

# 제5강 지연시간, 손실, 처리율

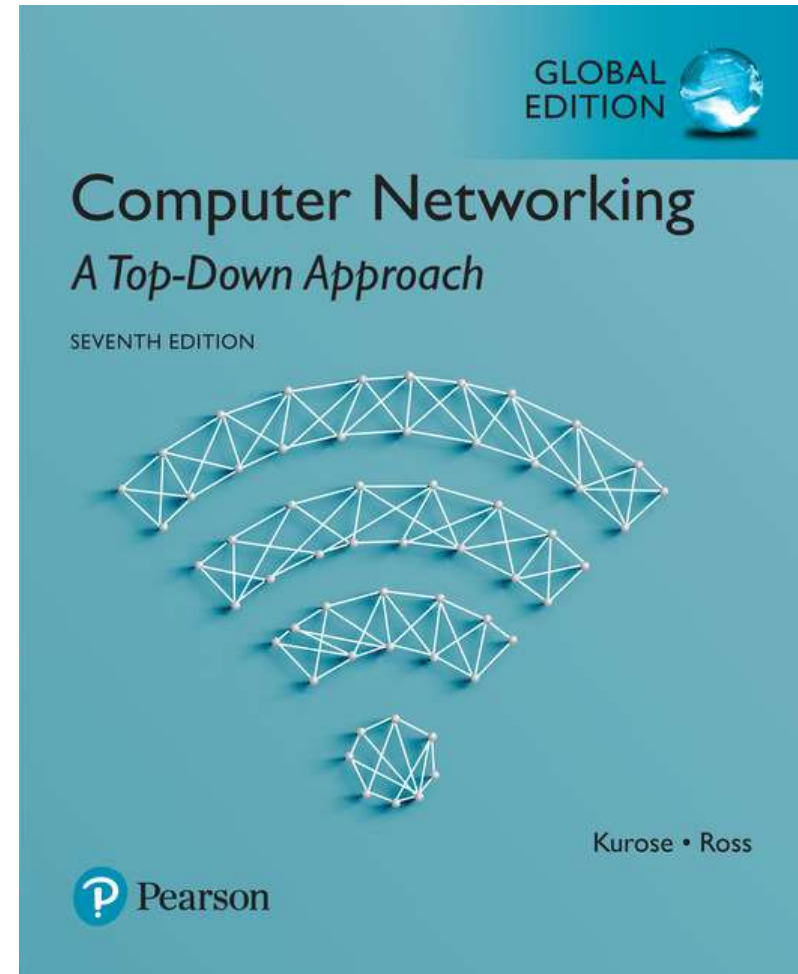
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## *Computer Networking: A Top Down Approach*

컴퓨터 네트워크  
(2019년 1학기)

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# Pre-study Test :

1) 다음 중 가변적인 지연시간은 어느 것인가?

- ① 처리 지연시간(processing delay)
- ② 큐잉 지연시간(queueing delay)
- ③ 전송 지연시간(transmission delay)
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2) 다음 중 패킷이 손실될 수 있는 곳이 아닌 곳은?

- ① 송신자
- ② 수신자
- ③ 경로상의 스위치
- ④ 경로상의 라우터

3) 다음 중 라우터에서 혼잡 상황이 발생하지 않는 경우는 어떤 경우인가?  
 $L$ 은 패킷의 길이,  $a$ 은 패킷 도착율,  $R$ 은 링크의 전송속도이다.

- ①  $L a / R \rightarrow 0$
- ②  $L a / R \rightarrow 1$
- ③  $L a / R = 1$
- ④  $L a / R > 1$

4) 전송속도가 각각  $R_a$ ,  $R_b$ ,  $R_c$ ,  $R_d$ 이고,  $R_a > R_b > R_c > R_d$  관계를 가지는 링크가 연결된 전송 경로가 있다. 해당 경로의 총 전송속도는 얼마인가?

- ①  $R_a$
- ②  $R_d$
- ③  $R_a + R_b + R_c + R_d$
- ④  $R_a * R_b * R_c * R_d$

5) 다음 중 최종 전송속도를 결정하는 링크는?

