Overview of ARM Architecture & Cortex-M Processors

Lecture 3

Yeongpil Cho

Hanyang University

Topics

- ARM Architecture
 - Overview
- Cortex-M Processors
 - Overview
 - Programmer's model

ARM Architecture Overview

ARM

ARM Powered Products



39u10 The ARM Architecture

ARM

- ARM is best known for its range of RISC processor core designs
 - Other products fabric IP, software tools, models, cell libraries to help partners develop and ship ARM-based SoCs (System-on-Chip)
- ARM does not manufacture silicon
 - Licensed to partners to develop and fabricate new microcontrollers
 - Snapdragon series of Qualcommm
 - A series of Apple
 - Exynos series of Samsung
 - Tegra series of Nvidia

– ...

ARM's IPs (Intellectual Properties) other than processors

Graphic IP

- ARM Mali-G71
- Mali-DP550 (Display processor)
- Mali-V550 (Video Processor)

Other IP

- ARM CoreLink CCI-550 (Cache Coherent Interconnect) • ARM Compiler 6
- CoreLink GIC-500 (Interrupt Controller)
- CoreLink MMU-500 (System Memory Management Unit)
- CoreLink TZC-400 (ARM TrustZone® Controller)
- CoreLink DMC-500/DMC-520 (Dynamic Memory Controller)
- ARM CoreSight M SoC-400 (Debug and Trace)

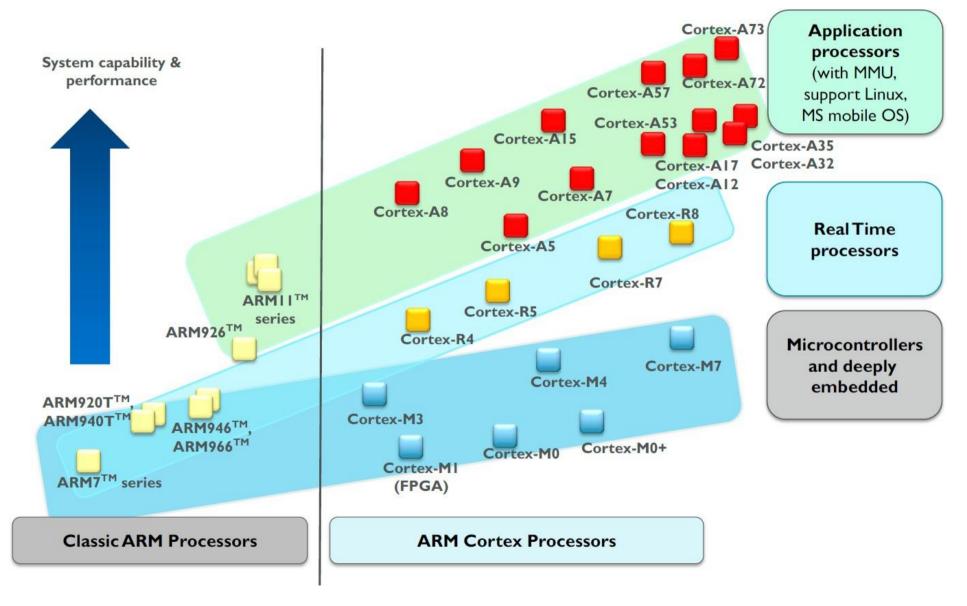
Tools

- ARM Development Studio
- Fixed Virtual Platforms
- ARM Fast Models
- ARM Versatile Express

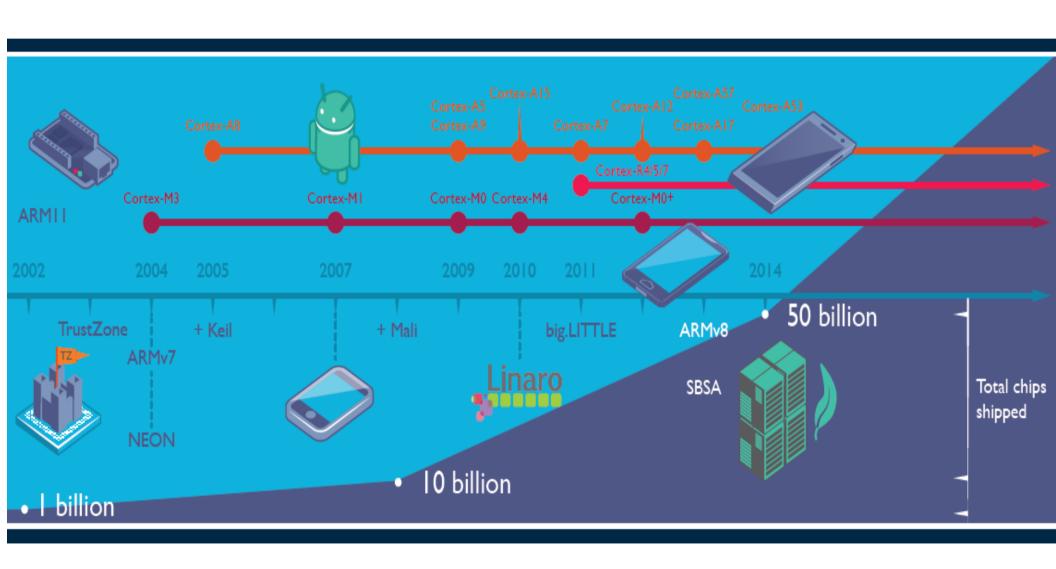
ARM Architecture

- Based upon RISC (Reduced Instruction Set) Architecture with enhancements to meet requirements of embedded applications
 - Fixed length instructions
 - Load-store architecture
 - Memory-to-register load instructions
 - Register-to-memory store instructions
 - A large uniform register file
 - 32-bit architecture (v1-v7), 64-bit architecture (v8-v9)
 - Good speed/power

