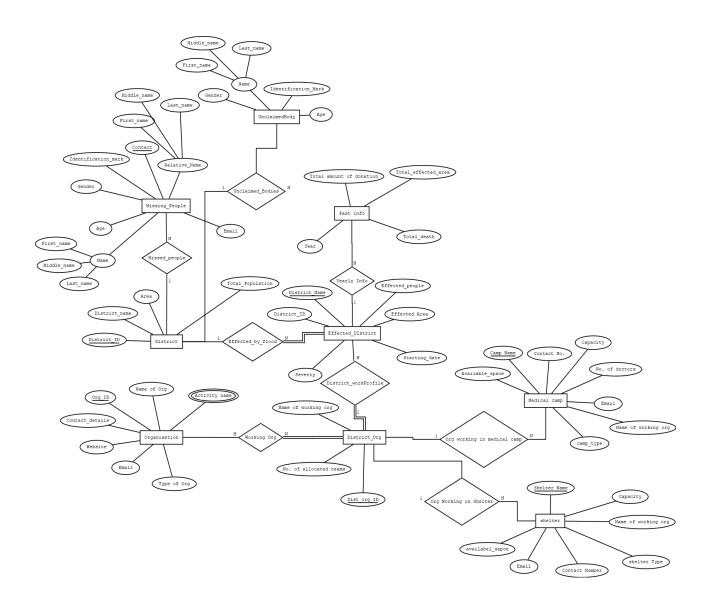
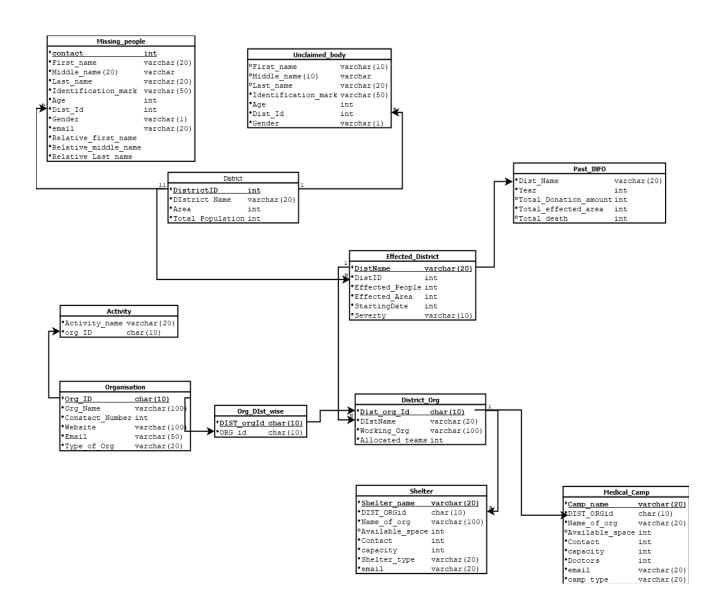
# ER DIAGRAM

# Tittle -flood rescue



# Relational Schema



# Proof of Normalisation

#### 1. District

District (districtid, district\_name, area, total\_population)

```
FDS:

districtid → district_name

district_name → area

area → total_population
```

total\_population → districtid

districtid + = (districtid, district\_name, area, total\_population)

Primary Key = districtid

BCNF Proof:

For every dependencies of minimal FD set, each determinant of relation is a candidate key. Hence the relation is BCNF.

### 2.Effected District

Effected District(district\_name, districtid, effected\_people, starting\_date, severty, effected\_area)

#### FDS:

```
district_name → districtid

districtid→ effected_people

district_name → starting_date
```

```
district_name → severty
effected_people → effected_area
effected area → district name
district_name + = (district_name, districtid, effected_people,
starting_date , severty, effected_area)
Primary Key = district_name
BCNF Proof:
For every dependencies of minimal FD set, each determinant of
relation is a candidate key. Hence the relation is BCNF.
3. Medical Camp
Medical Camp (dist_org_id, name_of_org,
camp_name,contact,capacity,available_space,camp_tpye,doctors,
email)
FDS:
camp_name → name_of_org
camp_name → dist_org_id
camp_name → contact
camp_name → capacity
camp_name → available_space
camp_name → camp_tpye
```

```
camp_name → doctors
camp_name → email
camp_name + = (dist_org_id, name_of_org,
camp_name,contact,capacity,available_space,camp_tpye,doctors,
email)
```

Primary Key = camp\_name

BCNF Proof:

For every dependencies of minimal FD set, each determinant of relation is a candidate key. Hence the relation is BCNF.

