Communicators and Groups

- A group is an ordered set of processes
- Each process in a group is associated with a unique integer rank
- Rank values start at zero and go to N-1, where N is the number of processes in the group
- In MPI, a group is represented within system memory as an object. It is accessible to the programmer only by a "handle"
- A group is always associated with a communicator object

- A communicator encompasses a group of processes that may communicate with each other
- All MPI messages must specify a communicator
- In the simplest sense, the communicator is an extra "tag" that must be included with MPI calls
- Like groups, communicators are represented within system memory as objects and are accessible to the programmer only by "handles"
- For example, the handle for the communicator that comprises all tasks is MPI_COMM_WORLD

Purpose of Group and Communicators

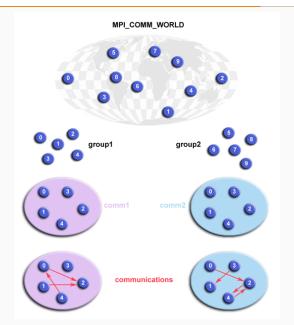
- Allow you to organize tasks, based upon function, into task groups
- Enable Collective Communications operations across a subset of related tasks
- Provide basis for implementing user defined virtual topologies
- Provide for safe communications

Programming Considerations and Restrictions

- Groups/communicators are dynamic they can be created and destroyed during program execution
- Processes may be in more than one group/communicator.
 They will have a unique rank within each group/communicator
- MPI provides over 40 routines related to groups, communicators, and virtual topologies

Usage of Groups and Communicators

- Extract handle of global group from MPI_COMM_WORLD using MPI_Comm_group
- Form new group as a subset of global group using MPI_Group_incl
- Create new communicator for new group using MPI_Comm_create
- Determine new rank in new communicator using MPI_Comm_rank
- Conduct communications using any MPI message passing routine
- When finished, free up new communicator and group (optional) using MPI_Comm_free and MPI_Group_free



```
#include "mpi.h"
      #include <stdio.h>
      #define NPROCS 8
 4
 5
      main(int argc, char *argv[])
      int rank, new rank, sendbuf, recvbuf, numtasks,
 6
                 ranks1[4]=\{0,1,2,3\}, ranks2[4]=\{4,5,6,7\};
      MPI Group orig group, new group; // required variables
 8
      MPI Comm new comm; // required variable
9
10
11
      MPI Init(&argc,&argv);
12
      MPI Comm rank (MPI COMM WORLD, &rank);
13
      MPI Comm size (MPI COMM WORLD, &numtasks);
14
15
      if (numtasks != NPROCS) {
        printf("Must specify MP PROCS= %d. Terminating.\n", NPROCS);
16
        MPI Finalize();
17
18
        exit(0);
19
20
21
      sendbuf = rank;
22
```

```
22
23
      // extract the original group handle
24
      MPI Comm group (MPI COMM WORLD, &orig group);
25
26
      // divide tasks into two distinct groups based upon rank
27
      if (rank < NPROCS/2) {
28
        MPI Group incl (orig group, NPROCS/2, ranks1, &new group);
29
30
      else {
31
        MPI Group incl (orig group, NPROCS/2, ranks2, &new group);
32
33
34
      // create new new communicator and then perform collective communicati
35
      MPI Comm create (MPI COMM WORLD, new group, &new comm);
36
      MPI Allreduce (&sendbuf, &recvbuf, 1, MPI INT, MPI SUM, new comm);
37
38
      // get rank in new group
39
      MPI Group rank (new group, &new rank);
40
      printf("rank= %d newrank= %d recvbuf= %d\n",rank,new rank,recvbuf);
41
42
      MPI Finalize();
43
```

Sample program output:

```
rank= 7 newrank= 3 recvbuf= 22
rank= 0 newrank= 0 recvbuf= 6
rank= 1 newrank= 1 recvbuf= 6
rank= 2 newrank= 2 recvbuf= 6
rank= 6 newrank= 2 recvbuf= 22
rank= 3 newrank= 3 recvbuf= 6
rank= 4 newrank= 0 recvbuf= 22
rank= 5 newrank= 1 recvbuf= 22
```

Virtual Topology

What Are They?

- In terms of MPI, a virtual topology describes a mapping/ordering of MPI processes into a geometric "shape"
- The two main types of topologies supported by MPI are Cartesian (grid) and Graph
- MPI topologies are virtual there may be no relation between the physical structure of the parallel machine and the process topology
- Virtual topologies are built upon MPI communicators and groups
- Must be "programmed" by the application developer

Virtual Topology

Why Use Them?