ARM Processor Family



Release dates of ARM cores



ARM Cortex Processors

- Cortex-A series (Application)
 - High performance processors capable of full Operating System
 (OS) support
 - Applications include smartphones, digital TV, smart books
- Cortex-R series (Real-time)
 - High performance and reliability for real-time applications
 - Applications include automotive braking system, powertrains
- Cortex-M series (Microcontroller)
 - Cost effectiveness and reliability for real-time applications
 - Applications include tiny IoT devices, smart sensors

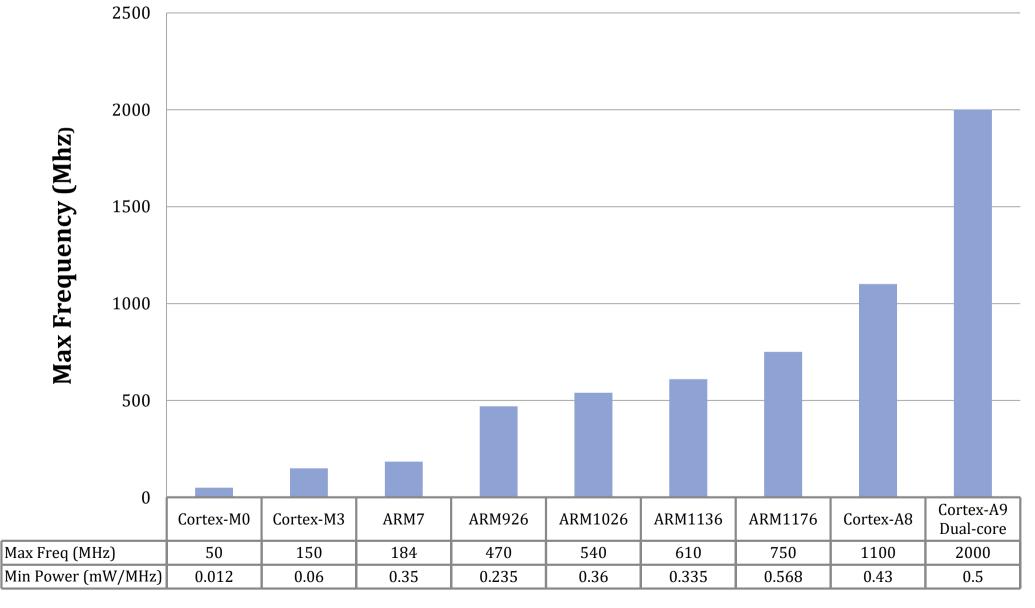
Summary of Processor Characteristics

	Application processors	Real-time processors	Microcontroller		
			processors		
Design	High clock frequency,	High clock frequency,	Short pipeline,		
	Long pipeline,	Long to medium	ultra low power,		
	High performance,	pipeline length,	Deterministic (low		
	Multimedia support (NEON	Deterministic (low	interrupt latency)		
	instruction set extension)	interrupt latency)			
System	Memory Management Unit	Memory Protection	Memory Protection Unit		
features	(MMU), virtualizing memor	γ Unit (MPU), cache	(MPU), Nested Vectored		
	cache memory,	memory, Tightly	Interrupt Controller		
	ARM TrustZone® security	Coupled Memory	(NVIC), Wakeup Interrupt		
	extension	(TCM)	Controller (WIC)		
Targeted	Mobile computing, smart	Industrial	Microcontrollers,		
markets	phones,	microcontrollers,	Deeply embedded		
	energy-efficient servers,	automotives,	systems (e.g. sensors,		
	high-end microprocessors	Hard disk controllers,	MEMS, mixed signal IC),		
		Baseband modem	Internet of Things (IoT)		

