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# **Normal Forms**

# **Functional Dependencies**

For a table [x,y] y is functionally dependent on x if, when  $t_1(x) = t_2(x)$ ,  $t_1(y) = t_2(y)$ 

Here, \$x\$ is called determinant and \$y\$ is called dependent

#### First Normal Form 1NF

- 1. Each attribute must be atomic (ie should not be divisible any further)
  - 1. Composite attributes get separated to their own distinct attributes
  - 2. Multi Valued attributes get their own tuple or row
- 2. A attribute should contain values of the same type
- 3. attributes must have unique names
- 4. No order towards the rows should affect the database
- By default if a ER diagram is converted into a relational schema, it is in 1NF

### Second Normal Form 2NF

- 1. Should be in 1NF
- 2. No Partial Dependency must be present in the table

#### Partial Dependency

- Partial Dependency is when A proper subset of the candidate key will determine a non prime attribute
- Prime attributes are ones that are part of Candidate key

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RCA, B, C, D) F.D: {AB -> CD, C-> A, D-> B3 taking closure of entire set, ABCD = { ABC D3; now versue to defendants. > ABQ = {ABCDJ; AB is a super key now check if sujex key's the subset are sujex keys . At = A only St= B only as its subsets aren't sujer keys, AB is a condidate boy To check for more condidate keys," if a prime attribute is the defendent of on other attribute, the cull be another R As both A, B are defendent on C, D respectivly. we will reflece them and reject which for condidate keep CB+= 2ABC D3 7 C+= 2C, A3 7 / 15 a Cx 1010) pune Atubbes D+ : 1 B5 ADT = { A DC 03 ( At-2Ay / 15 a Ck d= 2B,03

Third Normal Form 3NF

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- 1. Should be in 2NF
- 2. There shouldn't be any Transitive Dependency for Non Prime Attributes

A table is in 3NF if and only if for each of its non trivial functional dependencies, atleast one of the following conditions holds

- 1. LHS is Super Key
- 2. RHS is Prime Attribute

# Boyce Codd Normal Form BCNF

- 1. Should be in 3NF
- 2. For each Non trivial Functional dependency, the Determinant must be a super key

### Rules of Inference

- 1. Reflexivity Rule
  - 1. If \$Y \subseteq X\$, then \$X \rightarrow Y\$
  - 2. If Y is a subset of X, X functionally determines Y
  - 3. Ex:  $X = \{A,B\}$  then  $\{A,B\} \rightarrow \{A,B\}$  and  $\{A,B\} \rightarrow \{B\}$
- 2. Augmentation Rule
  - 1. If \$X \rightarrow Y\$ then \$XZ \rightarrow YZ\$ for every \$Z\$
  - 2. If X determines Y, then X combined with any other attribute Z will determine Y combined with Z
- 3. Transitivity Rule
  - 1. \$X \rightarrow Y\$ and \$Y \rightarrow Z\$ then \$X \rightarrow Z\$
  - 2. If X determines Y and Y Determines Z, then X will Determine Z
- 4. Union Rule
  - 1. if \$X \rightarrow Y\$ and \$X \rightarrow Z\$ then, \$X \rightarrow YZ\$
  - 2. Allows you to combine the results of separate FDs with the same Determinant
- 5. Decomposition Rule
  - 1. if \$X \rightarrow YZ\$, then \$X \rightarrow Y\$ and \$X \rightarrow Z\$
  - 2. If X Determines a composite set of Attributes YZ, it indivisually Determines Both or all subsets of the Composite Set
- 6. Psudo Transitivity Rule
  - 1. if \$X \rightarrow Y\$ and \$WY \rightarrow Z\$, then \$WX \rightarrow Z\$
  - 2. If X determines Y and WY determines Z, then combining X with W will determine Z
- 7. Identity Rule
  - 1. \$X \rightarrow X\$
  - 2. Any set of attributes are functionally dependent on itself