| 001. | | or inte | erception of data. | D |
|------------------|---|------------|-----------------------------------|---|
| | A repudiation | В | replaying | |
| | C modification | D | Snooping | |
| 002. | Find the solution of x^2 16 mod 23. | | | Α |
| | A $x = 4$ and 19 | В | x = 6 and 17 | |
| | C = 11 and 12 | | x = 7 and 16 | |
| በበ3 | Theof data can be threatened by sev | _ | | В |
| 00 3. | | ciai r | ands of attacks. Modification, | ט |
| | masquerading, replaying, and repudiation. | D | intogrity | |
| | A Availability | В | integrity Confidentiality | |
| 004 | C Authenticity The linear combination of god(252.108), 18 | D | Confidentiality | _ |
| υ υ 4. | The linear combination of gcd(252,198)=18 | | | С |
| | A 252*5-198*4 | | 252*4-198*2 | |
| 005 | C 252*4-198*5 | D | 252*2-198*5 | _ |
| 005. | In computer security, means that t | | formation contained in a computer | Α |
| | system can only be read by authorized pers | | | |
| | A Confidentiality | В | Integrity | |
| | C Availability | D | Authenticity | |
| 006. | $-5 \mod -3 =$ | | | С |
| | A 3 | В | 2 | |
| | C 1 | D | 5 | |
| 007. | In computer security, means that a | active | computer systems can only be | В |
| | modified by authorized persons. | | | |
| | A Confidentiality | В | Integrity | |
| | • | D | Authenticity | |
| 008 | The two types of attacks threaten the | _ | • | С |
| | analysis. | O1 11 | | • |
| | A Integrity | В | Availability | |
| | C confidentiality | D | • | |
| 000 | • | U | Authenticity | Þ |
| oos. | is an attack threatening availability | В | Denial of Service | В |
| | A Repudiation | | | |
| 040 | C Replaying | D :- | Masquerading | _ |
| U10. | The solution of linear congruence 4x5(mod | | | В |
| | A 6(mod 9) | В | 8(mod 9) | |
| | C 9(mod 9) | D | 10(mod 9) | _ |
| 011. | In, the attackers goal is just to obta | | | С |
| | A Slow attack | В | Active attack | |
| | C passive attack | D | sleep attack | |
| 012. | The value of 5 ²⁰⁰³ mod 7 is | | | Α |
| | A 3 | В | 4 | |
| | C 8 | D | 9 | |
| 013 | In, the sender of the message might la | _ | | Α |
| 010. | the receiver of the message might later den | | | ~ |
| | | ıy ша В | replaying | |
| | A repudiation C modification | D D | , , , | |
| 04.4 | | ט | Snooping | ^ |
| U14. | The inverse of 7 modulo 26 is | D | 4.4 | С |
| | A 12 | В | 14 | |
| 045 | C 15 | D | 20 | _ |
| 015. | • | | | D |
| | A Replaying | В | Snooping | |
| | C Traffic Analysis | D | Masquerading | |
| 016. | The inverse of 3 modulo 7 is | | | В |
| | A -1 | В | -2 | |
| | C -3 | D | -4 | |
| 017. | Attacks that threaten the integrity and availa | ability | are | В |
| | | | | |

| | А | Slow attack | В | Active attack | |
|--------------------------------------|--|--|--|--|-------------|
| | С | passive attack | D | sleep attack | |
| 018. | Expa | and ITU-T: | | • | D |
| | Α΄ | International Telecom Union-Telecom | В | International Telecom Union-Telecom | |
| | | Standardization Section | | Standardization Sector | |
| | С | International Telecommunication | D | International Telecommunication | |
| | | Union-Telecommunication | | Union-Telecommunication | |
| | | Standardization Section | | Standardization Sector | |
| 019 | | is designed to protect data from disc | losure | | С |
| 0.0. | Α | Data Integrity | В | nonrepudiation | • |
| | C | 5 , | D | Access Control | |
| 000 | _ | Data Confidentiality | _ | | |
| UZU. | | is designed to protect data from mo | Juliuca | ation, insertion, deletion, andreplaying | Α |
| | | n adversary | _ | 11. 41 | |
| | A | Data Integrity | В | nonrepudiation | |
| | С | Data Confidentiality | D | Access Control | |
| 021. | An - | attack may change the data or ha | rm the | e system | Α |
| | Α | active | В | passive | |
| | С | slow | D | fast | |
| 022. | A bir | nary operation takes inputs and c | reate | s output | C |
| | Α | 1, 2 | В | 2, 2 | |
| | С | 2, 1 | D | 2, 0 | |
| 023. | Atta | | ng an | nd traffic analysis, areattacks | D |