11

Relative Performance*





^{*}Represents attainable speeds in 130, 90, 65, or 45nm processes

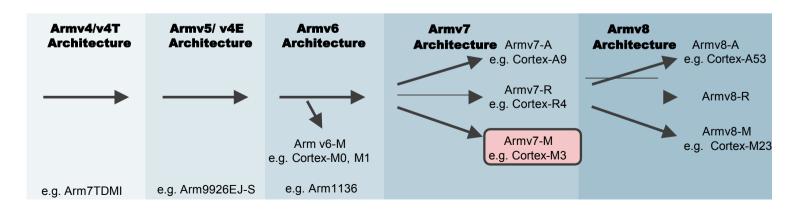
ARM processors vs. ARM architectures

ARM architecture

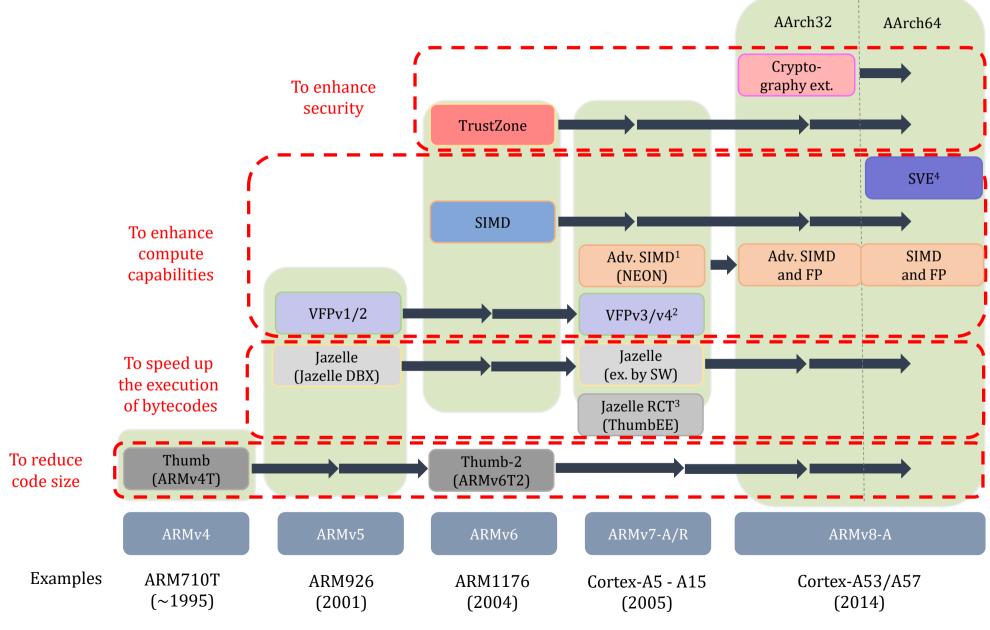
- Describes the details of instruction set, programmer's model, exception model, and memory map
- Documented in the Architecture Reference Manual

ARM processor

- Developed using one of the Arm architectures
- More implementation details, such as timing information
- Documented in processor's Technical Reference Manual

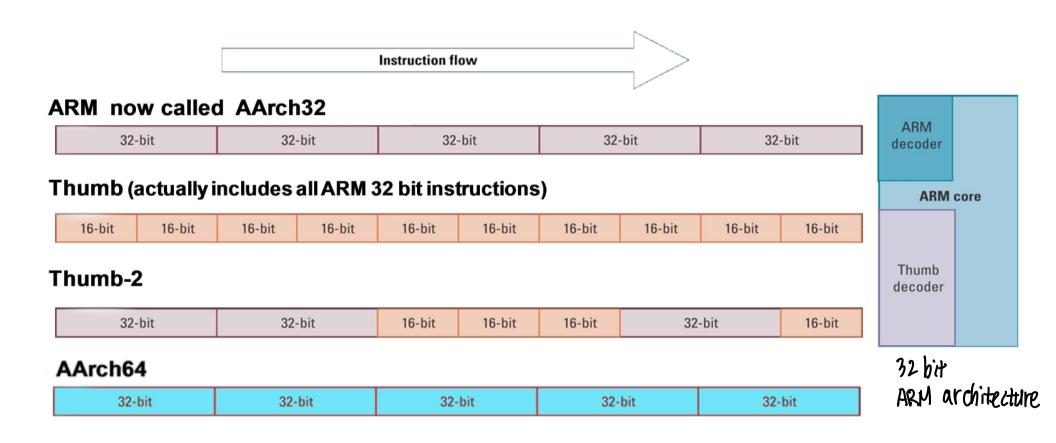


Functional Evolutions by Architectures



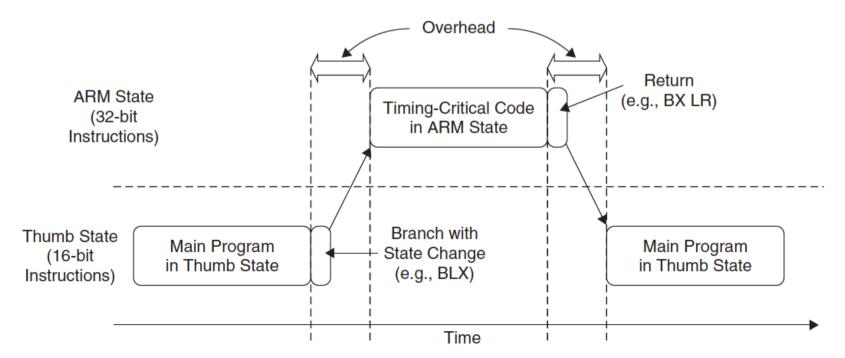
Remarks: See on the next slide.

