

A Resource-Efficient 1024-Point MSC FFT Using Time-Multiplexed Constant Multiplication

First A. Author, *Member, IEEE*, Second B. Author, *Fellow, IEEE*, and Third C. Author, *Student Member, IEEE*

Abstract—This paper

Index Terms—FFT, constant multiplication, time-multiplexing, FPGA, ASIC, low-complexity, signal processing.

I. INTRODUCTION

II. BACKGROUND AND RELATED WORK

III. PROPOSED ARCHITECTURE

IV. IMPLEMENTATION AND OPTIMIZATION

V. EXPERIMENTAL RESULTS AND COMPARISON

To evaluate the efficacy of the proposed 4-parallel 1024-point MSC FFT architecture, we implemented it on a Xilinx Virtex 7 XC7VX330TFFG1157 FPGA platform and compared it against five state-of-the-art designs [1]–[5].

We employ the same FFT size, word length (WL) and degree of parallelism for a fair comparison, i.e., $N = 1024$, $WL = 16\text{bit}$ and $P = 4$. Table I presents a detailed comparison of hardware resource utilization and performance characteristics between the proposed 4-parallel 1024-point MSC FFT architecture and contemporary designs. The key metrics include the number of slices, LUTs, flip-flops (FFs), DSPs, Block-RAMs, maximum clock frequency (f_{CLK}), throughput (Th.), latency (in cycles and microseconds), signal-to-quantization noise ratio (SQNR), power consumption (P), and normalized power (NP). All architectures listed in the table are implemented on either Virtex-6 (V6) or Virtex-7 (V7) FPGA platforms.

In terms of hardware resources utilization, the proposed MSC architecture achieves competitive area efficiency. It utilizes only 12 DSPs and 4 BRAMs, significantly fewer than those required by previous MDC-based architectures. [1]–[3]. Among recent high-radix FFT designs, the proposed MSC architecture achieves the lowest Slices count, reducing area by 8.5% compared to the prior MSC implementation [5] and by 44% relative to the CM-based design [4]. Although the LUTs count of proposed architecture is comparable to [5], the proposed design reduces flip-flops usage by 17.3%, indicating that the use of time-multiplexed constant multiplication contributes to a more resource-efficient implementation.

Regarding operating frequency and power efficiency, the proposed implementation operates at 493 MHz, outperforming all MDC-based counterparts [1]–[3] and recent MSC implementation [5]. This enables a throughput of 1820 MS/s, surpassing [5] by 8.3%. The latency is measured at 307 clock cycles (0.62 μs), offering a balanced trade-off between pipeline depth and processing speed. In terms of power, the design consumes 1.16 W, slightly higher than [5] (0.98 W) but 31% lower than [4] (1.68 W). Critically, the normalized power metric (NP = power per MHz) of proposed design is 2.35 mW/MHz,

nearly identical to [5] (2.33 mW/MHz), confirming that the optimization using time-multiplexed constant multiplication contributes to low-power operation across frequency scaling.

To further quantify the trade-offs between hardware cost, speed, and energy efficiency, we introduce two composite metrics: the resource-delay product ($\text{RDP} = \text{Slices} \cdot T_{\text{CLK}}$) and the power-delay product ($\text{PDP} = \text{Power} \cdot T_{\text{CLK}}$), where $T_{\text{CLK}} = 1/f_{\text{CLK}}$. These metrics jointly capture area–time and power–time efficiencies, respectively, providing a more holistic view of design quality beyond isolated figures. The proposed MSC architecture achieves the lowest RDP among all compared designs, with a value of 2995.94, representing a 22.1% reduction over the previous best MSC implementation [5] and a 29.2% improvement over the CM-based approach [4]. As for PDP, the proposed design reports 2352.94 mW·ns, which is comparable to [5] and 5% lower than [4] —note that lower PDP is better. The near-parity in PDP between our work and [5], combined with significantly better RDP, demonstrates that the time-multiplexed constant multiplication offers a more balanced optimization across resources, speed, and power dimensions.

With respect to numerical accuracy, the SQNR of the proposed MSC FFT (49.46 dB) is comparable to that of [5] (50.16 dB), differing by less than 1 dB, yet considerably higher than the 40.30 dB reported in [1]. This demonstrates that the time-multiplexed constant multiplication strategy employed in the MSC architecture effectively preserve signal fidelity despite resource reduction.

VI. CONCLUSION

This brief presents a 4-parallel 1024-point MSC FFT architecture optimized by time-multiplexed constant multiplication. Experimental results on a Xilinx Virtex-7 FPGA show that the proposed architecture outperforms several state-of-the-art FFT implementations in terms of resources efficiency, operating frequency, power consumption, and accuracy. The use of time-multiplexed constant multiplication proves to be an effective strategy for resource-efficient FFT implementations.

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TABLE I
COMPARISON OF HARDWARE RESOURCES AND PERFORMANCE FOR
4-PARALLEL 1024-POINT PIPELINED FFT IMPLEMENTED ON FPGA

	[1]	[2]	[3]	[4]	[5]	Proposed
Architecture	MDC	MDC	MDC	CM	MSC	MSC
Radix	2^5	2^2	2	2^5	2^5	2^5
WL	16	16	16	16	16	16
Slices	1420	1351	-	2631	1615	1477
LUTs	-	-	4116	-	4682	4629
FFs	-	-	1920	-	5910	4887
DSPs	16	48	72	12	12	12
Block-RAMs	12	12	0	0	4	4
f_{CLK} (MHz)	253	227	380	680	420	493
Th. (MS/s)	1012	910	1520	2720	1680	1820
Latency (cyc.)	265	285	767	394	300	307
Latency (μ s)	1.04	1.25	2.02	0.58	0.71	0.62
SQNR (dB)	40.30	-	-	-	50.16	49.46
P (W)	-	-	-	1.68	0.98	1.16
NP ($\frac{mW}{MHz}$)	-	-	-	2.47	2.33	2.35

TABLE II
COMPARISON OF RESOURCE-DELAY PRODUCT (RDP) AND
POWER-DELAY PRODUCT (PDP). $PDP = POWER \cdot T_{CLK} (mW \cdot ns)$, RDP
 $= SLICES \cdot T_{CLK} (SLICES \cdot ns)$

	[1]	[2]	[3]	[4]	[5]	Proposed
Architecture	MDC	MDC	MDC	CM	MSC	MSC
Radix	2^5	2^2	2	2^5	2^5	2^5
PDP(mW·ns)	-	-	-	2470.61	2333.38	2352.94
RDP	5612.69	5951.56	-	3869.15	3845.32	2995.94

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