- ① Client link
- ② Server link
- ③ Backbone link
- ④ Bottleneck link

6) 라우터에서 혼잡 상황이 발생하지 않게 하려면 어떻게 하면 될까?

- ① 입력 링크의 전송속도를 높인다.
- ② 패킷 버퍼의 크기를 키운다.
- ③ 라우터의 CPU 처리속도를 높인다.
- ④ 출력 링크의 전송속도를 높인다.

# Chapter 1: roadmap

1.1 what is the Internet?

1.2 network edge

- end systems, access networks, links

1.3 network core

- packet switching, circuit switching, network structure

1.4 delay, loss, throughput in networks

1.5 protocol layers, service models

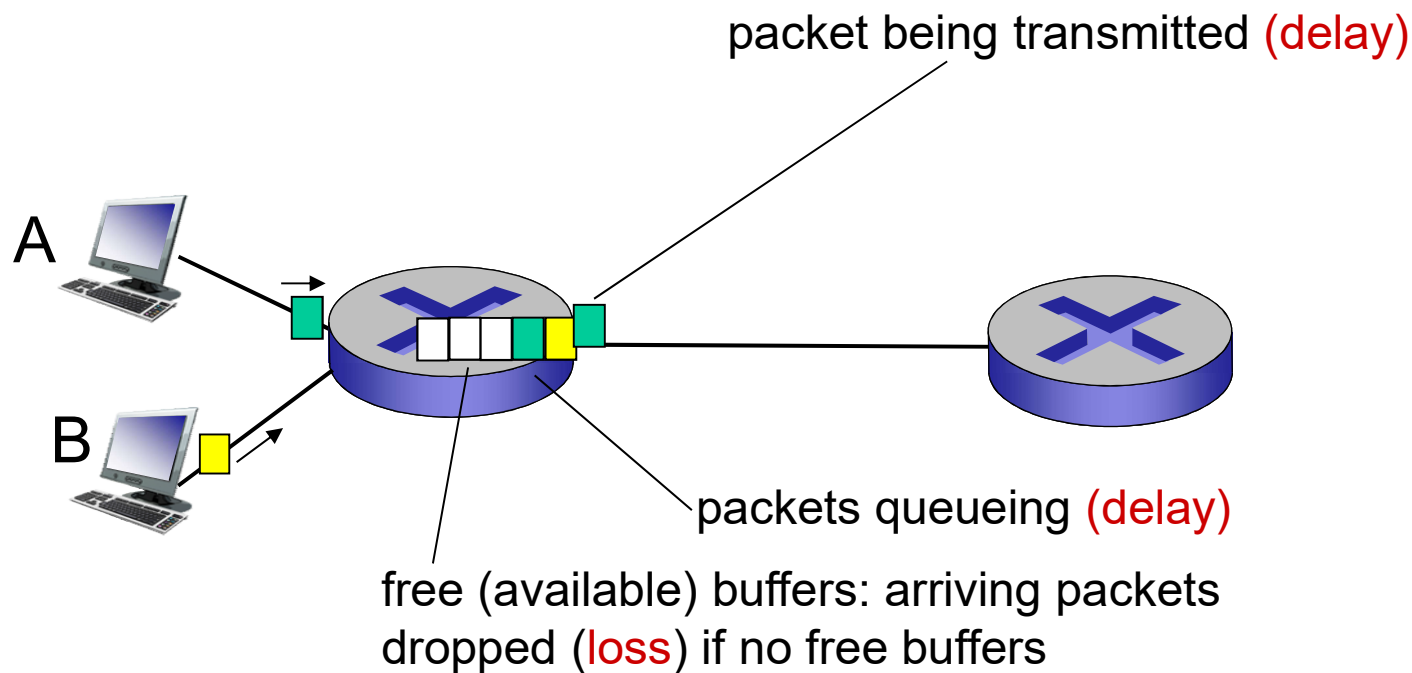
1.6 networks under attack: security

1.7 history

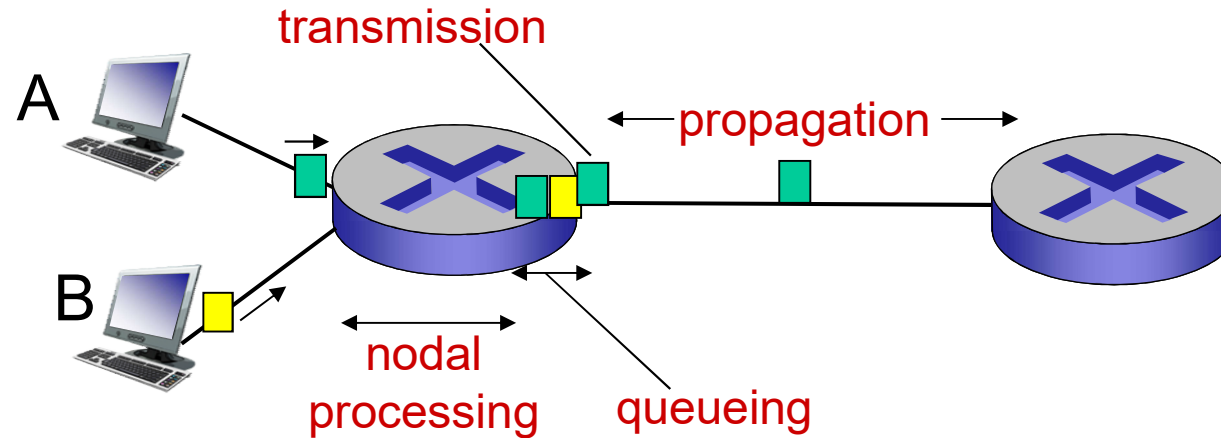
# How do loss and delay occur?

packets *queue* in router buffers

- packet arrival rate to link (temporarily) exceeds output link capacity
- packets queue, wait for turn



# Four sources of packet delay



$$d_{\text{nodal}} = d_{\text{proc}} + d_{\text{queue}} + d_{\text{trans}} + d_{\text{prop}}$$

