

# Demand Side Management for a Campus Infrastructure

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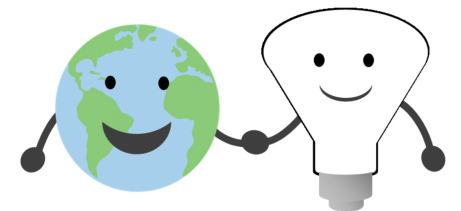
Batch : A6 | Place : In-House

# Abstract

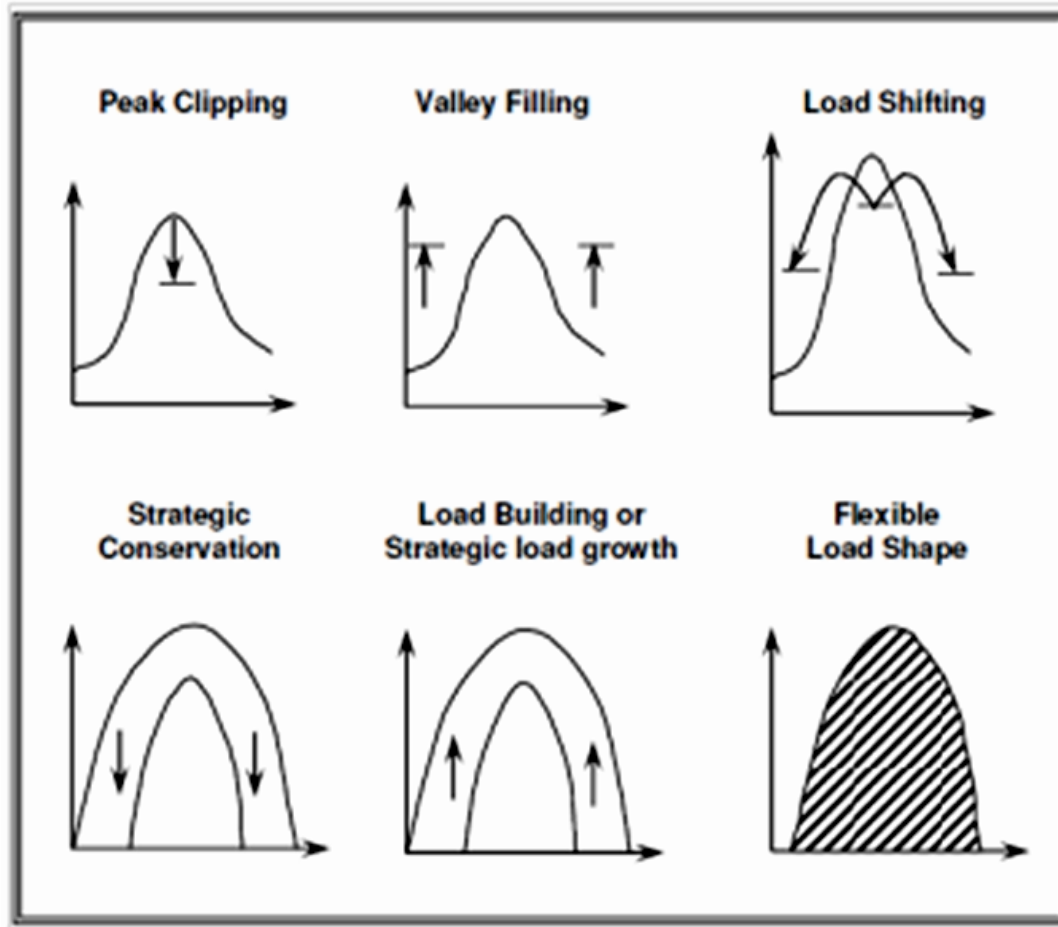
- We propose a novel optimization algorithm for handling loads using **Demand Side Management (DSM)**. DSM consists of planning, implementing, and monitoring activities of electrical utilities which encourage consumers to modify their level and pattern of electricity usage. Fortunately, DSM focuses only on cost reduction by flattening the load curve, and energy management is not considered. Introducing Fuzzy controlled energy management technique along with existing DSM strategies results in increased sustainability.
- Simulations are carried out for a **Campus infrastructure**, utilizing the **existing photo-voltaic array with battery storage**. These existing assets are treated as local **DC grid**, and plays a vital role in managing the overall cost. In addition to this, DSM also **encourage consumers to lower energy consumption** which leads to a significant drop in emission of harmful gases into the atmosphere thereby helping curtail the global warming process. The bulk proportion of **CO2 emissions** is from fossil fuel based power plants. Reducing environmental degradation and saving our future generation from this huge complication, reduction of CO2 footprint has also been accounted in our objective function.

