	•	•	•	•		•	•	•	•	•	•	•	•	•
A(x)					B(x)		/+. X	. •	٠	٠	٠	٠	٠	
n= H -=	= ,2 k	? ?P/)	7 29	5 M										
	•	•				٠	, 3	,	•	•	•	•	٠	•
$W_{h}^{0} = V$			1	W 4	= - 1		4	/						
FFT	(2)	(b)	·	·	·	·	,	, and the second	·	·	·	·	·	Ť
FFT	27/1 27/1		•			•	•	•	•	•	•	•	•	
.,														
WEA	•	•	•	•	•	•	٠	٠	•	•	•	•	•	•
$a^{2} = (2)$														
o, < (1)	•	•	•	•	•	•	•	•	•	•	•	•	•	•
y° < FFT	((2))	ج ر		•										•
y'efft														
			•			. 0(z	$A^{1}(x^{n})$	35 = 5	•	•	•	•	•	•
y = 2+(	1). 1 =[	$A_{\omega} = \hat{A}$	L, A	(x) = 1,	A(x) = A(	x)= AW	, ray res,	1						
y 3= 2 - 1	1.1=		•	•	•	•	•	٠	•	٠			•	•
y, = 2 +1	·1.=2	<u>)</u>									Z - Q	Li – i	+(-1)	
y4 = 2 - 1														
	· - 2	1	•			•		1 (2)		2/11	1 2 4 1	1/12	, ) . (13 -	)*
							7	(x)					1) (6,2	
FF T((1,1))	•	•	•	•	٠	•	· B	g(x)	("W"	2)(W)	1+i)	(w; 0)	) .(\v3, 1	- i)
a° < 1 a^ < 1														
y° = FFT (1)	•	•	•	•	•	•	. (	しり	(w, c	/ (W,	1+3 <i>i</i> J\	.**	) (W, 1	<del>-</del> 3i)
y'= FF7 ()		* 2	•		•	•			•	•	•	•	•	
$\int_{C} e^{-A^{\circ}(x)} + x A^{\prime}(x)$		(x)= 15			0 15		1+31	+ 6	1 1-2	; ] _	ລ			
	•	•	•		- = 4 (	6 7 6	•		+ 1-3		٠	•	•	•
3=1-1.1												= ?		
$y_{\lambda} = 1 + i \cdot \lambda$			•			_			[6-1+3				٠	•
4=1-1.1	= 1- 1			$C^{2}$	= 464	$(1+3i)\cdot W_{ij}$	+0+(1-	31) W., "	$\left[ \int = \frac{1}{4} \left[ \int d^2 x \right] \right]$	6-1-31+	0-1+3	]=1		
•	•	•	•			1					)		•	•
•		•	•	<u> </u>	(2)	3,1)	→ (	.(X) =	-2-	' 3 <sup>X</sup>	- ~	•	•	
•	•	•					٠		•	•		•	•	

FFT ((2)) FFT ((1))
return 2 return 1

$$P(x) = \sum_{j=0}^{h-1} c_j x^j \qquad g(x) = \sum_{i=0}^{h-1} c_i x^i \qquad w = e^{2\pi i/h} = \left[w^e, w^i, \dots, w^{h-1}\right]$$

$$\rho(x) = \rho(x^2) + \chi \rho(x^2)$$

$$\rho(-\chi) = \rho(\chi^2) - \chi \rho(\chi^2)$$

The Fast Fourier Transform (FFT): Most Ingenious Algorithm of the following problem of the following problem of the first point of the fast fourier Transform (FFT): Most Ingenious Algorithm of the fast fourier Transform (FFT): Most Ingenious Algorithm of the fast fourier Transform (FFT): Most Ingenious Algorithm of the fast fourier Transform (FFT): Most Ingenious Algorithm of the fast fourier Transform (FFT): Most Ingenious Algorithm of the fast fourier Transform (FFT): Most Ingenious Algorithm of the fast fourier Transform (FFT): Most Ingenious Algorithm of the fast fourier Transform (FFT): Most Ingenious Algorithm of the fast fourier Transform (FFT): Most Ingenious Algorithm of the fast fourier Transform (FFT): Most Ingenious Algorithm of the fast fourier Transform (FFT): Most Ingenious Algorithm of the fast fourier Transform (FFT): Most Ingenious Algorithm of the fast fourier Transform (FFT): Most Ingenious Algorithm of the fast fourier Transform (FFT): Most Ingenious Algorithm of the fast fourier Transform (FFT): Most Ingenious Algorithm of the fast fourier Transform (FFT): Most Ingenious Algorithm of the fast fourier Transform (FFT): Most Ingenious Algorithm of the fast fourier Transform (FFT): Most Ingenious Algorithm of the fast fourier Transform (FT): 
$$p(x) : p(x) : p$$

$$x_{k} = \frac{1}{n} \sum_{j=0}^{n-1} \chi_{j} W_{n}^{-jk}$$

- $x_k$  are the coefficients we want to recover,
- ullet  $X_j$  are the DFT values (the pointwise multiplication result),
- $\omega_n = e^{-2\pi i/n}$  is the primitive nth root of unity (inverse of  $\omega$  used in FFT),
- n is the number of points (in our case, 4)

Let's go through the Inverse FFT (IFFT) calculation in detail, starting from the pointwise multiplication results. The goal of the inverse FFT is to take the frequency-domain representation we obtained (after pointwise multiplication) and return it to the time-domain, or coefficient form, of the polynomial C(x).