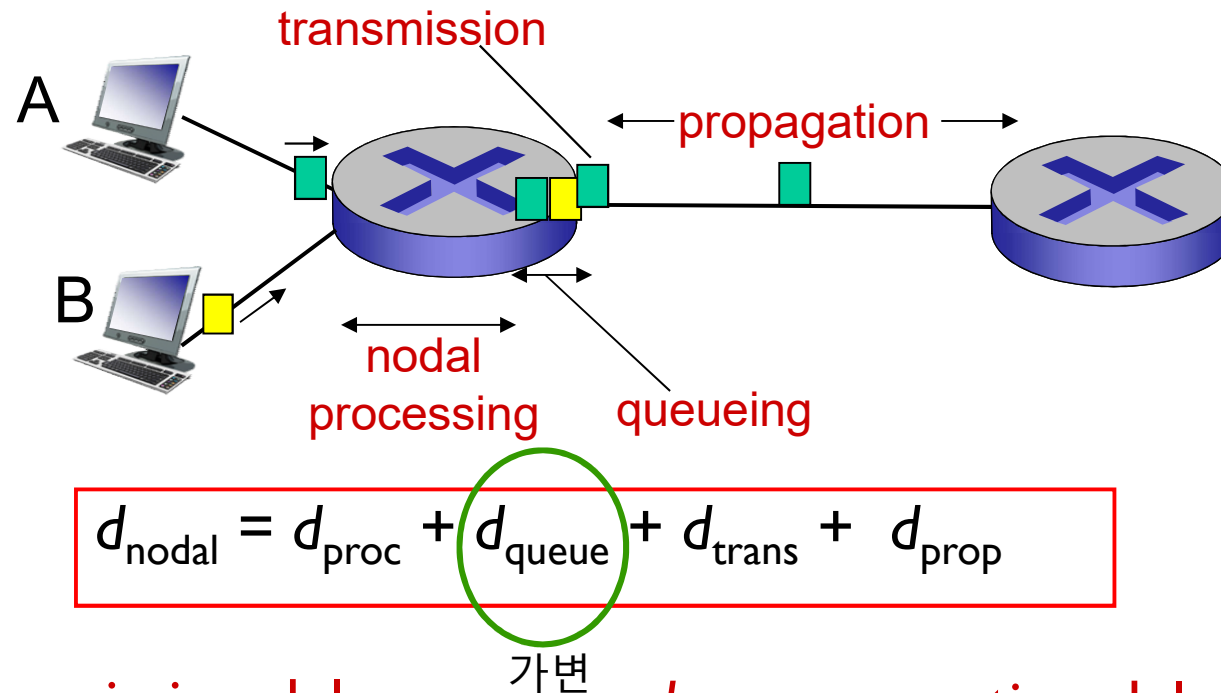
## $d_{\text{proc}}$ : nodal processing

- check bit errors
- determine output link
- typically < msec

## $d_{\text{queue}}$ : queueing delay

- time waiting at output link for transmission
- depends on congestion level of router

# Four sources of packet delay



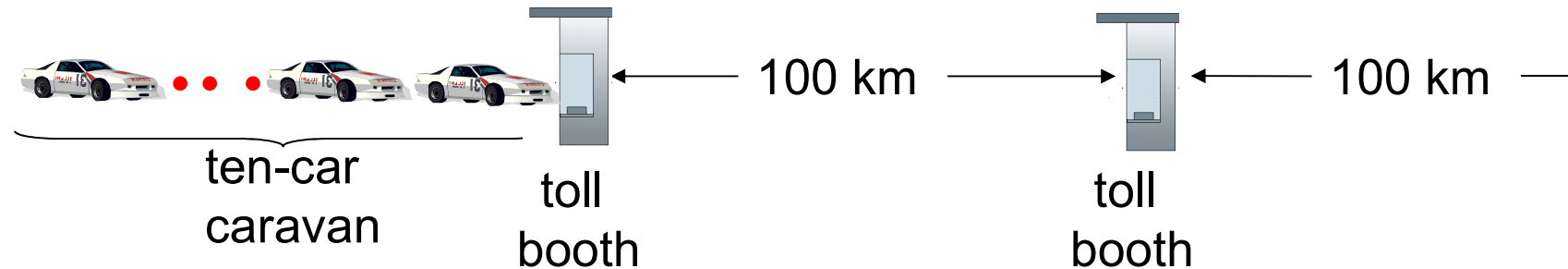
$d_{\text{trans}}$ : transmission delay:

- $L$ : packet length (bits)
  - $R$ : link bandwidth (bps)
  - $d_{\text{trans}} = L/R$
- ←  $d_{\text{trans}}$  and  $d_{\text{prop}}$  →  
very different

$d_{\text{prop}}$ : propagation delay:

- $d$ : length of physical link
- $s$ : propagation speed ( $\sim 2 \times 10^8$  m/sec)
- $d_{\text{prop}} = d/s$

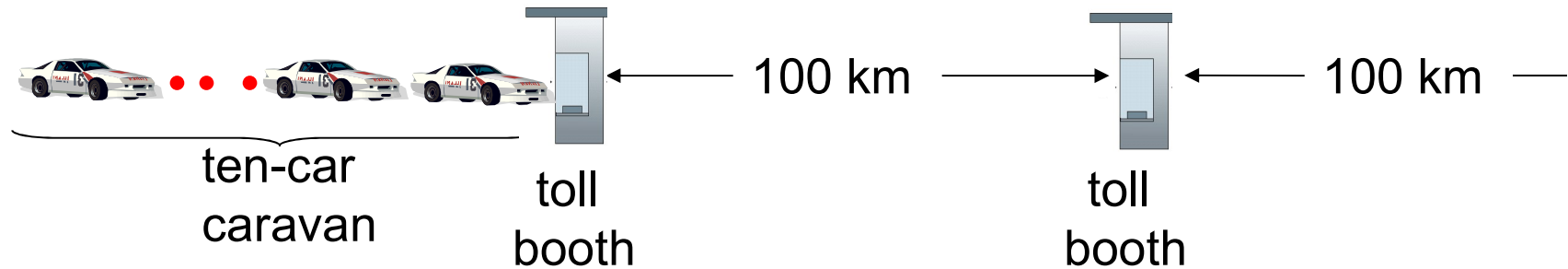
# Caravan analogy



- cars “propagate” at 100 km/hr
- toll booth takes 12 sec to service car (bit transmission time)
- car  $\sim$  bit; caravan  $\sim$  packet
- **Q: How long until caravan is lined up before 2nd toll booth?**
- time to “push” entire caravan through toll booth onto highway =  $12 * 10 = 120$  sec
- time for last car to propagate from 1st to 2nd toll booth:  
 $100\text{km} / (100\text{km/hr}) = 1$  hr
- **A: 62 minutes**