Instruction Sets

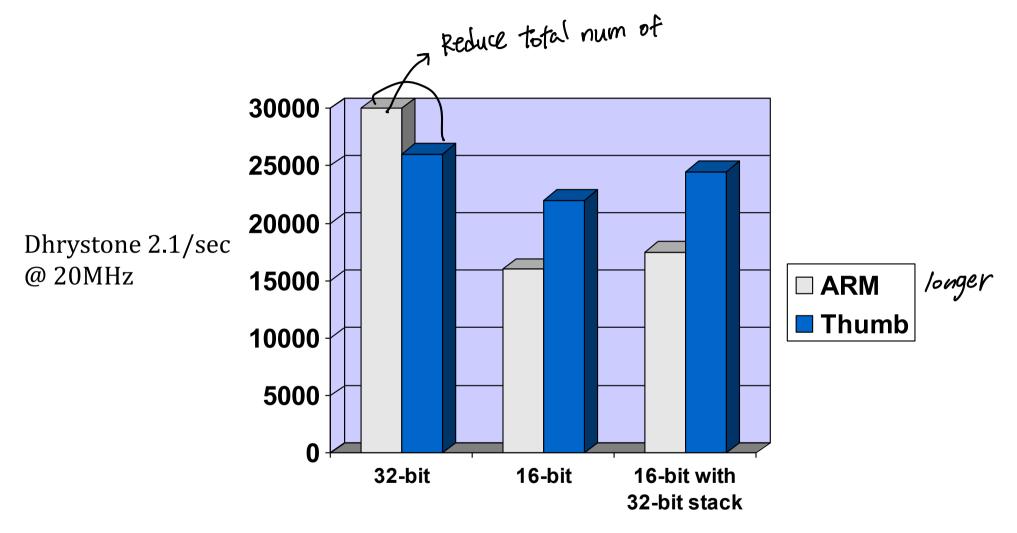


Instruction State Switches

- The current instruction state (between ARM and Thumb/Thumb-2) is determined depending on the LSB (Least Significant bit) of the branch address. (branches: function calls, jumps, ...)
 - LSB = $0 \rightarrow ARM$ instruction set
 - LSB = $1 \rightarrow$ Thumb/Thumb-2 instruction set



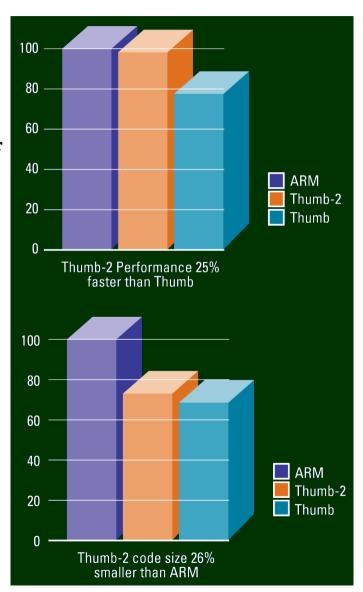
ARM and Thumb Performance



Memory width (zero wait state)

The Thumb-2 instruction set

- Variable-length instructions
 - ARM instructions are a fixed length of 32 bits
 - Thumb instructions are a fixed length of 16 bits
 - Thumb-2 instructions can be either 16-bit or 32-bit
- Thumb-2 gives approximately 25% improvement in performance over Thumb
- Thumb-2 gives approximately 26% reduction in code size over ARM



Cortex-M Processor Overview

Cortex-M Processor Family

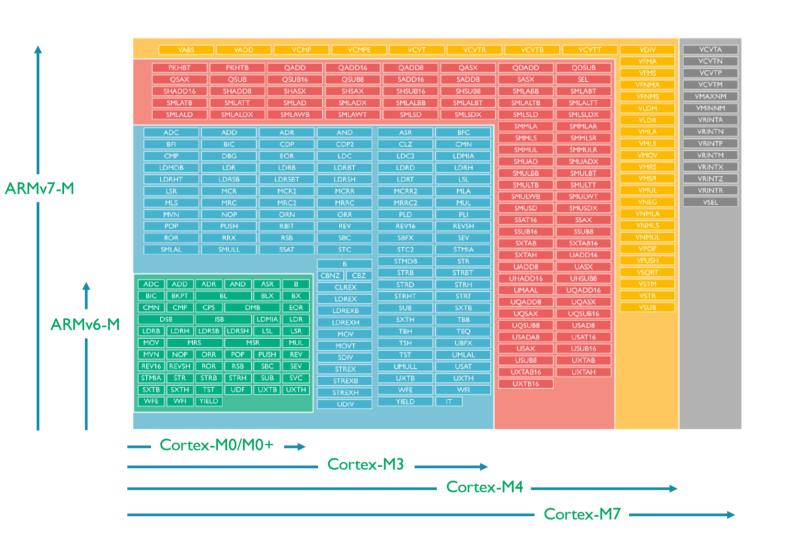
Processor	Descriptions
Cortex-M0	A very small processor (starting from 12K gates) for low cost, ultra low power microcontrollers and deeply embedded applications
	The most energy-efficient processor for small embedded system. Similar size and programmer's
Cortex-M0+	model to the Cortex-M0 processor, but with additional features like single cycle I/O interface and
	vector table relocations
Cortex-M1	A small processor design optimized for FPGA designs and provides Tightly Coupled Memory
Cortex-M1	(TCM) implementation using memory blocks on the FPGAs. Same instruction set as the Cortex-M0
	A small but powerful embedded processor for low-power microcontrollers that has a rich
Cortex-M3	instruction set to enable it to handle complex tasks quicker. It has a hardware divider and Multiply-
GOI LEX 1VIS	Accumulate (MAC) instructions. In addition, it also has comprehensive debug and trace features to
	enable software developers to develop their applications quicker
	It provides all the features on the Cortex-M3, with additional instructions target at Digital Signal
Cortex-M4	Processing (DSP) tasks, such as Single Instruction Multiple Data (SIMD) and faster single cycle
	MAC operations. In addition, it also have an optional single precision floating point unit that
	support IEEE 754 floating point standard
	High-performance processor for high-end microcontrollers and processing intensive
Cortex-M7	applications. It has all the ISA features available in Cortex-M4, with additional support for double-
	precision floating point, as well as additional memory features like cache and Tightly Coupled Memory (TCM)
	A small processor for ultra-low power and low cost designs, similar to the Cortex-M0+ processor,
Cortex-M23	but with various enhancements in instruction set and system-level features. It also supports the
Gortex 1425	TrustZone security extension.
	A mainstream processor design, similar to previous Cortex-M3 and Cortex-M4 processors, but
Cortex-M33	with much better flexibility in system design, better energy efficiency and higher performance. It
	also supports the TrustZone security extension .

