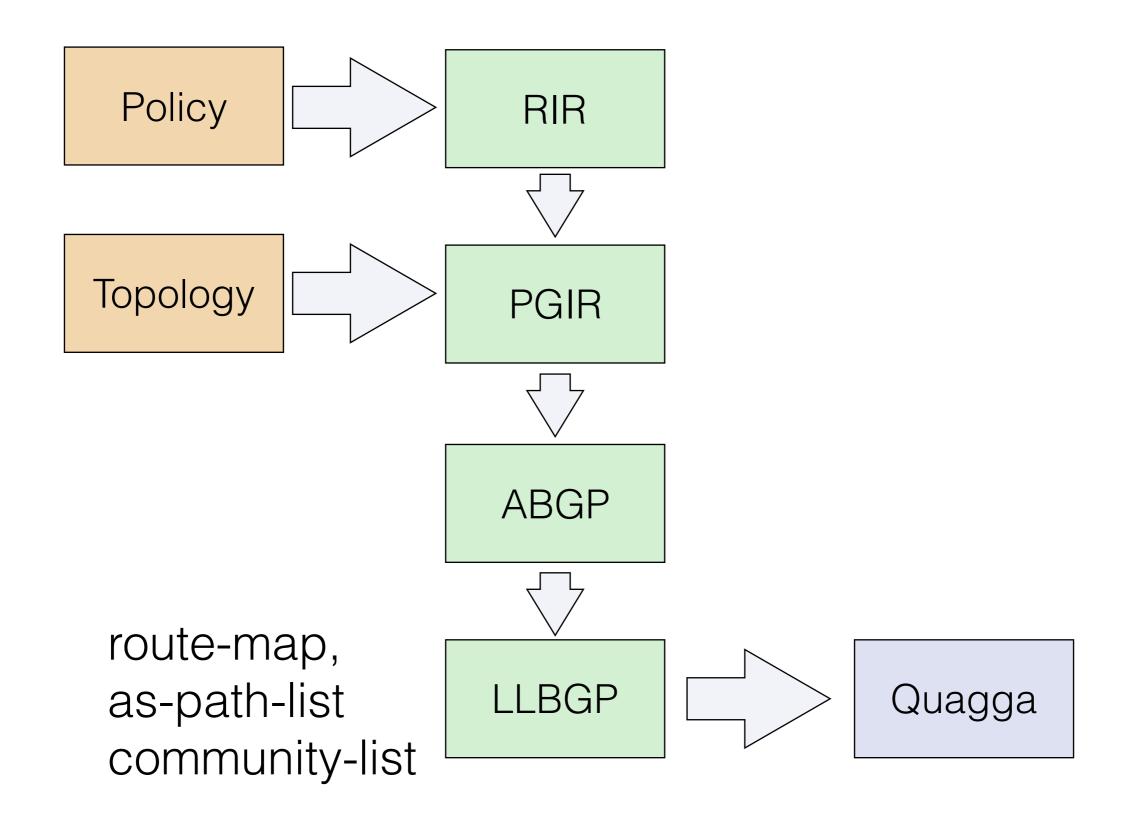
Implementation

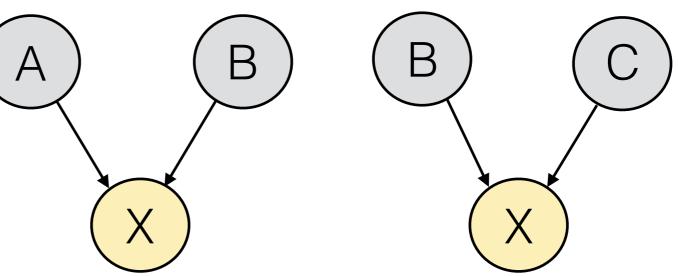


Quagga

```
ip prefix-list pl-1 permit 0.0.0.0/0 ge 0 le 32
ip community-list standard cl-1 permit 100:1
ip community-list standard cl-2 permit 100:2
ip as-path access-list path-1 permit (^103_ | ^102_)
route-map rm-in permit 10
 match community cl-1
 match ip address prefix-list pl-1
  set community additive 200:14
route-map rm-in permit 20
 match community cl-2
 match ip address prefix-list pl-1
 set local-preference 99
  set community additive 200:14
```