## Setup for Inverse FFT

After the pointwise multiplication, we have:

$$C' = [6, 1+3i, 0, 1-3i]$$

Since we zero-padded to length 4, we're working with a vector length n=4, and we'll use the 4th roots of unity in the inverse direction to perform the inverse FFT.

The inverse FFT formula for a sequence  $C' = \left[C'(0), C'(1), C'(2), C'(3)\right]$  is:

$$C(k) = rac{1}{n} \sum_{j=0}^{n-1} C'(j) \cdot \omega_n^{-jk}$$

Step 2: Calculate C(1)

$$C(1) = \frac{1}{4} \left( C'(0) \cdot \omega_0^{-1} + C'(1) \cdot \omega_1^{-1} + C'(2) \cdot \omega_2^{-1} + C'(3) \cdot \omega_3^{-1} \right)$$

Substitute the roots of unity values:

$$C(1) = rac{1}{4} \left( 6 + (-i)(1+3i) + (-1) \cdot 0 + i(1-3i) 
ight)$$

Calculating each term:

1. 
$$(-i)(1+3i) = -i+3 = 3-i$$

2. 
$$i(1-3i)=i-3=-3+i$$

So,

$$C(1) = rac{1}{4} \left( 6 + (3-i) + 0 + (-3+i) 
ight)$$

Simplify:

$$C(1)=rac{1}{4}\cdot 12=3$$

- n = 4.
- $\omega_n^{-1}$  is the inverse root of unity, specifically  $e^{-2\pi i/4}=i$ ,
- We need to divide by n=4 at the end to normalize.

The 4th roots of unity (for n=4) in the inverse FFT are:

- $\omega_0^{-1} = 1$
- $\omega_1^{-1}=-i$
- $\omega_2^{-1} = -1$
- $\omega_3^{-1}=i$

Let's calculate each C(k) for k=0,1,2,3.

Calculation of Each C(k)

Step 1: Calculate C(0)

$$C(0) = \frac{1}{4} \left( C'(0) \cdot \omega_0^{-1} + C'(1) \cdot \omega_0^{-1} + C'(2) \cdot \omega_0^{-1} + C'(3) \cdot \omega_0^{-1} \right)$$

Since  $\omega_0^{-1}=1$ , this becomes:

$$C(0) = rac{1}{4}(6+(1+3i)+0+(1-3i)) = rac{1}{4}(6+1+3i+1-3i)$$
  $C(0) = rac{1}{4} \cdot 8 = 2$ 

Step 3: Calculate C(2)

$$C(2) = rac{1}{4} \left( C'(0) \cdot \omega_0^{-2} + C'(1) \cdot \omega_1^{-2} + C'(2) \cdot \omega_2^{-2} + C'(3) \cdot \omega_3^{-2} 
ight)$$

Since  $\omega_2^{-2}=-1$ , this simplifies to:

$$C(2) = rac{1}{4}(6-(1+3i)+0-(1-3i))$$

Simplify further:

1. 
$$-(1+3i) = -1-3i$$

2. 
$$-(1-3i)=-1+3i$$

So:

$$C(2) = \frac{1}{4}(6-1-3i-1+3i) = \frac{1}{4}\cdot 4 = 1$$

Step 4: Calculate C(3)

$$C(3) = rac{1}{4} \left( C'(0) \cdot \omega_0^{-3} + C'(1) \cdot \omega_1^{-3} + C'(2) \cdot \omega_2^{-3} + C'(3) \cdot \omega_3^{-3} 
ight)$$

Substitute:

$$C(3) = rac{1}{4} \left( 6 + i(1+3i) + (-1) \cdot 0 - i(1-3i) 
ight)$$

Calculating:

1. 
$$i(1+3i)=i+3=3+i$$

2. 
$$-i(1-3i) = -i+3 = 3-i$$

Then.

$$C(3) = \frac{1}{4} (6 + (3+i) + 0 + (3-i)) = \frac{1}{4} \cdot 12 = 3$$

Thus, the coefficients of the polynomial  $C(x)=2+3x+x^2$  are:

- C(0)=2 (constant term),
- C(1) = 3 (coefficient of x),
- C(2) = 1 (coefficient of  $x^2$ ).

**Final Result** 

Thus, the coefficients of the polynomial  $C(x)=2+3x+x^2$  are:

where:

- C(0) = 2 (constant term),
- C(1) = 3 (coefficient of x),
- C(2) = 1 (coefficient of  $x^2$ ).

The Inverse FFT has given us the coefficients in the time domain, confirming that:

$$C(x) = 2 + 3x + x^2$$