ARM Cortex-M series family

Digital Gignal

Processor	Arm Architecture	Core Architecture	Thumb®	Thumb®-2	Hardware Multiply	Hardware Divide	Saturated Math	DSP Extensions	Floating Point
Cortex-M0	Armv6-M	Von Neumann	Most	Subset	1 or 32 cycle	No	No	No	No
Cortex-M0+	Armv6-M	Von Neumann	Most	Subset	1 or 32 cycle	No	No	No	No
Cortex-M3	Armv7-M	Harvard	Entire	Entire	1 cycle	Yes	Yes	No	No
Cortex-M4	Armv7E-M	Harvard	Entire	Entire	1 cycle	Yes	Yes	Yes	Optional
Cortex-M7	Armv7E-M	Harvard	Entire	Entire	1 cycle	Yes	Yes	Yes	Optional
Cortex-M23, 33	Armv8-M	Harvard	Entire	Entire	1 cycle	Yes	Yes	Yes	Optional

Instruction Set support in the Cortex-M processor



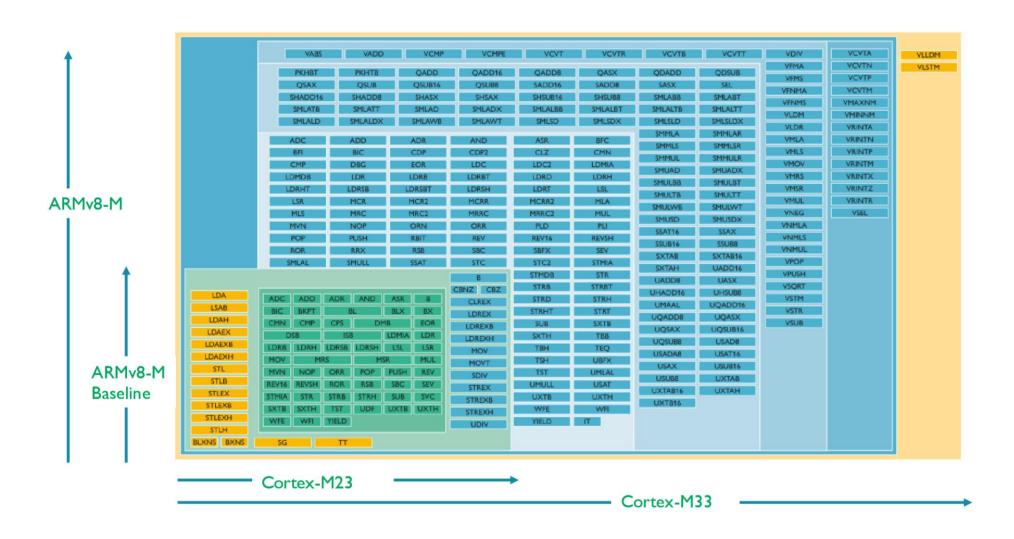
Floating Point

DSP (SIMD, fast MAC)

Advanced data processing bit field manipulations

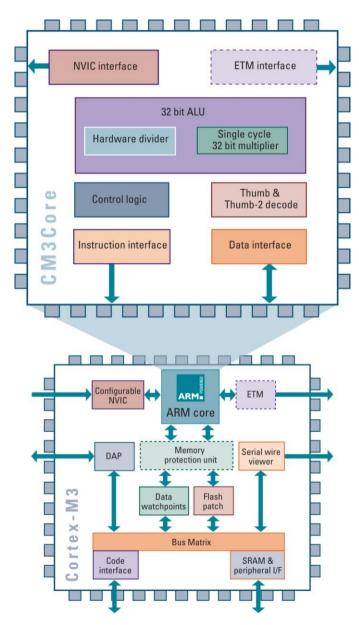
General data processing I/O control tasks

Instruction Set support in the Cortex-M processor

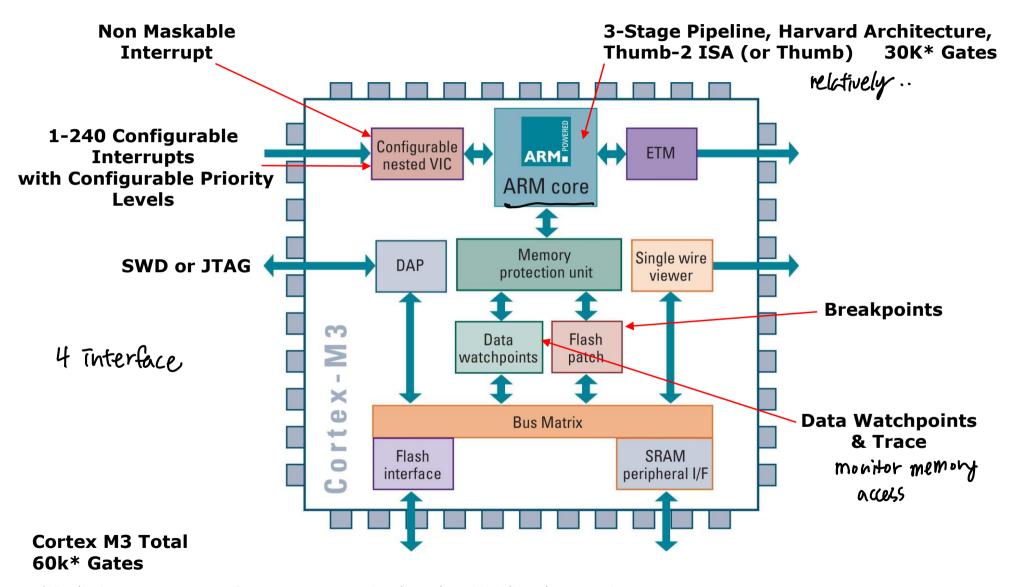


Cortex-M3 Processor(1/2)

