Mathematical Foundations of Data Sciences



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Chapter 13

Convex Optimization

The main references for this chapter are [10, 11, 5], see also [31, 4, 3]. We consider a general convex optimization problem

$$\min_{x \in \mathcal{H}} f(x) \tag{13.1}$$

where $\mathcal{H} = \mathbb{R}^N$ is a finite dimensional Hilbertian (i.e. Euclidean) space, and try to devise "cheap" algorithms with a low computational cost per iterations. The class of algorithms considered are first order, i.e. they make use of gradient information.

13.1 Gradient Descent Methods

We have already encountered the gradient descent method informally in Section ?? for the regularization of inverse problem. We now give a detailed analysis of the method.

13.1.1 Gradient Descent

The optimization program (10.26) is a example of unconstrained convex optimization of the form (13.1) where $f: \mathcal{H} \to \mathbb{R}$ is a \mathcal{C}^1 function with Lipschitz gradient (so-called "smooth" function). Recall that the gradient $\nabla f: \mathcal{H} \mapsto \mathcal{H}$ of this functional (not to be confound with the discretized gradient $\nabla x \in \mathcal{H}$ of f) is defined by the following first order relation

$$f(x+r) = f(x) + \langle f, r \rangle_{\mathcal{H}} + O(\|r\|_{\mathcal{H}}^2)$$

where we used $O(\|r\|_{\mathcal{H}}^2)$ in place of $o(\|r\|_{\mathcal{H}})$ (for differentiable function) because we assume here f is of class \mathcal{C}^1 (i.e. the gradient is continuous). Section 10.4.3 shows typical examples of gradient computation.

For such a function, the gradient descent algorithm is defined as

$$x^{(\ell+1)} \stackrel{\text{def.}}{=} x^{(\ell)} - \tau_{\ell} \nabla f(x^{(\ell)}), \tag{13.2}$$

where the step size $\tau_{\ell} > 0$ should be small enough to guarantee convergence, but large enough for this algorithm to be fast.

One also needs to quantify the smoothness of f. This is enforced by requiring that the gradient is L-Lipschitz, i.e.

$$\forall (x, x') \in \mathcal{H}^2, \quad \|\nabla f(x) - \nabla f(x')\| \leqslant L \|x - x'\|. \tag{\mathcal{R}_L}$$

In order to obtain fast convergence of the iterates themselve, it is needed that the function has enough "curvature" (i.e. is not too flat), which corresponds to imposing that f is μ -strongly convex

$$\forall (x, x'), \in \mathcal{H}^2, \quad \langle \nabla f(x) - \nabla f(x'), x - x' \rangle \geqslant \mu \|x - x'\|^2. \tag{S_{\mu}}$$

The following proposition express these conditions as constraints on the hessian for C^2 functions.

Proposition 32. Conditions (\mathcal{R}_L) and (\mathcal{S}_{μ}) imply

$$\forall (x, x'), \quad f(x') + \langle \nabla f(x), x' - x \rangle + \frac{\mu}{2} \|x - x'\|^2 \leqslant f(x) \leqslant f(x') + \langle \nabla f(x'), x' - x \rangle + \frac{L}{2} \|x - x'\|^2. \quad (13.3)$$

If f is of class C^2 , conditions (\mathcal{R}_L) and (\mathcal{S}_{μ}) are equivalent to

$$\forall x, \quad \mu \operatorname{Id}_{N \times N} \leq \partial^2 f(x) \leq L \operatorname{Id}_{N \times N}$$
 (13.4)

where $\partial^2 f(x) \in \mathbb{R}^{N \times N}$ is the Hessian of f, and where \leq is the natural order on symmetric matrices, i.e.

$$A \leq B \iff \forall x \in \mathcal{H}, \quad \langle Au, u \rangle \leqslant \langle Bu, u \rangle.$$

Proof. We prove (13.3), using Taylor expansion with integral remain

$$f(x') - f(x) = \int_0^1 \langle \nabla f(x_t), x' - x \rangle dt = \langle \nabla f(x), x' - x \rangle + \int_0^1 \langle \nabla f(x_t) - \nabla f(x), x' - x \rangle dt$$

where $x_t \stackrel{\text{def.}}{=} f + t(x' - x)$. Using Cauchy-Schwartz, and then the smoothness hypothesis (\mathcal{R}_L)

$$f(x') - f(x) \le \langle \nabla f(x), x' - x \rangle + \int_0^1 L \|x_t - f\| \|x' - x\| dt \le \langle \nabla f(x), x' - x \rangle + L \|x' - x\|^2 \int_0^1 t dt$$

which is the desired upper-bound. Using directly (S_{μ}) gives

$$f(x') - f(x) = \langle \nabla f(x), x' - x \rangle + \int_0^1 \langle \nabla f(x_t) - \nabla f(x), \frac{x_t - x}{t} \rangle dt \geqslant \langle \nabla f(x), x' - x \rangle + \mu \int_0^1 \frac{1}{t} \|x_t - x\|^2 dt$$

which gives the desired result since $||x_t - x||^2/t = t||x' - x||^2$.

The relation (13.3) shows that a smooth (resp. strongly convex) functional is bellow a quadratic tangential majorant (resp. minorant).

Condition (13.4) thus reads that the singular values of $\partial^2 f(x)$ should be contained in the interval $[\mu, L]$. The upper bound is also equivalent to $\|\partial^2 f(x)\|_{\text{op}} \leq L$ where $\|\cdot\|_{\text{op}}$ is the operator norm, i.e. the largest singular value. In the special case of a quadratic function $\mathcal Q$ of the form (10.24), $\partial^2 f(x) = A$ is constant, so that $[\mu, L]$ can be chosen to be the range of the singular values of A.