| | Α | fast | В | slow | |
| | C | active | D | passive | |
| 024 | _ | nodn) if and only if n (ab). | | paddivo | В |
| UZ-T. | A | n (a). | В | n (ab) . | |
| | C | n (a+b) . | D | n (ab) . n (a*b) . | |
| 005 | | • • • • | D | τη(a b). | Α |
| | | 4 (I | | | |
| 025. | | 1, then a = ± | _ | | • |
| 025. | Α | 1 | В | 0 | , |
| | A C | 1 2 | D | a | |
| | A C | 1 2 the greatest common divisor of 2740 a | D | a 760. | D |
| | A C | 1 2 | D | a | |
| | A C Find | 1 2 the greatest common divisor of 2740 a | D ind 17 | a 760. | |
| 026. | A C Find A C | 1 2 the greatest common divisor of 2740 a 25 | D ind 17 B D | a 760. 5 20 | |
| 026. | A C Find A C | 1 2 the greatest common divisor of 2740 a 25 10 | D ind 17 B D | a 760. 5 20 | D |
| 026. | A C Find A C A lin | 1 2 the greatest common divisor of 2740 a 25 10 ear Diophantine equation of two variab | D and 17 B D les is | a 760. 5 20 | D |
| 026. 027. | A C Find A C A lin A | 1 2 the greatest common divisor of 2740 a 25 10 ear Diophantine equation of two variab xy+a=y ax + by = c | D and 17 B D les is B D | a 760. 5 20 xy = c x+y=xy | D C |
| 026. 027. | A C Find A C A lin A | 1 2 the greatest common divisor of 2740 a 25 10 ear Diophantine equation of two variab xy+a=y ax + by = c | D and 17 B D les is B D | a 760. 5 20 xy = c | D C |
| 026. 027. | A C Find A C A lin A | 1 2 the greatest common divisor of 2740 a 25 10 ear Diophantine equation of two variab xy+a=y ax + by = c | D and 17 B D les is B D | a 760. 5 20 xy = c x+y=xy | D C |
| 026. 027. | A C Find A C A lin A C The | the greatest common divisor of 2740 a 25 10 ear Diophantine equation of two variab xy+a=y ax + by = c result of the modulo operation with modulo n | D Ind 17 B D les is B D dulus | a 760. 5 20 xy = c x+y=xy n is always an integer between 0 and- n 1 | D C |
| 026. 027. 028. | A C Find A C A lin A C The A C | the greatest common divisor of 2740 at 25 10 ear Diophantine equation of two variab xy+a=y ax + by = c result of the modulo operation with modulo n n+1 | D and 17 B D les is B D dulus B D | a 760. 5 20 xy = c x+y=xy n is always an integer between 0 and- n 1 n/2 | D C B |
| 026. 027. 028. | A C Find A C A lin A C The A C | the greatest common divisor of 2740 at 25 10 ear Diophantine equation of two variab | D and 17 B D les is B D dulus B D raffic | a 760. 5 20 xy = c x+y=xy n is always an integer between 0 and- n 1 n/2 analysis attack | D C |
| 026. 027. 028. | A C Find A C A lin A C The A C | the greatest common divisor of 2740 at 25 10 ear Diophantine equation of two variable xy+a=y ax + by = c result of the modulo operation with modulo operat | D and 17 B D les is B D dulus B D raffic B | a 760. 5 20 xy = c x+y=xy n is always an integer between 0 and- n 1 n/2 analysis attack nonrepudiation | D C B |
| 026. 027. 028. | A C Find A C A lin A C The A C A C | the greatest common divisor of 2740 at 25 10 ear Diophantine equation of two variab xy+a=y ax + by = c result of the modulo operation with modulo operatio | D and 17 B D les is B D dulus Paffic B D | a 760. 5 20 xy = c x+y=xy n is always an integer between 0 and- n 1 n/2 analysis attack nonrepudiation Access Control | D C B |
| 026. 027. 028. | A C Find A C A lin A C The A C | the greatest common divisor of 2740 at 25 10 ear Diophantine equation of two variab | D and 17 B D les is B D dulus Paffic B D | a 760. 5 20 xy = c x+y=xy n is always an integer between 0 and- n 1 n/2 analysis attack nonrepudiation Access Control | D C B |
| 026. 027. 028. | A C Find A C A lin A C The A C A C In deliv | the greatest common divisor of 2740 at 25 10 ear Diophantine equation of two variab xy+a=y ax + by = c result of the modulo operation with modulo operatio | D and 17 B D les is B D dulus B D raffic B D of da | a 760. 5 20 xy = c x+y=xy n is always an integer between 0 and- n 1 n/2 analysis attack nonrepudiation Access Control ta can later prove that data were | D C B |
