
    void insert(const char c) { // c in range [0, maxc)
17        Node *p = last;
        last = get_node(p->max_len + 1);
19        while (p && !p->edge[c])
            p->edge[c] = last, p = p->green;
21        if (!p) last->green = root;

```

```

    else {
23        Node *pot_green = p->edge[c];
        if ((pot_green->max_len) == (p->max_len + 1))
25            last->green = pot_green;
        else {
27            Node *wish = get_node(p->max_len + 1);
            wish->times = 0;
29            while (p && p->edge[c] == pot_green)
                p->edge[c] = wish, p = p->green;
31            for (int i = 0; i < maxc; i++)
                wish->edge[i] = pot_green->edge[i];
33            wish->green = pot_green->green;
            pot_green->green = wish;
35            last->green = wish;
        }
37    }
    }
39    Node *q[maxn * 2];
    int ql, qr;
41    void get_times(Node *p) {
        ql = 0, qr = -1, reg[0].in = 1;
43        for (int i = 1; i < top; i++) reg[i].green->in++;
        for (int i = 0; i < top; i++)
45            if (!reg[i].in) q[++qr] = &reg[i];
        while (ql <= qr) {
47            q[ql]->green->times += q[ql]->times;
            if (!(--q[ql]->green->in)) q[++qr] = q[ql]->green;
49            ql++;
        }
51    }
    void build(const string &s) {
53        top = 0;
        root = last = get_node(0);
55        for (char c : s) insert(c - 'a'); // change char id
        get_times(root);
57    }
    // call build before solve
59    int solve(const string &s) {
        Node *p = root;
61        for (char c : s)
            if (!(p = p->edge[c - 'a'])) // change char id
63                return 0;
        return p->times;
65    }
};

```

7.4. Cocke-Younger-Kasami Algorithm

```

1 struct rule {
    // s -> xy
3     // if y == -1, then s -> x (unit rule)
    int s, x, y, cost;
5 };
    int state;
7 // state (id) for each letter (variable)
    // lowercase letters are terminal symbols
9 map<char, int> rules;
    vector<rule> cnf;
11 void init() {
        state = 0;
13        rules.clear();
        cnf.clear();
15    }
    // convert a cfg rule to cnf (but with unit rules) and add
17 // it
    void add_to_cnf(char s, const string &p, int cost) {
19        if (!rules.count(s)) rules[s] = state++;
        for (char c : p)
21            if (!rules.count(c)) rules[c] = state++;
        if (p.size() == 1) {
23            cnf.push_back({rules[s], rules[p[0]], -1, cost});
        } else {
25            // length >= 3 -> split
            int left = rules[s];
27            int sz = p.size();
            for (int i = 0; i < sz - 2; i++) {
29                cnf.push_back({left, rules[p[i]], state, 0});
                left = state++;
31            }
            cnf.push_back(
33                {left, rules[p[sz - 2]], rules[p[sz - 1]], cost});
        }
35    }
37    constexpr int MAXN = 55;
    vector<long long> dp[MAXN][MAXN];
39    // unit rules with negative costs can cause negative cycles
    vector<bool> neg_INF[MAXN][MAXN];
41    void relax(int l, int r, rule c, long long cost,

```

```

43     bool neg_c = 0) {
44     if (!neg_INF[l][r][c.s] &&
45         (neg_INF[l][r][c.x] || cost < dp[l][r][c.s])) {
46         if (neg_c || neg_INF[l][r][c.x]) {
47             dp[l][r][c.s] = 0;
48             neg_INF[l][r][c.s] = true;
49         } else {
50             dp[l][r][c.s] = cost;
51         }
52     }
53 }
54 void bellman(int l, int r, int n) {
55     for (int k = 1; k <= state; k++)
56         for (rule c : cnf)
57             if (c.y == -1)
58                 relax(l, r, c, dp[l][r][c.x] + c.cost, k == n);
59 }
60 void cyk(const string &s) {
61     vector<int> tok;
62     for (char c : s) tok.push_back(rules[c]);
63     for (int i = 0; i < tok.size(); i++) {
64         for (int j = 0; j < tok.size(); j++) {
65             dp[i][j] = vector<long long>(state + 1, INT_MAX);
66             neg_INF[i][j] = vector<bool>(state + 1, false);
67         }
68         dp[i][i][tok[i]] = 0;
69         bellman(i, i, tok.size());
70     }
71     for (int r = 1; r < tok.size(); r++) {
72         for (int l = r - 1; l >= 0; l--) {
73             for (int k = l; k < r; k++)
74                 for (rule c : cnf)
75                     if (c.y != -1)
76                         relax(l, r, c,
77                             dp[l][k][c.x] + dp[k + 1][r][c.y] +
78                             c.cost);
79             bellman(l, r, tok.size());
80         }
81     }
82 }
83 // usage example
84 int main() {
85     init();
86     add_to_cnf('S', "aSc", 1);
87     add_to_cnf('S', "BBB", 1);
88     add_to_cnf('S', "SB", 1);
89     add_to_cnf('B', "b", 1);
90     cyk("abbbbc");
91     // dp[0][s.size() - 1][rules[start]] = min cost to
92     // generate s
93     cout << dp[0][5][rules['S']] << '\n'; // 7
94     cyk("acbc");
95     cout << dp[0][3][rules['S']] << '\n'; // INT_MAX
96     add_to_cnf('S', "S", -1);
97     cyk("abbbbc");
98     cout << neg_INF[0][5][rules['S']] << '\n'; // 1
99 }

```

7.5. Z Value

