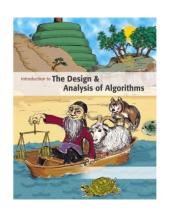




Introduction to

Algorithm Design and Analysis

[1] Model of Computation



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Course Information

- Syllabus
- Textbook
- Website

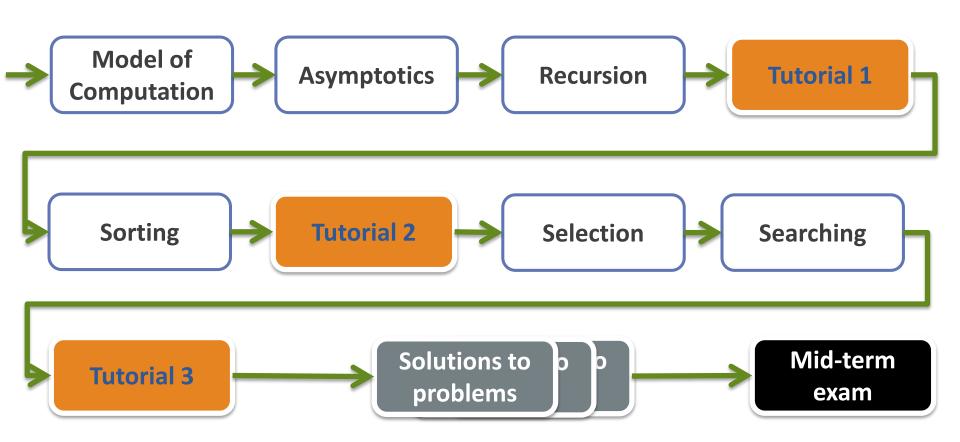


Model of Computation

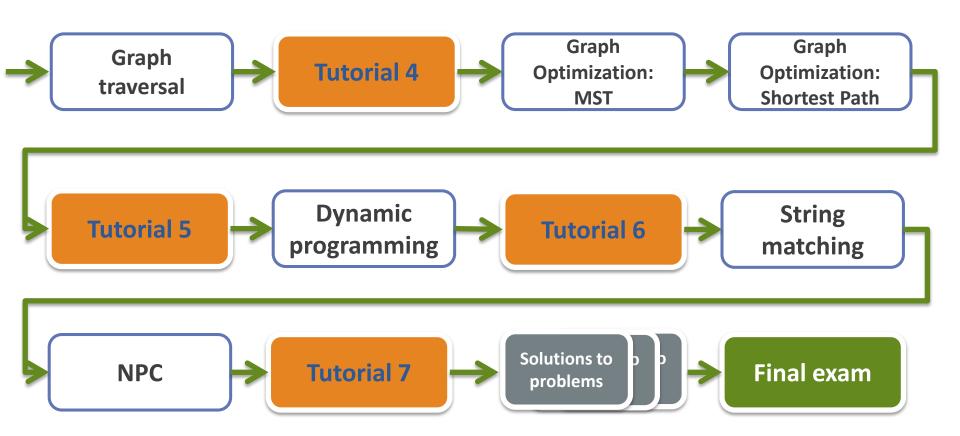
Algorithm
Design &
Analysis
Techniques

Computation Complexity











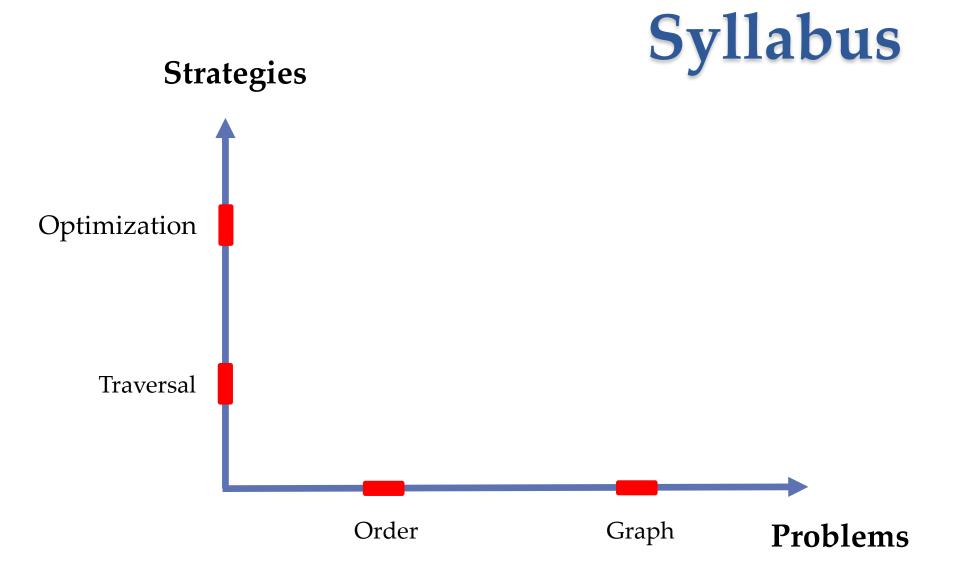
Strategies

Algorithm

Design & Analysis

Problems







Strategies Shortest Sorting Dynamic Divide Programming Path Selection & Optimization **MST** Conquer Greedy Searching **Shortest Path** Sorting DFS Brute Graph Selection **Traversal** Force Traversal **BFS** Searching Order Graph **Problems**

