

Property-based testing: an introduction

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Property-based testing: the basics

Unit testing and its pitfalls

- ▶ The most famous testing paradigm is **unit testing**, where the programmer will manually write test cases.
- ▶ For example, a test could assert that for a function `factorial`, a certain input 3 gives a certain output 5.
- ▶ Unfortunately, writing many test cases can become repetitive and unwieldy, especially if the participating data structures are large, complex, and annoying to instantiate by hand.
- ▶ Unit tests can also be hard to read and understand, as they consist of a set of example input and expected output, but no formal explanation as to why the test is to be expected to succeed.

An alternative: property-based testing

An alternative to unit testing is **property-based testing**, where the programmer will write down a list of *properties* that functions should satisfy, and these will be automatically tested against *random input* in search of a *counterexample*.

Slogan: instead of writing tests by hand, we generate them.

Example of properties

Suppose we have defined operations for addition and multiplication:

$(+) :: \text{Int} \rightarrow \text{Int} \rightarrow \text{Int}$

$(*) :: \text{Int} \rightarrow \text{Int} \rightarrow \text{Int}$

Then we could state properties that we expect them to hold, for example

- ▶ *(Commutative property for addition.)* $\forall x. \forall y. x + y = y + x$
- ▶ *(Distributive property.)* $\forall x. \forall y. \forall z. (x + y) * z = x * z + y * z$
- ▶ ...

If we can find values for x, y, z for which one of these properties fails to hold, then surely the code we have written must be wrong, i.e., the test fails!

Haskell lends itself to property-based testing in view of its lack of side-effects. The famous **QuickCheck** library implements property-based testing for Haskell. Today, we introduce the basics of its API.

A first peek at QuickCheck

Consider our two properties of $+$ and $*$ stated earlier. How would we verify them with QuickCheck?

```
import Test.QuickCheck
```

```
prop_CommutativeAddition :: Int -> Int -> Property
```

```
prop_CommutativeAddition x y = x + y == y + x
```

```
prop_Distributive :: Int -> Int -> Int -> Property
```

```
prop_Distributive x y z = (x + y) * z == x * z + y * z
```

Let us run the tests:

```
ghci> quickCheck prop_Distributive
```

```
+++ OK, passed 100 tests.
```

```
ghci> quickCheck prop_CommutativeAddition
```

```
+++ OK, passed 100 tests.
```

In each case, 100 random examples are chosen to check the properties we have defined.

This is not a proof that the properties hold, but it is *empirical evidence*.

A failing property

Let's see what happens if we write down a wrong property.

```
import Test.QuickCheck
```

```
prop_AssocFloat :: Double  
                -> Double  
                -> Double  
                -> Property
```

```
prop_AssocFloat x y z = (x + y) + z == x + (y + z)
```

```
ghci> quickCheck prop_AssocFloat
*** Failed! Falsified (after 6 tests and 6 shrinks):
-1.0
0.1
1.0
9.999999999999998e-2 /= 0.100000000000000009
```

The last line is the failing equality. The three lines preceding it are the values for the arguments (respectively x, y, z) that constitute the counterexample.

QuickCheck fundamentals

With a few examples in mind, we now dive into the workings of QuickCheck.

- ▶ A **Property** carry instructions for how to generate random test cases, how to test them, and what to print.
- ▶ A **Gen** a (called a generator) specifies how to randomly pick values of type **a**.
 - One can extract random values from a generator using the **generate** function. This is useful for testing your generators.

Working with generators

- ▶ `Gen` is a `Monad` in the natural way. In particular, `return x` is the generator that always yields `x`.
- ▶ The `Arbitrary` type class is used to specify how to randomly generate arbitrary values of a type.

```
class Arbitrary a where
  arbitrary :: Gen a
  -- ...
```

- `Arbitrary` is implemented for many basic types. Therefore, one can use `arbitrary` to get many simple generators. For example, to flip a coin, we can do

```
do x <- arbitrary :: Gen Bool
    if x then return "heads"
        else return "tails"
```

- ▶ `chooseInt` (a, b) is the generator that picks out an integer between a and b (inclusive).
- ▶ `elements :: [a] -> Gen` a produces the generator that picks out a random element of the list.
- ▶ `oneof :: [Gen a] -> Gen` a builds a generator that generates using a generator chosen at random from the list.
 - For example, `oneof [return 1, chooseInt (10, 20)]` is the generator that with probability $\frac{1}{2}$ generates 1 and otherwise picks a random integer in the range from 10 to 20.

