

# Introduction to Object-Oriented Programming

## Binary Search Trees

Christopher Simpkins

`chris.simpkins@gatech.edu`

Trees are everywhere.



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<sup>1</sup>Source:[http://commons.wikimedia.org/wiki/File:Winnersh\\_Meadows\\_Trees.jpg](http://commons.wikimedia.org/wiki/File:Winnersh_Meadows_Trees.jpg)

They're in our web browsers.

The screenshot shows a web browser window with the address bar displaying `www.cc.gatech.edu/~simpkins/teaching/gatech/cs1331/syllabus/cs1331-fall2013.html`. The page title is "CS 1331 Introduction to Object Oriented Programming". The main heading is "CS 1331 Introduction to Object Oriented Programming" followed by "Fall 2013 Syllabus for Sections A1-A4, B1-B4, and GR". The page content includes a last update date of 2013-12-03 at 20:10, a statement about the open-access syllabus, and several links to course policies. The browser's developer tools are open at the bottom, showing the DOM tree with the `h1` element selected. The DOM tree shows a document structure with a single `h1` element. The console, inspector, debugger, style editor, profiler, and network panels are also visible.

CS 1331 Introduction to Object Oriented Programming

GT | GT Login CS 1331 Introduction to Object ... CS 1331 (237 unread)

www.cc.gatech.edu/~simpkins/teaching/gatech/cs1331/syllabus/cs1331-fall2013.html

Most Visited Getting Started Work Personal Scala Diving

# CS 1331 Introduction to Object Oriented Programming

## Fall 2013 Syllabus for Sections A1-A4, B1-B4, and GR

Last updated on 2013-12-03 at 20:10.

This open-access part of the syllabus contains schedule and general information for my sections. All grade-related information are on T-Square.

There is also a general syllabus here: [cs1331-syllabus.pdf](#).

Please read and be sure you fully understand the [cs1331-course-policies.html](#).

Please read and be sure you fully understand the [cs1331-course-policies.html](#).

I'm serious! Please read and be sure you fully understand the [cs1331-course-policies.html](#) !!!

Console Inspector Debugger Style Editor Profiler Network

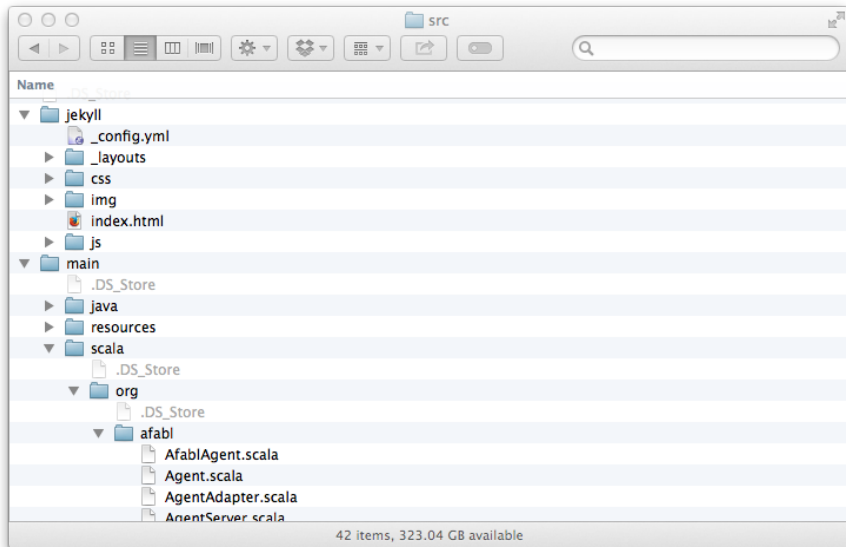
html body div#cs-1331-introduction-to-object-or h1.title

```
<!--?xml version="1.0" encoding="utf-8" ?-->
<!DOCTYPE html PUBLIC "-//W3C//DTD XHTML 1.0 Transitional//EN" "http://www.w3.org/TR/xhtml1/DTD/xhtml1-transitional.dtd">
<html lang="en" xml:lang="en" xmlns="http://www.w3.org/1999/xhtml">
  <head> ... </head>
  <body>
    <div id="cs-1331-introduction-to-object-oriented-programming" class="document"> ... </div>
  </body>
</html>
```

Rules Computed Fonts Box

element { inline

They're in our file systems.



