

Design and Analysis of Algorithms

L30: Knapsack problem Dynamic Programming

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Resources

- Text book 1: Levitin
 - Sec 8 . 4
- Text book 2: Horowitz
 - Sec 5 . 7
- R1: Introduction to Algorithms
 - Cormen et al.

Knapsack problem

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- Goal: solve the knapsack problem using dynamic programming.

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- Let $V[i, j]$ be the optimal solution to this instance
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- Approach: divide first i items into two categories:
 - Those that include i^{th} item, and
 - Those that don't include i^{th} item

DP Approach: Knapsack

	0		$j - w_i$		j		W
0	0	0	0	0	0	0	0
	0						
$i-1$	0						
i	0						
	0						
n	0						

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	0		$j - w_i$		j		W
0	0	0	0	0	0	0	0
	0						
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i	0						
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Table for solving knapsack problem using dynamic programming

DP Approach: Knapsack

	0		$j - w_i$		j		W
0	0	0	0	0	0	0	0
	0						
$i-1$	0						
i	0				$v[i, j]$		
	0						
n	0						

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	0		$j - w_i$		j		W
0	0	0	0	0	0	0	0
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	0						
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- Category 1: subsets that do not include i^{th} item.

DP Approach: Knapsack

	0		$j - w_i$		j		W
0	0	0	0	0	0	0	0
	0						
$i-1$	0						
i	0				$V[i, j]$		
	0						
n	0						

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	0		$j - w_i$		j		W
0	0	0	0	0	0	0	0
	0						
$i-1$	0				$V[i-1, j]$		
i	0				$V[i, j]$		
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	0		$j - w_i$		j		W
0	0	0	0	0	0	0	0
	0						
$i-1$	0		$V[i-1, j-w_i]$		$V[i-1, j]$		
i	0				$V[i, j]$		
	0						
n	0						

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- Category 1: subsets that do not include i^{th} item.
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 - Thus $j > w_i$ i.e. $j - w_i \geq 0$.

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	0		$j - w_i$		j		W
0	0	0	0	0	0	0	0
	0						
$i-1$	0		$V[i-1, j-w_i]$		$V[i-1, j]$		
i	0				$V[i, j]$		
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- Category 2: subsets that do include i^{th} item.
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 - Value of optimal subset is $v_i + V[i-1, j - w_i]$

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$$V[i, j] = \begin{cases} \max\{V[i-1, j], v_i + V[i-1, j - w_i]\} & \text{if } j - w_i \geq 0 \\ V[i-1, j] & \text{if } j - w_i < 0 \end{cases} \quad (1)$$

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 - **Values as**

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 $v_1=\$12, v_2=\$10, v_3=\$20, v_4=\15

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 - **Values as**
 $v_1=\$12, v_2=\$10, v_3=\$20, v_4=\15
 - **Need to compute $V[4, 5]$**
 - Max value with 4 items with knapsack capacity 5

Example Knapsack

Capacity→ wts, values↓		0	1	2	3	4	5
	0						
$w_1=2$ $v_1=12$	1						
$w_2=1$ $v_2=10$	2						
$w_3=3$ $v_3=20$	3						
$w_4=2$ $v_4=15$	4						

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$V[0, j] = 0$ **for** $0 \leq j \leq 5$

$V[i, 0] = 0$ **for** $0 \leq i \leq 4$

Example Knapsack

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$w_1=2 \quad v_1=12$	1	0					
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$w_3=3$ $v_3=20$	3	0					
$w_4=2$ $v_4=15$	4	0					

$V[1, 1] = V[1-1, 1]$ **since** $j=1 < w_1=2$

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2$ $v_1=12$	1	0					
$w_2=1$ $v_2=10$	2	0					
$w_3=3$ $v_3=20$	3	0					
$w_4=2$ $v_4=15$	4	0					

$V[1, 1] = V[1-1, 1]$ **since** $j=1 < w_1=2$
 $=0$

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2$ $v_1=12$	1	0	0				
$w_2=1$ $v_2=10$	2	0					
$w_3=3$ $v_3=20$	3	0					
$w_4=2$ $v_4=15$	4	0					

$V[1, 1] = V[1-1, 1]$ **since** $j=1 < w_1=2$
 $= 0$

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2$ $v_1=12$	1	0	0				
$w_2=1$ $v_2=10$	2	0					
$w_3=3$ $v_3=20$	3	0					
$w_4=2$ $v_4=15$	4	0					

$$V[1, 1] = V[1-1, 1] \text{ since } j=1 < w_1=2 \\ = 0$$

$$V[1, 2] = \max \{ V[0, 2], 12 + V[0, 2-2] \} ; j=2 \geq w_1=2$$

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2$ $v_1=12$	1	0	0				
$w_2=1$ $v_2=10$	2	0					
$w_3=3$ $v_3=20$	3	0					
$w_4=2$ $v_4=15$	4	0					

$$V[1, 1] = V[1-1, 1] \text{ since } j=1 < w_1=2 \\ = 0$$

$$V[1, 2] = \max\{V[0, 2], 12 + V[0, 2-2]\}; j=2 \geq w_1=2 \\ = \max\{0, 12 + V[0, 0]\} = 12$$

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2$ $v_1=12$	1	0	0	12			
$w_2=1$ $v_2=10$	2	0					
$w_3=3$ $v_3=20$	3	0					
$w_4=2$ $v_4=15$	4	0					

$$V[1, 1] = V[1-1, 1] \text{ since } j=1 < w_1=2 \\ = 0$$

$$V[1, 2] = \max\{V[0, 2], 12 + V[0, 2-2]\}; j=2 \geq w_1=2 \\ = \max\{0, 12 + V[0, 0]\} = 12$$

$$V[1, 3] = \max\{V[0, 3], 12 + V[0, 3-2]\}; j=3 \geq w_1=2$$

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2$ $v_1=12$	1	0	0	12			
$w_2=1$ $v_2=10$	2	0					
$w_3=3$ $v_3=20$	3	0					
$w_4=2$ $v_4=15$	4	0					

$$V[1, 1] = V[1-1, 1] \text{ since } j=1 < w_1=2 \\ = 0$$

$$V[1, 2] = \max\{V[0, 2], 12 + V[0, 2-2]\}; j=2 \geq w_1=2 \\ = \max\{0, 12 + V[0, 0]\} = 12$$

$$V[1, 3] = \max\{V[0, 3], 12 + V[0, 3-2]\}; j=3 \geq w_1=2 \\ = \max\{0, 12 + V[0, 1]\} = 12$$

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2$ $v_1=12$	1	0	0	12	12		
$w_2=1$ $v_2=10$	2	0					
$w_3=3$ $v_3=20$	3	0					
$w_4=2$ $v_4=15$	4	0					

$$V[1, 1] = V[1-1, 1] \text{ since } j=1 < w_1=2 \\ = 0$$

$$V[1, 2] = \max\{V[0, 2], 12 + V[0, 2-2]\}; j=2 \geq w_1=2 \\ = \max\{0, 12 + V[0, 0]\} = 12$$

$$V[1, 3] = \max\{V[0, 3], 12 + V[0, 3-2]\}; j=3 \geq w_1=2 \\ = \max\{0, 12 + V[0, 1]\} = 12$$