The following theorem ensure the convergence of the gradient descent with a linear speed.

Theorem 35. If f satisfy conditions (\mathcal{R}_L) and (\mathcal{S}_{μ}) , assuming there exists $(\tau_{\min}, \tau_{\max})$ such that

$$0 < \tau_{\min} \leqslant \tau_{\ell} \leqslant \tau_{\max} < \frac{2\mu}{L} \tag{13.5}$$

then there exists $0 \le \rho < 1$ such that

$$\|x^{(\ell)} - x^*\| \leqslant \rho^{\ell} \|x^{(0)} - x^*\| \tag{13.6}$$

where x^* is the unique solution to (??).

Proof. Since $\nabla f(x^*) = 0$, one has

$$x^{(\ell+1)} - x^* = (x^{(\ell)} - x^*) - \tau_{\ell}(\nabla f(x^{(\ell)}) - \nabla f(x^*)).$$

Hence, using strong convexity and Lipschitz gradient

$$\begin{aligned} \|x^{(\ell+1)} - x^{\star}\|^2 &= \|x^{(\ell)} - x^{\star}\|^2 - 2\tau_{\ell}\langle x^{(\ell)} - x^{\star}, \nabla f(x^{(\ell)}) - \nabla f(x^{\star})\rangle + \tau_{\ell}^2 \|\nabla f(x^{(\ell)}) - \nabla f(x^{\star})\|^2 \\ &\leqslant P(\tau_{\ell}) \|x^{(\ell)} - x^{\star}\|^2 \quad \text{where} \quad P(\tau) = 1 - 2\mu\tau + L^2\tau^2. \end{aligned}$$

Figure 13.1, left, shows visually the shape of the second order polynomial P, which shows that condition (13.11) on τ_{ℓ} implies

$$P(\tau_{\ell})^{\frac{1}{2}} \leqslant \rho \stackrel{\text{\tiny def.}}{=} \max(P(\tau_{\min}), P(\tau_{\max}))^{\frac{1}{2}} < 1,$$

which shows the desired result.

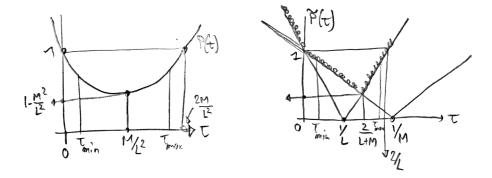


Figure 13.1: Contraction constant $P(\tau)$ and $\tilde{P}(\tau)$ for a gradient descent step in the generic case (left) and for a quadratic function (right).

The error decay rate (13.9), although it is geometrical $O(\rho^{\ell})$ is called a "linear rate" in the optimization literature. It is a "global" rate because it hold for all ℓ (and not only for large enough ℓ). The best (smallest) rate ρ is obtained when choosing

$$\tau_{\ell} = \frac{\mu}{L^2} \quad \Longrightarrow \quad \rho = 1 - \frac{\mu^2}{L^2}. \tag{13.7}$$

In the case of a quadratic functional of the form (10.24), one can sharpen the convergence proof because the iterates are computed in closed form using matrix multiplication

$$x^{(\ell)} - x^* = (\mathrm{Id}_N - \tau_\ell A)(x^{(0)} - x^*)$$

which leads to the following proposition (see also Figure 13.1, right, for the corresponding contraction constant involved as a function of τ).

Proposition 33. For $f(x) = \langle A, f \rangle - \langle b, x \rangle$ with the singular values of A upper-bounded by L, assuming there exists $(\tau_{\min}, \tau_{\max})$ such that

$$0 < \tau_{\min} \leqslant \tau_{\ell} \leqslant \tilde{\tau}_{\max} < \frac{2}{L} \tag{13.8}$$

then there exists $0 \leq \tilde{\rho} < 1$ such that

$$\|x^{(\ell)} - x^*\| \leqslant \tilde{\rho}^{\ell} \|x^{(0)} - x^*\|. \tag{13.9}$$

If the singular values are lower bounded by μ , then the best rate $\tilde{\rho}$ is obtained for

$$\tau_{\ell} = \frac{2}{L+\mu} \implies \tilde{\rho} \stackrel{\text{def.}}{=} \frac{L-\mu}{L+\mu}.$$
(13.10)

The maximum allowable step size $\tilde{\tau}_{\text{max}}$ in (13.11) is much larger than τ_{max} given in (13.11), and the optimal rate (13.10) is also much better (smaller) than the one in (13.7). In particular, if

$$\varepsilon \stackrel{\text{def.}}{=} M/L \ll 1$$

(which is the typical setup for ill-posed problems), then

$$\rho \sim 1 - \varepsilon^2$$
 and $\tilde{\rho} \sim 1 - 2\varepsilon$.

The quantity ε in some sense reflects the inverse-conditioning of the problem. For quadratic function, it indeed corresponds exactly to the inverse of the condition number (which is the ratio of the largest to smallest singular value). The condition number is minimum and equal to 1 for orthogonal matrices.

