

minimum and maximumAim:

To write a program to find the minimum and maximum elements in a given array of integers.

* Pseudocode

Start

Read N

Read array $a[0..N-1]$

$min = a[0]$

$max = a[0]$

for $i = 1$ to $N-1$ do // loop to find min and max

if $a[i] < min$

$min = a[i]$

if $a[i] > max$

$max = a[i]$

End for

Print min and max

End

* Input

$N = 8$

$a[] = \{5, 7, 3, 4, 9, 12, 6, 2\}$

* Output

$min = 2$

$max = 12$

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Aim: To write a program to find the minimum and maximum elements in a given array of integers sorted in ascending order.

* Pseudocode:

Start

Read N

Read array $a[0..N-1]$

$min = a[0]$

$max = a[0]$

for $i = 1$ to $N-1$ do // loop to find min and max

if $a[i] < min$

$min = a[i]$

if $a[i] > max$

$max = a[i]$

End for

Print min and max

END

* Input

N=8

AC₁ = {2, 4, 6, 8, 10, 12, 14, 18}

* Output

min=2

max=18

④

Merge Sort

Aim: To write a program to Sort an given unsorted array using the merge sort technique.

* Pseudocode:

Mergesort (arr, low, high)

If Low < high

mid = (low+high)/2

Mergesort (arr, low, mid) + n(n-1)/2

Mergesort (arr, mid+1, high)

merge (arr, low, mid, high)

End if

merge (arr, low, mid, high)

Create temporary arrays L and R

copy data to L and R

merge L and R back into arr

* Input

N=8

AC₁ = {31, 23, 38, 27, 11, 21, 15, 28}

* Output

11, 15, 21, 23, 27, 28, 31, 35

⑤

Merge Sort → Count the no. of Comparisons

Aim:

To Implement the merge sort algorithm, sort a given array and count the number of comparisons made during the sorting process.

* Pseudocode:

Count=0

Mergesort (aL, r)

If (l < r)

$m = \lfloor \frac{l+r}{2} \rfloor$
 mergeSort(a, l, m)
 mergeSort($a, m+1, r$)
 merge(a, l, m, r)
 y
 merge(a, l, m, r)
 while ($i < m \text{ } \& \text{ } j < r$)
 count++
 copy smaller of $a[i], a[j]$ to temp, y
 copy remaining elements from temp back to a
 y

* Input:

$N = 8$
 $a[7] = \{12, 4, 78, 23, 45, 67, 89, 13\}$

* Output:

Sorted array: $1, 4, 12, 23, 45, 67, 78, 89$
 No. of Comparisons : 16

(30)

Aim:

To Sort a given array using quick sort by choosing the first element as pivot, showing the array after each partition and recursive call

* Pseudocode

quicksort(a, l, r)

{
 if ($l < r$)

{
 p = partition(a, l, r)

quicksort($a, l, p-1$)

quicksort($a, p+1, r$)

y

partition(a, l, r) {
 pivot = $a[l]$

rearrange elements

return pivot position

* Input

$a[7] = \{38, 27, 43, 3, 9, 82, 102\}$

* Output

$3, 9, 10, 27, 38, 43, 82$

Quick Sort Algorithm

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Aim:

To Implement the Quick-Sort Algorithm using the middle element as the pivot & splits the array after each partition and merges it back and obtain the final sorted array.

* Pseudocode

Quicksort(a, l, r)

{

 p = partition(a, l, r)

 Quicksort(a, l, p-1) and Quicksort(a, p+1, r)

}

y

* Input

N=8

a[] = {19, 72, 35, 46, 58, 91, 22, 31}

* Output

19, 22, 31, 35, 46, 58, 72, 91

(38)

Binary Search Algorithm

Aim:

To implement the Binary Search algorithm, find the position of a given key element in a sorted array & count the no. of comparison.

* Pseudocode

BinarySearch(a, n, key) {

 l=0, r=n-1

 while(l <= r)

{

 mid = (l+r)/2

 if (a[mid] == key) return mid

 else if (key < a[mid]) r=mid-1

 else l=mid+1

}

* Input

N=7

a[] = {21, 32, 40, 54, 65, 76, 87}

Search key = 32

* Output

Position: 2

Comparisons: 2

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Binary Search Algorithm

Aim:

To Locate a given element in a sorted array using the Binary Search algorithm

* Pseudocode

BinarySearch(a, n, key)

```
l=0 ; r=n -1
    while(l <= r)
```

```
    mid = (l+r) / 2
```

```
    if (a[mid] == key) return mid + 1
```

```
    else if (key < a[mid]) r = mid - 1
```

```
    else l = mid + 1
```

y
y

* Input:

N = 7

a[] = {13, 19, 24, 29, 35, 41, 42}

Search key = 42

* Output:

position = 7

K-dosest points

40

Aim:
To find the K closest points to the origin ($0,0$) from a given set of points on the $x-y$ plane.

