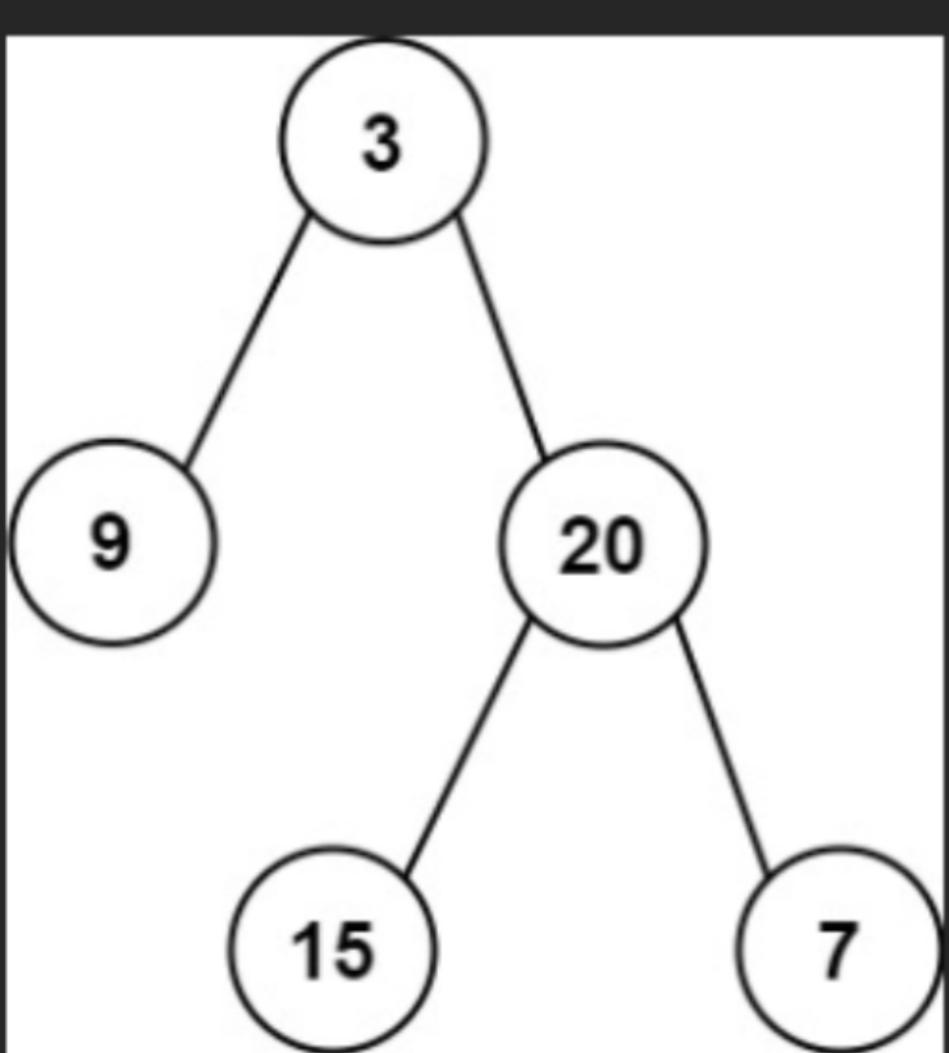


106. Construct Binary Tree from Inorder and Postorder Traversal

Medium Topics Companies

Given two integer arrays `inorder` and `postorder` where `inorder` is the inorder traversal of a binary tree and `postorder` is the postorder traversal of the same tree, construct and return the *binary tree*.

Example 1:



Input: `inorder = [9, 3, 15, 20, 7]`, `postorder = [9, 15, 7, 20, 3]`
Output: `[3, 9, 20, null, null, 15, 7]`

Example 2:

Input: `inorder = [-1]`, `postorder = [-1]`
Output: `[-1]`

inorder :

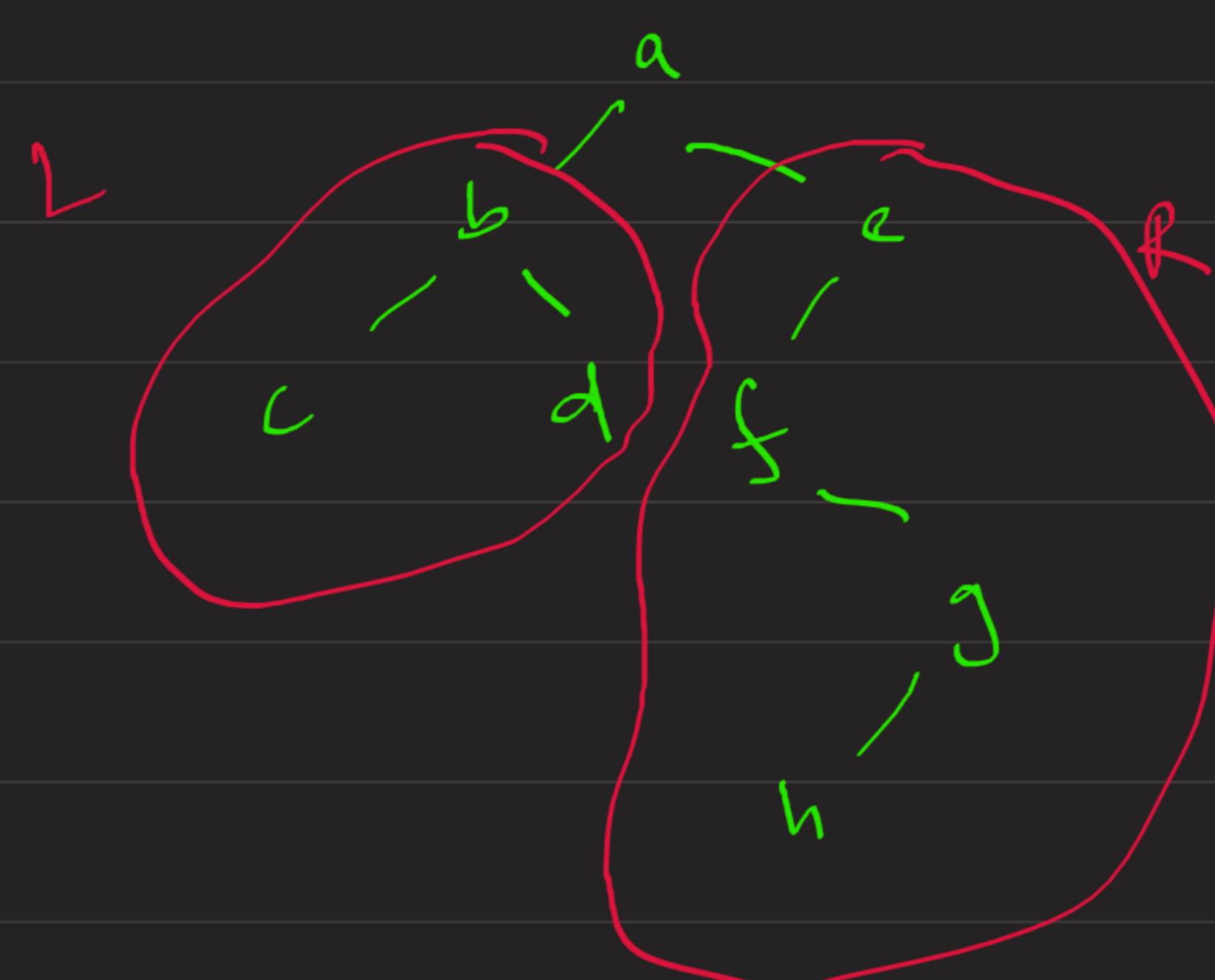
(9, 3, 15, 20, 7)

(L, Root, R)

(9, 15, 7, 20, 3)

(L, R, Root)

* Whenever you see postOrder, the last element will be our root...