| 026. 027. 028. | A C Find A C A lin A C The A C A C In A C In deliv A | the greatest common divisor of 2740 at 25 10 ear Diophantine equation of two variab xy+a=y ax + by = c result of the modulo operation with modulo operatio | D and 17 B D les is B D dulus B D raffic B D of da B | a 760. 5 20 xy = c x+y=xy n is always an integer between 0 and- n 1 n/2 analysis attack nonrepudiation Access Control ta can later prove that data were nonrepudiation | D C B |
| 026. 027. 028. 029. | A C Find A C A lin A C The A C In A C deliv A C | the greatest common divisor of 2740 at 25 10 ear Diophantine equation of two variab xy+a=y ax + by = c result of the modulo operation with modulo operatio | D and 17 B D les is B D dulus B D raffic B D of da B D | a 760. 5 20 xy = c x+y=xy n is always an integer between 0 and- n 1 n/2 analysis attack nonrepudiation Access Control ta can later prove that data were nonrepudiation Access Control | D C B |
| 026. 027. 028. 029. | A C Find A C A lin A C The A C A C In A C In | the greatest common divisor of 2740 at 25 10 ear Diophantine equation of two variab xy+a=y ax + by = c result of the modulo operation with modulo operatio | D and 17 B D les is B D dulus B D raffic B D of da B D | a 760. 5 20 xy = c x+y=xy n is always an integer between 0 and- n 1 n/2 analysis attack nonrepudiation Access Control ta can later prove that data were nonrepudiation Access Control | D C B |
| 026. 027. 028. 029. | A C Find A C A lin A C In A C In the state of the sta | the greatest common divisor of 2740 at 25 10 ear Diophantine equation of two variab xy+a=y ax + by = c result of the modulo operation with modulo and the confidentiality Data Confidentiality Data Confidentiality Data Integrity Data Integrity Data Integrity Data Integrity Data Confidentiality Data Confi | D and 17 B D les is B D dulus B D raffic B D of da B D er of t | a 760. 5 20 xy = c x+y=xy n is always an integer between 0 and- n 1 n/2 analysis attack nonrepudiation Access Control ta can later prove that data were nonrepudiation Access Control the data can later prove the identity of | D C B |
| 026. 027. 028. 029. | A C Find A C A lin A C In A C In the s A | the greatest common divisor of 2740 at 25 10 ear Diophantine equation of two variab xy+a=y ax + by = c result of the modulo operation with modulo at a Integrity Data Confidentiality Data Integrity Data Integrity Data Integrity Data Integrity Data Confidentiality Data Integrity Data Integrity Data Integrity | D and 17 B D les is B D dulus B D raffic B D of da B D er of t B | a 760. 5 20 xy = c x+y=xy n is always an integer between 0 and- n 1 n/2 analysis attack nonrepudiation Access Control ta can later prove that data were nonrepudiation Access Control he data can later prove the identity of nonrepudiation | D C B |
| 026. 027. 028. 029. 030. | A C Find A C A Iin A C The A C In A C In the s A C | the greatest common divisor of 2740 at 25 10 ear Diophantine equation of two variab xy+a=y ax + by = c result of the modulo operation with modulo operatio | D and 17 B D les is B D dulus B D raffic B D of da B D er of t | a 760. 5 20 xy = c x+y=xy n is always an integer between 0 and- n 1 n/2 analysis attack nonrepudiation Access Control ta can later prove that data were nonrepudiation Access Control the data can later prove the identity of | D C B |
| 026. 027. 028. 029. 030. | A C Find A C A Iin A C The A C In A C In the s A C | the greatest common divisor of 2740 at 25 10 ear Diophantine equation of two variab xy+a=y ax + by = c result of the modulo operation with modulo at a Integrity Data Confidentiality Data Integrity Data Integrity Data Integrity Data Integrity Data Confidentiality Data Integrity Data Integrity Data Integrity | D and 17 B D les is B D dulus B D raffic B D of da B D er of t B D | a 760. 5 20 Xy = C X+y=xy n is always an integer between 0 and- n 1 n/2 analysis attack nonrepudiation Access Control ta can later prove that data were nonrepudiation Access Control he data can later prove the identity of nonrepudiation Access Control | D C B |