### 4. Shelter

Shelter (dist\_org\_id, name\_of\_org, shelter\_name, contact, capacity, available\_space, shelter\_tpye, email)

#### FDS:

```
shelter_name → name_of_org
shelter_name → dist_org_id
shelter_name → contact
shelter_name → capacity
shelter_name → available_space
shelter_name → shelter_tpye
shelter name → email
```

```
shelter_name + = (dist_org_id, name_of_org,
shelter_name,contact,capacity,available_space,shelter_tpye ,
email)
```

Primary Key = shelter\_name

BCNF Proof:

For every dependencies of minimal FD set, each determinant of relation is a candidate key. Hence the relation is BCNF.

### 5. Past Information

Past Information(district\_name, year ,total\_amount\_donation ,total\_effected\_area, total\_death,)

#### FDS:

```
(district_name, year) → total_amount_donation
(district_name, year) → total_effected_area
(district_name, year) → total_death
```

```
(district_name, year) + = (district_name, year
,total_amount_donation ,total_effected_area, total_death,)
```

Primary Key = (district\_name, year)

BCNF Proof:

For every dependencies of minimal FD set, each determinant of relation is a candidate key. Hence the relation is BCNF.

## 6.Organisation

```
Organisation
(org_id,org_name,conatact_number,website,email,type_of_org)
FDS:
org_id → org_name
org_id → conatact_number
conatact_number → website
website → email
email → org_id
org_id → type_of_org
org_id + =
(org_id,org_name,conatact_number,website,email,type_of_org)
Primary Key = org_id
BCNF Proof:
```

For every dependencies of minimal FD set, each determinant of relation is a candidate key. Hence the relation is BCNF.

# 7. Organisation Ids district wise

```
organisation_ids _dist (org_id,dist_org_id)
FDS:
dist_org_id → org_id
dist_org_id + = (org_id,dist_org_id)
Primary Key = dist_org_id
BCNF Proof:
For every dependencies of minimal FD set, each determinant of
relation is a candidate key. Hence the relation is BCNF.
8. Activity
activity(org_id,activity_name)
In this table there is no FDs and primary key.
This table is in 1NF.
9. District Organisation
district_org(dist_org_id, district_name, working_org, allocated_team
s)
FDS:
dist_org_id → district_name
dist_org_id → working_org
```

dist\_org\_id → allocated\_teams

```
dist_org_id * =
  (dist_org_id,district_name,working_org,allocated_teams)
Primary Key = dist_org_id
BCNF Proof :
```

For every dependencies of minimal FD set, each determinant of relation is a candidate key. Hence the relation is BCNF.

## 10. Missing People

```
Missing_people(contact_number ,name_of_relative_first, name_of_relative_middle, name_of_relative_last, email,name_of_missing_person_first, name_of_missing_person_middle, name_of_missing_person_last,age,gender,identification_mark,distrc it_id)
```

### FDS:

```
contact_number → email

contact_number → name_of_relative_first

contact_number → name_of_relative_middle

contact_number → name_of_relative_last

contact_number → name_of_missing_person_first

contact_number → name_of_missing_person_middle

contact_number → name_of_missing_person_last

contact_number → name_of_missing_person_last

contact_number → age
```

```
contact_number → gender
contact_number → identification_mark
contact_number → district_id

contact_number + = (contact_number ,name_of_relative_first,
name_of_relative_middle, name_of_relative_last,
email,name_of_missing_person_first,
name_of_missing_person_middle,
name_of_missing_person_last,age,gender,identification_mark,distrc
it_id)

Primary Key = contact_number

BCNF Proof :
```

For every dependencies of minimal FD set, each determinant of relation is a candidate key. Hence the relation is BCNF.

## 11. Unclaimed Body

```
unclaimed_body (name_of_missing_person_first,
name_of_missing_person_middle,
name_of_missing_person_last,age,gender,identification_mark,distrc
it_id)
```

In this table there is no FDs and primary key.

This table is in 1NF.