

Chapter 2. Sets

▶ 1. Sets

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▶ 2. Relations

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1.1 Set Definitions

- ▶ *Set* = a collection of distinct unordered objects, and represented as A,B,C. The members of a set are called *elements*
ex) $A=\{1,2,3,4\}$, Graph $G=\{V, E\}$ // set of vertices and edges
- ▶ How to define a set?
 - ▶ Listing: (원소 나열법) Example: $A = \{1,3,5,7\}$
 - ▶ Predicates: 조건 제시법 (for large set or infinite set)
 - ▶ Example: $B = \{x \mid x = 2k + 1, 0 \leq k \leq 3\}$
 $A = \{x \mid x \text{ is a positive, even integer}\}$
 $S = \{1,2,3\} \rightarrow \{x \mid x \text{ is a positive integer less than } 4\}$

Kinds of Sets

▶ *Finite* sets

▶ Ex: $A = \{1, 2, 3, 4\}$, $B = \{x \mid x \text{ is an integer, } 1 \leq x \leq 4\}$

□ *Infinite* sets

□ Ex: $Z = \{\text{integers}\} = \{\dots, -3, -2, -1, 0, 1, 2, 3, \dots\}$

$S = \{x \mid x \text{ is a real number and } 1 \leq x \leq 4\} = [0, 4]$

▶ The *empty* set (*null set*), \emptyset has no elements. $\emptyset = \{ \}$

□ *Universal* set: the set of all elements about which we make assertions. (must be given explicitly, or inferred from context)

Ex: $U = \{\text{all real numbers}\}$, $U = \{x \mid x \text{ is a natural number and } 1 \leq x \leq 10\}$



Cardinality

- ▶ Cardinality of a set A (in symbols $|A|$) is the number of elements in A , $|A| = n$ 으로 표기

(a measure of how many different elements of S has)

- ▶ Examples:

If $A = \{1, 2, 3\}$ then $|A| = 3$

If $B = \{x \mid x \text{ is a natural number and } 1 \leq x \leq 9\}$ then $|B| = 9$



Subsets

- ▶ X is a **subset** of Y if every element of X is also contained in Y (in symbols $X \subseteq Y$)

- ▶ X is a **proper subset** of Y if $X \subseteq Y$ but $Y \not\subseteq X$
 - ▶ Observation: \emptyset is a subset of every set
 - ▶ Ex) $X = \{1, 2\}$ $Y = \{1, 2, 3\}$

- **Equality**: $X = Y$ if $X \subseteq Y$ and $Y \subseteq X$
(two sets are Equal iff each is a SUBSET of the other)
Ex) $A = \{x \mid x^2 + x - 6 = 0\}$, $B = \{2, -3\}$ then $A = B$.



Power set

- ▶ The power set of X is the set of all subsets of X , in symbols $P(X)$, i.e. $P(X) = \{A \mid A \subseteq X\}$

Ex) if $X = \{1, 2, 3\}$,

then $P(X) = \{\emptyset, \{1\}, \{2\}, \{3\}, \{1,2\}, \{1,3\}, \{2,3\}, \{1,2,3\}\}$

Ex) given set S , $S = \{1,2\}$, $P(S) = \{\emptyset, \{1\}, \{2\}, \{1,2\}\}$

□ Theorem: If $|X| = n$, then $|P(X)| = 2^n$.



1.2 Set operations:

- ▶ The union of X and Y is defined as the set

$$X \cup Y = \{ x \mid x \in X \text{ or } x \in Y \}$$

- The intersection of X and Y is defined as the set

$$X \cap Y = \{ x \mid x \in X \text{ and } x \in Y \}$$

The difference of two sets $X - Y = \{ x \mid x \in X \text{ and } x \notin Y \}$

- Symmetric difference $X \odot Y = (X - Y) \cup (Y - X)$

Ex) $X = \{a, b, c, e, f\}$, $Y = \{b, d, r, s\}$ $X \odot Y = \{a, c, e, f, d, r, s\}$

- **COMPLEMENT** : every element that not in A .

$$\tilde{A} = \{x \mid x \in U \text{ and } x \notin A\} = U - A,$$

ex) $A = \{1, 3, 5\}$, $U = \{1, 2, 3, 4, 5\}$, $\tilde{A} = \{2, 4\}$

Cartesian Product

- ▶ $A \times B$: a set of all ordered pairs (a,b) , where
 $a \in A$ and $b \in B$
- ▶ **‘Ordered’** \Rightarrow a = first element, and b = second element
 $A \times B \neq B \times A$, unless $A = \emptyset$ or $B = \emptyset$ or $A=B$
(not commutative)
- ▶ ex) Let $A=\{1,2,3\}$, $B=\{a, b\}$. Find $A \times B$
sol) $A \times B = \{(1,a),(1,b),(2,a),(2,b),(3,a),(3,b)\}$