| 026. 027. 028. 029. 030. | A C Find A C A Iin A C The A C In A C In the s A C | the greatest common divisor of 2740 at 25 10 ear Diophantine equation of two variab xy+a=y ax + by = c result of the modulo operation with modulo operatio | D and 17 B D les is B D dulus B D raffic B D of da B D er of t B | a 760. 5 20 xy = c x+y=xy n is always an integer between 0 and- n 1 n/2 analysis attack nonrepudiation Access Control ta can later prove that data were nonrepudiation Access Control he data can later prove the identity of nonrepudiation | D C B |

| 033. | In Zn, two numbers a and b are the multiplic | cative | inverse of each other if | Α |
|-------|---|--------------|--|---|
| | A a + b 1 (mod n) | В | a + b 0 (mod n) | |
| | C $a + b - 1 \pmod{n}$ | D | a + b k (mod n) | |
| 034. | The multiplicative Inverse of 1234 mod 432 | 1 is | , | Α |
| | A 3239 | В | 3213 | |
| | C 3242 | D | 3225 | |
| 035. | (1,723,345 + 2,124,945) mod 11= | | | C |
| | A 8 | В | 10 | |
| | C 6 | D | 9 | |
| 036. | What is 11 mod 7 and -11 mod 7? | | | D |
| | A 4 and 5 | В | 4 and 4 | |
| | C 5 and 3 | D | 4 and -4 | |
| 037. | What is the result of the following operation | ? Sub | | С |
| | A 6 | В | 0 | |
| | C 9 | D | 17 | _ |
| 038. | In Zn, two numbers a and b are additive inv | | | В |
| | A a + b 1 (mod n) | В | a + b 0 (mod n) | |
| | C a + b -1 (mod n) | D | a + b k (mod n) | _ |
| 039. | What is the result of the following operation | | | Α |
| | A 6 | В | 10 | |
| 0.40 | C 9 | D O Maria | 17 | _ |
| 040. | What is the result of the following operation | | | D |
| | A 6 C 9 | B D | 10 17 | |
| 044 | | _ | | D |
| 041. | the conder and the receiver to provent the | | - | D |
| | the sender and the receiver to prevent the c particular route. | pppoi | lent from eavesdropping on a | |
| | A Notarization | В | Routing control | |
| | C Traffic padding | D | Access control | |
| 042 | uses methods to prove that a user I | _ | | D |
| UTZ. | owned by a system. | ias a | coess fight to the data of resources | |
| | | R | Routing control | |
| | C Traffic padding | D | Access control | |
| 043. | refers to selecting a third trusted par | _ | | Α |
| 0 .0. | two entities. | ., . | | |
| | A Notarization | В | digital signature | |
| | C authentication | D | enciphering | |
| 044. | Inexchange, two entities exchange | som | , • | С |
| | each other | | , | |
| | A Notarization | В | digital signature | |
| | C authentication | D | enciphering | |
| 045. | -5 mod -3 = | | | C |
| | A 3 | В | 2 | |
| | C 1 | D | 5 | |
| 046. | means inserting some bogus data | into t | he data traffic to thwart the adversarys | C |
| | attempt to use the traffic analysis. | | | |
| | A Notarization | В | Routing control | |
| | C Traffic padding | D | Access control | |
| 047. | $GCD(a,b) = GCD(b,a \mod b)$ | | | A |
| | A true | В | false | |
| | C infinity | D | cantt defined | |
| 048. | 7x 6 mod 5. Then the value of x is | | | В |
| | A 2 | В | 3 | |
| | C = A | | E | |

| 049. | Find | the multiplicative inverse of 11 in Z26 | | | В |
|------------------|----------|--|------------|--|---|
| | Α | -7 | В | 19 | |
| | С | -2 | D | 26 | |
| 050. | A re | sidue matrix has a multiplicative invers | e if go | | С |
| | Α | 0 | В | Doesnt exist | |
| | С | 1 | D | -1 | _ |
| 051. | | follows | _ | | С |
| | A | Hash Algorithm | В | Caesars Cipher | |
| | C | Feistel Cipher Structure | D | SP Networks | _ |
| 052. | For | the group S _n of all permutations of n dis | stinct | symbols, what is the number of | D |
| | eiem | nents in S _n ? | | | |
| | Α | n | В | n-1 | |
| | С | 2n | D | n! | |
| 053. | The | integer a in Zn has a multiplicative inve | | | Α |
| | A | 1 (mod n) | В | -1 (mod n) | |
| | C | 0 (mod n) | D | k (mod n) | _ |
| 054. | | extended Euclidean algorithm finds the | mult | iplicative inverses of b in Zn when n | С |
| | _ | b are given and gcd (n, b) = | D | L | |
| | A C | 0 1 | B D | b | |
| 0EE | _ | • | _ | -1 | D |