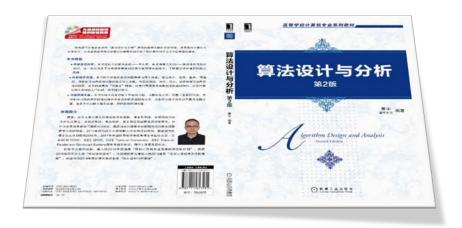


Textbooks

本页与视频略有不同, 以此pdf为准

- Course outline: LADA
 - Lectures on Algorithm Design & Analysis (slides)
- Course contents
 - Algorithm Design and Analysis

More info about the book: https://zhuanlan.zhihu.com/p/24150569

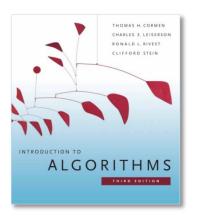


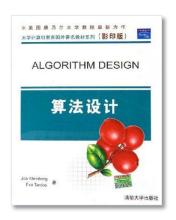
Textbooks

Further reading

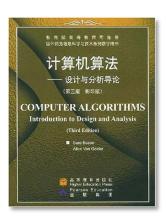
- o Introduction to Algorithms
- o Algorithm Design
- o Algorithms
- o Computer Algorithms*

See the "douban list" for more info: http://book.douban.com/doulist/1155824/









Algorithm – Design & Analysis

- Algorithm the spirit of computing
 - Model of computation
- Algorithm by example
 - o Greatest common divisor
 - o Sequential search
- Algorithm design & analysis
 - o Correctness
 - o Worst-case / average-case cost analysis



Computer and Computing

Problem 1

- o Why the computer seems to be able to do anything?
 - Scientific computing, document processing, computer games, ebooks, movies, computer games, ...













Computer and Computing

Problem 2

- o What can / cannot be efficiently done by a computer?
 - Manage millions of songs vs. music composition



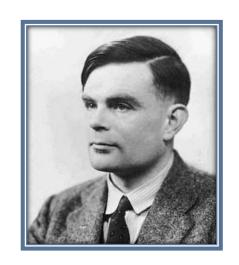




Computer and Computing

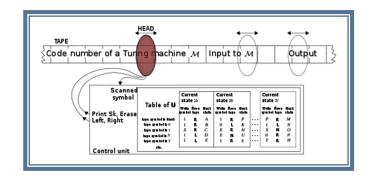
Computing

- Encoding everything into `0's and `1's
- o Operations over '1's and '0's
- o Decoding the '1's and '0's



Turing machine

An abstract/logical computer





Computing in Everyday Life





Algorithm





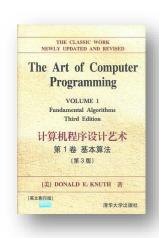
Algorithm

Algorithm is the spirit of computing

- o To solve a specific problem (so called an *algorithmic problem*)
- Combination of basic operations
 - in a precise and elegant way