# Objectives

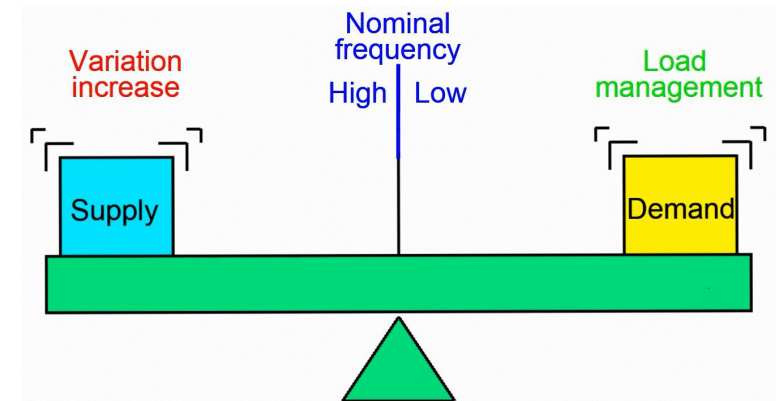
- The proposed project is an alternative to traditional load shedding, fall in system frequency, overall cost effectiveness, and also features reduction of Carbon Footprint. These are discussed in detail below.
  - **Load Shedding** - *The proposed project is an alternative to traditional load Shedding*
  - **Frequency Drop** - *Sudden fall of system frequency when demand on a power system is greater than the generation.*
  - **Cost Effectiveness** - *Develop efficient utilization strategy for electricity, and encourages consumers to help flattening their load curve.*
  - **Carbon Footprint** – *Reducing the consumption, inturn reduces the Co2 Emission's for Fossil Fuel based power generation units.*



