

When ClickHouse marries graph processing

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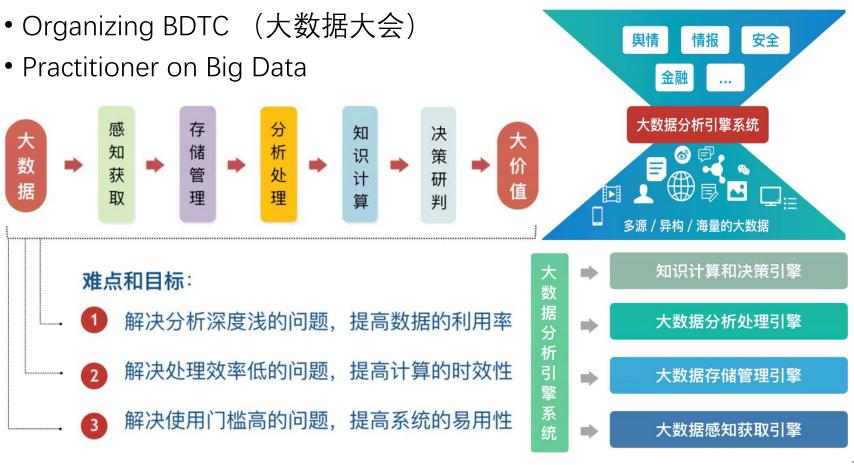
### About Me

- Active ClickHouse Contributor
  - ~70 valid PRs
  - ~40 Stack Overflow Answers
- Open Source Enthusiast (Hacked 20+ Projects)
  - DB: Impala, Greenplum, Cockroach, Citus, Kudu
  - Misc: emacs, tmux, gdb, fish-shell, tdesktop ...
- SQLGraph Author



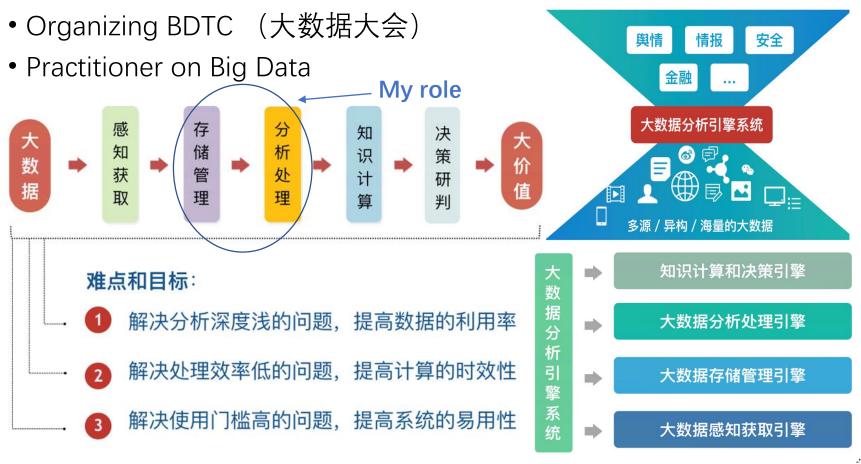
### About My Lab

• CCF Task Force on Big Data (大数据专家委员会)

