The Formal Semantics of Flix

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1 Model-theoretic Semantics

A Flix program P = (C, L) is a set of constraints C and a set of complete lattices L.

A constraint is a rule $A \Leftarrow A_1, \ldots, A_n$ where A is an atom (called the head of the rule) and A_1, \ldots, A_n are atoms (called the body of the rule). A fact is a rule with an empty body. An atom is of the form $p_\ell(t_1, \ldots, t_n)$ where p is a predicate symbol, $\ell \in L$ is the lattice associated with p, and t_1, \ldots, t_n are terms. A term is either a wildcard variable, a named variable or a constant value. The possible values are the unit value (), the booleans (true, false), the integers (-5, 3, 7), tagged unions of values (e.g. Tag v) and tuples of values (e.g. (1, true, 42)).

A complete lattice $l \in L$ is a 6-tuple $l = (E, \bot, \top, \sqsubseteq, \sqcup, \sqcup, \sqcap)$ where E is a set of elements $E \subseteq V, \bot \in E$ is the least element, $\top \in E$ is the greatest element, \sqsubseteq is the partial order on E, \sqcup is the least upper bound, and \sqcap is the greatest lower bound.

The Herbrand Universe \mathcal{U} of a Flix program P is the set of all possible ground terms. A ground term is non-wildcard, non-variable term. That is, the Herbrand Universe is exactly the set of values.

The Herbrand Base \mathcal{B} of P is the set of all possible ground atoms whose predicate symbols occur in P and where the arguments are drawn from the Herbrand Universe.

We introduce a partial order on ground atoms. Given two ground atoms $A=p_\ell(v_1,\ldots,v_n)$ and $B=p_\ell(v_1',\ldots,v_n')$ with the same (i) predicate symbol, (ii) associated lattice and (iii) number of ground terms, we define their partial order as: If n=1 then $A\sqsubseteq B$ when $v_1\sqsubseteq v_1'$. If n>1 then $A\sqsubseteq B$ when $v_1=v_1',\ldots,v_{n-1}=v_{n-1}'$ and $v_n\sqsubseteq v_n'$. This definition is implicitely stating that the last term of a ground atom is a lattice element.

An interpretation I of a Flix program P is a subset of the Herbrand Base \mathcal{B} . A ground atom A is true w.r.t. an interpretation if $\exists A' \in I$ such that $A \sqsubseteq A'$. A conjunction of atoms A_1, \dots, A_n is true

w.r.t. an interpretation if each atom is true in the interpretation. A ground rule is true if either the body conjunction is false, or the head is true.

A model M of P is an interpretation that makes each ground instance of a each rule in P true.

A model is *compact* if it has no two ground atoms $A, A' \in M$ such that $A \sqsubseteq A'$. A model is *reduced* if it has no two ground atoms $A, A' \in M$ such that $A \sqsubseteq (A \sqcup A')$. Finally, a model M_1 is *minimal* if it is compact and reduced, and there is no other model M_2 such that there are two ground atoms $A \in M_1$ and $A' \in M_2$ where $A' \sqsubseteq A$.

Example. The Flix program P with constraints:

$$A_{\ell}(\texttt{Even})$$
 $A_{\ell}(\texttt{Odd})$ $B_{\ell}(\texttt{Odd})$

and lattices $\{\ell=(\bot,\top,\mathtt{Even},\mathtt{Odd}\},\sqsubseteq,\sqcup,\sqcap)\}$ has the Herbrand Universe:

$$\mathcal{U} = \{\bot, \top, \mathsf{Even}, \mathsf{Odd}\}$$

and the Herbrand Base:

$$\mathcal{B} = \{A_{\ell}(\bot), A_{\ell}(\mathtt{Even}), A_{\ell}(\mathtt{Odd}), A_{\ell}(\top), \\ B_{\ell}(\bot), B_{\ell}(\mathtt{Even}), B_{\ell}(\mathtt{Odd}), B_{\ell}(\top)\}$$

An interpretation of P is a subset of \mathcal{B} . For example,

$$\begin{split} I_1 &= \{A_\ell(\top)\} \\ I_2 &= \{A_\ell(\top), B_\ell(\bot)\} \\ I_3 &= \{A_\ell(\top), B_\ell(\texttt{Odd}), B_\ell(\top)\} \\ I_4 &= \{A_\ell(\texttt{Even}), A_\ell(\texttt{Odd}), B_\ell(\texttt{Odd})\} \\ I_5 &= \{A_\ell(\top), B_\ell(\top)\} \\ I_6 &= \{A_\ell(\top), B_\ell(\texttt{Odd})\} \end{split}$$

The interpretation I_1 is not a model of P since it does not make $B_{\ell}(\texttt{Odd})$ true. I_2 is also not a model of p since it does not make $B_{\ell}(\texttt{Odd})$ true. I_3 is a model of p, but it is not compact. I_4 is a model of p, but it is not reduced. I_5 is a model of p, it is compact and reduced, but it is not minimal as evidenced by I_6 .

Example. The Flix program P with constraints:

$$A_\ell(1,\operatorname{Pos})$$
 $A_\ell(2,\operatorname{Pos})$ $A_\ell(2,\operatorname{Neg})$

and lattices $\{\ell=(\bot,\top,\texttt{Neg},\texttt{Zer},\texttt{Pos}\},\sqsubseteq,\sqcup,\sqcap)\}$ has the Herbrand Universe:

$$\mathcal{U} = \{1, 2, \perp, \top, \mathtt{Neg}, \mathtt{Zer}, \mathtt{Pos}\}$$

and the Herbrand Base:

$$\mathcal{B} = \{ A_{\ell}(1,1), A_{\ell}(1,2), A_{\ell}(1,\perp), A_{\ell}(1,\top), \cdots \\ A_{\ell}(2,1), A_{\ell}(2,2), A_{\ell}(2,\perp), A_{\ell}(2,\top), \cdots \\ A_{\ell}(\perp,1), A_{\ell}(\perp,2), A_{\ell}(\perp,\perp), A_{\ell}(\perp,\top), \cdots \\ \cdots \}$$

An interpretation of P is a subset of \mathcal{B} . For example,

$$\begin{split} I_1 &= \{A_\ell(1,\top)\} \\ I_2 &= \{A_\ell(1,\top), A_\ell(2,\top)\} \\ I_3 &= \{A_\ell(1,\operatorname{Pos}), A_\ell(1,\operatorname{Zer}), A_\ell(2,\top)\} \\ I_4 &= \{A_\ell(1,\operatorname{Pos}), A_\ell(2,\top)\} \end{split}$$

Here I_2 , I_3 and I_4 are models and I_4 is minimal.