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2$ $v_1=12$	1	0	0	12	12		
$w_2=1$ $v_2=10$	2	0					
$w_3=3$ $v_3=20$	3	0					
$w_4=2$ $v_4=15$	4	0					

$$V[1, 1] = V[1-1, 1] \text{ since } j=1 < w_1=2 \\ = 0$$

$$V[1, 2] = \max\{V[0, 2], 12 + V[0, 2-2]\}; j=2 \geq w_1=2 \\ = \max\{0, 12 + V[0, 0]\} = 12$$

$$V[1, 3] = \max\{V[0, 3], 12 + V[0, 3-2]\}; j=3 \geq w_1=2 \\ = \max\{0, 12 + V[0, 1]\} = 12$$

$$V[1, 4] = \max\{V[0, 4], 12 + V[0, 4-2]\}; j=4 \geq w_1=2$$

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2$ $v_1=12$	1	0	0	12	12		
$w_2=1$ $v_2=10$	2	0					
$w_3=3$ $v_3=20$	3	0					
$w_4=2$ $v_4=15$	4	0					

$$V[1, 1] = V[1-1, 1] \text{ since } j=1 < w_1=2 \\ = 0$$

$$V[1, 2] = \max\{V[0, 2], 12 + V[0, 2-2]\}; j=2 \geq w_1=2 \\ = \max\{0, 12 + V[0, 0]\} = 12$$

$$V[1, 3] = \max\{V[0, 3], 12 + V[0, 3-2]\}; j=3 \geq w_1=2 \\ = \max\{0, 12 + V[0, 1]\} = 12$$

$$V[1, 4] = \max\{V[0, 4], 12 + V[0, 4-2]\}; j=4 \geq w_1=2 \\ = \max\{0, 12 + V[0, 2]\} = 12$$

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2$ $v_1=12$	1	0	0	12	12	12	
$w_2=1$ $v_2=10$	2	0					
$w_3=3$ $v_3=20$	3	0					
$w_4=2$ $v_4=15$	4	0					

$$V[1, 1] = V[1-1, 1] \text{ since } j=1 < w_1=2 \\ = 0$$

$$V[1, 2] = \max\{V[0, 2], 12 + V[0, 2-2]\}; j=2 \geq w_1=2 \\ = \max\{0, 12 + V[0, 0]\} = 12$$

$$V[1, 3] = \max\{V[0, 3], 12 + V[0, 3-2]\}; j=3 \geq w_1=2 \\ = \max\{0, 12 + V[0, 1]\} = 12$$

$$V[1, 4] = \max\{V[0, 4], 12 + V[0, 4-2]\}; j=4 \geq w_1=2 \\ = \max\{0, 12 + V[0, 2]\} = 12$$

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2$ $v_1=12$	1	0	0	12	12	12	
$w_2=1$ $v_2=10$	2	0					
$w_3=3$ $v_3=20$	3	0					
$w_4=2$ $v_4=15$	4	0					

$$V[1, 1] = V[1-1, 1] \text{ since } j=1 < w_1=2 \\ = 0$$

$$V[1, 2] = \max\{V[0, 2], 12 + V[0, 2-2]\}; j=2 \geq w_1=2 \\ = \max\{0, 12 + V[0, 0]\} = 12$$

$$V[1, 3] = \max\{V[0, 3], 12 + V[0, 3-2]\}; j=3 \geq w_1=2 \\ = \max\{0, 12 + V[0, 1]\} = 12$$

$$V[1, 4] = \max\{V[0, 4], 12 + V[0, 4-2]\}; j=4 \geq w_1=2 \\ = \max\{0, 12 + V[0, 2]\} = 12$$

$$V[1, 5] = \max\{V[0, 5], 12 + V[0, 5-2]\}; j=5 \geq w_1=2$$

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2$ $v_1=12$	1	0	0	12	12	12	
$w_2=1$ $v_2=10$	2	0					
$w_3=3$ $v_3=20$	3	0					
$w_4=2$ $v_4=15$	4	0					

$$V[1, 1] = V[1-1, 1] \text{ since } j=1 < w_1=2 \\ = 0$$

$$V[1, 2] = \max\{V[0, 2], 12 + V[0, 2-2]\}; j=2 \geq w_1=2 \\ = \max\{0, 12 + V[0, 0]\} = 12$$

$$V[1, 3] = \max\{V[0, 3], 12 + V[0, 3-2]\}; j=3 \geq w_1=2 \\ = \max\{0, 12 + V[0, 1]\} = 12$$

$$V[1, 4] = \max\{V[0, 4], 12 + V[0, 4-2]\}; j=4 \geq w_1=2 \\ = \max\{0, 12 + V[0, 2]\} = 12$$

$$V[1, 5] = \max\{V[0, 5], 12 + V[0, 5-2]\}; j=5 \geq w_1=2 \\ = \max\{0, 12 + V[0, 3]\} = 12$$

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2$ $v_1=12$	1	0	0	12	12	12	12
$w_2=1$ $v_2=10$	2	0					
$w_3=3$ $v_3=20$	3	0					
$w_4=2$ $v_4=15$	4	0					

$$V[1, 1] = V[1-1, 1] \text{ since } j=1 < w_1=2 \\ = 0$$

$$V[1, 2] = \max\{V[0, 2], 12 + V[0, 2-2]\}; j=2 \geq w_1=2 \\ = \max\{0, 12 + V[0, 0]\} = 12$$

$$V[1, 3] = \max\{V[0, 3], 12 + V[0, 3-2]\}; j=3 \geq w_1=2 \\ = \max\{0, 12 + V[0, 1]\} = 12$$

$$V[1, 4] = \max\{V[0, 4], 12 + V[0, 4-2]\}; j=4 \geq w_1=2 \\ = \max\{0, 12 + V[0, 2]\} = 12$$

$$V[1, 5] = \max\{V[0, 5], 12 + V[0, 5-2]\}; j=5 \geq w_1=2 \\ = \max\{0, 12 + V[0, 3]\} = 12$$

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2 \quad v_1=12$	1	0	0	12	12	12	12
$w_2=1 \quad v_2=10$	2	0					
$w_3=3 \quad v_3=20$	3	0					
$w_4=2 \quad v_4=15$	4	0					

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2$ $v_1=12$	1	0	0	12	12	12	12
$w_2=1$ $v_2=10$	2	0					
$w_3=3$ $v_3=20$	3	0					
$w_4=2$ $v_4=15$	4	0					

$$V[2, 1] = \max\{V[1, 1], 10 + V[1, 1-1]\}; \quad j=1 \geq w_2=1$$

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2 \quad v_1=12$	1	0	0	12	12	12	12
$w_2=1 \quad v_2=10$	2	0					
$w_3=3 \quad v_3=20$	3	0					
$w_4=2 \quad v_4=15$	4	0					

$$V[2,1] = \max\{V[1,1], 10 + V[1,1-1]\}; \quad j=1 \geq w_2=1$$

$$= \max\{0, 10 + V[1,0]\} = 10$$

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2$ $v_1=12$	1	0	0	12	12	12	12
$w_2=1$ $v_2=10$	2	0	10				
$w_3=3$ $v_3=20$	3	0					
$w_4=2$ $v_4=15$	4	0					