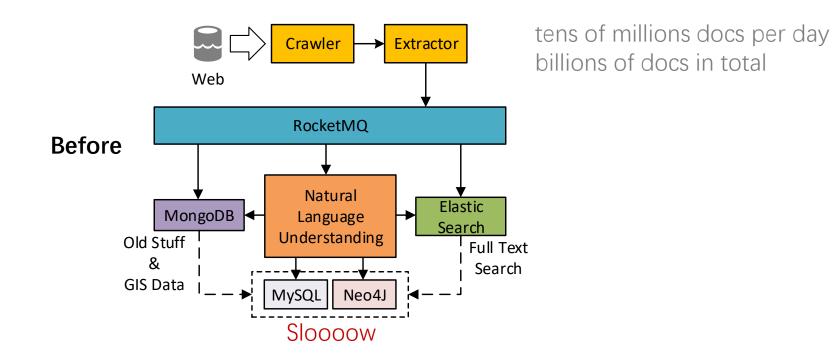


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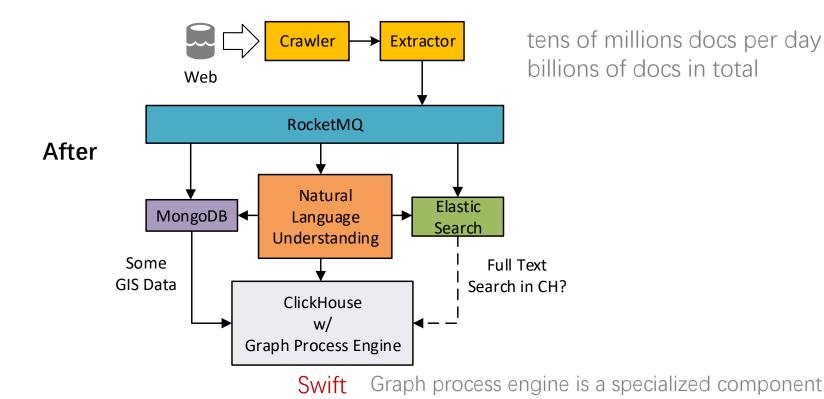


### How We Use ClickHouse



The NLU module outputs well structured, clean, (almost) immutable data.

### How We Use ClickHouse



to run analytical graph algorithms efficiently.

The NLU module outputs well structured, clean, (almost) immutable data.

## Why Graph Processing Engine

- Finding Important Nodes and Edges
  - PageRank
  - Personalized PageRank
  - Shortest Path



- Label Propagation
- Connection Components
- Louvain



- Collaborative Filtering
- Random Walk
- Neural Embeddings







## Why Graph Processing Engine in RDBMS

Hidden graph structures all over the place



- Graph traversal via joins is both slow and inconvenient
- Almost impossible to do iterative algorithms (recursion)
- Cumbersome graph extraction process if acting as a graph data source
- Powerful in data management and processing
  - Full fledged processing (Instead of awk, perl ...)
  - Mature concepts of metadata handling
  - Widely adopted and ease of use interface



Mining what exists, aiding what emerges

## Why Graph Processing Engine in ClickHouse

- From Architect's Perspective
  - Fast data storage and fast data processing
  - Can handle many data sources
  - Rich user interface
- From Developer's Perspective
  - Nice building blocks for performance critical applications
  - Versatile SQL pipelines for mixed SQL/graph processing
  - Good code quality

Think OLAP but in the graph field.

## What to Expect from an OLAP Graph DB

#### General Users

- Easy to use (think as a graph)
- Easy to understand (view as a graph)
- Faster than Non-Graph databases or OLTP graph databases

#### Graph Experts

- Extremely fast analysis
- Interactive graph algorithm designing
- Ability to do low level optimizations

Be efficient to end users and computer hardware.

### System Design Goals

### **Usability**

- 1. Visualization
- 2. Graphs Being First-Class Citizens
- 3. SQL accessibility

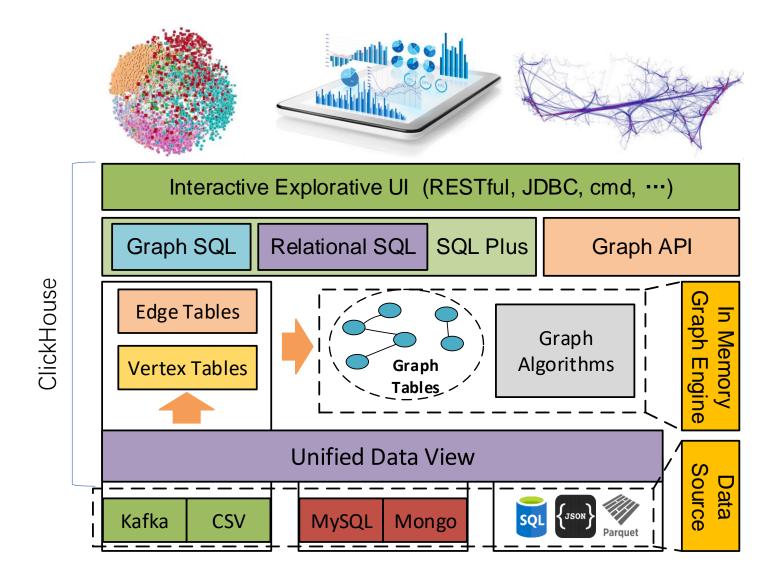
#### Versatile

- 1. Graph Algorithm SDK
- 2. Graph Inspector and Rich Operators
- 3. Reusable Graph Modeling

### **Performance**

- 1. Outperform graph compute engines
- 2. Quick development iterations
- 3. Fast graph data digestion

## SQLGraph



### Some Results

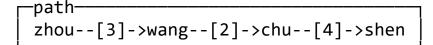
#### Calculate PageRank value per person