These two results are however complementary. Indeed, if the gradient descent converges, then ultimately $x^{(\ell)}$ is close to x^* , so that one can approximate up to second order $f(x) \approx f(x^*) + \langle Af, f \rangle - \langle f, b \rangle$ with $A = \partial^2 f(x^*)$ and $b = -\nabla f(x^*)$. So that the "local" rate, the one obtained after a large enough of iterations, is actually driven by $\tilde{\rho}$ and not ρ . It is thus important to distinguish between the global rate and the local rate. In practice, descent algorithm typically have two phase: a first "slow" phase govern by the global rate, and a second "fast" phase governed by the local rate. Unfortunately, the optimal step sizes τ_{ℓ} are in general different for the two phase, so that optimal adaptation of step size is a difficult problems. This is why more advanced users typically use various line search strategies (to find the optimal step size at each iteration) or use second order information using quasi-Newton technics (BFGS).

The convergence result of Proposition 33 does not requires strong convexity, while Theorem 35 does. In the general non-strongly convex case, it is still possible to prove convergence, but the rate is only sub-linear, and is only on the value of f, not on the iterate $x^{(\ell)}$ themselves. Note that in this case, the solution of the minimization problem is not necessarily unique. The proof is more technical.

Theorem 36. If f satisfy conditions (\mathcal{R}_L) , assuming there exists $(\tau_{\min}, \tau_{\max})$ such that

$$0 < \tau_{\min} \leqslant \tau_{\ell} \leqslant \tau_{\max} < \frac{2}{L},\tag{13.11}$$

then $x^{(\ell)}$ converges to a solution x^* of $(\ref{eq:converges})$ and there exists C>0 such that

$$f(x^{(\ell)}) - f(x^*) \leqslant \frac{C}{\ell + 1}.$$
 (13.12)

Proof. We only prove (13.12) since the proof that $x^{(\ell)}$ converges is more technical. Note indeed that if the minimizer x^* is non-unique, then it might be the case that the iterate $x^{(\ell)}$ "cycle" while approaching the set of minimizer, but actually convexity of f prevents this kind of pathological behavior. For simplicity, we do the proof in the case $\tau_{\ell} = 1/L$, but it extends to the general case. The L-smoothness property imply (13.3), which reads

$$f(x^{(\ell+1)}) \leqslant f(x^{(\ell)}) + \langle \nabla f(x^{(\ell)}), x^{(\ell+1)} - x^{(\ell)} \rangle + \frac{L}{2} \|x^{(\ell+1)} - x^{(\ell)}\|^2.$$

Using the fact that $x^{(\ell+1)} - x^{(\ell)} = -\frac{2}{L}\nabla f(x^{(\ell)})$, one obtains

$$f(x^{(\ell+1)}) \leqslant f(x^{(\ell)}) - \|\nabla f(x^{(\ell)})\|^2 \leqslant f(x^{(\ell)}) - \frac{1}{2L} \|\nabla f(x^{(\ell)})\|^2$$
(13.13)

This shows that $(f(x^{(\ell)}))_{\ell}$ is a decaying sequence. By convexity

$$f(x^{(\ell)}) + \langle \nabla f(x^{(\ell)}), x^* - x^{(\ell)} \rangle \leq f(x^*)$$

and plugging this in (13.13) shows

$$\begin{split} f(x^{(\ell+1)}) &\leqslant f(x^{\star}) - \langle \nabla f(x^{(\ell)}), \, x^{\star} - x^{(\ell)} \rangle - \frac{1}{2L} \| \nabla f(x^{(\ell)}) \|^2 \\ &= f(x^{\star}) + \frac{L}{2} \left(\| x^{(\ell)} - x^{\star} \|^2 - \| x^{(\ell)} - x^{\star} - \frac{1}{L} \nabla f(x^{(\ell)}) \|^2 \right) \\ &= f(x^{\star}) + \frac{L}{2} \left(\| x^{(\ell)} - x^{\star} \|^2 - \| x^{\star} - x^{(\ell+1)} \|^2 \right). \end{split}$$

Summing these inequalities for $\ell = 0, ..., k$, one obtains

$$\sum_{\ell=1}^{k} f(x^{(\ell+1)}) - kx^* \leqslant \frac{L}{2} \left(\|x^{(0)} - x^*\|^2 - \|x^{(k+1)} - x^*\|^2 \right)$$

and since $f(x^{(\ell+1)})$ is decaying $\sum_{\ell=1}^k f(x^{(\ell+1)}) \geqslant (k+1)f(x^{(k+1)})$, thus

$$f(x^{(k+1)}) - f(x^*) \le \frac{L\|x^{(0)} - x^*\|^2}{2(k+1)}$$

which gives (13.12) for $C \stackrel{\text{def.}}{=} L \|x^{(0)} - x^*\|^2 / 2$.

13.1.2 Sub-gradient Descent

The gradient descent (13.2) cannot be applied on a non-smooth function f. One can use in place of a gradient a sub-gradient, which defines the sub-gradient descent

$$x^{(\ell+1)} \stackrel{\text{def.}}{=} x^{(\ell)} - \tau_{\ell} q^{(\ell)} \quad \text{where} \quad q^{(\ell)} \in \partial f(x^{(\ell)}). \tag{13.14}$$

The main issue with this scheme is that to ensure convergence, the iterate should go to zero. One can easily convince oneself why by looking at the iterates on a function f(x) = |x|.

Theorem 37. If $\sum_{\ell} \tau_{\ell} = +\infty$ and $\sum_{\ell} \tau_{\ell}^{2} < +\infty$, then $x^{(\ell)}$ converges to a minimizer of f.

13.1.3 Projected Gradient Descent

We consider a generic constraint optimization problem as

$$\min_{x \in \mathcal{C}} f(x) \tag{13.15}$$

where $\mathcal{C} \subset \mathbb{R}^S$ is a closed convex set and $f : \mathbb{R}^S \to \mathbb{R}$ is a smooth convex function (at least of class \mathcal{C}^1).