Partition

- Let S be a set, then a partition of S is a set

$$\Pi = \{A_1, A_2, \dots, A_i, \dots, A_k\}$$

- (1) each A_i is a **nonempty** subset of S , $i=1, \dots, k$,
(2) the A_i **cover** S , in that $S = A_1 \cup A_2 \cup \dots \cup A_k$
(3) the A_i are **MUTUALLY) DISJOINT** in that

$$\text{If } i \neq j, \text{ then } A_i \cap A_j = \emptyset$$

Ex) Let $S = S_9$, Let $A_1 = \{1, 5\}$, $A_2 = \{2\}$, $A_3 = \{7, 8, 9\}$, $A_4 = \{3, 4, 6\}$

show that $\Pi = \{A_1, A_2, A_3, A_4\}$ is a partition of S

1) each set is nonempty

2)

3)



Properties of set operations

Theorem 2.1.10: Let U be a universal set, and A , B and C subsets of U . The following properties hold:

- a) Associativity: $(A \cup B) \cup C = A \cup (B \cup C)$, $(A \cap B) \cap C = A \cap (B \cap C)$
- b) Commutativity: $A \cup B = B \cup A$, $A \cap B = B \cap A$
- c) Distributive laws: $A \cap (B \cup C) = (A \cap B) \cup (A \cap C)$, $A \cup (B \cap C) = (A \cup B) \cap (A \cup C)$
- d) Identity laws: $A \cap U = A$ $A \cup \emptyset = A$
- e) Complement laws: $A \cup A^c = U$ $A \cap A^c = \emptyset$
- f) Idempotent laws: $A \cup A = A$ $A \cap A = A$
- g) Bound laws: $A \cup U = U$ $A \cap \emptyset = \emptyset$
- h) Absorption laws: $A \cup (A \cap B) = A$ $A \cap (A \cup B) = A$
- i) Involution law: $(A^c)^c = A$
- j) De Morgan's laws for sets: $(A \cup B)^c = A^c \cap B^c$ $(A \cap B)^c = A^c \cup B^c$



1.3 Sequences and strings

- ▶ A *sequence* is an **ordered list** of numbers, usually defined according to a formula:

$$s_n = \text{a function of } n = 1, 2, 3, \dots$$

- ▶ If s is a sequence $\{s_n \mid n = 1, 2, 3, \dots\}$,
 - ▶ s_1 denotes the first element, s_2 the second element,...
 - ▶ s_n the n^{th} element...

- finite: stops after n steps - infinite : continues infinitely
- countable: finite sets
- uncountable : ex) real number between 0, 1



SEQUENCES & Subsequence

- ▶ **Increasing sequence**, if $S_n \leq S_{n+1}$, $\forall n$, when $n=1,2,3..$

Ex) $S_n = 2n-1$, is increasing: $S_n = 1, 3, 5, \dots$

- ▶ **Decreasing sequence**, if $S_n \geq S_{n+1}$, $\forall n$, when $n=1,2,3..$

Ex) $S_n = 4 - 2n$, is decreasing:, $2, 0, -2, -4, -6, \dots$

ex) If we let S denote this sequence, $s_1 = 2, s_2 = 4, \dots s_n = 2n$

\Rightarrow { **infinite sequence, increasing sequence** }

- ▶ A **subsequence** of a sequences $S = \{S_n\}$ is a sequence $T = \{T_n\}$ that consists of certain elements of s retained in the original order they had in S .

Ex) let $S = \{S_n = n \mid n = 1, 2, 3, \dots\}$ $1, 2, 3, 4, 5, 6, 7, 8, \dots$

Let $T = \{T_n = 2n \mid n = 1, 2, 3, \dots\}$ $2, 4, 6, 8, 10, 12, 14, 16, \dots$ **T is a subsequence of S**



Sigma & Pi notation

If $\{a_n\}$ is a sequence, then the sum

$$\sum_{k=1}^m a_k = a_1 + a_2 + \dots + a_m$$

This is called the “**sigma notation**”, where the Greek letter Σ indicates a sum of terms from the sequence

If $\{a_n\}$ is a sequence, then the product

$$\prod_{k=1}^m a_k = a_1 a_2 \dots a_m$$

This is called the “**pi notation**”, where the Greek letter Π indicates a product of terms of the sequence



Strings (finite sequence)

- ▶ Let X be a nonempty set. A *string over X* is a finite sequence of elements from X .

Ex) if $X = \{a, b, c\}$ Then $\alpha = \text{bbaccc}$ is a string over X

Notation: $\text{bbaccc} = b^2ac^3$

The **length** of a string α is the number of elements of α and is denoted by $|\alpha|$. If $\alpha = b^2ac^3$ then $|\alpha| = 6$.

- ▶ order is important $\text{baac} \neq \text{abbc}$
- ▶ The *null string* is the string with no elements and is denoted by the Greek letter λ (lambda). It has length zero.



More on strings

- ▶ Let $X^* = \{\text{all strings over } X \text{ including } \lambda\}$
 - ▶ Let $X^+ = X^* - \{\lambda\}$, the set of all non-null strings
ex) $X = \{a,b\}$ $X^* \Rightarrow \lambda, a, b, abab, b^20a^{50}ba, \dots$
 - ▶ **Concatenation** of two strings α and β is the operation on strings consisting of writing α followed by β to produce a new string $\alpha\beta$
 - ▶ Ex: $\alpha = bbaccc$ and $\beta = caaba$,
then $\alpha\beta = bbaccccaaba = b^2ac^4a^2ba$
Clearly, $|\alpha\beta| = |\alpha| + |\beta|$
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2. Relations

2.1 Definitions

- ▶ What is a relation?

ex) comparison between two objects (bigger, better, faster,...)

