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Fast Fourier Transformation

Step-1 FFT

$$A(x) = a_0 x^0 + a_1 x^1 + ... + a_{n-1} x^{n-1}$$

Divide it into two polynomials, one with even coefficients(0) and the other with the odd coefficients(1):

$$A_0(x) = a_0 x^0 + a_2 x^1 + ... + a_{n-2} x^{n/2-1}$$

$$A_1(x) = a_1 x^0 + a_3 x^1 + \ldots + a_{n-1} x^{n/2-1}$$

Finally we can write original polynomial as combination of Ao and A1:

$$A(x) = A_0(x^2) + xA_1(x^2)$$

Now writing th original polynomial A(x) in point form for n-th root of unity:

$$y_k = y_k^0 + w_n^k y_k^1, \quad k = 0 \dots n/2 - 1$$

$$\begin{split} y_{k+n/2} &= A(w_n^{k+n/2}) = A_0(w_n^{2k+n}) + w_n^{k+n/2} A_1(w_n^{2k+n}) = A_0(w_n^{2k} w_n^n) A_1(w_n^{2k} w_n^n) \\ &= A(w_n^{2k}) - w_n^k A_1(w_n^{2k}) = y_k^0 - w_n^k y_k^1 \end{split}$$

Finally our **point form** will be:

$$y_k = y_k^0 + w_n^k y_k^1, \quad k = 0...n/2-1,$$

$$y_{k+n/2} = y_k^0 + w_n^k y_k^1, \quad k = 0 \dots n/2 - 1$$

FFT()

```
vector<complex<double> > fft(vector<complex<double> > a){
   int n = (int)a.size();
   if(n <= 1)
        return a;
   //Dividing a0 and a1 as even and odd polynomial of degree n/2
   vector<complex<double> > a0(n/2), a1(n/2);
   for(int i = 0;i<n/2;i++){
        a0[i] = a[2*i];
        a1[i] = a[2*i + 1];
   }</pre>
```



```
//Divide step
   //Recursively calling FFT on polynolmial of degree n/2
   a0 = fft(a0);
   a1 = fft(a1);
   double ang = 2*PI/n;
   //defining w1 and wn
   complex<double> w(1) , wn(cos(ang),sin(ang));
   for(int i = 0; i < n/2; i++){
       //for powers of k <= n/2
       a[i] = a0[i] + w*a1[i];
       //powers of k > n/2
       a[i + n/2] = a0[i] - w*a1[i];
       //Updating value of wk
       w *= wn;
   }
   return a;
}
```

Time Complexity: O(nlogn)

Step 2: Convolution

Multiply()

```
void multiply(vector<int> a, vector<int> b){
   vector<complex<double> > fa(a.begin(),a.end()), fb(b.begin(),b.end());
   int n = 1;
   //resizing value of n as power of 2
   while(n < max(a.size(),b.size()))
        n <<= 1;
   n <<= 1;
   //cout<<n<<endl;
   fa.resize(n);
   fb.resize(n);
   //Calling FFT on polynomial A and B
   //fa and fb denotes the point form
   fa = fft(fa);
   fb = fft(fb);</pre>
```



```
//Convolution step
for(int i = 0;i<n;i++){
    fa[i] = fa[i] * fb[i];
}
//Converitng fa from coefficient back to point form
fa = inv_fft(fa);
return;
}</pre>
```

Time Complexity: O(n) (exluding the time for calculating FFT())

Step 3: Inverse FFT

So, suppose we are given a vector $(y_0, y_1, y_2, ...y_{n-1})$. Now we need to restore it back to coefficient form.

$$\begin{pmatrix} w_n^0 & w_n^0 & w_n^0 & w_n^0 & w_n^0 & \cdots & w_n^0 \\ w_n^0 & w_n^1 & w_n^2 & w_n^3 & \cdots & w_n^{n-1} \\ w_n^0 & w_n^2 & w_n^4 & w_n^6 & \cdots & w_n^{2(n-1)} \\ w_n^0 & w_n^3 & w_n^6 & w_n^9 & \vdots & w_n^{3(n-1)} \\ \vdots & \vdots & \vdots & \vdots & \ddots & \vdots \\ w_n^0 & w_n^{n-1} & w_n^{2(n-1)} & w_n^{3(n-1)} & \cdots & w_n^{(n-1)(n-1)} \end{pmatrix} \begin{pmatrix} a_0 \\ a_1 \\ a_2 \\ a_3 \\ \vdots \\ a_{n-1} \end{pmatrix} = \begin{pmatrix} y_0 \\ y_1 \\ y_2 \\ y_3 \\ \vdots \\ y_{n-1} \end{pmatrix}$$

Then the vector $\{a_0, a_1, a_2, ... a_{n-1}\}$ can be found by multiplying the vector $\{y_0, y_1, y_2, ..., y_{n-1}\}$ by an inverse matrix:

$$\begin{pmatrix} w_n^0 & w_n^0 & w_n^0 & w_n^0 & \cdots & w_n^0 \\ w_n^0 & w_n^1 & w_n^2 & w_n^3 & \cdots & w_n^{n-1} \\ w_n^0 & w_n^2 & w_n^4 & w_n^6 & \cdots & w_n^{2(n-1)} \\ w_n^0 & w_n^3 & w_n^6 & w_n^9 & \vdots & w_n^{3(n-1)} \\ \vdots & \vdots & \vdots & \vdots & \ddots & \vdots \\ w_n^0 & w_n^{n-1} & w_n^{2(n-1)} & w_n^{3(n-1)} & \cdots & w_n^{(n-1)(n-1)} \end{pmatrix} \begin{pmatrix} a_0 \\ a_1 \\ a_2 \\ a_3 \\ \vdots \\ a_{n-1} \end{pmatrix} = \begin{pmatrix} y_0 \\ y_1 \\ y_2 \\ y_3 \\ \vdots \\ y_{n-1} \end{pmatrix}$$

