

实验报告

实验名称（相对简单 CPU 电路设计）

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实验目标

利用 VHDL 设计相对简单 CPU 的电路并验证。

实验要求

- 采用 VHDL 描述电路及其测试平台
- 采用时序逻辑设计电路
- 采用从 1 累加到 n 的程序进行测试

实验内容

相对简单 CPU 的设计需求

相对简单 CPU 的设计需求请详见课件，主要特征如下：

- 地址总线 16 位，数据总线 8 位

```
entity rscpu is
    port(
        clk: in std_logic;
        reset: in std_logic;
        addrbus: out std_logic_vector(15 downto 0);
        databus: inout std_logic_vector(7 downto 0);
        readl: out std_logic;
        writel: out std_logic
    );
end entity;
```

- 有一个 8 位累加寄存器 AC，一个 8 位通用寄存器 R，一个 1 位的零标志
- 有一个 16 位 AR 寄存器，一个 16 位程序计数器 PC，一个 8 位数据寄存

器 DR，一个 8 位指令寄存器 IR，一个 8 位临时寄存器 TR

```
signal pc: std_logic_vector(15 downto 0);
signal ac: std_logic_vector(7 downto 0);
signal r: std_logic_vector(7 downto 0);
signal ar: std_logic_vector(15 downto 0);
signal ir: std_logic_vector(7 downto 0);
signal dr: std_logic_vector(7 downto 0);
signal tr: std_logic_vector(7 downto 0);
signal z: std_logic;
```

- 有 16 条指令，每条指令 1 个或 3 个字节，其中操作码 8 位。3 字节的指令有 16 位的地址

```
constant RSNOP: std_logic_vector(7 downto 0) := "00000000";
constant RSLDAC: std_logic_vector(7 downto 0) := "00000001";
constant RSSTAC: std_logic_vector(7 downto 0) := "00000010";
constant RSMVAC: std_logic_vector(7 downto 0) := "00000011";
constant RSMOVR: std_logic_vector(7 downto 0) := "00000100";
constant RSJUMP: std_logic_vector(7 downto 0) := "00000101";
constant RSJMPZ: std_logic_vector(7 downto 0) := "00000110";
constant RSJPNZ: std_logic_vector(7 downto 0) := "00000111";

constant RSADD: std_logic_vector(7 downto 0) := "00001000";
constant RSSUB: std_logic_vector(7 downto 0) := "00001001";
constant RSINAC: std_logic_vector(7 downto 0) := "00001010";
constant RSCLAC: std_logic_vector(7 downto 0) := "00001011";
constant RSAND: std_logic_vector(7 downto 0) := "00001100";
constant RSOR: std_logic_vector(7 downto 0) := "00001101";
constant RSXOR: std_logic_vector(7 downto 0) := "00001110";
constant RSNOT: std_logic_vector(7 downto 0) := "00001111";
```

相对简单 CPU 设计方案

相对简单 CPU 的设计方案请详见课件，主要思路如下：

1. 指令执行过程分为取指、译码、执行三个阶段
2. 取指包括三个状态，FETCH1，FETCH2，FETCH3，FETCH4
3. 译码体现为从 FETCH4 状态到各指令执行状态序列的第一个状态

```

when fetch4 =>
  case ir is
    when RSNOP =>
      nextstate <= nop1;
    when RSCLAC =>
      nextstate <= clacl;
    when RSSTAC=>
      nextstate <= stacl;
    when RSLDAC=>
      nextstate <= ldac1;
    when RSINAC=>
      nextstate <= inacl;
    when RSMVAC=>
      nextstate <= mvac1;
    when RSADD=>
      nextstate <= add1;
    when RSSUB=>
      nextstate <= sub1;
    when RSJPNZ=>
      if z='0' then
        nextstate <= jpnzyl;
      else
        nextstate <= jpnznl;
      end if;
    when RSJMPZ=>
      if z='1' then
        nextstate <= jmpzyl;
      else
        nextstate <= jmpznl;
      end if;
    when others =>
      nextstate <= fetch1;
  end case;

when others =>
  nextstate <= fetch1;

