

# Algebraic Topology 1: Exercise 3

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## Exercise 3.1

(a) Recall the definition of geometric realization: For a simplicial set  $X$ , its geometric realization  $|X| = \coprod X_n \times \nabla^n / \sim$ , equipped with quotient topology inherited from  $\coprod X_n \times \nabla^n$ . Thus by the properties of quotient topology, a continuous map  $\coprod X_n \times \nabla^n \rightarrow T$  induces a continuous map  $|X| \rightarrow T$  if and only if it factors through the quotient map:

$$\begin{array}{ccc} \coprod X_n \times \nabla^n & \xrightarrow{\quad} & T \\ \downarrow q & \searrow & \uparrow \\ \coprod X_n \times \nabla^n / \sim & & \end{array}$$

Thus we only need to check that equivalent elements under  $\sim$  are sent to the same target.

Recall that the equivalence relation  $\sim$  is spanned by  $(x, \alpha_*(t)) \sim (\alpha^*(x), t)$  for all  $x \in X_m, t \in \nabla^n, \alpha : [n] \rightarrow [m]$ . Thus we only need to ensure that  $f_m(x)(\alpha_*(t)) = f_n(\alpha^*(x))(t)$ . However, this equation follows directly from the following commutative diagram, which depicts the definition of  $f : X \rightarrow \mathcal{S}(T)$  being a morphism. (recall that  $\alpha^*$  on the right is just precompose with  $\alpha_*$ ).

$$\begin{array}{ccc} X_m & \xrightarrow{f_m} & \mathcal{S}(T)_m \\ \downarrow \alpha^* & & \downarrow - \circ \alpha_* \\ X_n & \xrightarrow{f_n} & \mathcal{S}(T)_n \end{array}$$

Thus we conclude that each morphism  $f : X \rightarrow \mathcal{S}(T)$  induces a map  $\hat{f} : |X| \rightarrow T$ .

(b) For notation reasons, we start by denoting the map described in the previous subproblem as  $\phi : \text{Hom}_{sim}(X, \mathcal{S}(T)) \rightarrow \text{Hom}_{TOP}(|X|, T)$  by sending  $f \mapsto \phi(f) = \hat{f}$ . We want to show that  $\phi$  is bijective. We proceed by constructing an explicit inverse  $\psi : \text{Hom}_{TOP}(|X|, T) \rightarrow \text{Hom}_{sim}(X, \mathcal{S}(T))$  described as follows:

Given a function  $g : \coprod (X_n \times \nabla^n) / \sim \rightarrow T$ ,  $\psi(g)$  is a morphism  $X \rightarrow \mathcal{S}(T)$  such that  $\psi(g)_n$  sends  $x_n \in X_n$  to the map  $g(x_n, -) \in \mathcal{S}_n(T)$ .

Now we need to check that  $\phi$  and  $\psi$  are indeed inverse to each other.

$(\phi \circ \psi)$ : Given a map  $g : \coprod (X_n \times \nabla^n) / \sim \longrightarrow T$ , by definition  $\psi(g)_n : X_n \rightarrow \mathcal{S}(T)$  with  $x \mapsto g(x, -)$ . Then by composing  $\phi$ , we end up with a map  $\phi(\psi(g)) : \coprod (X_n \times \nabla^n) / \sim \longrightarrow T$  sending  $(x, t) \mapsto \psi(g)_n(x)(t) = g(x, t)$ . Thus this composition is identity.

$(\psi \circ \phi)$ : Given a morphism  $f : X \rightarrow \mathcal{S}(T)$ , we know by definition that  $\phi(f) : \coprod (X_n \times \nabla^n) / \sim \longrightarrow T$  with  $(x, t) \mapsto f_n(x)(t)$ . Thus by composing  $\psi$ , we have the map  $\psi(\phi(f)) : X \rightarrow \mathcal{S}(T)$  such that  $\psi(\phi(f))_n : X_n \rightarrow \mathcal{S}(T)_n$  by sending  $x \mapsto \phi(f)(x, -) = f_n(x)(-)$ , thus is the same map as  $f$ . Thus this composition is also identity.

Finally we conclude that the map  $\phi$  is a bijection.