$$V[2,1] = \max\{V[1,1], 10 + V[1,1-1]\}; \quad j=1 \geq w_2=1$$

$$= \max\{0, 10 + V[1,0]\} = 10$$

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2$ $v_1=12$	1	0	0	12	12	12	12
$w_2=1$ $v_2=10$	2	0	10				
$w_3=3$ $v_3=20$	3	0					
$w_4=2$ $v_4=15$	4	0					

$$V[2,1] = \max\{V[1,1], 10 + V[1,1-1]\}; \quad j=1 \geq w_2=1$$

$$= \max\{0, 10 + V[1,0]\} = 10$$

$$V[2,2] = \max\{V[1,2], 10 + V[1,2-1]\}; \quad j=2 \geq w_2=1$$

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2 \quad v_1=12$	1	0	0	12	12	12	12
$w_2=1 \quad v_2=10$	2	0	10				
$w_3=3 \quad v_3=20$	3	0					
$w_4=2 \quad v_4=15$	4	0					

$$V[2,1] = \max\{V[1,1], 10 + V[1,1-1]\}; \quad j=1 \geq w_2=1$$

$$= \max\{0, 10 + V[1,0]\} = 10$$

$$V[2,2] = \max\{V[1,2], 10 + V[1,2-1]\}; \quad j=2 \geq w_2=1$$

$$= \max\{12, 10 + 0\} = 12$$

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2$ $v_1=12$	1	0	0	12	12	12	12
$w_2=1$ $v_2=10$	2	0	10	12			
$w_3=3$ $v_3=20$	3	0					
$w_4=2$ $v_4=15$	4	0					

$$V[2,1] = \max\{V[1,1], 10 + V[1,1-1]\}; \quad j=1 \geq w_2=1$$

$$= \max\{0, 10 + V[1,0]\} = 10$$

$$V[2,2] = \max\{V[1,2], 10 + V[1,2-1]\}; \quad j=2 \geq w_2=1$$

$$= \max\{12, 10 + 0\} = 12$$

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2 \quad v_1=12$	1	0	0	12	12	12	12
$w_2=1 \quad v_2=10$	2	0	10	12			
$w_3=3 \quad v_3=20$	3	0					
$w_4=2 \quad v_4=15$	4	0					

$$V[2,1] = \max\{V[1,1], 10 + V[1,1-1]\}; \quad j=1 \geq w_2=1$$

$$= \max\{0, 10 + V[1,0]\} = 10$$

$$V[2,2] = \max\{V[1,2], 10 + V[1,2-1]\}; \quad j=2 \geq w_2=1$$

$$= \max\{12, 10 + 0\} = 12$$

$$V[2,3] = \max\{V[1,3], 10 + V[1,3-1]\}; \quad j=3 \geq w_2=1$$

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2$ $v_1=12$	1	0	0	12	12	12	12
$w_2=1$ $v_2=10$	2	0	10	12			
$w_3=3$ $v_3=20$	3	0					
$w_4=2$ $v_4=15$	4	0					

$$V[2,1] = \max\{V[1,1], 10 + V[1,1-1]\}; \quad j=1 \geq w_2=1$$

$$= \max\{0, 10 + V[1,0]\} = 10$$

$$V[2,2] = \max\{V[1,2], 10 + V[1,2-1]\}; \quad j=2 \geq w_2=1$$

$$= \max\{12, 10 + 0\} = 12$$

$$V[2,3] = \max\{V[1,3], 10 + V[1,3-1]\}; \quad j=3 \geq w_2=1$$

$$= \max\{12, 10 + V[1,2]\} = \max\{12, 22\} = 22$$

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2$ $v_1=12$	1	0	0	12	12	12	12
$w_2=1$ $v_2=10$	2	0	10	12	22		
$w_3=3$ $v_3=20$	3	0					
$w_4=2$ $v_4=15$	4	0					

$$V[2,1] = \max\{V[1,1], 10 + V[1,1-1]\}; \quad j=1 \geq w_2=1$$

$$= \max\{0, 10 + V[1,0]\} = 10$$

$$V[2,2] = \max\{V[1,2], 10 + V[1,2-1]\}; \quad j=2 \geq w_2=1$$

$$= \max\{12, 10 + 0\} = 12$$

$$V[2,3] = \max\{V[1,3], 10 + V[1,3-1]\}; \quad j=3 \geq w_2=1$$

$$= \max\{12, 10 + V[1,2]\} = \max\{12, 22\} = 22$$

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2$ $v_1=12$	1	0	0	12	12	12	12
$w_2=1$ $v_2=10$	2	0	10	12	22		
$w_3=3$ $v_3=20$	3	0					
$w_4=2$ $v_4=15$	4	0					

$$V[2,1] = \max\{V[1,1], 10 + V[1,1-1]\}; \quad j=1 \geq w_2=1$$

$$= \max\{0, 10 + V[1,0]\} = 10$$

$$V[2,2] = \max\{V[1,2], 10 + V[1,2-1]\}; \quad j=2 \geq w_2=1$$

$$= \max\{12, 10 + 0\} = 12$$

$$V[2,3] = \max\{V[1,3], 10 + V[1,3-1]\}; \quad j=3 \geq w_2=1$$

$$= \max\{12, 10 + V[1,2]\} = \max\{12, 22\} = 22$$

$$V[2,4] = \max\{V[1,4], 10 + V[1,4-1]\}; \quad j=4 \geq w_2=1$$

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2 \quad v_1=12$	1	0	0	12	12	12	12
$w_2=1 \quad v_2=10$	2	0	10	12	22		
$w_3=3 \quad v_3=20$	3	0					
$w_4=2 \quad v_4=15$	4	0					

$$V[2,1] = \max\{V[1,1], 10 + V[1,1-1]\}; \quad j=1 \geq w_2=1 \\ = \max\{0, 10 + V[1,0]\} = 10$$

$$V[2,2] = \max\{V[1,2], 10 + V[1,2-1]\}; \quad j=2 \geq w_2=1 \\ = \max\{12, 10 + 0\} = 12$$

$$V[2,3] = \max\{V[1,3], 10 + V[1,3-1]\}; \quad j=3 \geq w_2=1 \\ = \max\{12, 10 + V[1,2]\} = \max\{12, 22\} = 22$$

$$V[2,4] = \max\{V[1,4], 10 + V[1,4-1]\}; \quad j=4 \geq w_2=1 \\ = \max\{12, 10 + V[1,3]\} = \max\{12, 22\} = 22$$

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2$ $v_1=12$	1	0	0	12	12	12	12
$w_2=1$ $v_2=10$	2	0	10	12	22	22	
$w_3=3$ $v_3=20$	3	0					
$w_4=2$ $v_4=15$	4	0					

$$V[2,1] = \max\{V[1,1], 10 + V[1,1-1]\}; \quad j=1 \geq w_2=1 \\ = \max\{0, 10 + V[1,0]\} = 10$$

$$V[2,2] = \max\{V[1,2], 10 + V[1,2-1]\}; \quad j=2 \geq w_2=1 \\ = \max\{12, 10 + 0\} = 12$$

$$V[2,3] = \max\{V[1,3], 10 + V[1,3-1]\}; \quad j=3 \geq w_2=1 \\ = \max\{12, 10 + V[1,2]\} = \max\{12, 22\} = 22$$

$$V[2,4] = \max\{V[1,4], 10 + V[1,4-1]\}; \quad j=4 \geq w_2=1 \\ = \max\{12, 10 + V[1,3]\} = \max\{12, 22\} = 22$$