→ first we traverse L
 next R
 then root → a

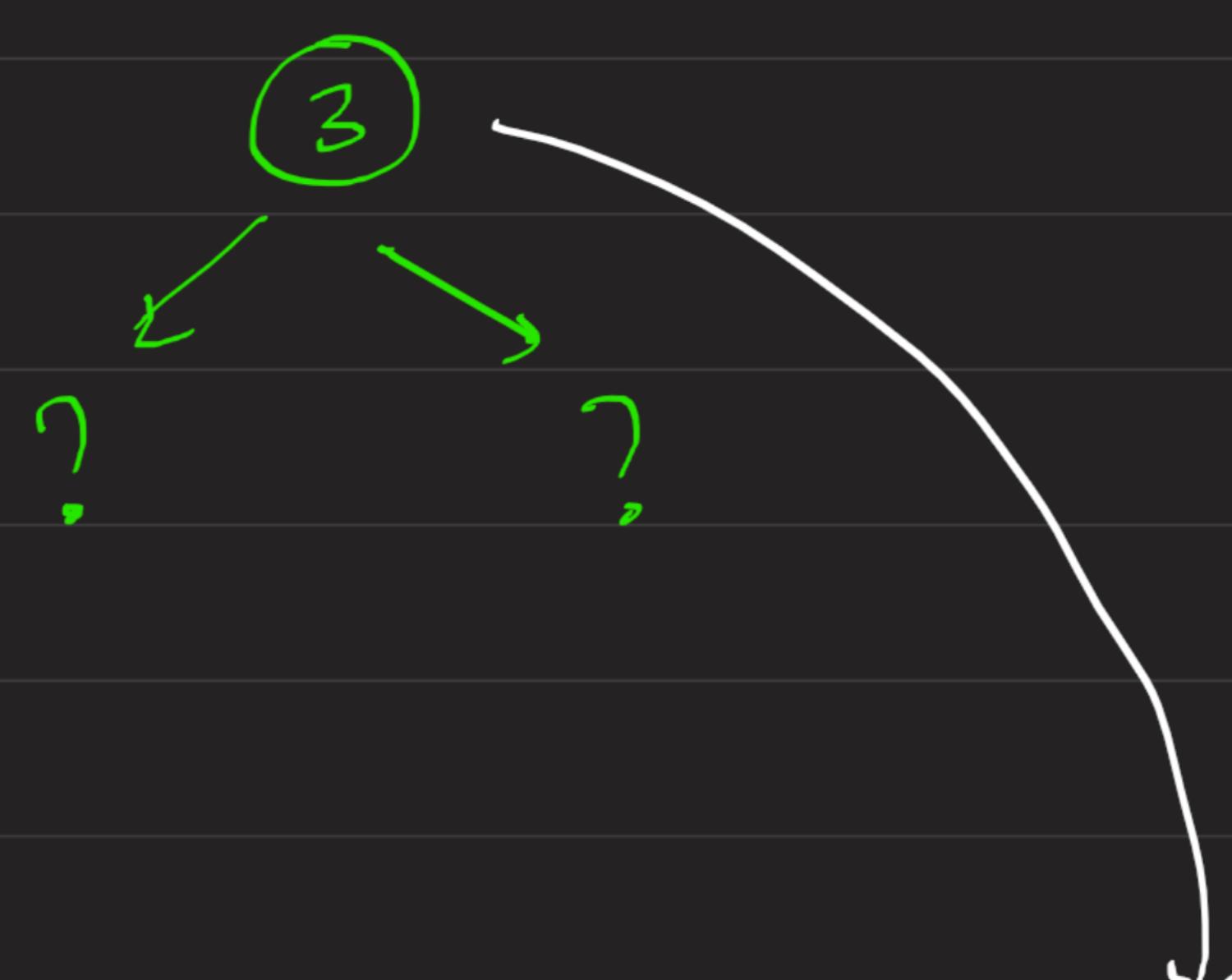
so, a will come at last

a is our root

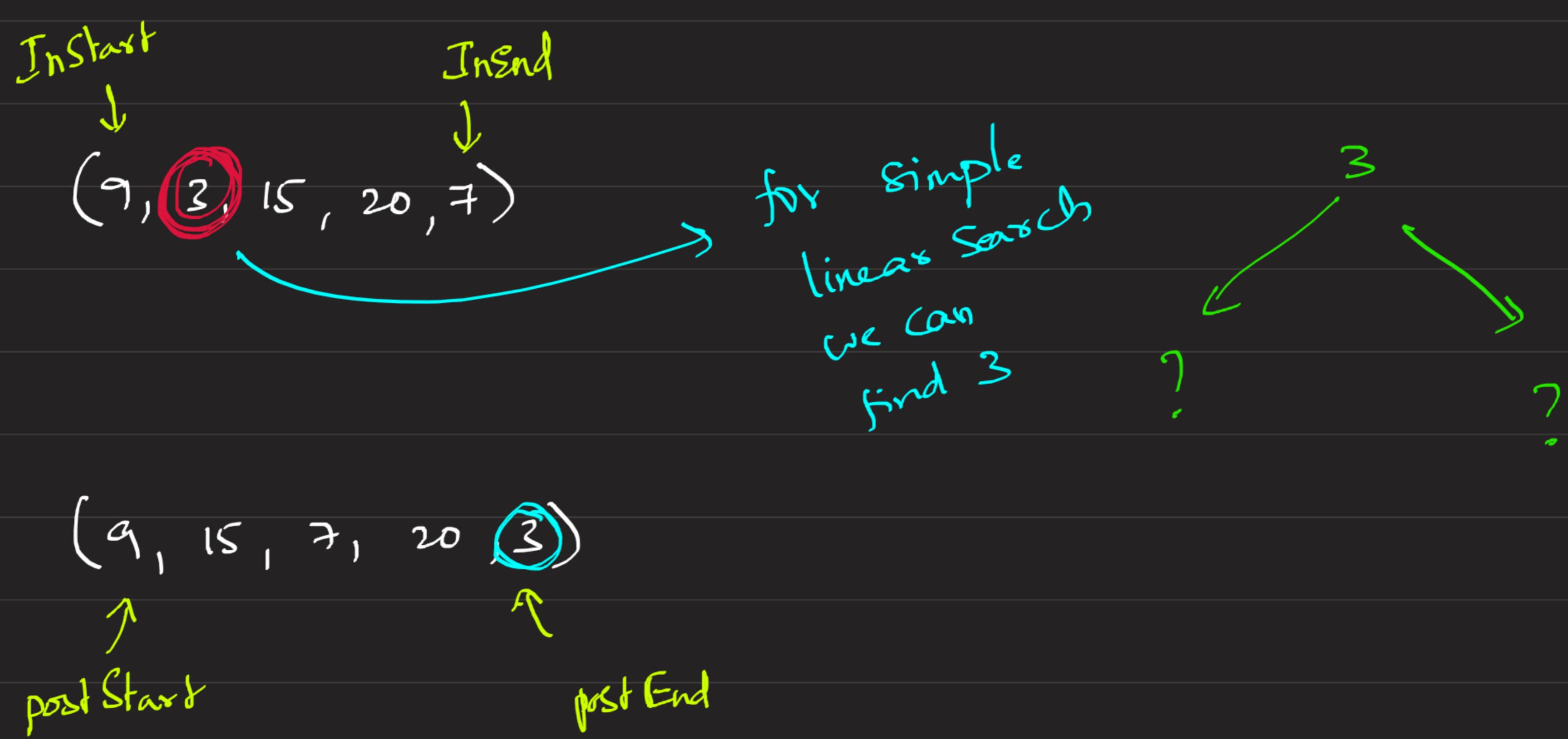
InStart
↓
(9, 3, 15, 20, 7)
InEnd
↓

* root = postOrder[postEnd]

(9, 15, 7, 20, 3)
↑
postStart
↑
postEnd

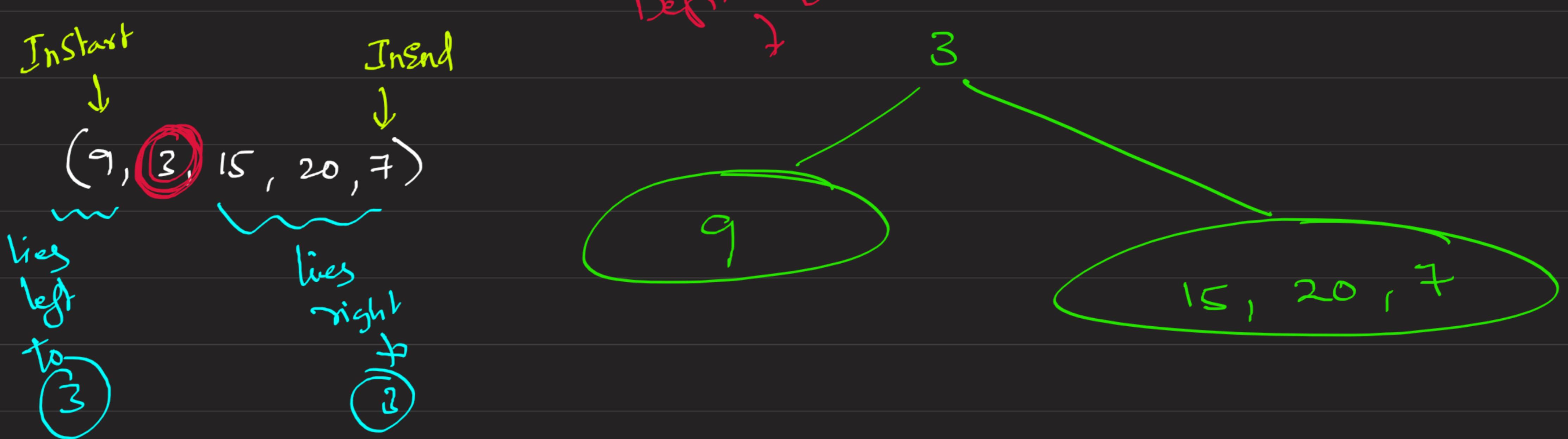


* Now Question ourselves than , where did this (3) located in inorder ...