(c) We want to show that  $\phi$  is natural in both diagram, i.e. we need to show that for all morphisms  $\alpha : Y \rightarrow X$  between simplicial sets, all continuous maps between topological spaces  $h : T \rightarrow T'$ , the following diagram commutes:

$$\begin{array}{ccc} \text{Hom}_{\mathcal{S}}(X, \mathcal{S}(T)) & \xrightarrow{\phi} & \text{Hom}_{TOP}(|X|, T) \\ \downarrow & & \downarrow \\ \text{Hom}_{\mathcal{S}}(Y, \mathcal{S}(T')) & \xrightarrow{\phi} & \text{Hom}_{TOP}(|Y|, T') \end{array}$$

We prove commutativity by directly chasing the diagram, starting from the morphism  $f : X \rightarrow \mathcal{S}(T)$ .

(Go down then go right): We first go down. The morphism  $f$  is mapped to a morphism  $Y \rightarrow \mathcal{S}(T')$ , sending  $y_n \in Y_n$  to  $h_* \circ f_n \circ \alpha_n(y, n) \in \mathcal{S}(T')_n$ . This means if we apply  $\phi$  (go right), we get a map  $\coprod Y_n \times \nabla^n / \sim \longrightarrow T'$ , sending  $(y_n, t_n) \mapsto h(f_n(\alpha_n(y_n))(t_n))$ .

(Go right then go down): We first apply  $\phi$ , which means we get a map  $\coprod X_n \times \nabla^n / \sim \longrightarrow T$  sending  $(x_n, t_n) \mapsto f_n(x_n)(t_n)$ . Then we go down, and we again get the map  $\coprod Y_n \times \nabla^n / \sim \longrightarrow T'$ , sending  $(y_n, t_n) \mapsto h(f_n(\alpha_n(y_n))(t_n))$ .

Thus we conclude that the above diagram commutes, and thus  $\phi$  is natural with respect to both variables, and  $| \cdot |$  and  $\mathcal{S}$  are adjoint functor pair.

**Exercise 3.2** For any topological space  $T$ , we have the following calculation based on the adjunction we have from problem 1).

$$\begin{aligned}
Hom_{TOP}(|\Delta[n] \times \Delta[1]|, T) &= Hom_{\mathcal{S}}(\Delta[n] \times \Delta[1], \mathcal{S}(T)) \\
&= Hom_{\mathcal{S}}(\Delta[n], \mathcal{S}(T)) \times Hom_{\mathcal{S}}(\Delta[1], \mathcal{S}(T)) \\
&= Hom_{TOP}(|\Delta[n]|, T) \times Hom_{TOP}(|\Delta[1]|, T) \\
&= Hom_{TOP}(|\Delta[n]| \times |\Delta[1]|, T)
\end{aligned}$$

Finally we apply Yoneda Lemma and we have  $|\Delta[n] \times \Delta[1]| \cong |\Delta[n]| \times |\Delta[1]|$ .

### Exercise 3.3

(a) Consider the following composition of functors (which we also denote by  $NI$ )

$$NI: \Delta^{\text{op}} \xrightarrow{\mathcal{C}} \mathcal{C}\text{at}^{\text{op}} \xrightarrow{\text{Hom}_{\mathcal{C}\text{at}}(-, I)} (\text{Sets}).$$

We first explain the notations involved. Here  $\mathcal{C}\text{at}$  is the category of all small categories. For every object  $[n]$  of  $\Delta$ , the functor  $\mathcal{C}$  sends  $[n]$  to a category  $\mathcal{C}_{[n]} := \mathcal{C}([n])$  defined as

- The set of objects is  $\text{Ob}(\mathcal{C}_{[n]}) = [n] = \{0, 1, \dots, n\}$ ;
- $\text{Hom}_{\mathcal{C}_{[n]}}(x, y) = \begin{cases} \{*\}, & \text{if } x \leq y \\ \emptyset, & \text{otherwise.} \end{cases}$  In other words, for  $x, y \in [n]$ , there is at most one morphism from  $x$  to  $y$  and the morphism exists if and only if  $x \leq y$ .