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2$ $v_1=12$	1	0	0	12	12	12	12
$w_2=1$ $v_2=10$	2	0	10	12	22	22	
$w_3=3$ $v_3=20$	3	0					
$w_4=2$ $v_4=15$	4	0					

$$V[2,1] = \max\{V[1,1], 10+V[1,1-1]\}; \quad j=1 \geq w_2=1 \\ = \max\{0, 10+V[1,0]\} = 10$$

$$V[2,2] = \max\{V[1,2], 10+V[1,2-1]\}; \quad j=2 \geq w_2=1 \\ = \max\{12, 10+0\} = 12$$

$$V[2,3] = \max\{V[1,3], 10+V[1,3-1]\}; \quad j=3 \geq w_2=1 \\ = \max\{12, 10+V[1,2]\} = \max\{12, 22\} = 22$$

$$V[2,4] = \max\{V[1,4], 10+V[1,4-1]\}; \quad j=4 \geq w_2=1 \\ = \max\{12, 10+V[1,3]\} = \max\{12, 22\} = 22$$

$$V[2,5] = \max\{V[1,5], 10+V[1,5-1]\}; \quad j=5 \geq w_2=1$$

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2$ $v_1=12$	1	0	0	12	12	12	12
$w_2=1$ $v_2=10$	2	0	10	12	22	22	
$w_3=3$ $v_3=20$	3	0					
$w_4=2$ $v_4=15$	4	0					

$$V[2,1] = \max\{V[1,1], 10+V[1,1-1]\}; \quad j=1 \geq w_2=1 \\ = \max\{0, 10+V[1,0]\} = 10$$

$$V[2,2] = \max\{V[1,2], 10+V[1,2-1]\}; \quad j=2 \geq w_2=1 \\ = \max\{12, 10+0\} = 12$$

$$V[2,3] = \max\{V[1,3], 10+V[1,3-1]\}; \quad j=3 \geq w_2=1 \\ = \max\{12, 10+V[1,2]\} = \max\{12, 22\} = 22$$

$$V[2,4] = \max\{V[1,4], 10+V[1,4-1]\}; \quad j=4 \geq w_2=1 \\ = \max\{12, 10+V[1,3]\} = \max\{12, 22\} = 22$$

$$V[2,5] = \max\{V[1,5], 10+V[1,5-1]\}; \quad j=5 \geq w_2=1 \\ = \max\{12, 10+V[1,4]\} = \max\{12, 22\} = 22$$

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2 \quad v_1=12$	1	0	0	12	12	12	12
$w_2=1 \quad v_2=10$	2	0	10	12	22	22	22
$w_3=3 \quad v_3=20$	3	0					
$w_4=2 \quad v_4=15$	4	0					

$$V[2,1] = \max\{V[1,1], 10+V[1,1-1]\}; \quad j=1 \geq w_2=1 \\ = \max\{0, 10+V[1,0]\} = 10$$

$$V[2,2] = \max\{V[1,2], 10+V[1,2-1]\}; \quad j=2 \geq w_2=1 \\ = \max\{12, 10+0\} = 12$$

$$V[2,3] = \max\{V[1,3], 10+V[1,3-1]\}; \quad j=3 \geq w_2=1 \\ = \max\{12, 10+V[1,2]\} = \max\{12, 22\} = 22$$

$$V[2,4] = \max\{V[1,4], 10+V[1,4-1]\}; \quad j=4 \geq w_2=1 \\ = \max\{12, 10+V[1,3]\} = \max\{12, 22\} = 22$$

$$V[2,5] = \max\{V[1,5], 10+V[1,5-1]\}; \quad j=5 \geq w_2=1 \\ = \max\{12, 10+V[1,4]\} = \max\{12, 22\} = 22$$

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2 \quad v_1=12$	1	0	0	12	12	12	12
$w_2=1 \quad v_2=10$	2	0	10	12	22	22	22
$w_3=3 \quad v_3=20$	3	0					
$w_4=2 \quad v_4=15$	4	0					

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2$ $v_1=12$	1	0	0	12	12	12	12
$w_2=1$ $v_2=10$	2	0	10	12	22	22	22
$w_3=3$ $v_3=20$	3	0					
$w_4=2$ $v_4=15$	4	0					

$V[3,1]=V[2,1] = 10;$ ($j=1 < w_3=3$)

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2$ $v_1=12$	1	0	0	12	12	12	12
$w_2=1$ $v_2=10$	2	0	10	12	22	22	22
$w_3=3$ $v_3=20$	3	0	10				
$w_4=2$ $v_4=15$	4	0					

$V[3,1]=V[2,1] = 10;$ ($j=1 < w_3=3$)

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2 \quad v_1=12$	1	0	0	12	12	12	12
$w_2=1 \quad v_2=10$	2	0	10	12	22	22	22
$w_3=3 \quad v_3=20$	3	0	10				
$w_4=2 \quad v_4=15$	4	0					

$V[3, 1] = V[2, 1] = 10;$ ($j=1 < w_3=3$)

$V[3, 2] = V[2, 2] = 12;$ ($j=2 < w_3=3$)

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2 \quad v_1=12$	1	0	0	12	12	12	12
$w_2=1 \quad v_2=10$	2	0	10	12	22	22	22
$w_3=3 \quad v_3=20$	3	0	10	12			
$w_4=2 \quad v_4=15$	4	0					

$V[3, 1] = V[2, 1] = 10;$ ($j=1 < w_3=3$)

$V[3, 2] = V[2, 2] = 12;$ ($j=2 < w_3=3$)

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2$ $v_1=12$	1	0	0	12	12	12	12
$w_2=1$ $v_2=10$	2	0	10	12	22	22	22
$w_3=3$ $v_3=20$	3	0	10	12			
$w_4=2$ $v_4=15$	4	0					

$V[3,1]=V[2,1] = 10;$ ($j=1 < w_3=3$)

$V[3,2]=V[2,2] = 12;$ ($j=2 < w_3=3$)

$V[3,3]=\max\{V[2,3], 20+V[2,3-3]\};$ ($j=3 \geq w_3=3$)

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2 \quad v_1=12$	1	0	0	12	12	12	12
$w_2=1 \quad v_2=10$	2	0	10	12	22	22	22
$w_3=3 \quad v_3=20$	3	0	10	12			
$w_4=2 \quad v_4=15$	4	0					

$$V[3,1]=V[2,1] = 10; \quad (j=1 < w_3=3)$$

$$V[3,2]=V[2,2] = 12; \quad (j=2 < w_3=3)$$

$$V[3,3]=\max\{V[2,3], \quad 20+V[2,3-3]\}; \quad (j=3 \geq w_3=3)$$

$$=\max\{22, \quad 20+0\} = 22$$

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2$ $v_1=12$	1	0	0	12	12	12	12
$w_2=1$ $v_2=10$	2	0	10	12	22	22	22
$w_3=3$ $v_3=20$	3	0	10	12	22		
$w_4=2$ $v_4=15$	4	0					

$$V[3,1]=V[2,1] = 10; \quad (j=1 < w_3=3)$$

$$V[3,2]=V[2,2] = 12; \quad (j=2 < w_3=3)$$

$$V[3,3]=\max\{V[2,3], 20+V[2,3-3]\}; \quad (j=3 \geq w_3=3)$$

$$=\max\{22, 20+0\} = 22$$

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2 \quad v_1=12$	1	0	0	12	12	12	12
$w_2=1 \quad v_2=10$	2	0	10	12	22	22	22
$w_3=3 \quad v_3=20$	3	0	10	12	22		
$w_4=2 \quad v_4=15$	4	0					

