Programming Assignment #2 - CS325

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Introduction

In this writeup, we will both assess time complexity and prove three algorithms designed to solve the Maximum Subarray problem.

According to Wikipedia:

'In computer science, the maximum subarray problem is the task of finding the contiguous subarray within a one-dimensional array of numbers (containing at least one positive number) which has the largest sum. For example, for the sequence of values 2, 1, 3, 4, 1, 2, 1, 5, 4; the contiguous subarray with the largest sum is 4, 1, 2, 1, with sum 6.'

Pseudocode

Brute Force:

```
def brute_force(x):
    max_sum = 0
    l = len(x)
    for i in xrange(l):
        new_sum = 0
        for j in range(i,l):
            new_sum += x[j]
            if new_sum > max_sum: max_sum = new_sum
    return max_sum
```

Divide and Conquer:

```
def div_and_conq(x):
  \max_{\text{sum}} = 0
  def inner(x, max_sum):
    if (len(x) \le 1): # base case
       if sum(x) > max_sum:
         \max_{\text{sum}} = \text{sum}(x)
       return max_sum
    mid = int(len(x)/2) \# split into two halves
    1 = x [: mid]
    r = x [mid:]
    \max_{sum} = \max(\max_{sum}, sum(1)) \# is \ left \ half \ max?
    \max_{sum} = \max(\max_{sum}, sum(r)) \# is \ right \ half \ max?
    \#\ does\ max\ consists\ of\ suffix+prefix?
    suffix_sum = 0
    max_suffix_sum = 0
    for i in range (\text{mid}-1,-1,-1):
       suffix_sum += x[i]
       max_suffix_sum = max(max_suffix_sum, suffix_sum)
    prefix_sum = 0
```

```
max_prefix_sum = 0
         for i in range(mid, len(x)):
            prefix_sum += x[i]
            max_prefix_sum = max(max_prefix_sum, prefix_sum)
         max_sum = max(max_sum, max_suffix_sum + max_prefix_sum)
         \# \ recursive \ calls
         ret = inner(l, max_sum)
         if (ret != None):
            \max_{\text{sum}} = \max(\text{ret}, \max_{\text{sum}})
         ret = inner(r, max_sum)
         if (ret != None):
           \max_{\text{sum}} = \max(\text{ret}, \max_{\text{sum}})
         return max_sum # end of inner function
       return inner(x, max_sum)
Dynamic Programming:
    def dynamic_prog(x):
       this\_sub\_arr\_sum = 0
       \max_{\text{sum}} = 0
       for i in x:
         if this_sub_arr_sum + i > 0:
            this_sub_arr_sum = this_sub_arr_sum + i
         else:
            this\_sub\_arr\_sum = 0
         if this_sub_arr_sum > max_sum:
           \max_{\text{sum}} = \text{this\_sub\_arr\_sum}
       return max_sum
```

Correctness Proofs

Divide and Conquer:

Dynamic Programming:

Asymptotic Analysis of Run Time

Brute Force:

Divide and Conquer:

Dynamic Programming:

Empirical Testing of Correctness

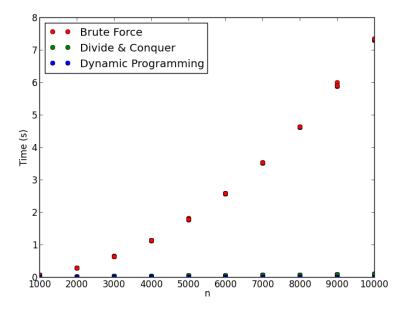
Correctness of the three algorithms was verified using the provided file of test cases $verify_2.txt$. See the function 'test_correctness' in main.py for details.

Here is the output from the final text input file (name.txt):

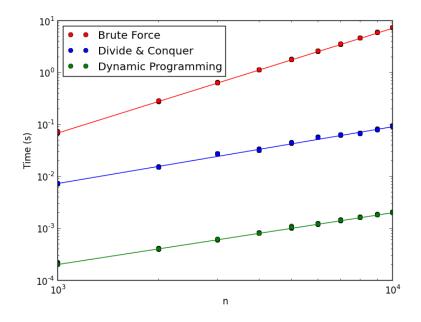
• OUTPUT HERE (TODO)

Empirical Analysis of Run Time

Linear Plot:



Log-log Plot:



Slope of lines in log-log plot:

The equation for the best fit line on the log-log plot (calculated using numpy.polyfit()) has the following form:

$$f(n) = e^{y-intercept} * n^{slope}$$

Brute Force:

• slope: 2.0177140011

Divide & Conquer:

 \bullet slope: 1.09279503259

Dynamic Programming:

 \bullet slope: .996887276493

Performance Comparison

'In your report, present and compare the empirical run time results of the three different algorithms. Provide a discussion of the comparative benefits and drawbacks of different algorithms.'