Every morphism  $\alpha: [n] \rightarrow [m]$  is sent by  $\mathcal{C}$  to a functor  $\mathcal{C}_{\alpha}$  from  $\mathcal{C}_{[m]}$  to  $\mathcal{C}_{[n]}$  defined as

- For every  $x \in [n]$ ,  $\mathcal{C}_{\alpha}(x) = \alpha(x)$
- A morphism  $x \rightarrow y$  in  $\mathcal{C}_{[n]}$  is sent to the unique morphism  $\alpha(x) \rightarrow \alpha(y)$ , which exists since  $\alpha$  is weakly monotone.

The second functor  $\text{Hom}_{\mathcal{C}\text{at}}(-, I)$  sends a small category  $\mathcal{D}$  to the set of all functors from  $\mathcal{D}$  to  $I$ , which is indeed a set since  $I$  is small. Its action on morphisms between small categories is natural.

Next we verify that this  $NI$  coincides with the  $NI$  defined in the problem for simplices by showing that giving a functor in  $\text{Hom}_{\mathcal{C}\text{at}}(\mathcal{C}_{[n]}, I)$  is equivalent to giving a composable  $n$ -tuple of morphisms in  $I$ . Note that every morphism in  $\mathcal{C}_{[n]}$  can be written as composition of morphisms of the form  $(j-1) \rightarrow j$  with  $j \in \{1, \dots, n\}$ . Thus giving images of all morphisms of  $\mathcal{C}_{[n]}$  is equivalent to only giving images of morphisms of form  $(j-1) \rightarrow j$ . Now If we have a functor  $F$  from  $\mathcal{C}_{[n]}$  to  $I$ , then we have for every  $x \in [n]$  an object  $F(x)$  of  $I$  and for every morphism  $(j-1) \rightarrow j$  in  $\mathcal{C}_{[n]}$  a morphism  $F_j: F(j-1) \rightarrow F(j)$  in  $I$ . Note that by definition the target of  $F_{j-1}$  equals the source of  $F_j$ , so  $(F_n, \dots, F_1)$  is a composable  $n$ -tuple of morphisms in  $I$ . Conversely, given a composable  $n$ -tuple of morphisms in  $I$ , say  $(F_n, \dots, F_1)$  as described in the problem, we may define  $F: \mathcal{C}_{[n]} \rightarrow I$  by

- $F(j) := \begin{cases} \text{source}(F_{j+1}), & \text{if } 0 \leq j < n \\ \text{target}(F_n), & \text{if } j = n. \end{cases}$

- $F((j-1) \rightarrow j) := F_j$ .

The functor is well-defined by composability of  $(F_n, \dots, F_1)$  and it clearly preserves identity and composition of morphisms. This finishes this verification. Moreover, the above argument shows that a functor in  $\text{Hom}_{\text{Cat}}(\mathcal{C}_{[n]}, I)$  is determined by the image of all morphisms  $(j-1) \rightarrow j$  under it.

Finally we verify that the action of  $d_i: [n-1] \rightarrow [n]$  on  $(NI)_n$  is as required. The case for  $s_i$  will be similar and thus omitted. First,  $d_i$  induces a functor  $\mathcal{C}_{d_i}: \mathcal{C}_{[n-1]} \rightarrow \mathcal{C}_{[n]}$ , whose action on objects and morphisms are explained in the first paragraph. This  $\mathcal{C}_{d_i}$  induces morphism of sets

$$\begin{aligned} \text{Hom}_{\text{Cat}}(\mathcal{C}_{[n]}, I) &\rightarrow \text{Hom}_{\text{Cat}}(\mathcal{C}_{[n-1]}, I) \\ \left( \begin{array}{c} \text{images of } (j-1 \rightarrow j) \\ 1 \leq j \leq n \end{array} \right) &\mapsto \left( \begin{array}{c} \text{images of } d_i(j-1) \rightarrow d_i(j) \\ 1 \leq j \leq n-1 \end{array} \right). \end{aligned}$$

Therefore, when  $j < i$ , the image of  $f_j$  is still  $f_j$ . When  $j = i$ , the image of  $f_j$  is  $f_{j+1} \circ f_j$  since  $d_i(i-1) \rightarrow d_i(i)$  is composition of  $(i-1) \rightarrow i$  and  $i \rightarrow (i+1)$ . When  $j > i$ , the image of  $f_j$  is  $f_{j+1}$ , as the morphism  $d_i(j-1) \rightarrow d_i(j)$  becomes  $j \rightarrow j+1$ . This finishes the verification and hence the proof.

