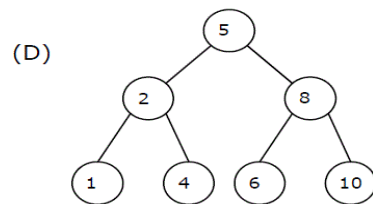
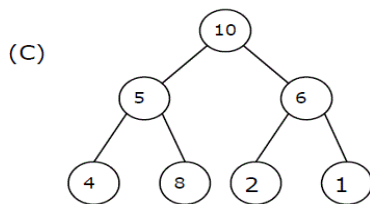
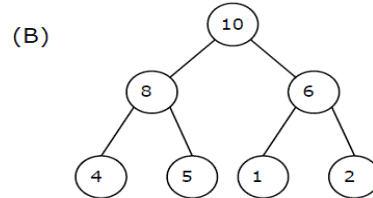
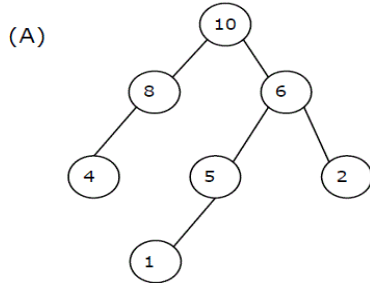


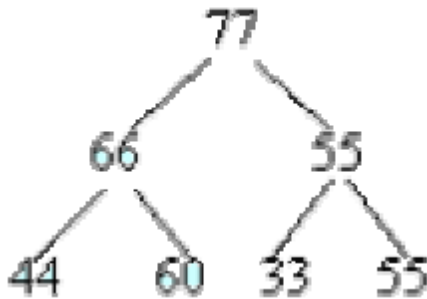
Handwriting Assignment #2 Solution

1. Which of the following is a max-heap?



Solution) (B)

2. Draw the following list of numbers as a heap with the first number as the root: 77, 66, 55, 44, 60, 33, 55
solution)



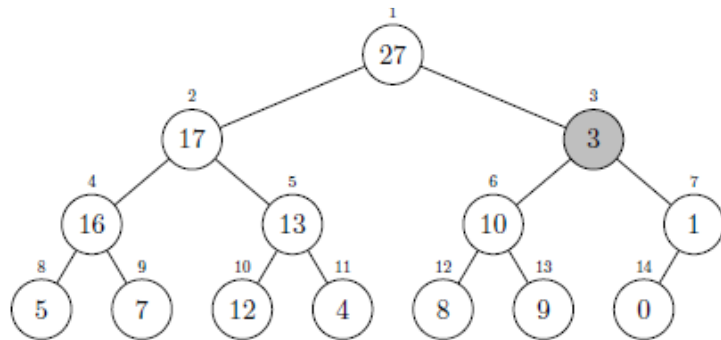
3. What is the minimum and maximum numbers of elements in a heap of height h ?
solution)

Minimum # of nodes happens in a heap in which the last level contains only one node.
Thus, minimum # of nodes = $2^0 + 2^1 + 2^2 + \dots + 2^{h-1} + 1 = 2^h$

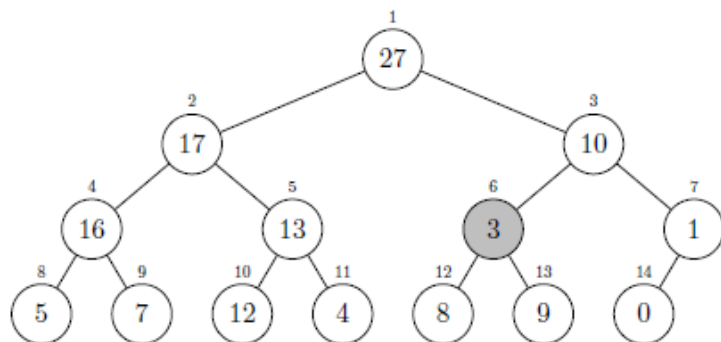
Maximum # of nodes happens in a heap in which the last level is full.
Thus, maximum # of nodes = $2^0 + 2^1 + 2^2 + \dots + 2^h = 2^{h+1} - 1$.