- ▶ `frequency :: [(Int, Gen a)] -> Gen a` is a variant of `oneof` that gives you control over probabilities.
 - For example,
`frequency [(1, return 5), (2, return 7)]` yields 5 with probability $\frac{1}{3}$ and 7 with probability $\frac{2}{3}$.
- ▶ `vectorOf :: Int -> Gen a -> Gen [a]` produces a generator for a list of `n` items.
 - For example, `vectorOf 3 (elements [1, 2])` produces lists like `[2, 1, 2]` and `[2, 2, 2]`.

The quickCheck function

```
quickCheck :: Testable prop
            => prop
            -> IO ()
```

where `Testable` is a type class:

```
class Testable prop where
  property :: prop -> Property
```

```
instance (Arbitrary a, Show a, Testable prop)
  => Testable (a -> prop) where
  property f = forAll arbitrary f
```

Take-away: Function arguments represent universal quantifiers.

An example: Morse code

Encoding and decoding Morse code

```
table :: [(Char, String)]
table = [('A', ".-"), ('N', "-."), ...]
encode :: String -> String
encode [] = []
encode (x:xs) = fromJust (lookup x table) ++ encode xs
decode :: String -> [String]
decode "" = [""]
decode m = concatMap
    (\(letter, code) ->
        case stripPrefix code m of
            Just rest -> map (letter:) (decode rest)
            Nothing -> []
    )
table
```

Properties of Morse code

We certainly expect the following property to hold for our program:

- (*) Encoding a string `str` and decoding it again yields `str` and perhaps some other strings.

Indeed, remember that Morse code is ambiguous, which is also why `decode` returns a list.

QuickCheck-ing Morse encoder/decoder

We may test (*) as follows:

```
import Test.QuickCheck

-- Note that 'Bool' implements 'Testable', so this
-- is 'quickCheck'-able.
prop_EncodeDecode :: String -> Bool
prop_EncodeDecode str = str `elem` decode (encode str)
```

Let's try to run it:

```
ghci> quickCheck prop_EncodeDecode
*** Failed! (after 2 tests and 1 shrink):
Exception:
  Maybe.fromJust: Nothing
  CallStack (from HasCallStack):
    error, called at libraries/base/Data/Maybe.hs:149:21
      in base:Data.Maybe
    fromJust, called at test/Test.hs:60:17 in main:Main
"a"
```

The issue: `encode` doesn't accept arbitrary strings, only all uppercase alphabetic strings.

We have to make a custom generator:

```
-- | A string of alphabetic, uppercase characters.
data UpperCaseString = UpperCaseString String

instance Arbitrary UpperCaseString where
  arbitrary = do
    n <- chooseInt (0, 5)
    str <- vectorOf n (elements ['A'..'Z'])
    return $ UpperCaseString str

prop_EncodeDecode :: UpperCaseString -> Bool
prop_EncodeDecode (UpperCaseString str) =
  str `elem` decode (encode str)
```

Let's try again:

```
ghci> quickCheck prop_EncodeDecode  
=== prop_EncodeDecode from test/Test.hs:84 ===  
+++ OK, passed 100 tests.
```

Testing Algebraic Data Types

An example

```
data Expr = Cst Int
          | Add Expr Expr
  deriving (Eq, Show)
```

```
eval :: Expr -> Int
```

```
eval (Cst n) = n
```

```
eval (Add x y) = eval x + eval y
```

```
prop_CommutativeAddition :: Expr -> Expr -> Property
```

```
prop_CommutativeAddition x y =
```

```
  eval (Add x y) == eval (Add y x)
```

Incomplete: We need to implement `Arbitrary` for `Expr`.