| 055. | | is a means by which the sender | | ectronically sign the data and the | В |
| | A | iver can electronically verify the signatu Notarization | ле. В | digital signature | |
| | Ĉ | authentication | D | enciphering | |
| 056 | _ | tography and steganography are used | _ | , | D |
| 000. | A | Notarization | В | digital signature | |
| | C | authentication | D | enciphering | |
| 057. | In th | e DES algorithm the round key is bi | t and | . • | Α |
| | Α | 48, 32 | В | 64,32 | |
| | С | 56, 24 | D | 32, 32 | |
| 058. | How | many numbers cannot be used in GF(| n) in 3 | 2 ⁿ where n=4? | C |
| | Α | 2 | В | 3 | |
| | | 5 | D | 1 | |
| 059. | In th | e DES algorithm the Round Input is 32 | bits, | which is expanded to 48 bits via | Α |
| | | | | · | |
| | Α | Scaling of the existing bits | В | Duplication of the existing bits | |
| | С | Addition of zeros | D | Addition of ones | |
| 060. | On r | nultiplying $(x^5 + x^2 + x)$ by $(x^7 + x^4 + x^3 + x^3)$ | $(^2 + x)$ | in GF(2 ⁸) with irreducible polynomial | D |
| | $(x^8 +$ | $x^4 + x^3 + x + 1$) we get | , | ` , | |
| | | $x^{12}+x^{7}+x^{2}$ | В | $x^5 + x^3 + x^3$ | |
| | | $x^5 + x^3 + x^2 + x$ | | $x^{5}+x^{3}+x^{2}+x+1$ | |
| 061 | | DES algorithm has a key length of | | X TX TX TXT | С |
| 001. | A | 128 Bits | В | 32 Bits | C |
| | C | 64 Bits | D | 16 Bits | |
| 062 | _ | GCD of $x^3 + x + 1$ and $x^2 + x + 1$ over G | _ | | Α |
| 002. | | | ` ' | | • |
| | A C | 1 | B D | x + 1 | |
| 000 | | ^ | | $x^2 + 1$ | _ |
| 063. | | DES Algorithm Cipher System consists | s of | rounds (iterations) each | ט |
| | | a round key | D | 10 | |
| | A C | 12 9 | B D | 18 16 | |
| 064 | _ | | _ | | D |
| 004 . | On r | nultiplying ($x^6+x^4+x^2+x+1$) by (x^7+x+1) | in GF | ·(2°) with irreducible polynomial (x°+ | В |

| | $x^4 + x$ | x ³ + x + 1) we get | | | |
|----------|-----------|--|---------|---------------------------------------|---|
| | | $x^7 + x^6 + x^3 + x^2 + 1$ | В | $x^7 + x^6 + 1$ | |
| | | $x^6 + x^5 + x^2 + x + 1$ | | $x^{7}+x^{6}+x+1$ | |
| 065. | | many keys does the Triple DES algorit | | | С |
| | | 2 | В | 3 | |
| | С | 2 or 3 | D | 3 or 4 | |
| 066. | How | many rounds does the AES-192 performance | m? | | В |
| | Α | 10 | В | 12 | |
| | С | 14 | D | 16 | _ |
| 067. | _ | t is the size of the key in the SDES algo | | | D |
| | A | 24 bits | В | 16 bits | |
| 068 | C | 20 bits many rounds does the AES-256 performance of the AES-256 performance | D m2 | 10 bits | С |
| 000. | A | 10 | В | 12 | C |
| | C | 14 | D | 16 | |
| 069. | • | number of unique substitution boxes in | _ | | Α |
| | Α | 8 | В | 4 | |
| | С | 6 | D | 12 | |
| 070. | AES | uses a bit block size ar | nd a k | key size of bits. | D |
| | Α | 128; 128 or 256 | В | 64; 128 or 192 | |
| | | 256; 128, 192, or 256 | D | 128; 128, 192, or 256 | _ |
| 071. | | ES, the Initial Permutation table/matrix | | | С |
| | A | 16x8 | В | 12x8 | |
| 072 | | 8x8 | D | 4x8 | ٨ |
| 012. | | GCD of x^3 - x + 1 and x^2 + 1 over GF(3) | | | Α |
| | A | 1 | В | X | |
| | | x + 1 | | $x^2 + 1$ | _ |
| 073. | | th of the 4 operations are false for each | | · , | В |
| | | s ii) Shift Columns iii) Mix Rows iv) XOI | _ | • | |
| | A C | i) only ii) and iii) | B D | ii) iii) and iv) only iv) | |
| 074 | | combination of the set and the operatio | _ | • , | Α |
| . | | s calledan/a | | at are applied to the elements of the | • |
| | Α | algebraic structure | В | group | |
| | С | commutative group | D | abelian group | |
| 075. | In AE | ES the 4x4 bytes matrix key is transforn | ned ir | nto a keys of size | D |
| | | 32 words | В | 64 words | |
| | | 54 words | D | 44 words | _ |
| 076. | | is a set of elements with a binary op | | | В |
| | | algebraic structure | В | group | |
| 077 | | commutative group | D | abelian group | ٨ |
| U//. | | 4x4 byte matrices in the AES algorithm States | В | Words | Α |
| | C | Transitions | D | Permutations | |
| 078. | _ | DES function has components | | Territations | С |