4. Illustrate the operation of Max-Heapify(A, 3) on the array
 $A = \langle 27, 17, 3, 16, 13, 10, 1, 5, 7, 12, 4, 8, 9, 0 \rangle$.
 solution)

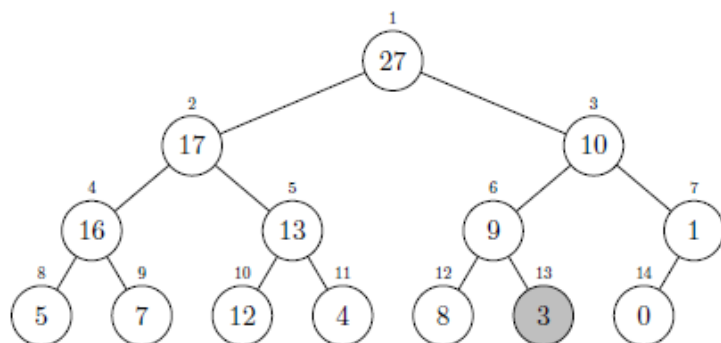
Step 0:



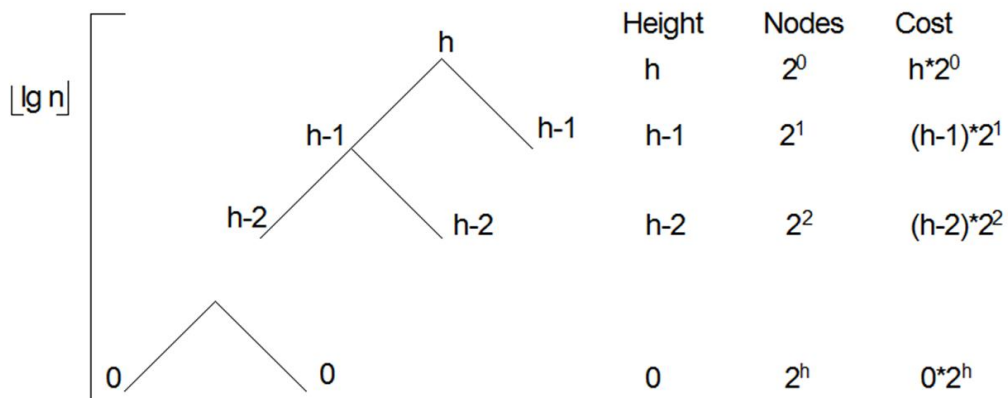
Step 1:



Step 2:



5. Build Heap is used to build a max(or min) binary heap from a given array. Build Heap is used in Heap Sort as a first step for sorting.
 What is the time complexity of Build Heap operation?



$$\begin{aligned}
\sum_{h=0}^{\lfloor \lg n \rfloor} h * 2^{\lfloor \lg n \rfloor - h} &= \sum_{h=0}^{\lfloor \lg n \rfloor} h * \frac{2^{\lfloor \lg n \rfloor}}{2^h} \leq \sum_{h=0}^{\lfloor \lg n \rfloor} h * \frac{2^{\lg n}}{2^h} = \sum_{h=0}^{\lfloor \lg n \rfloor} h * \frac{n^{\lg 2}}{2^h} \\
&= \sum_{h=0}^{\lfloor \lg n \rfloor} h * \frac{n}{2^h} = n \sum_{h=0}^{\lfloor \lg n \rfloor} \frac{h}{2^h} \leq n \sum_{h=0}^{\infty} \frac{h}{2^h} = 2n = O(n)
\end{aligned}$$

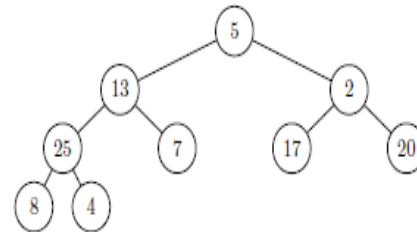
6. Illustrate the operation of Heapsort on the array

$A = \langle 5, 13, 2, 25, 7, 17, 20, 8, 4 \rangle$.