- ▶ relationship between objects belong to same set or different set
==> result is ORDERED PAIRS (순서쌍)

ex) $S = \{\text{set of students}\}$, $C = \{\text{set of courses}\}$ ==>
ordered pairs (s,c) from $S \times C$.

For each student there will be some pairs



Binary Relation

- ▶ Binary relation is relationship between elements of two sets using ordered pairs
- Let A, B be sets A binary relation from A to B is subsets R of $A \times B$

ex) Let $A = \{1, 2, 3\}$ $B = \{a, b\}$ Define some relations?

(답) $R1 = \{(1, a), (2, b), (3, a), (1, b)\}$

$R2 = \{(3, b)\}, \quad R3 = A \times B \text{ 전부}, \quad R4 = \emptyset$

- ▶ A relation R from a set X to a set Y is a subset of **cartesian product** $X \times Y$.
- ▶ If $(x, y) \in R$, we write **xRy** and say that **X IS RELATED TO Y** .



More binary relations

► Domain and Range

- * The set $\{x \in X \mid (x,y) \in R \text{ for some } y \in Y\} \Rightarrow \text{DOMAIN of } R$
- * The set $\{y \in Y \mid (x,y) \in R \text{ for some } x \in X\} \Rightarrow \text{RANGE of } R$

ex) Let $A = \{1,2,3\}$. R be the relation on A consisting of ordered pairs (a,b) . such that $a \geq b$. List elements of R ($A \times A$)

Ex) Let $X = \{2,3,4\}$ and $Y = \{3,4,5,6,7\}$ If we define a relation R from X to Y by “ $(x,y) \in R$ if x divides y (with zero remainder)”, list elements of R .

Ex) Set arising from relations

$$A = \{1,2,3\} \quad B = \{r,s\}, \quad R = \{(1,r)(2,s)(3,r)\}$$

► (sol) $\text{Dom}(R) = A \quad \text{Ran}(R) = B$

Relation Representation

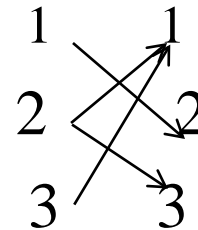
1) English style representation: ex) {x is greater than y}

2) List **ordered pairs** (순서쌍) ex) $A = \{1, 2, 3\}$, $A \times A$,

$$R = \{(1, 2)(2, 1)(2, 3)(3, 1)\}$$

3) Graph 표현

4) Matrix



Ex) Matrix representation

from \ To	c1	c2	c3
c1	0	140	100
c2	190	0	200
c3	110	180	0

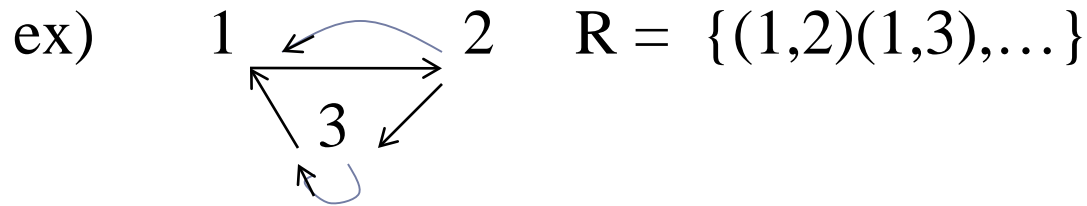
Relation R on set of cities

$$A = \{c1, c2, c3\}$$

Q) Find ' $C_i R C_j$ ' iff cost of going from C_i to C_j is less than \$150

Graph representation (Directed Graph)

- ▶ Graph representation 은 two copies of set 을 표현
- ▶ Directed graph 는 one copy of set 으로 표현



- ▶ Degree (indegree : 들어오는 갯수 (내차수)
(outdegree : 나가는 갯수 (외차수)
정점 3 의 차수는=> (내차수 2, 외차수 2)

ex) Draw the directed graph representation of the relation “ \leq ” on set $S_3 \Rightarrow S_3$ means (1,2,3)

(sol) $R =$

Matrix Representation (1)

► Let $A = \{a_1, a_2, \dots, a_n\}$ $B = \{b_1, b_2, \dots, b_m\}$, R be a relation between A and B