```

4. 执行根据指令的具体操作分为若干状态

```

constant fetch1: std_logic_vector(5 downto 0) := "000000";
constant fetch2: std_logic_vector(5 downto 0) := "000001";
constant fetch3: std_logic_vector(5 downto 0) := "000010";
constant fetch4: std_logic_vector(5 downto 0) := "000011";

constant nop1: std_logic_vector(5 downto 0) := "000100";

constant ldac1: std_logic_vector(5 downto 0) := "000101";
constant ldac2: std_logic_vector(5 downto 0) := "000110";
constant ldac3: std_logic_vector(5 downto 0) := "000111";
constant ldac4: std_logic_vector(5 downto 0) := "001000";
constant ldac5: std_logic_vector(5 downto 0) := "001001";

constant stac1: std_logic_vector(5 downto 0) := "001010";
constant stac2: std_logic_vector(5 downto 0) := "001011";
constant stac3: std_logic_vector(5 downto 0) := "001100";
constant stac4: std_logic_vector(5 downto 0) := "001101";
constant stac5: std_logic_vector(5 downto 0) := "001110";

constant mvac1: std_logic_vector(5 downto 0) := "001111";

constant movr1: std_logic_vector(5 downto 0) := "010000";

constant jump1: std_logic_vector(5 downto 0) := "010001";
constant jump2: std_logic_vector(5 downto 0) := "010010";
constant jump3: std_logic_vector(5 downto 0) := "010011";

constant jmpzn1: std_logic_vector(5 downto 0) := "010100";
constant jmpzn2: std_logic_vector(5 downto 0) := "010101";

constant jpnzn1: std_logic_vector(5 downto 0) := "010110";
constant jpnzn2: std_logic_vector(5 downto 0) := "010111";

constant jmpzy1: std_logic_vector(5 downto 0) := "011000";
constant jmpzy2: std_logic_vector(5 downto 0) := "011001";
constant jmpzy3: std_logic_vector(5 downto 0) := "011010";

constant jpnzy1: std_logic_vector(5 downto 0) := "011011";
constant jpnzy2: std_logic_vector(5 downto 0) := "011100";
constant jpnzy3: std_logic_vector(5 downto 0) := "011101";

constant add1: std_logic_vector(5 downto 0) := "011110";
constant sub1: std_logic_vector(5 downto 0) := "011111";
constant ina1: std_logic_vector(5 downto 0) := "100000";
constant cla1: std_logic_vector(5 downto 0) := "100001";
constant and1: std_logic_vector(5 downto 0) := "100010";
constant or1: std_logic_vector(5 downto 0) := "100011";
constant xor1: std_logic_vector(5 downto 0) := "100100";
constant not1: std_logic_vector(5 downto 0) := "100101";

```

5. 根据当前状态判断下一状态；执行的最后一个状态转移到 FETCH1 状态

```

if (state=clac1) then nextstate<=fetch1;
elsif (state=inacl) then nextstate<=fetch1;
--alu
elsif (state=addl) then nextstate<=fetch1;
elsif (state=subl) then nextstate<=fetch1;
elsif (state=andl) then nextstate<=fetch1;
elsif (state=orl) then nextstate<=fetch1;
elsif (state=xorl) then nextstate<=fetch1;
elsif (state=notl) then nextstate<=fetch1;
--mov
elsif (state=mvac1) then nextstate<=fetch1;
elsif (state=movr1) then nextstate<=fetch1;
--load
elsif (state=ldac1) then nextstate<=ldac2;
elsif (state=ldac2) then nextstate<=ldac3;
elsif (state=ldac3) then nextstate<=ldac4;
elsif (state=ldac4) then nextstate<=ldac5;
elsif (state=ldac5) then nextstate<=fetch1;
--store
elsif (state=stac1) then nextstate<=stac2;
elsif (state=stac2) then nextstate<=stac3;
elsif (state=stac3) then nextstate<=stac4;
elsif (state=stac4) then nextstate<=stac5;
elsif (state=stac5) then nextstate<=fetch1;

--jump
elsif (state=jump1) then nextstate<=jump2;
elsif (state=jump2) then nextstate<=jump3;
elsif (state=jump3) then nextstate<=fetch1;
--jmpz
elsif (state=jmpzy1) then nextstate<=jmpzy2;
elsif (state=jmpzy2) then nextstate<=jmpzy3;
elsif (state=jmpzy3) then nextstate<=fetch1;
elsif (state=jmpzn1) then nextstate<=jmpzn2;
elsif (state=jmpzn2) then nextstate<=fetch1;
--jpnz
elsif (state=jpnzy1) then nextstate<=jpnzy2;
elsif (state=jpnzy2) then nextstate<=jpnzy3;
elsif (state=jpnzy3) then nextstate<=fetch1;
elsif (state=jpnzn1) then nextstate<=jpnzn2;
elsif (state=jpnzn2) then nextstate<=fetch1;

