

Tutorial-3

Ans-1

while ($low \leq high$)

mid = $(low + high) / 2$;
 if ($arr[mid] == key$)
 return true;
 else if ($arr[mid] > key$)
 high = mid - 1;
 else
 low = mid + 1;

return false;

Ans-2

Iterative insertion sort:

~~Algorithm~~

for (int i=1; i<n; i++)

j = i - 1;
 x = arr[j];
 while (j > -1 && arr[j] > n)
 arr[j+1] = arr[j];
 j--;
 arr[j+1] = n;

Recursive insertion sort:

Void insertionSort(int arr[], int n)

if (n <= 1)
 return;
 insertionSort(arr, n-1);
 int last = arr[n-1];
 j = n-2;
 while (j >= 0 && arr[j] > last)
 arr[j+1] = arr[j];
 j--;
 arr[0] = last;

Inserion sort is online sorting
 because whenever a new
 element come, insertion sort
 define its right place.

Am-3

- Bubblesort → $O(n^2)$
 Insertion sort → $O(n^2)$
 Selection sort → $O(n^2)$
 Merge sort → ~~$O(n \log n)$~~ $O(n + \log n)$
 Quick sort → $O(n \log n)$
 Count sort → $O(n)$
 Bucket sort → $O(n)$

Am-4

Online sorting → Insertionsort

stable sorting → Merge sort, Insertion sort, Bubble sort.

Inplace sorting → Bubble sort, Insertion sort, Selection sort

Am-5

Iterative Binary Search:

$O(\log n)$

```

while (low <= high)
  int mid = (low+high)/2
  if (arr[mid] == key)
    return true;
  else if (arr[mid] > key)
    High = mid-1;
  else
    low = mid+1;
  }
  
```

Recursive Binary Search:

$O(\log n)$

```

while (low <= high)
  int mid = (low+high)/2
  if (arr[mid] == key)
    return true;
  else if (arr[mid] > key)
    Binary-search(arr, low, mid-1);
  else
    Binary-search(arr, mid+1, high);
  return false;
  
```

Am-6

$$T(n) = T(n/2) + T(n/2) + c$$

Am-7

```

for map<int,int> m;
    for (int i=0; i<arr.size(); i++)
        if (m.find(target - arr[i]) == m.end())
            m[arr[i]] = i;
        else
            cout << i << " " << arr[i];
    }
}

```

Am-8 Quicksort is the fastest general purpose sort. In

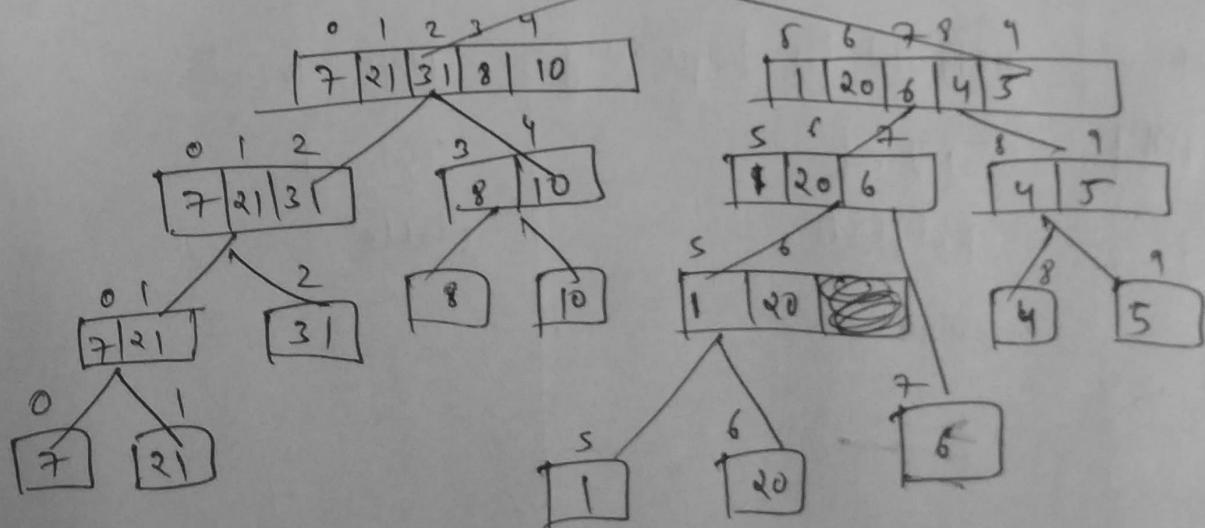
most practical situation, quicksort is the method of choice.

If stability is important and space is available, mergesort might be best.

Am-9 Inversion indicates how far or close the array is from being sorted

1	2	3	4	5	6	7	8	9
7	21	31	8	10	1	20	6	45

n=10



Inversions = 31

Ana to

Worst Case: The worst case occurs when the picked pivot is always an extreme (smallest or largest) element.

This happens when input array is sorted or reverse sorted and either first or last element is picked as pivot.
 $O(n^2)$.

Best Case: Best case occurs when pivot element is the middle element or near to the middle element.
 $O(n \log n)$

Ana II

Merge Sort: $T(n) = 2T\left(\frac{n}{2}\right) + O(n)$

Quick Sort: $T(n) = 2T\left(\frac{n}{2}\right) + n + 1$

Basic	Quick sort	Merge sort
• Partition	splitting is done in any ratio	array is partitioned into just 2 halves
• Works well on	smaller array	fine on any size of array.
• Additional space	Len (inplace)	None (Not inplace)
• Efficient	inefficient for large array	more efficient
• Sorting Method	Internal	External
• Stability	Not stable	stable

Anrit We will use MergeSort because we can divide the 4 GB data into 4 packets of 1 GB and sort them separately and combine them later.

- Internal Sorting: all the data to sort is stored in memory at all times while sorting is in progress.
- External Sorting: all the data is stored outside memory and only loaded into memory in small chunks.