The gradient descent algorithm (13.2) is generalized to solve a constrained problem using the projected gradient descent

$$x^{(\ell+1)} \stackrel{\text{def.}}{=} \operatorname{Proj}_{\mathcal{C}} \left(x^{(\ell)} - \tau_{\ell} \nabla f(x^{(\ell)}) \right), \tag{13.16}$$

where $\operatorname{Proj}_{\mathcal{C}}$ is the orthogonal projector on \mathcal{C}

$$\operatorname{Proj}_{\mathcal{C}}(x) = \underset{x' \in \mathcal{C}}{\operatorname{argmin}} \|x - x'\|$$

which is always uniquely defined because \mathcal{C} is closed and convex. The following proposition shows that all the convergence properties of the classical gradient descent caries over to this projected algorithm.

Theorem 38. Theorems 35 and 36 still holds when replacing iterations (13.2) by (13.16).

Proof. The proof of Theorem 35 extends because the projector is contractant, $\|\operatorname{Proj}_{\mathcal{C}}(x) - \operatorname{Proj}_{\mathcal{C}}(x')\| \le \|x - x'\|$ so that the strict contraction properties of the gradient descent is maintained by this projection.

The main bottleneck that often prevents to use (13.16) is that the projector is often complicated to compute. We are however lucky since for ℓ^1 minimization, one can apply in a straightforward manner this method.

13.2 Proximal Algorithm

For non-smooth functions f, it is not possible to perform an "explicit" gradient descent step because the gradient is not even defined. One thus needs to replace this "explicit" step by an "implicit" one, which is possible even if f is non-smooth.

13.2.1 Proximal Map

The implicit stepping of amplitude $\tau > 0$ is defined as

$$\forall x, \quad \text{Prox}_{\tau f}(x) \stackrel{\text{def.}}{=} \underset{x'}{\operatorname{argmin}} \frac{1}{2} ||x - x'||^2 + f(x').$$
 (13.17)

It amounts to minimize function f locally around x, in a ball of radius controlled by τ . This the involved function $\frac{1}{2}||x-\cdot||^2+f$ is strongly convex, this operator $\operatorname{Prox}_{\tau f}$ is well defined and single-valued.

When $f = \iota_{\mathcal{C}}$ is an indicator, the proximal map boils down to a projection $\operatorname{Prox}_{\iota_{\mathcal{C}}} = \operatorname{Proj}_{\mathcal{C}}$, it is thus in some sense a generalization of the projection to arbitrary function. And can also be interpreted as a projector on a level set of f. An interesting feature of the proximal map is that it is a contraction, thus generalizing the well-known property of projectors.

Proposition 34. One has $\|\operatorname{prox}_f(x) - \operatorname{prox}_f(y)\| \leq \|x - y\|$.

Examples The following proposition states a few simples examples.

Proposition 35. One has

$$\operatorname{Prox}_{\frac{\tau}{2}\|\cdot\|^{2}}(x) = \frac{x}{1+\tau}, \quad and \quad \operatorname{Prox}_{\tau\|\cdot\|_{1}} = \mathcal{S}_{\tau}^{1}(x), \tag{13.18}$$

where the soft-thresholding is defined as

$$S_{\tau}^{1}(x) \stackrel{\text{def.}}{=} (S_{\tau}(x_{i}))_{i=1}^{N} \quad where \quad S_{\tau}(r) \stackrel{\text{def.}}{=} \operatorname{sign}(r)(|r| - \lambda)_{+},$$

(see also (11.5)). For $A \in \mathbb{R}^{P \times N}$, one has

$$\operatorname{Prox}_{\frac{\tau}{2}||A - y||^2} = (\operatorname{Id}_P + \tau A^* A)^{-1} A^* = A^* (\operatorname{Id}_N + \tau A A^*)^{-1}$$
(13.19)

Proof. The proximal map of $\|\cdot\|_1$ was derived in Proposition 24. [ToDo: for the quadratic case]

Note that in some case, the proximal map of a non-convex function is well defined, for instance $\operatorname{Prox}_{\tau\|\cdot\|_0}$ is the hard thresholding associated to the threshold $\sqrt{2\tau}$, see Proposition 24.

13.2.2 Basic Properties

We recap some useful proximal-calculus.

Proposition 36. One has ToDo: ??

$$\operatorname{Prox}_{f+\langle y,\cdot\rangle} = \operatorname{Prox}_f(\cdot - y), \quad \operatorname{Prox}_{\lambda f} = \operatorname{Prox}_f(\cdot / \lambda), \operatorname{Prox}_{f(\cdot - y)} = y + \operatorname{Prox}_f(\cdot - y),$$

If $f(x) = \sum_{k=1}^{K} f(x_k)$ for $x = (x_1, \dots, x_K)$ is separable, then

$$Prox_{\tau f}(x) = (Prox_{\tau f_k}(x_k))_{k=1}^K.$$
(13.20)

The following proposition is very useful.

Proposition 37. If $A \in \mathbb{R}^{P \times N}$ is a tight frame, i.e. $AA^* = \mathrm{Id}_P$, then

$$\operatorname{Prox}_{f \circ A} = A^* \circ \operatorname{Prox}_f \circ A + \operatorname{Id}_N - A^* A.$$

In particular, if A is orthogonal, then $\operatorname{Prox}_{f \circ A} = A^* \circ \operatorname{Prox}_f \circ A$.