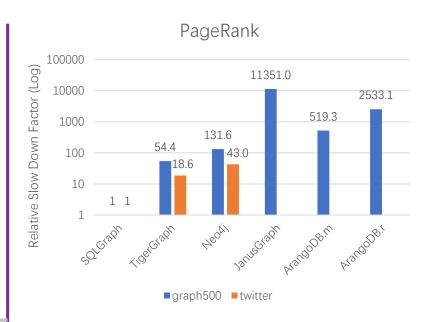
SELECT vp(v, 'name') AS name, pagerank FROM pagerank(wz)
ORDER BY pagerank DESC LIMIT 5

r-name-	<sub> </sub> pagerank-
li	0.06964007
zhao	0.063854694
qian	0.06365149
sun	0.06347877
shen	0.019985389

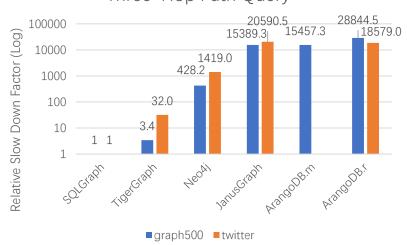
#### Find a longest path which ends at 'shen'

SELECT pp(p, 'name', 'w') AS path
FROM edgePath(wz)
WHERE vp(p[-1].2, 'name') = 'shen'
ORDER BY length(path) DESC
LIMIT 1





#### Three-Hop Path Query



# Let's Dive in

examples and implementations

## Graph Query Interface

- Graph Algorithms are table functions
- Graphs are special tables that used as the first argument
- Collection of functions to retrieve graph info
- Example: Get the top 5 pagerank value from graph "wz"

## Graph Query Output

- Three kinds of outputs : vertices, edges or graph info
- Vertices: PageRank, CC, BC, Radii, etc.
  - Has 'v : Vertices' column in the output list
  - Cover all vertices, no duplications
  - Can be inserted back into the graph as a global vertex property

```
SELECT
id,
groupArray(vp(v, 'name')) AS names

FROM cc(wz)

GROUP BY id

id—names
0 ['zhao','qian','sun','li']
4 ['zhou','wu','zheng','wang','feng','chen','chu','wei','shen']
14 ['han','yang']
12 ['jiang']
```

## Graph Query Output

- Three kinds of outputs : vertices, edges or graph info
- Edges: Hop, SSSP, CommonNeighbors, LinkCircle, etc.
  - Has 'e: Edges' column in the output list
  - Support function combinators: Out, In, All, Path
  - Can be used to derive subgraphs: create graph as select ...
  - Has 'Graph' format to support Non-Structural output (for visualization)

## Graph Query Output

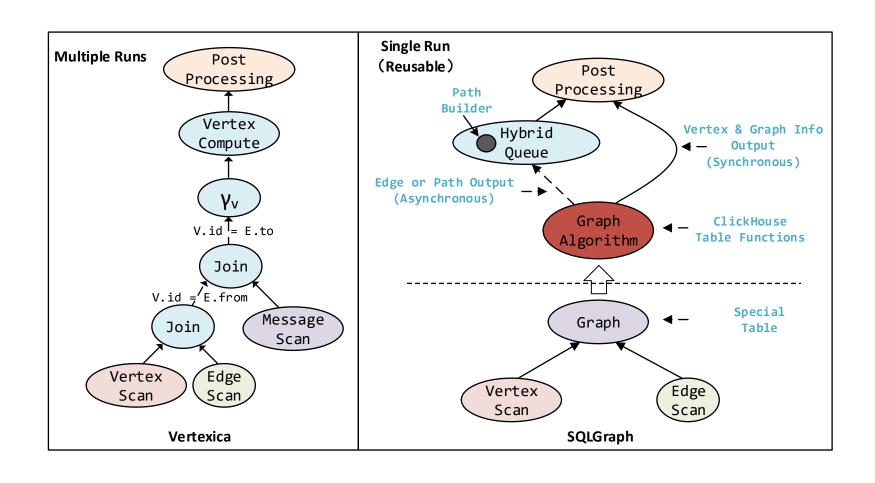
- Three kinds of outputs: vertices, edges or graph info
- Graph info: TriangleCount, MaxClique, ClusteringCoef, etc.
  - Results are graph characters

```
SELECT *
FROM triangle(wz)

-count-
2
```