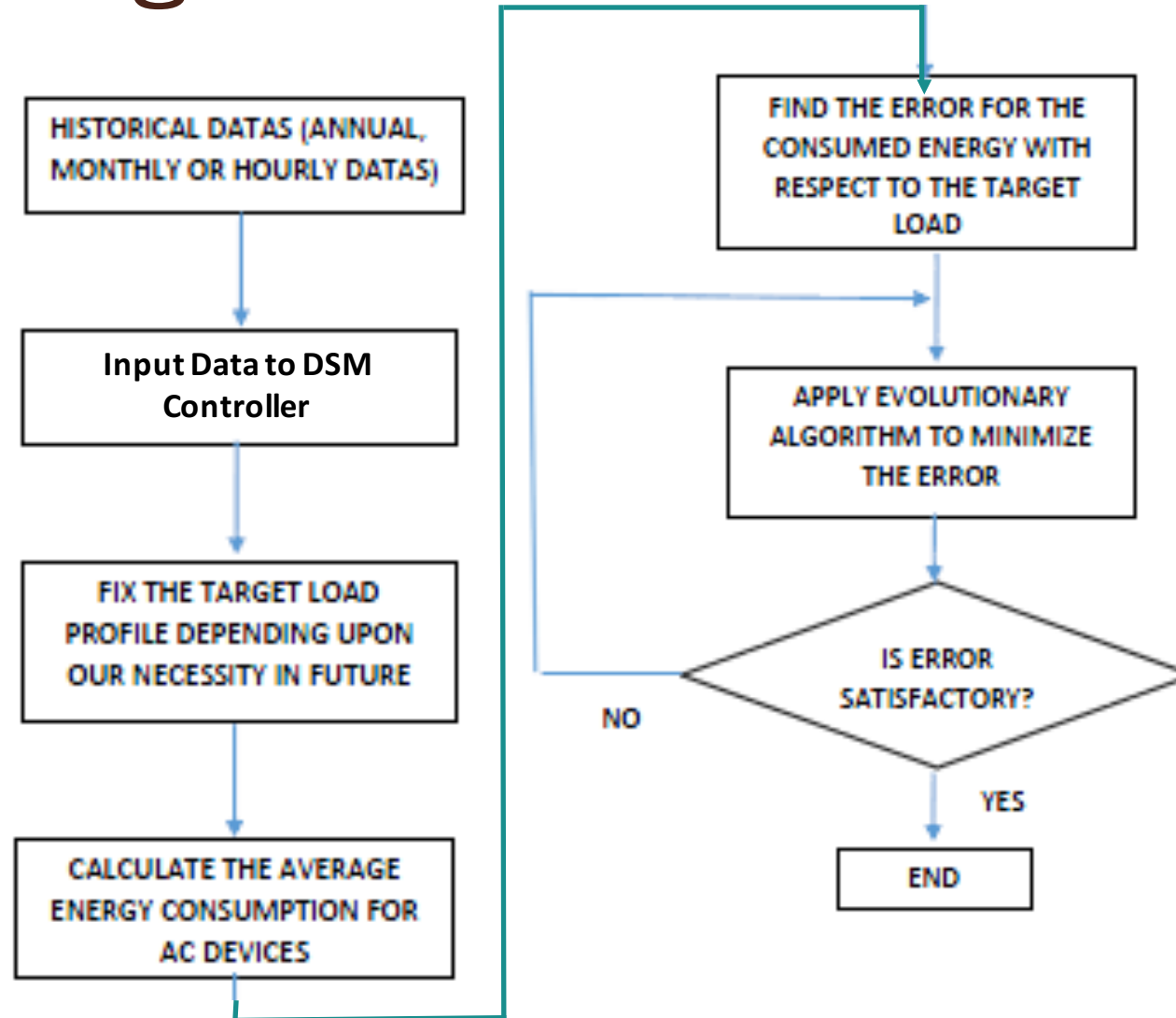
# DSM Strategy & Approach



- Adopted Strategy – Load Shifting
- Solution through Evolutionary Algorithm – GA



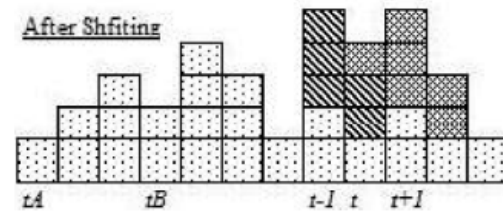
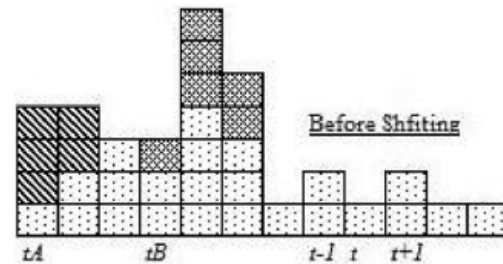
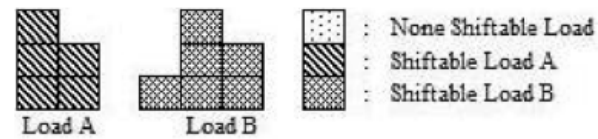
# Block Diagram



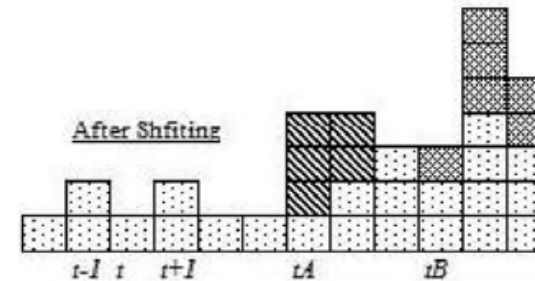
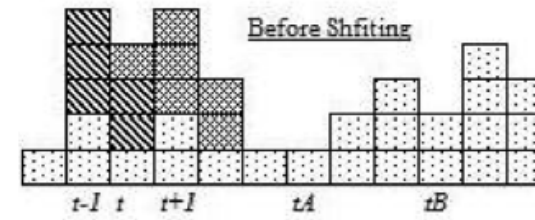
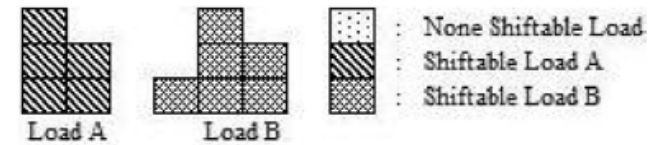
# The Current Progress