**(b)** For every  $[n] \in \Delta$ , the definition of  $(NF)_n$  is already given in the problem. So we only need to verify that for every morphism  $\alpha: [n] \rightarrow [m]$ , the diagram

$$\begin{array}{ccc} (NI)_m & \xrightarrow{(NF)_m} & (NJ)_m \\ \downarrow \alpha^* & & \downarrow \alpha^* \\ (NI)_n & \xrightarrow{(NF)_n} & (NJ)_n \end{array}$$

commutes. For this, let  $(f_m, \dots, f_1) \in (NI)_m$ . By the equivalence

$$\{\text{elements in } (NI)_m\} \leftrightarrow \{\text{functors in } \text{Hom}_{\text{Cat}}(\mathcal{C}_{[m]}, I)\} \leftrightarrow \left\{ \begin{array}{l} \text{image of morphisms } (j-1) \rightarrow j \\ \text{under a functor in } \text{Hom}_{\text{Cat}}(\mathcal{C}_{[m]}, I) \end{array} \right\} \quad (1)$$

established in part (a), we may write  $f_j = G((j-1) \rightarrow j)$  for all  $j \in [m]$  and some functor  $G \in \text{Hom}_{\text{Cat}}(\mathcal{C}_{[m]}, I)$ . Then by definition of maps in the diagram the image of  $f_j$  under  $\alpha^* \circ (NF)_m$  and  $(NF)_n \circ \alpha^*$  are both  $F \circ G(\alpha(j-1) \rightarrow \alpha(j))$ .

**(c)** The inverse morphism  $\psi: NI \times NJ \rightarrow N(I \times J)$  can be constructed as

$$\begin{aligned} (NI \times NJ)_n &= (NI)_n \times (NJ)_n \xrightarrow{\psi_n} N(I \times J)_n \\ (f_n, \dots, f_1) \times (g_n, \dots, g_1) &\mapsto (f_n \times g_n, \dots, f_1 \times g_1) \end{aligned}$$

By writing elements of  $(NI)_n$  as images of morphisms in  $\mathcal{C}_{[n]}$  under some functor from  $\mathcal{C}_{[n]}$  to  $I$ , as we did in the previous parts, we can show that for every  $\alpha \in \text{Mor}_\Delta([n], [m])$  the diagram

$$\begin{array}{ccc} (NI)_m \times (NJ)_m & \xrightarrow{\psi_m} & N(I \times J)_m \\ \alpha^* \downarrow & & \downarrow \alpha^* \\ (NI)_n \times (NJ)_n & \xrightarrow{\psi_n} & N(I \times J)_n \end{array}$$

commutes. Thus  $\psi$  is indeed a morphism between simplicial sets. By computing directly the action on elements, we can show that both the compositions

$$\begin{aligned} (NI \times NJ)_n &= (NI)_n \times (NJ)_n \xrightarrow{\psi_n} N(I \times J)_n \xrightarrow{(N_{\text{proj}_I}, N_{\text{proj}_J})_n} (NI)_n \times (NJ)_n = (NI \times NJ)_n \\ \text{and } N(I \times J)_n &\xrightarrow{(N_{\text{proj}_I}, N_{\text{proj}_J})_n} (NI)_n \times (NJ)_n = (NI \times NJ)_n \xrightarrow{\psi_n} N(I \times J)_n \end{aligned}$$

are identities. So  $(N_{\text{proj}_I}, N_{\text{proj}_J})$  is an isomorphism between simplicial sets.

