

## Majority Element - 1

A majority Element is

↳ A element with frequency  $> \frac{N}{2}$  ↗ where 'N' is size of array

arr[6] : { 1 2 1 6 1 1 }  $\Rightarrow$  frequency of '1' =  $(4 > \frac{6}{2})$

so, 1 is a majority Element

arr[9] : { 3 4 4 8 4 9 4 3 4 }  $\Rightarrow$  frequency of '4' =  $(5 > \frac{9}{2})$

so, '4' is a majority Element

Note : There always exists a majority Element, Code accordingly

Idea-1  $\rightarrow$  Use nested loops , TC :  $O(n^2)$

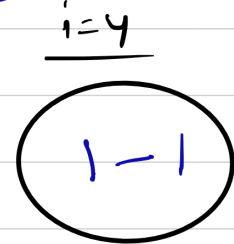
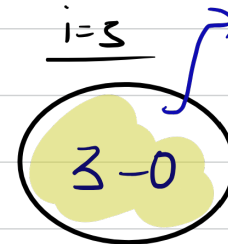
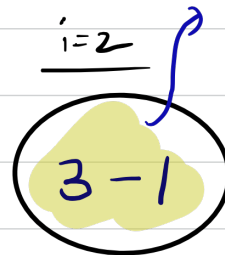
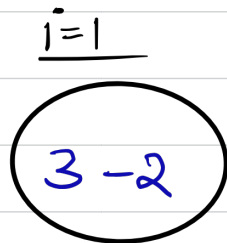
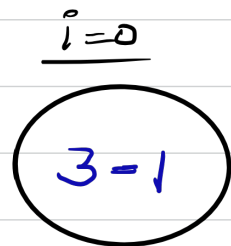
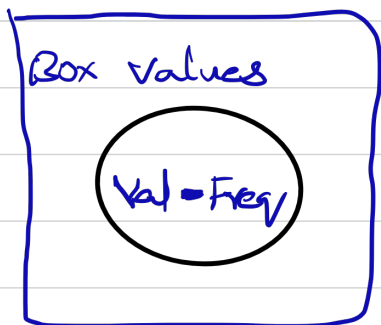
Idea-2  $\rightarrow$  Use Hashmap to Count , TC :  $O(n)$  S.C :  $O(n)$

Optimal approach

## Bayer Moore's Voting Algorithm

If we cancel out 2 distinct Elements one by one, the Uncancelled (or) remaining Element is going to be our majority Element

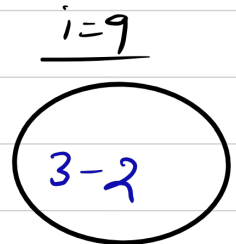
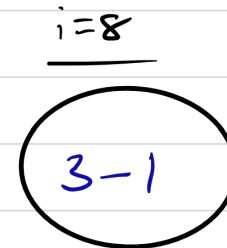
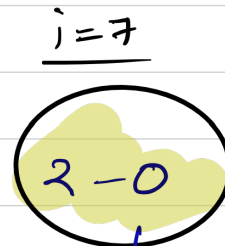
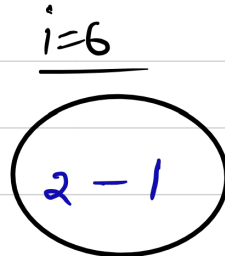
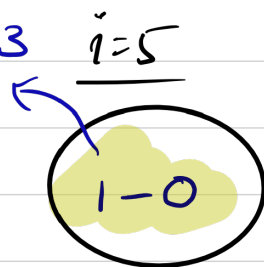
Ex: arr[11]: {<sup>0</sup>3 <sup>1</sup>3 <sup>2</sup>4 <sup>3</sup>6 <sup>4</sup>1 <sup>5</sup>3 <sup>6</sup>2 <sup>7</sup>5 <sup>8</sup>3 <sup>9</sup>3 <sup>10</sup>3}



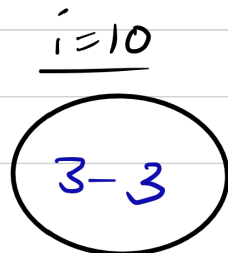
Cancel one 3 with 4

Cancel one 3 with 6

Cancel one 1 with 3



Cancel one 2 with 5



```

1 class Solution {
2     public int majorityElement(int[] nums) {
3         int n = nums.length;
4         int freq = 1;
5         int val = nums[0];
6
7         for(int i = 1; i < n; i++){
8             if(freq == 0){
9                 val = nums[i];
10                freq = 1;
11            }
12            else if(val == nums[i]) freq++;
13            else freq--;
14        }
15        return val;
16    }
17 }

```

See,

1st Box filled with 0th idx val & freq = 1

if freq → 0

then we have to assign new majority element

if opposite elements  
cancel out by  
reducing frequency

for understanding See Box ③ & ④

From Box - 3, the freq of 3 became '0', so in Box ④ we assigned  
1 as our new majority Element with freq 1