Stage-3 Report

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实验内容

step7

在Namer.visitBlock中,在访问块内部每条语句之前,创建一个局部作用域对象,再调用ctx.open以开启这个新的作用域,将其加入作用域栈顶。而访问块内部每条语句结束后,调用ctx.close将此作用域关闭,即退栈。

```
def visitBlock(self, block: Block, ctx: ScopeStack) -> None:
    ctx.open(Scope(ScopeKind.LOCAL))
    for child in block:
        child.accept(self, ctx)
    ctx.close()
```

实现了CFG.unreachable, 具体实现方式为:

新增了一个成员数组reachable,代表从0号基本块能否到达该号基本块。在_init_并对其做全False初始化后,从0号基本块开始进行DFS,将访问到的每个基本块的reachable数组对应项置为True,表示可达。

值得注意的是,由于**step8**中加入了循环语句,导致控制流图中出现了环路,因此,每个被访问到的基本块,在递归地访问其后继的各个基本块时,只访问那些尚未确定可达,即满足*reachable[i] == False*的基本块,以避免无穷调用。

如此,在初始化阶段,就确定了每个基本块是否可达。*CFG.unreachable*直接根据成员*reachable*数组的内容,返回对应项的值取反即可。

```
def __init__(self, nodes: list[BasicBlock], edges: list[(int, int)]) ->
None:
    self.nodes = nodes
    self.edges = edges

self.links = []

for i in range(len(nodes)):
    self.links.append((set(), set()))

for (u, v) in edges:
    self.links[u][1].add(v)
    self.links[v][0].add(u)

self.reachable = [False for _ in range(len(nodes))]
    self.dfs(0)
```

```
def dfs(self, root):
    self.reachable[root] = True
    for child in self.getSucc(root):
        if not self.reachable[child]:
            self.dfs(child)
```

```
def unreachable(self, id):
    return not self.reachable[id]
```

修改了BruteRegAlloc.accept,对控制流图中的每个基本块bb,先调用bb.unreachable,若其不可达,则跳过它。

```
def accept(self, graph: CFG, info: SubroutineInfo) -> None:
    subEmitter = self.emitter.emitSubroutine(info)
    for bb in graph.iterator():
        # you need to think more here
        # maybe we don't need to alloc regs for all the basic blocks
        if graph.unreachable(bb.id):
            continue
        if bb.label is not None:
            subEmitter.emitLabel(bb.label)
        self.localAlloc(bb, subEmitter)
        subEmitter.emitEnd()
```

step8

在tree.py中添加了For、DoWhile、Continue三个类,作为分别对应for、do while、continue语句的AST节点类,其中大部分方法的实现可参照While和Break类。

值得注意的是,for语句的init既可能是声明也可能是表达式,而初始化参数中除循环体body外,init、cond、update都可能为空。其中,若cond为空,对应的数据成员应为整数字面量1,即IntLiteral(1)。

```
class For(Statement):
   AST node of for statement.
   def __init__(
       self, init: Declaration | Expression, cond: Expression, update:
Expression, body: Statement) -> None:
       super().__init__('for')
       self.init = init
       self.cond = cond if not cond is NULL else IntLiteral(1)
       self.update = update
       self.body = body
   def __getitem__(self, key: int) -> Node:
       return (self.init, self.cond, self.update, self.body)[key]
   def __len__(self) -> int:
       return 4
   def accept(self, v: Visitor[T, U], ctx: T):
       return v.visitFor(self, ctx)
```

```
class Dowhile(Statement):
    """

AST node of do while statement.
    """

def __init__(self, body: Statement, cond: Expression) -> None:
    super().__init__('do while')
    self.body = body
    self.cond = cond

def __getitem__(self, key: int) -> Node:
    return (self.body, self.cond)[key]

def __len__(self) -> int:
    return 2

def accept(self, v: Visitor[T, U], ctx: T):
    return v.visitDowhile(self, ctx)
```

```
class Continue(Statement):
    """

AST node of continue statement.
    """

def __init__(self) -> None:
        super().__init__('continue')

def __getitem__(self, key: int) -> Node:
        raise __index_len_err(key, self)

def __len__(self) -> int:
        return 0

def accept(self, v: Visitor[T, U], ctx: T):
        return v.visitContinue(self, ctx)

def is_leaf(self):
    return True
```

而在Vistor类下添加了visitFor、visitDoWhile、visitContinue。

```
def visitFor(self, that: For, ctx: T) -> Optional[U]:
    return self.visitOther(that, ctx)

def visitDoWhile(self, that: DoWhile, ctx: T) -> Optional[U]:
    return self.visitOther(that, ctx)

def visitContinue(self, that: Continue, ctx: T) -> Optional[U]:
    return self.visitOther(that, ctx)
```

```
# Reserved keywords
reserved = {
    "return": "Return",
    "int": "Int",
    "if": "If",
    "else": "Else",
    "while": "While",
    'for': 'For',
    'do': 'Do',
    'continue': 'Continue',
    "break": "Break",
}
```