- Treating various stages of the issues, we are addressing these through Matlab Simulation Environment
- Stages completed
  - *Problem Formulation – DMS*
  - *Genetic Algorithm Tool – Single Objective*

# Problem Formulation

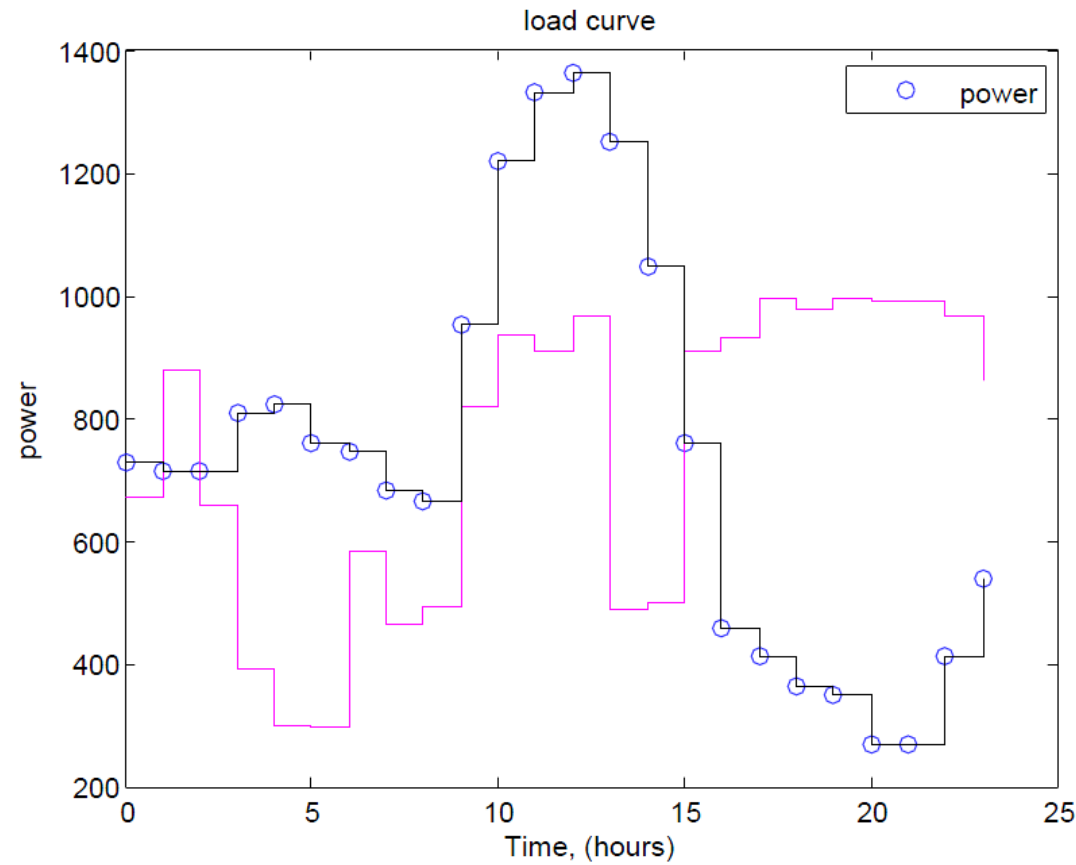


Disconnecting Loads



Connecting Loads

# Load Profile - Assumption



Time	Wholesale Price (ct/kWh)	Residential Microgrid
8hrs-9hrs	12.00	729.4
9hrs-10hrs	9.19	713.5
10hrs-11hrs	12.27	713.5
11hrs-12hrs	20.69	808.7
12hrs-13hrs	26.82	824.5
13hrs-14hrs	27.35	761.1
14hrs-15hrs	13.81	745.2
15hrs-16hrs	17.31	681.8
16hrs-17hrs	16.42	666.0
17hrs-18hrs	9.83	951.4
18hrs-19hrs	8.63	1220.9
19hrs-20hrs	8.87	1331.9
20hrs-21hrs	8.35	1363.6
21hrs-22hrs	16.44	1252.6
22hrs-23hrs	16.19	1046.5
23hrs-24hrs	8.87	761.1
24hrs-1hrs	8.65	475.7
1hrs-2hrs	8.11	412.3
2hrs-3hrs	8.25	364.7
3hrs-4hrs	8.10	348.8
4hrs-5hrs	8.14	269.6
5hrs-6hrs	8.13	269.6
6hrs-7hrs	8.34	412.3
7hrs-8hrs	9.35	539.1



*Proposed load shifting technique is mathematically formulated as follows.*

■ 
$$\sum_{t=1}^N (PLoad(t) - Objective(t))^2$$

- Where  $Objective(t)$  is the value of the objective curve at time  $t$
- $PLoad(t)$  is the actual consumption at time  $t$ .

