

* Making Change Problem.

1 dollar = 100 cents

1 quarter = 25 cents

1 dime = 10 cents

1 nickel = 5 cents.

1 Pennies = 1 cent.

$$C[i, 0] = 0.$$

1. If $i=1$ then $C[i, j] = 1 + C[1, j - d_1]$

2. If $j < d_i$ then $C[i, j] = C[i-1, j]$

3. otherwise $C[i, j] = \min(C[i-1, j], 1 + C[i, j - d_i])$

Ex. Solve making change problem for $d_1 = 1, d_2 = 4, d_3 = 6$, $n=3$ and $N=8$ units.

Solⁿ \Rightarrow Here, $d_1 = 1, d_2 = 4, d_3 = 6, n=3, N=8$ units.

So, we have to create a table having 3 rows and column ranging from 0 to 8.

Initially $C[i, 0] = 0 \quad \therefore C[1, 0] = C[2, 0] = C[3, 0] = 0.$

		j \rightarrow								
		0	1	2	3	4	5	6	7	8
i=1	$d_1 = 1$	0	1	2	3	4	5	6	7	8
i=2	$d_2 = 4$	0	1	2	3	1	2	3	4	2
i=3	$d_3 = 6$	0	1	2	3	1	2	1	2	2

$\Rightarrow C[1, 1]$ with $i=1, j=1, d_1=1.$

As $i=1$

Formula used : $C[1, 1] = 1 + C[1, j - d_1]$

$$= 1 + C[1, 0]$$

$$= 1 + 0 = 1$$

$$\therefore C[1, 1] = 1$$

$\Rightarrow C[1,2]$ with $i=1, j=2, d_1=1$.

As $j=1$

$$C[1,2] = 1 + C[1, j-d_1]$$

$$= 1 + C[1, 1]$$

$$= 1 + 1 = 2$$

$$\boxed{\therefore C[1,2] = 2}$$

$\Rightarrow C[1,3]$ with $i=1, j=3, d_1=1$

As $i=1$

$$C[1,3] = 1 + C[1, j-d_1]$$

$$= 1 + C[1, 2]$$

$$= 1 + 2 = 3$$

$$\boxed{\therefore C[1,3] = 3}$$

$\Rightarrow C[1,4]$ with $i=1, j=4, d_1=1$.

As $i=1$

$$C[1,4] = 1 + C[1, j-d_1]$$

$$= 1 + C[1, 3]$$

$$= 1 + 3 = 4$$

$$\boxed{\therefore C[1,4] = 4}$$

$\Rightarrow C[1,5]$ with $i=1, j=5, d_1=1$.

As $i=1$

$$C[1,5] = 1 + C[1, j-d_1]$$

$$= 1 + C[1, 4]$$

$$= 1 + 4 = 5$$

$$\boxed{\therefore C[1,5] = 5}$$

$\Rightarrow C[1,6]$ with $i=1, j=6, d_1=1$

As $i=1$

$$C[1,6] = 1 + C[1, j-d_1]$$

$$= 1 + C[1, 5]$$

$$= 1 + 5 = 6$$

$$\boxed{\therefore C[1,6] = 6}$$

$\Rightarrow C[1,7]$ with $i=1, j=7, d_1=1$.

As $i=1$

$$\begin{aligned} C[1,7] &= 1 + C[1, j-d_1] \\ &= 1 + C[1, 6] \\ &= 1 + 6 = 7 \end{aligned}$$

$$\therefore C[1,7] = 7$$

$\Rightarrow C[1,8]$ with $i=1, j=8, d_1=1$.

As $i=1$

$$\begin{aligned} C[1,8] &= 1 + C[1, j-d_1] \\ &= 1 + C[1, 7] \\ &= 1 + 7 = 8 \end{aligned}$$

$$\therefore C[1,8] = 8$$

- Now, let's move on to the second row.

$\Rightarrow C[2,1]$ with $i=2, j=1, d_2=4$

As $j < d_2$ i.e. $1 < 4$

Formula used:

$$\begin{aligned} C[2,1] &= C[i-1, j] \\ &= C[1,1] \\ &= 1 \end{aligned}$$

$$\therefore C[2,1] = 1$$

$\Rightarrow C[2,2]$ with $i=2, j=2, d_2=4$.

As $j < d_2$ i.e. $2 < 4$.

$$\begin{aligned} C[2,2] &= C[i-1, j] \\ &= C[1,2] \\ &= 2 \end{aligned}$$

$$\therefore C[2,2] = 2$$

$\Rightarrow C[2,3]$ with $i=2, j=3, d_2=4$.

As $j < d_2$ i.e. $3 < 4$,

$$C[2,3] = C[i-1, j]$$

$$= C[1,3]$$

$$= 3.$$

$$\therefore C[2,3] = 3$$

$\Rightarrow C[2,4]$ with $i=2, j=4, d_2=4$.

As $i \neq 1$ and $j \neq d_2$.

$$C[2,4] = \min(C[i-1, j], 1 + C[i, j-d_i])$$

$$= \min(C[1,4], 1 + C[2,0])$$

$$= \min(4, 1 + 0)$$

$$= \min(4, 1) = 1.$$

$$\therefore C[2,4] = 1$$

$\Rightarrow C[2,5]$ with $i=2, j=5$ and $d_2=4$.

As $i \neq 1$ and $j > d_2$

$$C[2,5] = \min(C[i-1, j], 1 + C[i, j-d_i])$$

$$= \min(C[1,5], 1 + C[2,1])$$

$$= \min(5, 1 + 1)$$

$$= \min(5, 2) = 2.$$

$$\therefore C[2,5] = 2$$

$\Rightarrow C[2,6]$ with $i=2, j=6, d_2=4$.