* Pseudocode

Kdosest(points, n, k)

```
Sort points by (x*x + y*y)
```

return first K points

y
y

* Input:

points = [(1, 3), (2, 2), (-8, 8), (0, 1)]

K = 2

* Output:

[(2, 2), (0, 1)]

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Aim:

To compute the number of tuples (i, j, k, l) such that

$$A[i] + B[j] + C[k] + D[l] = 0$$

* pseudocode:

```
fourSumCount(A, B, C, D) {
```

 Store sums of $A[i] + B[j]$ in map. For each $C[k] + D[l]$,
 add map. $\{ - (C[k] + D[l]) \}$ to count. Set pattern
 return count.

y

* Input:

$A = [1, 2]$

$B = [-2, -1]$

$C = [1, 2]$

$D = [0, 2]$

* Output:

2

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Aim:

To find the k -th smallest element in an unsorted array using the median of medians algorithm

* pseudocode:

```
Select (arr, n, k)
```

{

 if ($n <= 5$) return Sort (arr)[$k-1$]
 divide arr into groups of 5. Find medians of groups

 pivot = Select (medians, $m, m/2$) partition
 around pivot. recurse on required part

* Input:

arr = [12, 3, 5, 7, 19]

$k = 2$

* Output:

5

(43)

k -th smallest element

Aim: To find the k -th smallest element in an unsorted array using the median of medians algorithm

* pseudocode:

```
medianOfMedians (arr, n, k) {
```

 if ($n <= 5$) return Sort (arr)[$k-1$]

 divide arr into groups of 5

 pivot = median of medians

partition arr around pivot
recuse on required part

* Input

$$\text{arr} = [1, 2, 3, 4, 5, 6, 7, 8, 9, 10]$$

$$k=6$$

* Output

6

(4A)

Meet in the middle

Aim:

To find a subset whose sum is closest to a target using Meet in the middle technique.

* Pseudocode:

meetInMiddle(arr, n, target)
{

divide arr into Left and right halves

Compute all subset sums of left $\rightarrow L$

Compute all subset sums of right $\rightarrow R$
Sort R

for each sum in L

find closest sum in R to target - sum

update closest sum; if better return closest sum

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* Input:

Set = {45, 34, 4, 12, 5, 23}

target = 42

* Output:

Closest subset Sum = 41

(4B)

MITM technique

Aim: To determine whether there exists a subset of a given array that sums exactly to a target sum using Meet-in-the-middle technique.

* Pseudocode

meetInMiddle(arr, n, target) {

divide arr into Left and right halves

L = all subset sums of Left

R = all subset sums of right Sort R

for each sum in L

if target - sum exists in R

return true

return false

* Input

$$E = [3, 34, 4, 12, 5, 2]$$

$$\text{exact sum} = 15$$

Output

TRUE

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2x2 matrices

Aim:

To compute the products of two 2×2 matrices using Strassen's matrix multiplication algorithm and obtain the resulting matrix

$$C = A \times B.$$

* Pseudocode

read a,b,c,d,e,f,g,h

$$M_1 = (a+d) * (e+h); M_2 = (c+d) * e$$

$$M_3 = a * (f+h); M_4 = a * (g-e)$$

$$M_5 = (a+b) * h; M_6 = (c-a) * (e+f)$$

$$M_7 = (b-d) * (g+h)$$

$$C_{11} = M_1 + M_4 - M_5 + M_7$$

$$C_{12} = M_3 + M_5$$

$$C_{21} = M_2 + M_4$$

$$C_{22} = M_1 - M_2 + M_3 + M_6$$

print C₁₁, C₁₂, C₂₁, C₂₂

* Input

enter elements of matrix A:

a b

c d

enter elements of matrix B:

e f

g h

* Output

product matrix c:

C₁₁ C₁₂

C₂₁ C₂₂

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Karatsuba multiplication

To Compute the product of two integers using Karatsuba multiplication algorithm, which is the number of multiplications comparable to the traditional method.

* Pseudocode

Karatsuba(x, y):

: if $x < 10$ or $y < 10$ return $x * y$

split x into a, b ; y into c, d

$P_1 = \text{Karatsuba}(a, c)$

$P_2 = \text{Karatsuba}(b, d)$

$P_3 = \text{Karatsuba}(a+b, c+d)$

return $P_1 * 10^n + (P_3 - P_1 - P_2) * 10^{n/2} + P_2$

* Input

$x = 1234$

$y = 5678$

* Output

$z = 7010652$