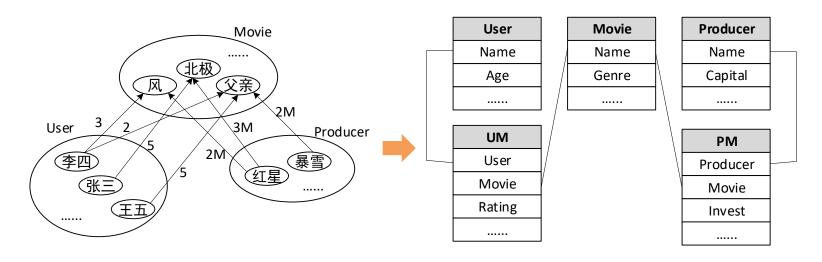
Mainly for graph mining algorithms

## Graph Query Output Pipeline (Compared)



### How to Get/Build a Graph

#### What does a graph look like in RDBMS?



#### How can we make the graph structure efficient?

Algorithm System	PageRank (5 iters)	Components
Ligra	6.69s	6.75s
Neo4j	131s	189s

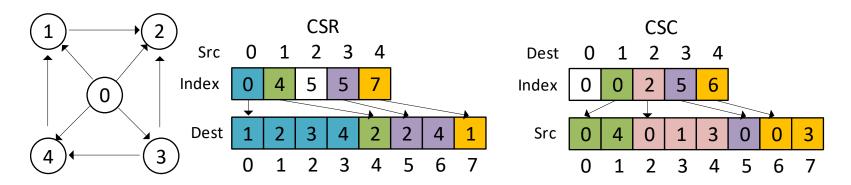
Table 1: Performance Comparison (Twitter-2010)

We need CSR/CSC to be fast!

CSR/C: Compressed Sparse Row/Column

### CSR/C Explained

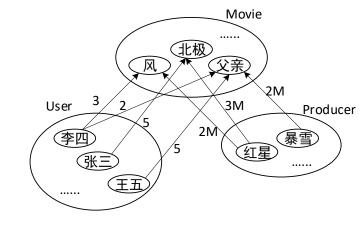
#### Compressed sparse storage



Similar to Arrays in ClickHouse

#### Vanilla edge list storage

Easy to store and modify (normal tables)



#### Define vertex tables to encode key columns

CREATE TABLE Producer (Name String key, Capital UInt64) ENGINE V; CREATE TABLE Movie (Name String key, Genre String) ENGINE V; CREATE TABLE User (Name String key, Age int) ENGINE V;

Key: HashMap
-> [0, P-1]

Property: Columns

Vertex Table: Producer

红星: 0 暴雪: 1 E.g.

Key: HashMap
-> [0, M-1]

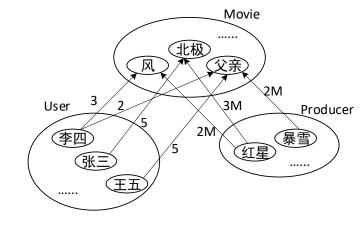
Property: Columns

Vertex Table: Movie

Movie
Name: Key
Genre
.....

风: 0 北极: 1 父亲: 2 E.g.

Key: HashMap<br/>-> [0, N-1]User<br/>Name: Key<br/>Age张三: 0<br/>李四: 1<br/>王五: 2Property: Columns......Vertex Table: User



#### Define edge tables to store relations

CREATE TABLE Producer\_Movie (Src VS(Producer), Dst VD(Movie), Invest UInt64) ENGINE E; CREATE TABLE User\_Movie (Src VS(User), Dst VD(Movie), Rating int) ENGINE E;

src: \_VS (Producer)
dst: \_VD (Movie)
order by <\_VS, \_VD>

Property: Columns

Producer_Movie		
Producer: VS		
Movie: VD		
Invest		

Edge Table: Producer\_Movie

红星-风: <0, 0> 红星-北极: <0, 1> 暴雪-父亲: <1, 1> E.g.

Property: Columns

src: VS (User)

dst: VD (Movie)

order by < VS, VD>

 User\_Movie
 张三-北极: <0, 1>

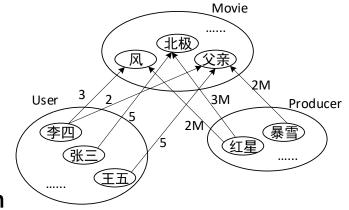
 User: VS
 李四-风: <1, 0>

 Movie: VD
 李四-父亲: <1, 2>

 Rating
 E.g.