- Convenience
 - Virtual topologies may be useful for applications with specific communication patterns - patterns that match an MPI topology structure
 - For example, a Cartesian topology might prove convenient for an application that requires 4-way nearest neighbor communications for grid based data

Virtual Topology

Why use them?

- Communication Efficiency
 - Some hardware architectures may impose penalties for communications between successively distant "nodes"
 - A particular implementation may optimize process mapping based upon the physical characteristics of a given parallel machine
 - The mapping of processes into an MPI virtual topology is dependent upon the MPI implementation, and may be totally ignored

Virtual Topology: Example

A simplified mapping of processes into a Cartesian virtual topology appears below:

| 0 | 1 (0,1) | 2 | 3 |
|-------------|----------|----------|-------------|
| (0,0) | | (0,2) | (0,3) |
| 4 | 5 | 6 | 7 |
| (1,0) | (1,1) | (1,2) | (1,3) |
| 8 | 9 | 10 | 11 |
| (2,0) | (2,1) | (2,2) | (2,3) |
| 12 (3,0) | 13 (3,1) | 14 (3,2) | 15 (3,3) |

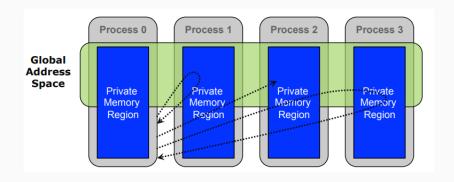
Virtual Topology: Example Code

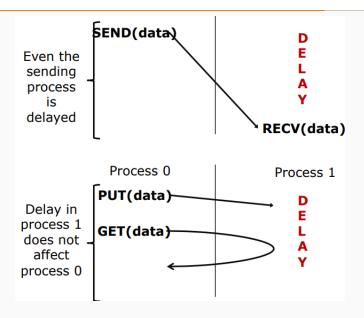
```
#include "mpi.h"
      #include <stdio.h>
 3
      #define SIZE 16
      #define UP
 4
      #define DOWN 1
6
      #define LEFT 2
      #define RIGHT 3
      main(int argc, char *argv[]) {
10
      int numtasks, rank, source, dest, outbuf, i, tag=1,
         inbuf[4]={MPI PROC NULL,MPI PROC NULL,MPI PROC NULL,MPI PROC NULL,},
         nbrs[4], dims[2]={4,4},
12
         periods[2]={0.0}, reorder=0, coords[2];
13
14
15
      MPI Request reqs[8];
16
      MPI Status stats[8];
      MPI Comm cartcomm; // required variable
17
18
19
      MPI Init(&argc,&argv);
20
      MPI Comm size (MPI COMM WORLD, &numtasks);
21
22
      if (numtasks == SIZE) {
         // create cartesian virtual topology, get rank, coordinates, neighbor rank
23
         MPI Cart create (MPI COMM WORLD, 2, dims, periods, reorder, &cartcomm);
24
         MPI Comm rank (cartcomm, &rank);
25
26
         MPI Cart coords (cartcomm, rank, 2, coords);
         MPI Cart shift(cartcomm, 0, 1, &nbrs[UP], &nbrs[DOWN]);
27
         MPI Cart shift(cartcomm, 1, 1, &nbrs[LEFT], &nbrs[RIGHT]);
28
29
```

Virtual Topology: Example Code

```
29
30
         printf("rank= %d coords= %d %d neighbors(u,d,l,r)= %d %d %d %d\n",
                rank, coords[0], coords[1], nbrs[UP], nbrs[DOWN], nbrs[LEFT],
31
32
                nbrs[RIGHT]);
33
34
         outbuf = rank;
35
36
         // exchange data (rank) with 4 neighbors
37
         for (i=0; i<4; i++) {
38
            dest = nbrs[i];
39
            source = nbrs[i];
40
            MPI Isend(&outbuf, 1, MPI INT, dest, tag,
41
                       MPI COMM WORLD, &regs[i]);
            MPI Irecv(&inbuf[i], 1, MPI_INT, source, tag,
42
                       MPI COMM WORLD, &regs[i+4]);
43
44
45
46
         MPI Waitall(8, regs, stats);
47
         printf("rank= %d
48
                                            inbuf(u,d,l,r) = %d %d %d %d n",
49
                rank.inbuf(UP).inbuf(DOWN).inbuf(LEFT).inbuf(RIGHT)); }
50
      else
51
         printf("Must specify %d processors. Terminating.\n", SIZE);
52
53
      MPI Finalize();
54
```

