

[Traversing async DAG](#)

[help](#)

[carlospereira1607](#) January 4, 2020, 1:46pm 1

Hello,

Currently I'm trying to get into Rust `async/await` and the Tokio crate. One of the things that I'm trying to do is implement an async version of a algorithm that I have that uses a DAG. This DAG that was created by me and under the hood there is a vec. I implemented the DAG as a vec because I want to reuse positions from "deleted" nodes of the DAG (a node is just softly deleted). I have 3 functions that update the graph: `add` (that might call `update`), `update` (which is recursive) and `remove`. When `add` is called, it checks the predecessors of that node in the graph (by trying to get the predecessors from the graph) and adds their indexes in the vec under hood to another vec that the struct node has.

What I did from my initial version without `async/await` was I transformed the functions `add`, `update` and `remove` as async methods and I wrapped the DAG on Tokio's `Mutex`. Whenever I want to mutate I call the `.lock().await` from the mutex, do the changes that I want and then I drop the variable so the lock is dropped.

The problem is traversing and mutating the graph. From my logs, if I add a node its predecessors (that already exists in the graph) are marked as not being there (even though the logs show that they are).

Another approach that I thought was having a `Mutex` at each node of the graph instead of wrapping the whole graph. This way, only one thread will have a lock at each node. Perhaps this second approach is better?

Regards,

[raidwas](#) January 4, 2020, 3:45pm 2

Without seeing the code the following is all guessing:

I assume the non-async version works?

Since the whole dag is inside a mutex there shouldn't be a race condition, so my guess would be that there is a problem in your data structure that gets triggered by a difference in call order from multiple threads compared to the single threaded version of your code.

Maybe try the `await` version with your single threaded algorithm?

About the other approach: without care that sounds very deadlock prone. You may want to read up on some of the required conditions for deadlocks. For example if you can make sure that all your resources (in this case nodes) have a unique id and every thread only allocates nodes in ascending order it would be impossible to create a deadlock.

Which approach is better (assuming better means faster, not simpler), it would depend on what you do with the data structure (for example how long does a thread lock the data structure (approach 1), how many mutually exclusive allocations would there be etc (approach 2)).

But approach certainly sounds simpler than 2.

[carlospereira1607](#) January 4, 2020, 5:50pm 3

Yes, the non-async version works (forgot to mention that). I did implement an async version that has the DAG struct wrapped in the `Mutex` and whenever its necessary to update it it calls `.lock().await`. Unfortunately, this version isn't that much better than my initial single threaded implementation. Since the functions that update the graph are still sequential and non concurrent, perhaps this makes sense.

In the end I just wanted to know if it would be possible to traverse a DAG asynchronously. The way I have the code it might not be possible to do it asynchronously, since I have to do a lot of comparisons between the node I'm currently visiting and its predecessors/successors.

[alice](#) January 4, 2020, 5:58pm 4

If you are not modifying it, you can share it with an `Arc` and just traverse it without any asynchronous functionality.

[raidwas](#) January 4, 2020, 5:58pm 5

If most of the time from the algorithm is spend in functions that work on `&must Dag` then it makes sense that the threaded version isn't much faster since most threads will just be waiting on the lock.

As for whether it is possible I would say yes, if it suits your needs is hard to say without knowing what you want to do with it.

[carlospereira1607](#) January 5, 2020, 12:25pm 6

What I want to do is keep adding nodes to the DAG and when a node has all of its predecessors in the graph, I want to remove it (by marking it as deleted so the position can be later used), but before I remove a node, I want to remove all of its predecessors. When I add a node to the graph I already know which nodes are its predecessors, but I don't know if they're already in the graph. If not, then the node will be in the graph until they are.

The idea of having the vec with the `Mutexes` doesn't sound very good to me, mainly because it will be shared by the threads and when a node were to be added to the graph, they would have to add it to this vec.

[carlospereira1607](#) January 5, 2020, 12:26pm 7

I do need to be able to mutate the graph, so wrapping it in an Arc doesn't work... But thanks for the suggestion!

[system](#) Closed April 4, 2020, 12:26pm 8

This topic was automatically closed 90 days after the last reply. New replies are no longer allowed.

Related topics

| Topic | | Replies | Views | Activity |
|---|----|---------|-------------------|----------|
| Safe-threading and async | | | | |
| HashMap | 5 | 164 | November 15, 2025 | |
| help | | | | |
| What makes async mutex more expensive than sync mutex? | 17 | 3059 | January 5, 2024 | |
| help | | | | |
| How to access vector from multiple threads without mutex? | 28 | 4151 | July 28, 2023 | |
| Mutable struct fields with async/await | 20 | 5751 | October 8, 2020 | |
| help | | | | |
| ``HashMap::retain()`` and async/await | 4 | 1936 | April 9, 2020 | |
| help | | | | |

- [Home](#)
- [Categories](#)
- [Guidelines](#)
- [Terms of Service](#)