

COSS HW-03 Solution Key

1. In direct mapped - main memory is 16 GB, cache with block size 4 KB and 10 bits in the tag. The tag directory size would be?

Answer:

Size of main memory = 16 GB = 2^{34} bytes

Number of bits in physical address = 34 bits

Number of bits in line number

= Number of bits in physical address – (Number of bits in tag + Number of bits in block offset)

= 12 bits

Total number of lines in cache = 2^{12} lines

Tag directory size = Number of tags x Tag size

= Number of lines in cache x Number of bits in tag

= $2^{12} \times 10$ bits

= **5120 bytes**

2. Consider a direct mapped cache having 512 cache lines machine with a byte addressable main memory of 2^{32} bytes divided into blocks of size 32 bytes. The size of the tag field in bits would be?

Answer:

Size of main memory = 2^{32} bytes

Number of bits in physical address = 32 bits

Block size = 32 bytes = 2^5 bytes

Number of bits in block offset = 5 bits

Total number of lines in cache = 512 lines = 2^9 lines

Number of bits in line number = 9 bits

No of bits in tag = No of bits in physical address – (No of bits in line number + No of bits in block offset)

= **18 bits**

3. In 8-way set associative mapped cache of size 512 KB with block size 1 KB & 7 bits in the tag. The size of main memory would be

Answer:

Block size = 1 KB = 2^{10} bytes

Number of bits in block offset = 10 bits

$$\begin{aligned}\text{Total number of lines in cache} &= \text{Cache size} / \text{Line size} \\ &= 512 \text{ KB} / 1 \text{ KB} = 512 \text{ lines}\end{aligned}$$

$$\begin{aligned}\text{Total number of sets in cache} &= \text{Total number of lines in cache} / \text{Set size} \\ &= 512 / 8 = 64 \text{ sets} = 2^6 \text{ sets}\end{aligned}$$

$$\text{Number of bits in set number} = 6 \text{ bits}$$

$$\begin{aligned}\text{No of bits in physical address} &= \text{No of bits in tag} + \text{No of bits in set number} + \text{No of bits in block offset} \\ &= 23 \text{ bits}\end{aligned}$$

$$\text{Size of main memory} = 2^{23} \text{ bytes} = \mathbf{8 \text{ MB}}$$

4. The size of main memory is 64 MB and 10 bits in the tag of a 4-way set associative mapped cache, the size of cache memory would be

Answer:

$$\text{Size of main memory} = 64 \text{ MB} = 2^{26} \text{ bytes}$$

$$\text{Number of bits in physical address} = 26 \text{ bits}$$

$$\text{Let Number of bits in set number field} = x \text{ bits}, \text{ Number of bits in block offset field} = y \text{ bits}$$

$$\text{No of bits in physical address} = \text{No of bits in tag} + \text{No of bits in set number} + \text{No of bits in block offset}$$

$$26 \text{ bits} = 10 \text{ bits} + x \text{ bits} + y \text{ bits}$$

$$x + y = 16$$

$$\text{Sum of number of bits of set number field and block offset field} = 16 \text{ bits}$$

$$\begin{aligned}\text{Cache memory size} &= \text{Number of sets in cache} \times \text{Number of lines in one set} \times \text{Line size} \\ &= 2^x \times 4 \times 2^y \text{ bytes} = 2^{2+x+y} \text{ bytes} = 2^{2+16} \text{ bytes} = 2^{18} \text{ bytes} = \mathbf{256 \text{ KB}}\end{aligned}$$

5. Cache access time is 100 ns & Memory access time is 500 ns. If the effective access time is 10% greater than the cache access time, hit ratio would be

Answer:

$$\text{cache access time} = 100 \text{ ns} \quad \text{memory access time} = 500 \text{ ns} \quad \text{Effective Access Time} = 110$$

$$\text{Effective Access Time} = \text{cache hit ratio} \times \text{cache access time} + \text{cache miss ratio} \times (\text{cache access time} + \text{main memory access time})$$

$$110 = h \times 100 \text{ ns} + (1-h) (100 + 500)$$

$$110 = 100h + 600 - 600h$$

$$500h = 490$$

$$h = 490/500 = .98 = \mathbf{98\%}$$

6) Consider a direct mapped cache of size 64 KB with block size 512 bytes. The size of main memory is 512 KB. Find Number of bits in tag.

Answer:

Number of Bits in Physical Address-

Size of main memory = 512 KB

= 2^{19} bytes

Thus, Number of bits in physical address = 19 bits.

We have, Block size = 512 bytes

= 2^9 bytes

Thus, Number of bits in block offset = 9 bits.

