Trampoline Training



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May 27, 2015

Note: All the software and documents are stored at http://www.irccyn.ec-nantes.fr/~bechenne/trampoline

1 Goal

The goal of this training is to become familiar with OSEK/VDX applications development process and with Trampoline. Trampoline is a Free Software implementation of the OSEK/VDX specification. Trampoline includes an OIL compiler which allows, starting from an OIL description, to generate OS level data structures of the application. In addition to the OIL description, the developer must provide the C sources of tasks and ISRs of the application. Trampoline runs on many hardware platforms and we will use it on the Cortex-M4 STM32F4 Discovery board. If you have not installed Trampoline yet, get the Trampoline Package and read the install document.

The source code is located in the labs/labs_stm32F4_discovery directory.

^{*}http://fancyferret.deviantart.com/art/Trampoline-Beaver-166211545, CC BY-NC 3.0 license, Josef Ek.

2 The board

We are going to use a demo board made by ST, the STM32F4 Discovery, with a Cortex M4 STM32F407 micro-controller. Here is a picture of the demo board:



There are 4 LEDs located below the micro-controller. LED3 is the orange one, LED4 is the green one, LED5 is the red one and LED6 is the blue one.

On the left, there is a blue button labelled User that can be used for user interaction. On the right, there is a black button which is the reset button.

The LEDs and the blue button are connected to the GPIO of the micro-controller. The GPIO is initialized with the LEDs as output and the button as input. The input corresponding to the button may be configured as an external interrupt line. The initialization is done by calling the <code>initBoard</code> function. The argument of this function may be <code>BUTTON_NOIT</code> to configure the corresponding GPIO input as a normal input or <code>BUTTON_IT</code> to configure the input as an external interrupt line. In summary <code>initBoard(BUTTON_NOIT)</code>; or <code>initBoard(BUTTON_IT)</code>; should be put in the main function before starting Trampoline as shown below:

```
FUNC(int, OS_APPL_CODE) main(void)
{
   initBoard(BUTTON_NOIT);
   StartOS(OSDEFAULTAPPMODE);
   return 0;
}
```

For the labs, functions are provided to switch on, switch off and toggle the LEDs. The unique argument is the LED identifier and should be LED3, LED4, LED5 LED6 or ORANGE, GREEN, RED, BLUE:

```
void ledOn(<led>) turns on LED <led>.
void ledOff(<led>) turns off LED <led>.
void ledToggle(<led>) toggles LED <led>.
```

A function gives the state of the User blue button. It returns BUTTON_PRESSED if the button is pressed and BUTTON_RELEASED if not.

```
ButtonState readButton(); returns the state of the User blue button.
```

At last a function called delay waits for an amount of time expressed in milliseconds.

```
void delay(<howManyMs>); waits <howManyMs> ms.
```

Will will use this function to slow down the application.

2.1 A word about memory sections

AUTOSAR defines a way to put objects: constants, variables and functions in memory sections in a portable way¹. For that, a set of macro are used along with a generated file: MemMap.h. Functions should be declared with the FUNC macro, variables with the VAR macro, constants with the CONST macro and pointers to variables, pointers to constant, constant pointers to variable and constant pointers to constant with P2VAR,

¹memory section declaration is not part of the C standard

P2CONST, CONSTP2VAR and CONSTP2CONST respectively. Sections are opened and close with a macro definition and the inclusion of the tpl_memmap.h file. For instance:

```
#define APP_Task_my_periodic_task_START_SEC_VAR_32BIT
#include "tpl_memmap.h"

VAR(int, AUTOMATIC) period;
VAR(int, AUTOMATIC) occurence;
#define APP_Task_my_periodic_task_STOP_SEC_VAR_32BIT
#include "tpl_memmap.h"
```

defines variables period and occurence in the variables section of task my_periodic_task.

```
#define APP_Task_my_periodic_task_START_SEC_CODE
#include "tpl_memmap.h"

TASK(my_periodic_task)
{
    ...
    TerminateTask();
}
#define APP_Task_my_periodic_task_STOP_SEC_CODE
#include "tpl_memmap.h"
```

defines the task my_periodic_task in the code section of task my_periodic_task. goil generates the sections for tasks according to the description.

3 Basic tasks

Go into the lab1 directory. There are 2 files:

lab1.oil the OIL description of the lab1 application;

lab1.c the C source for the lab1 task.

Edit the lab1.oil and look at the TRAMPOLINE_BASE_PATH attribute (in OS > BUILD attribute). TRAMPOLINE_BASE_PATH is set to "../..". If you move around the lab1 directory you will have to update this attribute.

lab1 is a very simple application with only 1 task called a_task. a_task starts automatically (AUTOSTART = TRUE in the OIL file). Look at the OIL file and the C source file.

```
To compile this application, go into the lab1 directory and type:

goil -t=thumb2/cortex-m4/STM32F4-Discovery lab1.oil
```

goil is the OIL compiler. It parses the OIL file and produces a set of C files. The -t option gives the target system. thumb2 is the instruction set of the target, cortex-m4 is the micro-controller core and STM32F4-Discovery is the board.

thumb2/cortex-m4/STM32F4-Discovery is a path inside the machines directory and in the templates directory. The OIL file gives the names of the C source files (with APP_SRC and the name of the executable file (with APP_NAME).