Generating ASTs: the naïve way

```
instance Arbitrary Expr where
  arbitrary = oneof [ fmap Cst arbitrary
                    , do x <- arbitrary
                      y <- arbitrary
                      return $ Add x y
                    ]
```

Issue: This can generate huge expressions!

Generating ASTs: the good way

Solution: Sized generators.

```
-- | A generator for 'Expr's with an upper bound on  
-- approximately the number of nodes.
```

```
exprN :: Int -> Gen Expr  
exprN 0 = fmap Cst arbitrary  
exprN n = oneof  
    [ fmap Cst arbitrary  
    , do x <- exprN (n `div` 2)  
        y <- exprN (n `div` 2)  
        return $ Add x y  
    ]  
  
instance Arbitrary Expr where  
    arbitrary = sized exprN
```

QuickCheck Mini for Erlang

We will use Quviq QuickCheck Mini for Erlang in this course.

You will need to install it on your computer. A guide is available on Absalon.

The `dict` module in Erlang implements a purely functional hash table.

- ▶ `dict:new()` creates an empty dictionary.
- ▶ `dict:store(Key, Value, OldDict)` creates a new dictionary from `OldDict` where `Value` is stored at `Key`.
- ▶ `dict:store(Key, Dict)` retrieves a value from a dictionary.
- ▶ ... (see documentation)

Let us write some tests for `dict` using QuickCheck Mini.

Generating dictionaries

```
key() ->
    eqc_gen:oneof([atom(), eqc_gen:int(), eqc_gen:real()])
value() ->
    eqc_gen:oneof([eqc_gen:int(), atom()]).
atom() ->
    eqc_gen:elements([a,b,c,d]).
dict_0() ->
    ?LAZY(
        eqc_gen:oneof([
            dict:new(),
            ?LET({K,V,D}, {key(), value(), dict_0()}),
            dict:store(K,V,D))
        ])
    ).
```

Some analogies to QuickCheck for Haskell

- ▶ `key()`, `value()`, and `atom()` in above example are generators in the same sense as QuickCheck for Haskell.
- ▶ `?LET(v, gen(), expr)` should be thought of as equivalent to

```
do v <- gen
    expr
```
- ▶ Non-generators are for the most part automatically coerced into generators with constant value. This is unlike QuickCheck for Haskell, where you have to manually use `return`.
- ▶ `?LAZY()` defers execution of code describing a generator. Without it in the `dict` example, we would get a stack overflow because Erlang has eager evaluation.

Testing some properties

```
no_duplicates(Lst) ->  
  length(Lst) == length(lists:usort(Lst)).  
  
prop_unique_keys() ->  
  ?FORALL(D, dict_0(),  
    no_duplicates(dict:fetch_keys(D))).
```

Let's try to run our test:


```

> eqc:quickcheck(d:prop_unique_keys()).
.....Failed! After 67 tests.
{dict,8,16,16,8,80,48,
  {[],[],[],[],[],[],[],[],[],[],[],[],[],[],[]},
  {[[-2.0|-13]], [],
    [[b|d]], [[c|c],[3.2|-6]], [[d|7]],
    [],[],[],[],[], [[-2|-4],[14|b]],
    [],[], [[-9|12]], [],[]}}}
Shrinking xxxx.xx.xxxxxxxxxxxxxxxxxxxxxxxxxxxxxxx(5 times)
{dict,2,16,16,8,80,48,
  {[],[],[],[],[],[],[],[],[],[],[],[],[],[],[]},
  {[[-2.0|0]], [],[],[],[],[],[],[],[],[],
    [[-2|0]], [],[],[],[],[]}}}
false

```

Symbolic testing

Issue: the counterexample is complete nonsense.