| 0.0. | | 2 | B | 3 | |
| | С | 4 | D | 5 | |
| 079. | Wha | t is the expanded key size of AES-1923 | ? | | C |
| | | 44 words | В | 60 words | |
| | С | 52 words | D | 36 words | |
| 080. | | DES was designed to increase the | ne siz | | В |
| | | Double | В | Triple | |
| | С | Quadruple | D | zero | |

| 081. | A ring involvesoperations | | | C |
|------|--|--------|---------------------------------------|---|
| | A one | В | three | |
| | C two | D | four | |
| 082. | Which of the following is a faulty S-AES ste | p fun | ction? | В |
| | A Add round key | В | Byte substitution | |
| | C Shift rows | D | Mix Columns | |
| 083. | What is the block size in the Simplified AES | 3 algo | orithm? | В |
| | A 8 bits | В | 40 bits | |
| | C 16 bits | D | 36 bits | |
| 084. | How many computation rounds does the si | mplifi | ed AES consists of? | В |
| | A 5 | В | 2 | |
| | C 8 | D | 10 | |
| 085. | If a subgroup of a group can be generated | using | the power of an element, the | Α |
| | subgroup is called thesubgroup. | | | |
| | A cyclic | В | finite | |
| | C infinite | D | closed | |
| 086. | On comparing AES with DES, which of the | follov | ving functions from DES does not have | C |
| | an equivalent AES function? | | | |
| | A permutation p | В | f function | |
| | C swapping of halves | D | xor of subkey with function f | |
| 087. | theorem relates the order of a group | | | D |
| | A Euclidian | В | tailor | |
| | C fermat | D | Lagranges | _ |
| 088. | Which function can be used in AES multipli | | | В |
| | A $m(x)=x^7+x^4+x^3$ | В | $m(x)=x^8+x^4+x^3+x+1$ | |
| | $C = m(x) = x^8 + x^3 + x^2 + x + 1$ | D | $m(x)=x^8+x^5+x^3+x$ | |
| 089. | Multiplication of polynomials in GF(2n) can | be a | chieved usingandexclusive-or | D |
| | operations. | | - | |
| | A AND | В | rotate | |
| | C shift-right | D | shift-left | |
| 090. | Polynomials representingbit words u | se tw | o fields: GF(2) and GF(2n). | C |
| | A 2n | В | 2 | |
| | C n | D | 8 | |
| 091. | A Galois field, $GF(p^n)$, is a finite field with | (| elements. | D |
| | A pn | В | n | |
| | Ср | D | p^n | |
| 092. | How many step functions do Round 1 and 2 | 2 eac | • | Α |
| | A 4 and 3 | В | Both 4 | |
| | C 1 and 4 | D | 3 and 4 | |
| 093. | Addition/subtraction in GF(2) is the same a | s the | operation | В |
| | A OR | В | XOR | |
| | C AND | D | NOR | |
| 094. | How many round keys are generated in the | AES | algorithm? | Α |
| | A 11 | В | 10 | |
| | C 8 | D | 12 | |
| 095. | DES uses a key generator to generate sixte | een _ | round keys. | В |
| | A 32-bit | В | 48-bit | |
| | C 54-bit | D | 42-bit | |
| 096. | A is invertible, but compression a | | • | С |
| | A P-Box | В | Compression P-Box | |
| | C Straight P-Box | D | Expansion P-boxes | _ |
| 097. | A(n) is a keyless substitution cipher | | | Α |
| | formula to define the relationship between | the in | put stream and the output stream. | |

| | Α | S-box | В | P-box | |
|------|--------|---|--------|------------------------------------|---|
| | С | T-box | D | B-box | |
| 098. | Whic | ch one of the following is not a possible | key l | ength for the Advanced Encryption | Α |
| | Stan | dard Rijndael cipher? | • | | |
| | Α | 56 bits | В | 128 bits | |
| | C | 192 bits | D | 256 bits | |
| nga | _ | parallels the traditional transposit | _ | | В |
| 033. | Α | S-Box | В | P-Box | |
| | | | _ | | |
| 400 | C | T-Box | D | N-Box | _ |
| 100. | | Advanced Encryption Standard (AES), | | nree different configurations with | С |
| | | ect to the number of rounds and | | | |
| | Α | data size | В | round size | |
| | С | key size | D | encryption size | |
| 101. | | are used when we need to permute | bits a | and the same time increase the | D |
| | num | ber of bits for the next stage | | | |
| | Α | P-Box | В | Compression P-Box | |
| | С | Straight P-Box | D | Expansion P-boxes | |
| 102. | To p | rovide security, AES uses type | es of | | D |
| | Α΄ | 2 | В | 3 | |
