

Algorithmics

Correction Final Exam #3 (P3)

UNDERGRADUATE 2nd YEAR - S3 – EPITA

January 5, 2021 - 9 : 30

Solution 1 (In the depth of the spanning forest – 3 points)

1. Spanning forest and extra-edges for the depth-first search of the graph in figure 1:

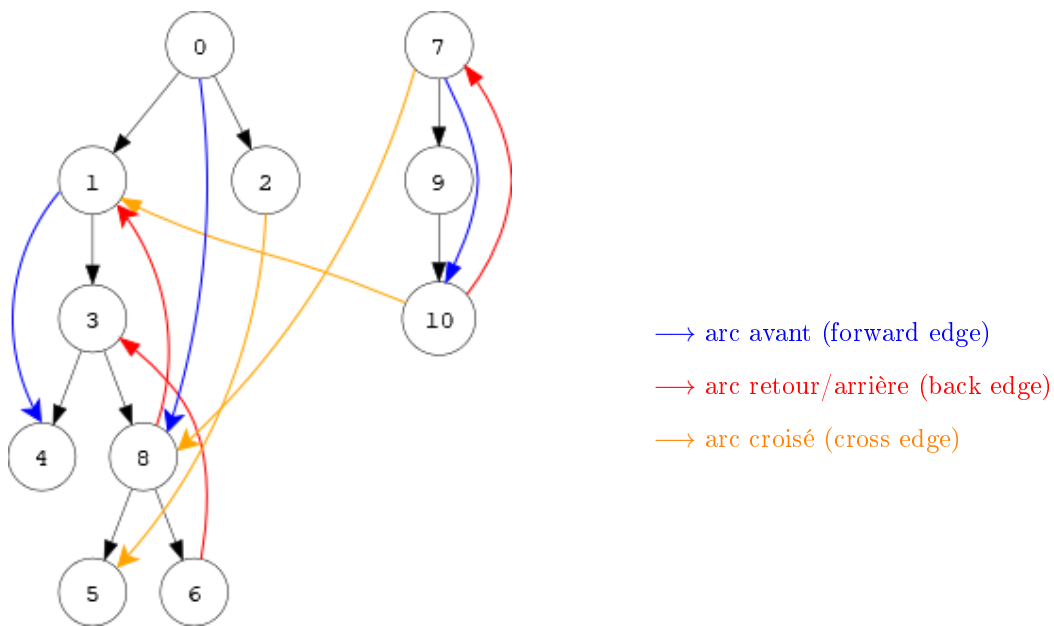


Figure 1: DFS: Spanning forest

2. Meeting orders in prefix **pref** and suffix **suff**:

	0	1	2	3	4	5	6	7	8	9	10
pref	1	2	14	3	4	7	9	17	6	18	19
suff	16	13	15	12	5	8	10	22	11	21	20

Solution 2 (Union-Find – 4 points)

1. Number of vertices of each connected component:

$C_1 : 4$ $C_2 : 6$ $C_3 : 4$

2. Edges to add: **two** among 5 – 8 8 – 12 5 – 12 for example...

3. Among the following chains, those which can not exist in G :

$\square 3 \leftrightarrow 7$ $\boxtimes 11 \leftrightarrow 6$ $\boxtimes 0 \leftrightarrow 13$ $\square 4 \leftrightarrow 9$

4. Vector p after adding the edge 7–4:

	0	1	2	3	4	5	6	7	8	9	10	11	12	13
p	5	8	5	8	12	-4	5	8	-10	12	8	12	8	8

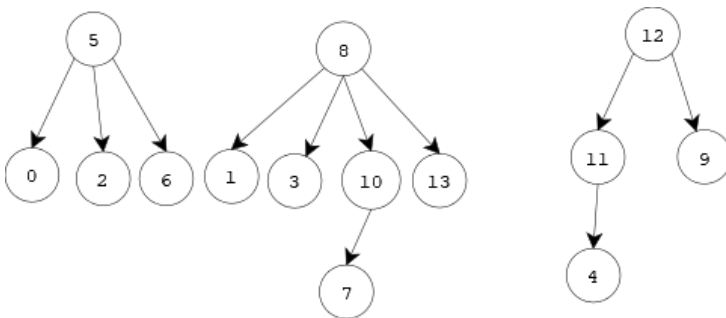


Figure 2: before

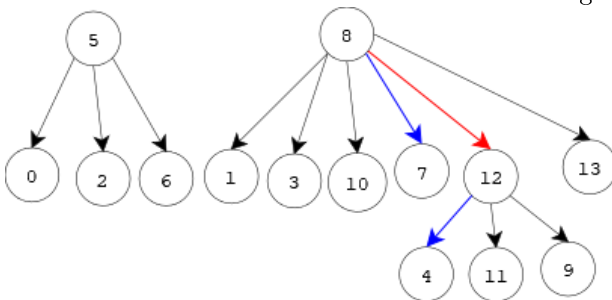


Figure 3: after edge 7 – 4 added

Solution 3 (Distance from start – 5 points)