They're even in pop culture.



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<sup>2</sup>Source:

<http://userserve-ak.last.fm/serve/500/44019065/Neon+Trees.png>

But they're not in Kansas.



3

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<sup>3</sup>Source:

[http://en.wikipedia.org/wiki/File:Wabaunsee\\_County\\_View.JPG](http://en.wikipedia.org/wiki/File:Wabaunsee_County_View.JPG)



# Binary Tree Nodes

The nodes of a binary tree have

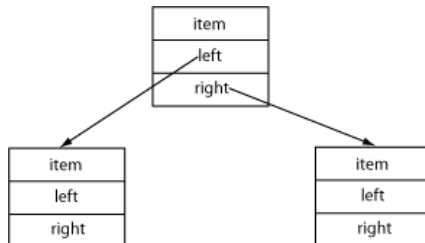
- a data item,
- a link to a left node, and
- a link to a right node.

```
private class Node<E> {  
    E item;  
    Node<E> left;  
    Node<E> right;  
  
    Node(E item, Node<E> left, Node<E> right) {  
        this.item = item;  
        this.left = left;  
        this.right = right;  
    }  
}
```

Just as in the other linked data structures we've studied, binary trees are recursive.

# Binary Tree Structure

- Every tree has a distinguished *root node* with no parent.
  - All other nodes have exactly one parent.
- Nodes which have no children are called *leaf nodes*.
- Nodes which have children are called *interior nodes*.
- Every node has 0, 1, or 2 children.
- Every node can be reached by a unique path from the root node.





# Binary Search Trees

A binary search tree (BST) encodes the binary search algorithm into its structure. The BST property: for any node,

- all the elements in the node's left subtree are less than the node's data item, and
- all the elements in the node's right subtree are equal to or greater than the node's data item.

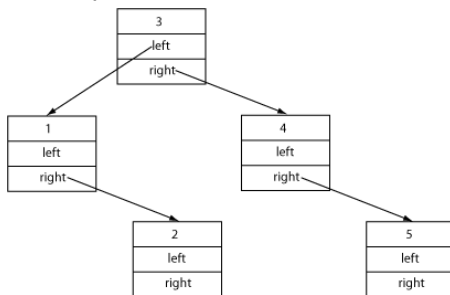
A BST is distinguished by this property, but its ADT is just like the others we've seen: add elements, find element's in the tree, and iterate over the elements in a tree.

# Maintaining The BST Property

To add a new value to binary tree and maintain the BST property, we

- insert new nodes for data items into the left subtree of a node if the new item is less than the node's item, or
- the right subtree otherwise.

Every new item creates a leaf node, which can later become an interior node after additional items have been added. Here's the structure of a BST after adding the sequence of numbers 3, 4, 1, 5, 2:



# Adding Elements to a BST

```
public class BinarySearchTree<E extends Comparable<? super E>>
    implements Iterable<E> {

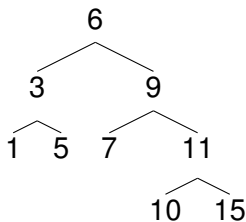
    protected class Node<E> { ... }

    protected Node<E> root;

    public void add(E item) {
        root = insert(item, root);
    }
    protected Node<E> insert(E newItem, Node<E> node) {
        if (node == null) {
            return new Node<E>(newItem, null, null);
        } else if (newItem.compareTo(node.item) < 0) {
            node.left = insert(newItem, node.left);
            return node;
        } else {
            node.right = insert(newItem, node.right);
            return node;
        }
    }
}
```

# Exercise: Insertion Locations

Given the following tree that conforms to the binary search tree property:



Where would 2, 4, 8, and 16 be inserted in the tree?

# Traversing a Binary Tree

There are three primary ways to traverse a binary tree:

Pre-order:

- Process node's item.
- Process left subtree.
- Process right subtree.

In-order:

- Process left subtree.
- Process node's item.
- Process right subtree.

Post-order:

- Process left subtree.
- Process right subtree.
- Process node's item.