```

1 int z[n];
2 void zval(string s) {
3     // z[i] => longest common prefix of s and s[i:], i > 0
4     int n = s.size();
5     z[0] = 0;
6     for (int b = 0, i = 1; i < n; i++) {
7         if (z[b] + b <= i) z[i] = 0;
8         else z[i] = min(z[i - b], z[b] + b - i);
9         while (s[i + z[i]] == s[z[i]]) z[i]++;
10        if (i + z[i] > b + z[b]) b = i;
11    }
12 }

```

7.6. Manacher's Algorithm

```

1 int z[n];
2 void manacher(string s) {
3     // z[i] => longest odd palindrome centered at i is
4     // s[i - z[i]] ... i + z[i]
5     // to get all palindromes (including even length),
6     // insert a '#' between each s[i] and s[i + 1]
7     int n = s.size();
8     z[0] = 0;
9     for (int b = 0, i = 1; i < n; i++) {
10        if (z[b] + b >= i)
11            z[i] = min(z[2 * b - i], b + z[b] - i);
12        else z[i] = 0;
13        while (i + z[i] + 1 < n && i - z[i] - 1 >= 0 &&
14            s[i + z[i] + 1] == s[i - z[i] - 1])

```

```

15        z[i]++;
16        if (z[i] + i > z[b] + b) b = i;
17    }
18 }

```

7.7. Minimum Rotation

```

1 int min_rotation(string s) {
2     int a = 0, n = s.size();
3     s += s;
4     for (int b = 0; b < n; b++) {
5         for (int k = 0; k < n; k++) {
6             if (a + k == b || s[a + k] < s[b + k]) {
7                 b = max(0, k - 1);
8                 break;
9             }
10            if (s[a + k] > s[b + k]) {
11                a = b;
12                break;
13            }
14        }
15    }
16    return a;
17 }

```

7.8. Palindromic Tree

```

1 struct palindromic_tree {
2     struct node {
3         int next[26], fail, len;
4         int cnt,
5             num; // cnt: appear times, num: number of pal. suf.
6         node(int l = 0) : fail(0), len(l), cnt(0), num(0) {
7             for (int i = 0; i < 26; ++i) next[i] = 0;
8         }
9     };
10    vector<node> St;
11    vector<char> s;
12    int last, n;
13    palindromic_tree() : St(2), last(1), n(0) {
14        St[0].fail = 1, St[1].len = -1, s.pb(-1);
15    }
16    inline void clear() {
17        St.clear(), s.clear(), last = 1, n = 0;
18        St.pb(0), St.pb(-1);
19        St[0].fail = 1, s.pb(-1);
20    }
21    inline int get_fail(int x) {
22        while (s[n - St[x].len - 1] != s[n]) x = St[x].fail;
23        return x;
24    }
25    inline void add(int c) {
26        s.push_back(c == 'a'), ++n;
27        int cur = get_fail(last);
28        if (!St[cur].next[c]) {
29            int now = SZ(St);
30            St.pb(St[cur].len + 2);
31            St[now].fail = St[get_fail(St[cur].fail)].next[c];
32            St[cur].next[c] = now;
33            St[now].num = St[St[now].fail].num + 1;
34        }
35        last = St[cur].next[c], ++St[last].cnt;
36    }
37    inline void count() { // counting cnt
38        auto i = St.rbegin();
39        for (; i != St.rend(); ++i) {
40            St[i->fail].cnt += i->cnt;
41        }
42    }
43    inline int size() { // The number of diff. pal.
44        return SZ(St) - 2;
45    }
46 };

```