But what is the need to find $\textcircled{3}$ in inorder..

* It Says ..

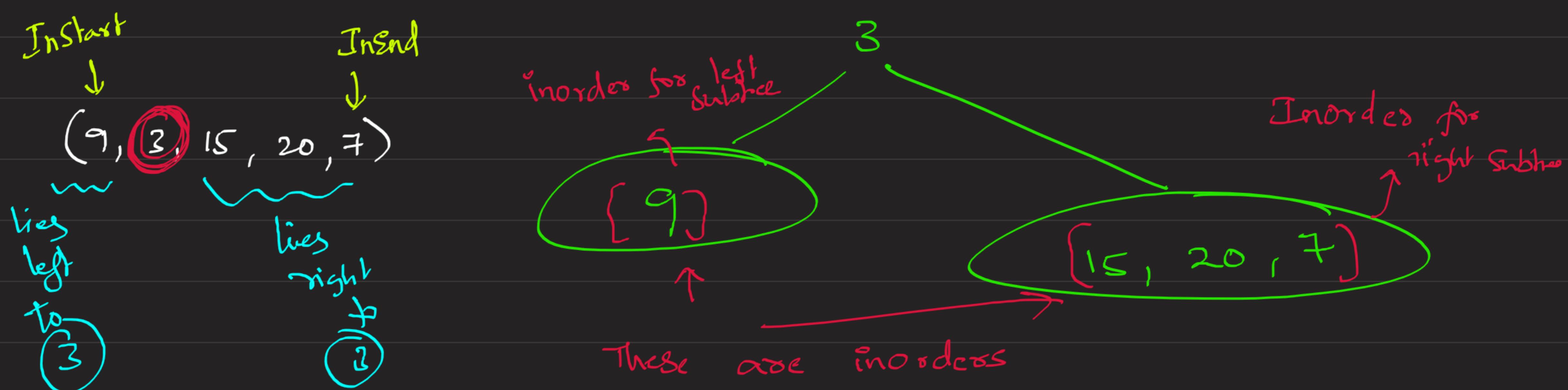


* A key point to remember here is,

We are solving this in recursion → so like we solved for 1st level we needed both inorder & postorder to do that

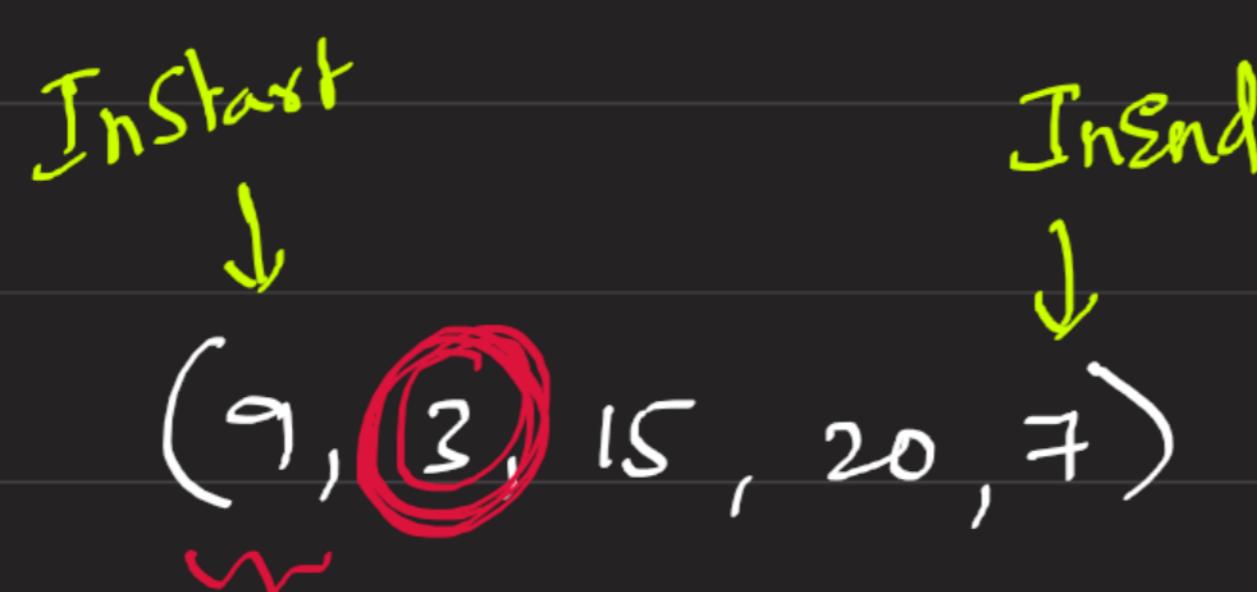
similarly one every level we need both inorder & postorders