| | С | 5 | D | 4 | |
| 103 | (231 | | _ | • | D |
| | Α Α | 230 | В | 60 | |
| | C | 80 | D | 120 | |
| 104 | _ | nonsynchronous stream cipher, the key | _ | _ | С |
| 107. | | ertext | , ucp | ends on entirer the plaintext of | U |
| | | | D | modern | |
| | A C | synchronous | B D | modern | |
| 105 | | nonsynchronous | _ | general | D |
| 105. | | hides the relationship between the | | · | ט |
| | | Key generator | В | Confusion | |
| 400 | С | Key schedule | D | Diffusion | _ |
| 106. | _ | hides the relationship between the ci | · | • | В |
| | A | Key generator | В | Confusion | |
| | С | Key schedule | D | Diffusion | _ |
| 107. | In a | stream cipher the key is inde | _ | | Α |
| | Α | synchronous | В | modern | |
| | С | asynchronous | D | general | |
| 108. | | stream cipher, each r-bit w | | | В |
| | | g an r-bit word in the key stream to crea | ate th | e corresponding r-bit word in the | |
| | ciphe | ertext stream. | | | |
| | Α | synchronous | В | modern | |
| | С | asynchronous | D | general | |
| 109. | The | inverse of 49 mod 37 is | | | D |
| | Α | 31 | В | 23 | |
| | С | 22 | D | 34 | |
| 110. | Find | the primitive roots of G= <z11*, x="">?</z11*,> | | | D |
| | Α | {2, 6, 8} | В | {2, 5, 8} | |
| | С | {3, 4, 7, 8} | D | {2, 6, 7, 8} | |
| 111. | | the number of primitive roots of G= <z1< th=""><th></th><th></th><th>С</th></z1<> | | | С |
| • • | Α | 5 | В, х | 6 | • |
| | C | 4 | D | 10 | |
| 112 | _ | the order of group G= <z20*, x=""></z20*,> | _ | . • | D |
| | A | 6 | В | 9 | |
| | C | 10 | D | 8 | |
| 112 | _ | x for the CRT when $x = 2 \mod 3$; $x = 3 \text{ r}$ | _ | | С |
| ııJ. | iiiu | A IOI LITE OIL I WITCH A- Z IIIOU 3, X= 3 I | nou t | , A – Z 1110U <i>i</i> | J |

| | Α | 33 | В | 22 | |
|------|------------------------------|---|---------|--|---|
| | С | 23 | D | 31 | |
| 114. | | | 38, ar | nd 50, does not have primitive roots in | В |
| | the g | $group G = \langle Zn^*, x \rangle$? | | | |
| | Α | 17 | В | 20 | |
| | С | 38 | D | 50 | _ |
| 115. | | | ents in | n the set {a: 0 <= a < n and gcd(a,n) = | В |
| | • | /hat is this function? | _ | | |
| | A | Primitive | В | Totient | |
| | С | Primality | D | Primitive, totient, primality | |
| 116. | | theoremstates that ifpis aprime nun | nber, | then for anyintegera, the numbera ^p - | D |
| | | n integer multiple ofp | | | |
| | Α | Euclidean | В | Eulers | |
| | С | Chinese remainder | D | Fermat 's little | |
| 117. | | x - 16 = 0 (mod 29) | | | В |
| | | $x = 6, 24 \pmod{29}$ | | $x = 9, 24 \pmod{29}$ | |
| | | $x = 9, 22 \pmod{29}$ | D | $x = 6, 22 \pmod{29}$ | |
| 118. | $17 x^2$ | ² 10 (mod 29) | | | С |
| | | x = 3, 22 (mod 29) | В | $x = 7, 28 \pmod{29}$ | |
| | С | $x = 2, 27 \pmod{29}$ | D | $x = 4, 28 \pmod{29}$ | |
| 119. | Six to | eachers begin courses on Monday Tue | sday | Wednesday Thursday Friday and | В |
| | Satu | rday, respectively, and announce their | inten | tions of lecturing at intervals of | |
| | 2,3,4 | 1,1,6 and 5 days respectively. Sunday le | ecture | es are forbidden. When first will all the | |
| | teach | ners feel compelled to omit a lecture? L | Jse C | RT. | |
| | Α | 354 | В | 371 | |
| | С | 432 | D | 213 | |
| 120. | The | number of primes of the form $ n ^2$ 6n + | 5l wh | here n is an integer is | C |
| | Α | 0 | B | 1 | |
| | С | 2 | D | 3 | |
| 121. | How | many primitive roots are there for 25? | | | D |
| | Α | 4 | В | 5 | |
| | С | | D | 8 | |
| 122. | wher | nmandnare coprime, (m*n) = | | | Α |
| | Α | (m)*(n) | В | (m)+(n) | |
| | С | (m)/(n) | D | (m)**(n) | |
| 123. | The | inverse of 37 mod 49 is - | | | С |
| | Α | 23 | В | 12 | |
| | С | 4 | D | 6 | |
| 124. | Does | s the set of residue classes (mod 3) for | m a g | roup with respect to modular addition? | Α |
| | Α | Yes | В | No | |
| | С | Cant Say | D | Insufficient Data | |
| 125. | Usin | g Differential Crypt-analysis, the minim | um co | omputations required to decipher the | D |
| | DES | algorithm is | | | |
