# **Pragmatics of Efficient Proof Generation**

### Stephan Schulz









### Introduction

$$\{\textit{A}_{1},\textit{A}_{2},\ldots,\textit{A}_{n}\} \models \textit{C}$$











$$\{A_1, A_2, \dots, A_n\} \models C$$

 $\{A_1, A_2, \dots, A_n, \neg C\}$  is unsatisfiable











$$\{A_1, A_2, \dots, A_n\} \models C$$

 $\{A_1,A_2,\ldots,A_n,\neg C\}$  is unsatisfiable iff

 $\operatorname{cnf}(\{A_1,A_2,A_n,\neg C\})$  is unsatisfiable











$$\{A_1, A_2, \dots, A_n\} \models C$$

$$\{A_1, A_2, \dots, A_n, \neg C\}$$
 is unsatisfiable iff

$$\operatorname{cnf}(\{A_1,A_2,A_n,\neg C\})$$
 is unsatisfiable iff

$$\operatorname{cnf}(\{A_1,A_2,A_n,\neg C\}) \stackrel{*}{\vdash} \Box$$











$$\{A_1,A_2,\ldots,A_n\} \models C$$
 iff 
$$\{A_1,A_2,\ldots,A_n,\neg C\} \text{ is unsatisfiable}$$
 iff 
$$\operatorname{cnf}(\{A_1,A_2,A_n,\neg C\}) \text{ is unsatisfiable}$$
 iff

 $\operatorname{cnf}(\{A_1, A_2, A_n, \neg C\}) \stackrel{*}{\vdash} \square$ 

Clausification

Refutation/ Saturation

## Ideal: Proofs as Sequences of Proof Steps

- A derivation is a list of steps
- Each step carries a clause/formula
- ► Each step is either...
  - Assumed (e.g. axioms, conjecture)
  - Logically derived from earlier steps
- ► A proof is a derivation that either...
  - derives the conjecture
  - derives a contradiction from the negated conjecture

#### Good mental model!

## Reality: Proofs as Sequences of Proof Steps

- Initial clauses/formulas
  - Axioms/Conjectures/Hypotheses
  - Justified by assumption
- ▶ Derived clauses/formulas
  - Justified by reference to (topologically) preceding steps
  - Defined logical relationship to predecessors
    - Most frequent case: theorem of predecessors
    - Exceptions: Skolemization, negation of conjecture, . . .
- ► (Introduced definitions)
  - Don't affect satisfiability/provability
  - Justified by definition

## Example

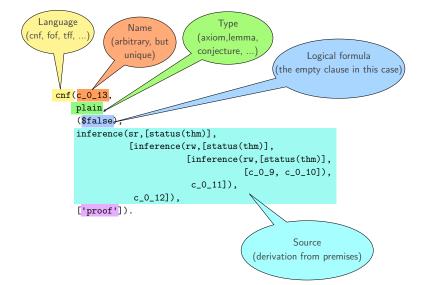
```
fof(c 0 0, conjecture, (?[X3]:(human(X3)&X3!=john)), file('humen.p', someone not john)).
fof(c 0 1, axiom, (?[X3]:(human(X3)&grade(X3)=a)), file('humen.p', someone got an a)).
fof(c 0 2, axiom, (grade(john)=f), file('humen.p', john failed)).
fof(c_0_3, axiom, (a!=f), file('humen.p', distinct_grades)).
fof(c 0 4, negated conjecture, (~(?[X3]:(human(X3)&X3!=john))),
    inference (assume negation, [status(cth)], [c 0 0])).
fof(c 0 5, negated conjecture, (![X4]:(~human(X4)|X4=john)),
    inference(variable rename,[status(thm)],[inference(fof nnf,[status(thm)],[c 0 4])])).
fof (c 0 6, plain, ((human(esk1 0)&grade(esk1 0)=a)),
    inference (skolemize, [status (esa)], [inference (variable rename, [status (thm)], [c 0 1])])).
cnf(c 0 7, negated conjecture, (X1=john|~human(X1)),
    inference(split conjunct,[status(thm)],[c 0 5])).
cnf(c 0 8.plain, (human(esk1 0)),
    inference(split_conjunct,[status(thm)],[c_0_6])).
cnf(c \ 0 \ 9, plain, (grade(esk1 \ 0)=a),
    inference(split conjunct,[status(thm)],[c 0 6])).
cnf(c_0_10, negated_conjecture, (esk1_0=john),
    inference(spm,[status(thm)],[c 0 7, c 0 8])).
cnf(c 0 11, plain, (grade(john)=f),
    inference(split conjunct,[status(thm)],[c 0 2])).
cnf(c 0 12, plain, (a!=f),
    inference(split conjunct,[status(thm)],[c 0 3])).
cnf(c 0 13.plain, ($false),
    inference (sr. [status(thm)], [inference(rw, [status(thm)],
    [inference(rw,[status(thm)],[c 0 9, c 0 10]), c 0 11]), c 0 12]), ['proof']).
```