■ 
$$Objective(t) = \frac{Pload_{avg}}{Pload_{max}} \times \frac{1}{Pload(t)} \times \sum_{s=1}^{24} Pload(s)$$

■ 
$$PLoad(t) = Forecast(t) + Connect(t) + Disconnect(t)$$

- Where  $Forecast(t)$  is the forecasted consumption at time  $t$ , and
- $Connect(t)$  and  $Disconnect(t)$  are the amount of loads to be connected and disconnected at time  $t$  respectively during the load shifting.

- *Connect(t)* is made up of two parts:

- The increment in the load at time  $t$  due to the connection times of devices shifted to time  $t$ .
- The increment in the load at time  $t$  due to the device connections scheduled for times that precede  $t$ .

- **$Connect(t) = \sum_{i=1}^N \sum_{k=1}^D X_{kit} \cdot P_{1k} + \sum_{l=1}^{j-1} \sum_{i=1}^{t-1} \sum_{k=1}^D X_{ki(t1)} P_{(1+l)k}$**

- Where  $X_{kit}$  is the number of devices of type  $k$  that are shifted from time step  $i$  to  $t$ .
- $D$  is the number of device types,  $P_{1k}$  and  $P_{(1+l)k}$  are the power consumptions at time steps 1 and  $(1+l)$

- Similarly, *Disconnect()* also consists of two parts:

- The decrement in the load due to delay in connection times of devices that were originally supposed to begin their consumption at time step  $t$ .
- The decrement in the load due to delay in connection times of devices that were expected to start their consumption at time steps that precede  $t$ .

- **$Disconnect(t) = \sum_{q=t+1}^{t+m} \sum_{k=1}^D X_{kqt} \cdot P_{1k} + \sum_{l=1}^{j-1} \sum_{q=t+1}^{t+m} \sum_{k=1}^D X_{kq(t-1)} P_{(1+l)k}$**

- Where  $X_{kqt}$  is the number of devices of type  $k$  that are delayed from time step  $t$  to  $q$
- $m$  is the maximum allowable delay.

- The number of devices shifted cannot be a negative value.

