

Stacks II

Adventures in Notation

The trouble with infix ...

- Rules for expression evaluation seem simple -- evaluate expression left to right, results of each sub-expression becoming operands to larger expression -- but ...
- All operators are not created equal -- multiplication & division take precedence over addition & subtraction ...

The trouble with infix ...

- So it isn't really all that simple -- we must first scan for higher-precedence operations, and evaluate these first -- and that's not so easy to program -- so ...
- In the calculator program, we rely on parentheses to tell us which operations to perform when -- hence the need for fully-parenthesized expression

Alternatives to infix -- prefix

- Prefix notation, a.k.a. Polish notation
- Operator precedes operands
 - infix expression: $A + B$
 - prefix expression: $+AB$
- Parentheses become unnecessary
 - infix expression: $(A + B) * C$
 - prefix expression: $* + A B C$

Converting from infix to prefix

- Write out operands in original order
- Place operators in front of their operands
- If there's a compound expression, the prefix expression may have two or more operators in a row
- If parentheses are not present, pay attention to precedence

Conversion examples

$A + B + C \ggggggg \quad + + A B C$

$A - B + C \ggggggg \quad + - A B C$

$A + (B - C) \ggggggg \quad + A - B C$

$A * ((B + C) - D) / E \gggg \quad / * A - + B C D E$

$A + B * C / D \ggggg \quad + A / * B C D$

$A * B + C - D / E \gggg \quad - + * A B C / D E$

Prefix evaluation

- scan left to right until we find the first operator immediately followed by pair of operands
- evaluate expression, and replace the “used” operator & operands with the result
- continue until a single value remains

Prefix Example

+ * / 4 2 3 9 // original expression

+ * 2 3 9 // 4/2 evaluated

+ 6 9 // 2*3 evaluated

15 // 6+9 evaluated

Another example

* - + 4 3 5 / + 2 4 3 // original expression
* - 7 5 / + 2 4 3 // 4+3 evaluated
* 2 / + 2 4 3 // 7-5 evaluated
* 2 / 6 3 // 2+4 evaluated
* 2 2 // 6/3 evaluated
4 // 2*2 evaluated

Prefix summary

- Operands (but often not operators) same order as infix
- Expression designated unambiguously without parentheses
- Improvement on infix, but still not quite as simple to evaluate as one might wish -- have to deal with exceptions

Alternative II: Postfix

- Postfix is also known as reverse Polish notation -- widely used in HP calculators
- In postfix notation, operators appear after the operands they operate on
- As with prefix, operand order is maintained, and parentheses are not needed
- Postfix expression is ***not*** merely a reverse of the equivalent prefix expression

Postfix expression examples

- Simple expression:
 - Original Expression: $A + B$
 - Postfix Equivalent: $A B +$
- Compound expression with parentheses:
 - original: $(A + B) * (C - D)$
 - postfix: $A B + C D - *$
- Compound expression without parentheses:
 - original: $A + B * C - D$
 - postfix: $A B C * + D -$

Postfix expression evaluation

- Read expression left to right
- When an operand is encountered, save it & move on
- When an operator is encountered, evaluate expression, using operator & last 2 operands saved, saving the result
- When entire expression has been read, there should be one value left -- final result

Postfix evaluation using stack

- Postfix evaluation can easily be accomplished using a stack to keep track of operands
- As operands are encountered or created (through evaluation) they are pushed on stack
- When operator is encountered, pop two operands, evaluate, push result

Translating infix to postfix

- Postfix expression evaluation is easiest type to program
- Next task is to take an infix expression and translate it into postfix for evaluation
- Some basic assumptions:
 - all operations are binary (no unary negative)
 - expressions are fully parenthesized

Translating infix to postfix

- General method:
 - move each operator to the right of its corresponding right parenthesis
 - eliminate parentheses

- Example:

$((((A + B) * C) - (E * (F + G))))$

$((((A B) + C) * (E (F G) +) *) -$

$A B + C * E F G + * -$

Pseudocode for translation program

```
Do {  
    if (left parenthesis) read & push  
    else if (operand) read & write to output string  
    else if (operator) read & push  
    else if (right parenthesis)  
        read & discard  
        pop operator & write to output string  
        pop & discard left parenthesis  
} while (expression to read)
```

OK -- but what if expression *isn't* fully parenthesized?

- We have to fall back on the rules for expression evaluation we know & love
 - do expression in parentheses first
 - do other operations in order of precedence -- in case of tie, leftmost sub-expression wins
- Example: $A - (B + C) * D - E$
- order: 3 1 2 4
- Postfix: $A B C + D * - E -$

Algorithm for expression conversion

Do

if (left parenthesis) read & push

else if (operand) read & write to file

else if (arithmetic operator)

// continued on next slide

...

Conversion algorithm continued

```
while (stack not empty && stack.peek( ) != '(' &&  
      op precedence lower than stack.peek( ))  
    pop stack & write to file  
    read op  
    push op  
else // character should be ')'  
    // next slide ...
```

Conversion algorithm continued

read & discard right paren

do

pop stack & write to file

while (stack.peek() != '(')

pop & discard left paren

while (expression to read) *// ends outer loop*

while (stack not empty)

pop & write to file