- Hierarchical processor integrating core and advanced system peripherals
- Cortex-M3 Processor (By chip manufacturers)
 - Cortex-M3 core
 - Configurable interrupt controller
 - NVIC: Nested Vectored Interrupt Controller
 - Bus matrix (ICode, Dcode, System Bus)
 - Advanced debug components(ETM...)
 - Optional MPU
- Cortex-M3 core (By ARM)
 - Harvard architecture
 - 3-stage pipeline prediction
 - Thumb®-2
 - ALU w. H/W divide and single cycle multiply



Cortex-M3 Processor(2/2)



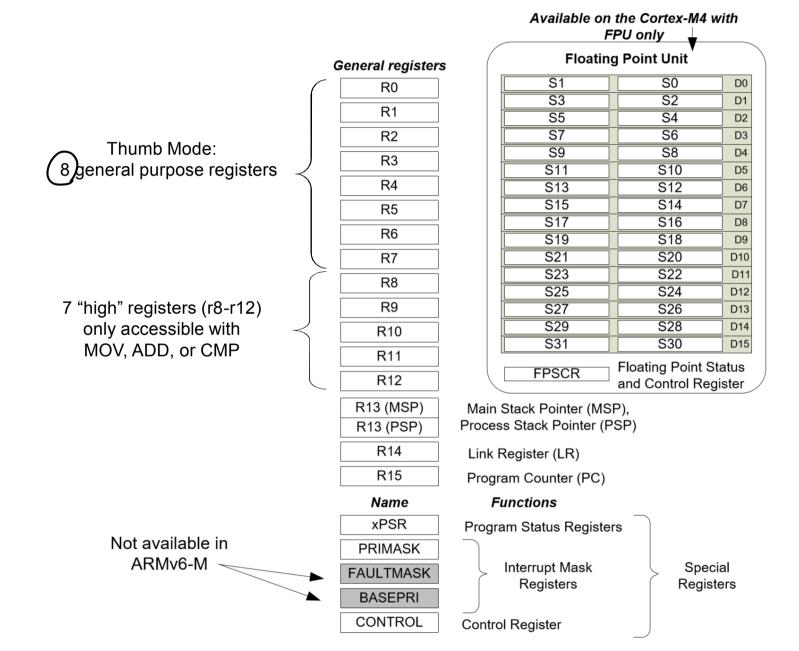
^{*} Preliminary gate counts & power consumption based on initial implementation Gate Counts are based on TSMC 0.18 at 50MHz Optional ETM & MPU gate counts not included

Programmer's Model

Cortex-M Programmer's Model

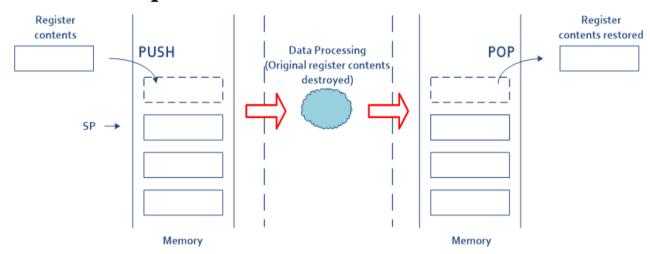
- Fully programmable in C
 - Assembly is not mandatory.
- Stack-based exception model
 - All corruptible registers automatically pushed onto stack upon exceptions in hardware
- Thumb & Thumb-2 instructions sets
- Only two processor modes
 - Thread Mode for User tasks
 - Handler Mode for exceptions
- Vector table contains addresses of exception handlers

Registers of Cortex-M Processors



Stack Pointers

R13 is the stack pointer.



- Two stack pointers are banked so that only one is visible at a time.
- The two stack pointers are:
 - Main Stack Pointer (MSP): This is the default stack pointer, used by the OS kernel, exception handlers, and privileged-mode programs
 - Process Stack Pointer (PSP): Used by the user application code

Program Counter

- R15 is the program counter (PC).
- When you read this register you will find that the value is different than the location of the executing instruction, due to the pipelining of the processor
 - Example:

```
0x1000: MOV RO, PC ; R0 = 0x1004
```

- When writing to the PC, it will cause a branch.
 - the LSB must be set to 1 to indicate the Thumb state operations (setting to 0 implies to switch to the ARM state, which will result in a fault exception in Cortex-M3)
- When reading the PC, the LSB (bit 0) is always 0.
 - Instructions are always aligned in 16-bit or 32-bit

Link Register

- •R14 is the link register (LR).
- LR is used to store the return program counter when a subroutine or function is called.
- e.g., when using the BL (Branch with Link) instruction:

```
main: ; Main program
...
BL foo ; Call foo using Branch with Link instruction:
; PC = foo
; LR = the next instruction in main
...
foo:
... ; Program code for foo
BX LR ; Return
```

Special Registers

- The special registers in the Cortex-M3 processor include:
 - Program Status Registers (PSRs)
 - Interrupt Mask Registers (PRIMASK, FAULTMASK, and BASEPRI)
 - Control Register (CONTROL)
- Can only be accessed via MSR and MRS instructions

```
E.g.,MRS < reg >, < special_reg > ; Read special registerMSR < special_reg >, < reg > ; write to special register
```

 Note: MSR and MRS cannot have memory addresses, only registers are allowed

Program Status Registers (PSRs)

- The program status registers are subdivided into three status registers:
- Application/Interrupt/Execution PSR (A/I/EPSR)

	31	30	29	28	27	26:25	24	23:20	19:16	15:10	9	8	7	6	5	4:0
APSR	Ν	Z	С	٧	Q											
IPSR	Exception Number															
EPSR						ICI/IT	Т			ICI/IT						

 When they are accessed as a collective item, the name xPSR is used (PSR used in program codes).