Edge Table: User\_Movie

Storing hidden columns \_vs, \_vd and use 'MergeTree order by \_vs, \_vd' as the underlying engine



#### Define MetaSQL to specify how to build a graph

CREATE GRAPH User\_Movie\_Producer AS SELECT \* FROM edgeGroup(User\_Movie, Producer\_Movie)

#### MetaSQL Defines the following workflow:

- 1. Adjusting IDs of vertices and edges
- 2. Aggregating edge properties
- 3. Grouping edges (also assigning Eid)

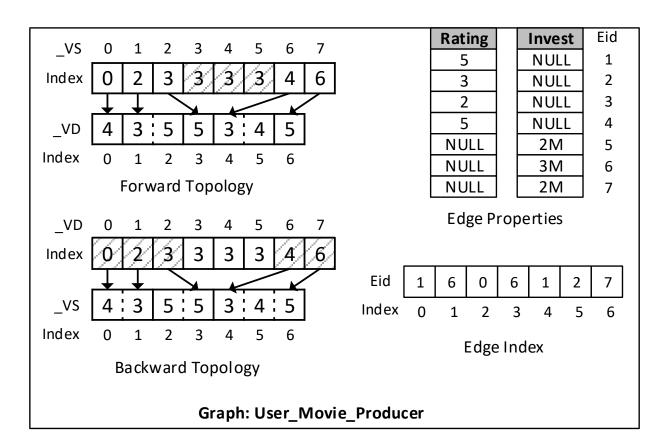
张三: 0	<u> </u>				<u>, , , , , , , , , , , , , , , , , , , </u>
孫二: 0 李四: 1	¦	User:	Movie:	Producer:	edge group
王五: 2		vertex num=N	vertex num=M	vertex num=P	\ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \
风: 3	Ш	key=Name	key=Name	key=Name	Vd
北极: 4	l'il	Property Handle	Property Handle	Property Handle	Rating
父亲: 5	ı'				i Invest
红星: 6	!		Vertex Table Pr	оху	, , , , , , , , , , , , , , , , , , ,
暴雪: 7	┞				- ''
E.g.	MetaSQL: User_Movie_Producer				

Eid 张三-北极: <0, 4> 1 李四-风: <1, 3> 2 李四-父亲: <1, 5> 3 王五-父亲: <2, 5> 4 红星-风: <6, 3> 5 红星-北极: <6, 4> 6 暴雪-父亲: <7, 5> 7 E.g.

#### 

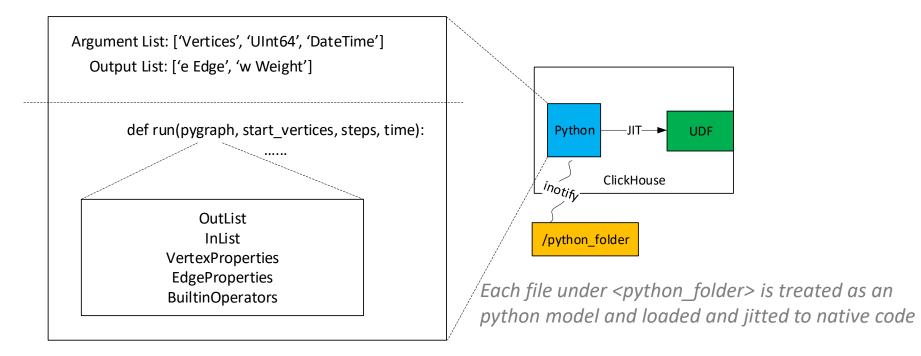
#### Execute MetaSQL to build a graph

REFRESH User\_Movie\_Producer [FULL]



### How to Write a Graph Algorithm

- Two ways of writing graph algorithms: python or c++
- Python: for ad-hoc queries, debugging and performance tuning



**Graph Algorithm NoteBook** 

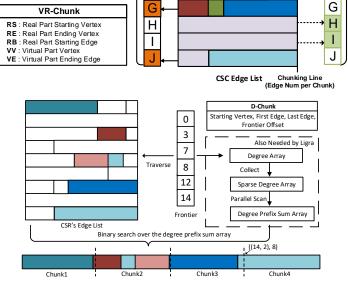
## How to Write a Graph Algorithm

Two ways of writing graph algorithms: python or c++

• C++: for long use/well known algorithms, efficiency

```
void PageRankFunction::run(Graph & graph) {
    ... // initialization code
Algorithm algo(
    [&](UInt32 s, UInt32 d) {// push
        atomicAdd(pr_new[d], pr[s]); },
    [&](UInt32* b, UInt32* e, UInt32 rd, UInt32 d) {// pull
    Float y = 0;
    while (b < e) y += pr[*b++];
    pr_new[d] = y; },
    [&](UInt32 rd, UInt32 d) {// pull reduce
        atomicAdd(pr_new[rd], pr_new[d]); }
);
while (!finish) {
    graph.run(algo);
    ... /* other related code */ }}</pre>
```

