

Welcome to the new Discuss.

This is a place to **post interview questions** (/category/5/interview-questions) or share solutions / ask questions related to OJ problems (https://leetcode.com/problems).

Generate Parentheses (/category/30/generate-parentheses)

/ An iterative method. 📄 (/topic/3474.rss)

▲
114
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● **left.peter** (/user/left-peter)

(/user/left-peter) Reputation: ★ 156

My method is DP. First consider how to get the result $f(n)$ from previous result $f(0) \dots f(n-1)$. Actually, the result $f(n)$ will be put an extra $()$ pair to $f(n-1)$. Let the "(" always at the first position, to produce a valid result, we can only put ")" in a way that there will be i pairs $()$ inside the extra $()$ and $n - 1 - i$ pairs $()$ outside the extra pair.

Let us consider an example to get clear view:

$f(0)$: ""

$f(1)$: "(" $f(0)$)"

$f(2)$: "(" $f(0)$)" $f(1)$, "(" $f(1)$)"

$f(3)$: "(" $f(0)$)" $f(2)$, "(" $f(1)$)" $f(1)$, "(" $f(2)$)"

So $f(n) = "("f(0)"f(n-1) , "("f(1)"f(n-2) "("f(2)"f(n-3) \dots "("f(i)"f(n-1-i) \dots "(f(n-1))"$

Below is my code:

```

public class Solution
{
    public List<String> generateParenthesis(int n)
    {
        List<List<String>> lists = new ArrayList<>();
        lists.add(Collections.singletonList(""));

        for (int i = 1; i <= n; ++i)
        {
            final List<String> list = new ArrayList<>();

            for (int j = 0; j < i; ++j)
            {
                for (final String first : lists.get(j))
                {
                    for (final String second : lists.get(i - 1 - j))
                    {
                        list.add("(" + first + ")" + second);
                    }
                }
            }

            lists.add(list);
        }

        return lists.get(lists.size() - 1);
    }
}

```

2 years ago (/topic/3474/an-iterative-method/1)

JAVA 7.4k (/tags/java)

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<https://leetcode.com/problems/generate-parentheses>



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weird (/user/weird)
Reputation: ★ 63

Nice! thanks.

2 years ago (/topic/3474/an-iterative-method/2)



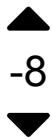
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jlu_chuang (/user/jlu_chuang)
Reputation: ★ 1

Nice ! ! ! ! ! ! ! !

2 years ago (/topic/3474/an-iterative-method/3)



-8



● **ky512** (/user/ky512)

(/user/ky512) Reputation: ★ 21

Thanks a lot ! Your solution is perfect!!!

about a year ago (/topic/3474/an-iterative-method/4)



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● **wenyan** (/user/wenyan)

(/user/wenyan) Reputation: ★ 0

very good solution!

about a year ago (/topic/3474/an-iterative-method/5)



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● **jingb** (/user/jingb)

(/user/jingb) Reputation: ★ 0

pretty awesome

about a year ago (/topic/3474/an-iterative-method/6)



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● **jackwhit3** (/user/jackwhit3)

(/user/jackwhit3) Reputation: ★ 0

I don't get this 😞

How do you know/prove this is correct?

about a year ago (/topic/3474/an-iterative-method/7)



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● **seekerwu** (/user/seekerwu)

(/user/seekerwu) Reputation: ★ 0

This post is deleted!

10 months ago (/topic/3474/an-iterative-method/8)



12



● **w0mTea** (/user/w0mtea)

(/user/w0mtea) Reputation: ★ 12

I think it's useful to prove this equation.

The equation is equivalent to the following one:

$$f(n) = (f(0))f(n-1) + (f(1))f(n-2) + \dots + (f(n-2))f(1) + (f(n-1))f(0)$$

First, let $f(n)$ to be a correct solution set when there is n pair of parentheses.

This means every combination in $f(n)$ is a valid combination, and any combination which isn't in $f(n)$ is not a valid combination for n .

And we can easily get the first three solution sets i.e. $f(0) = \{""\}$, $f(1) = \{ "()"\}$ $f(2) = \{ "(()())", "((()))"$.