$$V[3, 1] = V[2, 1] = 10; \quad (j=1 < w_3=3)$$

$$V[3, 2] = V[2, 2] = 12; \quad (j=2 < w_3=3)$$

$$V[3, 3] = \max\{V[2, 3], 20 + V[2, 3-3]\}; \quad (j=3 \geq w_3=3)$$

$$= \max\{22, 20+0\} = 22$$

$$V[3, 4] = \max\{V[2, 4], 20 + V[2, 4-3]\}; \quad (j=4 \geq w_3=3)$$

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2$ $v_1=12$	1	0	0	12	12	12	12
$w_2=1$ $v_2=10$	2	0	10	12	22	22	22
$w_3=3$ $v_3=20$	3	0	10	12	22		
$w_4=2$ $v_4=15$	4	0					

$$V[3,1]=V[2,1] = 10; \quad (j=1 < w_3=3)$$

$$V[3,2]=V[2,2] = 12; \quad (j=2 < w_3=3)$$

$$V[3,3]=\max\{V[2,3], 20+V[2,3-3]\}; \quad (j=3 \geq w_3=3)$$

$$=\max\{22, 20+0\} = 22$$

$$V[3,4]=\max\{V[2,4], 20+V[2,4-3]\}; \quad (j=4 \geq w_3=3)$$

$$=\max\{22, 20+V[2,1]\}=\max\{22, 30\}=30$$

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2$ $v_1=12$	1	0	0	12	12	12	12
$w_2=1$ $v_2=10$	2	0	10	12	22	22	22
$w_3=3$ $v_3=20$	3	0	10	12	22	30	
$w_4=2$ $v_4=15$	4	0					

$$V[3, 1] = V[2, 1] = 10; \quad (j=1 < w_3=3)$$

$$V[3, 2] = V[2, 2] = 12; \quad (j=2 < w_3=3)$$

$$V[3, 3] = \max\{V[2, 3], 20 + V[2, 3-3]\}; \quad (j=3 \geq w_3=3)$$

$$= \max\{22, 20+0\} = 22$$

$$V[3, 4] = \max\{V[2, 4], 20 + V[2, 4-3]\}; \quad (j=4 \geq w_3=3)$$

$$= \max\{22, 20 + V[2, 1]\} = \max\{22, 30\} = 30$$

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2 \quad v_1=12$	1	0	0	12	12	12	12
$w_2=1 \quad v_2=10$	2	0	10	12	22	22	22
$w_3=3 \quad v_3=20$	3	0	10	12	22	30	
$w_4=2 \quad v_4=15$	4	0					

$$V[3, 1] = V[2, 1] = 10; \quad (j=1 < w_3=3)$$

$$V[3, 2] = V[2, 2] = 12; \quad (j=2 < w_3=3)$$

$$V[3, 3] = \max\{V[2, 3], 20 + V[2, 3-3]\}; \quad (j=3 \geq w_3=3)$$

$$= \max\{22, 20+0\} = 22$$

$$V[3, 4] = \max\{V[2, 4], 20 + V[2, 4-3]\}; \quad (j=4 \geq w_3=3)$$

$$= \max\{22, 20 + V[2, 1]\} = \max\{22, 30\} = 30$$

$$V[3, 5] = \max\{V[2, 5], 20 + V[2, 5-3]\}; \quad (j=5 \geq w_3=3)$$

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2 \quad v_1=12$	1	0	0	12	12	12	12
$w_2=1 \quad v_2=10$	2	0	10	12	22	22	22
$w_3=3 \quad v_3=20$	3	0	10	12	22	30	
$w_4=2 \quad v_4=15$	4	0					

$$V[3, 1] = V[2, 1] = 10; \quad (j=1 < w_3=3)$$

$$V[3, 2] = V[2, 2] = 12; \quad (j=2 < w_3=3)$$

$$V[3, 3] = \max\{V[2, 3], 20 + V[2, 3-3]\}; \quad (j=3 \geq w_3=3)$$

$$= \max\{22, 20+0\} = 22$$

$$V[3, 4] = \max\{V[2, 4], 20 + V[2, 4-3]\}; \quad (j=4 \geq w_3=3)$$

$$= \max\{22, 20 + V[2, 1]\} = \max\{22, 30\} = 30$$

$$V[3, 5] = \max\{V[2, 5], 20 + V[2, 5-3]\}; \quad (j=5 \geq w_3=3)$$

$$= \max\{12, 20 + V[2, 2]\} = \max\{12, 20+12\} = 32$$

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2 \quad v_1=12$	1	0	0	12	12	12	12
$w_2=1 \quad v_2=10$	2	0	10	12	22	22	22
$w_3=3 \quad v_3=20$	3	0	10	12	22	30	32
$w_4=2 \quad v_4=15$	4	0					

$$V[3, 1] = V[2, 1] = 10; \quad (j=1 < w_3=3)$$

$$V[3, 2] = V[2, 2] = 12; \quad (j=2 < w_3=3)$$

$$V[3, 3] = \max\{V[2, 3], 20 + V[2, 3-3]\}; \quad (j=3 \geq w_3=3)$$

$$= \max\{22, 20+0\} = 22$$

$$V[3, 4] = \max\{V[2, 4], 20 + V[2, 4-3]\}; \quad (j=4 \geq w_3=3)$$

$$= \max\{22, 20 + V[2, 1]\} = \max\{22, 30\} = 30$$

$$V[3, 5] = \max\{V[2, 5], 20 + V[2, 5-3]\}; \quad (j=5 \geq w_3=3)$$

$$= \max\{12, 20 + V[2, 2]\} = \max\{12, 20+12\} = 32$$

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2$ $v_1=12$	1	0	0	12	12	12	12
$w_2=1$ $v_2=10$	2	0	10	12	22	22	22
$w_3=3$ $v_3=20$	3	0	10	12	22	30	32
$w_4=2$ $v_4=15$	4	0					

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2$ $v_1=12$	1	0	0	12	12	12	12
$w_2=1$ $v_2=10$	2	0	10	12	22	22	22
$w_3=3$ $v_3=20$	3	0	10	12	22	30	32
$w_4=2$ $v_4=15$	4	0					

$V[4,1]=V[3,1] = 10;$ ($j=1 < w_4=2$)

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2$ $v_1=12$	1	0	0	12	12	12	12
$w_2=1$ $v_2=10$	2	0	10	12	22	22	22
$w_3=3$ $v_3=20$	3	0	10	12	22	30	32
$w_4=2$ $v_4=15$	4	0	10				

$V[4,1]=V[3,1] = 10;$ ($j=1 < w_4=2$)

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2 \quad v_1=12$	1	0	0	12	12	12	12
$w_2=1 \quad v_2=10$	2	0	10	12	22	22	22
$w_3=3 \quad v_3=20$	3	0	10	12	22	30	32
$w_4=2 \quad v_4=15$	4	0	10				