This generate a Makefile for the application too. It has to be done only once. If you change something in the OIL file or in your C file, you do not need to rerun the goil compiler by hand because make will run it when needed. Then type:

make

The application and Trampoline OS are compiled and linked together. To load the application on the target, type:

make burn

The application may or may not start :-). Press the reset button if it does not start.

In this application, there is only one task called a_task which switches LED3 on.

```
TASK(a_task)
{
    ledOn(LED3);
    TerminateTask();
}
```

4 OS system calls and task launching

4.1 Task activation and scheduling

The ActivateTask() system call allows to activate another task of the application.

Go into the lab2 directory.

In lab2.oil and lab2.c, 2 tasks have been added: task_0 (priority 1) and task_1 (priority 8). task_0 toggles LED4 on and task_1 toggles LED5 on. Task a_task activates task_0 and task_1. All statements are separated by a busy-wait loop on the button so that by pressing the button we can control the execution. Examine the OIL and the C files.

Compile and execute. Why does task_1 execute before task_0 whereas it has been activated after?

4.2 Task chaining

The ChainTask() system call allows to chain the execution of a task to another one. This is roughly the same thing as calling ActivateTask and TerminateTask at the same

time.

Replace the call to TerminateTask by a ChainTask(task_1) at the end of task a_task. What is happening?

Chain to task_0 instead of task_1. What is happening?

Test the error code returned by ChainTask and correct your program to handle the error. ChainTask may return the following codes:

E_OS_ID the target task does not exist;

E_OS_RESOURCE the calling task holds a resource;

E_OS_CALLEVEL not called from a task;

E_OS_LIMIT too many activations of the target task.

4.3 Pre-task and Post-task hooks

Hook routines are used to insert application functions inside the kernel. Hook routines are called by the kernel when a particular event happens. The Pre-task hook is called when a task goes into the running state. The Post-task hook is called when a task leaves the running state. Hooks are useful for debugging purpose.

There are two boolean attributes in the OS object of the OIL to use Pre-task and Post-task hooks:

```
OS config {
STATUS = EXTENDED;
PRETASKHOOK = TRUE;
POSTTASKHOOK = TRUE;

...
```

When these hooks are used, the user have to write a hook functions:

```
1 FUNC(void, OS_CODE) PreTaskHook()
2 {
3    ...
4 }
```

for the pre-task hook and

```
1 FUNC(void, OS_CODE) PostTaskHook()
2 {
3    ...
4 }
```

Go into the lab3 directory and play with the application which is in it. It is the same application as in lab2 with the use of delay instead of busy-wait for the blue button. Pre and post task hooks are used to switch a led corresponding to the running task on.

4.4 Extended tasks and synchronization using events

Go into the lab4 directory.

This application has 2 tasks: a_task which is an extended task and task_0 which is a basic task. Unlike a basic task, an extended task may wait for an event. A task is extended because it has at least one event declared in the OIL file.

Look at the OIL file and at the C file. Task a_task activates task task_0 then goes into an infinite loop where it waits for event ev_0 and activate task task_0 again. When a_task runs, LED3 is switched on. When task_0 runs, LED4 is switched on. The application should run infinitely.

Question 1 Draw the Gantt diagram of the execution.

Question 2 However The application stops. What is happening? Correct the OIL file to have a proper behavior.

For the following application, we will use WaitEvent, GetEvent and ClearEvent.

Question 3 Extend the previous application by adding 1 task: task_1 (priority 1) and 1 event ev_1. a_task activates task_0 and task_1 and waits for one of the events. When one of the events is set, a_task activates the corresponding task again.

5 Alarms and periodic tasks

5.1 First application, basic use of alarms

Go into the lab5 directory.

Alarms are periodic activities. They are used to build periodic tasks. Alarms may activate a task or set an event. The application in the lab5 directory is a simple one that blinks the LED6 with a 500ms period. The SysTick is every 1ms. So the period of the alarm is 250ms. Examine the lab5.oil file and the lab5.c file. The underlying counter SystemCounter have a TICKSPERBASE attribute. This attribute is the number of SysTick needed to increment the counter by one. By default it is set to 1. Pre and Post tasks hooks are used. Add a delay(50); juste after ledToggle(BLUE); in order to see the execution of the task.

Question 4 Increase the TICKSPERBASE to 25 and change the ALARMTIME and CYCLETIME of the alarm blink_blink to keep the same period.

Question 5 Program an application using 2 periodic tasks, t1 and t2. t1 switches LED6 on and t2 switches LED6 off. Both tasks have the same period (1s) but have an offset so that LED6 is on during 150ms.

5.2 Second application, starting and stopping alarms

Go into the lab6 directory.

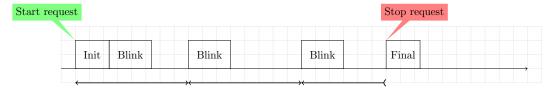
This application reads the **blue button** using a periodic task that is activated every 50ms. If the button is pressed, the LED6 is toggled.