* Whatever we have splitted above →

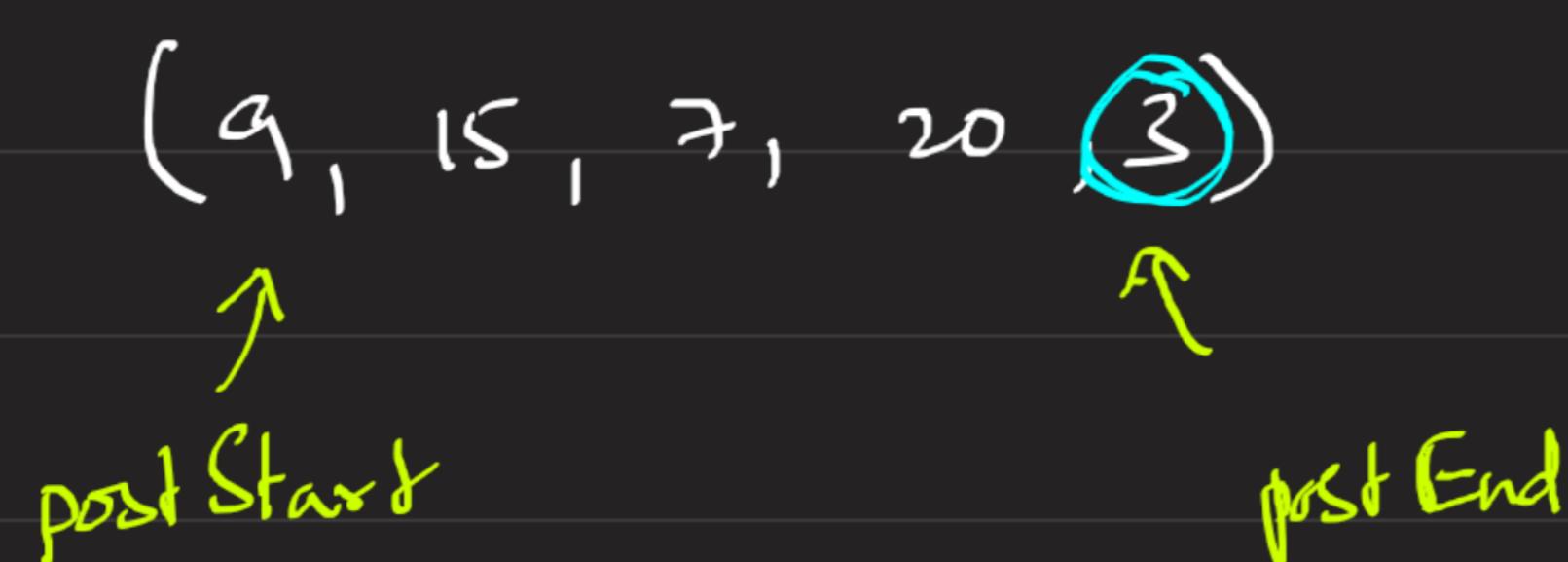


similarly, we need postOrder's to pass in our recursion

let's observe

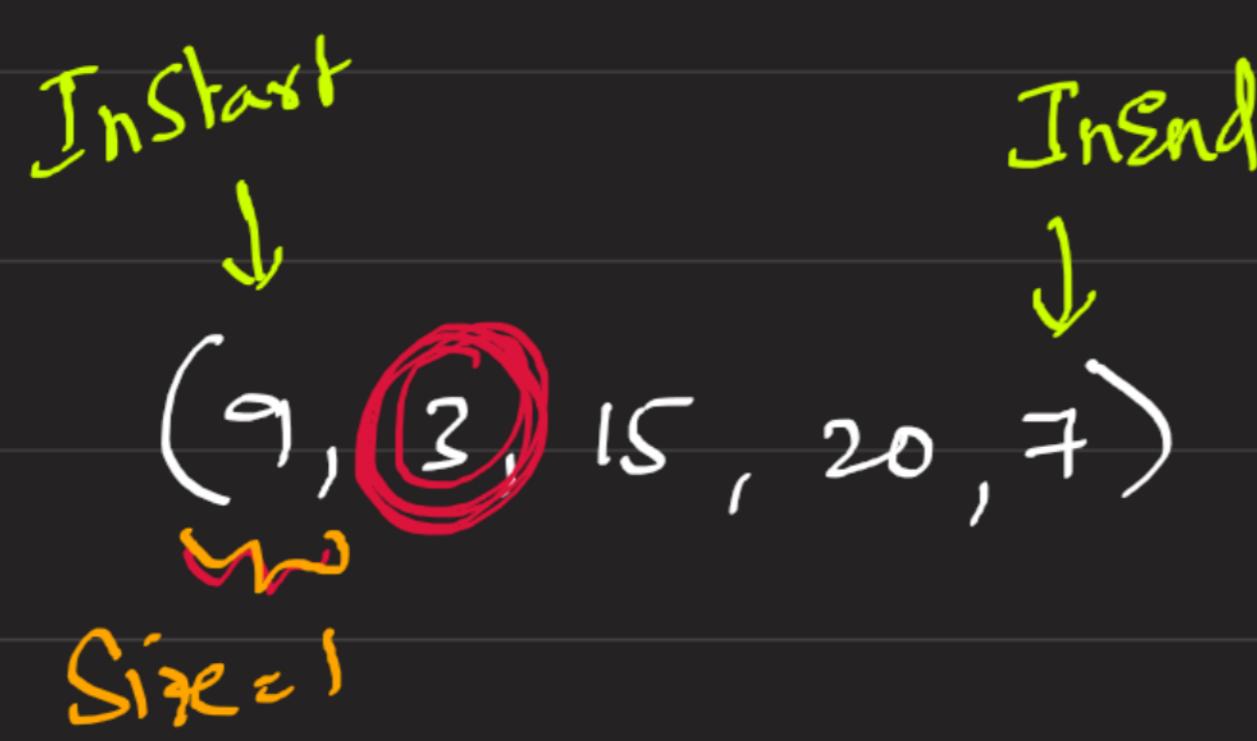


Inorder for left : [9]
right : [15, 20, 7]



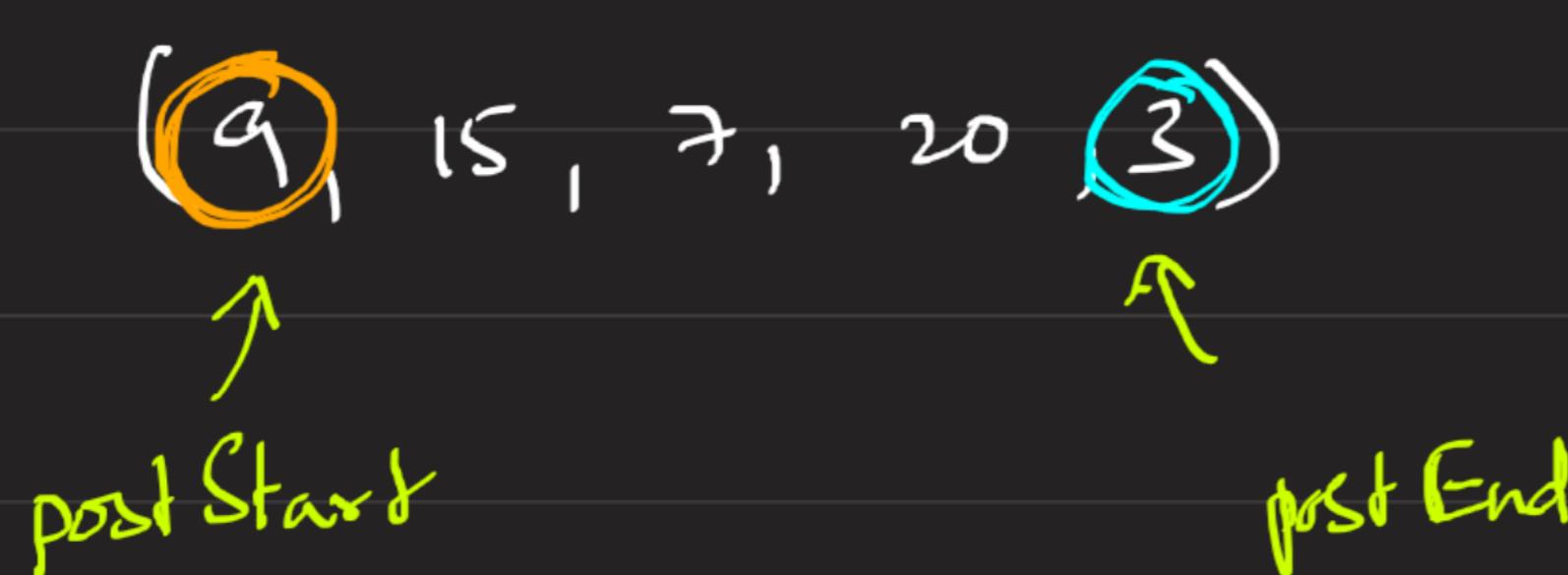
what will be postOrder for left = ?
right = ?

* Count sizes of left tree & right trees of InOrder... add that many nodes to our postOrder ..



Inorder for left : [9]
right : [15, 20, 7]

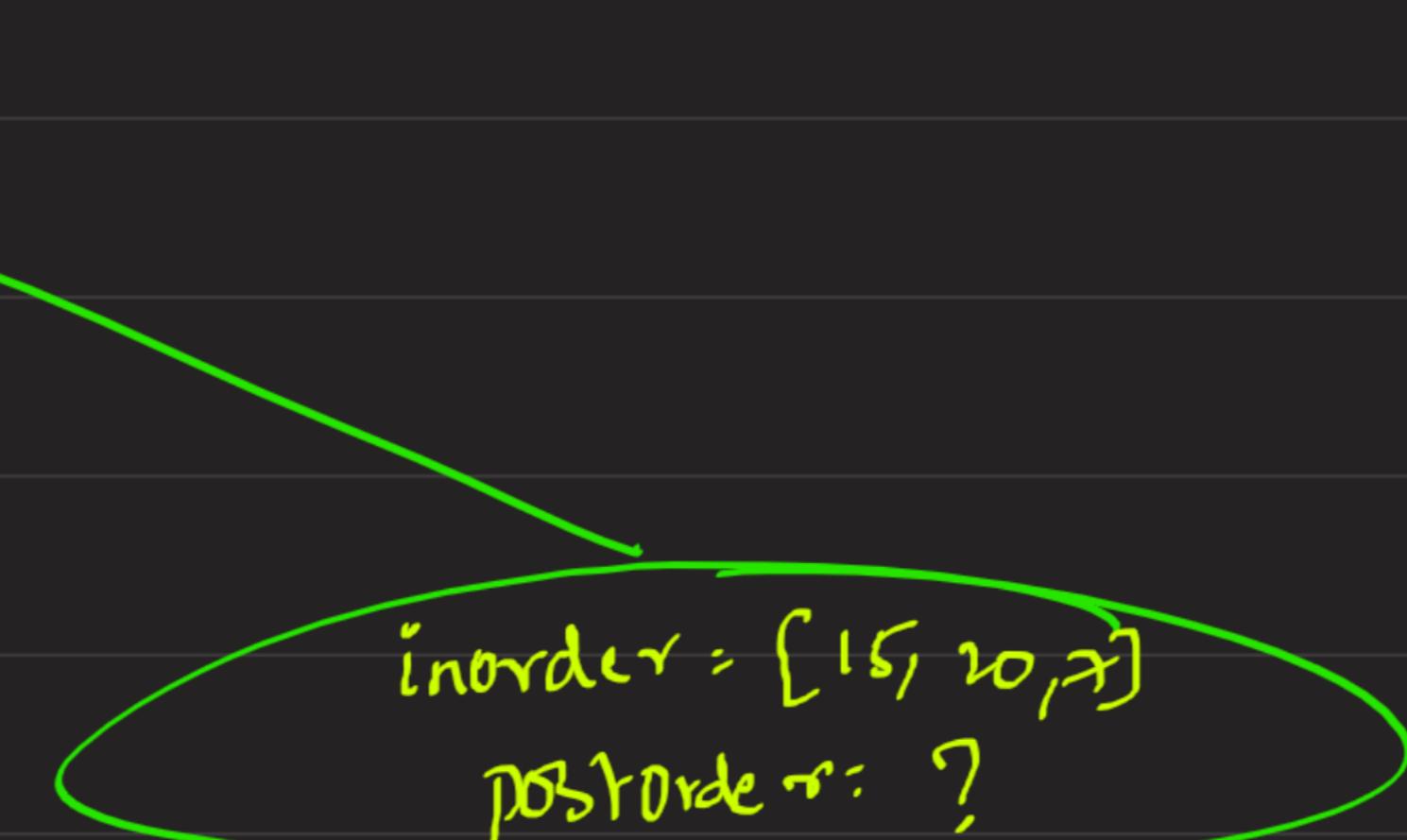
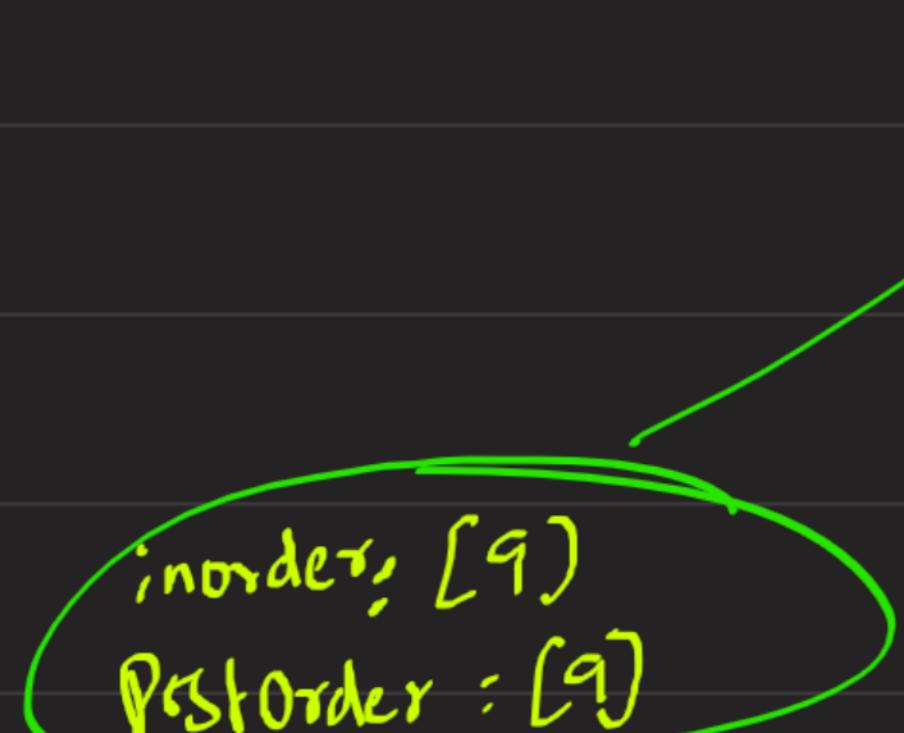
Count = size = 1