$V[4, 1] = V[3, 1] = 10; \quad (j=1 < w_4=2)$

$V[4, 2] = \max\{V[3, 2], 15 + V[3, 2-2]\}; \quad (j=2 \geq w_4=2)$

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2 \quad v_1=12$	1	0	0	12	12	12	12
$w_2=1 \quad v_2=10$	2	0	10	12	22	22	22
$w_3=3 \quad v_3=20$	3	0	10	12	22	30	32
$w_4=2 \quad v_4=15$	4	0	10				

$V[4, 1] = V[3, 1] = 10; \quad (j=1 < w_4=2)$

$V[4, 2] = \max\{V[3, 2], 15 + V[3, 2-2]\}; \quad (j=2 \geq w_4=2)$
 $= \max\{12, 15 + V[3, 0]\} = \max\{12, 15 + 0\} = 15$

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2 \quad v_1=12$	1	0	0	12	12	12	12
$w_2=1 \quad v_2=10$	2	0	10	12	22	22	22
$w_3=3 \quad v_3=20$	3	0	10	12	22	30	32
$w_4=2 \quad v_4=15$	4	0	10	15			

$V[4,1]=V[3,1] = 10; \quad (j=1 < w_4=2)$

$V[4,2]=\max\{V[3,2], 15+V[3,2-2]\}; \quad (j=2 \geq w_4=2)$
 $=\max\{12, 15+V[3,0]\}=\max\{12, 15+0\}=15$

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2 \quad v_1=12$	1	0	0	12	12	12	12
$w_2=1 \quad v_2=10$	2	0	10	12	22	22	22
$w_3=3 \quad v_3=20$	3	0	10	12	22	30	32
$w_4=2 \quad v_4=15$	4	0	10	15			

$$V[4, 1] = V[3, 1] = 10; \quad (j=1 < w_4=2)$$

$$V[4, 2] = \max\{V[3, 2], 15 + V[3, 2-2]\}; \quad (j=2 \geq w_4=2)$$

$$= \max\{12, 15 + V[3, 0]\} = \max\{12, 15 + 0\} = 15$$

$$V[4, 3] = \max\{V[3, 3], 15 + V[3, 3-2]\}; \quad (j=3 \geq w_4=2)$$

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2 \quad v_1=12$	1	0	0	12	12	12	12
$w_2=1 \quad v_2=10$	2	0	10	12	22	22	22
$w_3=3 \quad v_3=20$	3	0	10	12	22	30	32
$w_4=2 \quad v_4=15$	4	0	10	15			

$$V[4, 1] = V[3, 1] = 10; \quad (j=1 < w_4=2)$$

$$V[4, 2] = \max\{V[3, 2], 15 + V[3, 2-2]\}; \quad (j=2 \geq w_4=2)$$

$$= \max\{12, 15 + V[3, 0]\} = \max\{12, 15 + 0\} = 15$$

$$V[4, 3] = \max\{V[3, 3], 15 + V[3, 3-2]\}; \quad (j=3 \geq w_4=2)$$

$$= \max\{22, 15 + V[3, 1]\} = \max\{22, 15 + 10\} = 25$$

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2 \quad v_1=12$	1	0	0	12	12	12	12
$w_2=1 \quad v_2=10$	2	0	10	12	22	22	22
$w_3=3 \quad v_3=20$	3	0	10	12	22	30	32
$w_4=2 \quad v_4=15$	4	0	10	15	25		

$$V[4, 1] = V[3, 1] = 10; \quad (j=1 < w_4=2)$$

$$V[4, 2] = \max\{V[3, 2], 15 + V[3, 2-2]\}; \quad (j=2 \geq w_4=2)$$

$$= \max\{12, 15 + V[3, 0]\} = \max\{12, 15 + 0\} = 15$$

$$V[4, 3] = \max\{V[3, 3], 15 + V[3, 3-2]\}; \quad (j=3 \geq w_4=2)$$

$$= \max\{22, 15 + V[3, 1]\} = \max\{22, 15 + 10\} = 25$$

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2$ $v_1=12$	1	0	0	12	12	12	12
$w_2=1$ $v_2=10$	2	0	10	12	22	22	22
$w_3=3$ $v_3=20$	3	0	10	12	22	30	32
$w_4=2$ $v_4=15$	4	0	10	15	25		

$$V[4, 1] = V[3, 1] = 10; \quad (j=1 < w_4=2)$$

$$V[4, 2] = \max\{V[3, 2], 15 + V[3, 2-2]\}; \quad (j=2 \geq w_4=2)$$

$$= \max\{12, 15 + V[3, 0]\} = \max\{12, 15 + 0\} = 15$$

$$V[4, 3] = \max\{V[3, 3], 15 + V[3, 3-2]\}; \quad (j=3 \geq w_4=2)$$

$$= \max\{22, 15 + V[3, 1]\} = \max\{22, 15 + 10\} = 25$$

$$V[4, 4] = \max\{V[3, 4], 15 + V[3, 4-2]\}; \quad (j=4 \geq w_4=2)$$

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2 \quad v_1=12$	1	0	0	12	12	12	12
$w_2=1 \quad v_2=10$	2	0	10	12	22	22	22
$w_3=3 \quad v_3=20$	3	0	10	12	22	30	32
$w_4=2 \quad v_4=15$	4	0	10	15	25		

$$V[4, 1] = V[3, 1] = 10; \quad (j=1 < w_4=2)$$

$$V[4, 2] = \max\{V[3, 2], 15 + V[3, 2-2]\}; \quad (j=2 \geq w_4=2)$$

$$= \max\{12, 15 + V[3, 0]\} = \max\{12, 15 + 0\} = 15$$

$$V[4, 3] = \max\{V[3, 3], 15 + V[3, 3-2]\}; \quad (j=3 \geq w_4=2)$$

$$= \max\{22, 15 + V[3, 1]\} = \max\{22, 15 + 10\} = 25$$

$$V[4, 4] = \max\{V[3, 4], 15 + V[3, 4-2]\}; \quad (j=4 \geq w_4=2)$$

$$= \max\{30, 15 + V[3, 2]\} = \max\{30, 15 + 12\} = 30$$

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2 \quad v_1=12$	1	0	0	12	12	12	12
$w_2=1 \quad v_2=10$	2	0	10	12	22	22	22
$w_3=3 \quad v_3=20$	3	0	10	12	22	30	32
$w_4=2 \quad v_4=15$	4	0	10	15	25	30	

$$V[4, 1] = V[3, 1] = 10; \quad (j=1 < w_4=2)$$

$$V[4, 2] = \max\{V[3, 2], 15 + V[3, 2-2]\}; \quad (j=2 \geq w_4=2)$$

$$= \max\{12, 15 + V[3, 0]\} = \max\{12, 15 + 0\} = 15$$

$$V[4, 3] = \max\{V[3, 3], 15 + V[3, 3-2]\}; \quad (j=3 \geq w_4=2)$$

$$= \max\{22, 15 + V[3, 1]\} = \max\{22, 15 + 10\} = 25$$

$$V[4, 4] = \max\{V[3, 4], 15 + V[3, 4-2]\}; \quad (j=4 \geq w_4=2)$$

$$= \max\{30, 15 + V[3, 2]\} = \max\{30, 15 + 12\} = 30$$

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2$ $v_1=12$	1	0	0	12	12	12	12
$w_2=1$ $v_2=10$	2	0	10	12	22	22	22
$w_3=3$ $v_3=20$	3	0	10	12	22	30	32
$w_4=2$ $v_4=15$	4	0	10	15	25	30	