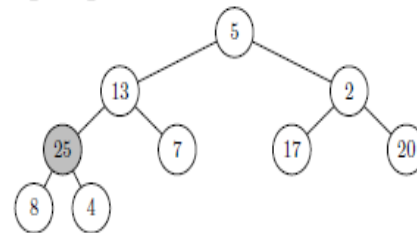
solution)

Since $A.length = 9$, the command $\text{MAX-HEAPIFY}(A, i)$ is called for $i = 4, 3, 2, 1$. The action of BUILD-MAX-HEAP is as follows (these first few diagrams are not required for a correct answer), with the nodes exchanged at each step shaded:

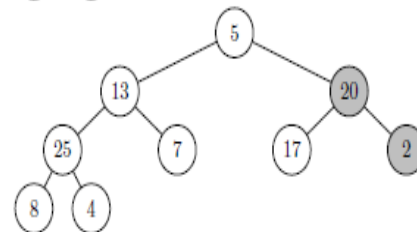
$\text{BUILD-MAX-HEAP}(A)$:



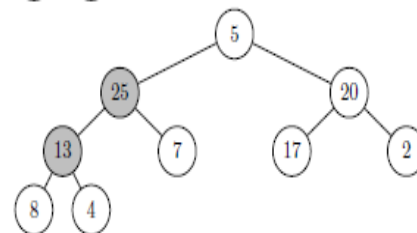
$\text{MAX-HEAPIFY}(A, 4)$:



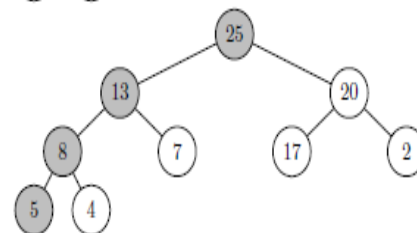
$\text{MAX-HEAPIFY}(A, 3)$:

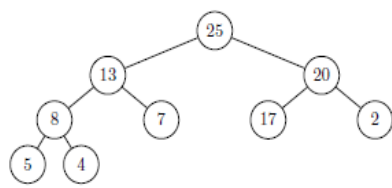


$\text{MAX-HEAPIFY}(A, 2)$:

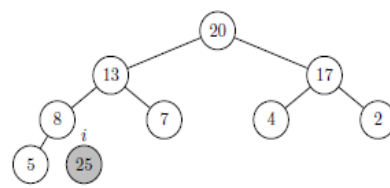


$\text{MAX-HEAPIFY}(A, 1)$:

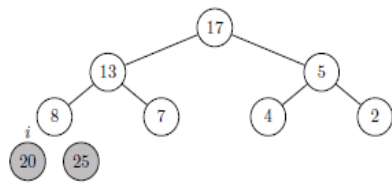




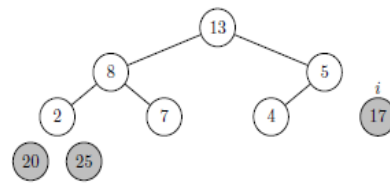
Step 0



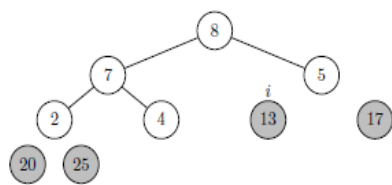
Step 1



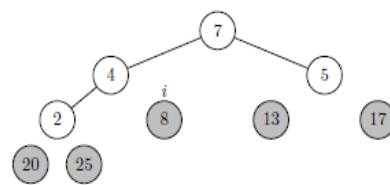
Step 2



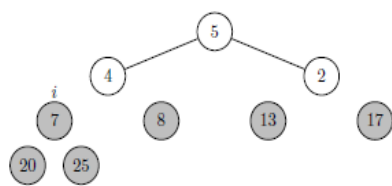
Step 3



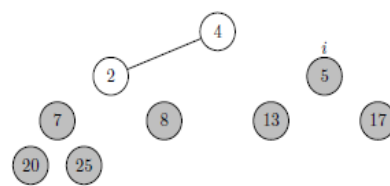
Step 4



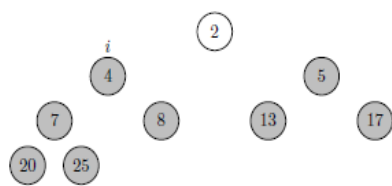
Step 5



Step 6



Step 7



Step 8

This gives a final sorted array $A = (2, 4, 5, 7, 8, 13, 17, 20, 25)$.



7. A priority queue can be implemented as a heap because
- The root can be easily identified as the topmost priority.
 - The heap is not always sorted so any value can be the top priority.
 - The heap always has a left bottom node that can be the top priority.
 - None of the above.
- solution) a.