- The basic idea of one-sided communication models is to decouple data movement with process synchronization
 - Should be able to move data without requiring that the remote process synchronize
 - Each process exposes a part of its memory to other processes
 - Other processes can directly read from or write to this memory





Advantages of RMA Operations

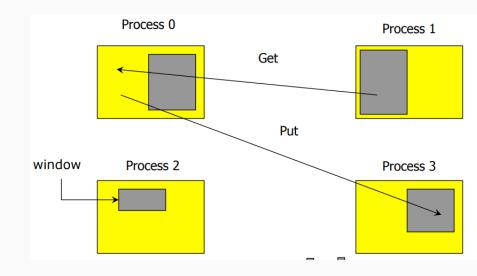
- Can do multiple data transfers with a single synchronization operation
- Bypass tag matching
- Some irregular communication patterns can be more economically expressed
- Can be significantly faster than send/receive on systems with hardware support for remote memory access

What we need to know in MPI RMA

- How to create remote accessible memory?
- Reading, Writing and Updating remote memory
- Data Synchronization
- Memory Model

Creating Public Memory

- Any memory created by a process is, by default, only locally accessible
- Once the memory is created, the user has to make an explicit MPI call to declare a memory region as remotely accessible
 - MPI terminology for remotely accessible memory is a window
- Once a memory region is declared as remotely accessible, all processes in the window object can read/write data to this memory without explicitly synchronizing with the target process



Creating Windows

Use MPI_WIN_CREATE to create windows

- Expose a region of the memory in an RMA window
- Only data exposed in a window can be accessed with RMA ops

int MPI_Win_create(void *base, MPI_int size, int disp_unit, MPI_Info info, MPI_Comm comm, MPI_Win *win)

- base: local data
- size: size of local data
- disp_unit: unit size for displacements
- info: info argument
- win: window object

```
int main(int argc, char ** argv)
{
    int *a; MPI Win win;
    MPI Init(&argc, &argv);
    /* create private memory */
    MPI Alloc mem(1000*sizeof(int), MPI INFO NULL, &a);
    /* use private memory like you normally would */
    a[0] = 1; a[1] = 2;
    /* collectively declare memory as remotely accessible */
    MPI Win create(a, 1000*sizeof(int), sizeof(int),
           MPI INFO NULL, MPI COMM WORLD, &win);
    /* Array 'a' is now accessibly by all processes in
     * MPI COMM WORLD */
   MPI Win free (&win);
   MPI Free mem(a);
   MPI Finalize(); return 0;
                                                              156
```

Window Allocate

Use MPI_WIN_ALLOCATE

Create a remotely accessible memory region in an RMA window

int MPI_Win_allocate(MPI_Aint size, int disp_unit,MPI_Info info,MPI_Comm comm, void *baseptr, MPI_Win *win)

Window allocate example

```
int main(int argc, char ** argv)
    int *a: MPI Win win:
   MPI Init(&argc, &argv);
   /* collectively create remote accessible memory in a window */
   MPI Win allocate (1000*sizeof(int), sizeof(int), MPI INFO NULL,
                     MPI COMM WORLD, &a, &win);
   /* Array 'a' is now accessible from all processes in
     * MPI COMM WORLD */
   MPI Win free (&win);
   MPI Finalize(); return 0;
```

Data Movement between windows

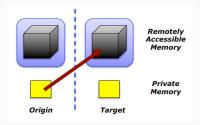
MPI provides ability to read, write and atomically modify data in remotely accessible memory regions

- MPLGET
- MPI_PUT

MPI_PUT: Put data in a window

MPI_Put(void *origin_addr, int origin_count, MPI_Datatype origin_dtype, int target_rank, MPI_Aint target_disp, int target_count, MPI_Datatype target_dtype, MPI_Win win)

- Move data from origin to target
- Separate data description triples for origin and target



MPI_GET: Get data from a window

MPI_Get(void *origin_addr, int origin_count, MPI_Datatype origin_dtype, int target_rank, MPI_Aint target_disp, int target_count, MPI_Datatype target_dtype, MPI_Win win)

Move data to origin, from target