	31	30	29	28	27	26:25	24	23:20	19:16	15:10	9	8	7	6	5	4:0
xPSR	N	Z	O	٧	Q	ICI/IT	Т			ICI/IT		Exc	eptic	n Nu	ımbe	r

Program Status Registers (PSRs)

EPSR and IPSR are read-only:

```
    MRS r0, APSR ; Read Flag state into R0
    MRS r0, IPSR ; Read Exception/Interrupt state
    MRS r0, EPSR ; Read Execution state
    MSR APSR, r0 ; Write Flag state
```

Accessing xPSR:

```
MRS r0, PSR ; Read the combined program status word
MSR PSR, r0 ; Write combined program state word
```

Program Status Registers (PSRs)

Bit	Description
N	Negative
Z	Zero
С	Carry/borrow
V	Overflow
Q	Sticky saturation flag
ICI/IT	Interrupt-Continuable Instruction (ICI) bits IF-THEN instruction status bit
Т	Thumb state, always 1; trying to clear this bit will cause a fault exception
Exception Number	Indicates which exception the processor is handling

cument exception

Bit Fields in Cortex-M3 Program Status Registers

PRIMASK, FAULTMASK and BASEPRI Registers

- These registers are used to disable exceptions.
 - NMI: Non-Maskable Interrupt

Register Name	Description
PRIMASK	A 1-bit register. When this is set, it allows NMI and the hard fault exception; all other interrupts and exceptions are masked; default is 0 (no masking)
FAULTMASK	A 1-bit register. When this is set, it allows only the NMI, and all interrupts and fault handling exceptions are disabled; default is 0
BASEPRI	A register of up to 9 bits. It defines the masking priority level. When this is set, it disables all interrupts of the same or lower level (larger priority value); default is 0

Cortex-M3 Interrupt Mask Registers

PRIMASK, FAULTMASK and BASEPRI Registers

- To access the PRIMASK, FAULTMASK, and BASEPRI registers, the MRS and MSR instructions are used.
- Example:

```
MRS r0, BASEPRI ; Read BASEPRI register into R0
MRS r0, PRIMASK ; Read PRIMASK register into R0
MRS r0, FAULTMASK ; Read FAULTMASK register into R0
MSR BASEPRI, r0 ; Write R0 into BASEPRI register
MSR PRIMASK, r0 ; Write R0 into FAULTMASK register

MSR FAULTMASK, r0 ; Write R0 into FAULTMASK register
```

- PRIMASK and BASEPRI are useful for temporarily disabling interrupts in timing-critical tasks
- FAULTMASK is used by the OS kernel which cleans up a crashed task
- The PRIMASK, FAULTMASK, and BASEPRI registers cannot be set in the user access level.

The Control Register

• The Control register is used to define the privilege level and the stack pointer selection. This register has two bits.

Bit	Function
CONTROL[1]	Stack status:
	1 = Alternate stack is used
	0 = Default stack (MSP) is used
	If it is in the Thread or base level, the alternate stack is the PSP. There is no alternate stack
	for handler mode, so this bit must be zero when the processor is in handler mode.
CONTROL[0]	0 = Privileged in Thread mode
	1 = User state in Thread mode
	If in handler mode (not Thread mode), the processor operates in privileged mode.

The Control Register

CONTROL[1]

- In Cortex-M3, the CONTROL[1] bit is always 0 (MSP) in handler mode.
- However, in the Thread mode, it can be either 0 or 1.
- This bit is writable only when the core is in Thread mode and privileged.

• CONTROL[0]

The CONTRL[0] bit is writable only in privileged level.

To access the Control register, the MRS and MSR instructions are used:

```
MRS r0, CONTROL ; Read CONTROL register into R0 MSR CONTROL, r0 ; Write R0 into CONTROL register
```

Operation Mode

- Two modes and two privilege levels.
- The operation modes determine whether the processor is running a normal program or running an exception handler.

	Privileged Level	Unprivileged(User) Level
When running an exception	Handler Mode	
When running main program	Thread Mode	Thread Mode

Operation Modes and Privilege Levels in Cortex-M3

Processor Mode

Handler Mode

- Used to handle exceptions.
- The processor returns to Thread mode when it has finished exception processing.
- In Handler mode, software execution is always privileged.

Thread Mode

- Used to execute application software.
- The processor enters Thread mode when it comes out of reset.
- In Thread mode, the CONTROL register controls whether software execution is privileged or unprivileged, see CONTROL register.

