

**EE236A Linear Programming**  
**Solutions of Quiz 1**  
**Tuesday October 13, 2020**

NAME: \_\_\_\_\_ UID: \_\_\_\_\_

This quiz has 3 questions, for a total of 20 points.

Open book.  
The exam is for a total of 1:00 hour. **Please, write your name and UID on the top of each sheet.**

**Good luck!**

Problem	Mark	Total
P1		6
P2		7
P3		7
Total		20

**Problem 1** (6 points) Let  $x$  be a real-valued random variable which takes values in  $\{a_1, a_2, \dots, a_n\}$  where  $0 < a_1 < a_2 < \dots < a_n$ , and  $\Pr(x = a_i) = p_i$ .

Consider the problem of determining the probability distribution that maximizes the expected value  $\mathbf{E}x$  subject to the constraint that  $\Pr(x \geq \alpha) = b$ , i.e.,

$$\begin{aligned} & \text{maximize} && \mathbf{E}x \\ & \text{subject to} && \Pr(x \geq \alpha) = b \end{aligned} \tag{1}$$

where  $\alpha$  and  $b$  are given ( $a_1 < \alpha < a_n$ , and  $0 \leq b \leq 1$ ). Write (1) as an LP.

**Solution:** Given  $a_i$ 's and  $\alpha$  select  $k = \arg\min_i \{a_i : a_i \geq \alpha\}$ . Then problem can be formulated as an LP as follows:

$$\begin{aligned} & \text{maximize} && p^T a \\ & \text{subject to} && \mathbf{1}^T p = 1 \\ & && \sum_{i=k}^n p_i = b \\ & && p_i \geq 0, \quad \forall i = 1, \dots, n \end{aligned} \tag{2}$$

**Problem 2** (7 points): Can the following problem be expressed using an LP? Explain your approach or why not possible. Consider the  $n$  dimensional real vectors  $x = [x_1, x_2, \dots, x_n]$  and  $z = [z_1, z_2, \dots, z_n]$ , we want to

$$\begin{aligned} & \text{minimize} && \|\alpha x\|_2^2 - \|z\|_1 \\ & \text{subject to} && \max_i x_i^2 \leq \beta, \quad i = 1 \dots n \\ & && -1 \leq z_i \leq 1, \quad i = 1 \dots n \end{aligned} \tag{3}$$

where  $\alpha$  and  $\beta$  are given real nonnegative constants.

**Solution:** First, we can change variables  $x_i^2$  to  $t_i$ . Also note that  $\max_i x_i^2 \leq \beta$  implies that  $x_i^2 \leq \beta$ . As a result we would have  $0 \leq t_i \leq \beta$ . Then we can note that in the optimal solution  $z_i$ 's would be positive so we can get rid of the  $l_1$ -norm and use lower bound  $0 \leq z_i$ . Another equivalent way is to define a new variable  $0 \leq k_i \leq 1$  and use it instead of  $|z_i|$ . In the end we can solve the following LP :

$$\begin{aligned} & \text{minimize} && \alpha^2 \sum_{i=1}^n t_i - \sum_{i=1}^n z_i \\ & \text{subject to} && 0 \leq z_i \leq 1, \quad i = 1 \dots n \\ & && 0 \leq t_i \leq \beta, \quad i = 1 \dots n \end{aligned} \tag{4}$$

**Problem 3** (7 points) Formulate the following problem as an LP. Four wireless basestations  $n_1, n_2, n_3$  and  $n_4$  are placed on the circumference of a circle, as depicted in Figure 1. When node  $i$  transmits, the two nodes closest to it cannot transmit, because they would cause interference. For example, when basestation 1 transmits, basestations 2 and 4 cannot transmit. Each basestation  $i$  transmits at a rate of  $r_i$  (the rates  $r_i$  are given constants) per time unit; moreover, each basestation needs to

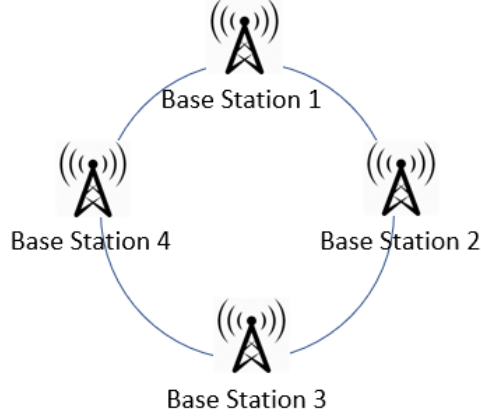


Figure 1: Wireless basestations positioned on the circumference of a circle

transmit for at least  $1/8$  of each time unit. Write an LP that maximizes the total amount of rate, transmitted from all four basestations during a time unit.

**Solution:** Let  $t_i$  ( $i = 1, 2, 3, 4$ ) be the fraction of unit time during which basestation  $i$  transmits. Then the LP can be formulated as follows.

$$\begin{aligned}
 & \text{maximize} && \sum_{i=1}^4 t_i r_i \\
 & \text{subject to} && t_i \geq \frac{1}{8}, \quad i = 1, 2, 3, 4 \\
 & && t_1 + t_2 \leq 1 \\
 & && t_1 + t_4 \leq 1 \\
 & && t_2 + t_3 \leq 1 \\
 & && t_3 + t_4 \leq 1
 \end{aligned} \tag{5}$$

Since  $t_1 = t_3$  and  $t_2 = t_4$  in the optimal solution, we can also formulate the problem as follows.

$$\begin{aligned}
 & \text{maximize} && t_1(r_1 + r_3) + t_2(r_2 + r_4) \\
 & \text{subject to} && t_i \geq \frac{1}{8}, \quad i = 1, 2 \\
 & && t_1 + t_2 \leq 1
 \end{aligned} \tag{6}$$

The solution of the LP can be found as follows.

If  $r_1 + r_3 \geq r_2 + r_4$ , then  $t_1^* = t_3^* = \frac{7}{8}$  and  $t_2^* = t_4^* = \frac{1}{8}$ .

If  $r_1 + r_3 < r_2 + r_4$ , then  $t_1^* = t_3^* = \frac{1}{8}$  and  $t_2^* = t_4^* = \frac{7}{8}$ .