

Algorithms

Algorithms: An *algorithm* is a sequence of unambiguous instructions for solving a problem.

Problems:

- No Solution
- One Solution
- **Multiple Solutions**

Multiple Solutions: Which solution amongst multiple is efficient to implement?

How to find efficiency??

Algorithm is associated with two types of efficiencies,

- **Time and Space.**
 - Time Efficiency: How fast an algorithm runs
 - Space efficiency: Extra space the algorithm requires

[One of the above is obsolete concept now. Can you recollect which one?
More importantly, why???

For Time efficiency, Is it running time of the algorithm??

Answer is **NO!!**

- It becomes machine dependent.
- Program dependent
- Compiler dependent etc.

Basic Operation is the measuring Unit.

What is it??

- **It the operation contributing the most to the total running time of the algorithm. Generally, statements in the innermost loops.**

Ex: 1

1. ALGORITHM Mystery(n)
2. //Input: A nonnegative integer n
3. $S \leftarrow 0$
4. for index $\leftarrow 1$ to n do
5. $S \leftarrow S + \text{index} * \text{index}$
6. return S

Basic Operation: $S \leftarrow S + \text{index} * \text{index}$ (Line no. 5)

Ex: 2

1. ALGORITHM *MaxElement*($\text{List}[0..n-1]$)
2. //Determines largest element in Array
3. //Input: An array $\text{list}[0..n-1]$ of real numbers
4. //Output: Value of largest element
5. $\text{max} \leftarrow \text{list}[0]$
6. for index $\leftarrow 0$ to $n-1$ do
7. if $\text{list}[\text{index}] > \text{max}$
8. $\text{max} \leftarrow \text{list}[\text{index}]$
9. return max

Basic Operation: if $\text{list}[\text{index}] > \text{max}$ (Line no. 7)

Ex: 3

1. ALGORITHM *SelectionSort*($\text{List}[0..n-1]$)
2. //Sorts a given array by selection sort
3. //Input: An array $\text{list}[0..n-1]$ of orderable elements
4. //Output: Array $\text{list}[0..n-1]$ sorted in ascending order
5. for index $\leftarrow 0$ to $n-2$ do
6. $\text{min} \leftarrow \text{index}$
7. for nextindex $\leftarrow \text{index}+1$ to $n-1$ do
8. if $\text{list}[\text{nextindex}] < \text{list}[\text{min}]$
9. $\text{min} \leftarrow \text{nextindex}$
10. swap $\text{list}[\text{index}]$ and $\text{list}[\text{min}]$

Basic Operation: if list[nextindex] < list[min]

We know how to identify the Basic Operation from a given algorithm. Let's know see how to calculate the efficiency of non-recursive algorithms!!

Ex 1:

1. ALGORITHM MaxElement(List[0..n-1])
2. //Determines largest element in Array
3. //Input: An array list[0..n-1] of real numbers
4. //Output: Value of largest element
5. currentmax ← list[0]
6. for index ← 0 to n-1 do
7. if list[index] > currentmax
8. currentmax ← list[index]
9. if list[index] > currentmax
10. currentmax ← list[index]
- 11.
12. return currentmax

Basic Operation: if list[index] > currentmax (Line no. 7)

1. Set up the summation:

$$C(n) = \sum_{\text{index}=0}^{n-1} 2 = 2 \sum_{\text{index}=0}^{n-1} 1 = 2n = O(n)$$

2. Solve the Summation

$$C(n) = \sum_{\text{index}=0}^{n-1} 1 = n-1+1 \in O(n)$$

$$\sum_{i=l}^u 1 = u - l + 1$$

$O(n)$

Efficiency of given algorithm is $O(n)$.

Ex 2:

1. ALGORITHM SelectionSort(List[0..n-1])
2. //Sorts a given array by selection sort
3. //Input: An array list[0..n-1] of orderable elements
4. //Output: Array list[0..n-1] sorted in ascending order
5. for index \leftarrow 0 to n-2 do
6. min \leftarrow index
7. for nextindex \leftarrow index+1 to n-1 do
8. if list[nextindex] < list[min]
9. min \leftarrow nextindex
10. swap list[index] and list[min]

1. Set up the summation:

$$C(n) = \sum_{i=0}^{n-2} \sum_{j=i+1}^{n-1} 1$$

2. Solve the Summation

$$C(n) = \sum_{i=0}^{n-2} \sum_{j=i+1}^{n-1} 1$$

$$= \sum_{i=0}^{n-2} n-1 - (i+1) + 1$$

$$= \sum_{i=0}^{n-2} n-1-i-1+1$$

$$i=0$$

$$= \sum_{i=0}^{n-2} n-1-i$$

$$= \sum_{i=0}^{n-2} a-b$$

$$= \sum_{i=0}^{n-2} a - \sum_{i=0}^{n-2} b$$

$$= \sum_{i=0}^{n-2} n-1 - \sum_{i=0}^{n-2} i$$

$$\sum_{i=1}^n i = 1+2+3+\dots+n = n(n+1)/2$$

$$= \sum_{i=0}^{n-2} n-1 - 0+1+2+\dots+n-2$$

$$= \sum_{i=0}^{n-2} n-1 - (n-2)(n-2+1)/2$$

$$= \sum_{i=0}^{n-2} n-1 - (n-2)(n-1)/2$$

$$= (n-1) \sum_{i=0}^{n-2} 1 - (n-2)(n-1)/2$$

$$\begin{aligned} &= (n-1)(n-2-0+1) - (n-2)(n-1)/2 \\ &= (n-1)(n-1) - (n-2)(n-1)/2 \\ &= (n-1)[(n-1) - ((n-2)/2)] \\ &= (n-1)((2n-2-n+2)/2) \\ &= (n-1)(n/2) = (n^2-n)/2 = n^2/2 - n/2 \\ &\in O(n^2) \end{aligned}$$

References:

Text Books:

1. Ananya Levitin – Introduction to design and analysis of Algorithms
2. Thomas H Cormen, Charles E Leiserson, Ronald L Rivest, Clifford Stein - Introduction to Algorithms
3. Goodrich and Tamassia - Algorithm Design
4. ELLIS AUTOR HOROWITZ, SARTAJ AUTOR SAHNI – Fundamentals of Computer Algorithms