| | Α | 256 | В | 243 | |
| | С | 255 | D | 247 | |
| 126. | Find | the set of quadratic residues in the set | - Z11 | 1* = { 1, 2, 3, 4, 5, 6, 7, 8, 9, 10} | D |
| | Α | QR set = {1, 2, 4, 5, 9} of Z11* | В | QR set = {1, 3, 6, 5, 9} of Z11* | |
| | С | QR set = {1, 3, 4, 9,10} of Z11* | D | QR set = {1, 3, 4, 5, 9} of Z11* | |
| 127. | _v 7 ₁₇ | 7 (mod 29) | | | В |
| | X 11 | (1110a <u>-</u> 0) | | | |
| | | | В | x = 8, 10, 12, 15, 18, 26, 27 (mod 29) | |
| | Α | x = 8, 9, 12, 13, 15, 24, 28 (mod 29) x = 8, 10, 12, 15, 17, 24, 27 (mod 29) | | | |
| | A C | $x = 8, 9, 12, 13, 15, 24, 28 \pmod{29}$ | D | $x = 8, 9, 13, 15, 17, 24, 28 \pmod{29}$ | В |

| | abou | ıt x? | | | |
|------|---------|---|----------|---|---|
| | Α | x has no solution in (mod 29 15) | В | x has exactly one solution in (mod 29 15) | |
| | С | x has more than one solution in (mod 29 15) | D | x has four solutions in (mod 29 15) because there are two equations | |
| 129 | How | many primitive roots are there for 19? | | bedause there are two equations | D |
| 120. | A | 4 | В | 5 | |
| | C | 3 | D | 6 | |
| 130. | _ | GCD of x5+x4+x3 - x2 - x + 1 and x3 + | _ | | С |
| | Α | 1 | B | X | |
| | | x + 1 | | $x^2 + 1$ | |
| 131 | | t is the Discrete logarithm to the base 1 | | | В |
| 151. | A | 14 | В | 1 | |
| | C | 8 | D | 17 | |
| 132. | _ | a number x between 0 and 28 with x ⁸⁵ | _ | | С |
| | _ | | _ ~ | | |
| | A C | 22 6 | B D | 12 18 | |
| 122 | _ | t is the Discrete logarithm to the base 1 | _ | _ | Α |
| 133. | A | 12 | В | 14 | ^ |
| | C | 8 | D | 11 | |
| 134. | _ | the solution of x ² 16 mod 23 | | | В |
| | | | В | v – 4 and 10 | |
| | A C | x = 6 and 17 x = 11 and 12 | D | x = 4 and 19 x = 7 and 16 | |
| 135 | _ | | D | X = 7 and 10 | Α |
| 133. | | mod 11 = | D | - | ^ |
| | A C | 3 | В | 5 | |
| 126 | • | 6 | Dorf | 10 | |
| 130. | | itive Polynomial is also called a i) ucible Polynomial iv) Imperfect Polynomial iv) | | ect Polynomiai ii) Prime Polynomiai iii) | A |
| | A | ii) and iii) | В | only iii) | |
| | Ĉ | iv) and ii) | D | only i) | |
| 137 | (41)= | , | D | Offig 1) | Α |
| 107. | Α | 40 | В | 18 | ^ |
| | C | 20 | D | 22 | |
| 138. | _ | theorem is used to solve a set | of co | | С |
| | | lifferent moduli that are relatively prime | | | |
| | Α | Euclidean | В | Eulers | |
| | С | Chinese remainder | D | Fermat 's little | |
| 139. | If the | e multiplicative inverse of 53 modulo 21 | exist | s, then which of the following is true? | В |
| | Α | GCD(53,21)=29 | В | GCD(53,21)=1 | |
| | С | GCD(53,21)=53 | D | GCD(53,21)=12 | |
| 140. | | mposite is a positive integer with at lea | st | | С |
| | Α | four | В | three | |
| | С | two | D | zero | _ |
| 141. | | t is the Discrete logarithm to the base 1 | _ | | D |
| | A | 3 | В | 7 | |
| 4.40 | С | 12 | D | 4 | _ |
| 142. | _ | many primitive roots are there for 25? | D | E | D |
| | A C | 4 | B D | 5 8 | |
| 1/2 | _ | 7 | _ | | ٨ |
| 143. | _ | a number x between 0 and 28 with x ⁸⁵ | _ | | Α |
| | A | 6 | В | 32 | |
| 444 | С | 8 | D | 28 | _ |
| 144. | If n is | s a prime, mod n = ± | | | В |

| | Α | 0 | В | 1 | |
|------|----------|---|--------|----------------------------------|---|
| | С | p | D | n | |
| 145. | Find | the order of the group $G = \langle Z12^*, x \rangle$? | | | Α |
| | Α | 4 | В | 5 | |
| | С | 6 | D | 2 | |
| 146. | How | many primitive roots does Z<19> have | ? | | D |
| | Α | 5 | В | 8 | |
| | С | 7 | D | 6 | |
| 147. | (27)= | = | | | D |
| | À | 6 | В | 12 | |
| | С | 26 | D | 18 | |
| 148. | | function finds the number of inte | gers t | that are both smaller than n and | В |
| | | vely prime to n. | J | | |
| | | Euclidean | В | Eulers totient | |
| | С | Chinese remainder | D | Fermat 's little | |
| 149. | 7^3 mc | od 19 = | | | В |
| | A | 18 | В | 1 | |
| | Ĉ | 14 | D | 12 | |
| 150 | _ | modulus in the Fermat theorem is a | U | 12 | _ |
| 130. | A | odd | В | 0,400 | C |
| | C | | D ח | even | |
| | C | prime | ט | integer | |