Using services GetAlarm, SetRelAlarm and CancelAlarm, build an application with the following requirements:

- Alarm blink_blink activates task blink. This alarm is not an AUTOSTART one.
- Task blink toggles LED4.
- If the button is pressed and alarm blink_blink is not yet started, start the alarm with a period of 100ms.
- If the button is pressed and alarm blink_blink is started, cancel the alarm.

5.3 Third application

Requirements change. Now blinking needs an Init code (runs once when the alarm is started) and a Final code (runs once when the alarm is stopped). This corresponds to the following diagram:



Question 6 Modify the application to take the new requirements into account. Use 3 basic tasks. Init and Final blink their own LED once (use delay to see the execution).

Question 7 Same question but with only one extended task.

5.4 Chase

Program a chase² with a 0.5s period. To do it, use 4 periodic tasks. Each periodic task manages a LED. The chase effect is done by using alarms with a time shift between them.

²chenillard in French

When the blue button is pressed, the chase direction changes.

6 Shared object access protection

6.1 Goal

To show resources usage, we will use a bad program that allows to corrupt a shared global variable which is not protected against concurrent writes. We will see different ways to prevent this wrong behavior by using resources (standard and internal) or other solutions (preemption and priority).

6.2 Application requirements

The application has 3 tasks and 2 **volatile** global variables: **val** and **activationCount** as shown in figure 1:

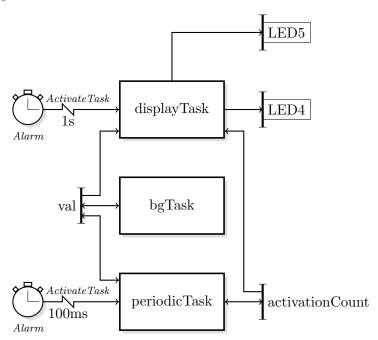


Figure 1: Application diagram

- a background task called bgTask, active at start (AUTOSTART = TRUE) and that never ends. In an infinite loop this task increments then decrements the global variable val. This task has a priority equal to 1.
- a periodic task called periodicTask, priority 10, that runs every 100ms. This periodic task increments the global variable activationCount which is initialized

to 0 at start. Then if activationCount is odd, val is incremented, otherwise it is decremented.

• a periodic task displayTask, priority 20, that runs every second. If val is inside interval [-1;2], LED4 is switched on and LED5 is switched off. Otherwise, LED5 is switched on and LED4 is switched off.

Describe the application in OIL and program it in C.

Question 8 Does the behavior correspond to what you expect? Why?

6.3 Global variable protection

Question 9 Update the OIL file and the C program to protect the access to the global variable val. Use a resource to do it.

The resource priority is automatically computed by goil according to the priorities of the tasks which use it.

Question 10 What priority will be given to the resource?

The OIL compiler (goil) generates many files in the directory bearing the same name as the oil file (less the .oil suffix). Among them 3 are interesting:

- tpl_app_define.h
- tpl_app_config.h
- tpl_app_config.c

The file tpl_app_config.c contains the tasks' descriptors as long as all other data structures. These structures are commented.

Question 11 For each task, find the priority computed by goil and the identifier. Is is the same as defined in the OIL file? if not is it a problem?

Question 12 What is the priority of the resource? Is it compliant with the PCP rule?

6.4 Protection with an internal resource

An internal resource is automatically taken when the task gets the CPU. Replace the standard resource by an internal resource in the OIL file. Remove the GetResource and ReleaseResource in the C file.

Question 13 What happens? Why?

Modify the task bgTask: instead of infinite loop, use a ChainTask to the bgTask (ie: the task chains to itself).

7 Adding a service into Trampoline

In this part we are going to see how to add semaphore services into Trampoline. We want to services: SemWait to lock a semaphore and SemPost to unlock a semaphore.

7.1 Adding a semaphore object

Let's start with the data.

The first thing to do is to provide an object type for semaphores in C. We need a struct with four members. token is the current number of token available. size is the number of tasks waiting in waiting_tasks. index is write index in waiting_tasks. waiting_tasks is a ring buffer to store the tasks waiting for the semaphore.

```
typedef struct {
  uint32   token;
  uint32   size;
  unit32   index;
  tpl_task_id waiting_tasks[TASK_COUNT];
} tpl_semaphore;
```

TASK_COUNT is computed by goil and is equal to the number of tasks in the application.

The second thing is to provide an object type for semaphores in the OIL description. This is done in the IMPLEMENTATION part of the OIL description. Normally the OIL standard does not allow to create new object types. This has been relaxed in goil.

```
1 SEMAPHORE [] {
2 UINT32 TOKEN;
3 };
```

The declare a new object type: SEMAPHORE with an UINT32 attribute, TOKEN. The [] means multiple instances of semaphore can be used. Now, semaphores can be declared in the OIL file:

```
1 SEMAPHORE sem1 { TOKEN = 3; };
2 SEMAPHORE sem2 { TOKEN = 1; };
```