13.2.3 Related Concepts

Link with sub-differential. For a set-valued map $U: \mathcal{H} \hookrightarrow \mathcal{G}$, we define the inverse set-valued map $U^{-1}: \mathcal{G} \hookrightarrow \mathcal{H}$ by

$$h \in U^{-1}(g) \iff g \in U(h)$$
 (13.21)

[ToDo: add picture] The following proposition shows that the proximal map is related to a regularized inverse of the sub-differential.

Proposition 38. One has $Prox_{\tau f} = (Id + \tau \partial f)^{-1}$.

Proof. One has the following equivalence

$$z = \operatorname{Prox}_{\tau f}(x) \Leftrightarrow 0 \in z - x + \tau \partial f(z) \Leftrightarrow x \in (\operatorname{Id} + \tau \partial f)(z) \Leftrightarrow z = (\operatorname{Id} + \tau \partial f)^{-1}(x)$$

where for the last equivalence, we have replace " \in " by "=" because the proximal map is single valued. \square

The proximal operator is hence often referred to the "resolvent" $\operatorname{Prox}_{\tau f} = (\operatorname{Id} + \tau \partial f)^{-1}$ of the maximal monotone operator ∂f .

Link with duality. One has the following fundamental relation between the proximal operator of a function and of its Legendre-Fenchel transform

Theorem 39 (Moreau decomposition). One has

$$\operatorname{Prox}_{\tau f} = \operatorname{Id} - \tau \operatorname{Prox}_{f^*/\tau}(\cdot/\tau).$$

This theorem shows that the proximal operator of f is simple to compute if and only the proximal operator of f^* is also simple. As a particular instantiation, since according to , one can re-write the soft thresholding as follow

$$\operatorname{Prox}_{\tau \|\cdot\|_1}(x) = x - \tau \operatorname{Proj}_{\|\cdot\|_{\infty} \leq 1}(x/\tau) = x - \operatorname{Proj}_{\|\cdot\|_{\infty} \leq \tau}(x) \quad \text{where} \quad \operatorname{Proj}_{\|\cdot\|_{\infty} \leq \tau}(x) = \min(\max(x, -\tau), \tau).$$

In the special case where $f = \iota_{\mathcal{C}}$ where \mathcal{C} is a closed convex cone, then

$$(\iota_{\mathcal{C}})^* = \iota_{\mathcal{C}^{\circ}} \quad \text{where} \quad \mathcal{C}^{\circ} \stackrel{\text{def.}}{=} \{ y \; ; \; \forall \, x \in \mathcal{C}, \langle x, \, y \rangle \leqslant 0 \}$$
 (13.22)

and \mathcal{C}° is the so-called polar cone. Cones are fundament object in convex optimization because they are invariant by duality, in the sense of (13.22) (if \mathcal{C} is not a cone, its Legendre transform would not be an indicator). Using (13.22), one obtains the celebrated Moreau polar decomposition

$$x = \operatorname{Proj}_{\mathcal{C}}(x) +^{\perp} \operatorname{Proj}_{\mathcal{C}^{\circ}}(x)$$

where " $+^{\perp}$ " denotes an orthogonal sum (the terms in the sum are mutually orthogonal). [ToDo: add drawing] In the case where $\mathcal{C} = V$ is a linear space, this corresponds to the usual decomposition $\mathbb{R}^N = V \oplus^{\perp} V^{\perp}$.

Link with Moreau-Yosida regularization. The following proposition shows that the proximal operator can be interpreted as performing a gradient descent step on the Moreau-Yosida smoothed version f_{μ} of f, defined in (12.11).

Proposition 39. One has

$$\operatorname{Prox}_{\mu f} = \operatorname{Id} - \mu \nabla f_{\mu}.$$

13.3 Primal Algorithms

We now describe some important algorithm which assumes some structure (a so-called "splitting") of the minimized functional to be able to apply proximal maps on sub-functions. Note that there is obviously many ways to split or structure a given initial problem, so there are many non-equivalent ways to apply a given proximal-based method to solve the problem. Finding the "best" way to split a problem is a bit like black magic, and there is no definite answer. Also all there algorithm comes with step size and related parameters, and there is no obvious way to tune these parameters automatically (although some insight might be gained by studying convergence rate).

13.3.1 Proximal Point Algorithm

One has the following equivalence

$$x^* \in \operatorname{argmin} f \iff 0 \in \partial f(x^*) \iff x^* \in (\operatorname{Id} + \tau \partial f)(x^*)$$
 (13.23)

$$\Leftrightarrow x^* = (\mathrm{Id} + \tau \partial f)^{-1}(x^*) = \mathrm{Prox}_{\tau f}(x^*). \tag{13.24}$$

This shows that being a minimizer of f is equivalent to being a fixed point of $\operatorname{Prox}_{\tau f}$. This suggest the following fixed point iterations, which are called the proximal point algorithm

$$x^{(\ell+1)} \stackrel{\text{\tiny def.}}{=} \operatorname{Prox}_{\tau_{\ell} f}(x^{(\ell)}). \tag{13.25}$$

On contrast to the gradient descent fixed point scheme, the proximal point method is converging for any sequence of steps.

Theorem 40. If $0 < \tau_{\min} \leqslant \tau_{\ell} \leqslant \gamma_{\max} < +\infty$, then $x^{(\ell)} \to x^*$ a minimizer of f.