Specifications: `dist_range(G , src , $dmin$, $dmax$)` returns the list of vertices that are at a distance between $dmin$ and $dmax$ from the vertex src in the graph G (with $0 < dmin \leq dmax$).

```

1 def dist_range(G, src, dmin, dmax):
2     dist = [None] * G.order
3     q = queue.Queue()
4     q.enqueue(src)
5     dist[src] = 0
6     L = []
7     while not q.isempty():
8         x = q.dequeue()
9         if dist[x] >= dmin:
10             L.append(x)
11         if dist[x] < dmax:
12             for y in G.adjlsts[x]:
13                 if dist[y] == None:
14                     dist[y] = dist[x] + 1
15                     q.enqueue(y)
16     return L

```

Solution 4 (Get cycle – 5 points)

Specifications:

the function `get_cycle(G)` returns a cycle of the undirected graph G , an empty list if G is acyclic.

Version 1: using parent vector

```
1 def __get_cycle(G, x, parent):           # DFS on G from x, interrupted at first back
    edge found
2     for y in G.adjlists[x]:             # parent: vertices marked with their parent
3         if parent[y] == None:           # return first back edge found (x, y) or None
4             get = __get_cycle(G, y, parent)
5             if get != None:
6                 return get
7         else:
8             if y != parent[x]:
9                 return (x, y)
10    return None
11
12 def get_cycle(G):
13     parent = [None] * G.order
14     s = 0
15     get = None
16     while s < G.order and get == None:
17         if parent[s] == None:
18             parent[s] = -1
19             get = __get_cycle(G, s, parent)
20         s += 1
21     L = []
22     if get != None:
23         (x, y) = get
24         L = [x]
25         while x != y:
26             x = parent[x]
27             L.append(x)
28         L.append(L[0])
29     return L
```

Version 2: the cycle is built by the recursive function in going up
Many ways to do it. The difficulty: no longer add vertices when the cycle is complete.

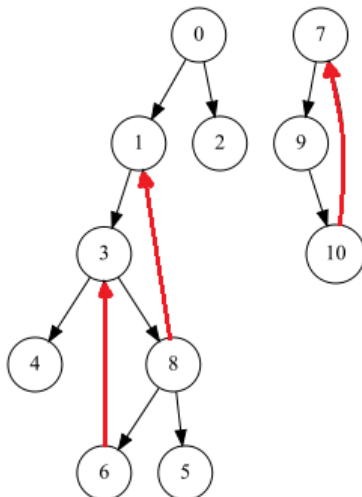
```

1 def __get_cycle2(G, x, M, p):
2     """
3     DFS on G from x
4     M: mark vector (boolean)
5     p: x's parent
6     return (cycle, done):
7     - cycle = the vertices of the first cycle found, [] if no cycle
8     - done: boolean: is the cycle completed?
9     """
10    M[x] = True
11    for y in G.adjlists[x]:
12        if not M[y]:
13            (cycle, done) = __get_cycle2(G, y, M, x)
14            if cycle:
15                if done:
16                    return (cycle, True)
17                if cycle[0] != y:
18                    cycle.append(y)
19                return (cycle, cycle[0] == y)
20            else:
21                if y != p:
22                    return ([y], False)
23    return ([], False)
24
25 def get_cycle_2(G):
26     M = [False] * G.order
27     for s in range(G.order):
28         if not M[s]:
29             cycle, done = __get_cycle2(G, s, M, -1)
30             if cycle:
31                 return cycle + [cycle[0]]
32    return []

```

Solution 5 (What is this? – 3 points)

1. The built graph (NG):



2. For each vertex s , during the traversal:

(a) What does $D[s]$ represent?

$D[s]$ is None if the vertex s has not been met. Otherwise it is the depth of s in the spanning forest of the DFS.

(b) What does $P[s]$ represent?

$P[s]$ indicates whether s was encountered in suffix.