- ▶ To get more readable counterexamples, it sometimes makes sense to construct random structures symbolically.
- ▶ In Erlang, we might represent a call as `{call, Module, Function, ArgumentList}`
 - Such representations of symbolic calls can be evaluated with the `eqc_symbolic:eval` function.

▶ Examples:

Function call	Symbolic call
<code>dict:new()</code>	<code>{call, dict, new, []}</code>
<code>dict:store(K,V,D)</code>	<code>{call, dict, store, [K,V,D]}</code>
<code>lists:usort([3,2])</code>	<code>{call, list, usort, [[3,2]]}</code>

Symbolic testing for dict

```
dict_1() ->
  ?LAZY(
    oneof([call(dict,new,[]),
          ?LET(D, dict_1(),
               {call(dict,store,[key(),value(),D])})])
  ).
```

```
prop_unique_keys() ->
  ?FORALL(D,dict_1(),
    no_duplicates(dict:fetch_keys(
      eqc_symbolic:eval(D))))).
```

Testing the quality of your generator

```
prop_measure() ->  
  ?FORALL(D,dict(),  
    collect(length(dict:fetch_keys(eval(D))),true)).  
  
prop_aggregate() ->  
  ?FORALL(D,dict(),  
    aggregate(eqc_symbolic:call_names(D), true)).
```

```
> eqc:quickcheck(d:prop_measure()).
```

```
.....  
OK, passed 100 tests
```

```
54% 0
```

```
28% 1
```

```
8% 2
```

```
7% 3
```

```
1% 12
```

```
1% 6
```

```
1% 4
```

```
true
```

```
> eqc:quickcheck(d:prop_aggregate()).
```

```
.....
```

```
OK, passed 100 tests
```

```
50.5% {dict,store,3}
```

```
49.5% {dict,new,0}
```

```
true
```

More improvements

- Statistics indicate that we often generate very small or empty dictionaries. Use `eqc_gen:frequency` to attain finer control over generation:

```
dict_2() ->
  ?LAZY(
    eqc_gen:frequency([
      {1, {call, dict, new, []}},
      {4, ?LET(D, dict_2(),
        {call, dict, store, [key(), value(), D]})}
    ])
  ).
```

- Better counterexamples with shrinking:

`dict_3() ->`

```
?LAZY(  
    eqc_gen:frequency([  
        {1,{call,dict,new,[]}},  
        {4,?LETSHRINK([D],[dict_3()],  
            {call,dict,store,[key(),value(),D]})}}]  
    ).
```

→ `?LETSHRINK([A, B, ...], [f1(), f2(), ...], expr)` is similar to `?LET([A, B, ...], [f1(), f2(), ...], expr)` but will try to “shrink” a counterexample to one of `A, B, ...`. Next lecture will have more on shrinking.

Returning to our property

```
no_duplicates(Lst) ->  
  length(Lst) == length(lists:usort(Lst)).  
  
prop_unique_keys() ->  
  ?FORALL(D,dict_3(),  
    no_duplicates(dict:fetch_keys(eval(D))))).
```

What is the problem?

```
> eqc:quickcheck(d:prop_unique_keys()).  
Failed! After 1 tests.  
{call,dict,store,  
  [0.0,b,{call,dict,store,[0,c,{call,dict,store,  
    [0,0,{call,dict,new,[]}]}}]}]  
Shrinking ..(2 times)  
{call,dict,store,  
  [0.0,a,{call,dict,store,[0,a,{call,dict,store,  
    [0,0,{call,dict,new,[]}]}}]}]  
false
```

In the dict module, two keys are different if they don't match according to `==`. However, `lists:usort/1` function uses `==` for equality.

```
> 1 == 1.0.
```

```
false
```

```
> 1 == 1.0.
```

```
true
```

Summary

- ▶ Instead of writing a lot of unit tests, write down a number of properties for your code and test them on random inputs.
- ▶ The QuickCheck library for Haskell (and QuickCheck Mini for Erlang) streamlines the process.
- ▶ To generate random data, you have to write so-called *generators*.
- ▶ Generators are usually built out of existing generators by use of the monad structure and other provided functions.