$$V[4, 1] = V[3, 1] = 10; \quad (j=1 < w_4=2)$$

$$V[4, 2] = \max\{V[3, 2], 15 + V[3, 2-2]\}; \quad (j=2 \geq w_4=2)$$

$$= \max\{12, 15 + V[3, 0]\} = \max\{12, 15 + 0\} = 15$$

$$V[4, 3] = \max\{V[3, 3], 15 + V[3, 3-2]\}; \quad (j=3 \geq w_4=2)$$

$$= \max\{22, 15 + V[3, 1]\} = \max\{22, 15 + 10\} = 25$$

$$V[4, 4] = \max\{V[3, 4], 15 + V[3, 4-2]\}; \quad (j=4 \geq w_4=2)$$

$$= \max\{30, 15 + V[3, 2]\} = \max\{30, 15 + 12\} = 30$$

$$V[4, 5] = \max\{V[3, 5], 15 + V[3, 5-2]\}; \quad (j=5 \geq w_4=2)$$

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2$ $v_1=12$	1	0	0	12	12	12	12
$w_2=1$ $v_2=10$	2	0	10	12	22	22	22
$w_3=3$ $v_3=20$	3	0	10	12	22	30	32
$w_4=2$ $v_4=15$	4	0	10	15	25	30	

$$V[4, 1] = V[3, 1] = 10; \quad (j=1 < w_4=2)$$

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$$V[4, 3] = \max\{V[3, 3], 15 + V[3, 3-2]\}; \quad (j=3 \geq w_4=2)$$

$$= \max\{22, 15 + V[3, 1]\} = \max\{22, 15 + 10\} = 25$$

$$V[4, 4] = \max\{V[3, 4], 15 + V[3, 4-2]\}; \quad (j=4 \geq w_4=2)$$

$$= \max\{30, 15 + V[3, 2]\} = \max\{30, 15 + 12\} = 30$$

$$V[4, 5] = \max\{V[3, 5], 15 + V[3, 5-2]\}; \quad (j=5 \geq w_4=2)$$

$$= \max\{32, 15 + V[3, 3]\} = \max\{32, 15 + 22\} = 37$$

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2$ $v_1=12$	1	0	0	12	12	12	12
$w_2=1$ $v_2=10$	2	0	10	12	22	22	22
$w_3=3$ $v_3=20$	3	0	10	12	22	30	32
$w_4=2$ $v_4=15$	4	0	10	15	25	30	37

$$V[4, 1] = V[3, 1] = 10; \quad (j=1 < w_4=2)$$

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$$= \max\{12, 15 + V[3, 0]\} = \max\{12, 15 + 0\} = 15$$

$$V[4, 3] = \max\{V[3, 3], 15 + V[3, 3-2]\}; \quad (j=3 \geq w_4=2)$$

$$= \max\{22, 15 + V[3, 1]\} = \max\{22, 15 + 10\} = 25$$

$$V[4, 4] = \max\{V[3, 4], 15 + V[3, 4-2]\}; \quad (j=4 \geq w_4=2)$$

$$= \max\{30, 15 + V[3, 2]\} = \max\{30, 15 + 12\} = 30$$

$$V[4, 5] = \max\{V[3, 5], 15 + V[3, 5-2]\}; \quad (j=5 \geq w_4=2)$$

$$= \max\{32, 15 + V[3, 3]\} = \max\{32, 15 + 22\} = 37$$

Example Knapsack: Optimal Subset

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2 \quad v_1=12$	1	0	0	12	12	12	12
$w_2=1 \quad v_2=10$	2	0	10	12	22	22	22
$w_3=3 \quad v_3=20$	3	0	10	12	22	30	32
$w_4=2 \quad v_4=15$	4	0	10	15	25	30	37

Example Knapsack: Optimal Subset

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2 \quad v_1=12$	1	0	0	12	12	12	12
$w_2=1 \quad v_2=10$	2	0	10	12	22	22	22
$w_3=3 \quad v_3=20$	3	0	10	12	22	30	32
$w_4=2 \quad v_4=15$	4	0	10	15	25	30	37

- Optimal subset

Example Knapsack: Optimal Subset

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	0	0	0	0	0	0	0
$w_1=2 \quad v_1=12$	1	0	0	12	12	12	12
$w_2=1 \quad v_2=10$	2	0	10	12	22	22	22
$w_3=3 \quad v_3=20$	3	0	10	12	22	30	32
$w_4=2 \quad v_4=15$	4	0	10	15	25	30	37

- Optimal subset
 - Backtrack from maximal value $V[4, 5]$ to prev. rows.

Example Knapsack: Optimal Subset

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2$ $v_1=12$	1	0	0	12	12	12	12
$w_2=1$ $v_2=10$	2	0	10	12	22	22	22
$w_3=3$ $v_3=20$	3	0	10	12	22	30	32
$w_4=2$ $v_4=15$	4	0	10	15	25	30	37

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 - Backtrack from maximal value $V[4, 5]$ to prev. rows.
 - Thus, optimal subsets are

Example Knapsack: Optimal Subset

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	0	0	0	0	0	0	0
$w_1=2$ $v_1=12$	1	0	0	12	12	12	12
$w_2=1$ $v_2=10$	2	0	10	12	22	22	22
$w_3=3$ $v_3=20$	3	0	10	12	22	30	32
$w_4=2$ $v_4=15$	4	0	10	15	25	30	37

- Optimal subset
 - Backtrack from maximal value $V[4, 5]$ to prev. rows.
 - Thus, optimal subsets are
 $V[4, 5] = 37$ ($\neq V[3, 5]$) implies $w_4=2$ is included

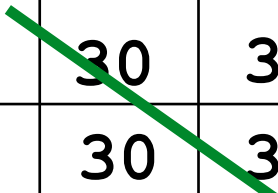
Example Knapsack: Optimal Subset

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2 \quad v_1=12$	1	0	0	12	12	12	12
$w_2=1 \quad v_2=10$	2	0	10	12	22	22	22
$w_3=3 \quad v_3=20$	3	0	10	12	22	30	32
$w_4=2$ $v_4=15$	4	0	10	15	25	30	37

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Example Knapsack: Optimal Subset

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2 \quad v_1=12$	1	0	0	12	12	12	12
$w_2=1 \quad v_2=10$	2	0	10	12	22	22	22
$w_3=3 \quad v_3=20$	3	0	10	12	22	30	32
$w_4=2$ $v_4=15$	4	0	10	15	25	30	37



- Optimal subset
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Example Knapsack: Optimal Subset

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2 \quad v_1=12$	1	0	0	12	12	12	12
$w_2=1 \quad v_2=10$	2	0	10	12	22	22	22
$w_3=3 \quad v_3=20$	3	0	10	12	22	30	32
$w_4=2$ $v_4=15$	4	0	10	15	25	30	37

- Optimal subset
 - Backtrack from maximal value $V[4, 5]$ to prev. rows.
 - Thus, optimal subsets are
 - $V[4, 5] = 37$ ($\neq V[3, 5]$) implies $w_4=2$ is included
 - $V[3, 3] = 22$ ($= V[2, 3]$) implies $w_3=3$ is not included