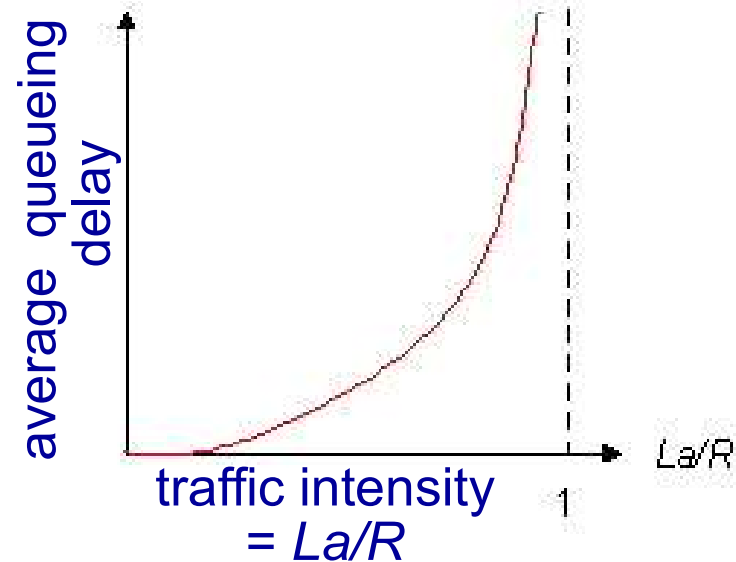
# Caravan analogy (more)



- suppose cars now “propagate” at 1000 km/hr
- and suppose toll booth now takes one min to service a car
- **Q: Will cars arrive to 2nd booth before all cars serviced at first booth?**
  - **A: Yes!** after 7 min, first car arrives at second booth; three cars still at first booth

# Queueing delay (revisited)

- $R$ : link bandwidth (bps)
- $L$ : packet length (bits)
- $a$ : average packet arrival rate



트래픽 강도 :

- $La/R \sim 0$ : avg. queueing delay small
- $La/R \rightarrow 1$ : avg. queueing delay large
- $La/R > 1$ : more “work” arriving than can be serviced, average delay infinite!



$La/R \sim 0$

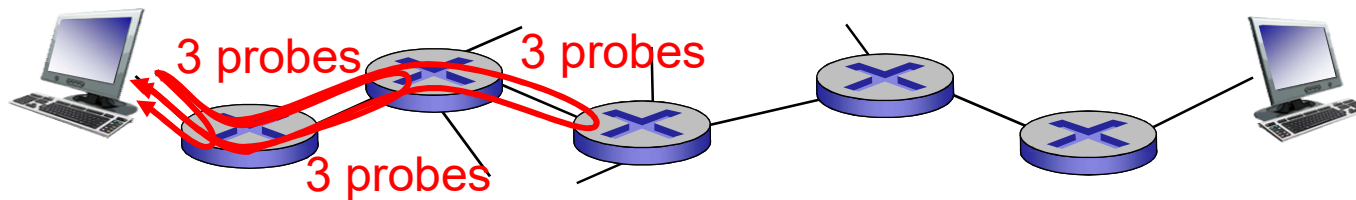


$La/R \rightarrow 1$

\* Check online interactive animation on queuing and loss

# “Real” Internet delays and routes


- what do “real” Internet delay & loss look like?
- **traceroute** program: provides delay measurement from source to router along end-end Internet path towards destination. For all  $i$ :
  - sends three packets that will reach router  $i$  on path towards destination
  - router  $i$  will return packets to sender(w/ router name & address)
  - sender times interval between transmission and reply.