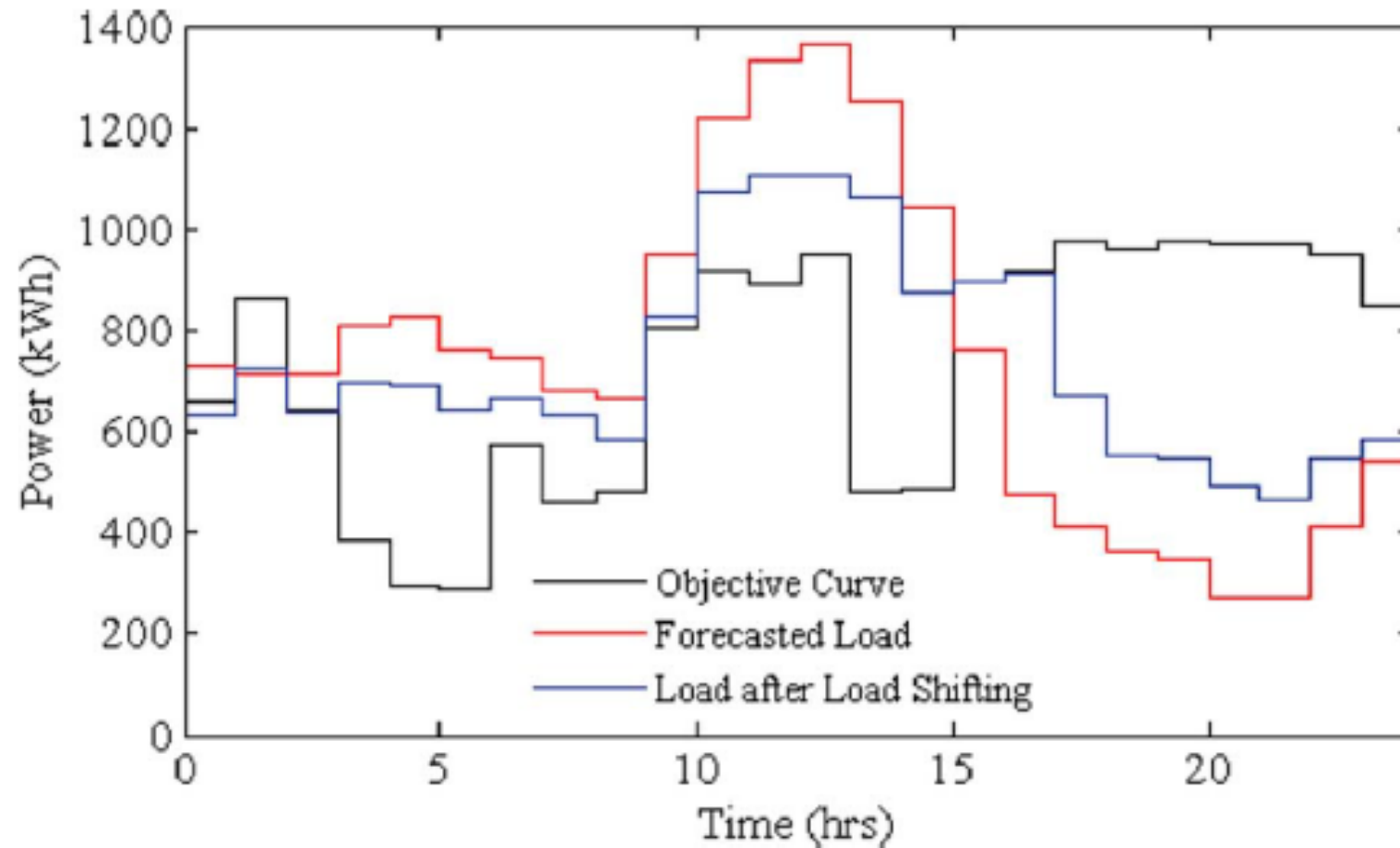
- $X_{kit} > 0$  for all  $i, j, k$ .

- The number of devices shifted away from a time step cannot be more than the number of devices available for control at the time step.

- $\sum_{t=1}^N X_{kit} \leq Ctrlable(i)$

- Where  $Ctrlable(i)$  is the number of devices of type  $k$  available for control at time step  $i$ .

# Expected Outcome



# Tool : Genetic algorithm

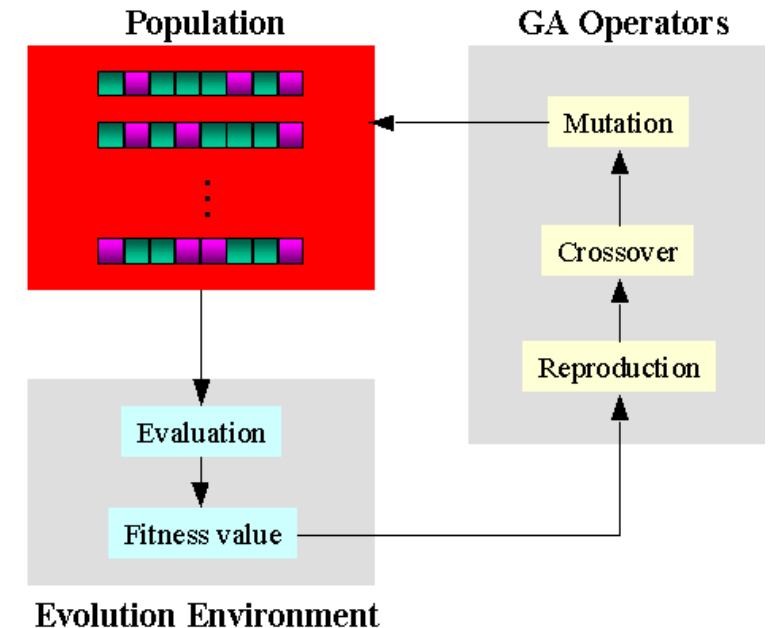
- The genetic algorithm is a method for solving both constrained and unconstrained optimization problems that is based on natural selection, that drives biological evolution.
- The genetic algorithm differs from a classical, derivative-based, optimization algorithm in two main ways