Example Knapsack: Optimal Subset

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2 \quad v_1=12$	1	0	0	12	12	12	12
$w_2=1 \quad v_2=10$	2	0	10	12	22	22	22
$w_3=3 \quad v_3=20$	3	0	10	12	22	30	32
$w_4=2 \quad v_4=15$	4	0	10	15	25	30	37

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 - $V[4, 5] = 37$ ($\neq V[3, 5]$) implies $w_4=2$ is included
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Example Knapsack: Optimal Subset

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2 \quad v_1=12$	1	0	0	12	12	12	12
$w_2=1 \quad v_2=10$	2	0	10	12	22	22	22
$w_3=3 \quad v_3=20$	3	0	10	12	22	30	32
$w_4=2 \quad v_4=15$	4	0	10	15	25	30	37

- Optimal subset
 - Backtrack from maximal value $V[4, 5]$ to prev. rows.
 - Thus, optimal subsets are
 - $V[4, 5] = 37$ ($\neq V[3, 5]$) implies $w_4=2$ is included
 - $V[3, 3] = 22$ ($= V[2, 3]$) implies $w_3=3$ is not included

Example Knapsack: Optimal Subset

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2 \quad v_1=12$	1	0	0	12	12	12	12
$w_2=1 \quad v_2=10$	2	0	10	12	22	22	22
$w_3=3 \quad v_3=20$	3	0	10	12	22	30	32
$w_4=2 \quad v_4=15$	4	0	10	15	25	30	37

- Optimal subset
 - Backtrack from maximal value $V[4, 5]$ to prev. rows.
 - Thus, optimal subsets are
 - $V[4, 5] = 37 (\neq V[3, 5])$ implies $w_4=2$ is included
 - $V[3, 3] = 22 (=V[2, 3])$ implies $w_3=3$ is not included
 - $V[2, 3] = 22 (\neq V[1, 3])$ implies $w_2=1$ is included

Example Knapsack: Optimal Subset

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2$ $v_1=12$	1	0	0	12	12	12	12
$w_2=1$ $v_2=10$	2	0	10	12	22	22	22
$w_3=3$ $v_3=20$	3	0	10	12	22	30	32
$w_4=2$ $v_4=15$	4	0	10	15	25	30	37

- Optimal subset
 - Backtrack from maximal value $V[4, 5]$ to prev. rows.
 - Thus, optimal subsets are
 - $V[4, 5] = 37$ ($\neq V[3, 5]$) implies $w_4=2$ is included
 - $V[3, 3] = 22$ ($= V[2, 3]$) implies $w_3=3$ is not included
 - $V[2, 3] = 22$ ($\neq V[1, 3]$) implies $w_2=1$ is included

Example Knapsack: Optimal Subset

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2$ $v_1=12$	1	0	0	12	12	12	12
$w_2=1$ $v_2=10$	2	0	10	12	22	22	22
$w_3=3$ $v_3=20$	3	0	10	12	22	30	32
$w_4=2$ $v_4=15$	4	0	10	15	25	30	37

- Optimal subset
 - Backtrack from maximal value $V[4, 5]$ to prev. rows.
 - Thus, optimal subsets are
 - $V[4, 5] = 37$ ($\neq V[3, 5]$) implies $w_4=2$ is included
 - $V[3, 3] = 22$ ($= V[2, 3]$) implies $w_3=3$ is not included
 - $V[2, 3] = 22$ ($\neq V[1, 3]$) implies $w_2=1$ is included

Example Knapsack: Optimal Subset

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
$w_1=2$ $v_1=12$	1	0	0	12	12	12	12
$w_2=1$ $v_2=10$	2	0	10	12	22	22	22
$w_3=3$ $v_3=20$	3	0	10	12	22	30	32
$w_4=2$ $v_4=15$	4	0	10	15	25	30	37

- Optimal subset
 - Backtrack from maximal value $V[4, 5]$ to prev. rows.
 - Thus, optimal subsets are
 - $V[4, 5] = 37$ ($\neq V[3, 5]$) implies $w_4=2$ is included
 - $V[3, 3] = 22$ ($= V[2, 3]$) implies $w_3=3$ is not included
 - $V[2, 3] = 22$ ($\neq V[1, 3]$) implies $w_2=1$ is included
 - $V[1, 2] = 12$ ($\neq V[0, 2]$) implies $w_1=2$ is included

Example Knapsack: Optimal Subset

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
w₁=2 v ₁ =12	1	0	0	12	12	12	12
w₂=1 v ₂ =10	2	0	10	12	22	22	22
w₃=3 v ₃ =20	3	0	10	12	22	30	32
w₄=2 v ₄ =15	4	0	10	15	25	30	37

- Optimal subset
 - Backtrack from maximal value $V[4, 5]$ to prev. rows.
 - Thus, optimal subsets are
 - $V[4, 5] = 37$ ($\neq V[3, 5]$) implies $w_4 = 2$ is included
 - $V[3, 3] = 22$ ($= V[2, 3]$) implies $w_3 = 3$ is not included
 - $V[2, 3] = 22$ ($\neq V[1, 3]$) implies $w_2 = 1$ is included
 - $V[1, 2] = 12$ ($\neq V[0, 2]$) implies $w_1 = 2$ is included

Example Knapsack: Optimal Subset

Capacity→ wts, values↓		0	1	2	3	4	5
	0	0	0	0	0	0	0
w₁=2 v ₁ =12	1	0	0	12	12	12	12
w₂=1 v ₂ =10	2	0	10	12	22	22	22
w₃=3 v ₃ =20	3	0	10	12	22	30	32
w₄=2 v ₄ =15	4	0	10	15	25	30	37

- Optimal subset
 - Backtrack from maximal value $V[4, 5]$ to prev. rows.
 - Thus, optimal subsets are
 - $V[4, 5] = 37$ ($\neq V[3, 5]$) implies $w_4 = 2$ is included
 - $V[3, 3] = 22$ ($= V[2, 3]$) implies $w_3 = 3$ is not included
 - $V[2, 3] = 22$ ($\neq V[1, 3]$) implies $w_2 = 1$ is included
 - $V[1, 2] = 12$ ($\neq V[0, 2]$) implies $w_1 = 2$ is included

Algorithm: Knapsack using DP

```
Algo DPKnapsack( $w[1..n]$ ,  $v[1..n]$ ,  $W$ )
    int  $V[0..n, 0..W]$ ,  $P[1..n, 1..W]$ ;
    for  $j=0$  to  $W$  do
         $V[0, j] = 0$ 
    for  $i=0$  to  $n$  do
         $V[i, 0] = 0$ 
    for  $i=1$  to  $n$  do
        for  $j=1$  to  $W$  do
            if  $w[i] \leq j$  and  $(v[i] + V[i-1, j-w[i]]) > V[i-1, j]$  then
                 $V[i, j] = v[i] + V[i-1, j-w[i]]$ ;
            else
                 $V[i, j] = V[i-1, j]$ 
    return  $V[n, W]$  (and the optimal subset by backtracing)
```

Efficiency of Knapsack

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- Time Efficiency: $\Theta(nW)$

Efficiency of Knapsack

- Time Efficiency: $\Theta(nW)$
- Space efficiency: $\Theta(nW)$

Summary

Summary

- Knapsack algorithm using dynamic programming

Summary

- Knapsack algorithm using dynamic programming
- Efficiency

Summary

- Knapsack algorithm using dynamic programming
- Efficiency
- Optimal subsets using backtracking