Total number of lines in cache

= Cache size / Line size = 64 KB / 512 bytes

= 2^{16} bytes / 2^9 bytes

= 2^7 lines

Thus, Number of bits in line number = 7 bits

Number of bits in tag

= Number of bits in physical address – (Number of bits in line number + Number of bits in block offset)

= 19 bits – (7 bits + 9 bits)

Thus, Number of bits in tag = 3 bits

= 17 bits – 14 bits

= 3 bits

Thus, Number of bits in tag = **3 bits**

7) Consider a direct mapped cache of size 256 KB with block size 512 Bytes. There are 7 bits in the tag. Find the size of main memory.

Answer:

Block size

$$= 512$$

$$= 2^9 \text{ bytes}$$

Thus, Number of bits in block offset = 9 bits

Total number of lines in cache

$$= \text{Cache size} / \text{Line size}$$

$$= 256 \text{ KB} / 512$$

$$= 2^9 \text{ lines}$$

Thus, Number of bits in line number = 9 bits

Number of bits in physical address

$$= \text{Number of bits in tag} + \text{Number of bits in line number} + \text{Number of bits in block offset}$$

$$= 7 \text{ bits} + 9 \text{ bits} + 9 \text{ bits}$$

$$= 25 \text{ bits}$$

Thus, Number of bits in physical address = 25 bits

Thus, Size of main memory

$$= 2^{25} \text{ bytes}$$

$$= \mathbf{32 \text{ MB}}$$

8) Consider a fully associative mapped cache of size 16 KB with block size 256 bytes. The size of main memory is 128 KB. Find Number of bits in tag.

Answer:

Size of main memory

$$= 128 \text{ KB}$$

$$= 2^{17} \text{ bytes}$$

Thus, Number of bits in physical address = 17 bits

Block size

= 256 bytes

= 2^8 bytes

Thus, Number of bits in block offset = 8 bits

Number of bits in tag

= Number of bits in physical address – Number of bits in block offset

= 17 bits – 8 bits

= 9 bits

Thus, Number of bits in tag = **9 bits**

9) Consider a fully associative mapped cache of size 4 MB with block size 4 KB. There are 17 bits in the tag. Find Size of main memory and tag size.

Answer:

Block size

= 4 KB

= 2^{12} bytes

Thus, Number of bits in block offset = 12 bits

Number of bits in physical address

= Number of bits in tag + Number of bits in block offset

= 17 bits + 12 bits

= 29 bits

Number of bits in physical address = 29 bits

Thus, Size of main memory

= 2^{29} bytes

= **512 MB**

10) A cache has an access time of 0.2 microseconds with memory access time of 3 microsecond. If the cache miss ratio is 25%, what is the average access time for a memory reference?

Answer:

Avg access time = Hit ratio * cache access time + miss ratio * (Cache access time + MM access time)

$$= (0.6 * 4) + 0.4 * (4 + 30)$$

$$= 2.4 + 13.6$$

=16 usec.

11. The memory access time is 2 nanosecond for a read operation with a hit in cache, 7 nanoseconds for a read operation with a miss in cache, 4 nanoseconds for a write operation with a hit in cache and 10 nanoseconds for a write operation with a miss in cache.

Execution of a sequence of instructions involves 160 memory operand read operations and 40 memory operand write operations. The cache hit-ratio is 0.9.

The average memory access time (in nanoseconds) in executing the sequence of instructions is _____.

Answer:

Time taken for 60 read operations = $160 * ((0.9 * 2) + (0.1 * 7))$

$$= 400 \text{ ns}$$

Time taken for 40 write operations = $40 * ((0.9 * 4) + (0.1 * 10))$

$$= 184 \text{ ns}$$

Total time taken for 200 operations is = $400 + 184 = 584 \text{ ns}$

Average time taken = time taken per operation = $584 / 200 \text{ ns}$

$$= \mathbf{2.92 \text{ ns}}$$

12. The main memory is structured into modules each with its own address register called _____

Answer:

Key: ABR stands for Address Buffer Register.

13. The number successful accesses to memory stated as a fraction is called as _____

Answer:

Key: The hit rate is an important factor in performance measurement.

14. 16K Bytes of main memory, How many address bits are required?

Answer:

Key: $16K = 2^4 * 2^{10} = 2^{14}$ Therefore 14 address bits are required

15. Convert Hexa A201 to Binary

Answer:

Key: Hex A = Binary 1010, Hex A = Binary 0010, Hex = Binary 0000, Hex 1 = Binary 0001

Hexa A201 = Binary 1010 0010 0000 0001