# “Real” Internet delays, routes

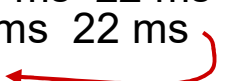
traceroute: gaia.cs.umass.edu to www.eurecom.fr

3 delay measurements from  
gaia.cs.umass.edu to cs-gw.cs.umass.edu




1 cs-gw (128.119.240.254) 1 ms 1 ms 2 ms  
2 border1-rt-fa5-1-0.gw.umass.edu (128.119.3.145) 1 ms 1 ms 2 ms  
3 cht-vbns.gw.umass.edu (128.119.3.130) 6 ms 5 ms 5 ms  
4 jn1-at1-0-0-19.wor.vbns.net (204.147.132.129) 16 ms 11 ms 13 ms  
5 jn1-so7-0-0-0.wae.vbns.net (204.147.136.136) 21 ms 18 ms 18 ms  
6 abilene-vbns.abilene.ucaid.edu (198.32.11.9) 22 ms 18 ms 22 ms  
7 nycm-wash.abilene.ucaid.edu (198.32.8.46) 22 ms 22 ms 22 ms  
8 62.40.103.253 (62.40.103.253) 104 ms 109 ms 106 ms  
9 de2-1.de1.de.geant.net (62.40.96.129) 109 ms 102 ms 104 ms  
10 de.fr1.fr.geant.net (62.40.96.50) 113 ms 121 ms 114 ms  
11 renater-gw.fr1.fr.geant.net (62.40.103.54) 112 ms 114 ms 112 ms  
12 nio-n2.cssi.renater.fr (193.51.206.13) 111 ms 114 ms 116 ms  
13 nice.cssi.renater.fr (195.220.98.102) 123 ms 125 ms 124 ms  
14 r3t2-nice.cssi.renater.fr (195.220.98.110) 126 ms 126 ms 124 ms  
15 eurecom-valbonne.r3t2.ft.net (193.48.50.54) 135 ms 128 ms 133 ms  
16 194.214.211.25 (194.214.211.25) 126 ms 128 ms 126 ms  
17 \* \* \*  
18 \* \* \*  
19 fantasia.eurecom.fr (193.55.113.142) 132 ms 128 ms 136 ms

trans-oceanic link



\* means no response (probe lost, router not replying)

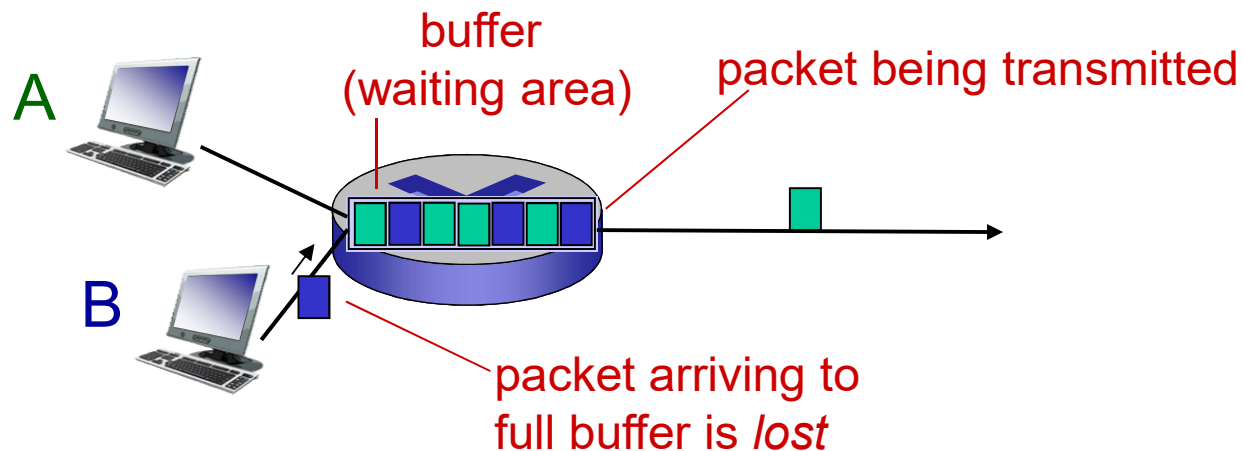


\* Do some traceroutes from exotic countries at [www.traceroute.org](http://www.traceroute.org)