This implicit step (13.25) should be compared with a gradient descent step (13.2)

$$x^{(\ell+1)} \stackrel{\text{def.}}{=} (\mathrm{Id} + \tau_{\ell} \nabla f)(x^{(\ell)}).$$

One sees that the implicit resolvent $(\mathrm{Id} - \tau_{\ell} \partial f)^{-1}$ replaces the explicit step $\mathrm{Id} + \tau_{\ell} \nabla f$. For small τ_{ℓ} and smooth f, they are equivalent at first order. But the implicit step is well defined even for non-smooth function, and the scheme (the proximal point) is always convergent (whereas the explicit step size should be small enough for the gradient descent to converge). This is inline with the general idea the implicit stepping (e.g. implicit Euler for integrating ODE, which is very similar to the proximal point method) is more stable. Of course, the drawback is that explicit step are very easy to implement whereas in general proximal map are hard to solve (most of the time as hard as solving the initial problem).

13.3.2 Forward-Backward

It is in general impossible to compute $\operatorname{Prox}_{\gamma f}$ so that the proximal point algorithm is not implementable. In oder to derive more practical algorithms, it is important to restrict the class of considered function, by imposing some structure on the function to be minimized. We consider functions of the form

$$\min_{x} \mathcal{E}(x) \stackrel{\text{def.}}{=} f(x) + g(x) \tag{13.26}$$

where $g \in \Gamma_0(\mathcal{H})$ can be an arbitrary, but f needs to be smooth.

One can modify the fixe point derivation (13.23) to account for this special structure

$$x^* \in \operatorname{argmin} f + g \quad \Leftrightarrow \quad 0 \in \nabla f(x^*) + \partial g(x^*) \quad \Leftrightarrow \quad x^* - \tau \nabla f(x^*) \in (\operatorname{Id} + \tau \partial g)(x^*)$$

 $\Leftrightarrow \quad x^* = (\operatorname{Id} + \tau \partial g)^{-1} \circ (\operatorname{Id} - \tau \nabla f)(x^*).$

This fixed point suggests the following algorithm, with the celebrated Forward-Backward

$$x^{(\ell+1)} \stackrel{\text{def.}}{=} \operatorname{Prox}_{\tau_{\ell}g} \left(x^{(\ell)} - \tau_{\ell} \nabla f(x^{(\ell)}) \right). \tag{13.27}$$

Derivation using surrogate functionals. An intuitive way to derive this algorithm, and also a way to prove its convergence, it using the concept of surrogate functional.

To derive an iterative algorithm, we modify the energy $\mathcal{E}(x)$ to obtain a surrogate functional $\mathcal{E}(x, x^{(\ell)})$ whose minimization corresponds to a simpler optimization problem, and define the iterations as

$$x^{(\ell+1)} \stackrel{\text{def.}}{=} \underset{\sim}{\operatorname{argmin}} \mathcal{E}(x, x^{(\ell)}).$$
 (13.28)

In order to ensure convergence, this function should satisfy the following property

$$\mathcal{E}(x) \leqslant \mathcal{E}(x, x')$$
 and $\mathcal{E}(x, x) = \mathcal{E}(x)$ (13.29)

and $\mathcal{E}(x) - \mathcal{E}(x, x')$ should be a smooth function. Property (13.29) guarantees that f is decaying by the iterations

$$\mathcal{E}(x^{(\ell+1)}) \leqslant \mathcal{E}(x^{(\ell)})$$

and it simple to check that actually all accumulation points of $(x^{(\ell)})_{\ell}$ are stationary points of f.

In order to derive a valid surrogate $\mathcal{E}(x, x')$ for our functional (13.26), since we assume f is L-smooth (i.e. satisfies (\mathcal{R}_L)), let us recall the quadratic majorant (13.3)

$$f(x) \le f(x') + \langle \nabla f(x'), x' - x \rangle + \frac{L}{2} ||x - x'||^2,$$

so that for $0 < \tau < \frac{1}{L}$, the function

$$\mathcal{E}(x, x') \stackrel{\text{def.}}{=} f(x') + \langle \nabla f(x'), x' - x \rangle + \frac{1}{2\tau} ||x - x'||^2 + g(x)$$
 (13.30)

satisfies the surrogate conditions (13.29). The following proposition shows that minimizing the surrogate functional corresponds to the computation of a so-called proximal operator.

Proposition 40. The update (13.28) for the surrogate (13.30) is exactly (13.27).

Proof. This follows from the fact that

$$\langle \nabla f(x'), x' - x \rangle + \frac{1}{2\tau} ||x - x'||^2 = \frac{1}{2\tau} ||x - (x' - \tau \nabla f(x'))||^2 + \text{cst.}$$

Convergence of FB. Although we impose $\tau < 1/L$ to ensure majorization property, one can actually show convergence under the same hypothesis as for the gradient descent, i.e. $0 < \tau < 2/L$, with the same convergence rates. This means that Theorem 38 for the projected gradient descent extend to FB.

Theorem 41. Theorems 35 and 36 still holds when replacing iterations (13.2) by (13.27).

Note furthermore that the projected gradient descent algorithm (13.16) is recovered as a special case of (13.27) when setting $J = \iota_{\mathcal{C}}$ the indicator of the constraint set, since $\operatorname{Prox}_{\rho J} = \operatorname{Proj}_{\mathcal{C}}$ in this case.

Of course the difficult point is to be able to compute in closed form $\text{Prox}_{\tau g}$ in (13.27), and this is usually possible only for very simple function. We have already seen such an example in Section 11.3.3 for the resolution of ℓ^1 -regularized inverse problems (the Lasso).