CLASSICAL ALGORITHM	GENETIC ALGORITHM
1.Generates a single point at each iteration. The sequence of points approaches an optimal solution.	1.Generates a population of points at each iteration. The best point in the population approaches an optimal solution.
2.Selects the next point in the sequence by a deterministic computation	2.Selects the next population by computation which uses random number generators

# The Process Flow

In the genetic algorithm process is as follows:

- Step 1. Determine the number of chromosomes, generation, and mutation rate and crossover rate value
- Step 2. Generate chromosome-chromosome number of the population, and the initialization value of the genes chromosome-chromosome with a random value
- Step 3. Process steps 4-7 until the number of generations is met
- Step 4. Evaluation of fitness value of chromosomes by calculating objective function
- Step 5. Chromosomes selection
- Step 5. Crossover
- Step 6. Mutation
- Step 7. New Chromosomes (Offspring)
- Step 8. Solution (Best Chromosomes)



Genetic Algorithm Evolution Flow

# Initialization

- The most common type of genetic algorithm works as follows
- A population is created with a group of individuals created randomly.
- The random no generation is possible with rand() function.
- The population is initialized keeping in account the **constraints** of the given problem.
- Chromosome is initialized based on the constraints with the help of following logic

$$chromosome(l,j)=\sum_{(i=1)}^n \sum_{(j=1)}^k (A(j)+(B(j)-A(j))*rand())$$

- NOTE: rand() generates values between 0 and 1



# Evaluation and Selection...

- The fitness function of individuals in the population are then evaluated.

$$Fitness = \frac{1}{1 + \sum_{t=1}^n (PLoad(t) - Objective(t))^2}.$$

( $i=1 \rightarrow n$ )  $Prob(i) = \sum_{i=1}^n (Fitness(i) / total)$ , Here  $total = \sum_{i=1}^n Fitness(i)$

( $i=1 \rightarrow n$ )  $Cummprob(i) = \sum_{i=1}^n Cummprob(i-1) + Prob(i)$

$Ra = rand()$  [Random no is generated ( $0 \rightarrow 1$ )]

if( $Ra(i) < cummprob(j)$ ) // Selection process..

for  $m=1:k$

$newc(i,m) = p(j,m);$

The higher the fitness, the higher the chance of chromosome being selected.



# Cross over

Three Steps involved in cross over

i).Parent Selection ii).N possibilities iii). Position for cross over

Step(i)

For  $k=1:n$

$R[k] \leftarrow \text{random}(0-1);$

if ( $R[k] < p_c$ ) then

select Chromosome[k] as parent;

Step (ii):

For example: If three parents are selected then possibilities of crossover are,

Chromosome[1]  $\times$  Chromosome[4]

Chromosome[4]  $\times$  Chromosome[5]

Chromosome[5]  $\times$  Chromosome[1]

Step(iii)

$m = \text{floor}((k-1) * \text{rand}() + 1);$  (random no between 1 to (lengthofchromo-1))

for  $j=1:k$

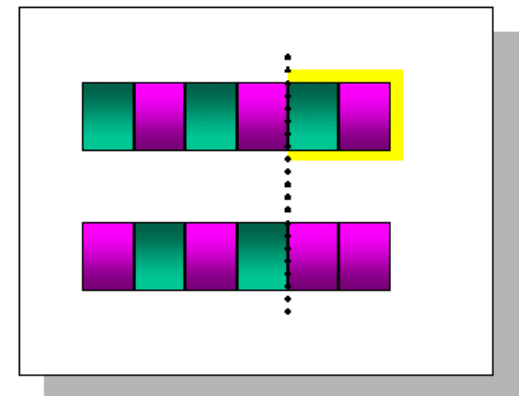
if ( $j \leq m$ )

child(i,j)=parent(i,j);

else

child(i,j)=parent(i+1,j);