Listing 1. Page Rank Implementation



В

( subgraph )

Real Part

**Virtual Part** 

D

SilverChunk Framework

## Graph Query Helper

- vp, ep, pp for outputting related properties
- v for getting the internal id(s) of a vertex/vertices
- edge, vertex for retrieving edges/vertices of a graph respectively
- show graphs/algos for global inspections
- desc graph/algo for detail info
- graphInfo as a binary interface to access graph info

. . . . . .

## Graph Query Helper (E.g.)

#### DESCRIBE GRAPH wz

	r-name	<sub>—</sub> value or type—	<sub>T</sub> —comment—————
Edge properties:			
	region	String	
	time	Date	
	W	Float64	
Vantay gnount	l W	1 100004	
Vertex group:			
Vertex name:	default.node		
Key name:	id	UInt32	vertex id range : 0 15
Vertex properties:			
	id	UInt32	
	age	UInt32	
	name	String	
		8	
Meta info:			
	vertex_num	16	
	edge_num	16	
	symmetric	false	
	load	true	

## Graph Query Helper (E.g.)

#### DESCRIBE ALGORITHM pagerank

	<del></del> name <del></del>	$_{ extsf{T}}$ value or type $-$	<sub>T</sub> -comment
Arguments:			
		Vertices UInt64 Float64 Float64	vertices: start vertices (default all) iter: iterations (default 5), 0 for unlimited damping: damping value (default 0.85) epsilon: epsilon value (default 0.001)
Outputs:	v pagerank	Vertices Float32	

## Graph Query Helper (E.g.)

#### SHOW ALGOS

r-name	—type——	-comment-
bfs	python	
dijkstra	python	
vertexfilter	python	
simplebfs	builtin	
cc	builtin	
bellmanford	builtin	
coregroup	builtin	
linkcircle	builtin	
pagerank	builtin	
bc	builtin	
hop	builtin	
ssp	builtin	
shortcutcc	builtin	
personalrank	builtin	
triangle	builtin	
neighbor	builtin	
closec	builtin	
path	builtin	
commonneighbors	builtin	
radiiest	builtin	
	• • • • •	

### Limitations

Single Machine, Main Memory

#### But we have good reasons

- Graph size aren't that huge for us (around multiple billions of edges)
  - also commonly seen (referring Sahu, Siddhartha, et al PVLDB 17)
- Main memory is big, very big (new Mac Pro supports 1.5TB RAM)
- Very popular in graph computing community (Ligra, Galois, Grazelle, etc.)
- Nearly impossible to have independent partitions for a graph
  - Commodity networks are too slow, InfiniBand is too expensive
  - Scalability! But at what COST? (McSherry HotOS15)

scalable system	cores	twitter	uk-2007-05
GraphChi [12]	2	3160s	6972s
Stratosphere [8]	16	2250s	-
X-Stream [21]	16	1488s	-
Spark [10]	128	857s	1759s
Giraph [10]	128	596s	1235s
GraphLab [10]	128	249s	833s
GraphX [10]	128	419s	462s
Single thread (SSD)	1	300s	651s
Single thread (RAM)	1	275s	-

scalable system	cores	twitter	uk-2007-05
Stratosphere [8]	16	950s	-
X-Stream [21]	16	1159s	-
Spark [10]	128	1784s	$\geq 8000s$
Giraph [10]	128	200s	$\geq 8000s$
GraphLab [10]	128	242s	714s
GraphX [10]	128	251s	800s
Single thread (SSD)	1	153s	417s

### **Future Works**

- External Property Partitions
  - Sequential access to properties stored in external storages
- Partial/Incremental Graph Loading
  - Static knowledge database (huge graph) plus loadable domain data
- After the Processor branch Landed
  - Vectorized Graph Queries
  - Streaming and Time Series Processing

### Lessons Learned While Extending CH

- Understanding of code
  - semantic tools, utilization of debugger, reading unit tests
  - trial and error
- Isolation of code
  - new modules > new subclasses > Pimpl isolation > in-place patching
  - dlopen for optional features (e.g. python)
- Compilation (linking) is slooooow
  - shared build might help
  - better to build a libclickhouse.so and write main functions per feature
- Contribute while extending

# Thank You!

Questions?