13.3.3 Douglas-Rachford

We consider here the structured minimization problem

$$\min_{x \in \mathbb{R}^N} f(x) + g(x), \tag{13.31}$$

but on contrary to the Forward-Backward setting studied in Section 13.3.2, no smoothness is imposed on f. We here suppose that we can compute easily the proximal map of f and g.

Example 7 (Constrained Lasso). An example of a problem of the form (13.31) where one can apply Douglas-Rachford is the noiseless constrained Lasso problem (11.11)

$$\min_{Ax=y} \|x\|_1$$

where one can use $f = \iota_{\mathcal{C}_y}$ where $\mathcal{C}_y \stackrel{\text{def.}}{=} \{x \; ; \; Ax = y\}$ and $g = \|\cdot\|_1$. As noted in Section 11.3.1, this problem is equivalent to a linear program. The proximal operator of g is the soft thresholding as stated in (13.18), while the proximal operator of g is the orthogonal projector on the affine space \mathcal{C}_y , which can be computed by solving a linear system as stated in (12.9) (this is especially convenient for inpainting problems or deconvolution problem where this is achieved efficiently).

The Douglas-Rachford iterations read

$$\tilde{x}^{(\ell+1)} \stackrel{\text{\tiny def.}}{=} \left(1 - \frac{\mu}{2}\right) \tilde{x}^{(\ell)} + \frac{\mu}{2} \operatorname{rProx}_{\tau g}(\operatorname{rProx}_{\tau f}(\tilde{x}^{(\ell)})) \quad \text{and} \quad x^{(\ell+1)} \stackrel{\text{\tiny def.}}{=} \operatorname{Prox}_{\tau f}(\tilde{x}^{(\ell+1)},)$$
(13.32)

We have used the following shortcuts

$$\operatorname{rProx}_{\tau f}(x) = 2 \operatorname{Prox}_{\tau f}(x) - x.$$

One can show that for any value of $\tau > 0$, any $0 < \mu < 2$, and any $\tilde{x}_0, x^{(\ell)} \to x^*$ which is a minimizer of f + g.

Note that it is of course possible to inter-change the roles of f and g, which defines another set of iterations.

More than two functions. Another sets of iterations can be obtained by "symetrizing" the algorithm. More generally, if we have K functions $(f_k)_k$, we re-write

$$\min_{x} \sum_{k} f_k(x) = \min_{X = (x_1, \dots, x_k)} f(X) + g(X) \quad \text{where} \quad f(X) = \sum_{k} f_k(x_k) \quad \text{and} \quad g(X) = \iota_{\Delta}(X)$$

where $\Delta = \{X ; x_1 = \ldots = x_k\}$ is the diagonal. The proximal operator of f is

$$\operatorname{Prox}_{\tau f}(X) = \operatorname{Proj}_{\Delta}(X) = (\bar{x}, \dots, \bar{x}) \text{ where } \bar{x} = \frac{1}{K} \sum_{k} x_k$$

while the proximal operator of f is easily computed from those of the $(f_k)_k$ using (13.20). One can thus apply DR iterations (13.32).

Handling a linear operator. One can handle a minimization of the form (13.34) by introducing extra variables

$$\inf_{x} f_{1}(x) + f_{2}(Ax) = \inf_{z=(x,y)} f(z) + g(z) \quad \text{where} \quad \left\{ \begin{array}{l} f(z) = f_{1}(x) + f_{2}(y) \\ g(z) = \iota_{\mathcal{C}}(x,y), \end{array} \right.$$

where $C = \{(x, y) ; Ax = y\}$. This problem can be handled using DR iterations (13.32), since the proximal operator of f is obtained from those of (f_1, f_2) using (13.20), while the proximal operator of g is the projector on C, which can be computed in two alternative ways as the following proposition shows.

Proposition 41. One has

$$\operatorname{Proj}_{\mathcal{C}}(x,y) = (x + A^* \tilde{y}, y - \tilde{y}) = (\tilde{x}, A\tilde{x}) \quad \text{where} \quad \begin{cases} \tilde{y} \stackrel{\text{def.}}{=} (\operatorname{Id}_P + AA^*)^{-1} (Ax - y) \\ \tilde{x} \stackrel{\text{def.}}{=} (\operatorname{Id}_N + A^*A)^{-1} (A^*y + x). \end{cases}$$
(13.33)

Remark 7 (Inversion of linear operator). At many places (typically to compute some sort of projector) one has to invert matrices of the form AA^* , A^*A , $\operatorname{Id}_P + AA^*$ or $\operatorname{Id}_N + A^*A$ (see for instance (13.33)). There are some case where this can be done efficiently. Typical examples where this is simple are inpainting inverse problem where AA^* is diagonal, and deconvolution or partial Fourier measurement (e.g. fMRI) for which A^*A is diagonalized using the FFT. If this inversion is too costly, one needs to use more advanced methods, based on duality, which allows to avoid trading the inverse A by the application of A^* . They are however typically converging more slowly.

13.4 Dual and Primal-Dual Algorithms

Convex duality, detailed in Section 12.2 (either from the Lagrange or the Fenchel-Rockafeller point of view – which are essentially equivalent), is very fruitful to derive new optimization algorithm or to apply existing algorithm on a dual reformulation.

