Short Assignment: The Substitution Model

Trace the evaluation of the expression involving the mystery procedure:

```
(mystery 7 2)
```

• Evaluate (mystery 7 2) where mystery is defined below.

• Evaluate (lambda) subexpression

```
(lambda ((a <number>) (b <integer>))
```

Since this is the evaluation of a compound expression, the EVAL rule says we must first evaluate the subexpressions. The value obtained by evaluating the <code>lambda</code> subexpression is a prodecure object, denoted:

```
[PROC ((x <number>) (b <number>))]
```

Evaluate cond expression

```
(cond ((zero? b) 0)
```

cond is a special form, so we handle it specially (has its own rule in tho model).

• Evaluate ((zero? b) 0)

```
(PROC ((b <integer>)) ((zero?) (b)) 0)
```

Evaluate 2

```
(PROC ((n <number>)) ((zero?) 2) 0)
```

We determine the value of the expression b is 2.

• Apply zero? procedure to 2

```
[PROC ((n <number>)) ((zero?) 2) 0]
```

We take the body of the procedure, zero? and subtitute the corresponding arguments where the parameters appear.

• Evaluate zero? procedure

```
(zero? 2)
```

 $\overline{\text{zero? x}}$ is a primitive procedure that returns #t if its numeric agreement x is equal to zero, #f otherwise. The above expression returns #f. The consequent is not evaluated.

• Evaluate cond expression

```
(cond ((odd? b)
```

We evaluate [cond], since the first test yields a false value, we evaluate the next test.

• Evaluate ((odd? b)

```
(PROC ((b <integer>)) ((odd?) b)))
```

• Evaulate 2

```
(PROC ((b <integer>)) ((odd?) 2)))
```

We determine the value of the expression b is 2.

• Apply odd? procedure to 2

```
[PROC ((n <number>)) ((odd?) 2) 0]
```

• Evaluate odd? procedure

odd? x returns #t if its integer argument x is odd, #f otherwise. The consequent is not evaluated.

• Evaluate the cond expression

We evaluate [cond], since the first test yields a false value, we evaluate the next test.

• Evaluate the else expression

The else clause is optional, but if it is present, and all the test expressions test1 through testn have evaluated to #f, then expr-e is evaluated, and its value becomes the value of the cond expression.

• Evaluate (mystery (+ a a) (quotient b 2))

Evaluate mystery

• Evaluate (+ a a)

```
(PROC (a <number>) (+ a a))
```

• Evaluate 7

```
(PROC (a <number>) (+ 7 7))
```

Apply procedure

```
[PROC (a <number>) (+ 7 7)]
```

Evaluate +

(+ 7 7)

Apply primitive add

```
[add 7 7]
```

Returns 14.

• Evaluate (quotient b 2)

```
(PROC (b <integer>) (quotient b 2))
```

• Evaluate 2

```
(PROC (b <integer>) (quotient 2 2))
```

• Apply quotient procedure

[PROC (b <integer>) (quotient 2 2)]

• Evaluate (quotient 2 2)

(quotient 2 2)

• Apply primitive quotient

[quotient 2 2]

quotient returns the quotient of a / b. In other words, if a = qb + r, with $0 \le r < b$, then (quotient a b) returns q. The above expression returns 1.0.

• Substitute the body of the procedure with values {14} and {1.0}.

(mystery 14 1.0)

• The process repeats, and so on and so forth.

((cond ((zero? b) 0)) returns #f. We evaluate the next test. (((odd? b)) returns #t. We evaluate the consequent.