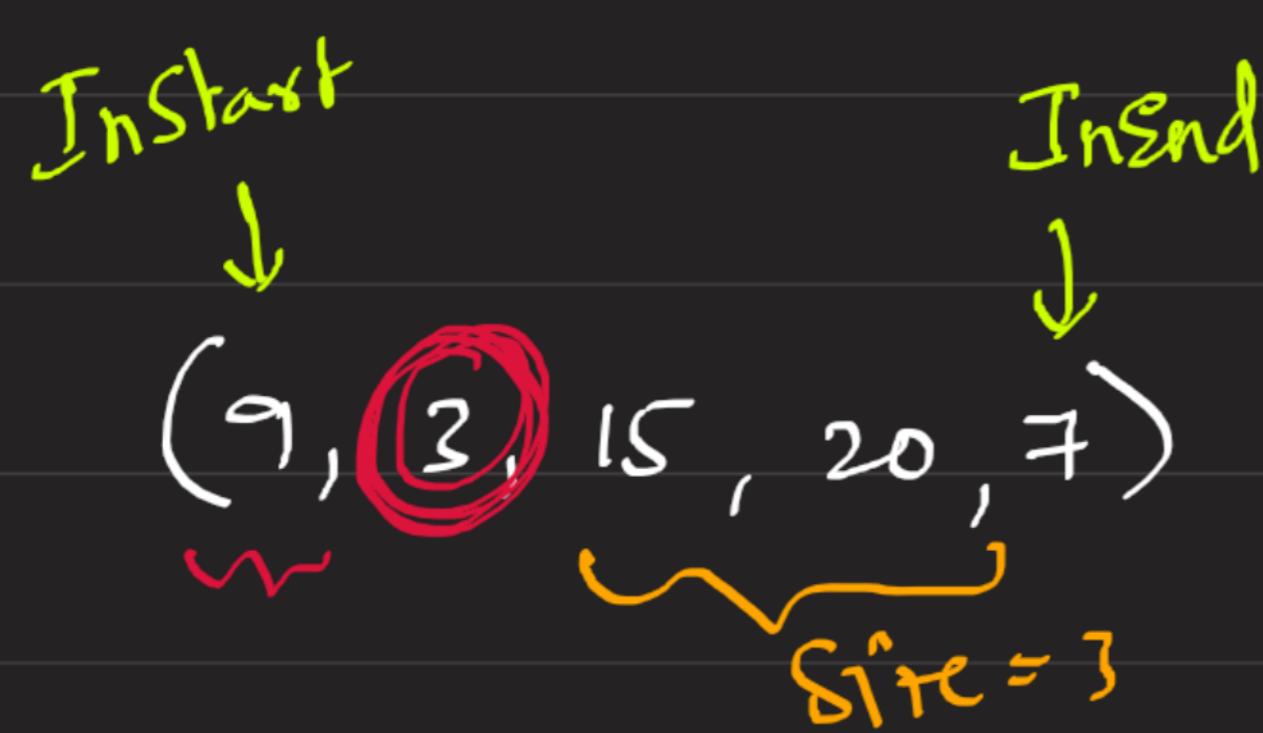


for postOrder's left array :

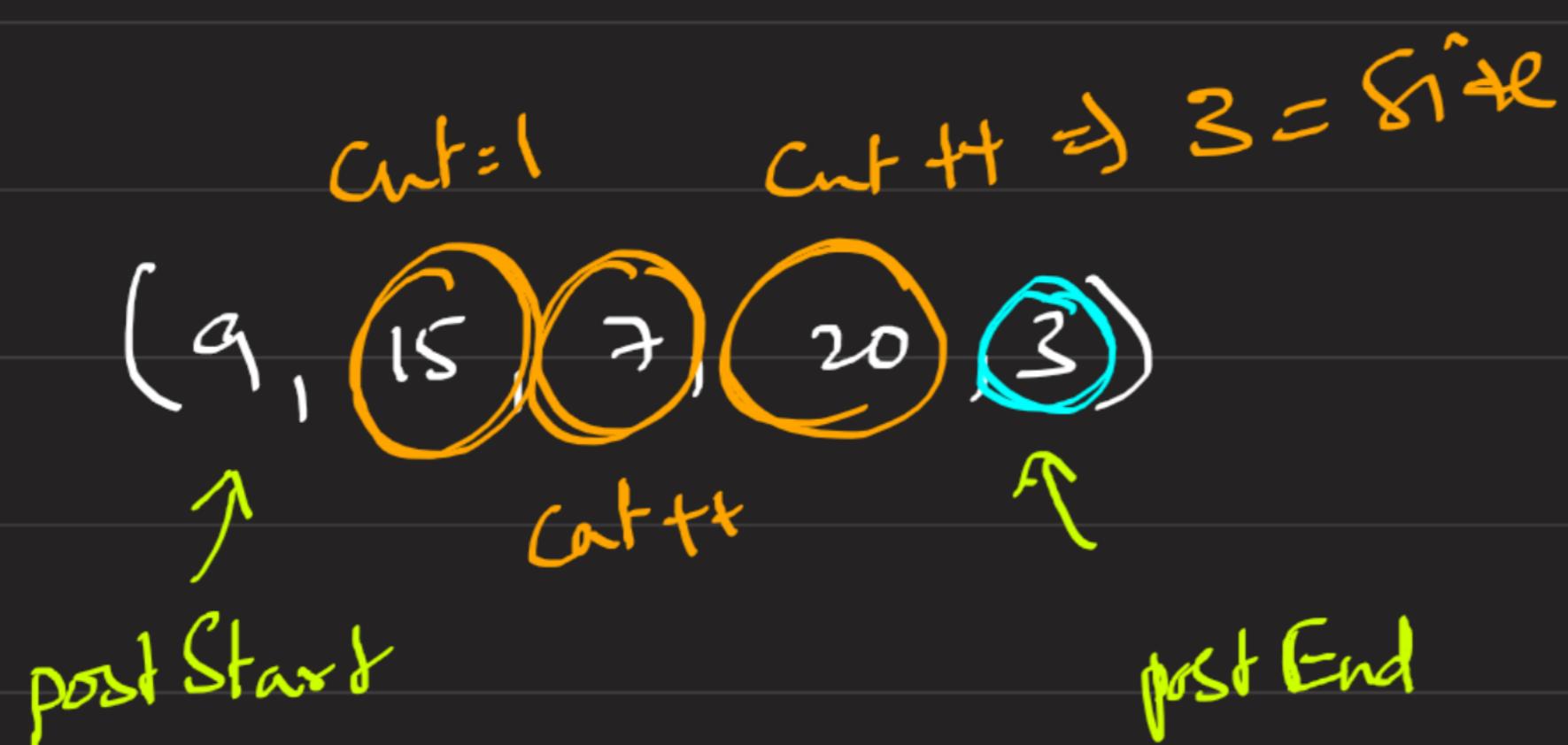
① take inorder's left : [9] \rightarrow size = 1