► Define $n \times m$ matrix M by $M(i,j) = \begin{cases} \text{false, if } (a_i, b_j) \notin R \\ \text{true, if } (a_i, b_j) \in R \end{cases}$

ex) $R = \{(1,2)(2,1)(2,3)(3,1)(3,3)\}$

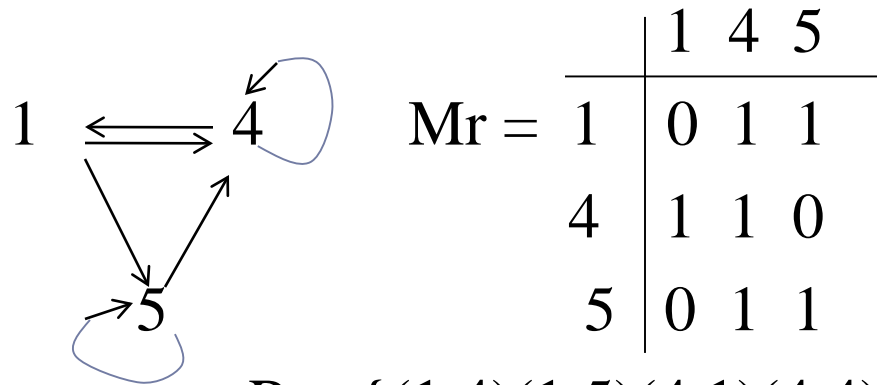
$$M = \begin{bmatrix} F & T & F \\ T & F & T \\ T & F & T \end{bmatrix} \text{ or } \begin{bmatrix} 0 & 1 & 0 \\ 1 & 0 & 1 \\ 1 & 0 & 1 \end{bmatrix} \qquad M = \begin{matrix} & F & T & F \\ T & F & T & T \\ T & F & T & T \end{matrix}$$

ex) Describe the relations corresponding to the matrix given above, using directed graph.



Matrix Representation (2)

ex) Let $A = \{1, 4, 5\}$ and graph is given below. Find M_r and R



$$M_r = \begin{array}{c|ccc} & 1 & 4 & 5 \\ \hline 1 & 0 & 1 & 1 \\ 4 & 1 & 1 & 0 \\ 5 & 0 & 1 & 1 \end{array}$$

$$R = \{(1.4)(1.5)(4.1)(4.4)(5.4)(5.5)\}$$

ex) Let $A = \{a, b, c, d\}$. Let R be a $M_r =$

Construct digraph of R , and list
in-degree and out-degree of all
vertices

1	0	0	0
0	1	0	0
1	1	1	0
0	1	0	1

Path

- ▶ R be a relation on A . a 에서 b 까지의 길이가 n 개인 path 는 a 부터 시작, b 에서 끝나는 유한 sequence (수열),
 $\Pi = a, x_1, x_2, \dots, x_{n-1}, b$ 를 말하는데, 각 원소들은 aRx_1 ,
 $x_1Rx_2, \dots, x_{n-1}Rb$ 이어야 한다.

(def)

$xR^n y$ = if n is fixed positive integer, x 에서 시작 y 에서 끝나는 길이 n 인 통로가 있다는 의미

$xR^\infty y$ = x 에서 시작, y 에서 끝나는 어떤 통로가 있다는 의미



2.2 Properties of Relations

- 1) Reflexive relation, Irreflexive relation
- 2) Symmetric relation, Anti-symmetric relation
- 3) Transitive relation

1) Reflexive Relation (반사관계)

정의) A relation R on a set X is called **REFLEXIVE**

if $(x,x) \in R$, for every $x \in X$.

$(\forall x \in X, xRx \Rightarrow \text{reflexive})$

* if for all $x \in X$, $x \not R x$ 이면, **irreflexive relation**



Reflexive relations

Ex) $X = \{1,2,3,4\}$

$R = \{(1,1)(1,2)(1,3)(1,4)(2,2)(2,3)(2,4)(3,3)(3,4)(4,4)\}$

$\text{dom}(R) = X \quad \text{ran}(R) = X$

(sol)

- R is reflexive, since each element $x \in X$, $(x,x) \in R$

Specifically, $(1,1)(2,2)(3,3)(4,4)$ are in R

- in Matrix representation, must have 1's on main diagonal

- or, in Digraph, each vertex has a LOOP.

► Ex) $A = \{1,2,3\} \quad R = \{(1,1)(1,2)\}$



Symmetric relation (대칭관계)

def: A relation R on set X is symmetric, if $\forall x, y \in X$,
if $(x, y) \in R$, then $(y, x) \in R$

$$\rightarrow R = R^{-1}$$

\rightarrow an element is related to second element iff second element is related to first element.

$$V \iff W$$

▶ if for all $x, y \in X$, $xRy = yRx$,

▶ matrix representation

$$\begin{matrix} 0 & 1 \end{matrix}$$

$$\begin{matrix} 1 & 0 \end{matrix}$$

ex) $X = \{a, b, c, d\}$ $R = \{(a, a)(b, c)(c, b)(d, d)\}$

ex) $X = \{1, 2, 3, 4\}$ $R = \{(1, 1)(1, 2)(1, 3)(1, 4)(2, 2)(2, 3)(2, 4)(3, 3)\}$



Anti-symmetric relation

Def: A relation R on a set X is called anti-symmetric,

- 1) if $\forall x, y \in X$, if $(x, y) \in R$ and $x \neq y$, then $(y, x) \notin R$.
- 2) if $(x, y) \in R$ & $(y, x) \in R$ and $x = y$, then anti-symmetric

- Property: The diagram of an anti-symmetric relation has a property that between any two vertices there is at most one directed edge.