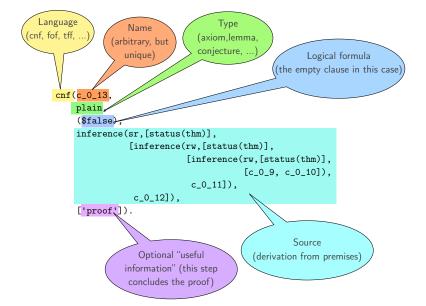
```
Language
(cnf, fof, tff, ...
     cnf(c_0_13,
         plain,
          ($false),
         inference(sr,[status(thm)],
                     [inference(rw,[status(thm)],
                                 [inference(rw,[status(thm)],
                                            [c_0_9, c_0_10]),
                                 c_0_11]),
                     c_0_12]),
          ['proof']).
```

```
Language
                     Name
(cnf, fof, tff, ...)
                 (arbitrary, but
                    unique)
     cnf(c_0_13,
          plain,
          ($false),
          inference(sr,[status(thm)],
                     [inference(rw,[status(thm)],
                                 [inference(rw,[status(thm)],
                                             [c_0_9, c_0_10]),
                                  c_0_11]),
                      c_0_12]),
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```

```
Language
                                        Type
                     Name
(cnf, fof, tff, ...)
                                    (axiom,lemma,
                 (arbitrary, but
                                    conjecture, ...
                    unique)
     cnf(c_0_13,
          plain
          ($false),
          inference(sr,[status(thm)],
                     [inference(rw,[status(thm)],
                                  [inference(rw,[status(thm)],
                                              [c_0_9, c_0_10]),
                                   c_0_11]),
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          ['proof']).
```

```
Language
                                         Type
                     Name
(cnf, fof, tff, ...
                                     (axiom,lemma,
                  (arbitrary, but
                                     conjecture, ..
                     unique)
                                                            Logical formula
                                                     (the empty clause in this case)
     cnf(c_0_13,
          plain
          ($false)
          inference(sr,[status(thm)],
                      [inference(rw,[status(thm)],
                                  [inference(rw,[status(thm)],
                                               [c_0_9, c_0_10]),
                                   c_0_11]),
                      c_0_12]),
          ['proof']).
```





```
Inference rule (sr:
      Simplify-reflect, rw:
        Rewriting, pm:
      Paramodulation, ...
cnf(c_0_13,
    plain,
    ($false),
    inference(sr,[status(thm)],
               [inference(rw,[status(thm)],
                           [inference(rw,[status(thm)],
                                       [c_0_9, c_0_10]),
                            c_0_11]),
                c_0_12]),
    ['proof']).
```

```
Inference rule (sr:
                                          "Useful
       Simplify-reflect, rw:
                                    information": logical
         Rewriting, pm:
                                     status (formula is
      Paramodulation, ...
                                   theorem of premises)
cnf(c_0_13,
    plain,
    ($false),
    inference(sr,[status(thm)],
                [inference(rw,[status(thm)],
                            [inference(rw,[status(thm)],
                                         [c_0_9, c_0_10]),
                             c_0_11]),
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```

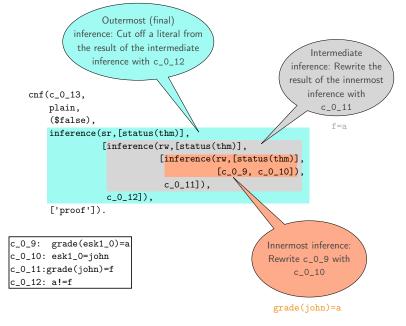
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                                         [c_0_9, c_0_10]),
                             c_0_11]),
                c_0_12]),
    ['proof']).
                                                   Names of the premises
```