Essential issues

- Model of computation
- o Algorithm design
- o Algorithm analysis



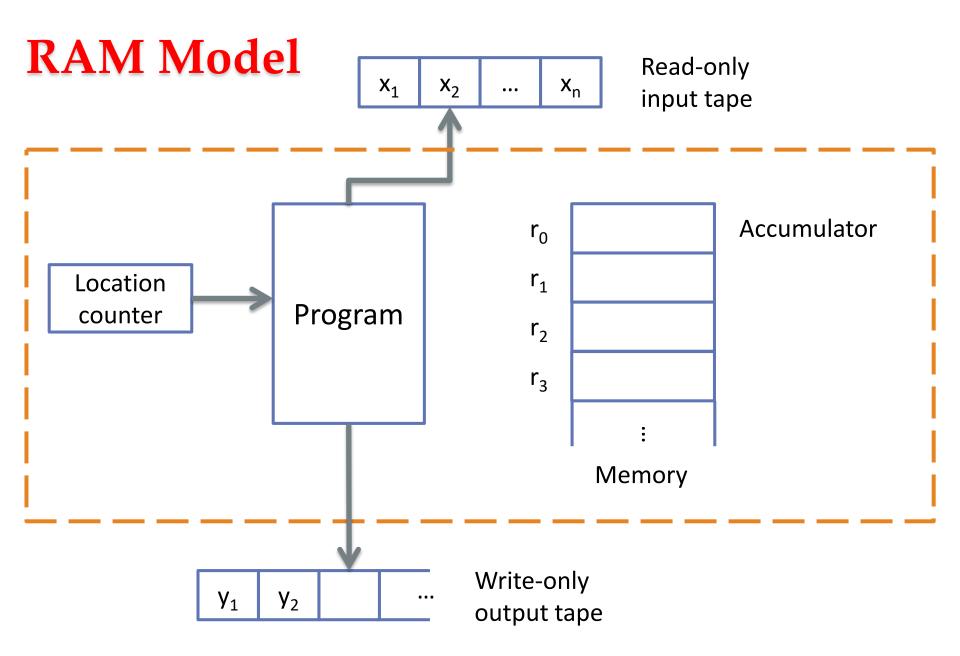


Model of Computation

Problems

- o Why the algorithms we learn can run almost everywhere?
- o Why the algorithms we learn can be implemented in any language?
- Machine- and language- independent algorithms, running on an abstract machine
 - o Turing machine: over-qualify
 - o RAM model: simple but powerful

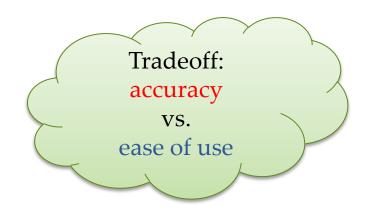






The RAM Model of Computation

- Each simple operation takes one time step
 - o E.g., key comparison, +/-, memory access, ...
- Non-simple operations should be decomposed
 - o Loop
 - o Subroutine
- Memory
 - Memory access is a simple operation
 - o Unlimited memory



Further Reading

"哼,你让他们成楔形攻击队形不就行了?"秦始皇轻蔑地看着冯·诺伊曼。牛顿不知从什么地方掏出六面小旗.三白三黑,冯·诺伊曼接过来分给三名士兵,每人一白一黑,说:"白色代表0,黑色代表1。好,现在听我说,出,你转身看着入1和入2,如果他们都举黑旗,你就举黑旗,其他的情况你都举白旗,这种情况有三种:入1白,入2黑;入1黑,入2白;入1、入2都是白。"

"不需要,我们组建一千万个这样的门部件,再将这些部件组合成一个系统,这个系统就能进行我们所需要的运算,解出那些预测太阳运行的微分方程。 这个系统,我们把它叫做……嗯,叫做……"

"计算机。"汪淼说。

"啊——好!"冯·诺伊曼对汪淼竖起一根指头,"计算机,这个名字好,整个系统实际上就是一部庞大的机器,是有史以来最复杂的机器!"

刘慈欣,《三体、牛顿、冯。诺依曼、秦始皇、三日连珠》,《三体》第一部



To Create an Algorithm

Algorithm design

o Composition of simple operations, to solve an algorithmic problem

Algorithm analysis

- o Amount of work done / memory used
 - In the worst/average case
- o Advanced issues
 - Optimality, approximation ratio, ...



Algorithm by Example

Algorithmic Problem 1

 Find the greatest common divisor of two nonnegative integers *m* and *n*

Algorithmic Problem 2

o Is a specific key *K* stored in array E[1..n]?



Probably the Oldest Algorithm

• Euclid Algorithm

Problem

 Find the greatest common divisor of two non-negative integers m and n

Specification

Input: non-negative integer m, n

Output: gcd(m, n)

Euclid algorithm

[E1] n divides m, the remainder -> r

[E2] if r = 0 then return n

[E3] n -> m; r-> n; goto E1

Euclid algorithm – recursive version

Euclid(m,n)

[E1] if n=0 then return m

[E2] else return Euclid(n, m mod n)



Sequential Search

Problem

 Search an array for a specific key

Specification

Input: K, E[1..n]

Output: Location of K (1,2,...,n; -1: K is

not in E[])

Sequential searchEuclid algorithm

```
Int seqSearch(int[] E, int n, int K)
  int ans, index;
  ans=-1;
  for (index=1; index<=n; index++)
    if (K==E[index])
      ans=index;
    break;
  Return ans;</pre>
```



Algorithm Design

Criteria

Defining correctness

Main challenge

o For proving correctness

Our strategy

Mathematical induction

0 ...