ex) $X = \{a, b, c, d\}$ $R = \{(a, a)(b, c)(c, b)(d, d)\}$

(sol) R is not antisymmetric since there is (b, c) & (c, b)

ex) $X = \{1, 2, 3, 4\}$ $R = \{(1, 1)(1, 2)(1, 3)(1, 4)(2, 2)(2, 3)(2, 4)(3, 3)(3, 4)(4, 4)\}$



Con't

ex) Let $A = \mathbb{Z}$, $R = \{(a,b) \in A \times A \mid a < b\}$

(sol) If $a < b$, then it's not true for $b < a$, so R is not symmetric

If $a \neq b$ then either $a < b$ or $b < a$, so R is antisymmetric

ex) Let $A = \{1,2,3,4\}$ and Let $R = \{(1,2)(2,2)(3,4)(4,1)\}$

ex) $R = \{(a, a)(b, b)(c, c)\}$

ex) $M1 = \begin{matrix} 1 & 0 & 1 \\ 0 & 0 & 1 \\ 1 & 1 & 1 \end{matrix}$ $M2 = \begin{matrix} 1 & 0 & 0 & 1 \\ 0 & 1 & 1 & 1 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{matrix}$



Transitive relation (추이관계)

► Def: A relation R on a set X is called TRANSITIVE

If $\forall x, y, z \in X$, if (x, y) and $(y, z) \in R$, then $(x, z) \in R$.

If $M_{ij}=1, M_{jk}=1$, then $M_{ik}=1$

ex) $A = \mathbb{Z}$, $R = \{(a, b) \in A \times A \mid a < b\}$

(sol) aRb and bRc , $a < b \rightarrow b < c$, therefore, $a < c \therefore aRc$

therefore, R is transitive

ex) $X = \{a, b, c, d\}$ $R = \{(a, a)(b, c)(c, b)(d, d)\}$

ex) $A = \{1, 2, 3, 4\}$ $R = \{(1, 1)(2, 2)(2, 3)(3, 2)(4, 2)(4, 4)\}$

ex) $A = \{0, 1, 2\}$ $R = \{(0, 0)(0, 1)(0, 2)(1, 2)\}$



exercises

ex) $X = \{1,2,3,4\}$

$R = \{(1,1)(1,2)(1,3)(1,4)(2,2)(2,3)(2,4)(3,3)(3,4)(4,4)\}$

(sol) R is transitive

If $x=y$, then it is transitive since
$$\frac{(x,y) (y,z) (x,z)}{(1,1) (1,1) (1,1)}$$

We just need to check for

$(1,2)(2,3)(1,3),$

$(1,2)(2,4)(1,4),$

$(1,3)(3,4)(1,4),$

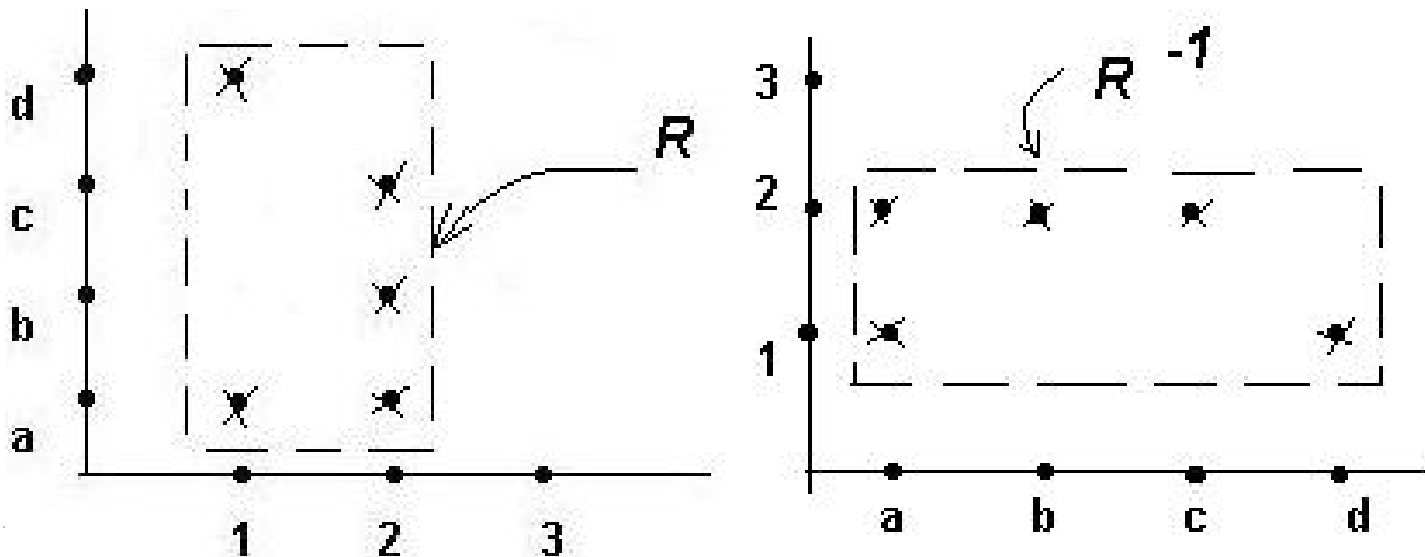
$(2,3)(3,4)(2,4)$



2.3 기타속성 - Inverse of a relation

Given a relation R from X to Y , its inverse R^{-1} is the relation from Y to X defined by $R^{-1} = \{ (y,x) \mid (x,y) \in R \}$