Thus we get the formula:

$$a_k = \frac{1}{n} \sum_{i=0}^{n-1} y_i w_n^{-kj}$$

Comparing it with the formula for y_k:

$$y_k = \sum_{j=0}^{n-1} a_j w_n^{kj},$$



Inverse_FFT()

```
vector<complex<double> > inv_fft(vector<complex<double> > y){
   int n = y.size();
   if(n <= 1)
      return y;
   vector<complex<double> > y0(n/2), y1(n/2);
   for(int i = 0; i < n/2; i++){
      y0[i] = y[2*i];
      y1[i] = y[2*i + 1];
   }
   y0 = inv_fft(y0);
   y1 = inv_fft(y1);
   double ang = 2 * PI/n * -1;
   complex<double> w(1), wn(cos(ang), sin(ang));
   for(int i = 0; i < n/2; i++){
      y[i] = y0[i] + w*y1[i];
      y[i + n/2] = y0[i] - w*y1[i];
      y[i] /= 2;
      y[i + n/2] /= 2;
      w *= wn;
   //each element of the result is divided by 2
   //assuming that the division by 2 will take place at each level of recursion
   //then eventually just turns out that all the elements are divided into n.
   return y;
}
```

Time Complexity: O(nlogn)

Ques: Very Fast Multiplicaiton

http://www.spoj.com/problems/VFMUL/

```
#include<iostream>
#include<vector>
#include<complex>
#include<algorithm>

using namespace std;
```



```
#define PI 3.14159265358979323846
vector<complex<double> > fft(vector<complex<double> > a){
   //for(int i = 0;i<a.size();i++)cout<<a[i]<<" ";cout<<endl;</pre>
   int n = (int)a.size();
   if(n <= 1)
      return a;
   vector<complex<double> > a0(n/2), a1(n/2);
   for(int i = 0; i < n/2; i++){
       a0[i] = a[2*i];
       a1[i] = a[2*i + 1];
       //cout<<n<<" "<<a0[i]<<" "<<a1[i]<<endl;
   }
   a0 = fft(a0);
   a1 = fft(a1);
   double ang = 2*PI/n;
   complex<double> w(1) , wn(cos(ang),sin(ang));
   for(int i = 0; i < n/2; i++){
       a[i] = a0[i] + w*a1[i];
       a[i + n/2] = a0[i] - w*a1[i];
       w *= wn;
       //cout<<a[i]<<" "<<a[i+n/2]<<endl;
   }
   return a;
}
vector<complex<double> > inv_fft(vector<complex<double>>y){
   int n = y.size();
   if(n <= 1)
        return y;
   vector<complex<double> > y0(n/2), y1(n/2);
   for(int i = 0; i < n/2; i++){}
       y0[i] = y[2*i];
       y1[i] = y[2*i + 1];
   y0 = inv_{fft}(y0);
```



```
y1 = inv_fft(y1);
   double ang = 2 * PI/n * -1;
   complex<double> w(1), wn(cos(ang), sin(ang));
   for(int i = 0;i< n/2;i++){
       y[i] = y0[i] + w*y1[i];
       y[i + n/2] = y0[i] - w*y1[i];
       y[i] /= 2;
       y[i + n/2] /= 2;
       w *= wn;
   }
   return y;
}
void multiply(vector<int> a, vector<int> b){
   vector<complex<double> > fa(a.begin(),a.end()), fb(b.begin(),b.end());
   int n = 1;
   while(n < max(a.size(),b.size()))</pre>
         n <<= 1;
   n \ll 1;
   //cout<<n<<endl;</pre>
   fa.resize(n);
   fb.resize(n);
   fa = fft(fa);
   fb = fft(fb);
   for(int i = 0;i<n;i++){</pre>
       fa[i] = fa[i] * fb[i];
       //cout<<fa[i]<<endl;</pre>
   }
   fa = inv_fft(fa);
   vector<int> res(n);
   for(int i = 0;i<n;i++){</pre>
       res[i] = int(fa[i].real() + 0.5);
       //cout<<res[i]<<endl;</pre>
   }
   int carry = 0;
   for(int i = 0;i<n;i++){</pre>
```

```
res[i] = res[i] + carry;
       carry = res[i] / 10;
       res[i] %= 10;
   }
   bool flag = 0;
   for(int i = n-1;i>=0;i-){if(res[i] || flag){printf("%d",res[i]);flag = 1;}}
       if(!flag)printf("0");
       printf("\n");
       return;
}
int main(){
    int n;
    scanf("%d",&n);
    while(n-){
       string x,y;
       vector<int> a,b;
       cin>>x>>y;
       for(int i = 0;i<x.length();i++){</pre>
            a.push_back(int(x[i] - '0'));
       }
       for(int i = 0;i<y.length();i++){</pre>
           b.push_back(int(y[i] - '0'));
       }
       reverse(a.begin(),a.end());
       reverse(b.begin(),b.end());
       multiply(a,b);
   }
   return 0;
}
```



SELF STUDY NOTES



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