```
c_0_9: grade(esk1_0)=a
c_0_10: esk1_0=john
c_0_11:grade(john)=f
c_0_12: a!=f
```

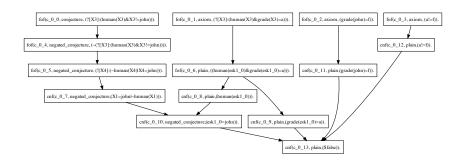
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        ($false),
       inference(sr,[status(thm)],
                  [inference(rw,[status(thm)],
                              [inference(rw,[status(thm)],
                                         [c_0_9, c_0_10]),
                               c_0_11]),
                   c_0_{12}),
        ['proof']).
c_0_9: grade(esk1_0)=a
                                                   Innermost inference:
c_0_10: esk1_0=john
                                                    Rewrite c_0_9 with
c_0_11:grade(john)=f
                                                         c_0_10
c_0_12: a!=f
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cnf(c_0_13,
       plain,
        ($false),
       inference(sr,[status(thm)],
                  [inference(rw,[status(thm)],
                              [inference(rw,[status(thm)],
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c_0_9: grade(esk1_0)=a
                                                   Innermost inference:
c_0_10: esk1_0=john
                                                    Rewrite c_0_9 with
c_0_11:grade(john)=f
                                                         c_0_10
c_0_12: a!=f
                                                     grade(john)=a
```

```
Intermediate
                                                           inference: Rewrite the
                                                          result of the innermost
                                                              inference with
   cnf(c_0_13,
        plain,
                                                                 c_0_11
        ($false),
                                                                   f=a
        inference(sr,[status(thm)],
                   [inference(rw,[status(thm)],
                               [inference(rw,[status(thm)],
                                           [c_0_9, c_0_10]),
                                c_0_11]),
                    c_0_{12}),
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c_0_9: grade(esk1_0)=a
                                                     Innermost inference:
c_0_10: esk1_0=john
                                                      Rewrite c_0_9 with
c_0_11:grade(john)=f
                                                           c_0_10
c_0_12: a!=f
                                                       grade(john)=a
```



## Proofs as Graphs



## **Proof Generation and Representation**

$$\{A_1,A_2,\ldots,A_n\} \models C$$
iff
 $\{A_1,A_2,\ldots,A_n,\neg C\}$  is unsatisfiable
iff
 $\operatorname{cnf}(\{A_1,A_2,A_n,\neg C\})$  is unsatisfiable
iff
 $\operatorname{cnf}(\{A_1,A_2,A_n,\neg C\}) \vdash \Box$ 

Clausification

Refutation

## Clausification and Saturation

- Clausification
  - Terminating
  - (Usually) deterministic
  - (Usually) non-destructive
  - Sometimes done by external tool
- ▶ Saturation
  - Many degrees of freedom
  - Arbitrary search time
  - Generating inferences
    - ► Create new clauses
    - Necessary for completeness
  - Simplifying inferences
    - ► Modify/remove existing clauses
    - Necessary for performance

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- ► Recording clausification is straightforward
  - ...but not always done
- Efficiently recording saturation is difficult
  - ...some settle for inefficient

## Deduction vs. Simplification

Superposition 
$$\frac{s \simeq t \vee S \quad u \not\simeq v \vee R}{\sigma(u[p \leftarrow t] \not\simeq v \vee S \vee R)}$$
 if  $\sigma = mgu(u|_p, s), [\dots]$ 

► Rewriting 
$$\frac{s \simeq t \quad u \not\simeq v \lor R}{s \simeq t \quad u[p \leftarrow \sigma(t)] \not\simeq v \lor R}$$
if  $u|_p = \sigma(s)$  and  $\sigma(s) > \sigma(t)$ 

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► Generating inferences create new clauses

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► Rewriting 
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 if  $u|_p = \sigma(s)$  and  $\sigma(s) > \sigma(t)$ 

- Generating inferences create new clauses
- ➤ Simplifying inferences replace or remove clauses

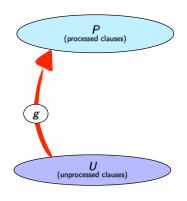
# Clauses vs. Clause Objects

Clauses	Clause Objects		
Logical abstraction	Real data structures		
Purely logical object	Can carry meta-information and history		
Immutable (changing a clause creates a completely new clause)	Mutable, changes are frequent		
Used in theoretical inference systems	Implemented in actual provers		

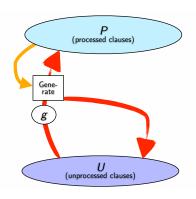
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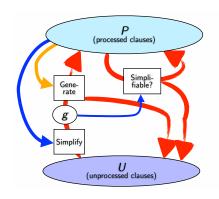
Clause objects represent a sequence of clauses over time (and possibly their history)