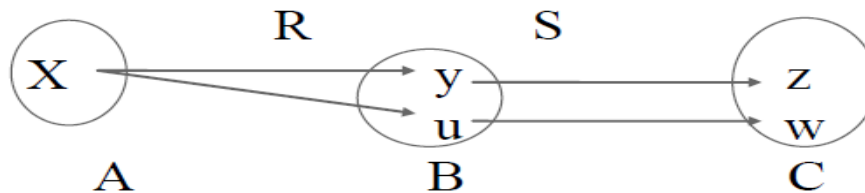
Ex) if $R = \{(1,a), (1,d), (2,a), (2,b), (2,c)\}$ then
 $R^{-1} = \{(a,1), (d,1), (a,2), (b,2), (c,2)\}$



Composition of Relation

- ▶ Create new relations from existing ones. Let $R: A \times B$, $S: B \times C$
- ▶ Composition of relation R and S is the relation between A and C .

$$S.R = \{(x,z) \mid x \in A, z \in C \text{ and } \exists y \in B, \text{ such that } xRy \text{ and } ySz\}$$



ex) $A = \{a,b,c\}$ $B = \{1,2\}$ $C = \{5,9,10\}$

$S (B \times C) = \{(1,5)(1,9)(2,9)\}$ $R(A \times B) = \{(a,1)(c,2)\}$ $S.R?$

* Composition with itself

ex) R on set $A = \{1,2,3\}$ $R = \{(1,1)(1,2)(1,3)(3,2)\}$ $R.R?$

2.4 Equivalence relations

- ▶ THM1) Let Q be a partition of a set X . Define xRy to mean that for some set S in Q , both x and y belong to S . Then R is reflexive, symmetric, transitive.
- ▶ Let X be a set and R a relation on X , R is an *equivalence relation* on $X \Leftrightarrow R$ is **reflexive, symmetric and transitive**.

Ex) Consider the partition $Q = \{ \{1,3,5\}, \{2,6\}, \{4\} \}$ of $X = \{1,2,3,4,5,6\}$.

The complete relation R is $= \{ (1,1)(1,3)(1,5) (3,1)(3,3)(3,5) \\ (5,1)(5,3)(5,5) (2,2)(2,6)(6,2)(6,6) (4,4) \}$



Con't

ex) Let $X = S_4$, $R = \{(1,1)(1,2)(1,3)(1,4)(2,2)(2,3)(2,4)(3,3)(3,4)(4,4)\}$

The diagraph of R is reflexive ? symmetric ? transitive ?

ex) Let $X = S_5$, $R = \{(1,1)(1,3)(1,5)(2,2)(2,4)(3,1)(3,3)(3,5)(4,2)(4,4)(5,1)(5,3)(5,5)\} \Rightarrow \text{E.R}$

ex) Let $A = \mathbb{Z}$, $R = \{(a,b) \in A \times A \mid a \leq b\}$

(sol) 1) since, $a \leq a$, reflexive. 2) if $a \leq b$, $b \leq a$, not SYM

3) since $a \leq b$ and $b \leq c$, transitive, NOT E.R.

► **Def: Equivalence Relation leads to Natural partitions of A into groups of like objects.**



Set of equivalence classes

def1: Let R be an equivalence relation on X . $\forall a \in X$,
let $[a] = \{x \in X \mid xRa\}$ then
 $S = \{[a] \mid a \in X\}$ is a partition of X .

- ▶ Let R be a equivalence relation on set X . The sets $[a]$ defined in def1 are called the equivalence Class of X given by relation R .

Ex) Let $R = \{(1,1)(1,3)(1,5)(3,1)(3,3)(3,5)(5,1)(5,3)(5,5)(2,2)(2,6)(6,2)(6,6)(4,4)\}$

The equivalence class $[1]$ containing 1 consists of all x , such that $(x,1) \in R$. Therefore, $[1] = \{1,3,5\}$

The remaining E.C are found similarly, $[3] = [5] = \{1,3,5\}$,

$[2] = [6] = \{2,6\}$ $[4] = \{4\}$

Exercises

ex) Let $S=S_6$. For $x, y \in S$. Let xRy if $x-y$ is divisible by 2.

- Show that R is an E.R. - Find E.C. relative to R .

(sol)

1) R is reflexive

2) R is SYM \Rightarrow since $x-y / 2 \text{ 이면 } \Leftrightarrow y-x / 2 \text{ 이기 때문에}$

3) Suppose, xRy and yRz , then there exist k and m , such that

$$x-y = 2k \text{ and } y-z = 2m. \quad // \text{ since it is divisible by 2}$$

$$x-z = (x-y) + (y-z) = 2k+2m=2(k+m). \text{ so } R \text{ is TRAN}$$

4) by 1)2)3), it is E.R on S .