# Simple In-Order Traversal

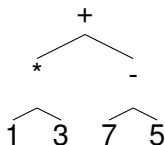
Traversal code follows the recursive structure of the tree:

```
public void printInOrder() {  
    printInOrder(root);  
}  
  
private void printInOrder(Node<E> node) {  
    if (node != null) {  
        printInOrder(node.left);  
        System.out.print(node.item + " ");  
        printInOrder(node.right);  
    }  
}
```

The code above prints the elements in ascending order. Let's add a `printDescending()` method to [BinarySearchTree.java](#).

# Exercise: Traversal Orders

Given the following tree:



If we processed each element by printing it, in what order would the elements be printed

- For a pre-order traversal:
- For an in-order traversal:
- For a post-order traversal:

# The Path to an Item

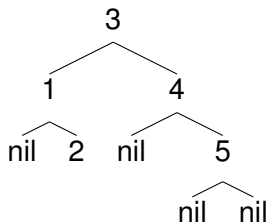
To find a path to an item in a BST:

- set the path to the empty list
- set the root node as the currentNode
- until we find the node containing the item or exhaust the BST:
  - if currentNode contains the item, add it to the path and return it
  - else if query item is less than the item in currentNode, add currentNode to path and set the left child as the new currentNode
  - else add add currentNode to path and set the right child as the new currentNode
- if the item wasn't found, set the path to the empty list



# Path Examples

Adding the elements [3, 4, 1, 5, 2] to a BST would result in the following structure:



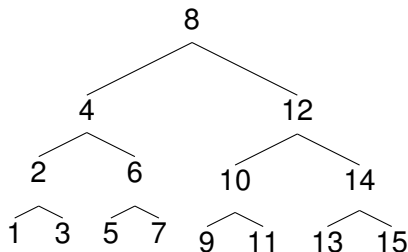
and the paths to each element in the tree would be:

- Path to 1: [3, 1]
- Path to 2: [3, 1, 2]
- Path to 3: [3]
- Path to 4: [3, 4]
- Path to 5: [3, 4, 5]

See `public List<E> path(E queryItem)` in [BinarySearchTree.java](#) for the code.

# Exercise: Paths

Given the following tree:



- What's the path to 1?
- What's the path to 11?

## Recursively Building a Result: `inOrderList()`

We can use the recursive accumulator idiom to collect the elements of the tree in an in-order traversal:

```
public List<E> toList() {  
    return inOrderList(root, new ArrayList<E>());  
}  
  
private List<E> inOrderList(Node<E> node, List<E> accum) {  
    if (null == node) {  
        return accum;  
    } else {  
        inOrderList(node.left, accum);  
        accum.add(node.item);  
        inOrderList(node.right, accum);  
    }  
    return accum;  
}
```

Again, the code follows the recursive structure of the tree.

# Imperative traversal: `inOrderImperative()`

Contrast the previous code for getting an in-order list of BST elements with an imperative version:

```
public List<E> inOrderImperative() {
    Node<E> curNode = root;
    Stack<Node<E>> fringe = new LinkedStack<>();
    List<E> accum = new ArrayList<E>();
    while ((curNode != null) || !fringe.isEmpty()) {
        while (curNode != null) {
            fringe.push(curNode);
            curNode = curNode.left;
        }
        curNode = fringe.pop();
        accum.add(curNode.item);
        curNode = curNode.right;
    }
    return accum;
}
```

We need extra bookkeeping variables to keep track of where we are in the tree so we can back-track. See [BinarySearchTree.java](#) for comments explaining the algorithm.

# Iterators

Iterators can free clients from having to implement traversal algorithms. We can even plug our data structures into Java's for-each loop by implementing `java.lang.Iterable`:

```
public interface Iterable<T> {  
    java.util.Iterator<T> iterator();  
}
```

As a reminder, `java.util.Iterator`:

```
public interface Iterator<E> {  
    boolean hasNext();  
  
    E next();  
  
    void remove();  
}
```

# Stateful In-Order Tree Traversal

In the traversal examples we saw earlier the traversal order was effected by the method call stack. A stateful iterator is much more challenging because:

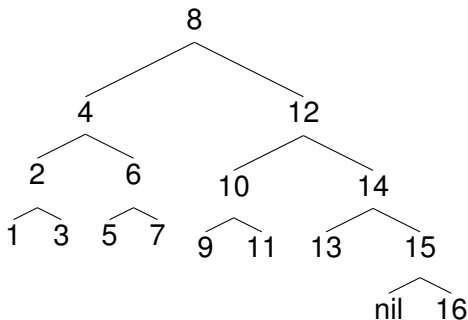
- The iterator must remember where it is in the tree
- The iterator must be able to back-track to parent nodes after processing child branches

The essential implementation idea is to use a stack to store nodes for back-tracking. Traditionally (at least in AI), this “to-do list” stack is called the *fringe*.

Let's look at [BinarySearchTree.java](#) again to see how we implement a stateful in-order iterator.

# Analysis of BSTs

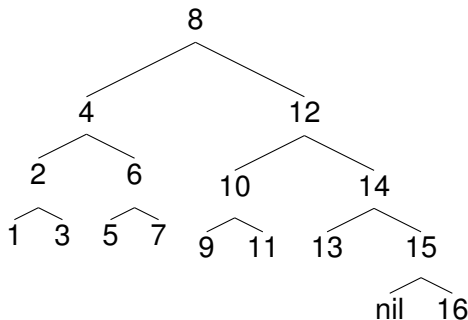
What is the Big-O of finding an element in a BST?



- Proportional to height of tree
- Height of tree is proportional to  $\log n$
- 16-element tree has height 4

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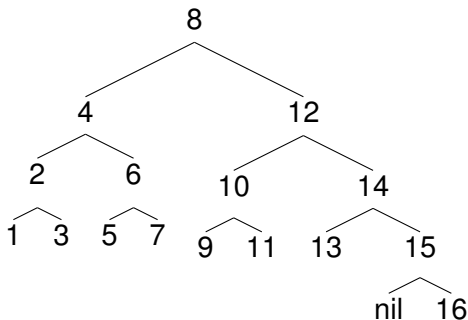


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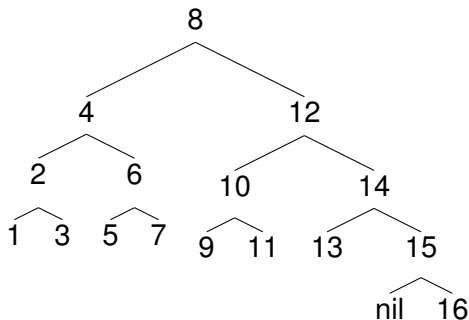
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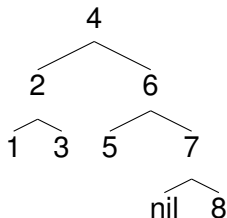
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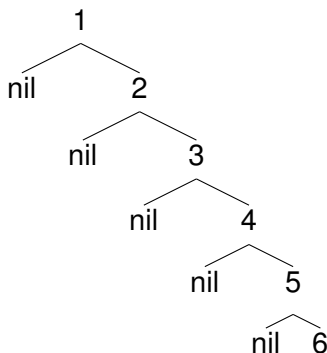
# Height of BST Proportional to $\log n$



8-element tree has height 3

# End of course cliff hanger ...

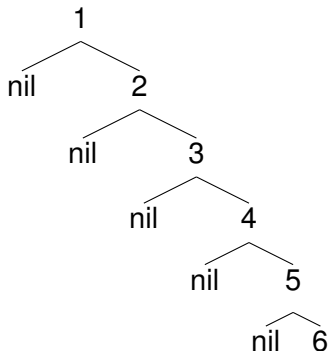
What if we add elements to a BST in order?



- We end up with a linked list!
- In your algorithms and data structures course you'll learn how to maintain a balanced binary search tree.

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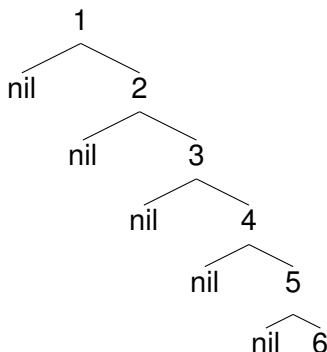
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