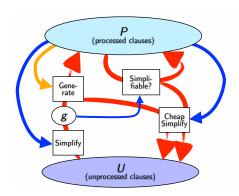
► Aim: Move everything from *U* to *P* 



- Aim: Move everything from U to P
- Invariant: All generating inferences with premises from P have been performed



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- ► Invariant: *P* is interreduced



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- Invariant: All generating inferences with premises from P have been performed
- Invariant: P is interreduced
- Clauses added to U are simplified with respect to P

## Typical Clause Object Lifecycle

- Generating inference creates a new clause object
  - Usually paramodulation (but may be equality factoring, equality resolution, . . .)
- Simplifying inferences modifies the clause object
  - Multiple rewrite steps
  - Possibly literal cutting, trivial literal removal, . . .
  - ➤ This modifies the existing clause object...
  - ≥ ≈10 modifications per clause on average (varies wildly)
- Deleting inference removes clause, but not (always) clause object
  - Subsumption
  - Tautology deletion
  - ► Typically ≈90% of all clauses

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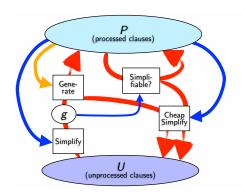
90% of clauses eventually deleted, 9 modified versions ⇒ 99% of (logical) clauses are not persistent

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- ► Generating inference *creates* a new clause object
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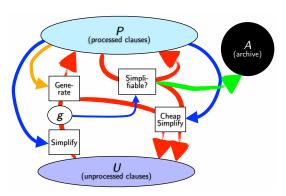
Storing all clauses is too expensive, but we don't know a-priori which clauses are needed!

# **Optimized Proof Object Construction**



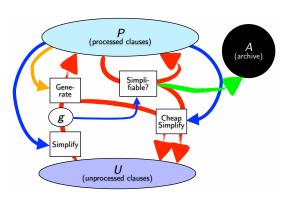
► Observation: Only clauses in *P* are premises!

# **Optimized Proof Object Construction**



- Observation: Only clauses in P are premises!
- Proof recording:
  - Simplified P-clauses are archived
  - Clauses record their history
    - Inference rules
    - ► P-clauses involved

# **Optimized Proof Object Construction**



- Observation: Only clauses in P are premises!
- Proof recording:
  - Simplified P-clauses are archived
  - Clauses record their history
    - ► Inference rules
    - P-clauses involved
- Proof extraction
  - Track parent relation
  - Topological sort
  - Print proof

# Real Clauses (excerpt)

```
typedef struct clause_cell
                      ident:
   long
                                 /* Hopefully unique ident for
                                     all clauses created during
                                    proof run */
   Eqn_p
                      literals: /* List of literals */
   FormulaProperties properties; /* Anything we want to note at
                                    the clause? */
   ClauseInfo_p
                      info; /* Currently about source in
                                    input. NULL for derived
                                    clauses */
   PStack<sub>-</sub>p
                      derivation; /* Derivation of the clause for
                                    proof reconstruction. */
}ClauseCell, *Clause_p;
```

```
cnf(c_0_{13},
    plain,
    ($false),
    inference(sr,[status(thm)],
               [inference(rw,[status(thm)],
                           [inference(rw,[status(thm)],
                                      [c 0 9, c 0 10]).
                           c_0_11]),
                c_0_{12},
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```

13	NULL		NULL	
ident	literals	properties	info	derivation

```
cnf(c_0_13,
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   inference(sr,[status(thm)],
           [inference(rw,[status(thm)],
                    [inference(rw,[status(thm)],
                             [c_0_9, c_0_10]),
                     c_0_11]),
            c_0_12]),
   ['proof']).
                                                                    [c_0_12]
                                                                       DCSR
                                                                    [c_0_11]
                                                                   DCRewrite
                                                                    [c_0_10]
                                                                   DCRewrite
                                                                     [c_0_9]
                                                                   DCCnfQuote
         13
                       NULL
                                                       NULL
                    literals
                                  properties
                                                                  derivation
       ident
                                                       info
```