Specification

Input: non-negative integer m, n

Output: gcd(m, n)

Main challenge

- The output is always correct, for any legal input.
- Infinite possible inputs

Mathematical induction

- Weak principle
- Strong principle



For Your Reference

Mathematical induction

The Weak Principle of Mathematical Induction

• If the statement p(b) is true and the statement p(n-1) => p(n) is true for all n>b, then p(n) is true for all integers n>=b.

The Strong Principle of Mathematical Induction

If the statement p(b) is true, and the statement {p(b) and p(b+1) and ... and p(n-1) => p(n)} is true, for all n>b, then p(n) is true for all integers n>=b.



Correctness of the Euclid Algorithm

• Induction on n

- o Base case
 - n = 0: for any m, Euclid(m, 0) = m;
 - n = 1: for any m, Euclid(m, 1) = 1;
 - n = 2: ...
- o Assumption
 - For any $n \le N_0$, Euclid(m, n) is correct;
- o Induction
 - Euclid(m, N_0+1) = Euclid(N_0+1 , m mod (N_0+1));

 $gcd(m, N_0+1) = gcd(N_0+1, m mod(N_0+1))$



Notes on Induction

"Notes on Structured Programming", E.W. Dijkstra

I have mentioned **mathematical induction** explicitly, because it is the only pattern of reasoning that I am aware of, that eventually enables us to cope with loops and recursive procedures.



- Criteria
 - o Performance metrics
- Worst case
 - o Best case?
- Average case
 - o Average cost?
- Advanced topics
 - o Lower bound, optimality, ...



How to measure

- o Not too general
 - Giving essential indication in comparison of algorithms
- o Not too precise
 - Machine independent
 - Language independent
 - Programming paradigm independent
 - Implementation independent



After two years of writing my book, I realized that its novelty was quantitatively determining how good a program is. I didn't just want to say one program was better than another. I wanted to say one is 13.8% better than another, and explain how to compare them.

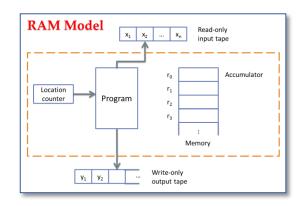
Author A would talk about algorithm A, and author B would talk about his competing algorithm B. And author A never wrote about algorithm B, and author B never wrote about algorithm A. Also, authors A and B used different computers. **As a neutral journalist**, I explained both from one point of view. Asking, "How good is an algorithm really?" is a fun problem. That's the analysis of algorithms.

D.E. Knuth "The Computer Scientist Who Can't Stop Telling Stories" Quanta Magazine



Criteria

- o Critical operation
- How many critical operation are conducted



For example

Algorithmic problem	Critical operation
Sorting, selection, searching String matching	Comparison (of keys)
Graph traversal	Processing a node/edge
Matrix multiplication	Multiplication



Amount of work done

- o usually depends on size of the input
- o usually does not depend on size of the input only





Worst-case Complexity

- W(n)
 - o Upper bound of cost
 - For any possible input

$$\circ W(n) = \max_{I \in D_n} f(I)$$



Average-case Complexity

• A(n)

- o Weighted average
- $o A(n) = \sum_{I \in D(n)} \Pr(I) f(I)$

A special case

- o Average cost
 - Total cost of all inputs, averaged over the input size

$$o Average(n) = \frac{1}{|D(n)|} \sum_{I \in D(n)} f(I)$$



Average-case Cost of SeqSearch

- Case 1: K is in E[]
 - o Assumptions:
 - 1. Assuming that K is in E[]
 - 2. Assuming no same entries in E[]
 - 3. Each possible input appears with equality (thus, K in the i^{th} location with probability $\frac{1}{n}$)

$$O A_{succ}(n) = \sum_{i=0}^{n-1} \Pr(I_i|succ) t(I_i)$$
$$= \sum_{i=0}^{n-1} \frac{1}{n} (i+1)$$
$$= \frac{n+1}{2}$$



Average-case Cost of SeqSearch

- Case 2: K may (or may not) be in E[]
 - o Assume that K is in E[] with probability q

How to make reasonable assumptions?

Advanced topics

- o Lower bound (Selection)
- o Optimality (Greedy, DP)
- Computation complexity
- o Approximate / online / randomized algorithms



Thank you!

Q & A

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Appendix

也许计算的设备会变得天翻 地覆,但是算 法终将是算法

- 黄宇 《算法设计 与分析讲义》





Appendix



How many people have read all of "Introduction to Algorithms"?

Many people suggest this book when asked about algorithms or programming in general. I have this book and it is BIG. I'm wondering how many people have actually read this cover to cover.



Thomas Cormen Professor of Computer Science at Dartmouth ...

Upvote • 1.6k upvotes by Barak Cohen, Hilawi W. Belachew, (more)

I have. And I mean truly cover to cover. I've even read the index.

Written Fri.