13.4.1 Forward-backward on the Dual

Let us illustrate first the idea of applying a known algorithm to the dual problem. We consider here the structured minimization problem associated to Fenchel-Rockafeller duality (12.12)

$$p^* = \inf_{x} f(x) + g(Ax),$$
 (13.34)

but furthermore assume that f is μ -strongly convex, and we assume for simplicity that both (f,g) are continuous. If f were also smooth (but it needs to be!), one could think about using the Forward-Backward algorithm (13.27). But the main issue is that in general $\operatorname{Prox}_{\tau g \circ A}$ cannot be computed easily even if one can compute $\operatorname{Prox}_{\tau g \circ A}$. An exception to this is when A is a tight frame, as exposed in Proposition 37, but in practice it is rarely the case.

Example 8 (TV denoising). A typical example, which was the one used by Antonin Chambolle [9] to develop this class of method, is the total variation denoising

$$\min_{x} \frac{1}{2} \|y - x\|^2 + \lambda \|\nabla x\|_{1,2}$$

where $\nabla x \in \mathbb{R}^{N \times d}$ is the gradient (a vector field) of a signal (d=1) or image (d=2) x, and $\|\cdot\|_{1,2}$ is the vectorial- ℓ^1 norm (also called $\ell^1 - \ell^2$ norm), such that for a d-dimensional vector field $(v_i)_{i=1}^N$

$$||v||_{1,2} \stackrel{\text{def.}}{=} \sum_{i} ||v_i||.$$

Here

$$f = \frac{1}{2}\|\cdot -y\|^2 \quad \text{and} \quad g = \lambda\|\cdot\|_{1,2}$$

so that f is $\mu = 1$ strongly convex, and one sets $A = \nabla$ the linear operator.

Applying Fenchel-Rockafeller Theorem 34 (since strong duality holds, all involved functions being continuous), one has that

$$p^* = \sup_{u} -g^*(u) - f^*(-A^*u).$$

But more importantly, since f is μ -strongly convex, one has that f^* is smooth with a $1/\mu$ -Lipschitz gradient. One can thus use the Forward-Backward algorithm (13.27) on (minus the energy of) this problem, which reads

$$u^{(\ell+1)} = \text{Prox}_{\tau_k g^*} \left(u^{(\ell)} + \tau_k A \nabla f^* (A^* u^{(\ell)}) \right).$$

To guarantee convergence, the step size τ_k should be smaller than 2/L where L is the Lipschitz constant of $A \circ \nabla f^* \circ A^*$, which is smaller than $||A||^2/\mu$.

Last but not least, one some (not necessarily unique) dual minimizer u^* is computed, the primal-dual relationships (12.17) ensures that one retrieves the unique primal minimizer x^* as

$$-A^*u^* \in \partial f(x^*) \quad \Leftrightarrow \quad x^* \in (\partial f)^{-1}(-A^*u^*) = \partial f^*(-A^*u^*) \quad \Leftrightarrow \quad x^* = \nabla f^*(-A^*u^*)$$

where we used here the crucial fact that f^* is smooth.

Example 9 (TV denoising). In the particular case of the TV denoising problem, one has

$$g^* = \iota_{\|\cdot\|_{\infty,2} \leqslant \lambda} \quad \text{where} \quad \|v\|_{\infty,2} \stackrel{\text{def.}}{=} \max_i \|v_i\| \quad \Longrightarrow \quad \operatorname{Prox}_{\tau g^*}(u) = \left(\min(\|v_i\|, \lambda) \frac{v_i}{\|v_i\|}\right)$$
$$f^*(h) = \frac{1}{2} \|h\|^2 + \langle h, y \rangle \quad \text{and} \quad \nabla f^*(h) = h + y.$$

Furthermore, $\mu = 1$ and $A^*A = \Delta$ is the usual finite difference approximation of the Laplacian, so that $||A||^2 = ||\Delta|| = 4d$ where d is the dimension.

13.4.2 Primal-Dual Splitting

We now comeback to the more general structure problem of the form (13.34), which we consider in primal-dual form as

$$\inf_{x} f(x) + g(Ax) = \sup_{x} \inf_{x} f(x) + \langle Ax, u \rangle - g^{*}(u), \tag{13.35}$$

but we do not suppose anymore that f is strongly convex.

A typical instance of such a problem is for the TV regularization of the inverse problem $\mathcal{K}x = y$, which corresponds to solving

$$\min_{x} \frac{1}{2} \|y - \mathcal{K}x\|^2 + \lambda \|\nabla x\|_{1,2}.$$

where $A = \nabla$, $f(x) = \frac{1}{2} \|y - \mathcal{K} \cdot \|^2$ and $g = \lambda \| \cdot \|_{1,2}$. Note however that with such a splitting, one will have to compute the proximal operator of f, which, following (13.19), requires inverting either $\mathrm{Id}_P + AA^*$ or $\mathrm{Id}_N + A^*A$, see Remark 7.

A standard primal-dual algorithm, which is detailed in [], reads

$$z^{(\ell+1)} \stackrel{\text{def.}}{=} \operatorname{Prox}_{\sigma g^*}(z^{(\ell)} + \sigma A(\tilde{x}^{(\ell)})$$

$$x^{(\ell+1)} \stackrel{\text{def.}}{=} \operatorname{Prox}_{\tau f}(x^{(\ell)} - \tau A^*(z^{(\ell+1)}))$$

$$\tilde{x}^{(\ell)} \stackrel{\text{def.}}{=} x^{(\ell+1)} + \theta(x^{(\ell+1)} - x^{(\ell)})$$

if $0 \leqslant \theta \leqslant 1$ and $\sigma \tau \|K\|^2 < 1$, then $x^{(\ell)}$ converges to a minimizer of (13.35) .

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