take 1 node in postOrder & pass it
to recursion as left subtree





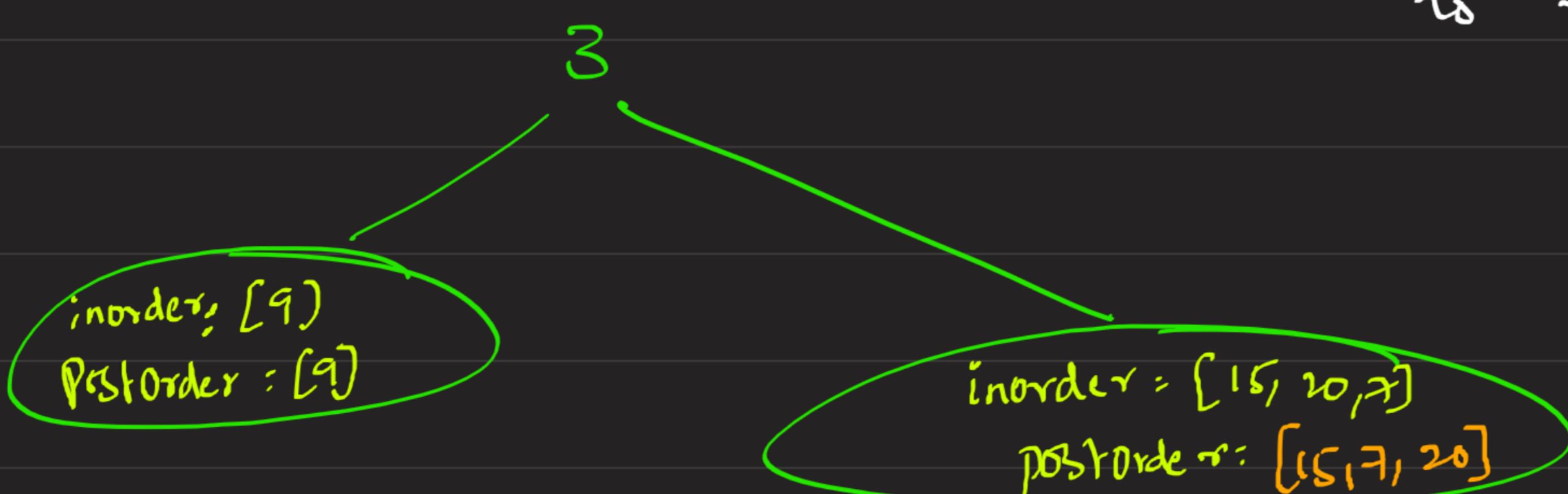
Inorder for left : [9]
right : [15, 20, 7]



for postOrder's right array :

① take inorder's right : [15, 20, 7] \Rightarrow size = 3

take 3 nodes in postOrder & pass it
to recursion as right subtree

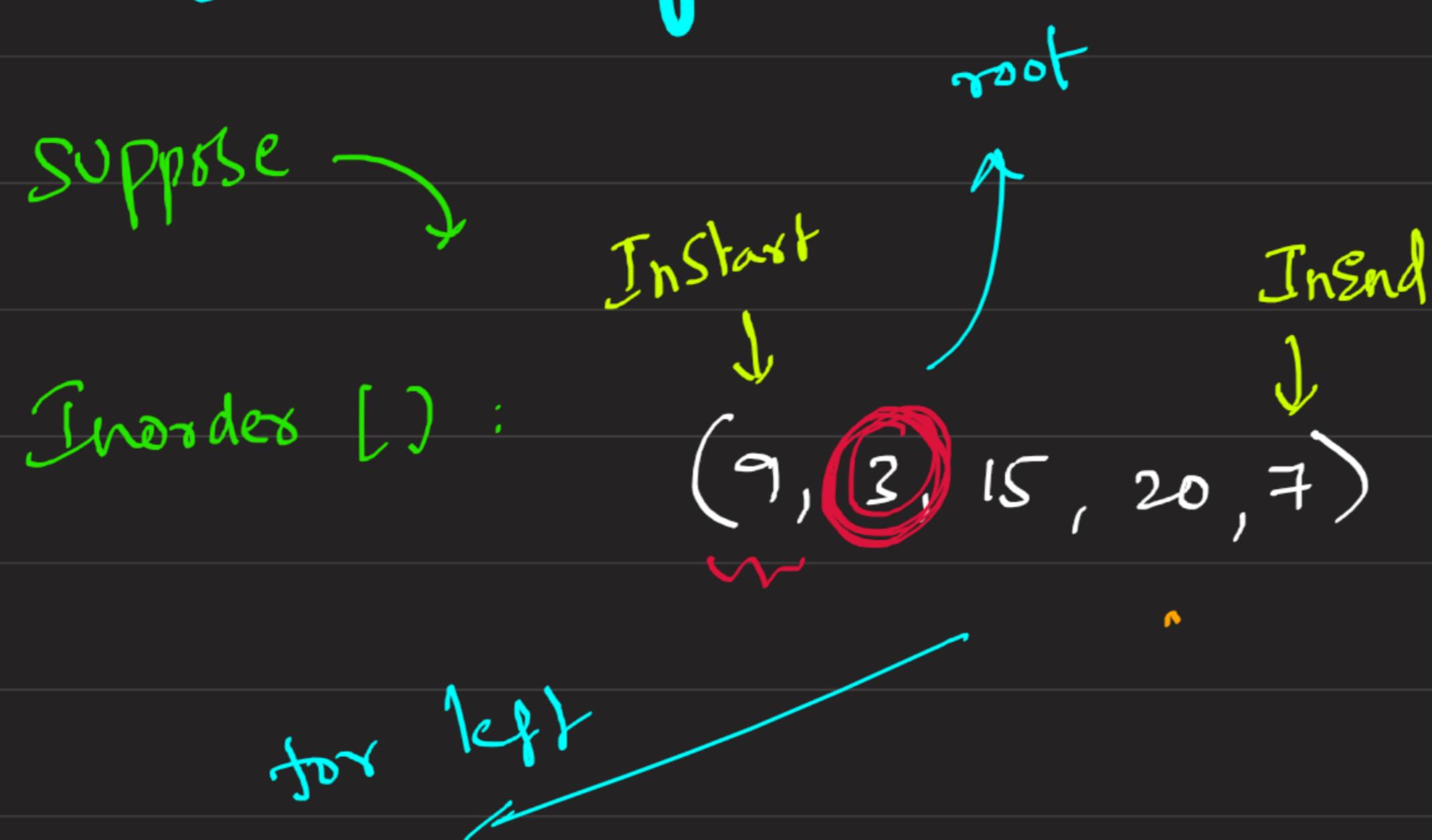


Now, solve these subproblem recursively...

* The main question here is, OK, I understood we need to partition the given array according to our needs, to solve the inorder, postorder