# Packet loss

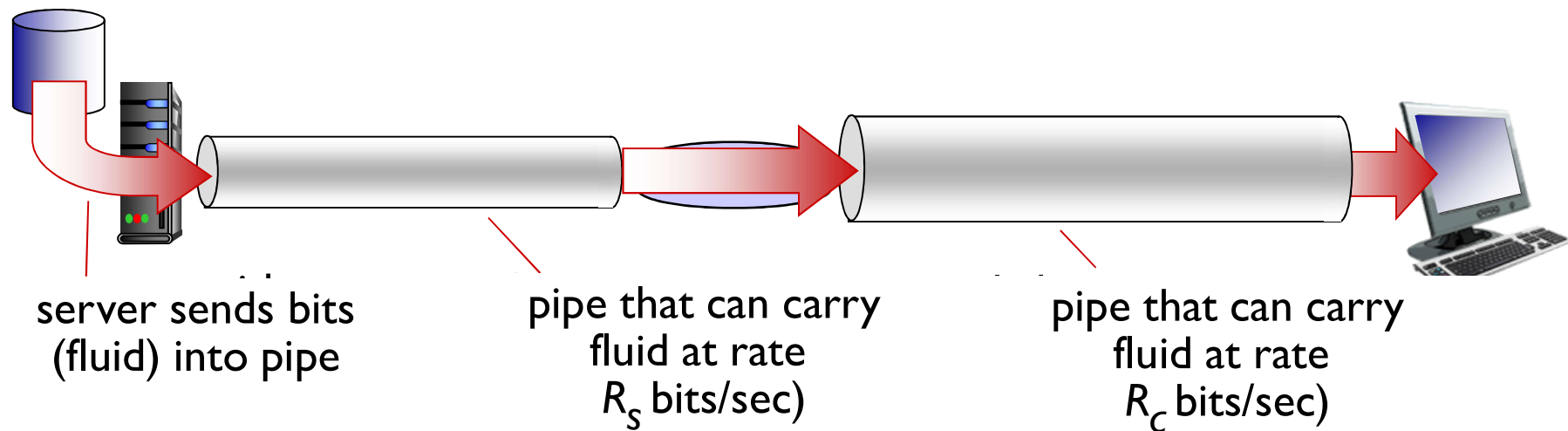
- queue (aka buffer) preceding link in buffer has finite capacity
- packet arriving to full queue dropped (aka lost)
- lost packet may be retransmitted by previous node, by source end system, or not at all

\* aka : also known as



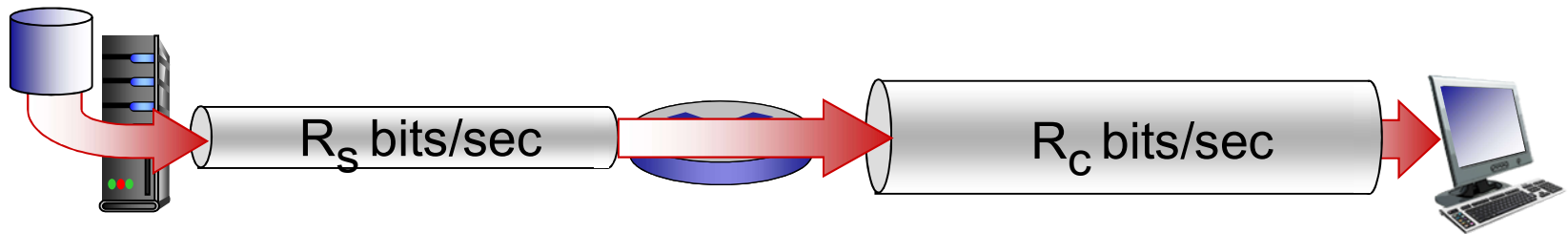
# Throughput

- *throughput*: rate (bits/time unit) at which bits transferred between sender/receiver
  - *instantaneous*: rate at given point in time
  - *average*: rate over longer period of time