Privilege Levels

Unprivileged (User)

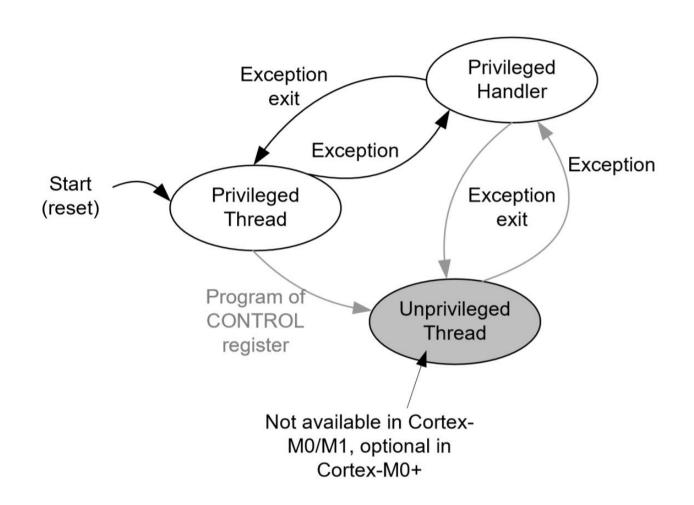
The software:

- has limited access to the MSR (Move Register to Special Register) and MRS (Move Special Register to Register) instructions, and cannot use the CPS (Change Processor State) instruction
- cannot access the system timer, NVIC, or system control block
- might have restricted access to memory or peripherals.
- Unprivileged software executes at the unprivileged level

Privileged

 The software can use all the instructions and has access to all resources.

Transition of Operation Mode



Vector Tables

- The vector table is an array of word data, with each representing the starting address of the handler for one exception/interrupt type.
 - In the Cortex-M3, vector addresses in the vector table should have their LSB set to 1 to indicate that they are Thumb code.
- The base address of the vector table is relocatable (set the relocation register in the NVIC); initially, the base address is 0x0.
- Example:
 - The reset is exception type 1. The address of the reset vector is 1 times 4, which equals 0x00000004; and NMI vector (type 2) is located in 2 * 4 = 0x00000008
 - The word stored at address 0x00000000 is used as the starting value for the MSP.

Vector Tables

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•
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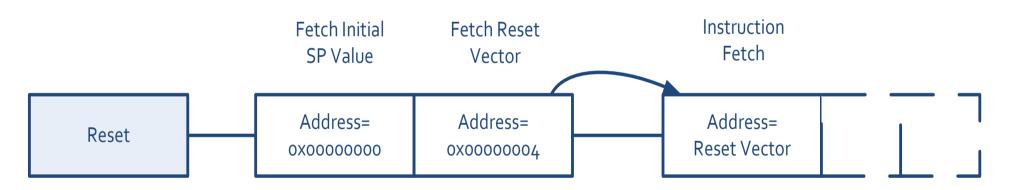
Address Offset	Exception Type	Exception Vector
0x3C	15	SYSTICK
0x38	14	PendSV
0x34	13	Reserved
0x30	12	Debug Monitor
0x2C	11	SVC
0x1C-0x28	7-10	Reserved
0x18	6	Usage fault
0x14	5	Bus fault
0x10	4	MemManage fault
0x0C	3	Hard fault
0x08	2	NMI
0x04	1	Reset
0x00	0	Starting value of the MSP

Vector Table Definition After Reset

debut Stack of pointer

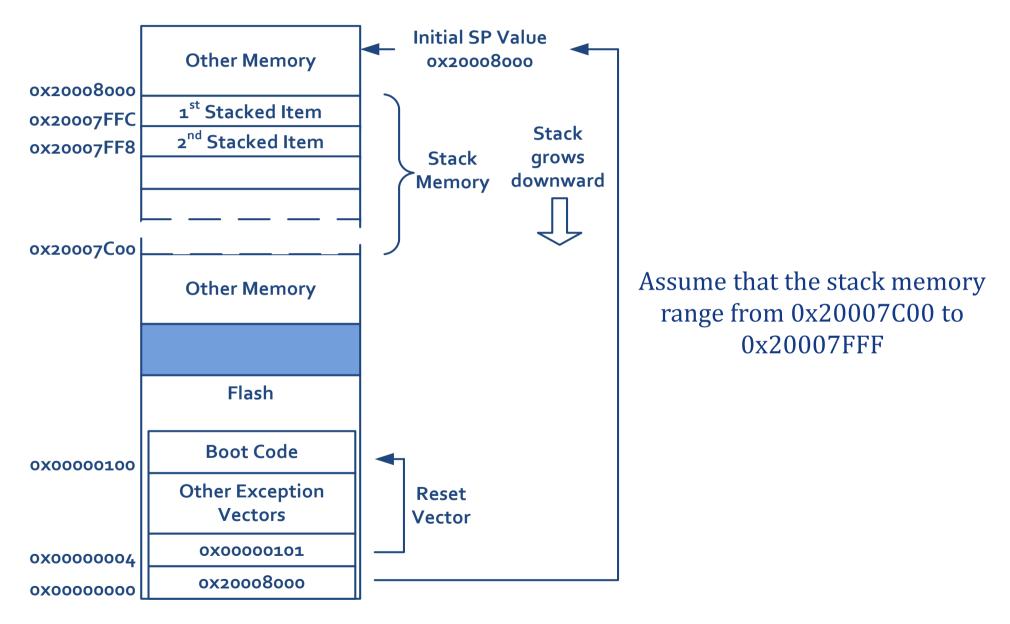
Reset Sequence

- After the processor exits reset, it will read two words from memory:
 - Address 0x00000000: default value of R13 (MSP)
 - Address 0x00000004: Reset vector (the starting address of startup program; LSB should be set to 1 to indicate Thumb state)



Time

Reset Sequence



Initial Stack Pointer Value and Initial Program Counter (PC) Value Example