5) E.C: $[1] = \{x \in S_6: x-1 \text{ is divisible by 2}\} = \{1,3,5\}$

$$[2] = \{x \in S_6: x-2 \text{ is divisible by 2}\} = \{2,4,6\}$$

Which leads to a natural partition of set S_6 into $[1]$ and $[2]$, and $[1] \cup [2] = S_6$.



More examples for E.C.

ex) Let $X = \{1, 2, \dots, 10\}$. Define xRy to mean that 3 divides $x-y$.

ex) If $R = \{(1,1)(1,3)(1,5)(2,2)(2,4)(3,1)(3,3)(3,5)(4,2) (4,4)(5,1)$
 $(5,3) (5,5)\}$

ex) If $R = \{(a,a)(b,b)(c,c)\}$



Partition and E.C.

▶ We need to show to be E.C.

- 1) For all $x \in S$, $[x] \neq \emptyset$
- 2) if $[x] \neq [y]$, then $[x] \cap [y] = \emptyset$
- 3) $\cup \{[x] \mid x \in S\} = S$

ex) When $R = \{(1,1)(1,2)(2,1)(2,2)(3,3)(3,4)(4,3)(4,4)\}$, Find E.C.

(sol) $[1] = \{1,2\}$, $[3] = \{3,4\}$ and $[1] \cap [3] = \emptyset$ 이므로,
 $[1]$ 과 $[3]$ 은 partition 을 이룬다. Partition $A = \{ \{1,2\}, \{3,4\} \}$

ex) Let Set $A = \{1,2,3,4\}$. Partition of A, $P = \{ \{1,2,3\} \{4\} \}$ 일때,
P 에 의해 결정되는 Set A 의 Equivalence Relation R 을 구하라.

ex) $S = \{a,b,c,d,e,f\}$ $P = \{ \{a,b\}, \{c,e,f\}, \{d\} \}$.

▶ Find Equivalence Relation R.

2.5 Closure

Relation은 반사적, 대칭적, 추이적 성질을 가지고 있지 않을 수도 있다.

- Relation 이 특정 성질을 가지고 있지 않으면, 그 성질을 갖아질 때까지 Relation 에 관련된 순서쌍들을 첨가할 수 있다
- 우리가 원하는 성질을 가지며, 관계 R 을 포함하는, 최소의 관계 R_1 을 찾는 것이 필요하다 => CLOSURE

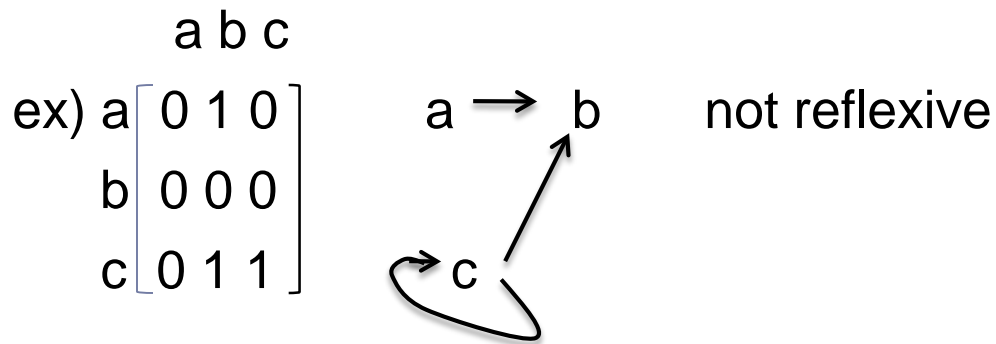
- ▶ 1) reflexive closure
- 2) symmetric closure
- 3) transitive closure



Reflexive closure

$R1 = R \cup \nabla$: **smallest** reflexive relation on A containing R.

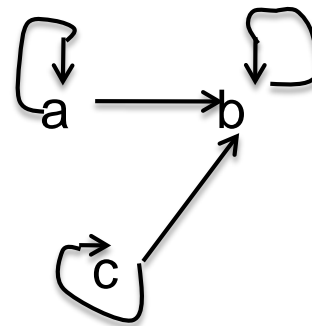
Reflexive closure of R is $R \cup \nabla$



(sol)

$$\text{So, } \begin{array}{c|ccc} & a & b & c \\ \hline a & 0 & 1 & 0 \\ b & 0 & 0 & 0 \\ c & 0 & 1 & 1 \end{array} \cup \begin{array}{c|ccc} & a & b & c \\ \hline a & 1 & 0 & 0 \\ b & 0 & 1 & 0 \\ c & 0 & 0 & 1 \end{array} = \begin{array}{c|ccc} & a & b & c \\ \hline a & 1 & 1 & 0 \\ b & 0 & 1 & 0 \\ c & 0 & 1 & 1 \end{array}$$

$R \qquad \nabla \qquad R1$



(identity matrix)

Symmetric Closure

- ▶ If R is relation on A which is not symmetric.

Then $\exists (x,y) \in R$, such that $(y,x) \notin R^{-1}$. So, if R is not symmetric, then we must add all pairs from R^{-1} .

→ Enlarge R to $R \cup R^{-1}$,

→ $R \cup R^{-1}$ is smallest symmetric Closure containing R .