m points to position of crossover



**One-Point Crossover**

# Mutation

- $\text{Total\_gen}(\text{noe}) = \text{number\_of\_gen\_in\_Chromosome} * \text{number of population}.$
- $\text{mut} = \text{pm} * \text{noe};$  where pm is the mutation rate.
- $\text{totalmut} = \text{round}(\text{mut});$
- $\text{pos} = \text{floor}((1 + (\text{noe} - 1) * \text{rand}(\text{totalmut}, 1)));$
- The value of mutated gens at mutation point is replaced by random number.
- for  $m = 1 : \text{totalmut}$
- $\text{if}((i == \text{row}(m)) \ \&\& \ (j == \text{col}(m)))$
- $\text{newc}(i, j) = s(j) + (r(j) - s(j)) * \text{rand}();$
- And after mutation process the next iteration starts and the process continues till we get the best chromosome...



Before: 1101101001101110  
After: 1101100001101110

# Future Proceedings

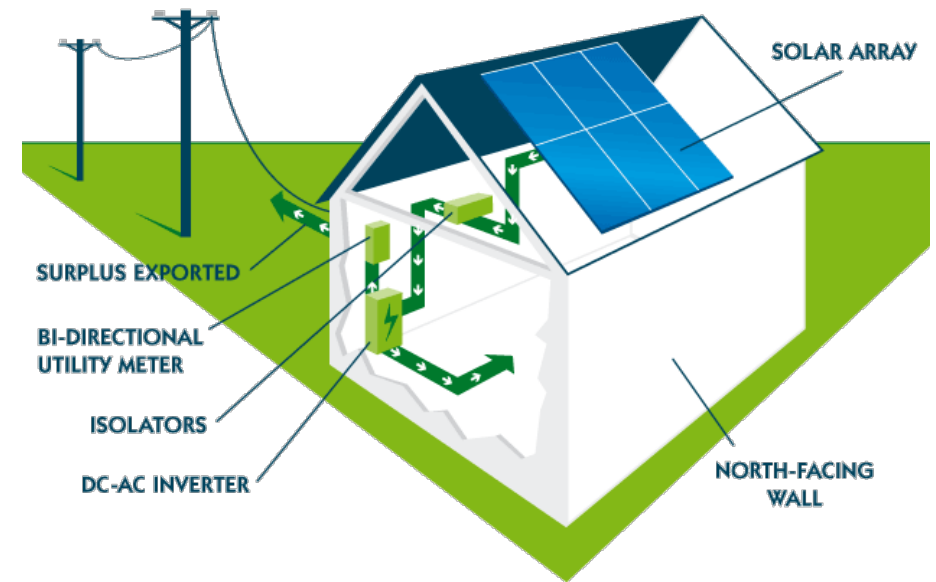
- In addition to the current work, solved as a single objective function, we add carbon footprint as a secondary objective together with DSM, treating it as a multi objective problem

$$E(t) = E_{Co2} \times Pload(t)$$

- **Multi objective Problem = DSM (Primary Objective) + Co2 Footprint (Secondary Objective)**
- Solution to this Multi objective problem can be solved using NGSA II

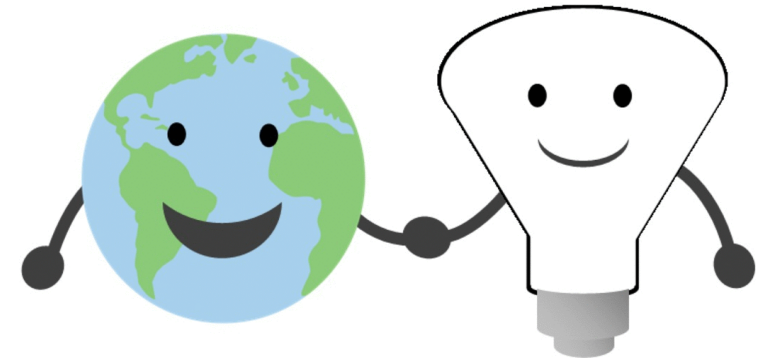
# Contribution to the Society:

- **Smartness:** Effectively managing loads in a power grid by adopting intelligent DSM Controllers.
- **Smart Pricing:** Automatic metering that allows consumers to make informed decisions regarding their energy consumption, and peak load pricing.
- **Virtual Power Plant:** Voluntarily lowering consumers demand for electricity
- **Carbon Footprint:** Reduced greenhouse gas emissions.



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Thank You !!