```

6. 控制器根据每个状态需要完成的操作产生相应的控制信号

```

case state is
  when fetch1 =>      --AR-PC      when fetch2 =>      --DR-M, PC-PC+1
    arload <= '1';      arload <= '0';
    arinc <= '0';      arinc <= '0';
    pcload <= '0';      pcload <= '0';
    pcinc <= '0';      pcinc <= '1';
    drload <= '0';      drload <= '1';
    trload <= '0';      trload <= '0';
    irload <= '0';      irload <= '0';
    rload <= '0';      rload <= '0';
    acload <= '0';      acload <= '0';
    acinc <= '0';      acinc <= '0';
    writel <= '0';      writel <= '0';
    readl <= '0';      readl <= '1';
    membus <= '0';      membus <= '1';
    pcbus <= '1';      pcbus <= '0';
    drbus <= '0';      drbus <= '0';
    trbus <= '0';      trbus <= '0';
    rbus <= '0';      rbus <= '0';
    acbus <= '0';      acbus <= '0';
    cle <= '0';      cle <= '0';
    busmem <= '0';      busmem <= '0';

  when fetch3 =>      --IR-DR      when fetch4 =>      --AR-PC
    arload <= '0';      arload <= '1';
    arinc <= '0';      arinc <= '0';
    pcload <= '0';      pcload <= '0';
    pcinc <= '0';      pcinc <= '0';
    drload <= '0';      drload <= '0';
    trload <= '0';      trload <= '0';
    irload <= '1';      irload <= '1';
    rload <= '0';      rload <= '0';
    acload <= '0';      acload <= '0';
    acinc <= '0';      acinc <= '0';
    writel <= '0';      writel <= '0';
    readl <= '0';      readl <= '0';
    membus <= '0';      membus <= '0';
    pcbus <= '0';      pcbus <= '1';
    drbus <= '1';      drbus <= '0';
    trbus <= '0';      trbus <= '0';
    rbus <= '0';      rbus <= '0';
    acbus <= '0';      acbus <= '0';
    cle <= '0';      cle <= '0';
    busmem <= '0';      busmem <= '0';

```

其余状态同理，根据具体的需要执行的操作，产生相应的控制信号。

7. 根据控制信号执行相应的操作

```

-- address and data bus
addrbus <= ar;
databus <= dr when busmem='1' else "ZZZZZZZZ";

--the bus
thebus<=pc
      "00000000"&databus      when pcbus='1'      else
      "00000000"&r            when membus='1'      else
      "00000000"&ac           when rbus='1'        else
      dr&tr                  when acbus='1'        else
      "00000000"&dr           when (trbus='1' and drbus='1') else
      "ZZZZZZZZZZZZZZZZZZ";  when (drbus='1' and trbus/='1') else

alu_result <= std_logic_vector(unsigned(ac)-unsigned(thebus(7 downto 0))) when state=sub1 else
               std_logic_vector(unsigned(ac)+unsigned(thebus(7 downto 0))) when state=add1 else
               std_logic_vector(ac and thebus(7 downto 0)) when state=and1 else
               std_logic_vector(ac or thebus(7 downto 0)) when state=or1 else
               std_logic_vector(ac xor thebus(7 downto 0)) when state=xor1 else
               std_logic_vector(not ac) when state=not1 else
               thebus(7 downto 0);

update_regs: process(clk)
begin
    if(rising_edge(clk)) then

        if(arload='1') then
            ar <= thebus;
        end if;
        if(pcload='1') then
            pc <= thebus;
        end if;
        if(drload='1') then
            dr <= thebus(7 downto 0);
        end if;
        if(trload='1') then
            tr <= dr;
        end if;
        if(irload='1') then
            ir <= dr;
        end if;
        if(rload='1') then
            r <= thebus(7 downto 0);
        end if;

        if(acload='1') then
            ac <= alu_result;
            if(state=ADD1 or state=SUB1 or state=NOT1 or state=AND1 or state=OR1 or state=XOR1) then
                if(alu_result="00000000") then
                    z<='1';
                else
                    z<='0';
                end if;
            end if;
        end if;

        if(arinc='1') then
            ar <= std_logic_vector(unsigned(ar)+1);
        end if;
        if(pcinc='1') then
            pc <= std_logic_vector(unsigned(pc)+1);
        end if;
        if(acinc='1') then
            ac <= std_logic_vector(unsigned(ac)+1);
            if(ac="11111111") then
                z<='1';
            else
                z<='0';
            end if;
        end if;
        if(cle='1') then
            ac <= "00000000";
            z <='1';
        end if;
    end if;
end if;