\therefore But how to partition

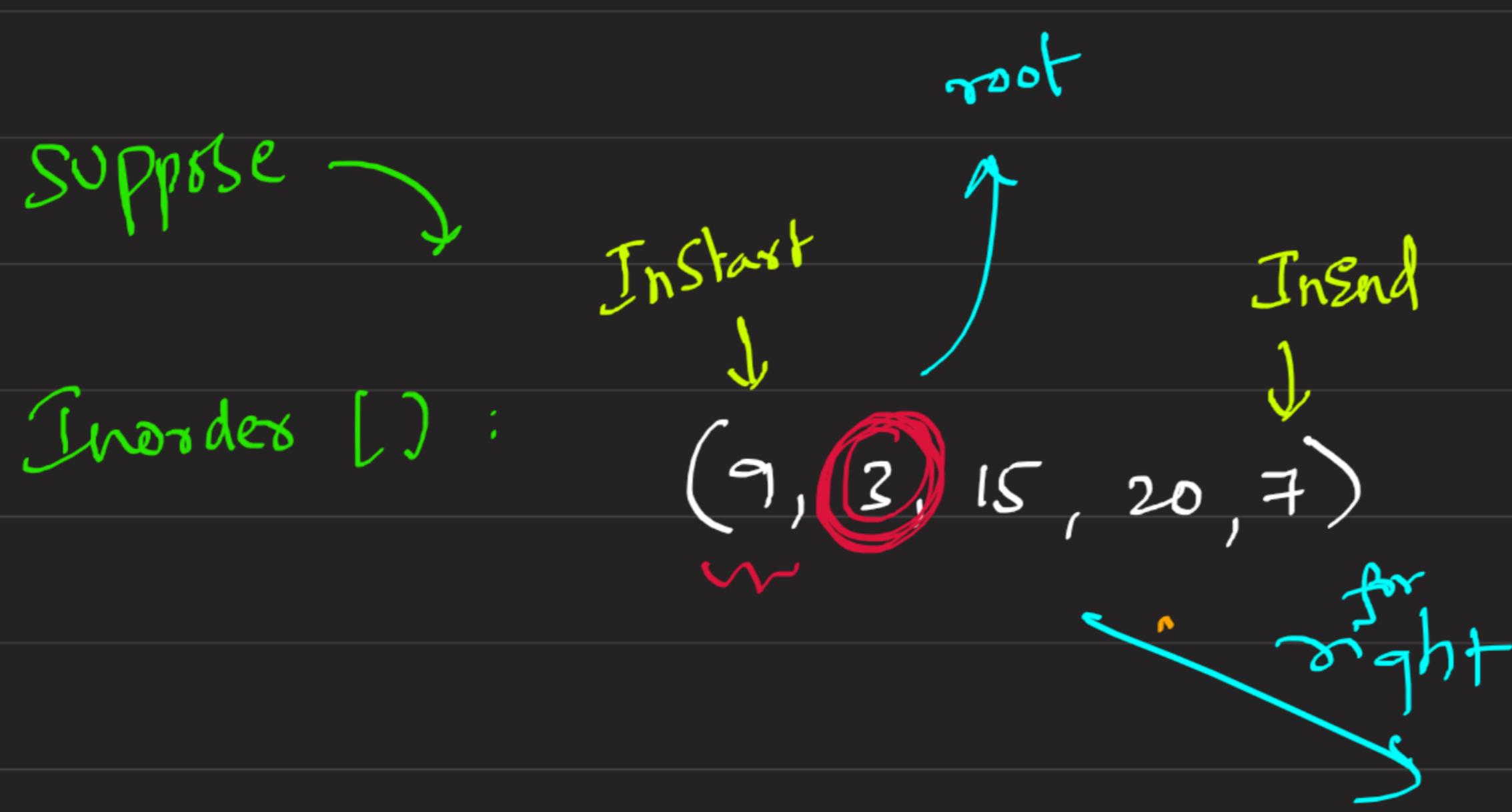
Simply... By indexing



Instart \rightarrow stays at Instart

Inend \rightarrow becomes the node just behind root (if root's index is i)

$$\text{InEnd} = i - 1$$



$inStart \rightarrow$ Becomes the node just after root (if root's index is i)

$$inStart = i + 1$$

$inEnd \rightarrow$ stays at $inEnd$

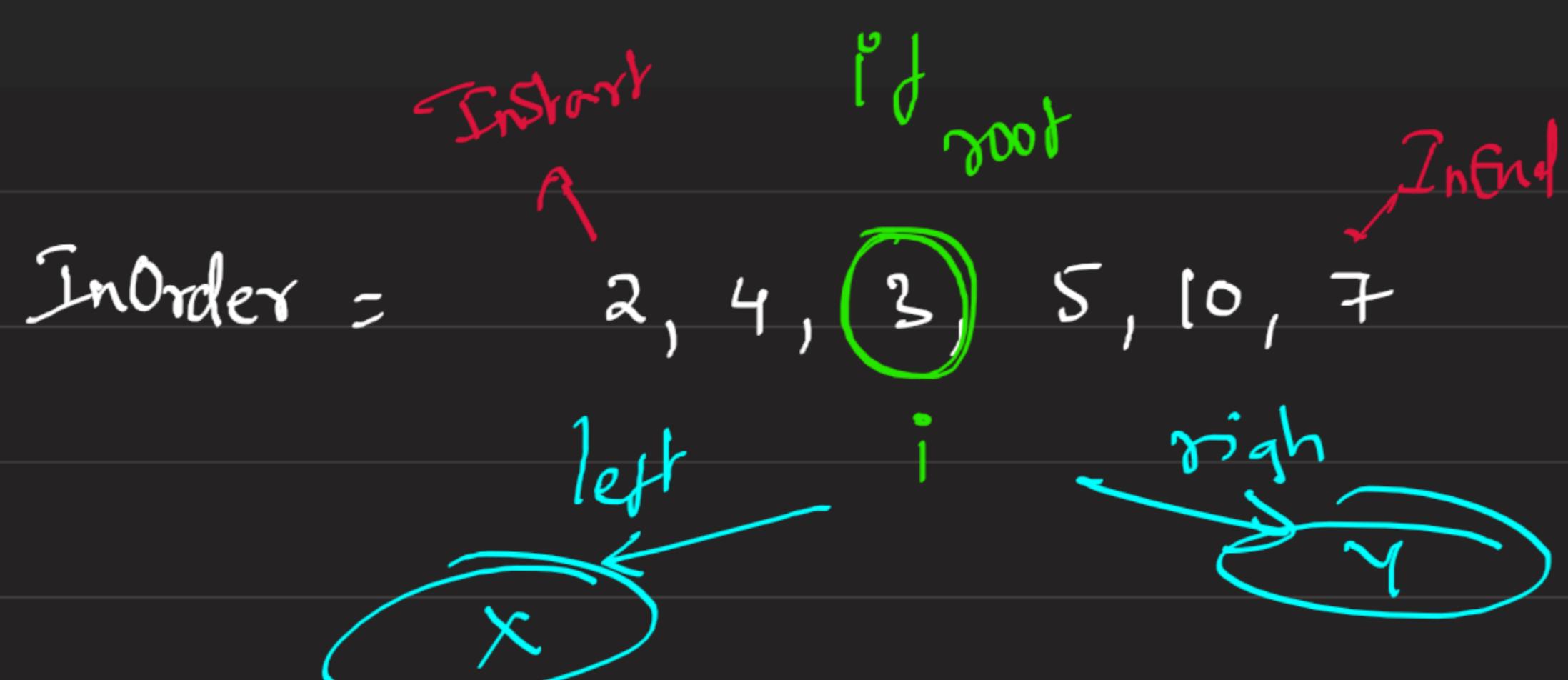
```
class Solution {
    public TreeNode solve(int[] inorder, int[] postorder, int inStart, int inEnd, int postStart, int postEnd) {
        //we know that our root will be last index value in our postorder in every recursive call
        if(inStart > inEnd) return null;
        TreeNode root = new TreeNode(postorder[postEnd]);

        //traverse in inorder and check at which index, our root present in inorder
        int i;
        for( i = inStart;i<=inEnd;i++){
            if(inorder[i] == root.val)break;
        }
        //Now i is at the root node for this current level...
        //we keep trust in recursion that answer for my self....i solved

        //now recursion we give me answers for left subtree and right subtree
        //left-subtree answer

        //to do that...we need size of left inorder tree and right as well
        int leftSize = i-inStart;
        int rightSize = inEnd-i;
        root.left = solve(inorder,postorder,inStart,i-1,postStart,postStart+leftSize-1);
        root.right = solve(inorder,postorder,i+1,inEnd,postStart+leftSize,postEnd-1);
        return root;
    }

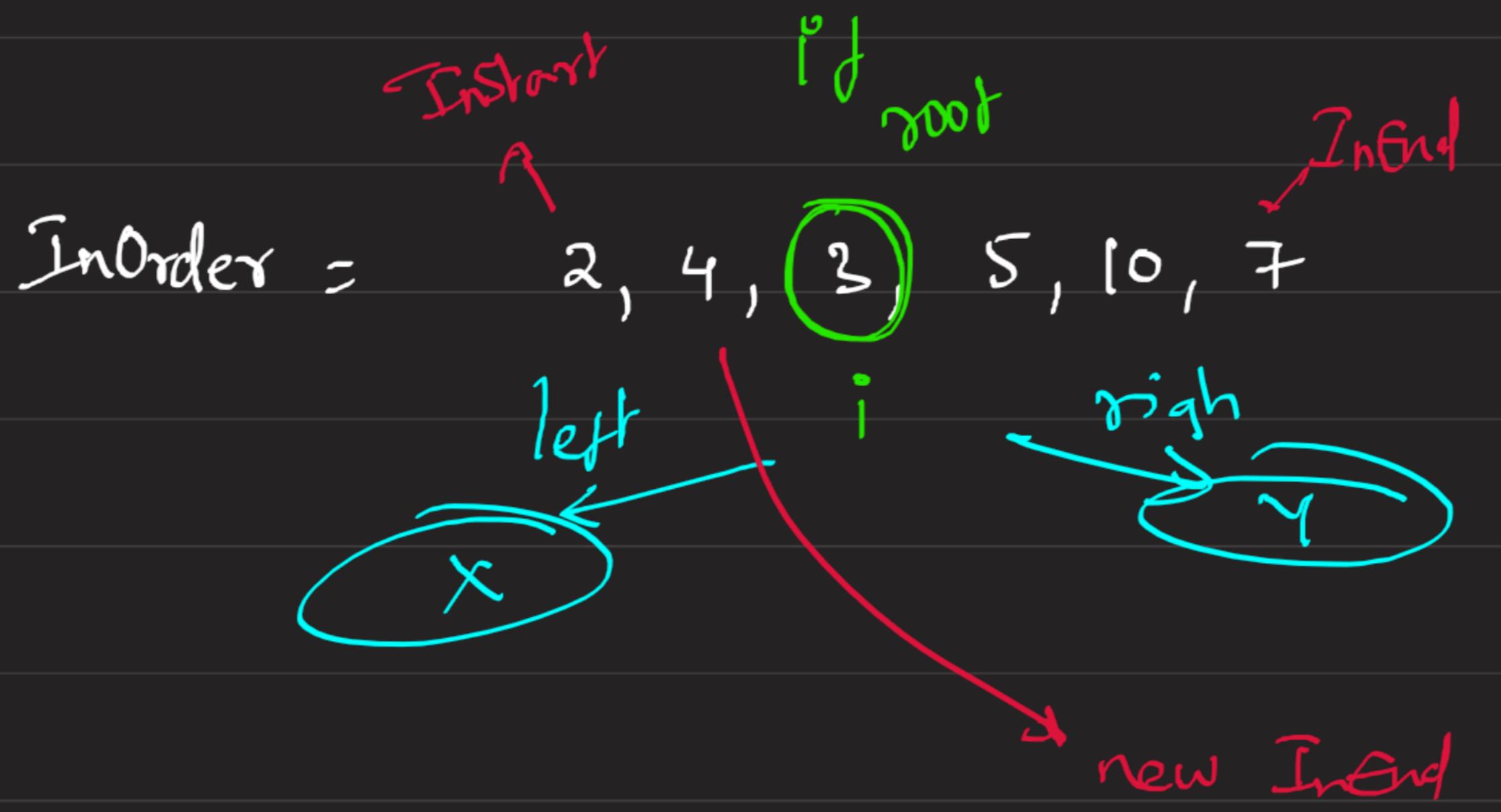
    public TreeNode buildTree(int[] inorder, int[] postorder) {
        return solve(inorder,postorder,0,inorder.length-1,0,postorder.length-1);
    }
}
```



```
root.left = solve(inorder,postorder,inStart,i-1,postStart,postStart+leftSize-1);
root.right = solve(inorder,postorder,i+1,inEnd,postStart+leftSize,postEnd-1);
```