```
cnf(c_0_13,
   plain,
   ($false).
   inference(sr.[status(thm)].
           [inference(rw,[status(thm)],
                    [inference(rw,[status(thm)],
                             [c 0 9, c 0 10]).
                     c_0_11]),
            c_0_12]),
   ['proof']).
                                     ...clause 12
                                                                      [c_0_{12}]
                      ...and looses a literal by cutting with...
                                                                        DCSR
                                    ...clause 11...
                                                                      [c_0_11]
                          ...and further rewritten with...
                                                                     DCRewrite
                             ...with clause number 10
                                                                      [c_0_10]
                                                                     DCRewrite
                          It was modified by rewriting...
                               ... of clause number 9
                                                                      [c_0_9]
                                                                    DCCnfQuote
                             Original clause is copy...
         13
                       NUI.I.
                                                        NULL
       ident
                     literals
                                   properties
                                                        info
                                                                   derivation
```

# **Encoding Operations (except)**

```
typedef enum
   DONop.
   DOQuote,
   DORewrite,
   DOApplyDef.
   DOSR.
   DONegateConjecture,
   DOFofSimplify,
   DOFNNF,
   DOShiftQuantors,
   DOSkolemize.
   DOParamod,
   DOSimParamod.
   DOOrderedFactor,
}OpCode:
```

# **Encoding Operations (except)**

```
typedef enum
typedef enum
                             Arg1Fof = 1 < < 8,
   DONop.
                             Arg1Cnf = 1 << 9,
   DOQuote,
                             Arg1Num = 1 << 10,
   DORewrite,
                             Arg2Fof = 1 << 11
   DOApplyDef.
   DOSR.
                          }ArgDesc;
   DONegateConjecture,
   DOFofSimplify,
   DOFNNF.
   DOShiftQuantors,
   DOSkolemize.
   DOParamod,
   DOSimParamod,
   DOOrderedFactor,
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                             Arg1Fof = 1 < < 8,
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   DOQuote,
                             Arg1Num = 1 << 10,
   DORewrite.
                             Arg2Fof = 1 << 11
   DOApplyDef.
   DOSR.
                          }ArgDesc:
   DONegateConjecture,
                         typedef enum
   DOFofSimplify,
   DOFNNF.
                             DCNop
                                        = DONop.
   DOShiftQuantors,
                             DCCnfQuote = DOQuote | Arg1Cnf,
   DOSkolemize.
                             DCFofQuote = DOQuote | Arg1Fof.
   DOParamod.
                             DCRewrite = DORewrite | Arg1Cnf,
   DOSimParamod.
   DOOrderedFactor,
                             DCParamod = DOParamod | Arg1Cnf | Arg2Cnf,
}OpCode:
```

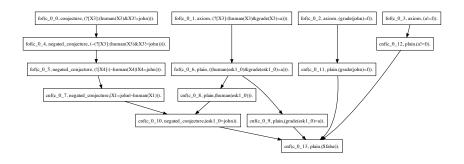
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- DerivationCodes inspired by machine instructions
  - Operation (OpCode)
  - Arguments
- Supports flexible processing
  - Proof structure is independent of exact inferences

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## Proofs as Graphs



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Derivation stacks allow a compact and efficient (in data and code) representation of clause history.

#### Challenges

# Unambiguous Inferences

- ► Complete inference records
  - Add inference positions (more arguments to OpCodes)
  - Add unifiers (if neccessary, e.g. HO)
- ► Complete clausification records
  - Clause simplification as rewriting (?)
  - Mini-scoping as rewriting (?)
  - Step-by-step skolemization

# Theoretically managable, but practically difficult – especially retroactively

#### **Proof Expansion**

- Calculus level expansion
  - Explicit results of each inference
  - Good for semantic proof checking
  - Good for understanding the structure of the proof
  - Potentially good for machine learning
- Primitive inferences
  - Convert inferences into primitive operations
  - For superposition:
    - Instantiation
    - Lazy conditional term replacement
    - Deleting trivial and duplicated literals
  - Uniform proof format for different provers/calculi
  - Uniform post-processing (proof checking, proof presentation, ...)

#### Conclusion

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- ► Efficient proof generation is non-trivial, but possible
- Clause objects and derivation stacks can represent internal inference sequences compactly
- TPTP v3 is a useful and used standard for external proof representation
- ► Proof objects are useful for trust building and learning
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Efficient and flexible proof representation is a key to useful proof output.

#### Ceterum Censeo...

- ▶ Bug reports for E should include:
  - The exact command line leading to the bug
  - All input files needed to reproduce the bug
  - A description of what seems wrong
  - ➤ The output of eprover --version