For any $n > 2$, each combination of $f(n)$ can be divided into two parts p_0 and p_1 .
 p_0 and p_1 has several properties:

1. Parentheses in both p_0 and p_1 can match well
2. p_0 should be as short as possible but not empty. This means that p_0 belongs to $(f(l_0-1))$ where l_0 is the number of pairs in p_0 .
This property can be proved easily. Shortest means the first left parenthesis in this combination always matches the last right parenthesis.
So without these two, what left is also a legal combination.

Now, let's reorganize $f(n)$ by p_0 .

Put those combinations with same length of p_0 into a same set, and then $f(n)$ is divided into several subsets.

Each combination in subset s whose p_0 has l_0 pair of parentheses also belongs to the set $(f(l_0-1))f(n-l_0)$.

So we can get $f(n)$ belongs to $(f(0))f(n-1) + (f(1))f(n-2) + \dots + (f(n-2))f(1) + (f(n-1))f(0)$.

OK, the only thing to prove now is $(f(0))f(n-1) + (f(1))f(n-2) + \dots + (f(n-2))f(1) + (f(n-1))f(0)$ also belongs to $f(n)$.

Notice that each combination in $(f(i))f(n-1-i)$ is a legal combination for n , and we've declared before that each legal combination for n belongs to $f(n)$.

So each combination in the left side of equation belongs to $f(n)$, and the left side as a whole set belongs to $f(n)$.

Prove complete.

10 months ago (/topic/3474/an-iterative-method/9)



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● **Csnerds** (/user/csnerds)

(/user/csnerds) Reputation: ★ 0

what is the time complexity for this one? $O(n^4)$?

10 months ago (/topic/3474/an-iterative-method/10)



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● **mehere** (/user/mehere)

(/user/mehere) Reputation: ★ 11

Another iterative method with complexity $O(n^2)$

```

LinkedList<String> queueBracket = new LinkedList<>();
queueBracket.add("(");

// 0 means # of Left brackets; 1 means # of right brackets
LinkedList<List<Integer>> queueBracketNum = new LinkedList<>();
queueBracketNum.add(Arrays.asList(new Integer[]{1, 0}));

for (int i = 1; i <= n * 2 - 1; i++) {
    while (queueBracket.peek().length() == i) {
        String bracket = queueBracket.remove();
        List<Integer> bracketNum = queueBracketNum.remove();

        if (bracketNum.get(0) < n) {
            queueBracket.add(bracket + "(");
            queueBracketNum.add(Arrays.asList(new Integer[]{bracketNum.get(0) + 1,
        }

        if (bracketNum.get(0) > bracketNum.get(1) && bracketNum.get(1) < n) {
            queueBracket.add(bracket + ")");
            queueBracketNum.add(Arrays.asList(new Integer[]{bracketNum.get(0), bra
        }
    }
}

return queueBracket;

```

8 months ago (/topic/3474/an-iterative-method/11)



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dragonfly9113 (/user/dragonfly9113) Reputation: ★ 0

Thanks for the explanations!

5 months ago (/topic/3474/an-iterative-method/12)



2



wfxr (/user/wfxr) Reputation: ★ 27

Too many loops and containers. How about this:

```
vector<string> generateParenthesis(int n) {
    vector<string> result;
    if (!n) return result;

    string s(n, '(');
    s.append(n, ')');

    for (;;) {
        auto l = n, r = n;
        result.push_back(s);
        for (;;) {
            if (s.back() == ')') --r;
            else if (l < r + 2) --l;
            else break;
            s.pop_back();
            if (s.empty()) return result;
        }

        s.back() = ')';
        s.append(n - (l - 1), '(');
        s.append(n - (r + 1), ')');
    }
}
```

4 months ago (/topic/3474/an-iterative-method/13)



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neer1304 (/user/neer1304)

Reputation: ★ 0

(/user/neer1304)

Can you explain your approach ?

3 months ago (/topic/3474/an-iterative-method/14)



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hongyi5 (/user/hongyi5)

Reputation: ★ 0

(/user/hongyi5)

I got a "Memory Limit Exceeded" using this method

19 days ago (/topic/3474/an-iterative-method/15)

JAVA 7.4k (/tags/java)

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