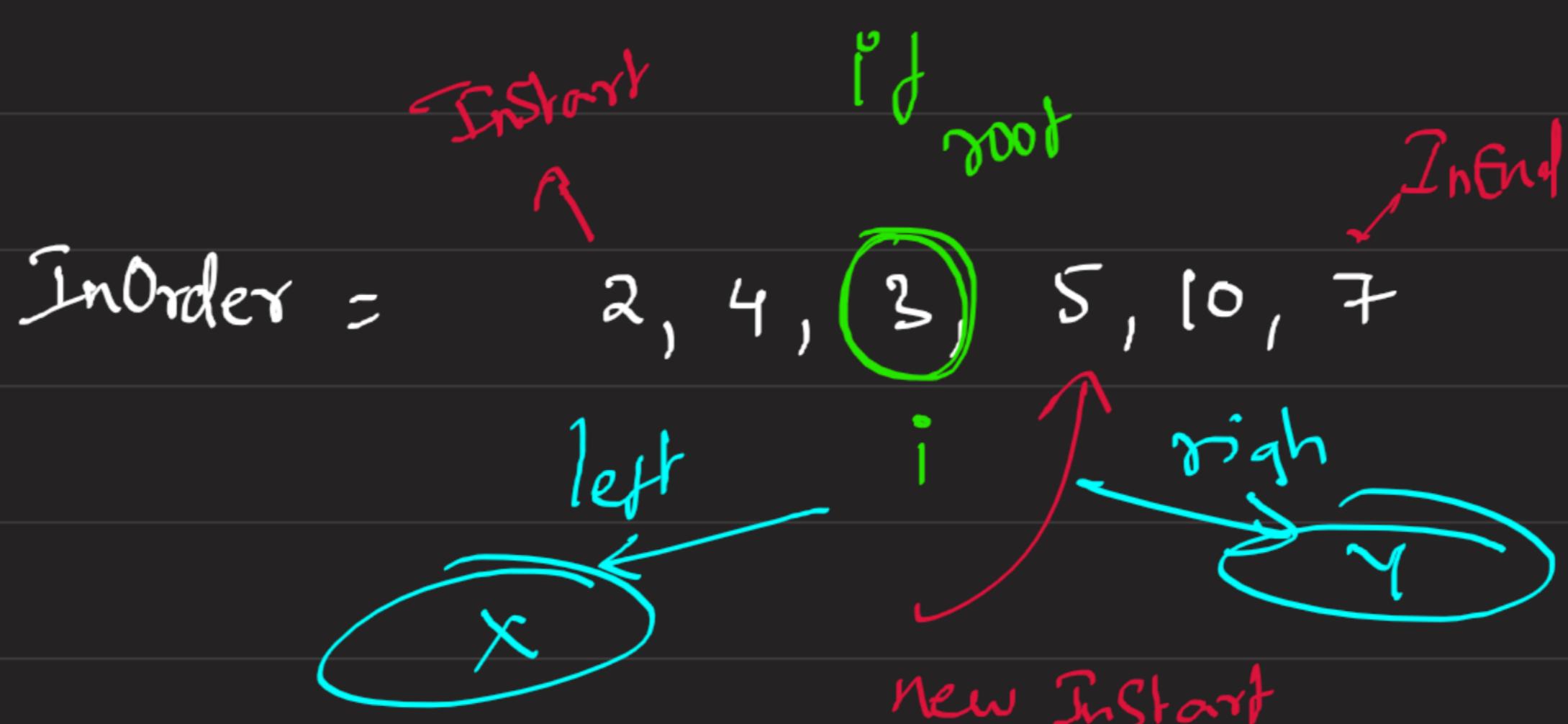
to solve X , we need inorder array & preOrder array

inorder will be left half of ③



InStart of left = InStart

InEnd of left = $i - 1$



InStart of right = $i + 1$

InEnd of right = InEnd

Similarly, to find postOrder of left & right

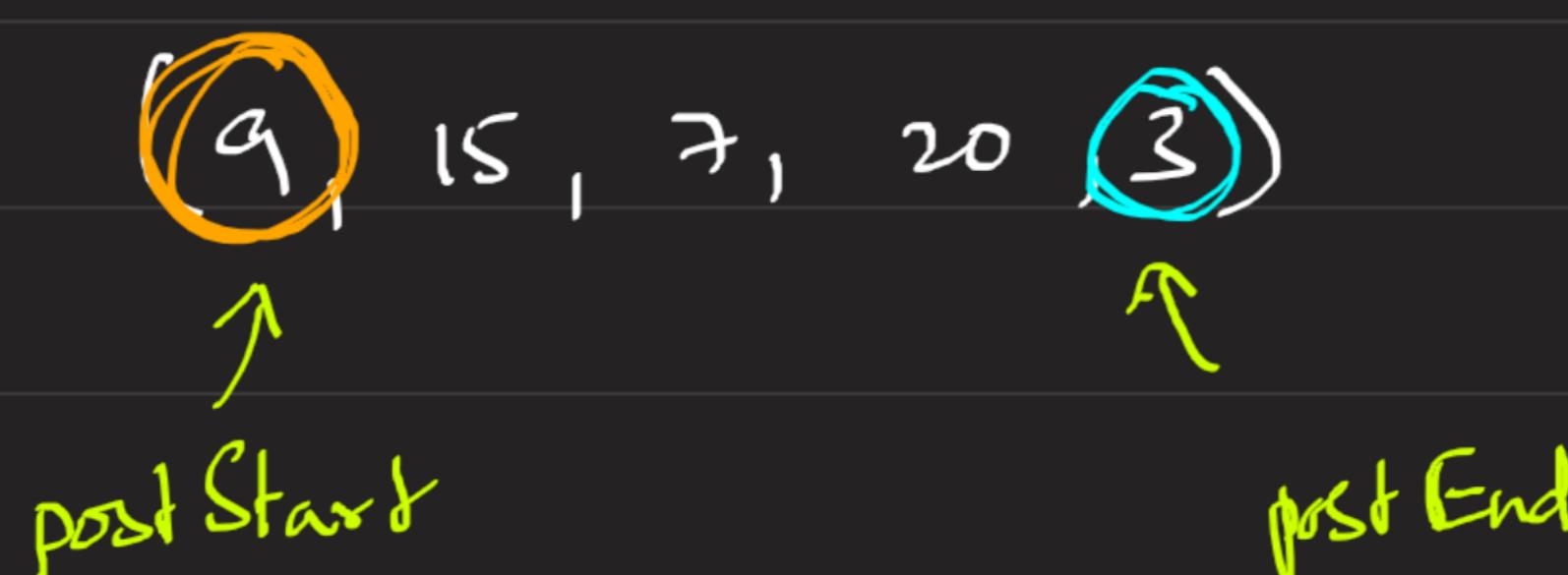
* We need to Count no. of nodes present in left inorder

& Count same no. of nodes from 0th index in postOrder &

pass it as recursion

Similarly in right as well

Cnt = 1



postStart = postStart only

postEnd = postStart + leftSize - 1

If left size of inorder left tree = 1

right size

= 3



postStart = next node of this
postEnd = postEnd only.