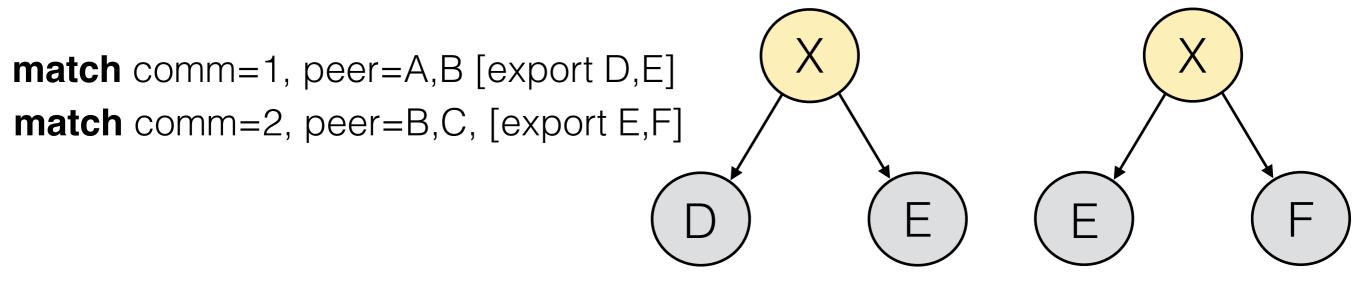
Quagga

match comm=1, peer=A,B, ...
match comm=2, peer=B,C, ...



```
ip as-path access-list peerAB permit (^A_ | ^B_) ip as-path access-list peerBC permit (^B_ | ^C_) ip community-list standard cl-1 permit 100:1 ip community-list standard cl-2 permit 100:2 route-map rm-in permit 10 match as-path peerAB match community cl-1 set community additive 100:3 !
```

Quagga



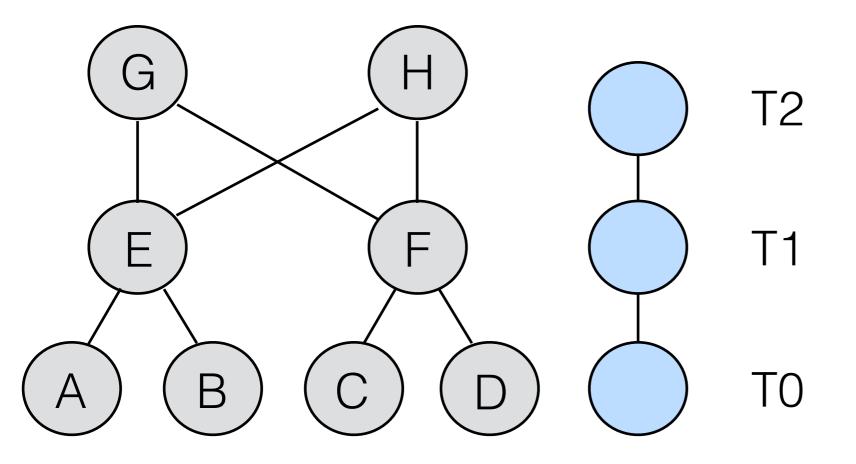
How do we know when to export to a peer? Need a per-peer export policy

- (1) New route-map for every peer
- (2) Assign community for at import filter for each peer.

Idea:

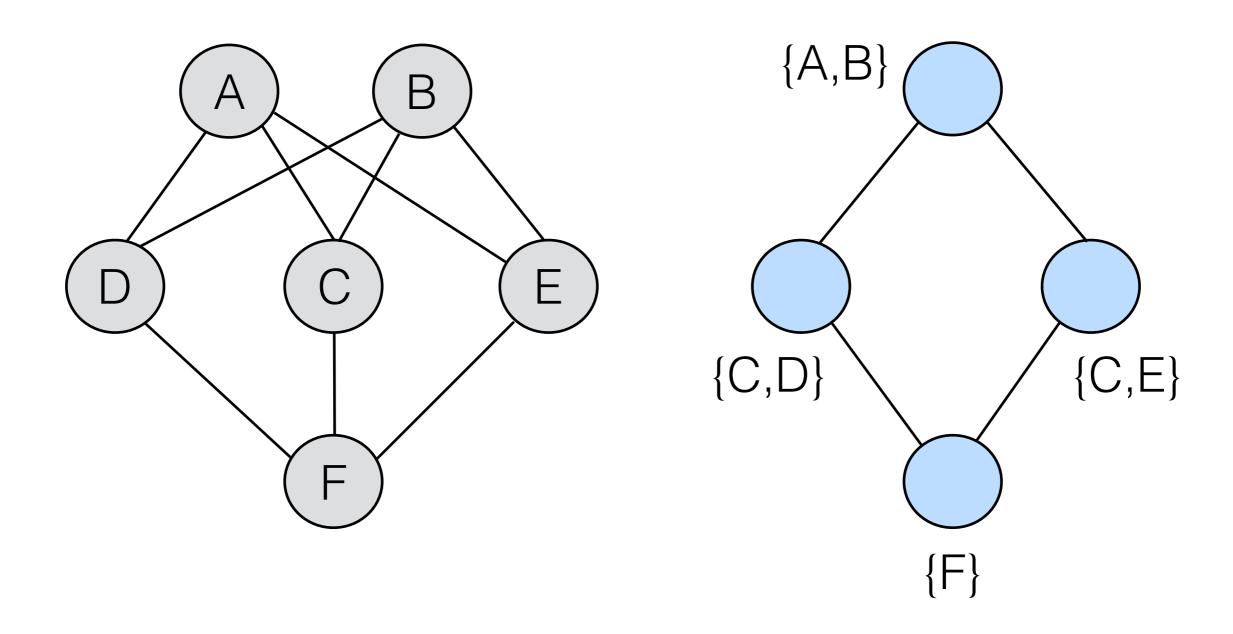
Each abstract node corresponds to a set of concrete nodes

Example: T2; T1; T0 (G + H); (E + F); (A + B + C + D)



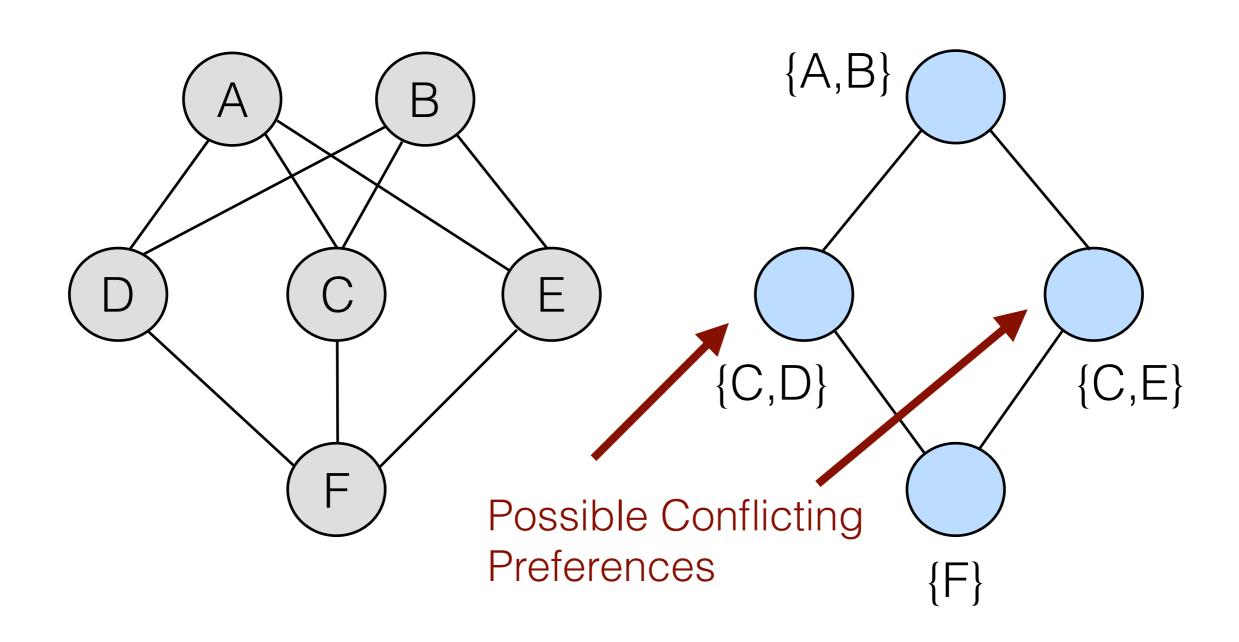
Attempt 1:

Each abstract node corresponds to a set of concrete node



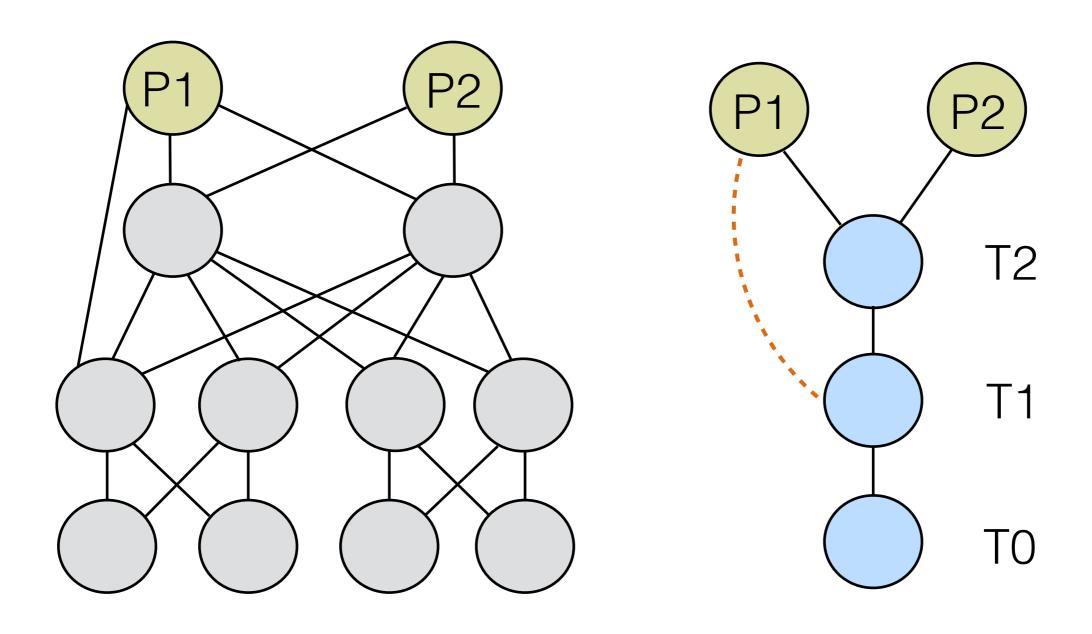
Attempt 1:

Each abstract node corresponds to a set of concrete node



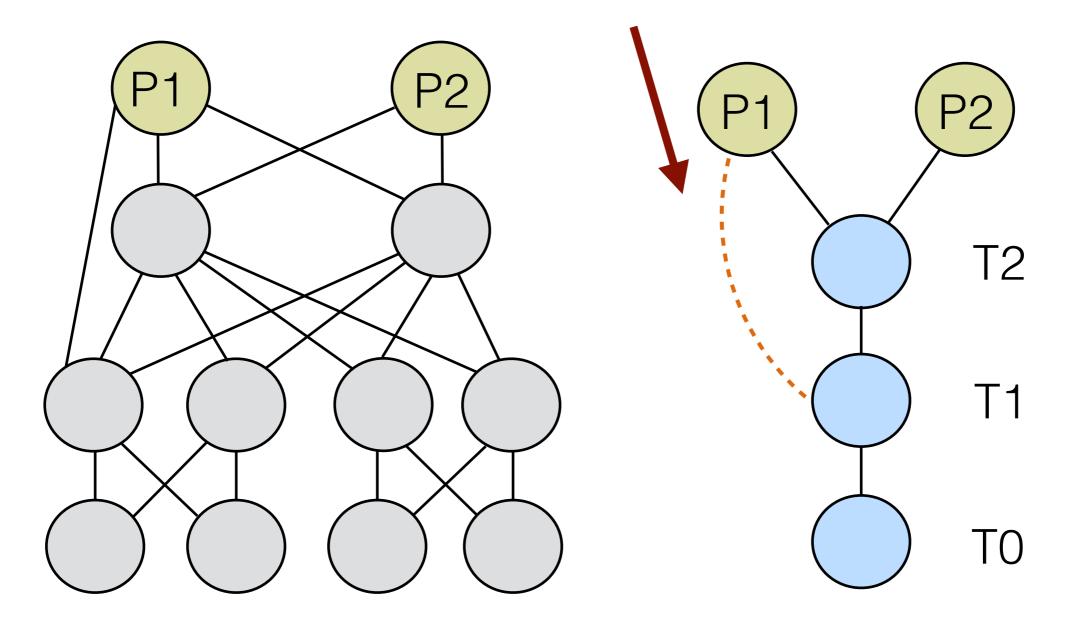
Attempt 2:

Not including a link in the abstract topology



Policy: out; T1; T0

Path exists in concrete network but won't be allowed in abstract network



Criteria:

```
Concrete Topology Tc: (Vc, Ec)
Abstract Topology Ta: (Va, Ea)
```

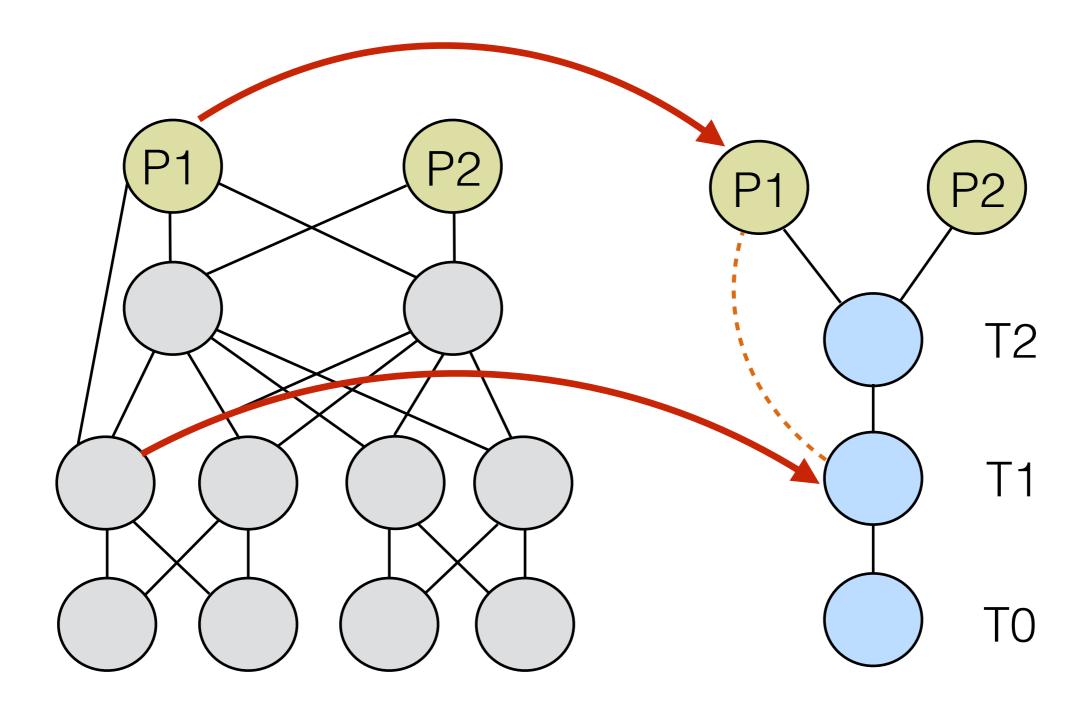
```
Looking for a mapping f: Vc —> Va
s.t. if v1,...,vn in Tc then f(v1),...,f(vn) in Ta
```

Graph Homomorphism

```
mapping f: Vc \longrightarrow Va
s.t. if (x,y) in Ec then f(x),f(y) in Ea
```

Satisfies the above criteria

Policy: out; T1; T0

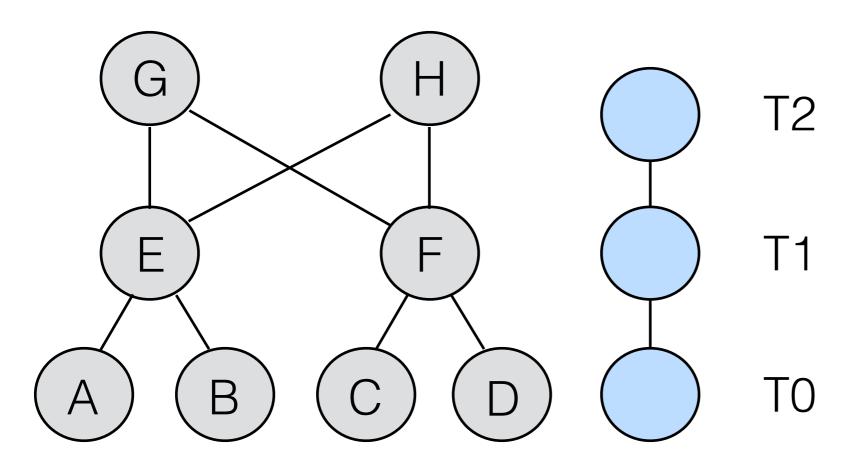


Correctness:

Each abstract node corresponds to a set of concrete nodes

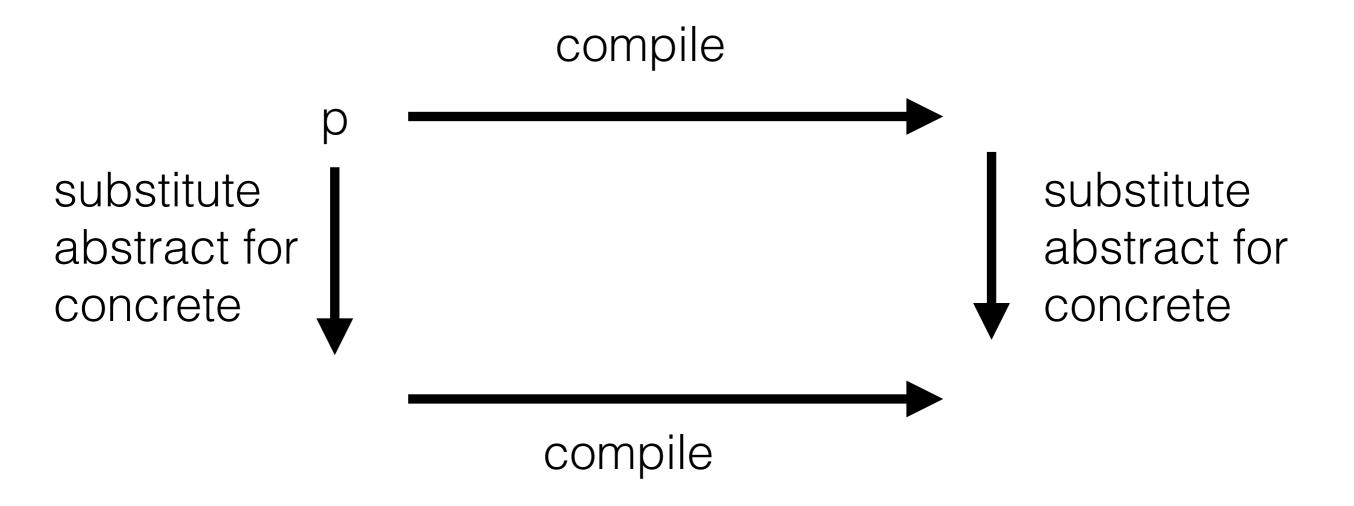
Example: T2; T1; T0

$$(G + H); (E + F); (A + B + C + D)$$



Correctness:

Concrete alphabet, Σc Abstract alphabet Σa Abstract policy p compile($p[x_i / \cup f^{-1}(x_i)]$)



Correctness:

Concrete alphabet, Σc Abstract alphabet Σa Abstract policy p

```
compile(p[xi / \cup f<sup>-1</sup>(xi)]) = compile(p)[xi / \cup f<sup>-1</sup>(xi)])

Policy substitution

ABGP substitution
```

Observations:

For every path through Tc, there is path through Ta Every path through PGc is a valid path through Tc Every path through PGa is a valid path through Ta For every path through PGc, there is a path through PGa

Idea:

Show that whenever certain properties hold in the abstract PG, they also hold in the concrete PG

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Show that whenever certain properties hold in the abstract PG, they also hold in the concrete PG

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Show that whenever certain properties hold in the abstract PG, they also hold in the concrete PG

Property X:

```
Given a PG that compiles and nodes X1, X2 such that shadows(X1,X2) and X1 \geq X2 for every path p = a_1,...,a_i, X1, a_{i+1},...,a_n either (1) there exists a path q = b_1,...,b_j, X2, a_{i+1},...,a_n where b_1,...,b_j, X is simple (2) there is no simple path to X2 (3) there exists a path q = b_1,...,b_{j-k}, a_{i+1+h},...,a_n with rank(p) \geq rank(q)
```

Property X —— Completeness

Property X (PGa) — Property X (PGc)