```

```

        if(reset='1') then
            pc <= X"0000";
            state <= fetch1;
        else
            --pc <= nextpc;
            state <= nextstate;
        end if; |
    end if;

```

测试

测试平台

相对简单 CPU 电路在如下机器上进行了测试：

操作系统：Windows10

综合软件：Quartus II Help Version 9.0

仿真软件：Quartus II Help Version 9.0

测试输入

我们采用从 1 累加到 n 的程序作为测试输入：

```

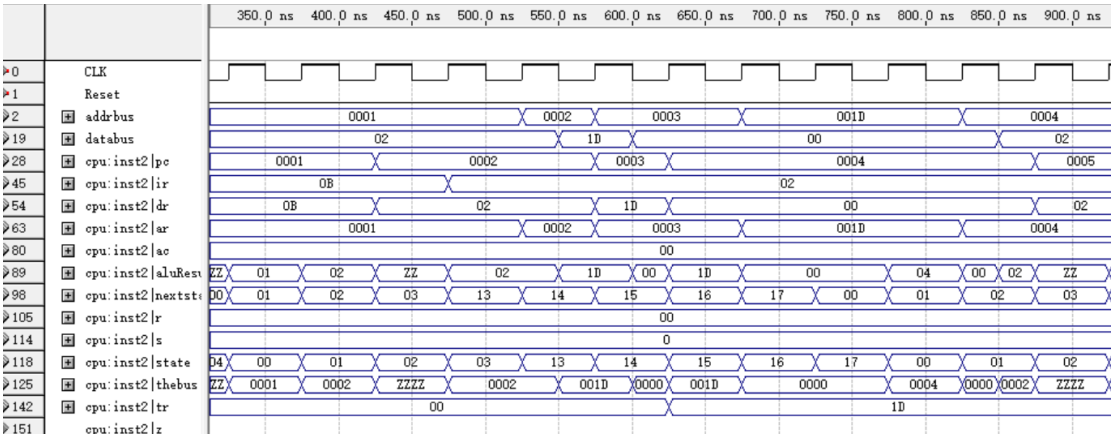
        CLAC
        STAC total } total = 0, i = 0
        STAC i
Loop:   LDAC i      }
        INAC        } i = i + 1
        STAC i
        MVAC
        LDAC total } total = total + i
        ADD
        STAC total
        LDAC n      }
        SUB          } IF i ≠ n THEN GOTO Loop
        JPNZ Loop
total:
i:

```

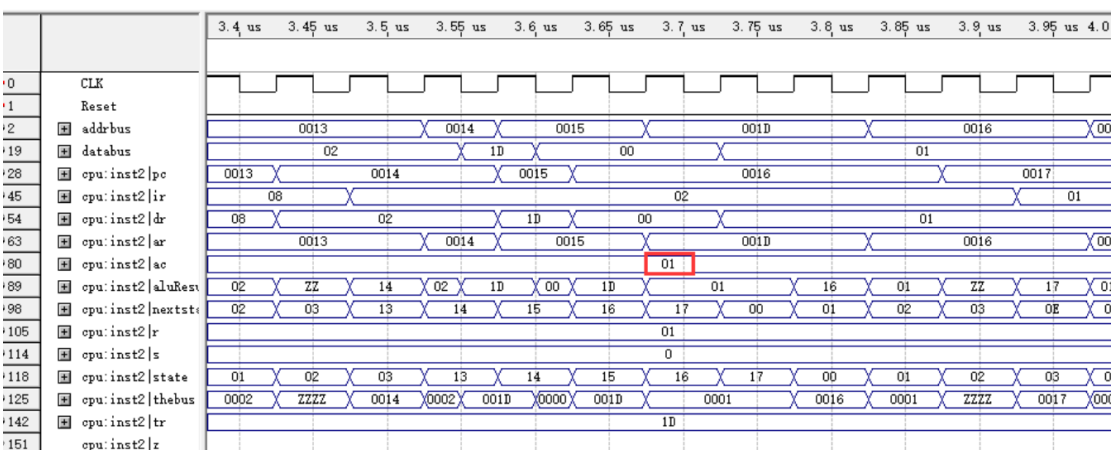

测试记录

相对简单 CPU 运行测试程序的 PC 寄存器、IR 寄存器、AC 累加器等信号波形截图如下：

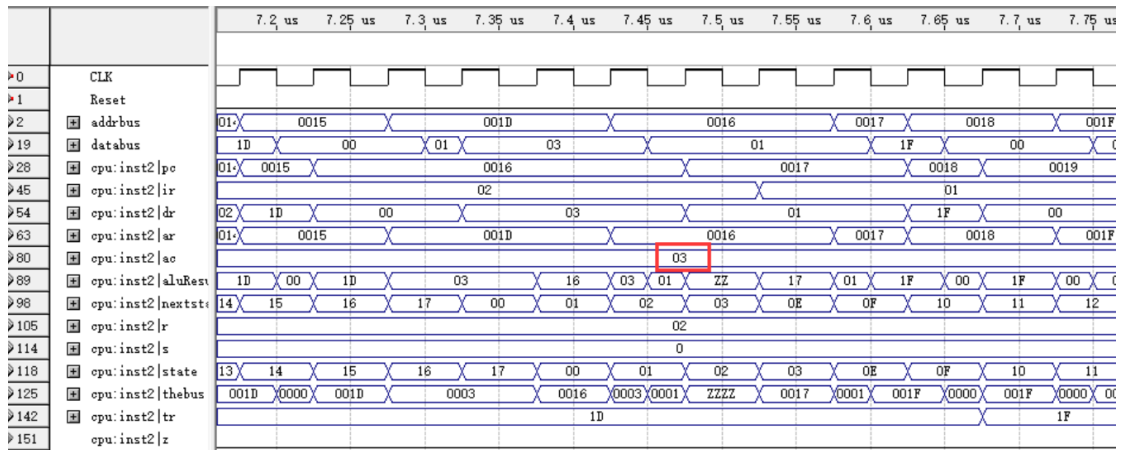
初始状态：



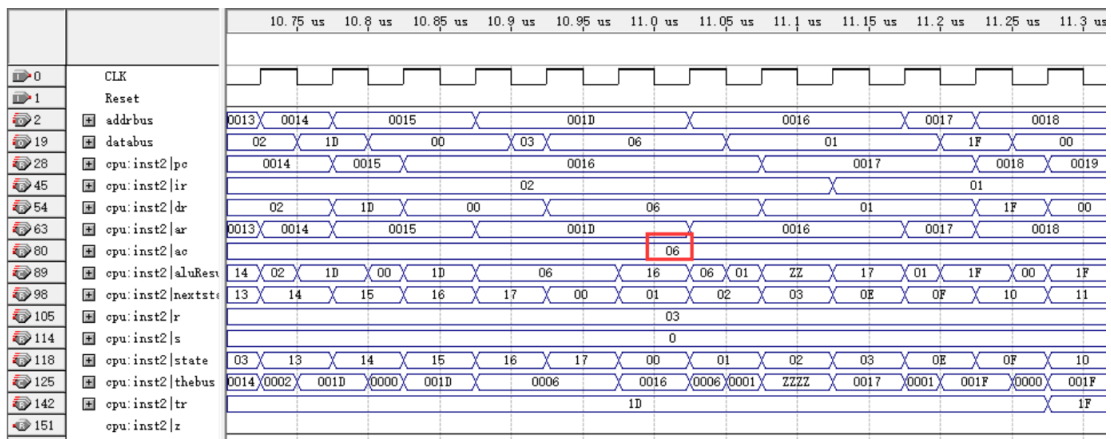
i=1,total=0 ->total=total+i=1



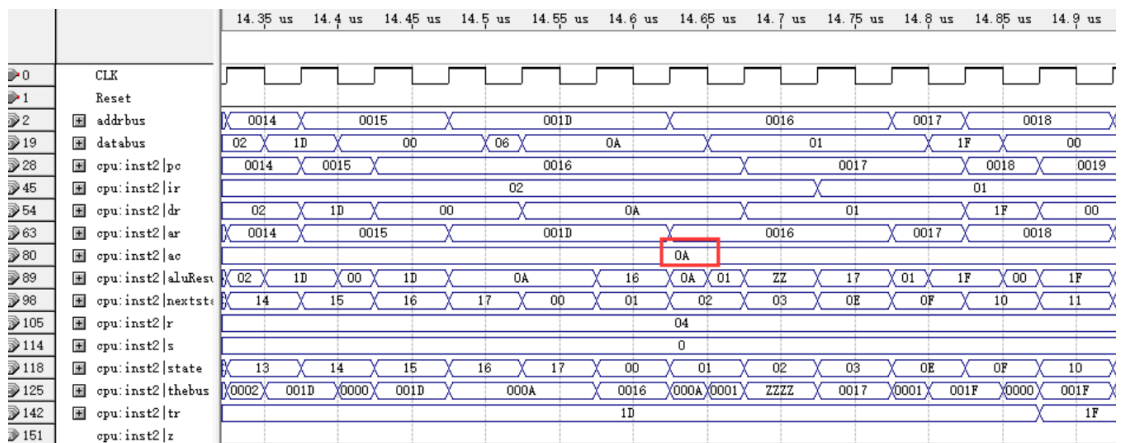
i=2,total=1 ->total=total+i=3



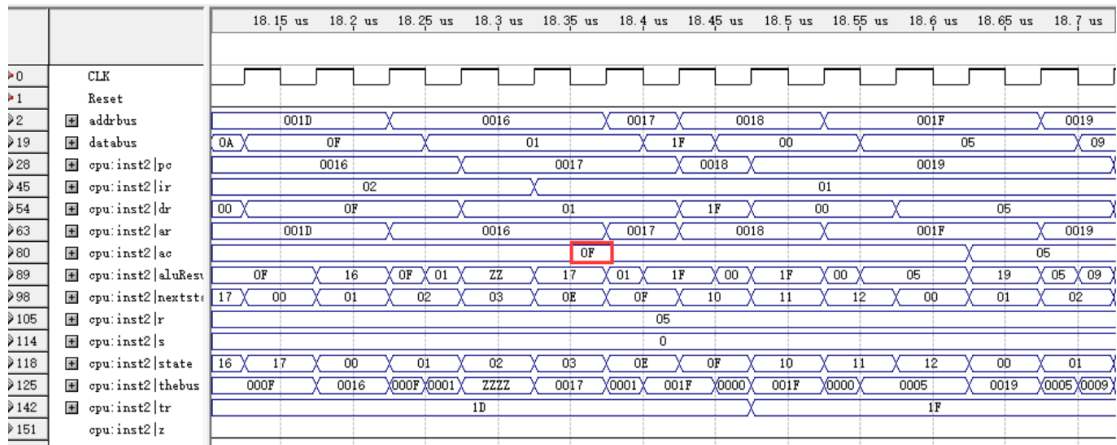
i=3,total=3 ->total=total+i=6



i=4,total=6 ->total=total+i=10(0xa)



i=5,total=10 ->total=total+i=15(0xf)



分析和结论

从测试记录来看，相对简单 CPU 实现了对测试程序指令的读取、译码和执行，得到的运算结果正确。

根据分析结果，可以认为所设计的相对简单 CPU 实现了所要求的功能，完成了实验目标。

这次实验也让我更加熟悉了计算机系统、指令执行的过程、VHDL 语言。虽然之前也有设计过 CPU，但和这次实验也有所不同，这次实验实现的过程和之前不太一样，之前都是先设计好每个模块，再用模块图将每个模块连接起来。这次实验更注重对整体实现的流程，运行过程的把握，更用了一种纯 VHDL 语言的描述方式。在完成实验的过程中，不断学习，了解并熟悉运行过程，更有大家一起讨论交流自己遇到的问题，解决方法，实现过程。