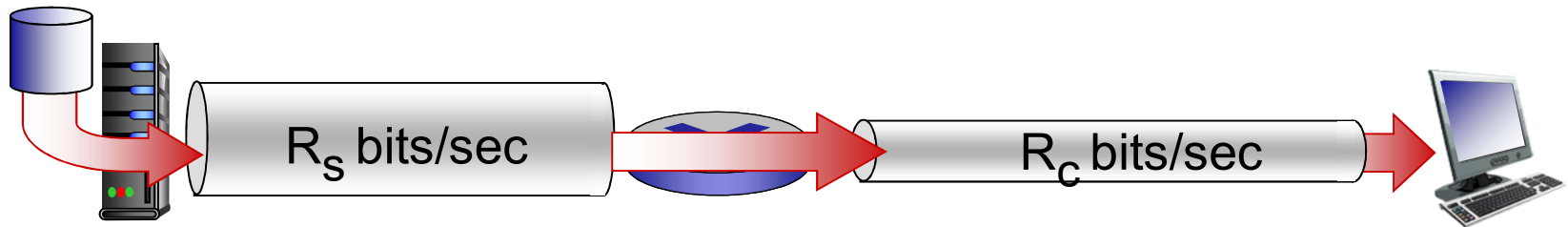


# Throughput (more)

- $R_s < R_c$  What is average end-end throughput?



- $R_s > R_c$  What is average end-end throughput?

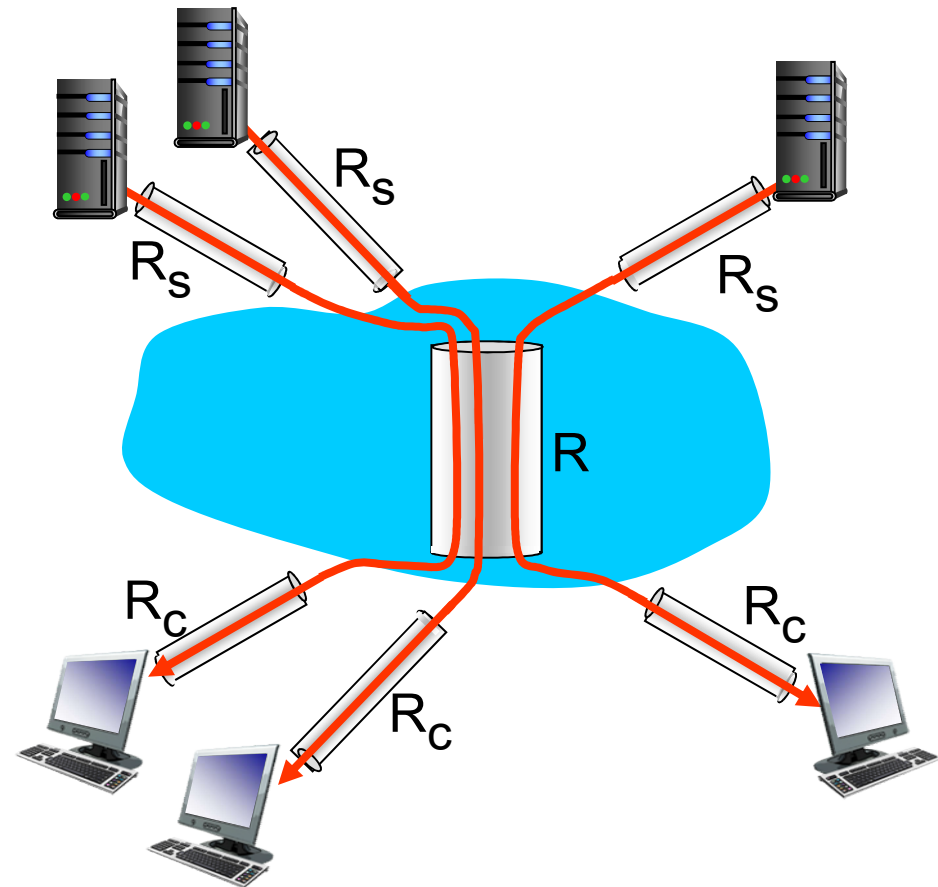


*bottleneck link*

link on end-end path that constrains end-end throughput

# Throughput: Internet scenario

- per-connection end-end throughput:  
 $\min(R_c, R_s, R/I)$
- in practice:  $R_c$  or  $R_s$  is often bottleneck



10 connections (fairly) share  
backbone bottleneck link  $R$  bits/sec



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