As $i \neq 1$ and $j > d_2$

$$C[2,6] = \min(C[i-1, j], 1 + C[i, j-d_i])$$

$$= \min(C[1,6], 1 + C[2,2])$$

$$= \min(6, 1 + 2)$$

$$= \min(6, 3) = 3$$

$$\therefore C[2,6] = 3$$

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$\Rightarrow C[2,7]$ with $i=2, j=7, d_2=4$.

As $i \neq 1$ and $j > d_2$ i.e. $7 > 4$,

$$\begin{aligned} C[2,7] &= \min(C[i-1, j], 1 + C[i, j-d_i]) \\ &= \min(C[1,7], 1 + C[2,3]) \\ &= \min(7, 1+3) \\ &= \min(7, 4) = 4. \end{aligned}$$

$$\therefore C[2,7] = 4$$

$\Rightarrow C[2,8]$ with $i=2, j=8, d_2=4$.

As $i \neq 1$ and $j > d_2$

$$\begin{aligned} C[2,8] &= \min(C[i-1, j], 1 + C[i, j-d_i]) \\ &= \min(C[1,8], 1 + C[2,4]) \\ &= \min(8, 1+1) \\ &= \min(8, 2) = 2. \end{aligned}$$

$$\therefore C[2,8] = 2$$

- Now, let's complete third row.

$\Rightarrow C[3,1]$ with $i=3, j=1, d_3=6$.

As $i \neq 1$ and $j < d_3$ i.e. $1 < 6$.

$$\begin{aligned} C[3,1] &= C[i-1, j] \\ &= C[2,1] \\ &= 1 \end{aligned}$$

$$\therefore C[3,1] = 1$$

$\Rightarrow C[3,2]$ with $i=3, j=2, d_3=6$.

As $i \neq 1$ and $j < d_3$

$$\begin{aligned} C[3,2] &= C[i-1, j] \\ &= C[2,2] \\ &= 2 \end{aligned}$$

$$\therefore C[3,2] = 2$$

$\Rightarrow c[3,3]$ with $i=3, j=3, d_3=6$.

As $i \neq 1, j < d_3$.

$$c[3,3] = c[i-1, j]$$

$$= c[2,3]$$

$$= 3$$

$$\boxed{\therefore c[3,3] = 3}$$

$\Rightarrow c[3,4]$ with $i=3, j=4, d_3=6$.

As $i \neq 1, j < d_3$

$$c[3,4] = c[i-1, j]$$

$$= c[2,4]$$

$$= 1$$

$$\boxed{\therefore c[3,4] = 1}$$

$\Rightarrow c[3,5]$ with $i=3, j=5, d_3=6$

As $i \neq 1, j < d_3$

$$c[3,5] = c[i-1, j]$$

$$= c[2,5]$$

$$= 2$$

$$\boxed{\therefore c[3,5] = 2}$$

$\Rightarrow c[3,6]$ with $i=3, j=6, d_3=6$

As $i \neq 1, j \neq d_3$

$$c[3,6] = \min(c[i-1, j], 1 + c[i, j-d_i])$$

$$= \min(c[2,6], 1 + c[3,0])$$

$$= \min(3, 1+0)$$

$$= \min(3, 1)$$

$$= 1$$

$$\boxed{\therefore c[3,6] = 1}$$

$\Rightarrow c[3,7]$ with $i=3, j=7, d_3=6$.

As $i \neq 1, j > d_3$.

$$\begin{aligned} c[3,7] &= \min(c[i-1, j], 1 + c[i, j-d_i]) \\ &= \min(c[2,7], 1 + c[3,1]) \\ &= \min(4, 1+1) \\ &= \min(4, 2) = 2. \end{aligned}$$

$$\therefore c[3,7] = 2$$

$\Rightarrow c[3,8]$ with $i=3, j=8, d_3=6$.

As $i \neq 1, j > d_3$.

$$\begin{aligned} c[3,8] &= \min(c[i-1, j], 1 + c[i, j-d_i]) \\ &= \min(c[2,8], 1 + c[3,2]) \\ &= \min(2, 1+2) \\ &= \min(2, 3) = 2. \end{aligned}$$

$$\therefore c[3,8] = 2$$

\rightarrow The value $c[n, N]$ i.e. $c[3,8] = 2$ represents the minimum number of coins required to get the sum 8 units. Hence we require 2 coins for getting 8 units. The coins are:

$$\begin{aligned} 4 \text{ units} &= 1 \text{ coin} \\ + 4 \text{ units} &= 1 \text{ coin} \end{aligned}$$

$$8 \text{ units} = 2 \text{ coins.}$$

* Algorithm of Making Change Problem

Algorithm Number-of-coins (N)

{

for ($i \leftarrow 1$ to n) do

for ($i \leftarrow 1$ to N) do

$c[i, 0] \leftarrow 0$.

For ($i \leftarrow 1$ to n) do

{

For ($d \leftarrow 1$ to N) do

{

if ($i = 1 \wedge j < d[i]$) then

$c[i][j] = \text{infinity}$;

else if ($i = 1$) then

$c[d][j] = 1 + c[1, j - d[i]]$;

else if ($j < d[i]$) then

$c[i][j] = c[i-1, j]$;

else

$c[i][j] = \min(c[i-1, j], 1 + c[i, j - d[i]])$;

}

}

return $c[n, N]$