在ply_parser.py下添加了for循环、do while循环、continue语句的对应文法。由于for循环的init、cond、update都可能为空,故产生式右侧使用了非终结符opt_expression。

```
def p_for(p):
    """
    statement_matched : For LParen opt_expression Semi opt_expression Semi
opt_expression RParen statement_matched
    statement_matched : For LParen declaration Semi opt_expression Semi
opt_expression RParen statement_matched
    statement_unmatched : For LParen opt_expression Semi opt_expression Semi
opt_expression RParen statement_unmatched
    statement_unmatched : For LParen declaration Semi opt_expression Semi
opt_expression RParen statement_unmatched
    """
    p[0] = For(p[3], p[5], p[7], p[9])
```

```
def p_do_while(p):
    """
    statement_matched : Do statement_matched while LParen expression RParen Semi
    statement_unmatched : Do statement_unmatched while LParen expression RParen
Semi
    """
    p[0] = Dowhile(p[2], p[5])
```

```
def p_continue(p):
    """
    statement_matched : Continue Semi
    """
    p[0] = Continue()
```

实现了Namer.visitFor, 即:

- ①为for循环开启一个新的局部作用域;
- ②若init非空,则访问之;访问cond;若update非空,则访问之;
- ③调用ctx.openLoop记录循环层数;
- ④访问body;
- ⑤调用ctx.closeLoop, 代表该层循环结束;

```
def visitFor(self, stmt: For, ctx: ScopeStack) -> None:
    1. Open a local scope for stmt.init.
    2. Visit stmt.init, stmt.cond, stmt.update.
    3. Open a loop in ctx (for validity checking of break/continue)
    4. Visit body of the loop.
    5. Close the loop and the local scope.
    ctx.open(Scope(ScopeKind.LOCAL))
    if not stmt.init is NULL:
        stmt.init.accept(self, ctx)
    stmt.cond.accept(self, ctx)
    if not stmt.update is NULL:
        stmt.update.accept(self, ctx)
    ctx.openLoop()
    stmt.body.accept(self, ctx)
    ctx.closeLoop()
    ctx.close()
```

实现了Namer.visitDoWhile, 即:

- ①调用ctx.openLoop记录循环层数;
- ②访问body;
- ③调用ctx.closeLoop, 代表该层循环结束;
- ④访问cond;

```
def visitDowhile(self, stmt: Dowhile, ctx: ScopeStack) -> None:
    """
    1. Open a loop in ctx (for validity checking of break/continue)
    2. Visit body of the loop.
    3. Close the loop.
    4. Visit the condition of the loop.
    """
    ctx.openLoop()
    stmt.body.accept(self, ctx)
    ctx.closeLoop()
    stmt.cond.accept(self, ctx)
```

实现了Namer.visitContinue,和Namer.visitBreak行为一致,只需通过检查当前循环层数,确保continue 语句位于某循环内,否则抛出DecafContinueOutsideLoopError。

```
def visitContinue(self, stmt: Continue, ctx: ScopeStack) -> None:
    """
    1. Refer to the implementation of visitBreak.
    """
    if not ctx.inLoop():
        raise DecafContinueOutsideLoopError()
```

前两方法同TACGen.visitWhile实现类似,调整循环体、控制表达式、跳转标签等部分的位置即可。

TACGen.visitContinue也与TACGen.visitBreak类似。

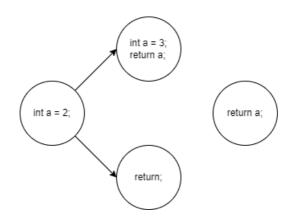
```
def visitFor(self, stmt: For, mv: FuncVisitor) -> None:
        beginLabel = mv.freshLabel()
        loopLabel = mv.freshLabel()
        breakLabel = mv.freshLabel()
        mv.openLoop(breakLabel, loopLabel)
        if not stmt.init is NULL:
            stmt.init.accept(self, mv)
        mv.visitLabel(beginLabel)
        stmt.cond.accept(self, mv)
        mv.visitCondBranch(tacop.CondBranchOp.BEQ, stmt.cond.getattr('val'),
breakLabel)
        stmt.body.accept(self, mv)
        mv.visitLabel(loopLabel)
        if not stmt.update is NULL:
            stmt.update.accept(self, mv)
        mv.visitBranch(beginLabel)
        mv.visitLabel(breakLabel)
        mv.closeLoop()
```

```
def visitDowhile(self, stmt: Dowhile, mv: FuncVisitor) -> None:
    beginLabel = mv.freshLabel()
    loopLabel = mv.freshLabel()
    breakLabel = mv.freshLabel()
    mv.openLoop(breakLabel, loopLabel)
    mv.visitLabel(beginLabel)
    stmt.body.accept(self, mv)
    mv.visitLabel(loopLabel)
    stmt.cond.accept(self, mv)
    mv.visitCondBranch(tacop.CondBranchOp.BEQ, stmt.cond.getattr('val'),
    breakLabel)
    mv.visitBranch(beginLabel)
    mv.visitLabel(breakLabel)
    mv.visitLabel(breakLabel)
    mv.closeLoop()
```

```
def visitContinue(self, stmt: Continue, mv: FuncVisitor) -> None:
    mv.visitBranch(mv.getContinueLabel())
```

思考题

step7



step8

设循环体执行了n次。

则两种翻译方式下*cond*和*body*部分执行次数一样,因此在这两部分,两种翻译方式下*CPU*执行的指令条数一样多。

除去cond和body部分指令,第一种翻译方式下,CPU执行的指令条数为2n+1;而第二种翻译方式下,CPU执行的指令条数为n+1。

由于循环体执行次数一定的情况下,执行的指令条数越少越好,而第二种翻译方式少执行n条指令,且考虑到实际情况中,一段程序中所有while循环的循环体执行次数均为0的概率较低,因此第二种翻译方式更好。