- ▶ ex) $A = \{a,b,c,d\}$. $R = \{(a,b)(b,c)(a,c)(c,d)\}$

Then $R^{-1} = \{(b,a)(c,b)(c,a)(d,c)\}$

so symmetric closure of R is

$$R \cup R^{-1} = \{(a,b)(b,c)(a,c)(c,d)(b,a)(c,b)(c,a)(d,c)\}$$



Transitive Closure

- Let R be a relation on A . Then R^* is transitive closure of R .

ex) Let R be the relation $R = \{(a,b)(b,c)(c,a)\}$ on set $S = \{a, b, c\}$.

Find Transitive closure of R

(sol) 1) $(a,b)(b,c)(c,a) \in R^*$, since they are in R

2) Since $aRb \ \& \ bRc \Rightarrow aR^*c$ similarly, $bRc \ \& \ cRa \Rightarrow bR^*a$,
 $cRa \ \& \ aRb \Rightarrow cR^*b$

Transitive closure $R^* = \{(a,b)(b,c)(c,a)(a,c)(b,a)(c,b)\}$

ex) $A = \{1 \ 2 \ 3 \ 4\}$ $R = \{(1 \ 2)(2 \ 3)(3 \ 4)(2 \ 1)\}$ Find Transitive closure of R

(sol) $R^* = R^1 \cup R^2 \cup R^3 \dots \cup R^n$

$R^* = \{(1 \ 1)(1 \ 2)(1 \ 3)(1 \ 4)(2 \ 1)(2 \ 2)(2 \ 3)(2 \ 4)(3 \ 4)\}$



(sol) matrix

$$Mr = \begin{pmatrix} 0 & 1 & 0 & 0 \\ 1 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{pmatrix}$$

boolean multiplication

$$(Mr) \bullet^2 = \begin{pmatrix} 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix}$$

$$(Mr) \bullet^3 = \begin{pmatrix} 0 & 1 & 0 & 1 \\ 1 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix} \dots$$

$$(Mr) \bullet^n = \begin{cases} Mr & \text{when } n=1, \\ Mr^2 & \text{when } n=\text{even}, \\ Mr^3 & \text{when } n=\text{odd} \end{cases}$$

$$\text{Therefore } Mr \bullet^* = Mr \cup Mr^2 \cup Mr^3 = \begin{pmatrix} 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{pmatrix}$$

Warshall's algorithm

- ▶ Graphic, Matrix: impractical for large sets, not systematic, inefficient
- ▶ Use WARSHALL's algorithm (find transitive closure for $n \times n$ matrix)

(Warshall's algorithm) : FIND W_k from W_{k-1}

- 1) Move all 1's in W_{k-1} into W_k
- 2) W_{k-1} 의 열(column) k 에서 값이 1인곳의 위치에 p_1, p_2, \dots 를 기록하고
 W_{k-1} 의 행(row) k 에서 값이 1인곳의 위치에 q_1, q_2, \dots 를 기록한다.
- 3) W_k 의 p_i, q_j 의 위치에 1을 쓴다.

Ex) 위의 예제 이용, $W_0 = M_r = \begin{bmatrix} 0 & 1 & 0 & 0 \\ 1 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{bmatrix}$ 이고 $n=4$ 이다

1 0 1 0

0 0 0 1

0 0 0 0



2.6 관계의 응용 - Relational databases

- ▶ A **binary** relation R is a relation among two sets X and Y , already defined as $R \subseteq X \times Y$.
 - ▶ An **n -ary relation R** is a relation among n sets X_1, X_2, \dots, X_n , (i.e. a subset of the Cartesian product, $R \subseteq X_1 \times X_2 \times \dots \times X_n$.)
 - ▶ A **database** is a collection of records that are manipulated by a computer. They can be considered as n sets $X_1 \dots X_n$, each of which contains a list of items with information.
 - ▶ **Database management systems (DBMS)** are programs that help access and manipulate information stored in databases
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Operators Examples

- ▶ The *selection* $T1 := \text{PLAYER}[\text{Position} = \text{Guard}]$
- ▶ The *projection* $T2 := \text{PLAYER}[\text{Name}, \text{ID}]$
- ▶ The *join* $T3 := \text{PLAYER}[\text{ID} = \text{PID}] \text{ ASSIGNMENT}$

PLAYERS

ID	Name	Position	Age
23	M.Jordan	Guard	40
32	S. O'Neal	Center	33
3	A.Iverson	Guard	29
8	K.Bryant	Guard	26
33	G. Hill	Forward	32

ASSIGNMENT

PID	Team
32	Heats
3	76ers
8	Lakers



Operators Examples

T1:=PLAYER[Position = Guard]

T2:=PLAYER[Name,ID]

T3:=PLAYER[ID = PID] ASSIGNMENT

T1

ID	Name	Position	Age
23	M.Jordan	Guard	40
3	A.Iverson	Guard	29
8	K.Bryant	Guard	26

T2

ID	Name
23	M.Jordan
32	S. O'Neal
3	A.Iverson
8	K.Bryant
33	G. Hill

T3

ID	Name	Position	Age	Team
32	S. O'Neal	Center	33	Heats
3	A.Iverson	Guard	29	76ers
8	K.Bryant	Guard	26	Lakers

