Near Complete Formal Semantics of X86-64

Abstract

ToDo

1. Introduction

2. Challenges

2.1. Using Strata Results

Following are the challenges in using Strata [1] (or Stoke) formula as is.

 Stoke uses C+-functions which define the semantics of instructions. For example, following is the function to define the semantics of add instruction. The functions are generic in the sense that they can be used to obtain obtain the concrete semantics of any instruction like add %rax, %rbx

The untested assumption here is the generic formula will behave identically for all the variants. We have tested all the formula for each instruction variant.

- Strata gives the concrete semantics for a concrete instructions. For other variants it generalize from the concrete semantics. Assumption is the generalization is correct. Test all the generalization.
- While porting to K rule, we generalize the from a concrete semantics that strata provides. Is this generalization faithful? For instruction like xchg, xadd, cmpxchg, the formula is different for different operands. So the general K rule we obtain from xchgl a, b may not represent the semantics for xchgl a, a. Fortunately there exists different instruction variants if the their semantics might be different and accordingly we might have different K rules. For example, xchgl_r32_eax and xchgl_r32_r32. But even for xchgl_r32_32 semantics could be different for cases r1 != r2 and r1 == r2. Idea: Once lifted as K rule, test the instruction for all variants.

Lets consider xaddb SRC, DEST, as per manual the semantics is as follows:

```
S1. Temp = Src + Dest
S2. Src = Dest
S3. Dest = Temp
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The point to note here is that the register updates follow an order. Strata uses xaddb %rax, %rbx, to obtain the semantics and it happened that the ordering is maintained and hence

strata can generalize the semantics of xaddb R1, R1. But even if the ordering is not maintained the semantics is going to be the same for the case R1! = R2, but the generalization for the R1 == R1 case will mess up. We cannot trust the above generalization by strata. We need to test the K rule for all possible operands.

3. Implementation

3.1. Porting Formulas for stratified instructions to K Rules

For the purpose, we could have directly converted the Strata formulas to K rule assuming that the Strata's symbolic execution over the stratified instruction sequence is correct.

Given that fact the K's symbolic execution engine is more trusted as that has been used extensively in language-agnostic manner to perform symbolic execution, we decided to use ...

- Implementing the base instructions semantics in K and testing them.
- Symbolic execution of the stratified instruction sequences.
- Dealing with scratch pad registers.
- Simplification of formulas. Simplification generates simpler K rule (than the corresponding Strata formula) in some cases.
- Check the the ported K rules are equivalent to the strata formulas to begin with. This helped us in finding a bug in Strata's symbolic execution engine which is anyway fixed by the authors in their recent versions.

3.2. Supporting un-stratified instructions & Porting their formulas to K Rules

3.2.1. Supporting un-stratified instructions

Instruction support status

3.3. Porting to K Rules

Strata could output the internal AST, used to model a register state formula, in different formats. Supported backend are SmtLib and Prefix notation. We have added another backend to generate K rule. We need some way to validate the backend.

Validate the Backend The K rules generated using the backend are matched against the ones we already obtained via symbolic execution on stratified instructions. This has an added benefit that in order to get the exact match we need to port all the verification lemma from K to strata code, which will later help in generating simplified K rules for non-stratified instructions.

Two main challenges in getting an exact match are:

• Strata rules uses *extract* to extract portion of a bit-vector. The high and low indices of *extract* are obtained considering

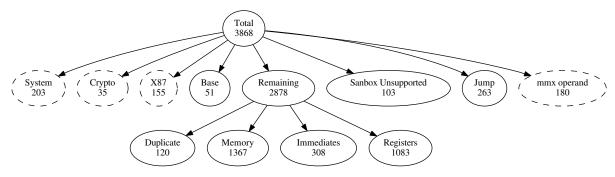


Figure 1: Instruction classification

LSB at index 0, whereas K uses *extractMInt* for the same purpose, but uses MSB at index zero.

• Strata uses flags as Bool, whereas they are treated as Bitvector in our semantics.

References

[1] Stefan Heule, Eric Schkufza, Rahul Sharma, and Alex Aiken. Stratified synthesis: Automatically learning the x86-64 instruction set. In *Proceedings of the 37th ACM SIGPLAN Conference on Programming Language Design and Implementation*, PLDI '16, pages 237–250, New York, NY, USA, 2016. ACM.