(d) We claim that there is an isomorphism  $\varphi: \Delta[1] \rightarrow N[1]$  between simplicial sets, which will be constructed later. If this holds, we have the following diagram (identify  $NI$  with  $NI \times \Delta[0]$ )

$$\begin{array}{ccccc} NI & \xrightarrow{1 \times d^1} & NI \times \Delta[1] & \xleftarrow{1 \times d^0} & NI \\ & & \downarrow \cong & & \\ & & NI \times N[1] & & \\ & & \downarrow \cong & & \\ & & N(I \times [1]) & & \\ & \swarrow NF & \downarrow NH & \searrow NG & \\ & & NJ & & \end{array}$$

where  $NH$  is the morphism constructed from  $H$  as in part (b). The commutativity of this diagram can be showed by routine verifications on simplices, using the definition of each maps and the equivalence (1). And the existence of such a commutative diagram is just the definition of a simplicial homotopy between  $NF$  and  $NG$ .

Now we show our claim by briefly explaining the construction of the isomorphism  $\varphi: \Delta[1] \rightarrow N[1]$ . For every  $n \in \mathbb{N}$ , define

$$\begin{aligned} \varphi_n: (\Delta[1])_n &\rightarrow (N[1])_n \\ (g: [n] \rightarrow [1]) &\mapsto (g_n, \dots, g_1), \text{ with } g_j = \begin{cases} \text{Id}_0, & \text{if } g(j) = g(j-1) = 0 \\ f: 0 \rightarrow 1, & \text{if } g(j) = 1, g(j-1) = 0 \\ \text{Id}_1, & \text{if } g(j) = g(j-1) = 1. \end{cases} \end{aligned}$$

For every  $\alpha: [n] \rightarrow [m]$ , The diagram

$$\begin{array}{ccc} (\Delta[1])_m & \xrightarrow{\alpha^*} & (\Delta[1])_n \\ \downarrow \varphi_m & & \downarrow \varphi_n \\ (N[1])_m & \xrightarrow{\alpha^*} & (N[1])_n \end{array}$$

commutes because if we choose some  $(g: [m] \rightarrow [1]) \in (\Delta[1])_m$ , then by computing with the definition of maps and the equivalence (1), we find that

$$\alpha^* \circ \varphi_m(g) = \varphi_n \circ \alpha^*(g) = (g_n, \dots, g_1)$$

with  $g_j = \begin{cases} \text{Id}_0, & \text{if } g \circ \alpha(j) = g \circ \alpha(j-1) = 0 \\ f: 0 \rightarrow 1, & \text{if } g \circ \alpha(j) = 1, g \circ \alpha(j-1) = 0 \\ \text{Id}_1, & \text{if } g \circ \alpha(j) = g \circ \alpha(j-1) = 1. \end{cases}$

So  $\varphi$  is a morphism. Now we describe its inverse  $\psi: N[1] \rightarrow \Delta[1]$ . Given some composable  $n$ -tuple  $(g_n, \dots, g_1) \in (N[1])_n$ , there are three possibilities (by composability)

- $g_n = \dots = g_1 = \text{Id}_0$
- There is one  $j$  such that  $g_n = \dots = g_{j+1} = \text{Id}_1$ ,  $g_j = f$  and  $g_{j-1} = \dots = g_1 = \text{Id}_0$
- $g_n = \dots = g_1 = \text{Id}_1$

In the first case, we let  $\psi_n(g_n, \dots, g_1)$  sends  $[n]$  to 0. In the second case,  $\psi_n(g_n, \dots, g_1)$  send  $1, \dots, j-1$  to 0 and  $j, \dots, n$  to 1. In the third case,  $\psi_n(g_n, \dots, g_1)$  sends  $[n]$  to 1. It is then straightforward to verify that  $\psi$  is indeed a morphism and that both  $(\psi \circ \varphi)_n$  and  $(\varphi \circ \psi)_n$  are identities for all  $n \in \mathbb{N}$ .

(e) Let  $F: I \rightarrow J$  and  $G: J \rightarrow I$  be functors such that there exist natural transformations  $\tau: GF \rightarrow \text{Id}_I$  and  $\tau': FG \rightarrow \text{Id}_J$ . The process in part (d) shows that  $\tau$  yields a simplicial homotopy between  $N(GF) = NG \circ NF$  and  $N(\text{Id}_I) = \text{Id}_{NI}$ . Similarly,  $\tau'$  yields a simplicial homotopy between  $N(FG) = NF \circ NG$  and  $N(\text{Id}_J) = \text{Id}_{NJ}$ . Thus  $NI$  and  $NJ$  are homotopy equivalent simplicial sets.